

Lecture 13

Last time

- ▷ Affine invariance
- ▷ A new guarantee for Newton's.

Today

- ▷ Back to IPM
- ▷ A "complete" IPM
- ▷ Answering our questions

Back to IPM

Recall the n problem

$$\min_{x \in \mathbb{R}^d} f_n(x) \text{ with } f_n(x) = \eta c^T x + B(x)$$

and our informal template:

IPM (Informal)

- ▷ Pick $x_0 \in \text{int } P$ sufficiently close to $x^*(0)$ and pick $\eta_0 > 0$ small.
- ▷ Loop $k = 0, 1, \dots, T$:
 - ▷ Find an approximate minimizer x_{k+1} of f_{n_k} using x_k for initialization
 - ▷ Increase $\eta_{k+1} = q \eta_k$ with $q > 1$.

▷ Return x_T

and we had several questions

Q₁: What B function to use?

Q₂: How to find x_0 ? What is sufficiently close?

Q₃: What method to use to find x_{k+1} ?

Q₄: How to pick η ?

Q₅: How to show that the method finds an approximate solution to the original LP in poly time?

← Intuition

Description
New questions

Answering questions Today we answer all these questions except Q₅ (Exercise). Our ultimate goal is to find

$$c^T x - c^T x^* \leq \epsilon \text{ with } x \in P.$$

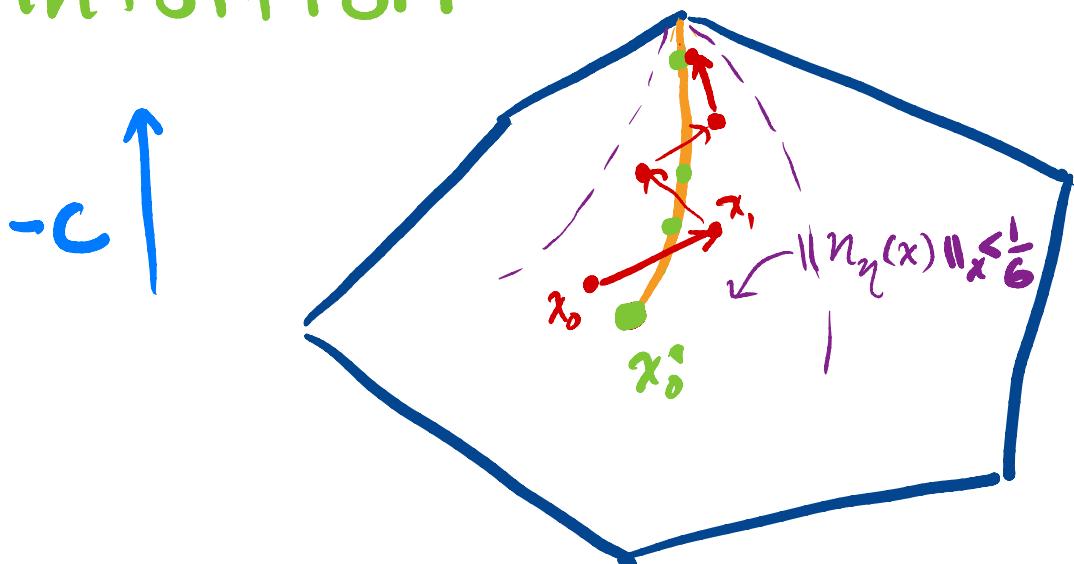
We will like to follow the

central path because we know $x_n^* \rightarrow x^*$. However, we don't need to solve for $x_{n_k}^*$ exactly, until the very end. Once we know

$$c^T x_{n_k}^* - c^T x^* \lesssim \epsilon.$$

So we could loosely track $x_{n_k}^*$ and then at the very end solve more accurately for $x_{n_k}^*$.

Intuition



Given what we proved last class it is natural

to expect that

$$\|n_n(x)\|_x \leq \frac{1}{6} \quad \text{with}$$

$$n_n(x) = -[\nabla^2 f_n(x)]^{-1} \nabla f(x).$$

is a natural condition to expect from all our iterates. This motivates

IPM (Full)

- ▷ Pick $x_0 \in \text{int } P$ s.t. $\|n_{n_0}(x_0)\| \leq \frac{1}{6}$
- ▷ Loop $k = 0, 1, \dots, T-1$: so that
 - ▷ $x_{k+1} \leftarrow x_k + n_{n_k}(x_k)$ $n_T > \frac{m}{2\epsilon}$
 - ▷ $n_{k+1} \leftarrow n_k \left(1 + \frac{1}{20m}\right)$
 - ▷ Increase $n_{k+1} = q n_k$ with $q > 1$.
- ▷ Run 2 steps of Newton for f_{n_T} starting from x_T .

$$A_0: B(x) = -\log \sum_{i=1}^m \log(b_i - a_i^T x).$$

The reason for this is that B is SC.

Lemma: The log barrier function

$$B(x) = -\sum_{i=1}^m \log(b_i - a_i^T x)$$

is self concordant. $r_i(x)$

Proof: Recall that

$$H(x) = \sum_{i=1}^m \frac{a_i a_i^T}{s_i(x)^2}$$

Let $\delta = \|y - x\|_2 < 1$, then

$$\delta^2 = (y - x)^T H(x) (y - x) = \sum_{i=1}^m \left(\frac{a_i^T (y - x)}{s_i(x)} \right)^2.$$

Hence, each individual term

$$\left(\frac{s_i(y) - s_i(x)}{s_i(x)} \right)^2 = \left(\frac{a_i^T (y - x)}{s_i(x)} \right)^2 \leq \delta^2.$$

Therefore,

$$(1-\delta)|s_i(x)| \leq |s_i(y)| \leq (1+\delta)|s_i(x)|$$

which implies

$$\frac{(1+\delta)^{-2}}{s_i(x)^2} \leq \frac{1}{s_i(y)^2} \leq \frac{(1-\delta)^{-2}}{s_i(x)^2}.$$

Thus

$$\frac{(1+\delta)^{-2}}{s_i(x)^2} a_i a_i^T \leq \frac{a_i a_i^T}{s_i(y)^2} \leq \frac{(1-\delta)^{-2}}{s_i(x)^2} a_i a_i^T.$$

Summing over all i and using the fact that

$$(1-3\delta) \leq (1+\delta)^{-2} \leq (1-\delta)^{-2} \leq 1+3\delta$$

for all $\delta \in [0, 1]$, yields the result. \square

This means that the theory we developed in Lecture 12 applies. Let

$$n_n(x) = -[\nabla^2 f_n(x)]^{-1} \nabla f(x).$$

Lemma +: For a given x , let $x_+ = x + n_n(x)$. Then if $\|n_n(x)\|_x \leq \frac{1}{6}$ we have

$$\|n_n(x_+)\|_{x_+} \leq 3 \|n_n(x)\|_x^2 \leq \frac{1}{12} +$$

A_1 : We do not answer how to get x_0 . But we notice it will suffice for it to satisfy

$$\|n_{x_0}(x_0)\|_{x_0} \leq \frac{1}{6}.$$

A_2 : We will run Newton's method for how many iterations are necessary to ensure $\|n_{x_{k+1}}(x_{k+1})\|_{x_{k+1}} \leq \frac{1}{6}$.

At the last iteration we might run it for longer to ensure

$$|C^T \hat{x} - C^T x^*| < \varepsilon.$$

Lemma 7: For every $x \in \text{int } P$ and $\eta'; \eta > 0$, we have

$$\|n_{\eta'}(x)\|_x \leq \frac{\eta'}{\eta} \|n_\eta(x)\|_x + \sqrt{m} \left| \frac{\eta'}{\eta} - 1 \right|.$$

Before proving this result,

notice that, it tells us how to increment n , since

$$\begin{aligned} \|n_{\eta_{k+1}}(x_{k+1})\|_{x_{k+1}} &\leq q \|n_{\eta_k}(x_{k+1})\|_{x_{k+1}} \\ &\leq \frac{q}{12} + \sqrt{m} |q-1|. \end{aligned}$$

A₃: We set $q = \left(1 + \frac{1}{20\sqrt{m}}\right)$

$$\begin{aligned} \|n_{\eta_{k+1}}(x_{k+1})\|_{x_{k+1}} &\leq \frac{1}{12} + \frac{1}{240\sqrt{m}} + \frac{1}{20} \\ (\heartsuit) \quad \text{One step suffices} &\rightarrow \leq \frac{1}{6}. \end{aligned}$$

Proof of Lemma n: We use $H(x) = \nabla^2 B(x)$ and $g(x) = \nabla B(x)$. Then

$$\begin{aligned} n_{\eta_1}(x) &= -H(x)^{-1}(\eta' c + g(x)) \\ &= -\frac{\eta'}{\eta} \underbrace{H(x)^{-1}(\eta c + g(x))}_{n_H(x)} \end{aligned}$$

$$+ \left(1 - \frac{n'}{n}\right) H(x)^{-1} g(x)$$

Taking the $\| \cdot \|_x$ and applying triangle inequality yields

$$\| n_{n'}(x) \|_x \leq \frac{n'}{n} \| n_n(x) \|_x + \left| \frac{n'}{n} - 1 \right| \| H(x)^{-1} g(x) \|_x.$$

It suffices to bound $\| H(x)^{-1} g(x) \|_x \leq \sqrt{m}$ for any $x \in \text{int } P$.

Applying Cauchy-Schwarz:

$$\begin{aligned} \| z \|_x^2 &= g(x)^T H(x)^{-1} g(x) \\ &= z^T g(x) \\ &= \sum_{i=1}^m 1 \cdot \frac{(z^T a_i)^2}{(b_i - a_i^T x)^2} \\ &\leq \sqrt{m} \sqrt{\sum_{i=1}^m \frac{(z^T a_i)^2}{(b_i - a_i^T x)^2}} \end{aligned}$$

$$\begin{aligned}
 &= \sqrt{m} \sqrt{z^T \left(\sum_i \frac{a_i a_i^T}{(b_i - a_i^T x)^2} \right) z} \\
 &= \sqrt{m} \|z_x\|. \quad H(x)
 \end{aligned}$$

We conclude $\|z\|_x \leq \sqrt{m}$. \square

Finally we need to decide how many steps to run for the last execution of Newton so that

$$C^T x_{T+1} - C^T x^* \leq \epsilon.$$

Proposition (B) Suppose $x \in \text{int } P$ and $n > 0$ s.t. $\|n_n(x)\|_x < \frac{1}{24}$, then

$$C^T x - C^T x^* \leq \frac{m}{n} (1 - 4 \|n_n(x)\|_x).$$

Notice that this proposition implies that if

$$\|n_n(x)\|_x \leq \frac{1}{24}$$

$$\frac{m}{2\epsilon} \leq n$$

Then, we obtain

$$c^T x - c^T x^* \leq \epsilon.$$

Recall that $\|n_{n_T}(x_T)\|_{x_T} < \frac{1}{6}$.

Therefore after 2 iterations of Newton

$$\begin{aligned} \|n_{n_T}(\hat{x})\|_{\hat{x}} &\leq 3 \|n_{n_T}(x_T^+)\|_{x_T^+}^2 \\ &\leq 27 \|n_{n_T}(x_T)\|_{x_T}^4 \\ &\leq \frac{1}{48} \end{aligned}$$

Thus,

$$C^T \hat{x} - C^T x^* \leq \frac{12}{11} \frac{m}{\eta_T} < \epsilon.$$

In turn, if we establish Proposition (B), we will have an answer for A4.

Theorem A4: IPM (Full) using $T = O(\sqrt{m} \log \frac{m}{\epsilon \eta_0})$ iterations, outputs a point $\hat{x} \in \text{int } P$ satisfying

$$C^T \hat{x} - C^T x^* \leq \epsilon.$$

Further, each step involves solving a linear system, which can be executed in polynomial time. \dagger