

# Minimum subset sum difference

Problem link - <https://practice.geeksforgeeks.org/problems/minimum-sum-partition3317/1>

## Flow

1. Problem statement
2. Similarity
3. Solve using previous concept

**Problem Statement** - Given an array `arr` of size `n` containing non-negative integers, the task is to divide it into two sets `S1` and `S2` such that the absolute difference between their sums is minimum and find the minimum difference

Example 1:

Input: `N = 4, arr[] = {1, 6, 11, 5}`

Output: 1

Explanation:

Subset1 = {1, 5, 6}, sum of Subset 1 = 12

Subset2 = {11}, sum of Subset2 = 11

Example 2:

Input: `N = 2, arr[] = {1, 4}`

Output: 3

Explanation:

Subset1 = {1}, sum of Subset1 = 1

Subset2 = {4}, sum of Subset2 = 4

Array - { } iska do paration kar deta hai partition 1 and partition 2 and let say partition 1 ka sum `s1` hai and partition 2 ka sum `s2`.

To , we have to find the minimum difference between both the sum `s1` and `s2` ,

That is ( $s1 - s2 = \min$ ).

## Similarity

This problem is similar to the equal sum partition problem. Where we find  
The  $s1 - s2 = 0$ , or  $s1 == s2$  and one more thing this equal sum partition question  
based on the subset sum problem. And here in this question we have to find the  
 $s1 - s2 == \min$

### How to approach this question

Here, we don't know the exact value of  $s1$  and  $s2$ , but we know that  $s1$  and  $s2$  lie in  
some range that is  $s1, s2 (0, \text{range})$  where  $\text{range} = \text{sum of all the numbers}$ .

Let's suppose array -  $\{1, 2, 7\}$ ;

Now we have a range from 0 to 10, can all the numbers between 0 to 10 be  
candidates for  $s1$  and  $s2$ ? Let's find out 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10.

Can 0 is candidate Yes, because empty subset is  $\{\}$  can be value of  $s1$

Can 1 is a candidate, Yes because  $\{1\}$  one of the subsets from the given array is 1.

Can 4 is a candidate, No because no subset can form with sum 4.

Can 9 is a candidate, yes because  $\{2, 7\}$  can form with sum 9.

So we filter some values from the range. Now we got 1, 2, 3, 7, 8, 9, 10. These  
numbers are candidates for  $s1$  and  $s2$ .

Ans one more thing here we consider  $s1$  as min and  $s2$ , so the range for  $s1$  is half  
the range we got **1, 2, 3** | 7, 8, 9, 10

Now we know that  $s1 + s2 = \text{Range}$

If we find the value of  $s1$ ,  $s2 = \text{range} - s1$ . And in the question asking for  $s1 - s2 = \min$   
, so, putting the value of  $s2$  in this  $(\text{range} - s1 - s1)$  that becomes  $\min = \text{range} - 2s1$ .

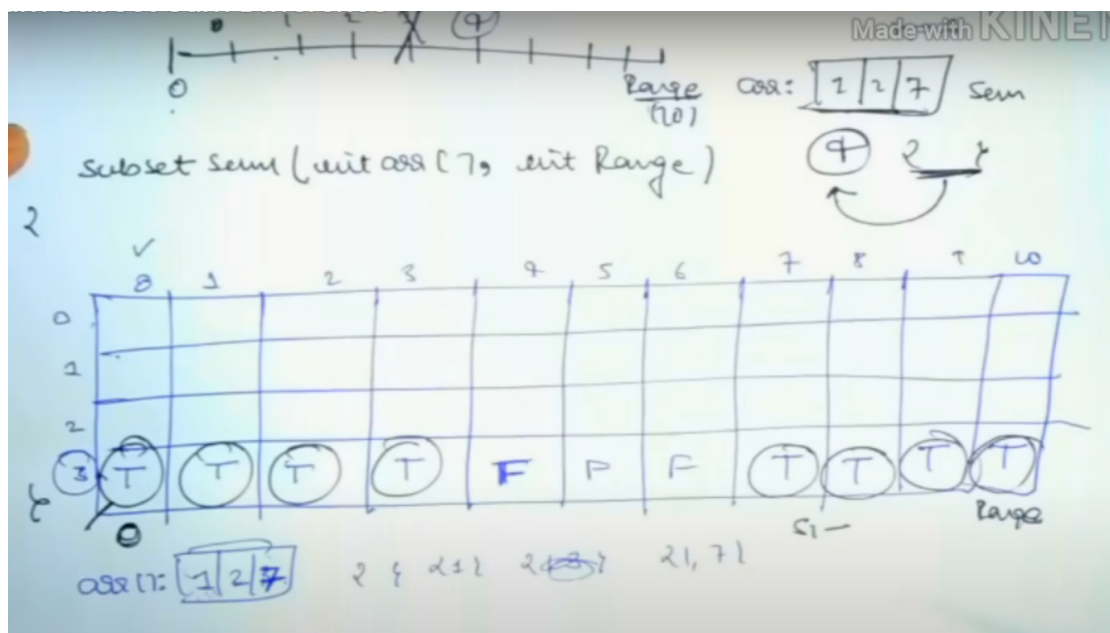
Now the value of  $s1 = 1, 2, 3$  and  $\text{min} = \text{range} - 2s1$  now if you put  $s1 = 1$  we get  $\text{min} = 10 - 1 = 9$ ,  $s1 = 2$  we get  $10 - 4 = 6$ ,  $s1 = 3$  we get  $10 - 6 = 4$ .  
So the minimum value of  $s1 - s2 = 4$ .

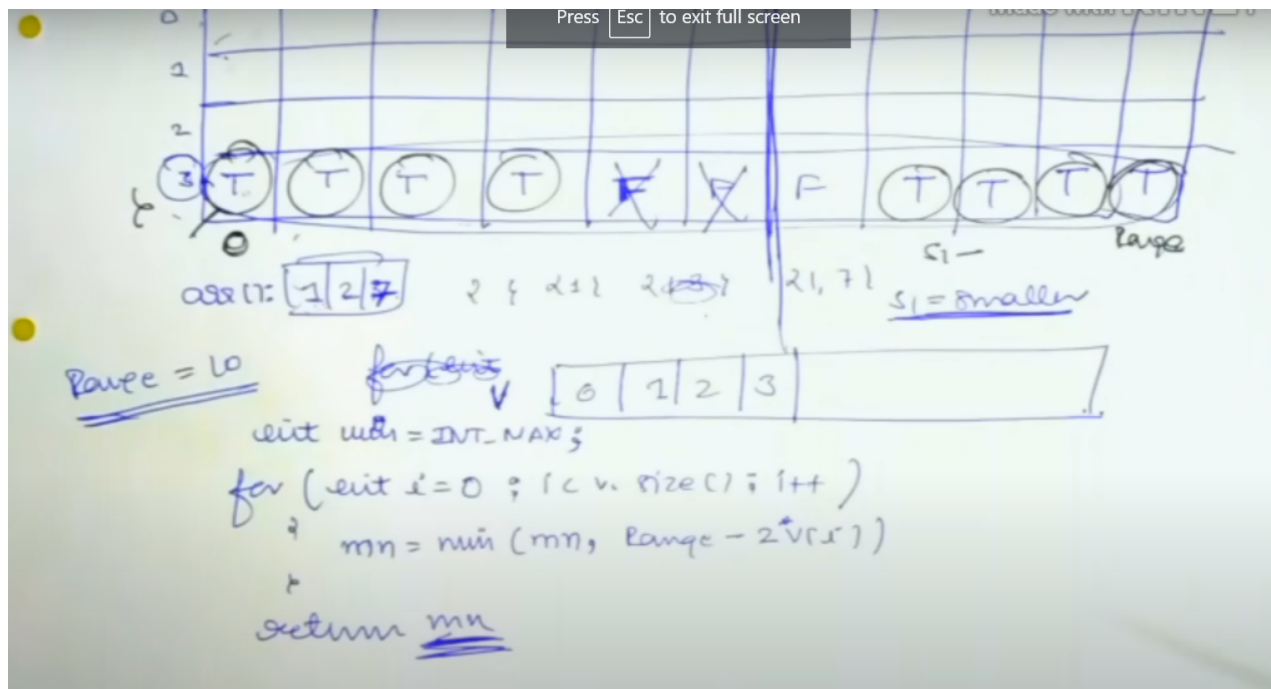
## // Code explanation

If you remember that the subset sum the last row in 2d, the boolean vector denotes that if the subset exists or not with the sum if true then exists else false.

So we check the last row of that matrix and check if true then put in a vector of int.

And for calculating minimum we already derive the formula that is  $\text{range} - 2s1$ .






```
public:
    int minDifference(int nums[], int n) {
        //finding the range or sum of all the numbers
        long long sum = 0;
        for(int i = 0; i < n; i++){
            sum += nums[i];
        }
        vector<int> ans;
        //creating a 2d vector so that we can filter the candidate
        vector<vector<bool>> dp(n+1, vector<bool>(sum/2+1, false));
        for(int i = 0; i <= n; i++){
            dp[i][0] = true;
        }
        for(int i = 1; i <= n; i++){
            for(int j = 1; j <= sum/2; j++){
                if(nums[i-1] <= j){
                    dp[i][j] = dp[i-1][j-nums[i-1]] || dp[i-1][j];
                }else{
                    dp[i][j] = dp[i-1][j];
                }
            }
        }
    }
}
```

```

        }
    }
}
//here filtering the candidate for s1
for(int i =0; i<=sum/2; i++){
    if(dp[n][i] == true){
        ans.push_back(i);
    }
}
//here finding the min value
int mn = INT_MAX;
for(int i =0; i<ans.size(); i++){
    int a = sum - (2 * ans[i]);
    mn = min(mn , a);
}
return mn;
}
};

```

If you still have donuts we can check out the videos -

 10 Minimum Subset Sum Difference