



# The algorithm showdown: A **theoretical** and **experimental** comparison of popular sorting algorithms

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## **Abstract**

Perhaps one of the most fundamental problems in Computer Science, algorithms that solve the problem of sorting are critical to anyone starting out in the area of Algorithms and Data Structures. They are a showcase of clever algorithmic thinking, as well as an application of common algorithm design techniques.

This research paper compares popular sorting algorithms such as Bubble Sort, Selection Sort, QuickSort, and more, using a theoretical as well as a practical, experimental approach.

Generally, this paper aims to emphasize that there is no “superior” sorting algorithm and that each algorithm has its ideal use case, its strengths and weaknesses. Such weaknesses are showcased, and moreover potential optimizations are discussed, in order to ensure algorithm reliability and stability.

The findings of this paper aim to guide developers, researchers in choosing the algorithm that suits their sorting problem the best.

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## Part I

# Introduction

## 1 An everyday sorting problem

Consider the following problem:

I have a T-shirt for every day of the week (7 *T-shirts*). I decide to wash them all together at the end of the week.

After they are all washed and dried, I have to put them back in the closet, in the order I am going to wear them (*ascending order of numbers*).

How can I do this in the **shortest time possible**, and using the **least amount of space possible** (*separating groups of T-shirts as little as possible*)?

**Note.** For analogy purposes, assume each T-shirt is laid facing-down, and its number can only be seen if picked up. Only one T-shirt can be picked up at a time.

The above is an over-simplified example of a **sorting problem** that an average person might encounter in everyday life.

In short, the sorting problem in computer science consists of organizing a collection of data, called *keys*, into a specific order (*in our case, in ascending order of numbers*).

### 1.1 The importance of sorting

Although at first, sorting might seem like a rare problem one might encounter, whether that be in their everyday life or in software development, there is a reason why it is considered the most fundamental problem in the study of algorithms:

- A Sorting is essential in data analysis, as it helps data scientists/analysts organize, find, visualize and understand data in order to make better, reliable decisions
- Many algorithms use sorting as a key subroutine, or require a sorted data structure as a prerequisite<sup>[1]</sup>.

Referring back to the T-shirt problem, suppose one has to search for the T-shirt number  $x$ ,  $1 \leq x \leq 7$ . Assuming the T-shirts are **sorted** in increasing order, the Binary Search algorithm can be applied, with a worst-case scenario runtime of  $O(\log n) \stackrel{n=7 \text{ (days)}}{=} O(\log 7) \approx O(3)$ , which means the search would require at most 3 steps, depending on the value of  $x$ .

- Since there is no “superior” sorting algorithm that can be used and relied on for any and every sorting problem, knowledge of various sorting algorithms and their respective strengths and weaknesses is critical.

The fastest sorting algorithm for a specific situation depends on many factors, such as the input data type (*key type*), how the data is stored (the *record*), prior knowledge about the data to be sorted<sup>[1]</sup>, and many more.

- The vast catalog of sorting algorithms that exist today employ a rich set of techniques. In fact, teaching various sorting algorithms also teaches various algorithm design techniques, along with algorithmic thinking.

This paper will discuss such design techniques in each case.

## 1.2 Possible solutions

- One solution is the so-called “brute-force” approach:  
Throw the pile of  $n = 7$  T-shirts in the air, and lay them out, eventually checking if they are sorted.

If they are sorted, celebrate.

If not, repeat until they become sorted.

This sorting algorithm is called **BogoSort**, and has an average-case complexity of  $\Theta(n \times n!)$ <sup>[2]</sup>, which in the case of our example evaluates to  $\Theta(n \times n!) = \Theta(7 \times 7!) = \Theta(7 \times 5040) = \Theta(35280)$ . Which means, one would have to repeat the process on average 35280 times. *Not to mention the fact that this process could repeat forever.*

- Another solution is laying out each T-shirt, and swapping each T-shirt with its neighbouring one, repeating the process  $n = 7$  times, until the T-shirts are laid out in ascending order.

This sorting algorithm is called **Bubble Sort**, and it will be explored in further detail later in this paper. It has a worst-case time complexity of  $O(n^2)$ , which in the case of our example evaluates to  $O(n^2) = O(7^2) = O(49)$ . Which means one would have to repeat the process at most 49 times.

As expected, these approaches do not ensure the shortest execution time.

**Transitioning to formal definitions.** Moving from a simple analogy to more complex, formal definitions and processes, we shall explore more time-efficient solutions to our sorting problem, as well as experimenting with larger, more complex types of inputs to be sorted.

## Declaration of originality

I, Raul-Andrei Arition, under the West University of Timișoara hereby present my research paper, titled *The algorithm showdown: A theoretical and experimental comparison of popular sorting algorithms*, and I declare that it is a product of my own research and investigation. Wherever deemed necessary, and to the best of my abilities, I have referenced and thus acknowledged the sources where I have derived ideas or extracts. Furthermore, I state that this paper has not been submitted anywhere else for publication.

The experiment process is done completely by myself, thus the experiment results are all original and not fabricated. Moreover, the plots, that visually represent my experiment results are all created using my code and mine only.

## 2 Reading Instructions

For the rest of the paper, the formal definition of the sorting problem will be presented, as well as various solutions of this problem (sorting algorithms).

For each solution, a short theoretical analysis will be conducted, proving its correctness, termination as well as its execution time complexity.

Afterwards, a general (*pseudocode*) implementation of the solution will be presented.

Throughout this paper, solutions will be categorized based off of the techniques they use, and for each such category, all of its solutions will be compared under various test cases (*input data types*), to then declare the most performant solution in that category.

Finally, all “winner” solutions of their respective category will be compared, in order to declare a “final winner”, i.e. solution which generally performs well in any test case.

### 3 The problem of sorting in Computer Science

Consider the following problem:

**Input:** A sequence of  $n$  numbers  $\langle a_1, a_2, \dots, a_n \rangle$

**Output:** A permutation (reordering)  $\langle a'_1, a'_2, \dots, a'_n \rangle$  of the inputs such that

$$a'_1 \preceq a'_2 \preceq \dots \preceq a'_n$$

where  $\preceq$  is an order relation (predicate).

The above describes the sorting problem in Computer Science<sup>[1]</sup>.

#### 3.1 Elements of a sorting problem

A sorting problem may consist of<sup>[1]</sup>:

- The **record**, the collection of data to be sorted.

*In the theoretical analyses of algorithms, the record is considered to be an array. In the implementations and experiments, however, Python lists will be used, which are essentially dynamic arrays, which grow along with the insertion of elements.*

- The **key**, the value to be sorted, or in other words, the value that determines the reordering of the record.

*In the theoretical analyses of algorithms, the keys are considered to be numerical values (numbers). In the experiments, however, various types of keys will be used.*

- The **satellite data**, which are values carried around with the key, but which do not determine the reordering of the record. In practice, when reordering keys, a sorting algorithm must permute the satellite data along with its respective key.

*The algorithm implementations in this paper are not implemented for such cases. Entries in the record consist purely of keys (the values that will be permuted)*

#### 3.2 Solutions to the sorting problem

The sorting problem in Computer Science is solved using sorting algorithms. Generally, an algorithm must be:

- Correct, which means:
  - its input satisfies the problem's preconditions<sup>1[3]</sup>
  - it terminates after a finite number of steps

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<sup>1</sup>conditions that the input data of the problem must satisfy.

- its output satisfies the problem's postconditions<sup>2[3]</sup>
- Efficient, i.e. it needs to utilize a *reasonable* amount of resources (*time, space in memory*).

**A note on space complexity.** This paper focuses mainly on comparing algorithms based on their execution time rather than the space they use in memory.

This is rather a personal choice, as machine memory as a resource is getting more and more affordable as a result of many impressive advances in the area of Computer Architecture.

However, time is a much more important resource to account for when analyzing and implementing algorithms.

Wherever possible, mentions of the algorithm's space complexity will be made.

Now that the formalities have been defined and explained, we shall move on to categorizing sorting algorithms and analyzing them one-by-one.

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<sup>2</sup>required properties of the problem's output.



## Part II

# The Algorithms

## 1 Classification of sorting algorithms

Below follow some terminologies that not only allow better organization of this paper, but also serve as a guide to researchers as well as learners to differentiate between the different types of sorting algorithms. Special acknowledgments to [4] for the extensive coverage on this topic.

A sorting algorithm can be categorized based on the following criteria:

- Comparison sorting or Non-Comparison Sorting
  - Comparison-based algorithms sort using a comparison operator between two elements in the array, which is defined using a function (*either built-in or user-defined*). The best known time complexity of a comparison-based sorting algorithm is  $O(n \log n)$ .
  - Algorithms which are not based on comparison of elements sort based on assumptions about the input data, thus can only be applied to input data which satisfy specific conditions. In a very ideal case, a Non-Comparison sorting algorithm can achieve a complexity of  $O(n)$ .
- Stability In the case where two values in the array are equal, their relative order, i.e. the order in which they appear in the original array order, must be preserved to consider a sorting algorithm *stable*.

If this relative order is not preserved, the algorithm is unstable, i.e. not stable

Few sorting algorithms are *naturally* stable, however most of the time any unstable sorting algorithm can be modified in order to become stable<sup>[3]</sup>.

- In-place or Out-of-place Sorting
  - An algorithm sorts in-place if only a constant number of elements of the input array get stored outside the array<sup>[1]</sup>. In practice, no copies of the array values are made outside of it, and the sorting is performed directly on the original array in memory.
  - An algorithm sorts out-of-place if it makes copies (also called *buffers*) of the array, or parts of it, and performs the sorting process on those copies, and not directly on the original array in memory.

## 2 In-place comparison-based sorting algorithms

### 2.1 Bubble Sort

**Main idea.** The array to be sorted is traversed repeatedly, adjacent elements are compared and swapped if they are out of order. The traversal stops when during a traversal no swapping has taken place (*i.e. when the array becomes sorted*).

**Correctness and complexity analysis**

**Implementation**

### 2.2 Insertion Sort

**Main idea.** The algorithm starts sorting the elements starting from the right-most element to the left-most element.

Starting with the *second element* in the array (**we consider the first element as sorted**), we compare it with every previous element (*for example, in the first instance, the second element gets compared with the first*) and proceed moving leftwards in the array, until we find an element such that if we place the current element after it, the sorting order will hold.

**Correctness and complexity analysis**

**Implementation**

### 2.3 Selection Sort

**Main idea.** Starting with the first element in the array, for each element at position  $i$  we search for a **minimum value** in the sub-array  $x[i \dots n]$ , where  $n$  is the size of the array.

If the minimum value found in the sub-array is less than the element at position  $i$ , we swap the element at position  $i$  with that minimum value (*hence* IF  $arr[x] < arr[i]$  THEN swap( $x[i], x[k]$ )).

This process is repeated  $n-1$  times, because after all the iterations, the last element will be **naturally sorted** (*hence* FOR  $i := 1, n-1$  DO).

**Correctness and complexity analysis**

**Implementation**

## 3 Not-in-place comparison-based sorting algorithms

### 3.1 Quick Sort

**Quick Sort** is different from what we've encountered so far. It uses the “**divide and conquer**” (also known as the “divide and rule”) technique. *Who knew that a policy used in ancient times to maintain power, by great rulers, could be implemented in solving a sorting problem?* Additionally, Quick Sort does not sort in place like all sorting algorithms mentioned so far.

**Main idea.**

**Corectness and complexity analysis**

**Implementation**

### 3.2 Merge Sort

**Main idea.**

**Corectness and complexity analysis**

**Implementation**

## **Part III**

# **Case Study**

- 1 Methodology**
- 2 In-place comparison-based sorting algorithms, compared**
- 3 Not-in-place comparison-based sorting algorithms, compared**
- 4 The “superior” sorting algorithm**

## **Part IV**

# **Conclusions and Ending**

## **1 Related Work**

## **2 Conclusions and Future Work**

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