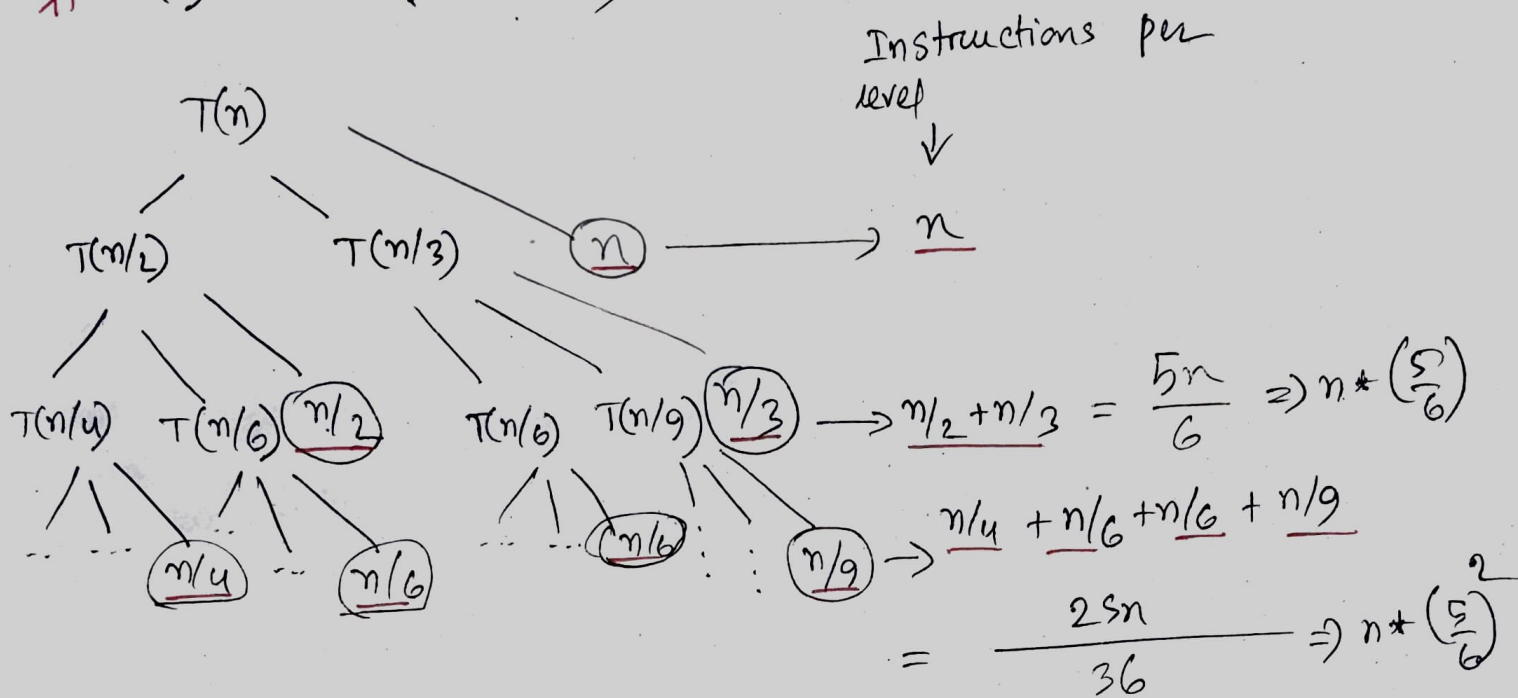


$$\# T(n) = T(n/2) + T(n/3) + O(n)$$



next level $\Rightarrow n * \left(\frac{5}{6}\right)^3$

so, at k^{th} level $\Rightarrow n * \left(\frac{5}{6}\right)^k$

total instructions $\Rightarrow n + \frac{5}{6} * n + \left(\frac{5}{6}\right)^2 * n + \left(\frac{5}{6}\right)^3 * n + \dots + \left(\frac{5}{6}\right)^k * n$

$\Rightarrow n \left(1 + \frac{5}{6} + \left(\frac{5}{6}\right)^2 + \left(\frac{5}{6}\right)^3 + \dots + \left(\frac{5}{6}\right)^k \right)$

[G.P. series formula $\Rightarrow a + ax + ax^2 + ax^3 + \dots + ax^n$
 $\Rightarrow \frac{a(x^{n+1} - 1)}{x - 1}$ / or, $\frac{a(1 - x^{n+1})}{1 - x}$]

depending if $x > 1$ or < 1 ; we use 2nd one since $x < 1$; $a = 1$, $x = \frac{5}{6}$

So, total instructions,

$$n * \left[\frac{1 - (5/6)^k}{(1 - 5/6)} \right] \text{ --- (1)}$$

now, since, the recursive function was, $T(n) = T(n/2) + T(n/3) + n$
we can safely assume that $T(n/3)$ will branch out and go to the base case long before the $T(n/2)$ term. So the number of level (k) depends on the $T(n/2)$ mostly.,

$$\text{So, } T(n/2^k) = 1 \quad [\text{last level, } k^{\text{th}} \text{ level}]$$

$$\Rightarrow k = \log_2 N.$$

$$\text{Putting it in eqn 1} \Rightarrow T(n) = n * \left[\frac{1 - (5/6)^{\log_2 N}}{1 - 5/6} \right]$$

we can neglect the constant $(1 - 5/6)$ & the value $(5/6)^{\log_2 N}$ will be very small. so we decide to ignore this as a constant term as it will have almost no effect compared to n .

$$\text{So, } T(n) = n * O(1)$$

$$\Rightarrow O(n) \quad \checkmark$$