1 Cryptography: An Overview

1.2 Enigma

Problem 4. Consider the notion of "evolution" that we introduced in this section. We can formulate this notion mathematically as follows. For a given initial setting of the machine, as determined, for example, by the daily key, let P_m be the permutation executed by the machine after m keystrokes. We considered a case in which, for some unknown letter x, $P_0x = Q$ and $P_3x = E$. This is what we meant by saying that Q "evolves" into E. Evidently the permutation that expresses this evolution is $P_3P_0^{-1}$. (Applying P_0^{-1} to Q gives x, and then applying P_3 gives P_3 . Show that the cycle lengths of the "evolution" $P_3P_0^{-1}$ are indeed independent of the plugboard permutation P_3 , as was noted by Rejewski. (Note that the permutation P_m is given by P_3 , with appropriate values of P_3 , and P_3 .)

Intuitively, the cycle lengths won't be affected by any setting of the plugboard as the plugboard statically shuffles the characters the rotor-reflector assembly receives and provides. The problem can be rigorously outlined as it was by Rejewski using the following equation, where A represents the involution of the plugboard $(A = A^{-1})$, S^n represents the effect of rotating n times, R_i represents the i-th rotor, and B represents the reflector.

$$A^{-1}(S^{-n}R_1S^n)^{-1}B(S^{-n}R_1S^n)A$$

The equation above represents the permutations a one-rotor Enigma machine would act on input, where n is incremented once per letter modulo 26. In group theory¹ the closed operation $A^{-1}(...)A$ represents a congugacy where the plugboard's involution doesn't affect the cycle length of the permutation it conjugates, i.e. $P^m = A^{-1}R^mA$. Therefore Rejewski concludes that the cycle lengths are consistent across all plugboard configurations, significantly reducing the number of combinations. After documenting the cycles, one can simply treat this iteration of the Enigma machine as a substitution cipher, to my knowledge.

Remark. I'd be curious to know how other students are rigorously outlining this answer; I may be reading into this question too much or I missed the lecture on congugacy. I remember this idea being briefly outlined in the introductory Linear Algebra course via matrix similarity, but I don't know if the average student would make the connection or be able to say much more than "because that's how it is with matrices" which is really just a case within non-abelian groups. The intuitive answer is too obvious and the rigorous answer too rigorous for this text, in my humble opinion. Ignore this if it was explained explicitly in class; I need to be present more!

Problem 5. Consider the Enigma machine with a certain initial setting of the rotors and plugboard. With this initial setting, let P_0 be the permutation the machine applies to the first letter of the plaintext, and let P_3 be the permutation that it applies to the fourth letter of the plaintext. Recall the following two facts about the permutations P_0 and P_3 :

(i)
$$P_0^{-1} = P_0$$
 and $P_3^{-1} = P_3$;

¹I think this is right; I haven't formally studied group theory and am working off of my first impressions from Wikipedia.

(ii) P_0 does not send any letter to itself, and neither does P_3 .

These facts will be useful in this problem.

We have seen how cryptanalysts were able to crack Enigma by considering the lengths of the cycles of the permutation $P_3P_0^{-1}$. Let y_1, y_2, \ldots, y_m be a cycle of this permutation. That is, the y_i 's are m distinct letters of the alphabet, and $P_3P_0^{-1}y_1 = y_2, P_3P_0^{-1}y_2 = y_3, \ldots, P_3P_0^{-1}y_m = y_1$.

- 1. Show that $P_0y_m, P_0y_{m-1}, \ldots, P_0y_1$ is also a cycle of $P_3P_0^{-1}$.
- 2. Show that the cycle defined in part (a) consists entirely of letters that do not appear in the original cycle y_1, y_2, \ldots, y_m . It follows that the cycle lengths always come in matching pairs.

Blah blah blah heres my answer...

1.3 A Review of Modular Arithmetic and \mathbb{Z}_n

Problem 1. Write out the addition and multiplication tables for \mathbb{Z}_3 , \mathbb{Z}_4 , and \mathbb{Z}_7 .

First, the addition tables (\mathbb{Z}_n^+) .

		$\mathbb{Z}_7^+ \mid 0 1 2 3 4 5 6$
	$77 + \begin{vmatrix} 0 & 1 & 0 & 2 \end{vmatrix}$	0 0 1 2 3 4 5 6
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	1 1 2 3 4 5 6 0
	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$2 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 0 \mid 1$
	$2 \mid 2 \mid 3 \mid 0 \mid 1$	3 3 4 5 6 0 1 2
		$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
'	$3 \mid 3 0 1 2$	5 5 6 0 1 2 3 4
		$6 \mid 6 \mid 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5$

Next, the multiplication tables (\mathbb{Z}_n^{\times}) .

$\mathbb{Z}_7^{ imes}$	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
	1						
	1						
	0 1 2 3 4 5	0 0 1 0 2 0 3 0 4 0 5 0	0 0 0 1 0 1 2 0 2 3 0 3 4 0 4 5 0 5	0 0 0 0 1 0 1 2 2 0 2 4 3 0 3 6 4 0 4 1 5 0 5 3	0 0 0 0 0 1 0 1 2 3 2 0 2 4 6 3 0 3 6 2 4 0 4 1 5 5 0 5 3 1	0 0 0 0 0 0 1 0 1 2 3 4 2 0 2 4 6 1 3 0 3 6 2 5 4 0 4 1 5 2 5 0 5 3 1 6	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

Problem 4. Show that addition and multiplication modulo n are well defined. In other words, show that if $a \equiv b \pmod{n}$ and $c \equiv d \pmod{n}$, then $a + c \equiv b + d \pmod{n}$ and $ac \equiv bd \pmod{n}$.

In order to show that the addition and multiplication operations endowed to the \mathbb{Z}_n ring are well defined, we need to show that both $(\mathbb{Z}, +)$ and (\mathbb{Z}_n, \cdot) used in a mapping won't produce two different, valid results in the range, i.e., for $f: a \mapsto b$ we need a unique value $b \in \mathbb{Z}_n$ given $a \in \mathbb{Z}_n$.

To start, we can rewrite the statements $a \equiv b \pmod{n}$ as a = b + in and $c \equiv d \pmod{n}$ as c = d + jn for some $i, j, n \in \mathbb{Z}$. For the addition operation, we can say that

$$a+c = (b+in) + (d+jn) \implies a+c = b+d+n(i+j)$$

 $\implies a+c = b+d \pmod{n}.$

a unique element in \mathbb{Z}_n . For the multiplication operation, we can say

$$a \cdot c = (b+in) \cdot (d+jn) \implies bd + b(jn) + d(in) + (in)(jn)$$

 $\implies a \cdot c = b \cdot d + n(bj + di + ijn)$
 $\implies a \cdot c = b \cdot d \pmod{n}$

which is a unique element in \mathbb{Z}_n . Thus, both addition and multiplication are well defined binary operations in \mathbb{Z}_n meaning their results do not depend on the choice of representatives in each equivalence class.

Matrix Practice

Consider the matrix

$$M = \begin{pmatrix} 2 & 3 & 9 \\ 3 & -1 & 2 \\ 4 & -4 & 7 \end{pmatrix}.$$

Calculate each of the matrix products below and use words to describe the product. For example, you would write

$$\begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix} M = \begin{pmatrix} 3 & -1 & 2 \\ 2 & 3 & 9 \\ 4 & -4 & 7 \end{pmatrix}$$

and say: "multiplication on the left by $\begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$ switches the first two rows of M."

Problem 1.
$$\begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix} M$$

Problem 2.
$$M \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

- **Problem 3.** $\begin{pmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix} M$
- **Problem 4.** $M \begin{pmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix}$
- **Problem 5.** $\begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} M$
- **Problem 6.** $M \begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$
- **Problem 7.** $\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 3 \end{pmatrix} M$
- **Problem 8.** $M \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 3 \end{pmatrix}$
- **Problem 9.** $\begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} M$
- **Problem 10.** $M \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$
- **Problem 11.** $\begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix} M$
- **Problem 12.** $M \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}$

Problem 13.
$$\begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} M$$

Problem 14.
$$M \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$

Problem 15.
$$\begin{pmatrix} 0 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix} M \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

Problem 16.
$$\begin{pmatrix} 0 & 0 & 1 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} M \begin{pmatrix} 0 & 0 & 1 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$