

CS2040 Lab 6

Priority Queue

Lab 6 – Priority Queue

- Java provides a PriorityQueue class, which uses a binary heap
 - To create a new priority queue:
 - `PriorityQueue<Integer> pq = new PriorityQueue<Integer>();`
- Is a min heap by default; use `Comparator.reverseOrder()` to change it to a max heap (for elements that implement Comparable):
 - `PriorityQueue<Integer> pq = new PriorityQueue<Integer>(Comparator.reverseOrder());`
- PriorityQueue enforces some ordering on its elements. As a result, any object stored must either implement Comparable, or the constructor must be given a Comparator

Lab 6 – PriorityQueue

Method name	Description	Time
.add(YourClass element)	Adds <i>element</i> to the PriorityQueue	$O(\log n)$
.clear()	Empties the PriorityQueue	$O(n)$
.contains(Object o)	Checks if <i>o</i> is in the PriorityQueue, based off the object's equals() method	$O(n)$
.peek()	Returns the top element of the PriorityQueue	$O(1)$
.poll()	Removes and returns the top element of the PriorityQueue	$O(\log n)$
.remove(Object o)	Removes <i>o</i> if it is in the PriorityQueue, based off the object's equals() method	$O(n)$
.size()	Returns the number of elements in the list	$O(1)$

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PQ for thought

What if I want a PQ using binary heaps that can

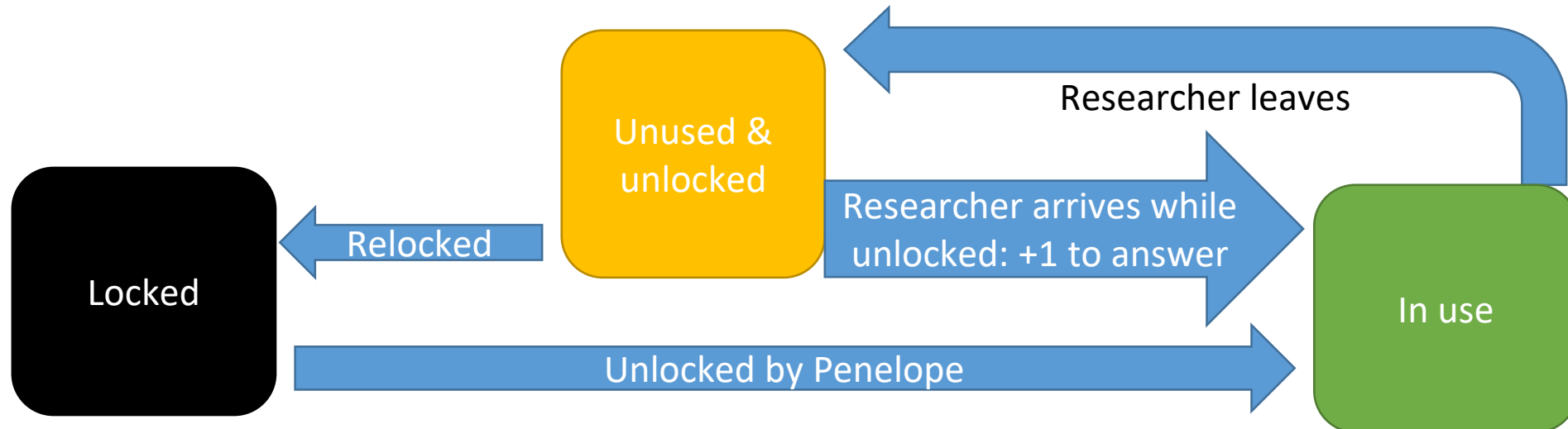
- Remove any element in $O(\log n)$?
- Merge with another pq in $O(n)$?
- Remove the top and bottom element in $O(\log n)$?
- Update the priority of an entry? (idea used later in semester)

(Not in syllabus) What if I want a PQ that can

- Merge with another pq in $O(\log n)$?
- Remove top and bottom element in $O(\log n)$, without using any HashMap / multiple data structures?

One-Day Assignment 5 – Assigning Workstations

- Workstations can be in one of three different states:
 - Locked (so Penelope has to unlock this workstation)
 - In use (so this workstation cannot be assigned to a different researcher yet)
 - Unused, but unlocked (this workstation can be assigned to a researcher, so Penelope does not need to unlock a new workstation)



One-Day Assignment 5 – Assigning Workstations

Legend - Locked						Legend - In use						Legend - Unused & unlocked					

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	
	Researcher 1																				
						Researcher 2															
														Researcher 3							

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
	Researcher 1					Researcher 2								Researcher 3						

The above is an illustration of Sample Input 1

The first version requires Penelope to unlock a new workstation every time a new researcher enters

The second version requires Penelope to only unlock a workstation 1 time instead of 3 times, thereby saving **2** unlocks

Assigning Workstations – Sample Input 2

Assigning Workstations – Sample Input 2									Locked				In use				Unused & unlocked				
1 2 – First researcher, no choice but to unlock new WS	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
		Rschr 1																			
2 6 – First relocking WS (WS1) still in use at time 3, unlock new WS	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
		Rschr 1																			
			Researcher 2																		
3 9 – First relocking WS (WS1) is unused and unlocked at time 3, assign it	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
		Rschr 1		Researcher 3																	
			Researcher 2																		
15 6 – First relocking WS (WS2) is unused and unlocked at time 15, assign it	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
		Rschr 1		Researcher 3																	
			Researcher 2												Researcher 4						
17 7 – First relocking WS (WS1) is unused and unlocked at time 17, assign it	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
		Rschr 1		Researcher 3													Researcher 5				
			Researcher 2												Researcher 4						