Note: unrequired parts of multi-part problems are listed obfuscated as ... to recognize they are multi-parted.

## 1 Arithmetic

## 1.3 Division Algorithm

**Problem 5.** Use old-fashioned long division to implement the Division Algorithm and write a = bq + r,  $0 \le r < b$ , for each a and b listed below:

- (a) a = 20, b = 3.
- (b) a = 54, b = 7.

Starting with the first subproblem, we can start by finding the maximal q in the set of possible divisors S. By examination we can quickly deduce

$$20 = 3(6) + 2, 0 \le 2 < 3.$$

We can do the same for the second subproblem:

$$\boxed{54 = 7(7) + 5, 0 \le 5 < 7.}$$

**Problem 8.** Show that when the square of an odd integer is divided by 8, the remainder is 1. (Hint: remember 2|n(n+1).)

We will restate the problem as a proposition and prove it using the Division Algorithm.

**Proposition.** Given the square of an odd integer, it's remainder is always 1 when divided by 8.

*Proof.* We start with a simple expression of any odd integer, 2n + 1. We can restate our proposition as  $\exists q \in \mathbb{Z}_{\geq 0}$  such that  $(2n + 1)^2 = 8q + 1 \ \forall n \in \mathbb{Z}_{\geq 0}$ . We can rearrange the statement to determine whether q is an integer:

$$(2n+1)^2 = 8q + 1$$
$$4n^2 + 4n + 1 = 8q + 1$$
$$4n(n+1) = 8q$$
$$\frac{1}{2}n(n+1) = q$$

Because n(n+1) is even (or 0) we now know that  $\forall n \in \mathbb{Z}_{\geq 0} \implies \exists q \in \mathbb{Z}_{\geq 0}$  given a remainder of 1. This is important as it ensures our divisor is an integer. Thus  $(2n+1)^2/8$  will yield a remainder of  $1 \ \forall n \in \mathbb{Z}_{\geq 0}$ .

Remark. Is this too discursive? There is probably a more elegant solution.

**Problem 10.** Let  $n, m \in \mathbb{N}$  with  $m \neq 1$ . Show n can be uniquely written in the form  $n = \sum_{k=0}^{N} a_k m^k$  for some  $N \in \mathbb{Z}_{\geq 0}$  and  $a_k \in \{0, 1, \dots, (m-1)\}$  with  $a_N \neq 0$ . Hint: Use induction on n and begin by choosing the largest  $N \in \mathbb{Z}_{\geq 0}$  so that  $m^N \leq n$ . Use the Division Algorithm to write  $n = a_N m^N + r$  and then apply the inductive hypothesis to r.

**Proposition.** Given two positive integers m, n where  $m \neq 1$ , there is a unique representation of  $n = \sum_{k=0}^{N} a_k m^k$  for some  $N \in \mathbb{Z}_{\geq 0}$  while  $a_k \in \{0, 1, \dots, m-1\}$  and  $a_N \neq 0$ .

*Proof.* Suppose  $m, n \in \mathbb{N}$  and  $m \geq 2$ . We'll first show that the representation exists. Starting with the base case of n = 0, the representation  $0 = n = \sum_{k=0}^{N} = a_k m^k = a_0 m^0 = a_0$  which is true as  $a_0 \in \{0, \dots, m-1\}$  and N = 0. Next with the inductive step, we choose an N such that  $m^N \leq n \leq m^{N+1}$ . Using the Division Algorithm's expression of  $n = a_N m^N + r$ , we know that r is a recursive iteration of this expression and can thus be expressed as  $r = \sum_{k=0}^{N-1} a_N m^N$ . Combining our base case and r, we see that

$$n = a_N m^N + \sum_{k=0}^{N-1} a_N m^N \implies n = \sum_{k=0}^{N} a_N m^N.$$

Next, to prove this representation is unique, consider two expressions of n that are supposedly equal in the non-unique case:

$$\sum_{k=0}^{N} a_k m^k = \sum_{k=0}^{M} b_k m^k$$

$$\sum_{k=0}^{N} a_k m^k - \sum_{k=0}^{M} b_k m^k = 0$$

$$\sum_{k=0}^{\max(M,N)} (a_k - b_k) m^k = 0.$$

The third line of the above derivation can only be true if  $(a_k - b_k) = 0 \implies a_k = b_k$ . Thus the representation is unique.

## 1.4 Divisors

**Problem 12.** List all of the divisors of the following:

- (a) 52,
- (b) ...

The divisors of 52 are  $\pm \{1, 2, 4, 13, 26, 52\}$ .

**Problem 14.** Evaluate the following:

(a) (42,56),

(b) ...

The expression (42,56), otherwise written as gcd(42,56), is equivalent to  $max(A \cap B)$  where  $A = \{1,2,3,6,7,14,21,42\}$ , the divisors of 42, and  $B = \{1,2,4,7,8,14,28,56\}$ , the divisors of 56. Thus the greatest common divisor of these two numbers is 14.

**Problem 17.** Let  $b, q, r \in \mathbb{Z}$  and let a = bq + r with a and b not both 0.

- (a) Show a common divisor of a and b is a divisor of r and that a common divisor of b and r is a divisor of a.
- (b) Conclude that (a,b) = (r,b)

**Proposition.** A common divisor of  $a, b \in \mathbb{Z}$  divides  $r \wedge a$  common divisor of  $b, r \in \mathbb{Z}$  divides  $a \implies (a, b) = (r, b)$  given a = br + r,  $a, b \neq 0$ .

*Proof.* Suppose  $b, q, r \in \mathbb{Z}$  such that a = bq + r and that a, b have a common divisor y. The given equation can be rewritten as iy = jyq + r where iy = a and jy = b. We can rearrange this equation to isolate r:

$$r = iy - jyq$$

$$r = y(i - jq)$$

$$r = yk$$

knowing k = (i - jq). We can then say that y|r using the multiple  $k \in \mathbb{Z}$ . Similarly let us suppose b, r have a common divisor z; they can be rewritten as r = mz and b = nz. Using our given expression, we can say

$$a = nzq + mz$$
$$a = z(nq + m)$$
$$a = zl$$

knowing l = (nq + m). Thus z|a using the multiple  $l \in \mathbb{Z}$ . Given that a, b|r and r, b|a we know that the set of common divisors between a, b and r, b are the same and hence their maxes are the same. Thus  $\gcd(a, b) = \gcd(r, b)$ .

**Problem 19.** Let  $a, b \in \mathbb{Z}$ , not both 0.

(a) If 
$$(a,b) = d$$
, then  $(\frac{a}{d}, \frac{b}{d}) = 1$ . Hint: Write  $ax + by = d$  so that  $\frac{a}{d}x + \frac{b}{d}y = 1$ .

(b) ...

We know via Theorem 1.10 from the text that the greatest common divisor of two integers is their minimal linear combination, i.e.  $gcd(a, b) = d \iff d = min(\{d|ax + by = d, x, y \in \mathbb{Z}\})$ . We can then write

$$ax + by = d$$
$$\frac{a}{d}x + \frac{b}{d}y = 1.$$

Because 1 is the smallest outcome of any linear combination assuming a, b are not both zero, we can say that  $\frac{a}{d}$  and  $\frac{b}{d}$  have a greatest common divisor of 1 and are thus relatively prime. This result is intuitive as the greatest common divisor is 'divided' out of each number, leaving them relatively prime.