

Computational restrictions on iterative prosodic processes

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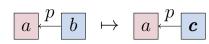
Overview

- **Goal**: Computational characterizations lead to restrictive, testable, and learnable theories of phonology (Heinz, 2018).
- **Question**: what are the computational requirements of **iterative prosody**?
- **Result**: Local (with recursion)!

Logical maps

• Logic: output defined over input (Courcelle, 1994)

$$c(x) \stackrel{\text{def}}{=} b(x) \wedge a(p(x))$$



- Local = Quantifier Free: no ∃, ∀ (Chandlee and Lindell, forthcoming)
- Recursion = Least Fixed Point (LFP) operators (Libkin 2004)
- Local recursion = QF-LFP reference information from the output string
- Implicit Recursion = *implicit definitions*; Rogers 1997
- Use either predecessor (p) or successor (s)

Quantifier Free syllabification

Simple non-directional syllabification is local (Strother-Garcia, 2019)

- To find a nucleus \rightarrow use a window of size 4
- Window of segment + predecessor + 2 segments that follow
- Don't need any quantifiers (QF \sim ISL functions)
- Contrast with global non-local OT (Prince and Smolensky, 2004)

Least Fixed Point: Local recursion

Example: $baaa \mapsto bbbb$

- = a's following a b are outputted as b
- Not local on the input
- Local on output: $a \rightarrow b$ when b before a on output

Implicit definition: $b'(x) \stackrel{d}{=} b(x) \lor b'(p(x))$

- Given an input element x,
- it is mapped to a b in the output
- when it is a b in the input or
- it is preceded by a b in the output.

Iterative stress

• Murinbata: stress 1st σ and every other σ $\sigma\sigma\sigma\sigma\sigma\sigma\sigma\mapsto \acute{\sigma}\sigma \dot{\sigma}\sigma\dot{\sigma}\sigma\dot{\sigma}\sigma\dot{\sigma}$

Formalize...

- $\dot{\sigma}(x) \stackrel{d}{=} first(x) \vee \dot{\sigma}(p(p(x)))$

Iterative syllabification

Arabic dialects: different epenthesis sites in CC* clusters (Ito, 1989)

- 3C: insert V after C_1 in Iraqi, and C_2 in Cairene.
- 4C: insert V after C_2 .

Why?

• Iraqi syllabifies R-to-L, while Cairene L-to-R + a V is added based on a CVC template.

		*
Iraqi	<katab-t-l-u></katab-t-l-u>	<katab-t-l-ha></katab-t-l-ha>
(R-to-L)	.ka.ta.bit.lu.	.ka.tab.til.ha.
Cairene	<katab-t-l-u></katab-t-l-u>	<katab-t-l-ha></katab-t-l-ha>
(L-to-R)	.ka.tab.ti.lu	.ka.tab.til.ha

QFLFP Characterization

- L'(x) and R'(x) determine L- and R-edges of σ 's before resyllabification.
- Resyllabification is only apparent in L-to-R.
- We only show R-to-L parsing

$L'(x) \stackrel{\mathrm{d}}{=}$	$[C(x) \wedge V(s(x))] \vee$	select C in <u>C</u> V
	$[C(x) \wedge C(s(x)) \wedge L(s(s(x)))]$	select C in <u>C</u> C[_L
$i'(x_2) \stackrel{\mathrm{d}}{=}$	$C(x) \wedge L(x) \wedge C(s(x))$	add V in $[_LC_C]$

For example...

Input	k	a	t	a	b	t	l	и
L' is true at								
Iteration 0	1		✓				✓	
Iteration 1	1		✓		✓		✓	
Interim Output:	k_L	a	t_L	a	b_L	t	l_L	и
					/			
$i'(x_2)$ Output:	k	a	t	a	<i>b</i> i	t	l	и

Discussion

- Provides a testable hypothesis for iterative phonological functions based on computational power
- Highlights output orientation of iterative functions
- What about feet? What about more patterns?

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