A photograph of the University College London building, featuring a large portico with columns and a dome. The building is surrounded by trees and a paved area with people walking. The text is overlaid on the image.

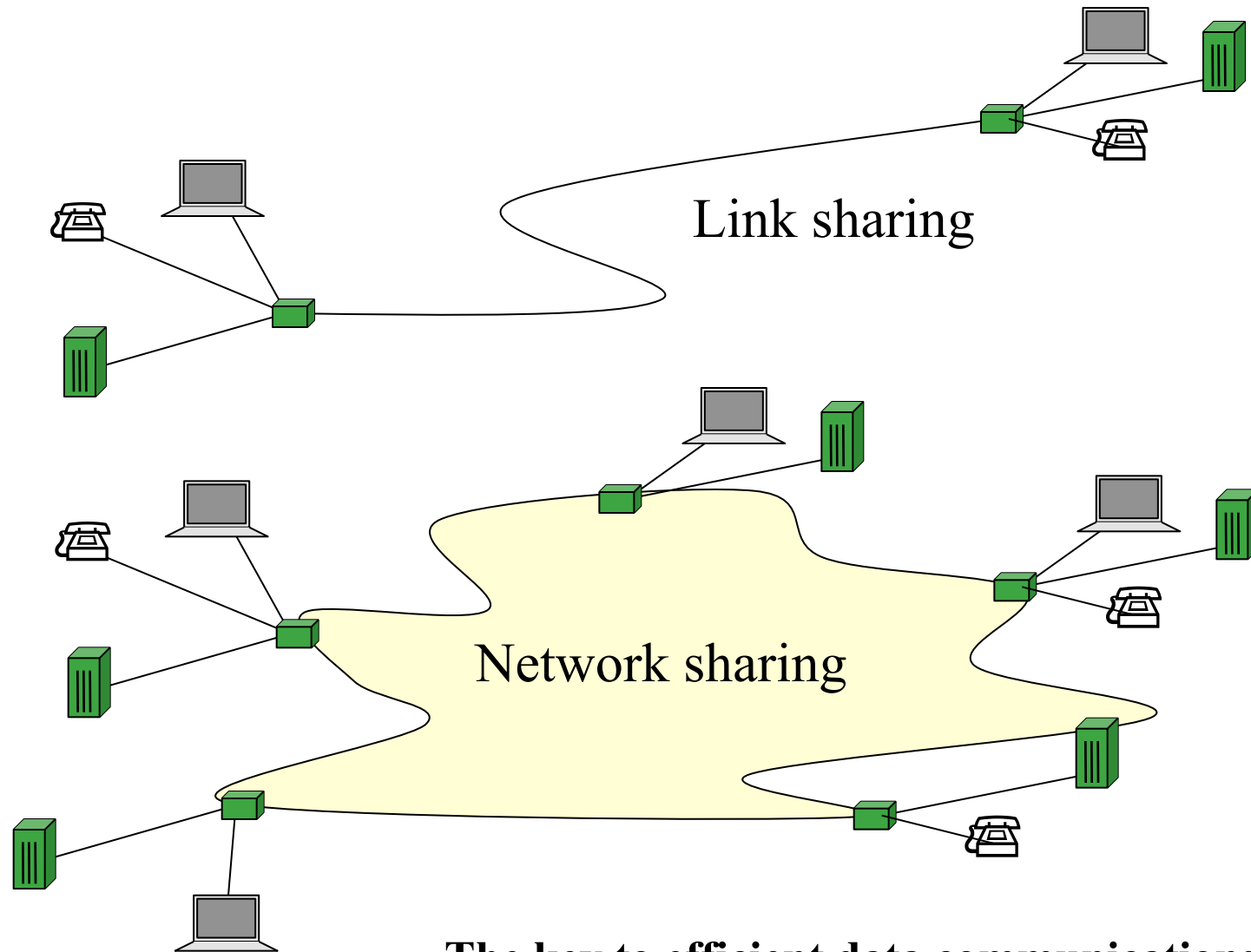
**COMP2011 -- Networks,
Databases and Graphics**

3. Sharing Resources

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Sharing



The key to efficient data communications is the efficient sharing of physical resources.

A little history

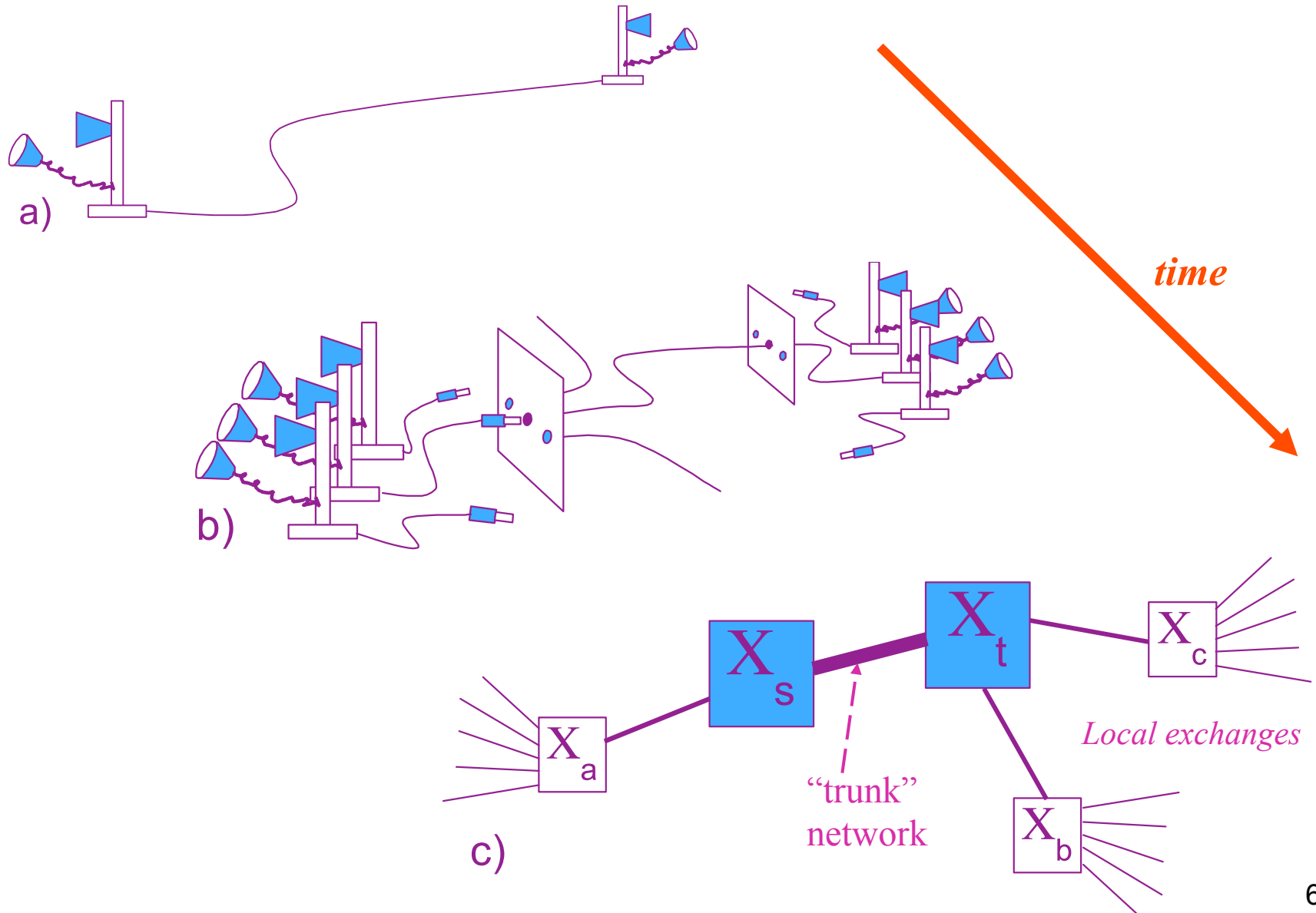
- 1837 the telegraph (Cooke and Wheatstone)
- 1838 Morse code (Samuel Morse)
 - Electronic digital communication
- 1843 Facsimile transmission
 - Data communication
- 1865 Commercial facsimile transmission
- 1876 The telephone (Alexander Graham Bell)
 - Analogue voice
- 1878 First commercial telephone exchange
- Analogue telephone network grew and grew ...
 - Telecommunications

A little more history

- 1960s communication between computers
 - Over quite long distance, Wide Area Networks (WANs)
 - Used analogue telephone infrastructure (circuit-switching)
 - MODEMs (digital-over-analogue-channel)
 - Low capacity, noisy, error-prone
- 1961 Packet-switching
 - Kleinrock and Baran (US), Davies and Scantlebury (UK), and Pouzin (France)
- 1970s Mini-computers, 1980s Micro-computers (desktop workstations and PCs)
 - Local Area Networks (LAN)
 - Share media network, independent of public network
- 1972 Metcalfe's Harvard PhD Thesis
 - Ethernet LAN and Metropolitan Area Network (MAN)

Part 1. Circuit-switching

Circuits and Circuit-switching



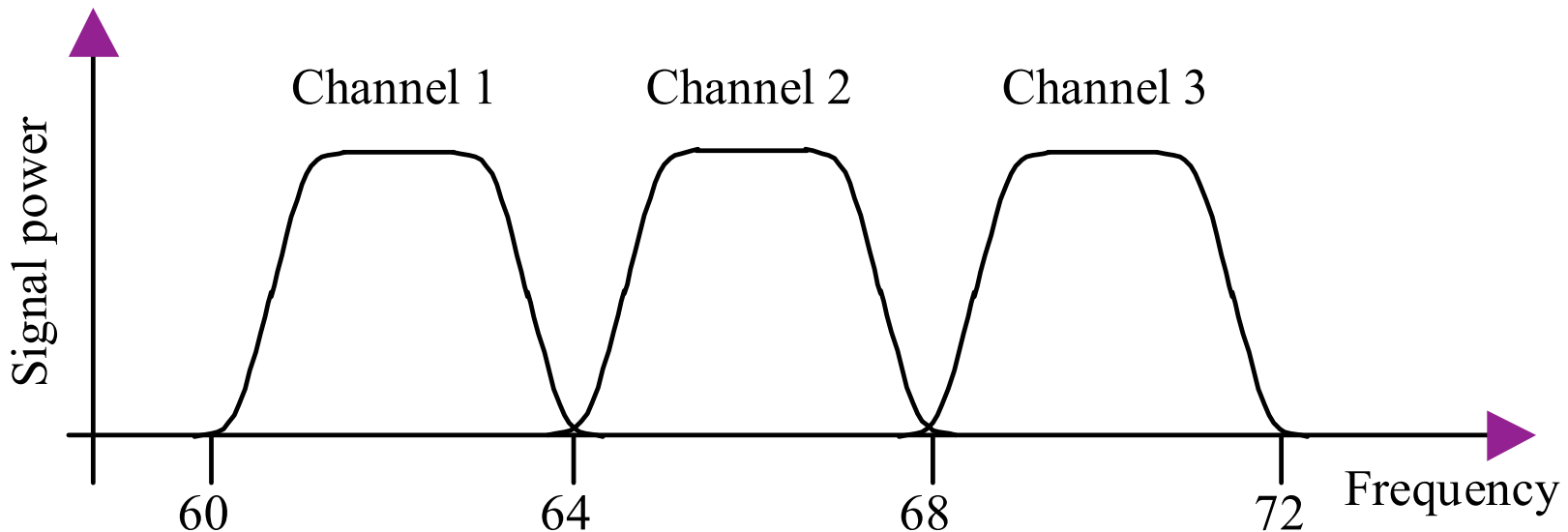
Eight characteristics of telephone service

- 1.Any telephone can communicate with any other telephone;
- 2.A telephone can only communicate with one other telephone at a time;
- 3.Other telephones may be "blocked" - unable to communicate because a wire they need is in use;
- 4.Communication between two telephones (once established) is unaffected by activity elsewhere in the network.
- 5.Whilst they are communicating, the two telephones have exclusive use of the wires that link them;
- 6.Words are delivered in the same order in which they are spoken.
- 7.Time intervals between words are preserved
- 8.We can distinguish three phases in the life of a circuit; setting it up, using it and tearing it down

Sharing the link..... Early solution to improve use of capacity of trunk network

Frequency Division Multiplexing (FDM)

- To provide sub-channels over a single piece of wire
- 12 carrier frequencies at 4KHz intervals
 - from 60 KHz to 108 KHz; carrier spreads ± 1.5 KHz
- Voice waves filtered to approx. 3KHz bandwidth
- Voice modulated on carrier



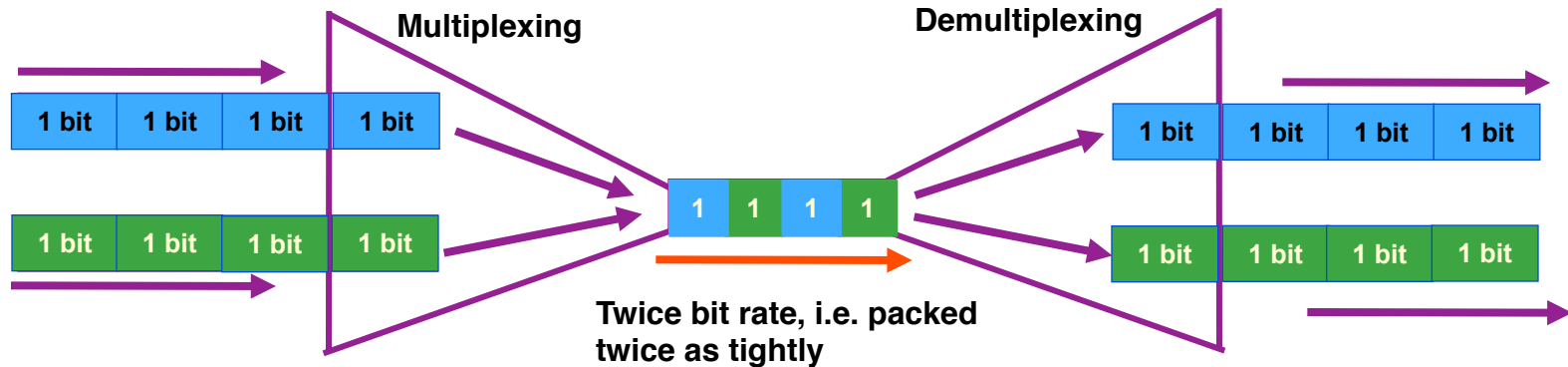
High “capacity” not High “speed”

- Higher speed vehicles travel faster
- Bits always go at the same speed
 - $\sim 2 \times 10^8$ m/s in copper, $\sim 3 \times 10^8$ in fibre
- “High-speed link” means more bits are injected and delivered in a given interval
 - It is really “high-capacity”

File transfer can complete more quickly over a higher capacity link!

Time Division Multiplexing

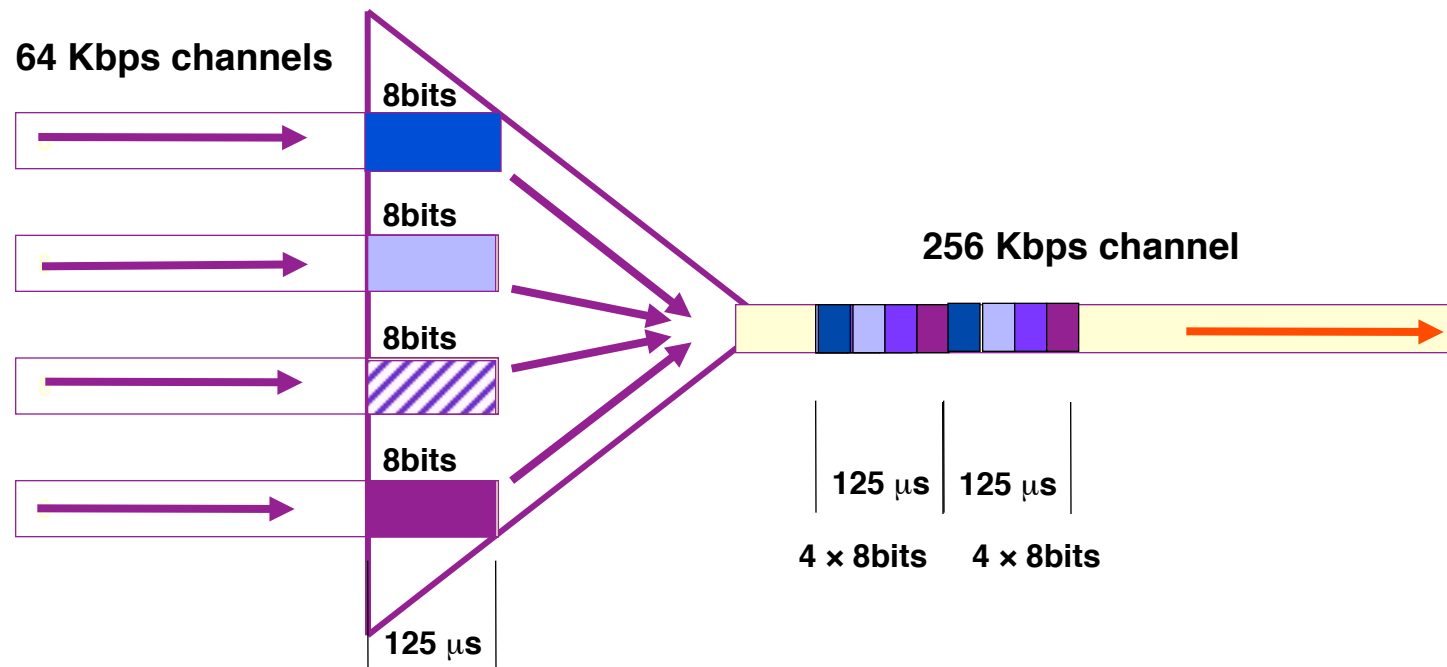
- TDM interleaves bits (sometimes bytes) from several streams onto a higher bit-rate stream



- Timings between bits of a stream are preserved
- The effect is as if there are two separate “bit-pipes”
- Easy to turn TDM system into a switch
 - to connect arbitrary inputs and outputs

Time Division Multiplexing

- Voice digitised by a CODEC
 - 64 Kbps pulse code modulation (PCM)
- Multiple digitised channels multiplexed on to “high capacity channel”
- Inevitable delays



Modern switched circuit

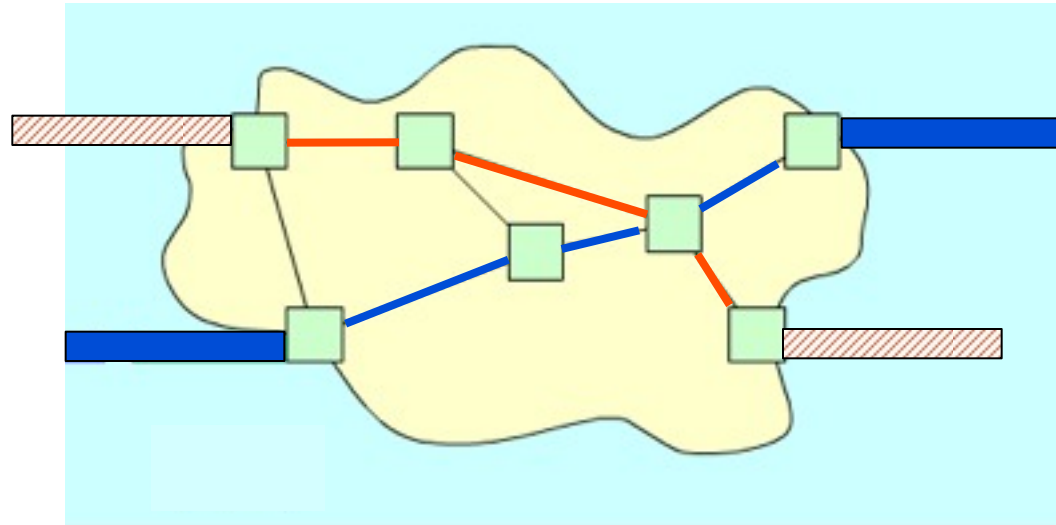
- Circuit – constant bit-rate pipe across network
 - Typically based on TDM through circuit-switches

1. Switched circuits

- Set up dynamically using signalling protocols
- Charges based on duration
- Examples: PSTN and ISDN

2. Leased circuits

- Leased lines
- Set up “administratively”
- Monthly charge



Circuit-switching is inefficient for data communications

- A host can only communication with one (or a small fixed number of) host at a time
 - But we'd like to to communicate with multiple computers concurrently
- Time intervals between bits are preserved
 - But we'd like to communicate between dissimilar devices
- The technology used to digitise calls inside a modern telecomm network results in a continuous stream of bits at 64Kbps
 - But data is “bursty”, we then have to open a circuit with capacity sufficient for the peak data-rate, while for much of the time the opened circuit will be idle.
- Fixed allocations
 - If we want 65Kbps, we may have to ask for 128 Kbps
- **Data communication needs a different approach**

Part 2. Packet-switching

TDM

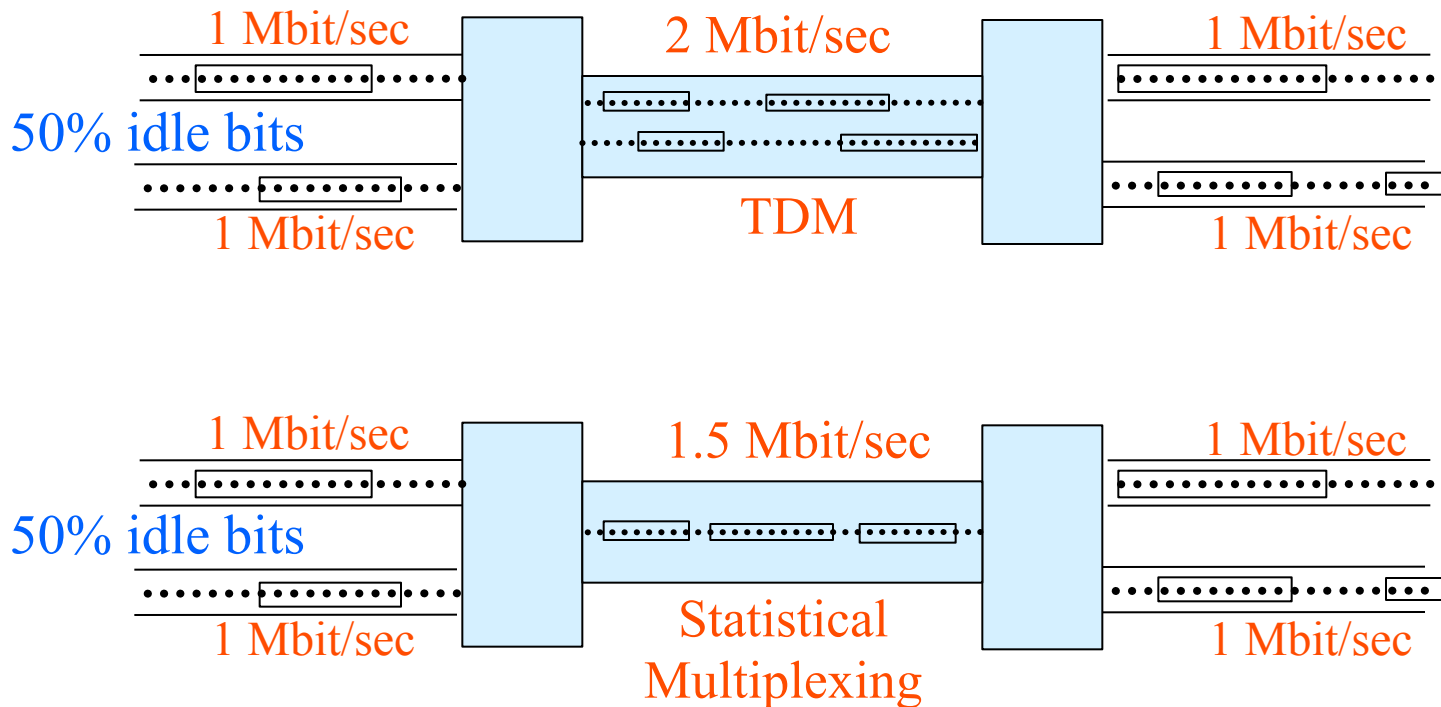
- TDM
 - works at bit (or byte) level
 - a stream of bits at a constant rate
- However,
 - most application generate data in bursts interspersed with periods of idleness.
 - in data communication not all bits are useful, e.g. idle pattern in bit-oriented synchronous framing

Sharing the link..... packet solution to improve use of channel capacity

Statistical Multiplexing

- Data sent in “lumps” called “frames”, “packets” etc.
- Lumps may vary in size
- Inter-arrival times are random
- Statistical multiplexor: packet-switches, Internet routers
- Usually processing is FIFO (First In, First Out)
 - Quality of service (QoS) needed
- Can overload an output => queues needed
 - Data discarded and additional delay

TDM vs Statistical Multiplexing



- Statistical multiplexing treats the data bits and ignore the idle bits
 - lower link capacity requirement
 - higher utilisation rate

Delays in Statistical Multiplexing

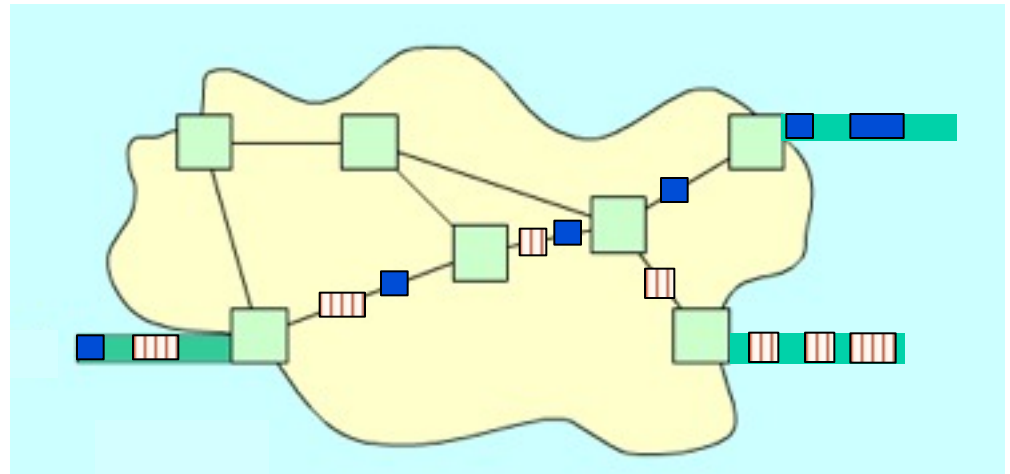
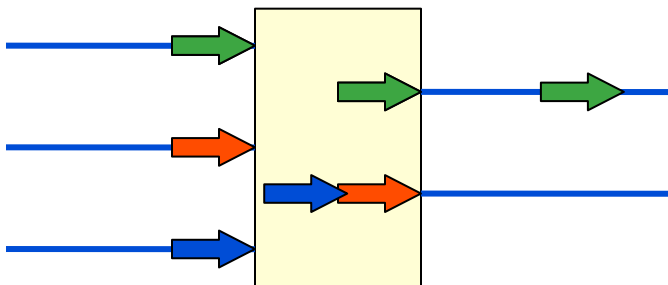
- Statistical multiplexing has longer, more variable delays than TDM
 - A complete frame must be received before transmission begins
 - A frame of 15,000 bits has delay of 10ms for 1.5Mbps
 - Queuing to solve possible contention when two frames arrive simultaneously
 - Highly variable, cause jitter
 - A frame may be discarded, cause data loss
 - Additional overhead
 - Some bits in the frames must carry control information, e.g. to indicate to which channels they belong
- Still the best choice for most data communications
 - Major advantage: enhanced utilisation for bursty data

Packet

- Packet -- a collection of bits
 - Packet lengths vary
 - mostly less than 1500 bytes
 - objects often need to be split into a series of packets
 - carry control information, overhead

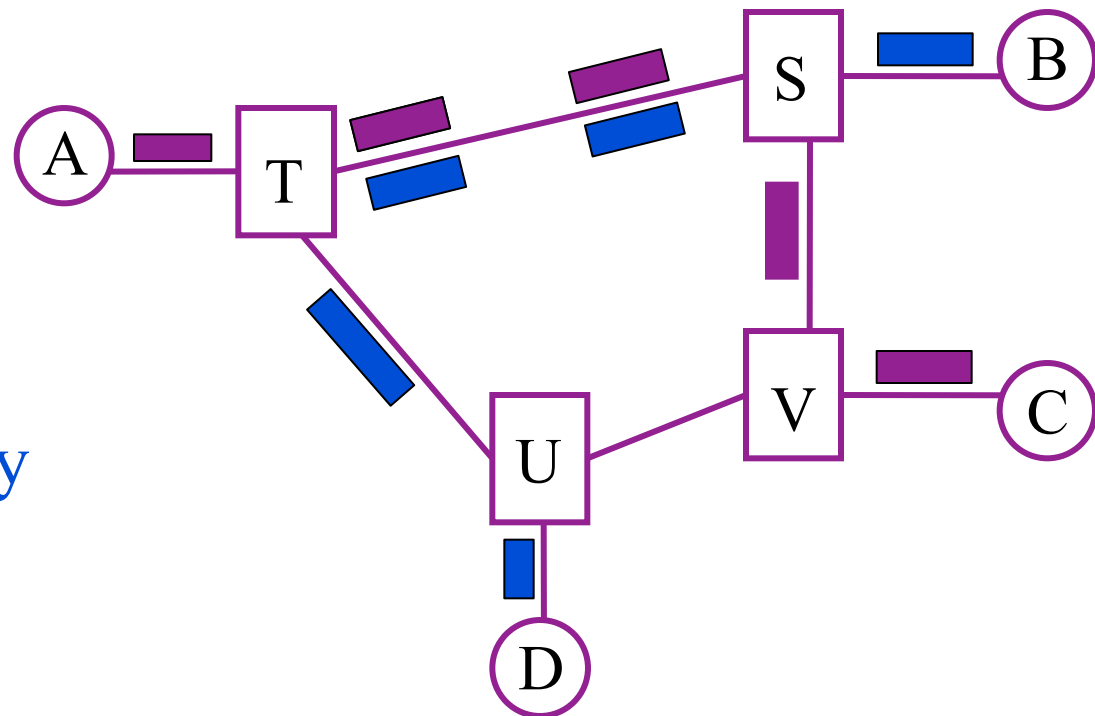
Packet-switching

- Based on statistical multiplexing
- “Packet-switches” – store and forward
 - packets from different sources are mixed on a link
 - packets multiplexed to many destinations
 - only send a packet when there is data to send
 - packets arrive at random intervals
- Allocate network capacity on-demand with fine time granularity



Mesh-type packet networks

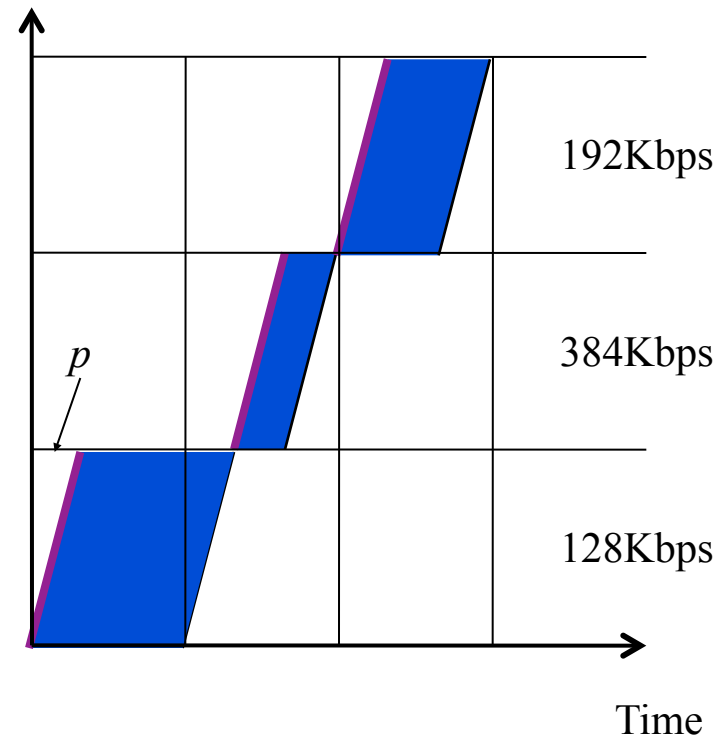
- Topology similar to circuit-switched network
 - “Packet-switch”, “router”
 - Store and forward operation



- X.25
- Frame relay
- ATM
- Ethernet

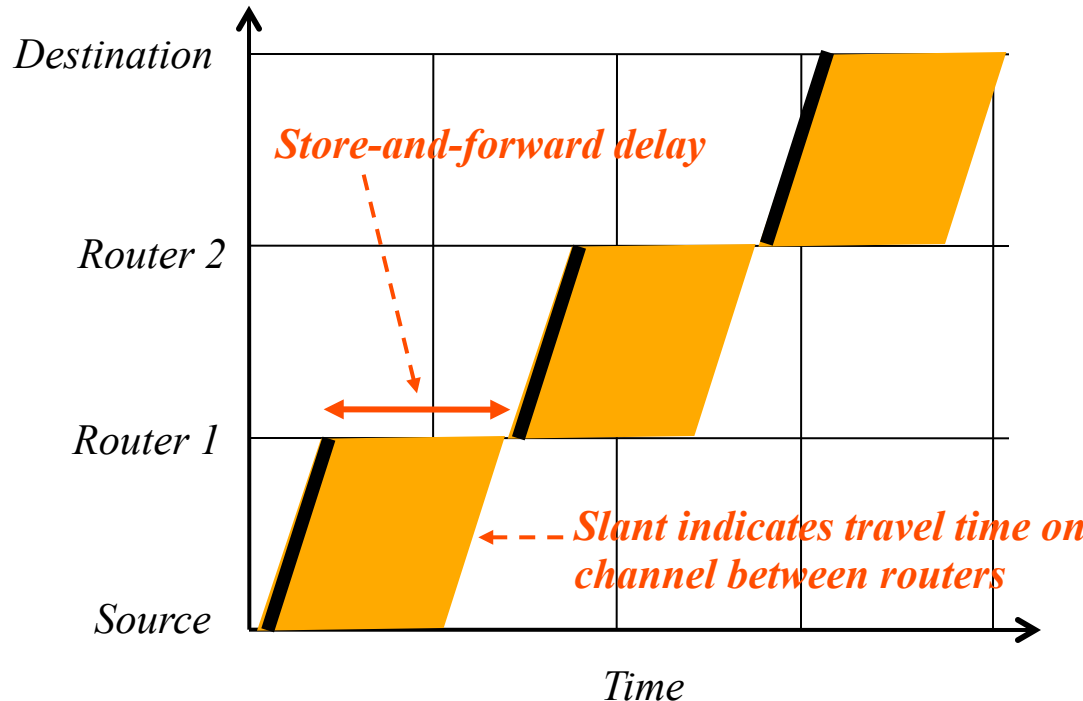
Packet-switching advantages

- Hosts with nothing to send send nothing
- Links are statistically multiplexed
- Easy for hosts to multiplex packets to many destinations
- Source and destination can operate at different bit-rates, as can intervening links

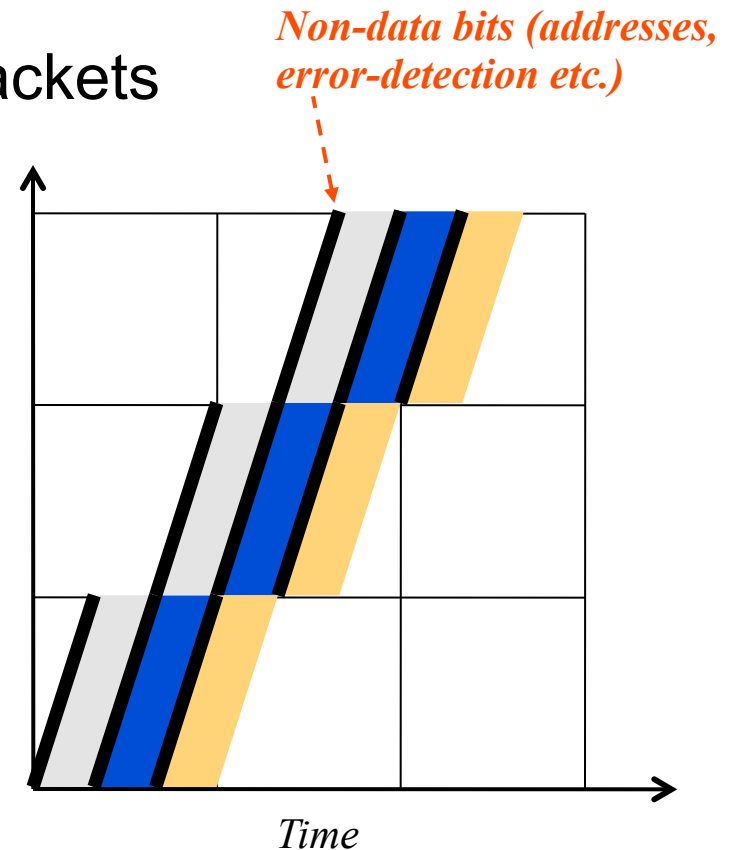


Packet-switching disadvantages (1)

- Store and forward delay
 - Can be minimised by small packets



a) A single large packet traversing 3 links with the same channel data rate.



b) Three small packets (containing the same data) traversing the same 3 links

Packet-switching disadvantages (2)

- Switches may become congested
 - Queues at packet switches
 - Queuing delay
 - Discard packets
- End-to-end delay is unpredictable
 - Jitter (variable delays in transmission time due to waiting in queues)
- Packet header overhead
 - Sequence numbers
 - Addresses
 - Connection identifiers

Packet variants

- Cell-switching
 - Cells: small, fixed-size packets (relatively large overhead)
 - Used in Asynchronous Transfer Mode (ATM) networks, 53-byte cell
 - Now used mainly inside switches (gigabit ethernet hubs)
- Frame-relay, frame-switching
 - Frames: unit of data transferred in data-link layer
 - Particular technologies transfer frames through frame-switches
- Message-switching
 - Email messages, files
 - Very large store and forward delay

Circuits or Packets?

	Circuits	Packets
1. Any host can communicate with any other host;	1: Yes	1: Yes
2. A host can only communicate with a small, fixed number of hosts at a time;	2: Yes	2: No
3. Other hosts may be "blocked" - unable to communicate because the wire they need is in use;	3: Yes	3: No
4. Communication between two hosts (once established) is unaffected by activity elsewhere in the network.	4: Yes	4: No
5. Whilst they are communicating, the two hosts have exclusive use of the sub-channel which links them;	5: Yes	5: No
6. Bits are delivered in the same order in which they are sent.	6: Yes	6: ?
7. Time intervals between bits are preserved	7: Yes	7: No
8. We can distinguish three phases in the life of a circuit; setting it up, using it and tearing it down	8: Yes	8: ?

For packet-switching, 2 and 3 are advantages; 4, 5 and 7 are potential disadvantages; 6 and 8 depend on whether we use “virtual circuits” or “datagrams” (to be discussed).

Sharing Summary

- Sharing links and networks
- Circuit-switching
 - Fixed allocations – FDM, TDM
- Packet-switching
 - Statistical multiplexing

The End