

## DAA DAY 9

1. There are  $3n$  piles of coins of varying size, you and your friends will take piles of coins as follows: In each step, you will choose any 3 piles of coins (not necessarily consecutive). Of your choice, Alice will pick the pile with the maximum number of coins. You will pick the next pile with the maximum number of coins. Your friend Bob will pick the last pile. Repeat until there are no more piles of coins. Given an array of integers `piles` where `piles[i]` is the number of coins in the  $i$ th pile. Return the maximum number of coins that you can have.

```
def maxCoins(piles):
```

```
    piles.sort(reverse=True)
```

```
    max_coins = 0
```

```
    for i in range(1, len(piles) * 2 // 3, 2):
```

```
        max_coins += piles[i]
```

```
    return max_coins
```

```
piles = [2,4,1,2,7,8]
```

```
print(maxCoins(piles))
```

2. You are given a 0-indexed integer array `coins`, representing the values of the coins available, and an integer `target`. An integer  $x$  is obtainable if there exists a subsequence of coins that sums to  $x$ . Return the minimum number of coins of any value that need to be added to the array so that every integer in the range  $[1, \text{target}]$  is obtainable. A subsequence of an array is a new non-empty array that is formed from the original array by deleting some (possibly none) of the elements without disturbing the relative positions of the remaining elements.

```
def minCoins(coins, target):
```

```
    coins.sort()
```

```
    count = 0 # The number of coins we need to add
```

```
    current_sum = 0 # This represents the current maximum sum we can obtain
```

```
    for coin in coins:
```

```
        while current_sum + 1 < coin:
```

```
            count += 1
```

```

    current_sum += current_sum + 1

    if current_sum >= target:

        return count

    current_sum += coin

    if current_sum >= target:

        return count

while current_sum < target:

    count += 1

    current_sum += current_sum + 1

return count

coins = [1, 3]

target = 6

print(minCoins(coins, target))

```

3. You are given an integer array jobs, where jobs[i] is the amount of time it takes to complete the ith job. There are k workers that you can assign jobs to. Each job should be assigned to exactly one worker. The working time of a worker is the sum of the time it takes to complete all jobs assigned to them. Your goal is to devise an optimal assignment such that the maximum working time of any worker is minimized. Return the minimum possible maximum working time of any assignment.

```

def canAssign(jobs, k, max_time, workers):

    if not jobs:

        return True

    current_job = jobs.pop()

    for i in range(k):

        if workers[i] + current_job <= max_time:

            workers[i] += current_job

            if canAssign(jobs, k, max_time, workers):

```

```

        return True

    workers[i] -= current_job

    if workers[i] == 0:

        break

    jobs.append(current_job)

    return False

def minimumTimeRequired(jobs, k):

    left, right = max(jobs), sum(jobs)

    jobs.sort(reverse=True)

    while left < right:

        mid = (left + right) // 2

        if canAssign(jobs[:], k, mid, [0] * k):

            right = mid

        else:

            left = mid + 1

    return left

jobs = [3, 2, 3]

k = 3

print(minimumTimeRequired(jobs, k))

```

4. We have  $n$  jobs, where every job is scheduled to be done from  $startTime[i]$  to  $endTime[i]$ , obtaining a profit of  $profit[i]$ . You're given the  $startTime$ ,  $endTime$  and  $profit$  arrays, return the maximum profit you can take such that there are no two jobs in the subset with overlapping time range. If you choose a job that ends at time  $X$  you will be able to start another job that starts at time  $X$ .

```

from bisect import bisect_right

def jobScheduling(startTime, endTime, profit):

    jobs = sorted(zip(startTime, endTime, profit), key=lambda x: x[1])

```

```

dp = [0] * len(jobs)
dp[0] = jobs[0][2]
for i in range(1, len(jobs)):
    current_profit = jobs[i][2]
    index = bisect_right([jobs[j][1] for j in range(i)], jobs[i][0]) - 1
    if index != -1:
        current_profit += dp[index]
    dp[i] = max(dp[i-1], current_profit)
return dp[-1]

startTime = [1, 2, 3, 4, 6]
endTime = [3, 5, 10, 6, 9]
profit = [20, 20, 100, 70, 60]

print(jobScheduling(startTime, endTime, profit)) # Output: 150

```

5. Given a graph represented by an adjacency matrix, implement Dijkstra's Algorithm to find the shortest path from a given source vertex to all other vertices in the graph. The graph is represented as an adjacency matrix where  $graph[i][j]$  denote the weight of the edge from vertex  $i$  to vertex  $j$ . If there is no edge between vertices  $i$  and  $j$ , the value is Infinity (or a very large number).

```

import heapq
import sys

def dijkstra(graph, src):
    n = len(graph)
    dist = [sys.maxsize] * n
    dist[src] = 0
    pq = [(0, src)] # Priority queue to get the vertex with the minimum distance
    while pq:
        current_dist, u = heapq.heappop(pq)

```

```

    if current_dist > dist[u]:
        continue
    for v in range(n):
        if graph[u][v] and dist[u] + graph[u][v] < dist[v]:
            dist[v] = dist[u] + graph[u][v]
            heapq.heappush(pq, (dist[v], v))
    return dist

graph = [
    [0, 10, 20, sys.maxsize, sys.maxsize],
    [10, 0, 30, 50, 10],
    [20, 30, 0, 20, 33],
    [sys.maxsize, 50, 20, 0, 2],
    [sys.maxsize, 10, 33, 2, 0]
]

source_vertex = 0

distances = dijkstra(graph, source_vertex)

print(distances)

```

6. Given a graph represented by an edge list, implement Dijkstra's Algorithm to find the shortest path from a given source vertex to a target vertex. The graph is represented as a list of edges where each edge is a tuple (u, v, w) representing an edge from vertex u to vertex v with weight w.

```

import heapq
import sys
from collections import defaultdict

def dijkstra(edges, n, src, target):
    graph = defaultdict(list)
    for u, v, w in edges:

```

```

graph[u].append((v, w))
graph[v].append((u, w))
dist = {i: sys.maxsize for i in range(n)}
dist[src] = 0
pq = [(0, src)]
while pq:
    current_dist, u = heapq.heappop(pq)
    if u == target:
        return current_dist
    for v, weight in graph[u]:
        if current_dist + weight < dist[v]:
            dist[v] = current_dist + weight
            heapq.heappush(pq, (dist[v], v))
return -1 # Return -1 if the target is unreachable

edges = [
    (0, 1, 10), (0, 2, 5),
    (1, 2, 2), (1, 3, 1),
    (2, 1, 3), (2, 3, 9),
    (2, 4, 2), (3, 4, 4),
    (4, 0, 7), (4, 3, 6)
]

n = 5
src = 0
target = 3
print(dijkstra(edges, n, src, target))

```

7. Given a set of characters and their corresponding frequencies, construct the Huffman Tree and generate the Huffman Codes for each character.

```
import heapq

from collections import defaultdict, Counter

class Node:

    def __init__(self, char=None, freq=0):

        self.char = char

        self.freq = freq

        self.left = None

        self.right = None

    def __lt__(self, other):

        return self.freq < other.freq

def huffman_codes(frequencies):

    heap = [Node(char, freq) for char, freq in frequencies.items()]

    heapq.heapify(heap)

    while len(heap) > 1:

        left = heapq.heappop(heap)

        right = heapq.heappop(heap)

        merged = Node(freq=left.freq + right.freq)

        merged.left = left

        merged.right = right

        heapq.heappush(heap, merged)

    def generate_codes(node, prefix="", codebook={}):

        if node:

            if node.char is not None:

                codebook[node.char] = prefix

            generate_codes(node.left, prefix + "0", codebook)

            generate_codes(node.right, prefix + "1", codebook)

    generate_codes(heap[0])

    return codebook
```

```

        generate_codes(node.left, prefix + "0", codebook)

        generate_codes(node.right, prefix + "1", codebook)

    return codebook

root = heapq.heappop(heap)

return generate_codes(root)

frequencies = Counter({'a': 5, 'b': 9, 'c': 12, 'd': 13, 'e': 16, 'f': 45})

huffman_code = huffman_codes(frequencies)

print(huffman_code)

```

8. Given a Huffman Tree and a Huffman encoded string, decode the string to get the original message.

**class Node:**

```

    def __init__(self, char=None, left=None, right=None):

        self.char = char

        self.left = left

        self.right = right

def decode_huffman(root, encoded_str):

    decoded_str = []

    current = root

    for bit in encoded_str:

        if bit == '0':

            current = current.left

        else: # bit == '1'

            current = current.right

    if current.char is not None:

        decoded_str.append(current.char)

        current = root # Reset to the root for the next character

```



```

    return ''.join(decoded_str)

root = Node()

root.left = Node()

root.right = Node('A')

root.left.left = Node('B')

root.left.right = Node('C')

encoded_str = "0110" # Represents "BC"

decoded_message = decode_huffman(root, encoded_str)

print(decoded_message)

```

9. Given a list of item weights and the maximum capacity of a container, determine the maximum weight that can be loaded into the container using a greedy approach. The greedy approach should prioritize loading heavier items first until the container reaches its capacity.

```

def max_weight_greedy(weights, max_capacity):

    weights.sort(reverse=True) # Sort weights in descending order

    total_weight = 0

    for weight in weights:

        if total_weight + weight <= max_capacity:

            total_weight += weight

        else:

            break

    return total_weight

weights = [10, 20, 30, 40, 50]

max_capacity = 100

print(max_weight_greedy(weights, max_capacity))

```

10. Given a list of item weights and a maximum capacity for each container, determine the minimum number of containers required to load all items using a greedy approach. The greedy approach should prioritize loading items into the current container until it is full before moving to the next container.

```
def min_containers(weights, capacity):  
  
    weights.sort(reverse=True) # Sort weights in descending order  
  
    containers = []  
  
    for weight in weights:  
  
        placed = False  
  
        for i in range(len(containers)):  
  
            if containers[i] + weight <= capacity:  
  
                containers[i] += weight  
  
                placed = True  
  
                break  
  
        if not placed:  
  
            containers.append(weight) # Need a new container  
  
    return len(containers)  
  
weights = [10, 20, 30, 40, 50]  
  
capacity = 100  
  
print(min_containers(weights, capacity))
```

11. Given a graph represented by an edge list, implement Kruskal's Algorithm to find the Minimum Spanning Tree (MST) and its total weight.

```
import heapq  
  
class UnionFind:  
  
    def __init__(self, size):  
  
        self.parent = list(range(size))  
  
        self.rank = [0] * size
```

```

def find(self, u):
    if self.parent[u] != u:
        self.parent[u] = self.find(self.parent[u])
    return self.parent[u]

def union(self, u, v):
    rootU = self.find(u)
    rootV = self.find(v)
    if rootU != rootV:
        if self.rank[rootU] > self.rank[rootV]:
            self.parent[rootV] = rootU
        elif self.rank[rootU] < self.rank[rootV]:
            self.parent[rootU] = rootV
        else:
            self.parent[rootV] = rootU
            self.rank[rootU] += 1

def kruskal(n, edges):
    uf = UnionFind(n)
    edges.sort(key=lambda x: x[2])
    mst_weight = 0
    mst_edges = []
    for u, v, weight in edges:
        if uf.find(u) != uf.find(v):
            uf.union(u, v)
            mst_edges.append((u, v, weight))
            mst_weight += weight
    return mst_edges, mst_weight

```

```
edges = [  
    (0, 1, 10), (0, 2, 6), (0, 3, 5),  
    (1, 3, 15), (2, 3, 4)  
]  
  
n = 4  
  
mst_edges, mst_weight = kruskal(n, edges)  
  
print("Edges in MST:", mst_edges)  
  
print("Total weight of MST:", mst_weight)
```