

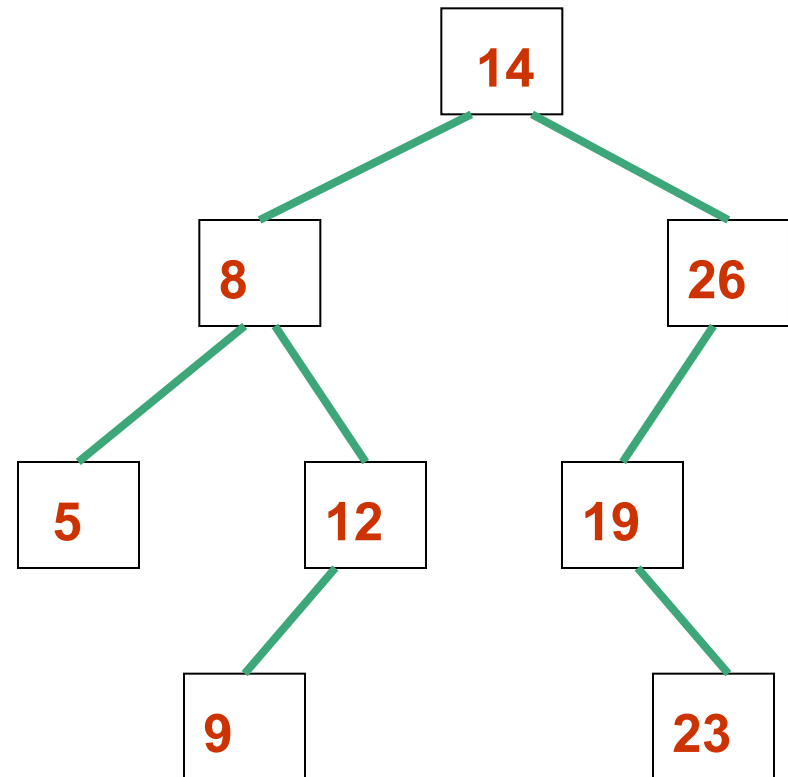
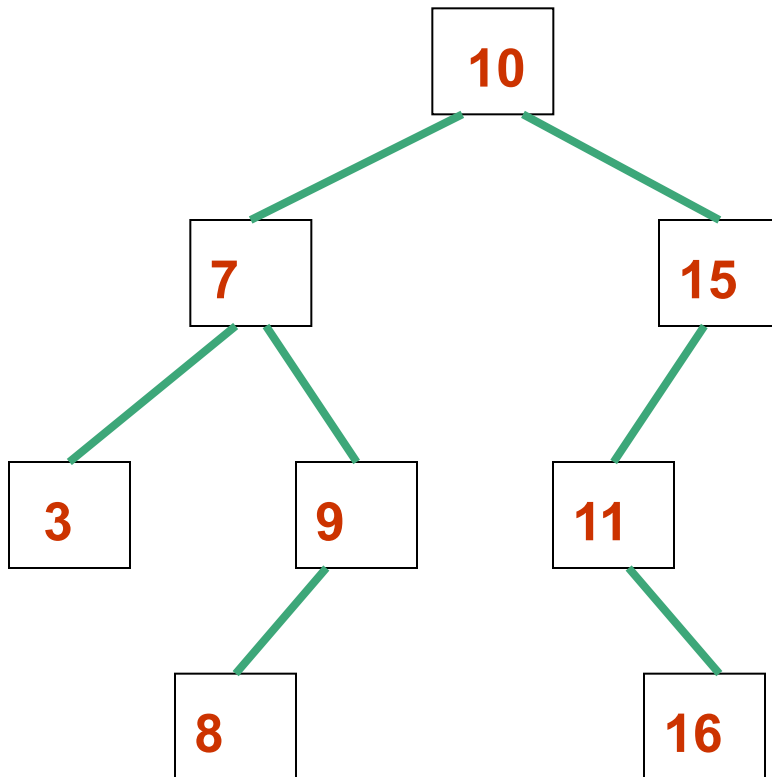
Topic 15

The Binary Search Tree ADT

Binary Search Tree

- A **binary search tree (BST)** is a binary tree with an **ordering** property of its elements, such that the data in any internal node is
 - **Greater than** the data in any node in its left subtree
 - **Less than** the data in any node in its right subtree
- **Note:** this definition does not allow duplicates; some definitions do, in which case we could say “**less than or equal to**”

Examples: are these Binary Search Trees?



Discussion

- Observations:
 - What is in the leftmost node?
 - What is in the rightmost node?

BST Operations

- A binary search tree is a special case of a binary tree
 - So, it has all the operations of a binary tree
- It also has *operations specific to a BST*:
 - *add* an element (requires that the BST property be maintained)
 - *remove* an element (requires that the BST property be maintained)
 - *remove the maximum* element
 - *remove the minimum* element

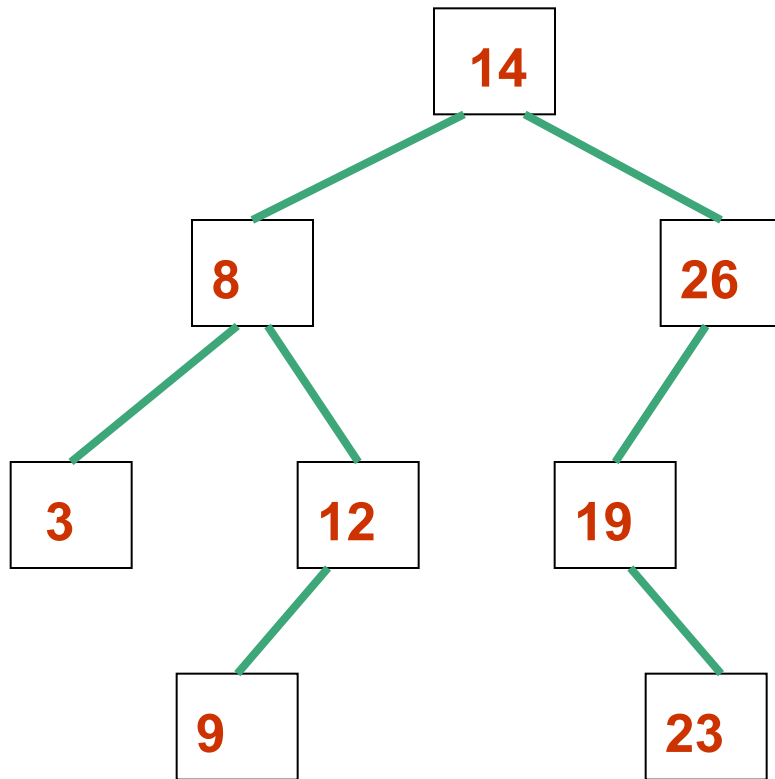
Searching in a BST

- Why is it called a binary **search** tree?
 - Data is stored in such a way, that it can be more **efficiently** found than in an ordinary binary tree

Searching in a BST

- **Algorithm to *search* for an item in a BST**
 - Compare data item to the root of the (sub)tree
 - If data item = data at root, found
 - If data item < data at root, go to the left; if there is no left child, data item is not in tree
 - If data item > data at root, go to the right; if there is no right child, data item is not in tree

Search Operation – a Recursive Algorithm



*To search for a value k ; returns **true** if found or **false** if not found*

If the tree is empty, return **false**.

If $k ==$ value at root

return **true**: we're done.

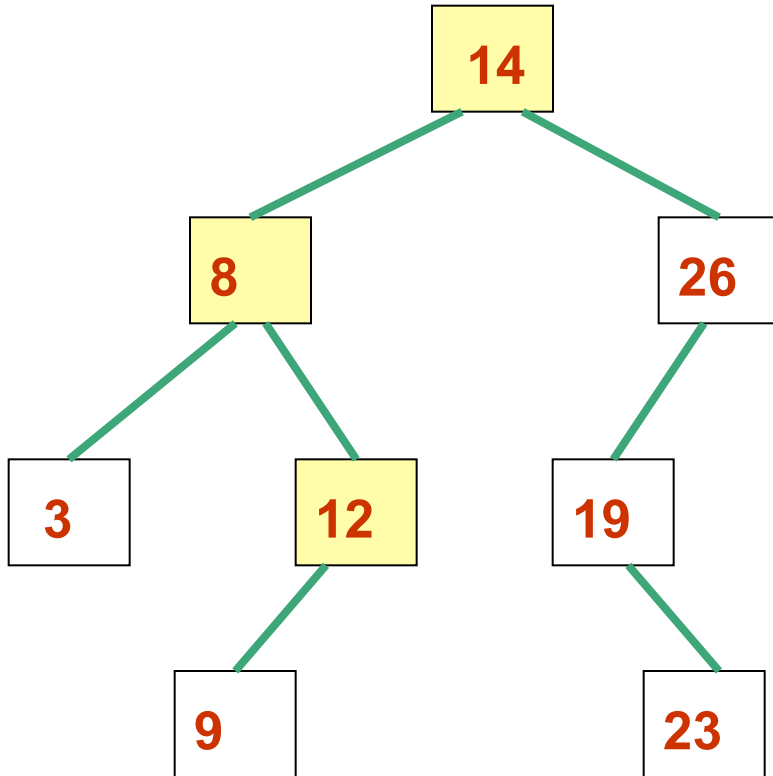
If $k <$ value at root

return result from **search for k** in the left subtree

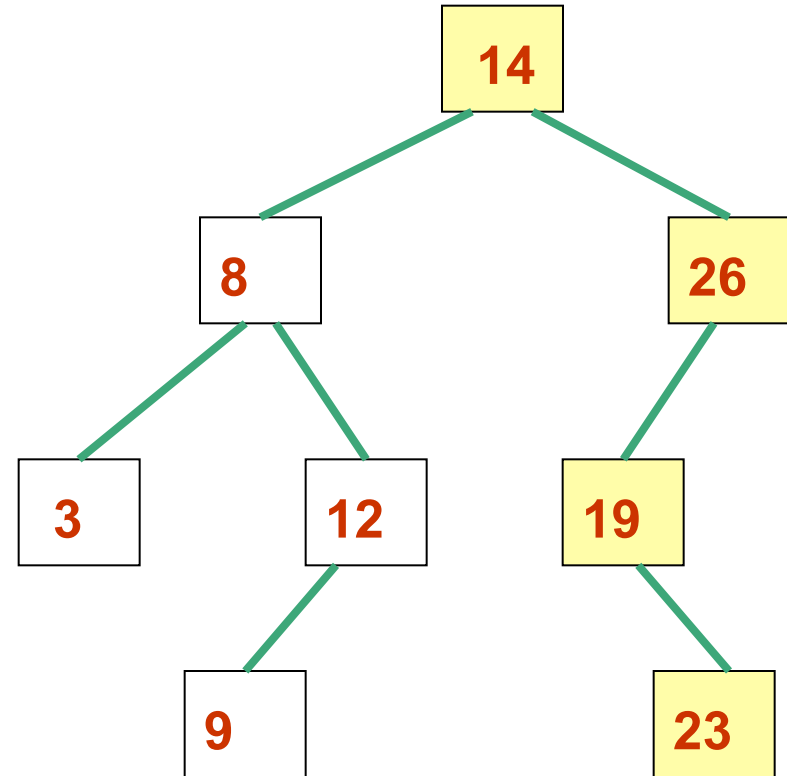
Else

return result from **search for k** in the right subtree.

Search Operation



Search for 13: visited nodes are coloured yellow; return false when node containing 12 has no right child

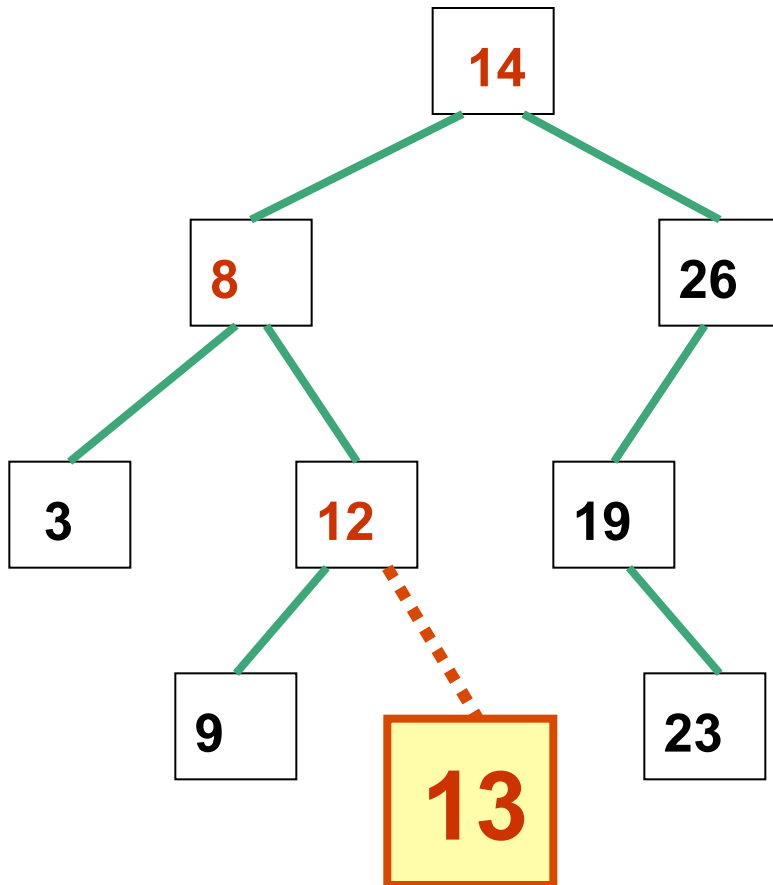


Search for 22: return false when node containing 23 has no left child

BST Operations: **add**

- To **add** an item to a BST:
 - Follow the algorithm for searching, until there is no child
 - Insert at that point
- So, new node will be added as a leaf
- (We are assuming no duplicates allowed)

Add Operation



To insert 13:

Same nodes are visited as when *searching* for 13.

Instead of returning *false* when the node containing 12 has no right child, build the new node, attach it as the right child of the node containing 12, and return *true*.

Add Operation – an Algorithm

*To insert a value k into a tree, returning **true** if successful and **false** if not*

Build a new node for k .

If tree is empty

 add new node as root node, return **true**.

If $k ==$ value at root

 return **false** (no duplicates allowed).

If $k <$ value at root

 If root has no left child

 add new node as left child of root, return **true**

 Else **insert k** into left subtree of root.

If $k >$ value at root

 If root has no right child

 add new node as right child of root, return **true**

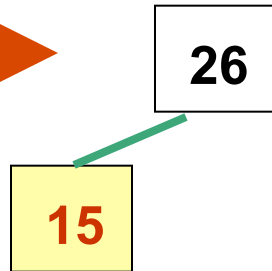
 Else **insert k** into the right subtree of root.

Example: Adding Elements to a BST

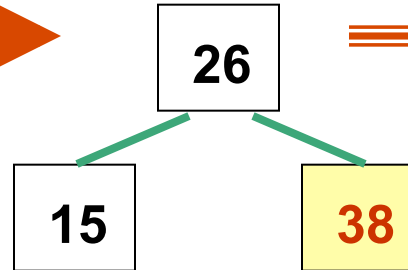
1: Add 26



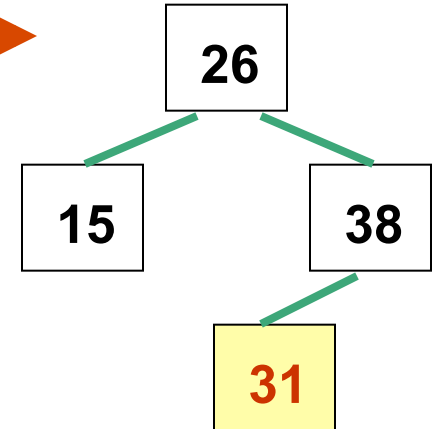
2: Add 15



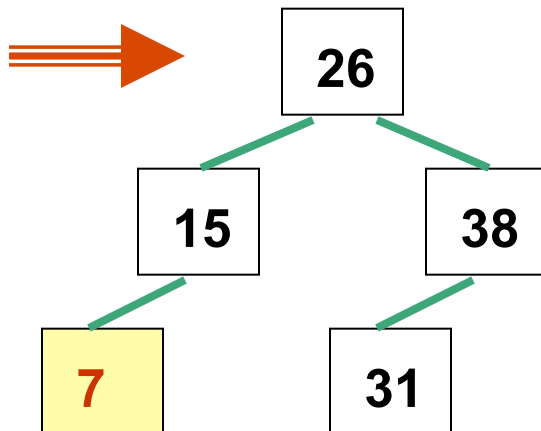
3: Add 38



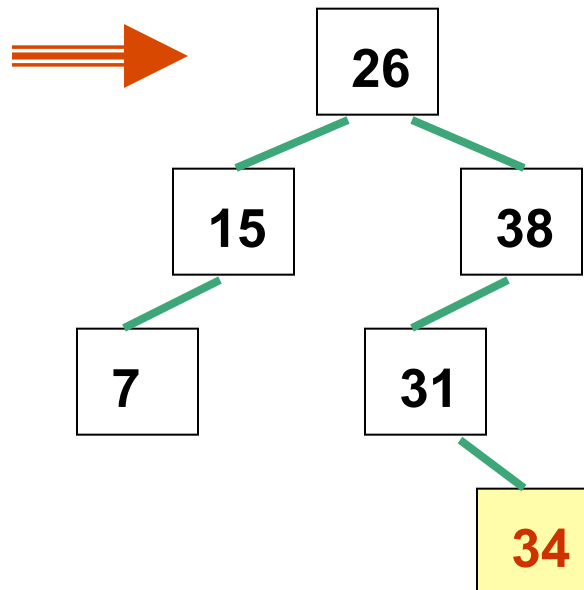
4: Add 31



5: Add 7



5: Add 34

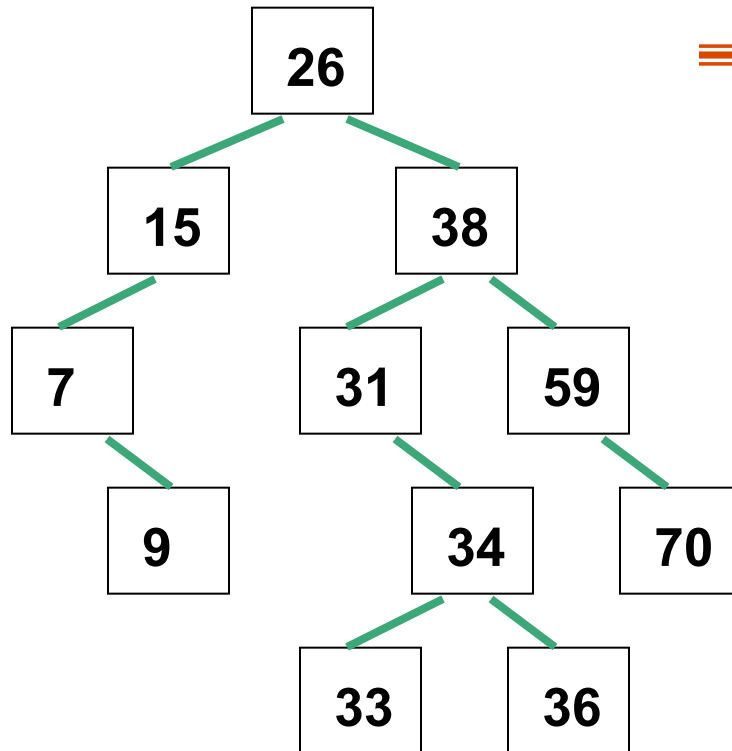


BST Operations: Remove

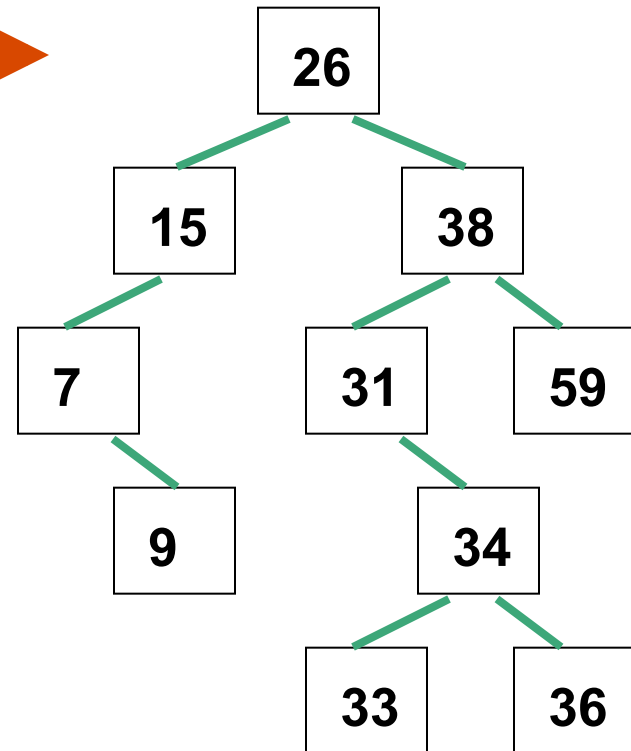
- **Case 1:** *value to be removed is in a leaf node*
 - Node can be removed and the tree needs no further rearrangement
- **Case 2:** *value to be removed is in an interior node*
 - Why can't we just change the link from its parent node to a successor node?
 - We *can* replace the node with its *inorder* predecessor (or successor)
 - Complex, and we will not implement this

Example: Removing BST Elements

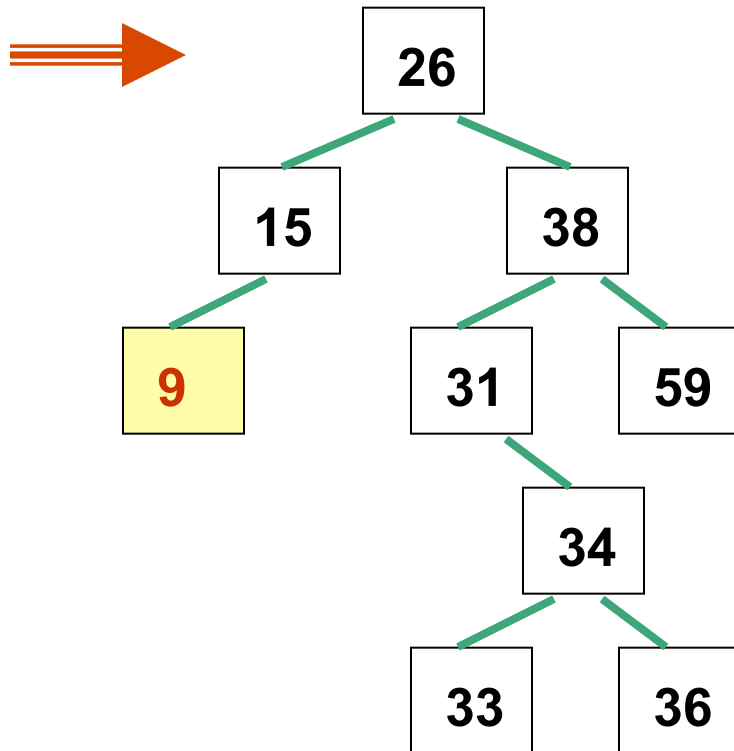
1: Initial tree



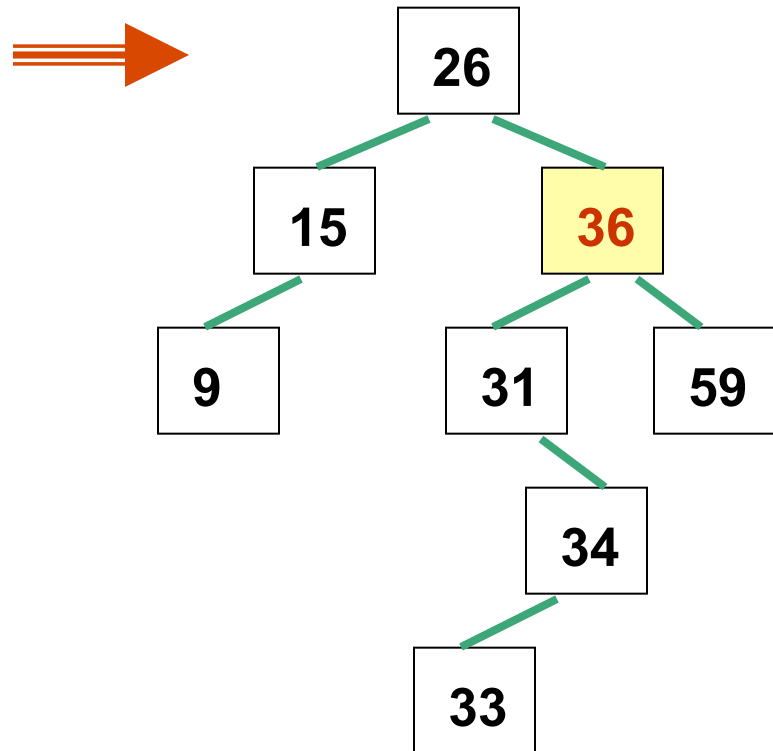
2: Remove 70



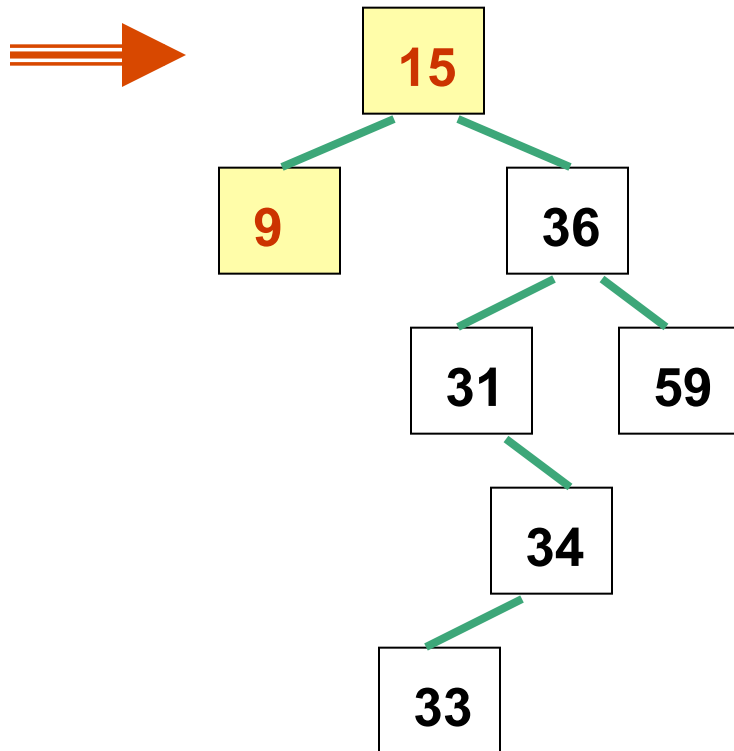
3: Remove 7



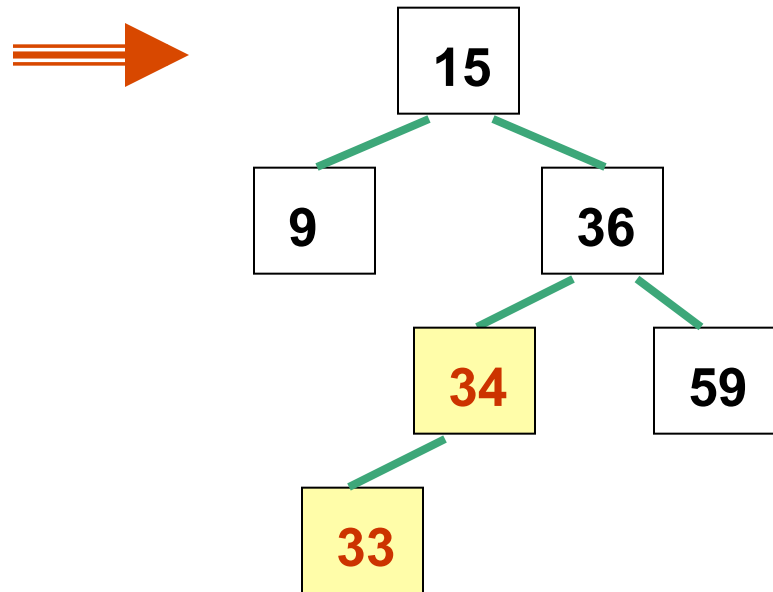
4: Remove 38



5: Remove 26



6: Remove 31



BST Operations: Remove Minimum

- Recall that *leftmost node* contains the minimum element
- Three cases:

1)root has no left child (so, root is minimum)

- its right child becomes the root

2)leftmost node is a leaf

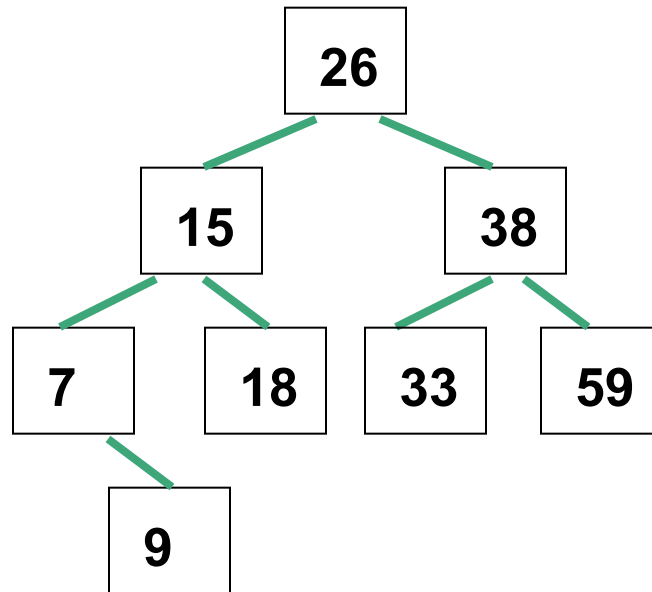
- set its parent's left child to null

3)leftmost node is internal

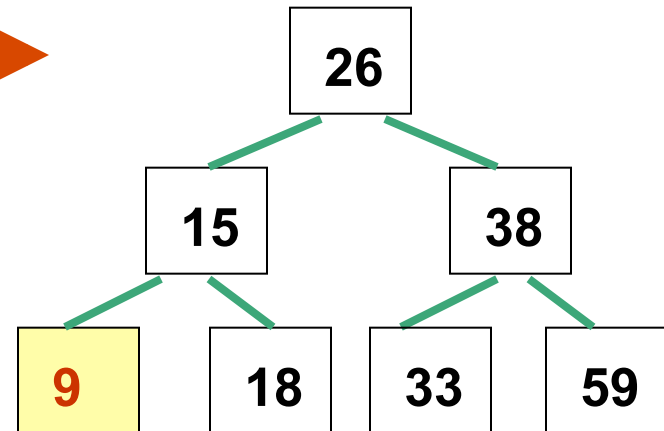
- the right child of the node to be removed becomes the parent's left child

Example: Removing Minimum BST Element

1: Initial tree

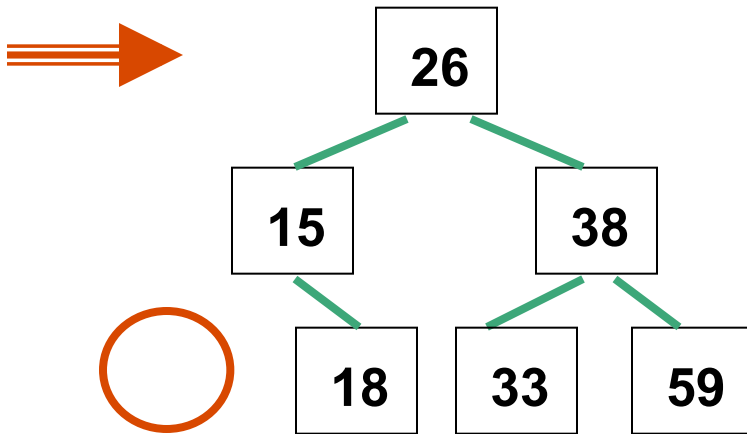


2: Remove minimum



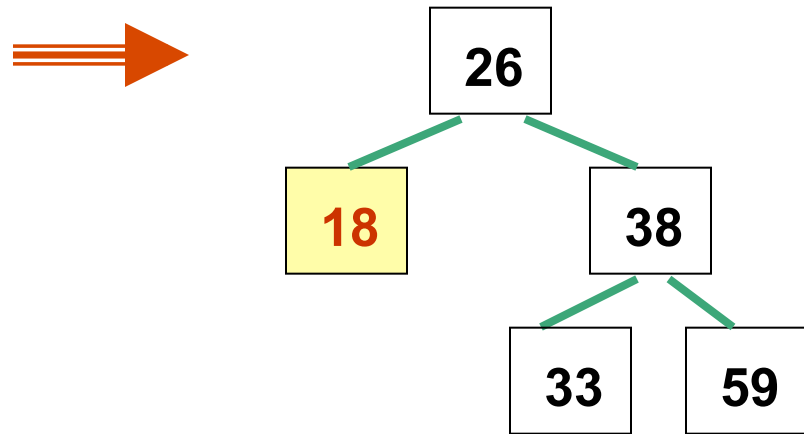
Case 3: internal node removed

3: Remove minimum



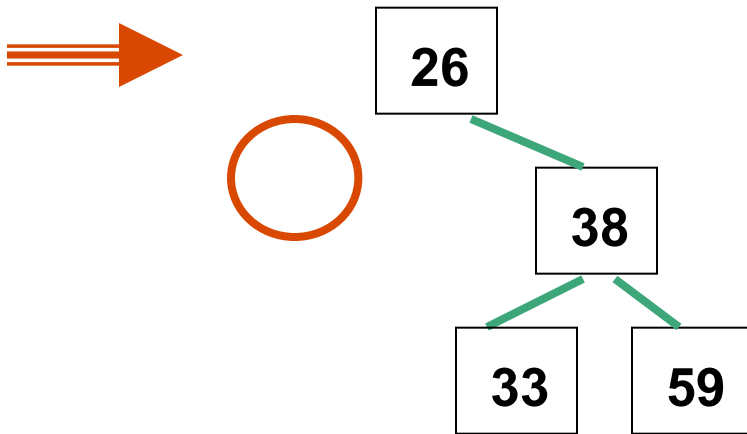
Case **2**: leaf node removed

4: Remove minimum



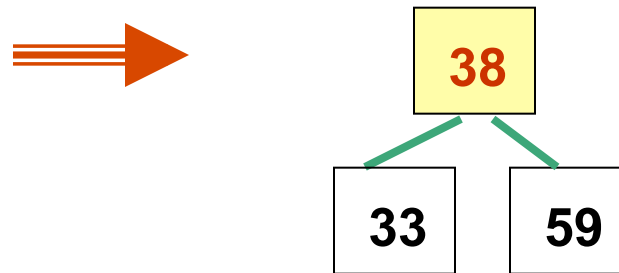
Case **3**: internal node removed

5: Remove minimum



Case **2**: leaf node removed

6: Remove minimum

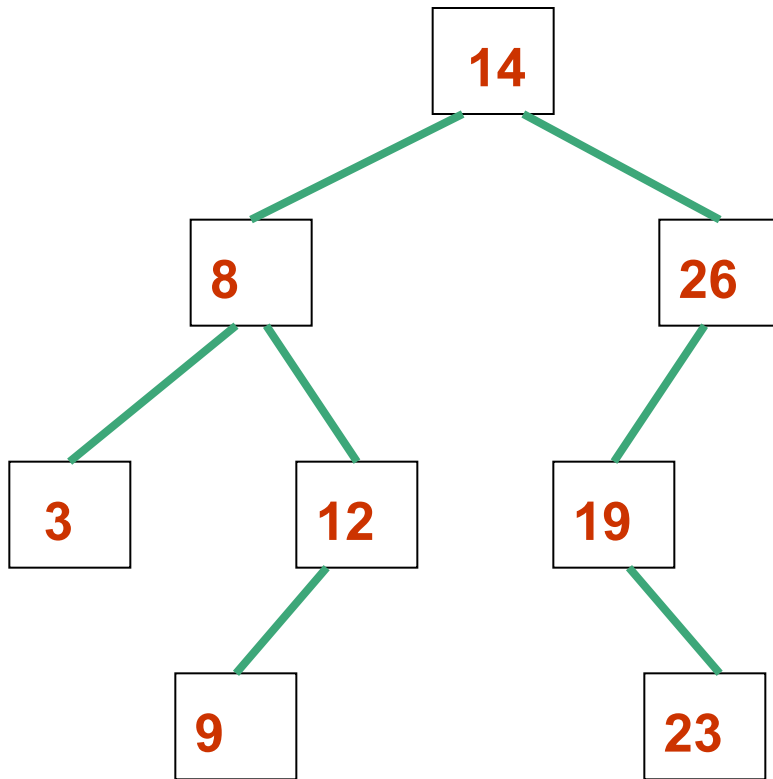


Case **1**: root node removed

Binary Search Tree Traversals

- Consider the traversals of a binary search tree: preorder, inorder, postorder, level-order
- Try the traversals on the example on the next page
 - Is there anything special about the *order of the data* in the BST, for each traversal?
- *Question*: what if we wanted to visit the nodes in *descending* order?

Binary Search Tree Traversals



Try these traversals:

- preorder
- inorder
- postorder
- level-order

Binary Search Tree ADT

- A BST is just a binary tree with the ordering property imposed on all nodes in the tree
- So, we can define the **BinarySearchTreeADT** interface as an *extension* of the **BinaryTreeADT** interface


```
public interface BinarySearchTreeADT<T> extends
                                   BinaryTreeADT<T> {
    public void addElement (T element);

    public T removeElement (T targetElement);

    public void removeAllOccurrences (T targetElement);

    public T removeMin( );

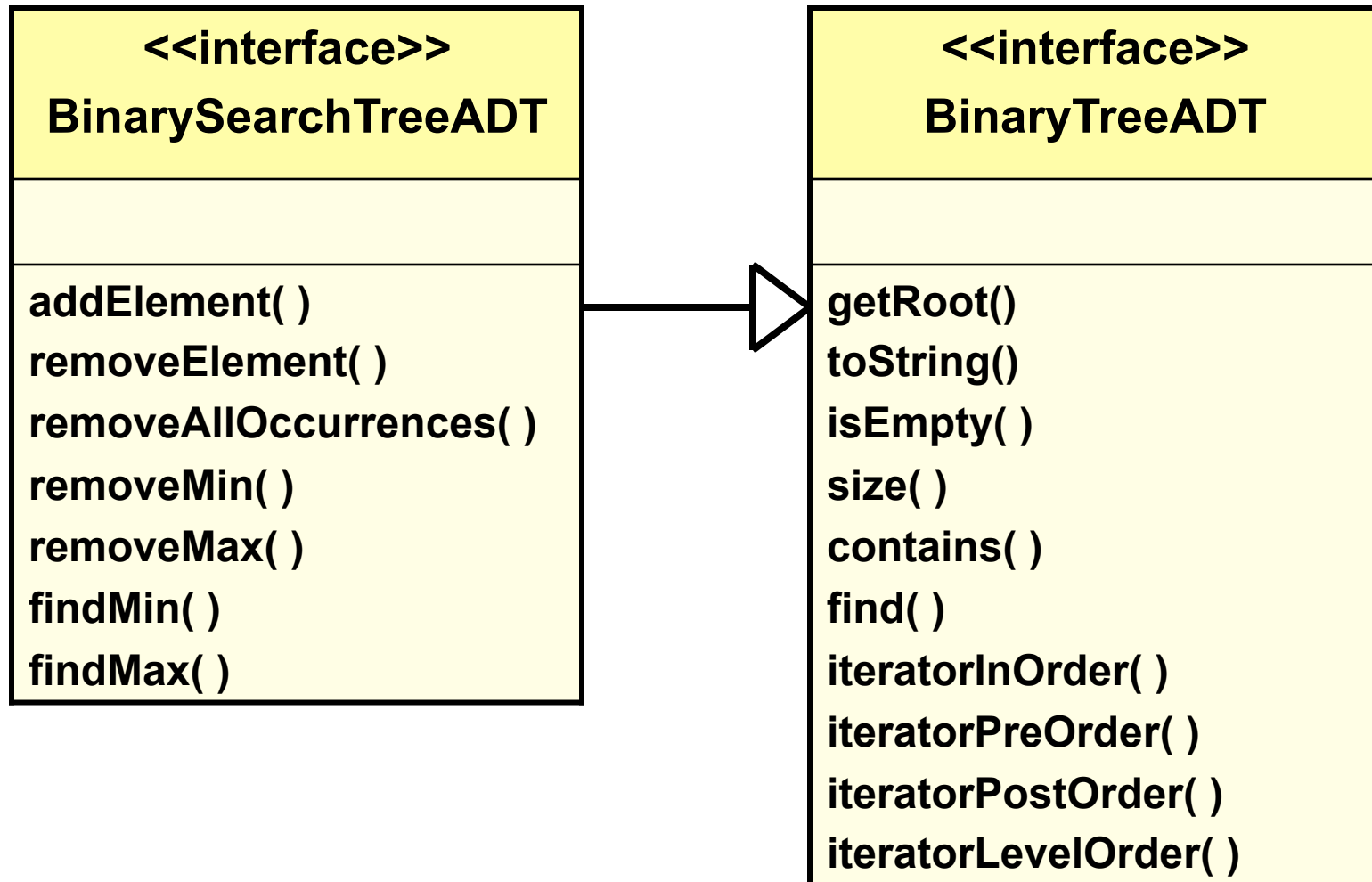
    public T removeMax( );

    public T findMin( );

    public T findMax( );
}
```

The
BinarySearchTreeADT
interface

UML Description of BinarySearchTreeADT



Implementing BSTs using Links

- See *LinkedBinarySearchTree.java*
 - Constructors: use `super()`
 - **addElement** method
 - *(does not implement our recursive algorithm of p.12; also, allows duplicates)*
 - note the use of **Comparable**: so that we can use `compareTo` method to know where to add the new node
 - **removeMin** method
 - essentially implements our algorithm of p. 18

Implementing BSTs using Links

- The special thing about a Binary Search Tree is that **finding a specific element is efficient!**
 - So, **LinkedBinarySearchTree** has a **find** method that **overrides** the **find** method of the parent class **LinkedBinaryTree**
 - It only has to search the appropriate side of the tree
 - It uses a recursive helper method **findAgain**
 - Note that it does not have a **contains** method that overrides the **contains** of **LinkedBinaryTree** – why not?
 - It doesn't need to, because **contains** just calls **find**

Using Binary Search Trees: Implementing Ordered Lists

- A BST can be used to provide *efficient* implementations of other collections!
- We will examine an implementation of an **Ordered List ADT** as a **binary search tree**
- Our implementation is called **BinarySearchTreeList.java**
(naming convention same as before: this is a BST implementation of a List)

Using BST to Implement Ordered List

- BinarySearchTreeList *implements* OrderedListADT
 - Which extends ListADT
 - So it also implements ListADT
 - So, what operations do we need to *implement*?
 - add
 - removeFirst, removeLast, remove, first, last, contains, isEmpty, size, iterator, toString
 - But, for which operations do we actually need to write code? ...

Using BST to Implement Ordered List

- **BinarySearchTreeList** *extends* our binary search tree class
LinkedBinarySearchTree
 - Which extends **LinkedBinaryTree**
 - So, what operations have we *inherited* ?
 - addElement, removeElement, removeMin, removeMax, findMin, findMax, find
 - getRoot, isEmpty, size, contains, find, toString, iteratorInOrder, iteratorPreOrder, iteratorPostOrder, iteratorLevelOrder

Discussion

- First, let's consider some of the methods of the **List ADT** that we do *not* need to write code for:
 - **contains** method: we can just *use* the one from the **LinkedBinaryTree** class
 - What about the methods
 - **isEmpty**
 - **size**
 - **toString**

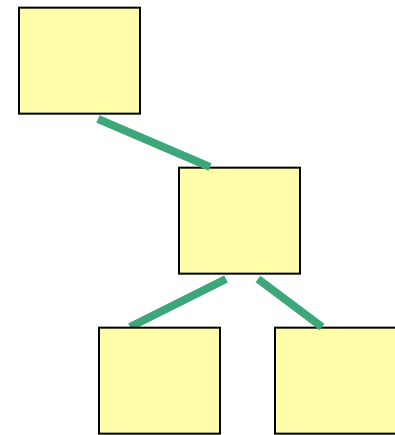
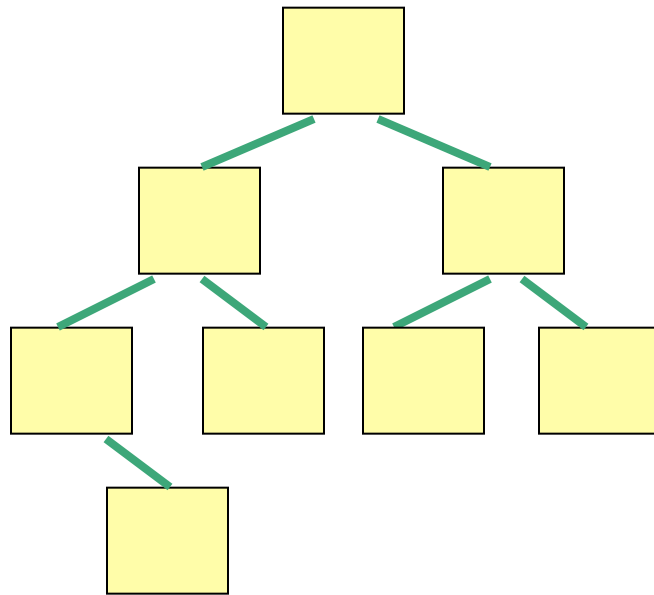
Discussion

- To implement the following methods of the **OrderedListADT** , we can *call* the appropriate methods of the **LinkedBinarySearchTree** class (*fill in the missing ones*)
 - **add** *call* **addElement**
 - **removeFirst** **removeMin**
 - **removeLast**
 - **remove**
 - **first**
 - **last**
 - **iterator**

Balanced Trees

- *Our definition*: a **balanced tree** has the property that, for any node in the tree, the height of its left and right subtrees can *differ by at most 1*
 - Note that conventionally the height of an empty subtree is **-1**

Balanced Trees



Which of these trees is a balanced tree?

Analysis of BST Implementation

- We will now compare the **linked list** implementation of an ordered list with its **BST** implementation, making the following important assumptions:
 - The BST is a ***balanced*** tree
 - The maximum level of any node is **$\log_2(n)$** , where **n** is the number of elements stored in the tree

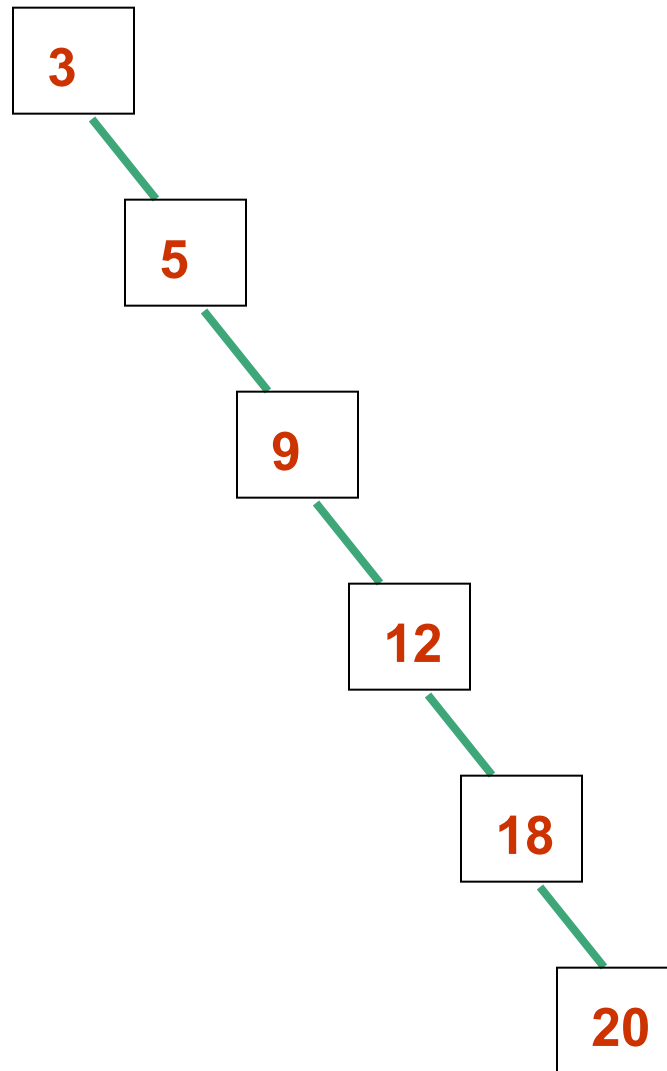
Analysis of Ordered List Implementations: Linked List vs. Balanced BST

Operation	LinkedList	BinarySearchTreeList
removeFirst	$O(1)$	$O(\log_2 n)$
removeLast	$O(n)$	$O(\log_2 n)$
remove	$O(n)$	$O(\log_2 n)$ *but may cause tree to become unbalanced
first	$O(1)$	$O(\log_2 n)$
last	$O(n)$	$O(\log_2 n)$
contains	$O(n)$	$O(\log_2 n)$
isEmpty	$O(1)$	$O(1)$
size	$O(1)$	$O(1)$
add	$O(n)$	$O(\log_2 n)$ *

Discussion

- Why is our balance assumption so important?
 - Look at what happens if we insert the following numbers in this order without rebalancing the tree:

3 5 9 12 18 20



Degenerate Binary Trees

- The resulting tree is called a *degenerate* binary tree
 - Note that it looks more like a linked list than a tree!
 - But it is actually less efficient than a linked list (*Why?*)

Degenerate Binary Trees

- Degenerate BSTs are far less efficient than balanced BSTs
 - Consider the worst case time complexity for the **add** operation:
 - **$O(n)$** for degenerate tree
 - **$O(\log_2 n)$** for balanced tree

Balancing Binary Trees

- There are many approaches to balancing binary trees
 - But they will not be discussed in this course ...