

Lecture 16

Last time

- ▷ Fixed points
- ▷ Set-valued mappings
- ▷ Krasnoselskij - Mann iteration.

Today

- ▷ Maximal monotone
- ▷ Douglas Rachford
- ▷ Augmented Lagrangian

Maximal monotone operator

Recall that given a monotone operator $T: E \rightarrow E$ we defined the resolvent as $R := (I + \alpha T)^{-1}$.

When $T = \partial f$ we showed that R was well-defined.

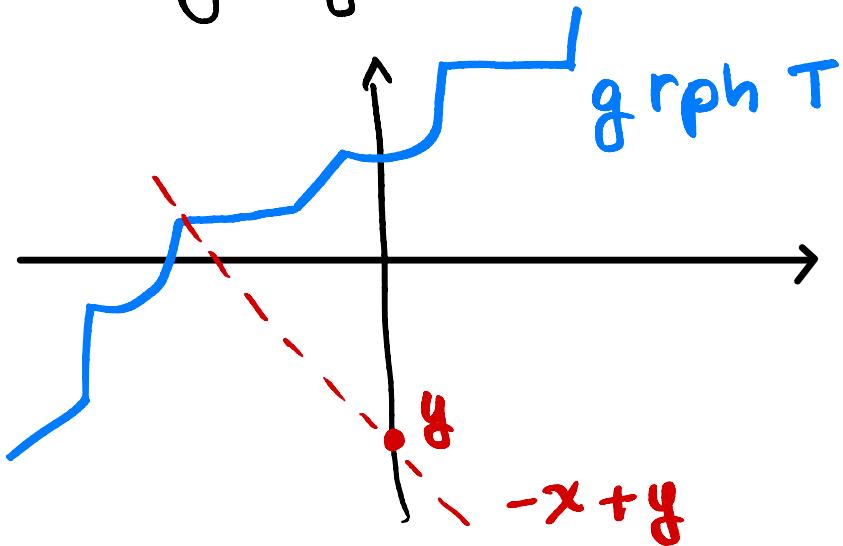
Q: How do we know $R(x) \neq \emptyset$?

Let's think about this

in $E = \mathbb{R}$. Monotonicity in 1D makes the graph

$$\text{grph } T = \{(x, y) \mid y \in T(x)\}$$

look like a monotonously non-decreasing function (with jumps)



Then, the following is solvable by

$$y - x \in T(x) \quad (\Leftrightarrow)$$

if, and only if, the graph of
 $x \mapsto -x + y$ intersects $\text{grph } T$
for all y . Thus, we want
 T to be as "large" as

possible.

Proposition: Let $T: E \rightrightarrows E$ a monotone operator and $\alpha > 0$ be arbitrary. Assume (\approx) is solvable Hg. Then, T is maximal.

i.e., if $T': E \rightrightarrows$ monotone $\xrightarrow{\text{s.t.}}$
 $T(x) \subseteq T'(x) \quad \forall x \Rightarrow T = T'$!

Proof: Suppose $y \in T'(x)$, we want to show $y \in T(x)$. Note that $\alpha y \in \alpha T'(x)$ and by assumption

$$(x + \alpha y) - z \in \alpha T(z) \subseteq \alpha T'(z).$$

By monotonicity

$$\begin{aligned} -\|x - y\|^2 &= \langle (x + \alpha y - z) - \alpha y, z - x \rangle \\ &\geq 0. \end{aligned}$$

Therefore $z = x$, and so
 $\alpha y = (x + \alpha y) - x \in \alpha T(x)$.

Proving the result. \square

A direct implication of this result is:

Corollary: For a closed, convex, proper $f: E \rightarrow \mathbb{R} \cup \{\infty\}$, we have that ∂f is maximal monotone.

In fact the opposite also holds. \dashv

Theorem (Minty 1962). Let $T: E \rightrightarrows E$ be a monotone operator and $\alpha > 0$. Then, T is maximal if, and only if, (\Leftarrow)

is solvable for all y . +

We will not prove this result.
It is usually simpler to show
directly that R is well-defined.

Baby steps: Smooth optimization

Consider the problem of minimizing
 $f: E \rightarrow \mathbb{R}$ a smooth convex
function with L -Lipschitz
gradient. A natural strategy
is gradient descent

$$x_{k+1} \leftarrow x_k - \alpha \nabla f(x_k),$$

Theorem: Let f be an L -smooth
convex function. Then, we have
 $\langle x - y, \nabla f(x) - \nabla f(y) \rangle \geq \frac{1}{L} \|\nabla f(x) - \nabla f(y)\|^2 \quad (\infty)$
 $\forall x, y$

Therefore,

$$T_\alpha(x) = x - \alpha \nabla f(x)$$

is non expansive if $\alpha \in [0, \frac{2}{L}]$ and it is averaged if $\alpha \in (0, \frac{2}{L})$.

Thus, if $\alpha \in (0, \frac{2}{L})$, gradient descent converges.

Proof: We derived (x) in Lecture 6 of Non linear 1 (see notes).

Then, we can prove nonexpansiveness

$$\begin{aligned} & \|x - \alpha \nabla f(x) - y + \alpha \nabla f(y)\|^2 \\ &= \|x - y\|^2 - 2\alpha \langle x - y, \nabla f(x) - \nabla f(y) \rangle \\ &\quad + \alpha^2 \|\nabla f(x) - \nabla f(y)\|^2 \\ &\leq \|x - y\|^2 - \left(2\frac{\alpha}{L} - \alpha^2\right) \|\nabla f(x) - \nabla f(y)\|^2 \end{aligned}$$

$\overbrace{\alpha}^{\text{(x)}}$ $\underbrace{2\frac{\alpha}{L} \text{ iff } \alpha \in [0, \frac{2}{L}]}$

Moreover, if $\alpha \in (0, \frac{2}{L})$

$$T_\alpha(x) = \theta x + (1-\theta) \left(x - \frac{\alpha}{(1-\theta)} \nabla f(x) \right)$$

We can pick θ so that
 $\frac{\alpha}{(1-\theta)} \in [0, \frac{2}{L}]$.

Convergence follows immediately via the KM iteration result. \square

Augmented Lagrangian method
Let us go back to the problem of solving

$$\begin{aligned} & \min f(x) \quad \text{smooth and convex.} \\ & \text{s.t. } g_i(x) \leq 0 \quad \forall i \in [m]. \end{aligned}$$

Before presenting the algorithm, let us motivate it. Notice

$$\begin{aligned} & \inf_x \{ f(x) \mid g(x) \leq 0 \} \\ &= \inf_{\substack{x \\ z \geq 0}} \{ f(x) \mid g(x) + z = 0 \} \end{aligned}$$

$$= \inf_{\substack{x \\ z \geq 0}} \{ f(x) + \frac{1}{2} \|g(x) + z\|^2 \mid g(x) + z = 0 \}$$

$$= \inf_{\substack{x \\ z \geq 0}} \sup_{\lambda} \left\{ f(x) + \frac{1}{2} \|g(x) + z\|^2 + \lambda^T(g(x) + z) \right\}$$

$$\geq \sup_{\lambda} \inf_{\substack{x \\ z \geq 0}} \left\{ f(x) + \frac{1}{2} \|g(x) + z\|^2 + \lambda^T(g(x) + z) \right\}$$

(Why?)

$$= \sup_{\lambda} \inf_{\substack{x \\ z \geq 0}} \left\{ f(x) + \frac{1}{2} \|(\lambda + g(x))_+\|^2 - \frac{1}{2} \|\lambda\|^2 \right\}$$

Augmented Lagrangian $\bar{L}(x; \lambda)$

Recall that we defined, the standard Lagrangian in lecture 8

$$L(x; z) = f(x) + z^T g(x).$$

Proposition: Suppose $f, g_i : E \rightarrow \mathbb{R}$ are convex, and $\lambda \geq 0$, $z = (z + g(x))_+$. Then, x minimizes $\bar{L}(\cdot, \lambda)$ if, and only if, x minimizes $L(\cdot, z)$.

Proof: Exercise.

□

This motivates a natural alternating algorithm where we minimize $\bar{L}(\cdot; \lambda)$ and then maximize $\bar{L}(x; \cdot)$. Let's understand how the maximizer of $\bar{L}(x; \cdot)$ looks like via:

Claim: If $h: E \rightarrow \mathbb{R}$ convex then so is h_+^2 and $\partial(h_+^2)(x) = 2h_+(x)\partial h(x)$. \rightarrow

For any fixed $x \in E$, the function

$$\lambda_i \mapsto \frac{1}{2} (\lambda_i + g_i(x))_+^2 - \frac{1}{2} \lambda_i^2$$

is concave and by first order optimality conditions it is maximized at

$$(\lambda_i + g_i(x))_+ - \lambda_i = 0.$$

Thus λ is a fixed point of $T(\lambda) = (\lambda + g(x))_+$. With this we can now introduce the Augmented Lagrangian method.

Augmented Lagrangian Method (ALM)

Pick $x_0 \in E$, $\lambda_0 \in \mathbb{R}^m$.

Loop $K \geq 0$:

$$x_{k+1} \leftarrow \underset{x}{\operatorname{argmin}} \bar{L}(x, \lambda_k)$$

$$\lambda_{k+1} \leftarrow (\lambda_k + g(x_{k+1}))_+$$

Let's make a couple of remarks:

1) What did we gain? We reduce a constrained problem into a sequence of unconstrained ones.

If we have an effective solver for $\bar{L}(\cdot; \lambda_k)$, this might be a reasonable algorithm.

2) Notice that in general, x_{k+1} is not well-defined since

the minimizer might not be well defined.

However, if, we assume that $x(\lambda) \leftarrow \operatorname{argmin}_x L(x; \lambda)$ is well-defined^{*} for all the λ 's we see, then we can analyze this method via the KM iteration. Consider the function

$$\varphi(\lambda) = \begin{cases} -\inf_x L(x, \lambda) & \text{if } y \geq 0, \\ +\infty & \text{otherwise.} \end{cases}$$

Proposition: Suppose that x_{k+1} is well-defined for all λ_k in an execution of ALM. Then,

$$\lambda_{k+1} \leftarrow \operatorname{prox}_{\varphi}(\lambda_k) \quad \forall k \geq 0.$$

Consequently, λ_{k+1} converges. +

The proof is left as an exerci
se. The final conclusion follows
from our KM iteration result
since φ is convex. One can
also show that any cluster point
of the iterates $\{x_k\}$ is a solution
to the original problem.

Warning These results are some-
what unsatisfactory since we needed
to assume the well-posedness of the
primal iterates. Next we will see methods
that do not suffer from this issue.