Task 01

I solve this problem using Dijkstra's algorithm. I initialize a dictionary (distances) to store the shortest distances from the sounce node to each node. Initially, setting all the distances to iso except sounce node which is set to 0. I use a heap and explore nodes based on their current shortest prdistance and iteratively explore nodes and update the shortest distances to nodes and update the shorter path is found. After, neighbouring nodes if a shorter path is found. After, neighbouring all reachable nodes it finally return disexploring all reachable nodes it finally return disexploring all reachable nodes it finally return disexploring all print the output. If a node is not tances and print the output.

In this task, I apply the same thing like task of.
But here I colculate shortest distances from both
Africe's starting node and Bob's starting node. Then
Africe's starting node and Bob's starting nodes to find
compare distances from both starting nodes to find
the aormnon node where Africe and Bob can meet
the common node where Africe and Bob can meet
the aormnon node where time. It calculates
in the minimum amount of time. It calculates
in the minimum of the two shortest distances for
the maximum of the two shortest distances for
the node and selects the node with the minimum
each node and selects

maximum distance as the meeting point. If no armon node is found, it neturns impossible.

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Same as Task of I use Dijkstra Algorithm and heapy for this problem too. I explore nodes to neighbouring nodes, updating the danger level of each node based on the maximum of its current danger level and the weight of the edge to its neighbour. If the recolculated danger level is lower than the previous one, I updates the distance and add the neighbour to the Equeue for further exploration.

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