

Accessing fields in records

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- Compiler tracks both the types and the relative addresses of the fields of a record.
- Keeping this info. in symbol table entries for the field names → routine for looking up names can also be used for field names.
- $t \rightarrow$ pointer to the symbol table entry for a record type
 $\text{record}(t) \rightarrow$ returned as T.type.
- $p \uparrow.\text{info} + 1 \left\{ \begin{array}{l} p \rightarrow \text{pointer to a record with a field name 'info'} \\ \text{type}(p) \rightarrow \text{pointer}(\text{record}(t)) \\ \text{type}(p \uparrow) \rightarrow \text{record}(t) \end{array} \right.$
- 'info' is looked up in the symbol table pointed to by t .

Boolean Expressions

- Used as conditional expressions in statements that alter the flow of control.
- operators - 'and', 'or', 'not' applied to boolean variables or relational expressions.
- $E \rightarrow E \text{ or } E \mid E \text{ and } E \mid \text{not } E \mid (E) \mid \text{id rel op id} \mid \text{true} \mid \text{false}$
or, and - left associative precedence: ~~!~~, and, or.

Two methods for translating boolean expressions

(1) Numerical Representation

$a \text{ or } b \text{ and not } c$

$t_1 = \text{not } c$

$t_2 = b \text{ and } t_1$

$t_3 = a \text{ or } t_2$

$a < b \iff \begin{cases} \text{if } a < b \text{ then } 1 \\ \text{else } 0 \end{cases}$

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100: if $a < b$ goto 103

101: $t := 0$

102: goto 104

103: $t := 1$

104:

$E \rightarrow E_1 \text{ or } E_2 \quad \{ E.\text{place} := \text{newtemp};$
 $\text{emit}(E.\text{place} := 'E_1.\text{place 'or' } E_2.\text{place}')$

$E \rightarrow E_1 \text{ and } E_2 \quad \{ E.\text{place} := \text{newtemp};$
 $\text{emit}(E.\text{place} := 'E_1.\text{place 'and' } E_2.\text{place}')$

$E \rightarrow \text{not } E_1 \quad \{ E.\text{place} := \text{newtemp};$
 $\text{emit}(E.\text{place} := 'not' E_1.\text{place}')$

$E \rightarrow (E_1) \quad \{ E.\text{place} := E_1.\text{place} \}$

$E \rightarrow id_1 \text{ relop } id_2 \quad \{ E.\text{place} := \text{newtemp};$
 $\text{emit}('if' id_1.\text{place relop.op } id_2.\text{place}$
 $\text{'goto' nextstat+3});$
 $\text{emit}(E.\text{place} := '0');$
 $\text{emit}('goto' nextstat+2);$
 $\text{emit}(E.\text{place} := '1'); \}$

$E \rightarrow \text{true} \quad \{ E.\text{place} := \text{newtemp};$
 $\text{emit}(E.\text{place} := '1'); \}$

$E \rightarrow \text{false} \quad \{ E.\text{place} := \text{newtemp};$
 $\text{emit}(E.\text{place} := '0'); \}$

100: if $a < b$ goto 103

101: $t_1 = 0$

102: goto 104

103: $t_1 = 1$

104: if $c < d$ goto 107

105: $t_2 = 0$

106: goto 108

107: $t_2 = 1$

108: if $e < f$ goto 111

109: $t_3 = 0$

110: goto 112

111: $t_3 = 1$

112: $t_4 = t_2$ and t_3

113: $t_5 = t_1$ or t_4

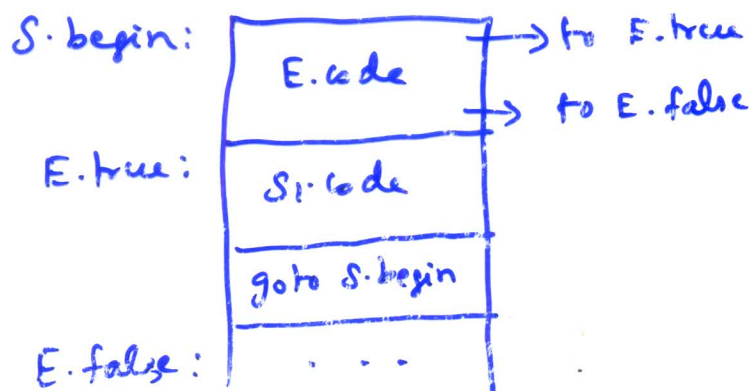
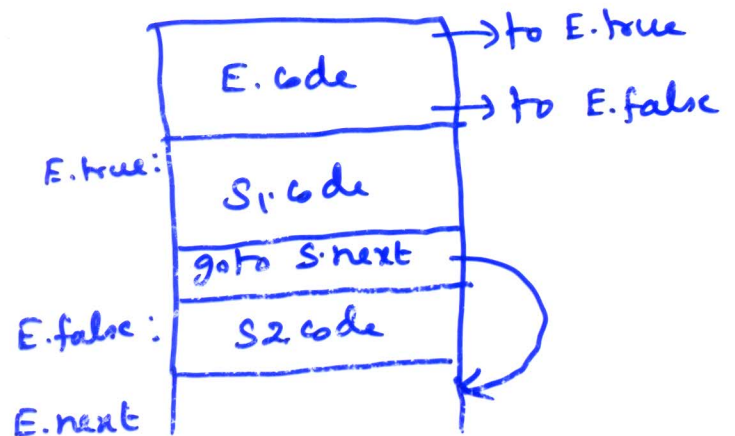
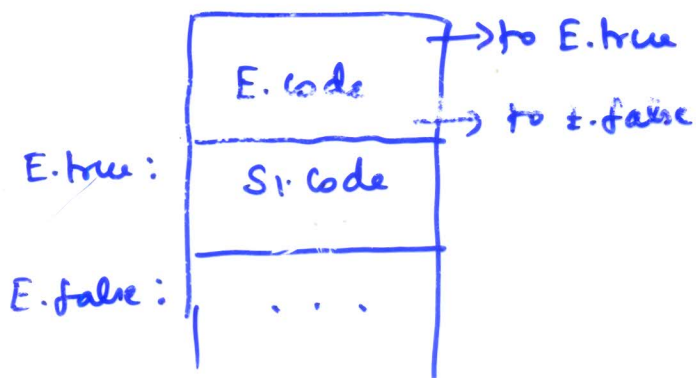
$\Rightarrow a < b$ or $c < d$ and $e < f$

'Short-circuiting'

t_1 need not be stored as the control flow through 101 or 103 determines t_1

Flow of Control Statements

$S \rightarrow$ if E then S_1 | if E then S_1 else S_2 |
while E do S_1



$S \rightarrow \text{if } E \text{ then } S_1$

$E.\text{true} = \text{newlabel};$

$E.\text{false} = S.\text{next};$

$S_1.\text{next} = S.\text{next}$

$S.\text{code} = E.\text{code} \parallel \text{gen}('E.\text{true}':') \parallel S_1.\text{code}$

$S \rightarrow \text{if } E \text{ then } S_1 \text{ else } S_2$

$E.\text{true} = \text{newlabel};$

$E.\text{false} = \text{newlabel};$

$S_1.\text{next} = S.\text{next};$

$S_2.\text{next} = S.\text{next};$

$S.\text{code} = E.\text{code} \parallel \text{gen}('E.\text{true}':') \parallel$

$S_1.\text{code} \parallel \text{gen}('goto' S.\text{next}) \parallel$

$\text{gen}('E.\text{false}':') \parallel S_2.\text{code}$

$S \rightarrow \text{while } E \text{ do } S_1$

$S.\text{begin} = \text{newlabel};$

$E.\text{true} = \text{newlabel};$

$E.\text{false} = S.\text{next};$

$S_1.\text{next} = S.\text{begin};$

$S.\text{code} = \text{gen}('S.\text{begin}':') \parallel E.\text{code} \parallel$

$\text{gen}('E.\text{true}':') \parallel S_1.\text{code} \parallel$

$\text{gen}('goto' S.\text{begin})$

Example:

$a < b \text{ or } c < d \text{ and } e < f$

$L_{\text{true}} / L_{\text{false}} \rightarrow \text{true and false exits for the entire expression}$

if $a < b$ goto L_{true}
goto L_1

L_2 : if $e < f$ goto L_{true}
goto L_{false}

L_1 : if $c < d$ goto L_2
goto L_{false}

$E \rightarrow E_1 \text{ or } E_2$

$E_1.\text{true} = E.\text{true};$
 $E_1.\text{false} = \text{newlabel};$
 $E_2.\text{true} = E.\text{true};$
 $E_2.\text{false} = E.\text{false};$
 $E.\text{code} = E_1.\text{code} \parallel \text{gen}(E_1.\text{false} ':') \parallel$
 $E_2.\text{code}$

$E \rightarrow E_1 \text{ and } E_2$

$E_1.\text{true} = \text{newlabel};$
 $E_1.\text{false} = E.\text{false}$
 $E_2.\text{true} = E.\text{true}$
 $E_2.\text{false} = E.\text{false}$
 $E.\text{code} = E_1.\text{code} \parallel \text{gen}(E_1.\text{true} ':') \parallel E_2.\text{code}$

$E \rightarrow \text{not } E_1$

$E_1.\text{true} = E.\text{false};$
 $E_1.\text{false} = E.\text{true};$
 $E.\text{code} = E_1.\text{code};$

$E \rightarrow (E_1)$

$E_1.\text{true} = E.\text{true};$
 $E_1.\text{false} = E.\text{false};$
 $E.\text{code} = E_1.\text{code};$

$E \rightarrow \text{id}_1 \text{ relop } \text{id}_2$ $E.\text{code} = \text{gen}('if' \text{ id}_1.\text{place relop.op}$
 $\text{id}_2.\text{place 'goto' } E.\text{true}) \parallel$
 $\text{gen}('goto' E.\text{false})$

$E \rightarrow \text{true}$

$E.\text{code} = \text{gen}('goto' E.\text{true})$

$E \rightarrow \text{false}$

$E.\text{code} = \text{gen}('goto' E.\text{false})$

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while a < b do
    if c < d then
        x = y + z
    else
        x = y - z

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L1: if a < b goto L2
    goto Lnext

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L2: if c < d goto L3
    goto L4

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L3: t1 = y + z
    x = t1
    goto L1

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L4: t2 = y - z
    x = t2
    goto L1

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Lnext:

Mixed mode boolean expressions

$E \rightarrow E + E \mid E \text{ and } E \mid E \text{ relop } E \mid \text{id}$

$E.\text{type} = \text{arith}$

if $E_1.\text{type} = \text{arith}$ and $E_2.\text{type} = \text{arith}$ then

$E.\text{place} = \text{newtemp}$

$E.\text{code} = E_1.\text{code} \parallel E_2.\text{code} \parallel \text{gen}(E.\text{place} = 'E_1.\text{place} + E_2.\text{place}')$

end

else if $E_1.\text{type} = \text{arith}$ and $E_2.\text{type} = \text{bool}$ then