



## Assignment Submission

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**Subject:** Computer Networks

**Class:** IT Morning

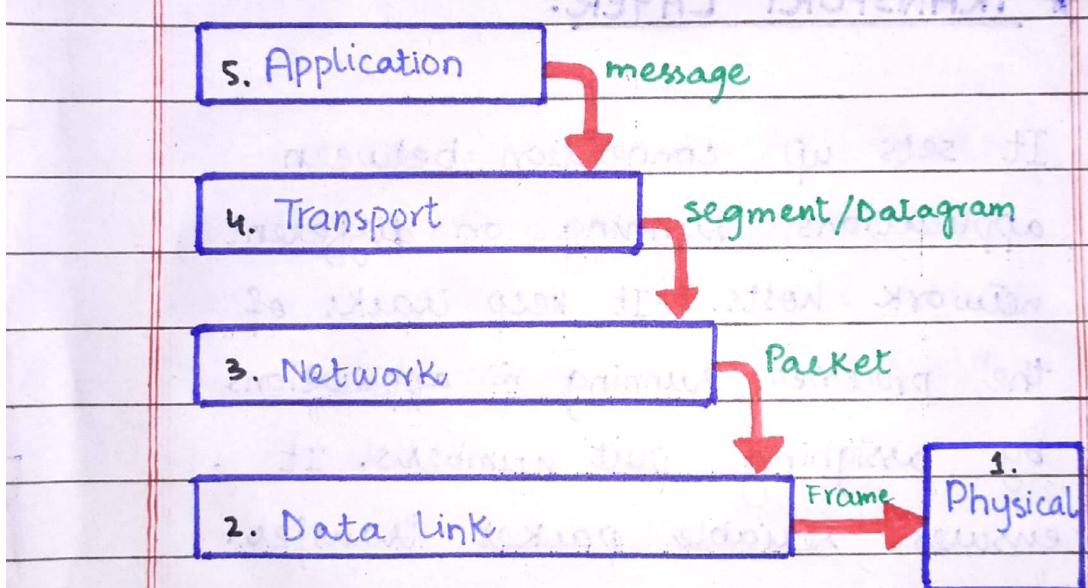
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Q No1:

## INTERNET ARCHITECTURE

DIAGRAM:



### \* APPLICATION LAYER:

It is a layer which contains all those applications that require standard internet/network connection.

These applications use Transport layer to send requests to connect themselves with remote host networks.

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## Protocols:

- HTTP (Hypertext Transfer Protocol)
- SMTP (Simple Mail Transfer Protocol)

## \* TRANSPORT LAYER:

It sets up connection between applications running on different network hosts. It keeps track of the processes running in applications.

by assigning port numbers. It ensures reliable packet transfer.

## Protocols:

- TCP (Transmission Control Protocol)
- UDP (User Datagram Protocol)

## \* NETWORK LAYER:

It is responsible for creation of packets that move across networks.

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### Protocols:

- IP (Internet Protocol)
- ARP (Address Resolution Protocol)

### \* DATA LINK LAYER:

It creates frames which enfolds packets and moves across the network host. It uses MAC address to identify source and destination of packets.

(A)

### Protocols:

- Token ring
- PPP (point to point protocol)
- Ethernet

### \* PHYSICAL LAYER:

It encodes/decodes the bits of data in frames and select the

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hardware that drives and receives signals on the network.

### Protocols:

- IEEE 802.11 (DATA LINK LAYER)
- USB 3.0 (Universal Serial Bus)
- Bluetooth

### Q No 3:

(A)

- (a) Coaxial cable
- (b) Ethernet cable
- (c) Optical Fiber

(B)

Coaxial cable and Ethernet cable are made of copper.

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Q NO 4:

## PROTOCOL:

A Protocol is a set of rules or procedures which are used for transmission of data and for communication between electronic devices (such as computers).

It is a format in which messages are transformed between sender and receiver.

## LIST OF PROTOCOLS & THEIR USES:

Name of Protocol	Uses
------------------	------

- (1) Transmission control Protocol (TCP)
- used for communication over a network host.
  - divides messages into a series of packets

to send from source  
to destination.

### (2) Internet Protocol (IP)

→ IPs are assigned to devices which help packets route through different nodes and reach the destination.

### (3) Post Office Protocol (POP)

→ used for receiving incoming Emails.

### (4) File Transfer Protocol (FTP)

→ allows user to transfer files from one machine to other machines.  
Files could be programs, text files, documents, multimedia, etc.

Q No9:

## AVERAGE THROUHPUT

Throughput refers to the rate at which packets of data are transmitted over a communication channel. To measure the total amount of time in which packets transferred, difference between the timestamps between first and last packet sent, are taken.

### Formula:

Suppose, packets of 'R' bits are transmitted to a host network in total time 'T'.

The average throughput will be calculated as:

$$\text{Average Throughput} = \frac{R}{T} \text{ (bits/sec)}$$

Unit: bps (bits per second)

Q No7:

## NETWORK LATENCY

### Definition:

- The time taken by a data packet to be transmitted from its source to destination, is called Network Latency. In technical words, we say the round trip time (RTT) from browser to server.
- The desired value of network latency is approximate to zero '0'.

### Causes:

The causes of network latency are:

- 1) low space in memory

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- 2) Presence of multiple routers
- 3) Poorly optimized backend database

## How to MEASURE?

Latency can be measured in:

1. Round Trip Time (RTT)
2. Time to first byte (TTFB)

### RTT:

RTT means measurement of time taken when a client sends a request to server and server gives response to client.

### TTFB:

TTFB means measurement of time taken when a client

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sends a request to sever  
and sever gives first byte  
of data to client.

## Factors:

A number of factors that contribute to network latency are:

- \* Propagation Delay
- \* Routing and switching
- \* Queuing and Buffering

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Q No2:

## CIRCUIT SWITCHING VS. PACKET SWITCHING

- |  |   |
|--|---|
| <ul style="list-style-type: none"><li>• data unit knows the entire path (address) by source.</li><li>• more reliable.</li><li>• data processed at source system only.</li><li>• more wastage of resources.</li><li>• not a store and forward technique.</li><li>• Examples:<ul style="list-style-type: none"><li>→ Optical Mesh</li><li>→ Analog Telephone</li></ul></li></ul> | <ul style="list-style-type: none"><li>• data unit knows only final destination decided by routers.</li><li>• less reliable.</li><li>• data processed at all intermediate nodes.</li><li>• less wastage of resources.</li><li>• Store and forward technique.</li><li>• Examples:<ul style="list-style-type: none"><li>→ School web servers</li></ul></li></ul> |
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## SUITABLE SITUATION FOR CIRCUIT SWITCHING:

Circuit switching is favourable in situations where connection is required for continuous long period of time, such as long distance communication.

## WHY INTERNET USE PACKET SWITCHING?

It is because packet switching is efficient enough to find its destination using routers when packets links are lost.

Also, it provide less wastage of bandwidth.

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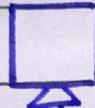
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## NUMERICALS

Q

No 10:

Host A



Host B



$$R_1 = 200 \text{ kbps}$$
$$D_1 = 15 \text{ Km}$$

$$R_2 = 5 \text{ Mbps}$$
$$D_2 = 100 \text{ Km}$$

Router

Given Data:

Propagation speed =  $s = 3 \times 10^8 \text{ m/sec}$

Processing delay =  $D_{proc} = 20 \text{ m/sec}$

Queuing delay =  $D_{queue} = 25 \text{ m/sec}$

(a) When one packet sent;

File size =  $L = 1500 \text{ bytes}$

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Average throughput = 200 kbps

Here we have assumed  
that packets are being  
transferred in "back-to-back"  
manner.

## Total end-to-end Delay:

As we know,

$$\text{avg delay} = \frac{1}{\text{avg throughput}}$$

$$D_{\text{Total}} = D_{\text{prop}} + D_{\text{trans}} + D_{\text{proc}} + D_{\text{queue}}$$

- First, we'll calculate:  $D_{\text{prop}}$

$$D_{1\text{ prop}} = \frac{d_1}{c} = \frac{15 \times 10^3}{3 \times 10^8} = 5 \times 10^{-5} \text{ m/sec}$$

$$D_{2\text{ prop}} = \frac{d_2}{c} = \frac{100 \times 10^3}{3 \times 10^8} = 3.33 \times 10^{-4}$$

$$D_{\text{total prop}} = \frac{d_1}{c} + \frac{d_2}{c}$$

$$= 5 \times 10^{-5} + 3.33 \times 10^{-4}$$

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$$D_{\text{total prop.}} = 3.83 \times 10^{-4} \text{ m/sec}$$

- Now, we'll calculate  $D_{\text{trans}}$ :

$$D_{\text{Trans}} = \underline{L}$$

(these are also R<sub>in</sub> and R<sub>out</sub> just now) (d)

$$= \underline{1500}$$

$$D_{\text{prop}} = \frac{200 \times 10^3}{1500} = 0.133 \text{ mm}$$

adjusted

$$D_{\text{Trans}} = 7.5 \times 10^{-3} \text{ sec}$$

- Put values to calculate  $D_{\text{total}}$ :

$$D_{\text{total}} = 3.83 \times 10^{-4} + 7.5 \times 10^{-3} + 20 \times 10^{-3}$$
$$+ 25 \times 10^{-3}$$

$$D_{\text{total}} = 0.052883$$

## Throughput:

$$\text{Throughput} = \frac{\underline{L}}{D_{\text{total}}}$$

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$$\text{Throughput} = \frac{1500}{0.052883}$$

$$\boxed{\text{Throughput} = 28364.50277 \text{ bps}}$$

(b) When four packets are sent;

$$\text{File size} = L = 1500 \times 4 = 6000 \text{ bytes}$$

Total end-to-end Delay:

$$D_{\text{total}} = D_{\text{prop}} + D_{\text{trans}} + D_{\text{proc}} + D_{\text{queue}}$$

$$D_{\text{trans}} = \frac{L}{R_s} = \frac{6000}{200 \times 10^3}$$

$$\boxed{D_{\text{trans}} = 0.03}$$

We have already calculated

$D_{\text{prop}}$  in part (a).

Put in  $D_{\text{total}}$ .

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$$D_{\text{Total}} = 3.83 \times 10^{-4} + 0.03 + 20 \times 10^{-3}$$
$$+ 25 \times 10^{-3}$$

$$D_{\text{Total}} = 0.075383$$

## Throughput:

$$\text{Throughput} = \frac{L}{D_{\text{Total}}}$$

$$= \frac{6000}{0.075383}$$

$$\text{Throughput} = 79593.54 \text{ bps}$$

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Q No 8:

GIVEN DATA:

Service Time = 10ms (sum of queue time and processing time)

Packet size. =  $L$  = 2000 bits

Propagation delay =  $D_{prop}$  = 150ms

Latency =  $D_{Total}$  = 200 ms

Minimum capacity =  $R$  = ?

SOLUTION:

As we have formula:

$$D_{Total} = D_{prop} + D_{Trans} + (D_{proc} + D_{queue})$$

$$D_{Total} = D_{prop} + \frac{L}{R} + \text{service time}$$

$$200 = 150 + \frac{2000}{R} + 10$$

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$$200 = 160 + \frac{2000}{R}$$

$$40 = \frac{2000}{R}$$

$$R = \frac{2000}{40}$$

$$R = 50 \text{ ms}$$

So, minimum capacity of link connecting two hosts is 50ms.

Q NOS:

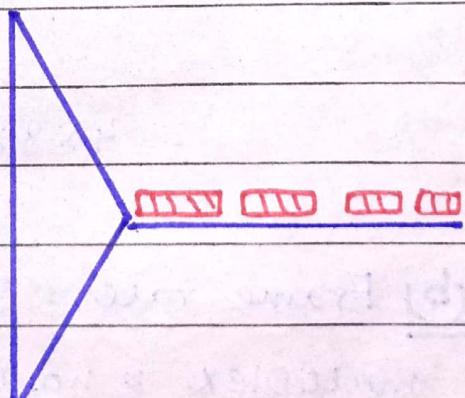
1. 1 2 3 4 5 6 7 8

2. 1 2 3 4 5 6 7 8

3. 1 2 3 4 5 6 7 8

⋮ ⋮

8. 1 2 3 4 5 6 7 8



(a) Frame rate = 4 F/s

multiplex = ?

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$$\text{Frame rate} = \frac{\text{no. of channels}}{\text{multiplex}}$$

$$\text{multiplex} = \frac{\text{no. of channel}}{\text{Frame rate}} = \frac{8}{4}$$

$$\text{multiplex} = 2$$

• data rate of single TDM channel =

$$\frac{\text{no. of frame} \times \text{single slot size}}{8 \times 8 \times 8}$$

$$= 4 \times 8 = 32 \text{ bits/sec.}$$

• data rate of TDM line =

$$\frac{\text{no. of slots} \times \text{single slot size}}{8 \times 8 \times 8}$$

$$= 4 \times 8 \times 8 = 256 \text{ bits/sec.}$$

(b) Frame rate = 8 F/s

$$\text{multiplex} = \frac{\text{no. of channel}}{\text{Frame rate}} = \frac{8}{8} = 1$$

• data rate of single TDM channel =

$$8 \times 8 = 64 \text{ bits/sec}$$

- data rate of TDM line =  $8 \times 8 \times 8$   
= 512 bits/sec

(c)

$$\text{Frame rate} = 4000 \text{ F/S}$$
$$\text{multiplex} = \frac{8}{4000} = 2 \times 10^{-3}$$

- data of single TDM channel =

$$4000 \times 8$$
$$= 32000 \text{ bits/sec}$$

- data rate of TDM line =  $4000 \times 8 \times 8$   
= 256000  
bits/sec

(d) Frame rate = 4000 F/s

$$\text{Data rate in TDM Line} = 8 \times 8 \times 4000$$
$$= 256000 \text{ bits/sec}$$

$$\text{Data size} = 6000 \text{ KB}$$
$$= 6000 \times 1000 \times 8 = 48000000$$
$$\text{bits}$$

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$$\text{Time required} = \frac{48000000}{256000}$$

$$= 187.5 \text{ sec}$$

$$= 187.5 / 60 \text{ min } (2)$$

$$\boxed{\text{Time} = 3.125 \text{ mins}}$$

Q No 6:

$$\text{Link rate / multiplex} = 1.68 \text{ Mbps}$$

$$= 1.68 \times 10^6 \text{ bps}$$

$$\text{Individual session rate} = 28 \text{ Kbps}$$

$$= 28 \times 10^3 \text{ bps}$$

$$\text{Total no. of sessions} = \frac{1.68 \times 10^6}{28 \times 10^3}$$

Total no. of sessions = 60 Links can carry.

$$L = 3 \text{ MB} = (3 \times 10^6 \text{ B} \times 8) \text{ bits}$$

$$D_{\text{total}} = ?$$

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SOLUTION:

$$D_{\text{Total}} = D_{\text{prop}} + D_{\text{Trans}} + D_{\text{proc}} + D_{\text{queue}}$$

Neglecting  $D_{\text{prop}}$ ,  $D_{\text{proc}}$  &  $D_{\text{queue}}$

So,

$$D_{\text{Total}} = D_{\text{Trans}}$$

$$\begin{aligned} D_{\text{Total}} &= \frac{L}{R} \\ &= \frac{3 \times 10^6 \times 8}{28 \times 10^3} = \frac{24000000}{28 \times 10^3} \\ &= 857.14 \text{ s} \end{aligned}$$

But, it also takes  
150 ms before establishing  
circuit.

So,

$$D_{\text{Total}} = 857.14 + 150 \times 10^{-3}$$

$$D_{\text{Total}} = 857.29 \text{ s}$$