

Homework #2

(Due: Apr 16)

Task 1. [150 Points] Average-Case Analysis of Median-of-3 Quickselect

Consider the selection algorithm given in Figure 1.

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MEDIAN-OF-3-QUICKSELECT(  $A[q : r]$ ,  $k$  )

Input: An array of distinct elements, and an integer  $k \in [1, r - q + 1]$ .
Output: An element  $x$  of  $A[q : r]$  such that  $\text{rank}(x, A[q, r]) = k$ .

1.  $n \leftarrow r - q + 1$ 
2. if  $n < 1$  then return NIL
3. elif  $n < 4$  then
4.   sort  $A[q : r]$  in increasing order of value
5.   return  $A[q + k - 1]$ 
6. else
7.    $x \leftarrow$  median of  $A[q]$ ,  $A[q + 1]$  and  $A[q + 2]$ 
8.   rearrange the numbers of  $A[q : r]$  such that
       •  $A[t] = x$  for some  $t \in [q, r]$ ,
       •  $A[i] < x$  for each  $i \in [q, t - 1]$ ,
       •  $A[i] > x$  for each  $i \in [t + 1, r]$ ,
9.   if  $k = t - q + 1$  then return  $A[t]$ 
10.  else if  $k < t - q + 1$  then return MEDIAN-OF-3-QUICKSELECT(  $A[q : t - 1]$ ,  $k$  )
11.  else return MEDIAN-OF-3-QUICKSELECT(  $A[t + 1 : r]$ ,  $k$  )

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Figure 1: [Task 1] Return the k^{th} smallest number in the input array $A[q : r]$, where $k \in [1, r - q + 1]$.

Given an input array A of size n , in this task we will analyze the average number of element comparisons (i.e., comparisons between two numbers of A) performed by this algorithm over all $n!$ possible permutations of the numbers in A and all n possible values of k . We will assume that the partitioning algorithm (in line 8) is *stable*, i.e., if two numbers y and z end up in the same partition and y appears before z in the input, then y must also appear before z in the resulting partition.

- (a) [15 Points] Show how to implement steps 7 and 8 of Figure 1 to get a stable partitioning of the numbers in $A[1 : n]$ using only $n - \frac{1}{3}$ element comparisons on average, where the average is taken over all $n!$ possible permutations of the input numbers.
- (b) [25 Points] Let t_n be the average number of element comparisons performed by the algorithm given in Figure 1 on an input array A of size n , where $n \geq 0$ and the average

is taken over all $n!$ possible permutations of the numbers in A and all n possible values of integer k . Argue that

$$t_n = \begin{cases} 0 & \text{if } n < 2, \\ 1 & \text{if } n = 2, \\ 3 & \text{if } n = 3, \\ n - \frac{1}{3} + \frac{6}{n^2(n-1)(n-2)} \sum_{k=1}^n (k-1)(n-k)((k-1)t_{k-1} + (n-k)t_{n-k}) & \text{otherwise.} \end{cases}$$

(c) [**50 Points**] Let $T(z)$ be a generating function for t_n :

$$T(z) = t_0 + t_1 z + t_2 z^2 + \dots + t_n z^n + \dots \dots$$

Show that

$$z(1-z)^2 T^{(4)}(z) + 3(1-z)^2 T^{(3)}(z) - 12z T^{(2)}(z) - 12 T^{(1)}(z) = \frac{8(z^2 + 8z + 6)}{(1-z)^4},$$

where, $T^{(k)}(z) = \frac{d^k}{dz^k} (T(z))$ for integer $k > 0$.

(d) [**25 Points**] Use the result from part (c) to show that

$$z(1-z)^2 T^{(3)}(z) + 2(1-z) T^{(2)}(z) + 2(1-6z) T^{(1)}(z) = \frac{8(z^2 + 3z + 1)}{(1-z)^3} - 4.$$

(e) [**25 Points**] Use the result from part (d) to show that

$$n^2(n-2)t_n = \frac{1}{2}(n+1) [(n-1)(n-6)t_{n-1} + n(n+1)t_{n+1}] - 10n(n-1) - 4.$$

(f) [**10 Points**] What bound on t_n can you get from part (e)?

<p>COMP-S(n) {Integer $n \geq 0$}</p> <ol style="list-style-type: none"> 1. if $n = 0$ then return 1 2. else 3. $s \leftarrow \text{COMP-S}(n - 1)$ 4. $t \leftarrow \text{COMP-T}(n - 1)$ 5. return $s + 2t$ 	<p>COMP-T(n) {Integer $n \geq 0$}</p> <ol style="list-style-type: none"> 1. if $n = 0$ then return 0 2. else 3. $s \leftarrow \text{COMP-S}(n)$ 4. $t \leftarrow \text{COMP-T}(n - 1)$ 5. return $3s + t$
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Figure 2: [Task 2] Two mutually recursive functions.

Task 2. [90 Points] Two Mutually Recursive Functions

Figure 2 shows two mutually recursive functions COMP-S and COMP-T. Both functions accept a nonnegative integer as the sole input.

Now answer the following questions.

- (a) [45 Points] Use generating functions to find the values returned by COMP-S(n) and COMP-T(n).
- (b) [45 Points] Use generating functions to find the running times of COMP-S(n) and COMP-T(n).

<p>RANDOM-FLIPS(A, p, q)</p> <ol style="list-style-type: none"> 1. $n \leftarrow q - p + 1$ 2. if $n \leq 4$ then return 3. else 4. for $i \leftarrow 1$ to $n^{\log_2(\frac{\sqrt{61}-1}{6})}$ do 5. $r \leftarrow \text{RANDOM}(p, q)$ {$r \leftarrow$ an integer chosen uniformly at random from $[p, q]$} 6. $A[r] \leftarrow 1 - A[r]$ 7. $t \leftarrow \text{RANDOM}(1, 10)$ 8. if t is divisible by 2 then RANDOM-FLIPS(A, p, q) 9. elif t is divisible by 3 then RANDOM-FLIPS($A, p, p + \frac{n}{4} - 1$) 10. elif t is divisible by 5 then RANDOM-FLIPS($A, q - \frac{n}{4} + 1, q$) 11. elif t is divisible by 7 then RANDOM-FLIPS($A, p + \frac{n}{4}, q - \frac{n}{4}$) 12. else return 	<p>{Input is a bit array $A[p : q]$, where $p, q \in \mathbb{Z}^+$, $q > p$, and $n = q - p + 1 = 2^k$ for some integer $k \geq 0$.}</p>
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Figure 3: [Task 3] Recursive random bit flips.

Task 3. [60 Points] A Recursive Random Bit Flipping Algorithm

The recursive algorithm given in Figure 3 flips the entries of a bit array at random. Write down a recurrence relation describing the expected running time of the algorithm and solve it.