Lab report # 9: CPU Performance Monitoring Email Address: 17-0024@student.gutech.edu.om

1. on page 2: By looking at the last column id (idle time), which system is under heavy load?

To know which system has heavy load, we need to look at four factors:

- a. id (idle time): We can see from the illustration given that most of the instances are with idle time of *zero*. In other words, the second system has so much less idle than the first system, which indicates that the second system is executing the CPU more than the first system.
- b. us (CPU utilization): Executing the processes of the second system led to more *US* than the first system. In other words, the CPU utilization in the second system is higher.
- c. in (interrupts): There are exceedingly more interrupts in the second system than the first system.
- d. cs (context switches): System two has a low number of context switches in its cases compared to the system one.

From the four factors above, we can conclude that the <u>second system</u> is the one under heavy load.

2. on page 3: A case study analysis (Sustained CPU Utilization).

Sustained CPU utilization means the CPU is completely utilized (highest CPU usage).

To analyze this case study, we need to consider multiple factors:

- a. r (Run Queue): We can notice that the average length of the run queue is about 2 threads. In other words, an average of 2 threads wait for the CPU to execute them but the CPU is busy executing other threads. However, we can see that the last case, in the illustration given, has a run queue length of *zero*, which means the CPU has been able to execute them without leading to starvation. In general, the run queue length, in this case study system, is suitable and does not affect performance negatively. It is important to mention that there is only one run queue for each processor, and usually if the length of the run queue is less than 10 in a processor, then that is appropriate, or if the average length of run queue in most cases is not more than 3, then it is more efficient.
- b. b (Processes Blocked): We can see from the illustration given that there are zero blocked processes in all cases. In other words, these processes are either done with the I/O wait or do not need I/O. This leads to a higher CPU utilization.
- c. in (interrupts): This is the number of interrupts per second. We can notice that the system has a relatively high number of interrupts in most cases. This means that the system is under heavy load.
- d. cs (Context Switches): This is the number of context switches per second. The less context switches there are, the less heavy the system. In this case, the context

Lab report # 9: CPU Performance Monitoring Email Address: 17-0024@student.gutech.edu.om

switches are relatively low. Which means there are less interrupts or switching between processes, which means the CPU is executing the process in large chunks. We can also observe that, in each run, the context switches are incredibly low compared to the interrupts, which means better performance because fewer context switches mean better process utilization, or better CPU performance.

- e. us (CPU Utilization) and sy (kernel Utilization): We notice that the sum of the user processes (us) and the kernel/interrupts processes (sy) results in almost 100% in most cases and it is divided almost 70/30 between user and kernal. Which means that the CPU is fully utilized, giving time for both the kernel and the user processes.
- f. Id (CPU idle time): Most cases are zero. In other words, the CPU is fully utilized, better performance. However, this can affect computers with small RAM, but in general, the system is working properly.

3. on page 4: Overloaded Scheduler analysis

- a. Let us first compare the context switches with the interrupts: we can observe that, in each run, the interrupt is less than the context switch. In other words, the result of difference between the context switches and the interrupts lead CPU overload. To explain, since the interrupt are lower than the context switches, this means between interrupts, there is relatively a high number of context switches for each process, and the more switches there are, the lower the performance and efficiency. This leads to a decrease in the overall performance.
- b. us (CPU utilization of user processes): to calculate the total utilization of the user processes, we need to do the following:

Total user processes: 4 + 8 + 9 + 3 + 7 + 2 + 22 + 12 + 6 + 5 + 6 = 84

Then calculate the total utilization in percentage:

Total user processes * (number of runs/100)

Average user utilization = 84/11 = 7.63%

Conclusion:

The average is less than the total, which means the utilization of user processes is appropriate.

c. r (run queue): The average run time length is 2+2+3+4+5+4+6=26

The average run time length = 26/11 = 2.4

We can notice that the least run time is 0 and the highest is 6. The average 2.4 is between these two values. This means that it is relatively normal. However, for better performance, the length of the run queue is better to be less than 3 in each run and less than 10 threads per run.

Lab report # 9: CPU Performance Monitoring Email Address: 17-0024@student.gutech.edu.om

d. Sy (kernel and interrupt utilization): We can observe that the system kernel processes are taking more of the CPU than the user processes in most runs. In other words, the CPU is overloaded and trying to reschedule priorities.

4. on page 4&5: Practice section

Part 1:

Step1: Run vmstat on your machine with a 1 second delay between updates. Let it run for approximately 10 seconds. Notice the CPU utilization by looking at the (id) column Hint: To stop vmstat, CTRL + C

г	awa	ın@	rawan-\	/irtualB	ox:~\$ \	/mstat -	t 1 10	9										
ΡI	oc	S		mem	огу		SW	ар	i	0	-syst	tem			-срі	J		timestamp
	Г	Ь	swpd	free	buff	cache	si	so	bi	bo	in	cs	us	sy	id	wa	st	+04
(9	0	0	612740	59868	749876	0	0	9492	385	736	7905	46	15	37	2	0	2021-01-06 12:48:47
(9	0	0	612732	59876	749924	0	0	0	52	289	955	10	4	85	1	0	2021-01-06 12:48:48
1	l	0	0	612732	59876	749924	0	0	0	0	341	967	б	1	93	0	0	2021-01-06 12:48:49
1	L	0	0	624844	59876	749920	0	0	0	0	235	848	14	2	84	0	0	2021-01-06 12:48:50
(9	0	0	621560	59876	751080	0	0	0	0	564	2372	31	7	62	0	0	2021-01-06 12:48:51
(9	0	0	621560	59876	751080	0	0	0	0	697	2063	52	2	46	0	0	2021-01-06 12:48:52
1	L	0	0	609576	59876	750984	0	0	0	0	369	1102	37	7	56	0	0	2021-01-06 12:48:53
(9	0	0	609576	59876	750948	0	0	0	0	618	2400	7	5	88	0	0	2021-01-06 12:48:54
(9	0	0	609576	59876	750948	0	0	0	0	201	522	5	1	94	0	0	2021-01-06 12:48:55
1	L	0	0	604032	59888	750948	0	0	0	84	96	205	3	1	96	0	0	2021-01-06 12:48:56

Illustration 1.1: The system <u>vmstat</u> analysis.

Step2: Run ex00.py script (Ubuntu\$ python ex00.py). Then run vmstat in another terminal.

awa	an@ı	rawan-\	/irtualB	ox:~\$ \	/mstat -	t 1 10	9		_										
ГО	cs ·		mem	огу		SW	ар	i	0	-syst	em			cpu			timesta	mp	
Γ	Ь	swpd	free	buff	cache	si	so	bi	bo	in	CS	us	sy	id 1	wa	st		+04	
4	0	0	523788	60804	807764	0	0	6371	250	602	5302	45	11	43	1	0	2021-01-06 1	2:49:32	
1	0	0	531352	60836	809516			1520	0	651	1151	98	2	0	0	0	2021-01-06 1	2:49:33	
1	0	0	531352	60836	809516	0	0	0	0	580	592	99	1	0	0	0	2021-01-06 1	2:49:34	
1	0	0	531352	60844	809508	0	0	0	24	506	280	100	0	0	0	(2021-01-06	12:49:35	
1	0	0	532596	60844	809516	0	0	0	0	614	706	100	0	0	0	(2021-01-06	12:49:36	
1	0	0	532620	60844	809516	0	0	0	0	466	226	100	0	0	0	(2021-01-06	12:49:37	
1	0	0	532620	60844	809516	0	0	0	64	709	1269	100	0	0	0	(2021-01-06	12:49:38	
1	0	0	532620	60844	809516	0	0	0	0	489	215	100	0	0	0	(2021-01-06	12:49:39	
1	0	0	532620	60844	809516	0	0	0	0	660	821	99	1	0	0	0	2021-01-06 1	2:49:40	
1	0	0	533344	60844	809516	0	0	0	0	494	352	100	0	0	0	(2021-01-06	12:49:41	

Illustration 1.2: The system <u>vmstat</u> analysis after running ex00.py

Q1. What do you notice in this case?

The <u>r (run time queue)</u> length increases slightly indicating a process waiting.

The \underline{id} ($\underline{idle\ time\ of\ the\ CPU}$) has dropped to zero in most runs indicating that the CPU – from ex00.py code – \underline{pass} for one second each run and allowing it to be fully engaged every run because it is executing the code.

The <u>us (User process utilization)</u> resembles in this case the code in ex00.py and it engages the CPU fully 100% in most runs, while the kernel processes give up the CPU to allow the code to run.

Q2. Without opening ex00.py, is there anything you can deduce regarding the script It will be looping infinitely and engaging the CPU for the user process in the script.

Lab report # 9: CPU Performance Monitoring Email Address: 17-0024@student.gutech.edu.om

Q3. Have a look inside ex00.py

```
while True:
pass
```

Illustration 1.3: the script inside ex00.py

Q4. Why CPU utilization is around 25% (not 100%)? Hint: run mpstat

In a quad-core we can run 4 threads at once. In our script, we only need one thread, which means we are only using quarter of our entire CPU. This is the reason for most modern computers. However, in our case, the user programs are consuming much more.

Part2:

Step1: Open ex01.py and run it with a parameter less than 1000 (exp: Ubuntu\$ python ex01.py 900), and then more than 1000.

```
rawan@rawan-VirtualBox:~$ python3 ex01.py 900 factortal(1) factorial(2) factorial(3) factorial(4) factorial(5) factorial(6) factorial(6) factorial(7) factorial(8) factorial(9) factorial(9) factorial(10)
```

```
Factorial(895)
Factorial(896)
Factorial(897)
Factorial(898)
Factorial(899)
Factorial(900)
Wawangrawan-VirtualBox:~$ python3 ex01.py 1100
```

Illustration 2.1: Running ex01.py with a parameter of 900

```
rawan@rawan-virtualBox:~$ python3 ex01.py 1100 factorial(1) factorial(2) factorial(3) factorial(4) factorial(5) factorial(6) factorial(7) factorial(8)
```

```
factorial(997)

factorial(998)

fraceback (most recent call last):

File "ex01.py", line 14, in <module>
factorial(i)

File "ex01.py", line 7, in factorial

return n * factorial(n-1)

file "ex01.py", line 7, in factorial

return n * factorial(n-1)

File "ex01.py", line 7, in factorial

return n * factorial(n-1)

File "ex01.py", line 4, in factorial

File "ex01.py", line 4, in factorial

if n == 0:

RecursionError: maximum recursion depth exceeded in comparison
```

Illustration 2.2: Running ex01.py with a parameter of 1100

Q1. What is the problem?

"RecursionError: maximum recursion depth exceeded in comparison"

Lab report # 9: CPU Performance Monitoring Email Address: 17-0024@student.gutech.edu.om

Since the script in ex01.py is a recursive function of factorial. The recursion usually stores the current value or state in the system's stack. There is a limit to the system stack, if it is gets full, no more values can be stored. Therefore, factorial(998) got stored in the system stack, but any further recursion will cause the program to fail. When we ran ex01.py with the parameter 1100, we noticed that the program failed. This is because there were too many recursion calls that the system stack got full. There is an overflow.

Step2: Run ex01b.py with a large number (more than 1000). Use vmstat and mpstat to check the CPU behaviour.

```
rawan@rawan-VirtualBox:~$ python3 ex01b.py 1200
factorial(1)
factorial(2)
factorial(3)
factorial(4)
factorial(5)
factorial(6)
factorial(7)
factorial(8)
factorial(9)
factorial(10)
factorial(10)
factorial(11)
```

factorial(1195)
factorial(1196)
factorial(1197)
factorial(1198)
factorial(1199)

factorial(1200)
rawan@rawan-VirtualBox:~\$

Illustration 2.3: Running ex01b.py with a parameter of 1200

raw	an@	rawan-\	/irtualE	ox:~\$ \	/mstat 1	l											
рго	CS		mer	югу		swa	p	io-		-syst	tem			-срі	J		
Г	Ь	swpd	free	buff	cache	si	so	bi	bo	in	cs	us	sy	id	wa	st	
0	0	0	451484	63328	818540	0	0	177	15	343	400	64	1	35	0	0	
0	0	0	451224	63328	818568	0	0	Θ	0	535	2214	47	7	46	0	0	
0	0	0	451224	63328	818584	0	0	0	0	339	1322	36	2	62	0	0	
2	0	0	448696	63328	818584	0	0	0	0	153	436	11	1	88	0	0	
0	0	0	450216	63328	818584	0	0	0	0	421	8296	68	4	28	0	0	
0	0	0	450216	63328	818584	0	0	0	0	381	1960	13	9	78	0	0	
0	0	0	450216	63328	818584	0	0	Θ	0	117	326	4	1	95	0	0	
1	0	0	450216	63328	818600	0	0	0	0	352	1264	44	3	53	0	0	
1	0	0	450216	63328	818600	0	0	0	0	111	277	3	1	96	0	0	
0	0	0	450216	63328	818600	0	0	0	0	102	232	2	0	98	0	0	
0	0	0	450216	63328	818600	0	0	0	0	108	257	2	2	96	0	0	
0	0	0	450216	63328	818600	0	0	0	0	107	244	3	1	96	0	0	
0	0	0	450216	63328	818600	0	0	0	0	98	225	1	0	99	0	0	
0	0	0	450216	63328	818600	0	0	0	0	109	266	3	0	97	0	0	
0	0	0	450216	63328	818600	0	0	0	0	109	327	3	0	97	0	0	
0	0	0	450216	63328	818600	0	0	0	0	266	1159	7	6	86	0	0	
2	0	0	450216	63328	818600	0	0	0	0	437	1651	21	4	74	0	0	
0	0		450246	62220	010630	0	۵		0	674	2742	гг	_	42	0	0	

Illustration 2.4: vmstat command to check the CPU behavior.

	@rawan-V					Box)	06 20	ينا, 21	_x86	_64_	(1 0		
+04 0	2:00:21	م	CPU	%usr	%nice	%sys	%iowait	%irq	%soft	%steal	%guest	%gnice	%idle
04 0	2:00:22	م											
04 0	2:00:23	_											
04 0	2:00:24	م											
04 0	2:00:25	م											
04 0	2:00:26	م											
04 0	2:00:27	م											
04 0	2:00:28	م											
04 0	2:00:29	م	all	2.04							0.00		
04 0	2:00:30	م											
04 0	2:00:31	م											
04 0	2:00:32	م											
	2:00:33												
04 0	2:00:34	م											
04 0	2:00:35	۰	all	4.17		3.12							

Illustration 2.5: mpstat command to check the CPU behavior.

Lab report # 9: CPU Performance Monitoring Email Address: 17-0024@student.gutech.edu.om

Q2. Explain what is happening.

The script in the ex01b.py contains a function that is iterative and not recursive. Iterative functions do not use the system stack to store values. Iterative functions can take large values of n and they are fast. We notice that we took a parameter of 1200 in this script and the script ran properly without any errors. That is due to the fact the code is an iterative function. We can observe, from illustration 2.4 and 2.5, that the CPU became more engaged during the execution of the code. In other words, the idle time percentage decreased.

Step3: Remove line 14 (the print call) and then run the program.

Illustration 2.6: removing or commenting the 14th line of ex01b.py (the print statement)

Q3. Analyze the behavior now

raw	co. Ariary ze the behavior now. -awan@rawan-VirtualBox:~\$ mpstat 1 -inux 5.4.0-58-generic (rawan-VirtualBox)													
+04	02:11:39	م	CPU	%usr	%nice	%sys	%iowait	%irq	%soft	%steal	%guest	%gnice	%idle	
+04	02:11:40	۵.	all	50.51		3.03							46.46	
+04	02:11:41	۰,	all	28.00									72.00	
+04	02:11:42	۰	all	8.25		3.09								
+04	02:11:43	۰,	all	44.55									53.47	
+04	02:11:44	م	all	7.07										
+04	02:11:45	م	all	6.06										

Illustration 2.7: removing or commenting the 14th line of ex01b.py (the print statement)

We notice that since we removed the print statement, there will be less instructions for the CPU, therefore more idle time. Comparing mpstat in this case with the previous one (with the print statement), we can see an improvement. We can observe that the CPU has became more idle. Which means less engaged with instructions.

Step4: Run both ex00.py and ex01b.py simultaneously and use vmstat and mpstat to monitor what is going on.

Lab report # 9: CPU Performance Monitoring Email Address: 17-0024@student.gutech.edu.om

inu	ix 5.4.0-58	-gener	ric (rawa	n-Virtual	Box)	06 20	ينا, 21	_x86	_64_	(1 CP		
94	م 02:23:44	CPU	%usr	%nice	%sys	%iowait	%irq	%soft	%steal	%guest	%gnice	%idle
4	م 02:23:45	all	58.16		7.14							34.69
4	م 02:23:46	all	11.58		3.16							
4	م 92:23:47	all	27.84		3.09							69.07
4	م 02:23:48	all		0.00	5.00		0.00	0.00	0.00	0.00	0.00	0.00
4	م 92:23:49	all										
4	م 02:23:50	all										
4	م 02:23:51	all										
4	م 02:23:52	all										
4	م 02:23:53	all										
4	م 02:23:54	all										
4	م 02:23:55	all										
4	م 02:23:56	all	99.00	0.00	1.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00
4	م 02:23:57	all	68.37		1.02							30.61
)4	م 02:23:58	all	7.69	0.00	1.10	0.00	0.00	0.00	0.00	0.00	0.00	

Illustration 2.8: mpstat command to check the CPU behavior.

ā	IWa	ın@		/irtualB														
ı	oc	S		mem	югу		swa	p	io)	-syst	tem		0	pu			
r		Ь	swpd	free	buff	cache	si	so	bi	bo	in	cs	us	sy i	Ld۱	wa	st	
6		0	0	383716	63988	824272	0	0	134	13	288	387	50	1 5	50	0	0	
4		0	0	383700	63988	824300	0	0	0	0	694	2944	51	9 4	10	0	0	
1		0	0	383700	63988	824336	0	0	0	0	811	3515	56	8 3	36	0	0	
e		0	0	383700	63988	824336	0	0	0	0	215	850	10	2 8	38	0	0	
3		0	0	389756	63988	824300	0	0	0	0	269	790	39	4 5	57	0	0	
2		Ø	Ø	389252	63988	824328	Ō	Ø	4	Ø	575	7951	96	4	Ø	Ø	Ø	
1		0	0	391780	63988	824328	0	0	0	0	500	554	99	1	0	0	0	
3		0	0	391780	63988	824328	0	0	0	0	778	1818	99	1	0	0	0	
1		0	0	391780	63988	824328	0	0	0	0	506	408	100	0	0	0	0	
1		0	0	391780	63996	824320	0	0	0	48	476	249	100	0	0	0	0	
1		0	0	391780	63996	824328	0	0	0	0	485	267	100	0	0	0	0	
1		0	0	391780	63996	824328	0	0	0	0	575	818	100	0	0	0	0	
1		0	0	391780	63996	824328	0	0	0	0	549	491	100	0	0	0	0	
2		0	0	391780	63996	824328	0	0	0	0	762	1830	98	2	0	0	0	
1		0	0	391780	63996	824340	0	0	0	0	633	2162	60	0 4	10	0	0	
3		0	0	391780	63996	824340	0	0	0	0	379	1557	17	2 8	31	0	0	
2		0	0	391780	63996	824376	0	0	0	0	797	3160	54	3 4	13	0	0	

Illustration 2.9: vmstat command to check the CPU behavior.

From illustration 2.8 and 2.9, we can observe that when the two scripts ran together, the CPU was utilized almost completely. Both are taking more CPU. However, after ex01b.py finishes its execution and the infinite loop in ex00.py is still running, this causes a decrease in the CPU utilization.

.