

Concurrent Programming (Part II) Lecture 6: Semaphores, Nested Monitors and Deadlock

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Course Web Site on Moodle

http://moodle.ucl.ac.uk/course/view.php?id=753

Enrolment Key: ATOMIC

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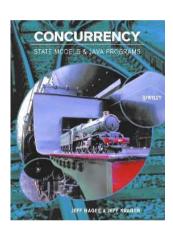
Overview of Lecture

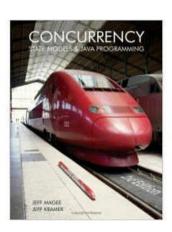
- We will introduce the Semaphore this is a key concurrency abstractions within a number of systems, for example:
 - Within Unix/POSIX (do "man sem_init").
 - Within Berkley SystemV Unix (do "man semop")
 - Within PERL (do "man Thread::Semaphore")
 - ... and now within Java (java.util.concurrent.Semaphore)
- We will examine how we can create a Java
 Semaphore using conditional synchronization.
- We will then reimplement the BoundedBuffer using our Java Semaphore.
- ... which will lead us into lots of problems!

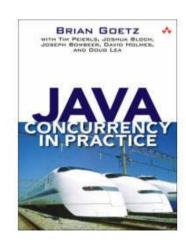
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Semaphores – you might have noticed a train theme on concurrency books ...

- Semaphores were introduced by Dijkstra in 1968 as one of the first mechanisms for inter-process synchronization (in the humbly named "THE multiprogramming system")
- The term comes from the railroad –
 a railway semaphore is a signal flag
 stopping multiple trains being on
 the same track.



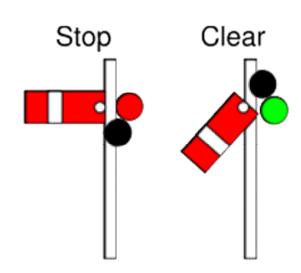




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Railway Semaphores



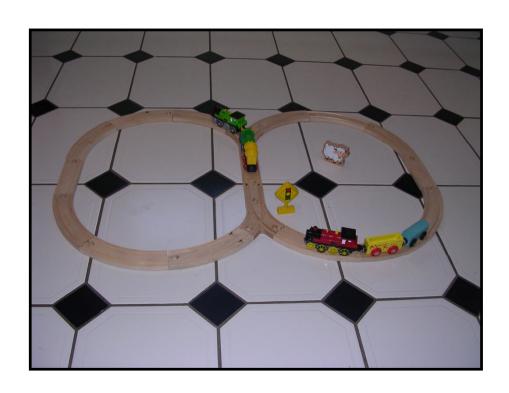


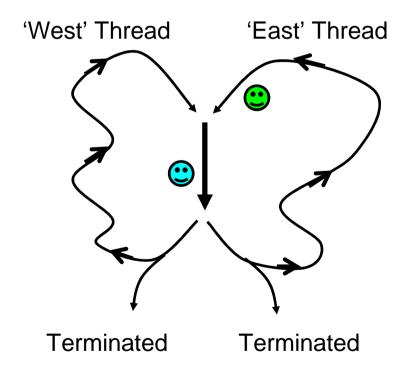
- Trains are real-world objects with concurrent behaviour (true parallelism to be precise ...)
- Semaphores are used to control the concurrent behaviour of trains by getting them to wait at particular points (thus avoiding multiple trains being on *critical sections of track*)



Transferring the railway concept to CS ...

See the analogy ... semaphores ... critical sections of code ...







Semaphores – Definition

Semaphores are widely used for dealing with inter-process synchronization in operating systems. Semaphore **s** is an integer variable that can take only non-negative values.

The only operations permitted on **s** are **up(s)** and **down(s)**.

Semaphores – Implementation

Semaphores are passive objects, therefore implemented as **monitors**.

(In practice, semaphores are a low-level mechanism often used in implementing the higher-level monitor construct.)

```
public class Semaphore {
  private int value;
  public Semaphore (int initial)
    {value = initial;}
  synchronized public void up() {
     ++value;
     notify();
  synchronized public void down()
      throws InterruptedException {
    while (value== 0) wait();
    --value;
```



Or we could just use the Java 5 implementation (java.util.concurrent.Semaphore)

```
public class Semaphore {
   public Semaphore(long permits);
   public Semaphore(long permits, boolean fair);
   public void acquire() throws InterruptedException;
   public void acquire(long permits) throws InterruptedException;
   public boolean tryAcquire();
   public boolean tryAcquire(long timeout, TimeUnit unit);
   public release();
   public release(long permits);
```

Using Semaphores



- A semaphore can be used to restrict access to a shared resource.
- The count is set to the number of threads that are permitted access simultaneously.

```
class ResourceHandler {
  Semaphore access = new Semaphore(10);
  public void accessResource() {
     try {
           access.down();
           // Use resource ...
      } finally {
           access.up();
```

Binary Semaphores



- The previous Semaphores are also known as 'Counting Semaphores'
- This distinguishes them from 'Binary Semaphores' which are initialized with a count of 1.
- Binary Semaphores behave like mutex (mutually exclusive) locks but are not re-entrant (unlike Java locks which are).

```
class CriticalCode {
  // Binary Semaphore with a value of 1.
  Semaphore access = new Semaphore(1);
  public void doSomethingWithSharedData() {
     try {
           access.down();
           // Critical section of code.
      } finally {
           access.up();
```