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HW4
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The Tiling Game

1.

```
#include <stdio.h>
#include <stdlib.h>
#include <time.h>
#include <math.h>
//gcc --std=c99 -o hw4-1 hw4-1.c
//./hw4-1.c n //(2^n)

struct tri{
    int x,y;
};
//choose the X location randomly
struct tri *get_x(int side){
    printf("Matrix:%d*%d\n",side,side);
    srand(time(NULL));
    struct tri *tri = malloc(sizeof(struct tri));
    tri->x = rand()%side;
    tri->y = rand()%side;
    printf("X is in [%d, %d] \n", tri->x, tri->y);
    return tri;
}

//create a array to show the game's board
char **getborad(struct tri *tri,int side){
    int i,j;
    printf("Get Board \n");
    char **board=malloc(sizeof(char*)*side);
    for(int i=0;i<side;i++){
        board[i] = malloc(sizeof(char)*side);
    }
    for(i=0;i<side;i++)
    {
        for(j=0;j<side;j++)
        {
            if(i == tri->x && j==tri->y)
            {
                board[i][j] = 'X';
                printf("%c\t",board[i][j]);
            }
            else
            {
                board[i][j] = '_';
                printf("%c\t",board[i][j]);
            }
        }
        printf("\n");
    }
}
```

```

printf("\n");
return board;
}
//start to solve problem
int put_tri(char **board,int x,int y,int n,int side,int markx,int marky){
    int center=side/2;
    int sx=x+center-1;
    int bx=x+center;
    int sy=y+center-1;
    int by=y+center;
    if (side == 2)
    {
        board[x][y] = n;
        board[x][y + 1] = n;
        board[x + 1][y] = n;
        board[x + 1][y + 1] = n;
        board[markx][marky] = -1;
        return n += 1;
    }
    else{
        if (markx < bx && marky < by)//bottom right
        {
            n = put_tri(board,x,y,n,center,markx,marky);
            n = put_tri(board,bx,y,n,center,bx,sy);
            n = put_tri(board,bx,by,n,center,bx,by);
            n = put_tri(board,x,by,n,center,sx,by);

            board[sx][by] = n;
            board[bx][by] = n;
            board[bx][sy] = n;
        }
        else if (markx >= bx && marky < by)//top right
        {
            n = put_tri(board,x,y,n,center,sx,sy);
            n = put_tri(board,bx,y,n,center,markx,marky);
            n = put_tri(board,bx,by,n,center,bx,by);
            n = put_tri(board,x,by,n,center,sx,by);

            board[sx][sy] = n;
            board[bx][by] = n;
            board[sx][by] = n;
        }
        else if (markx >= bx && marky >= by)//top left
        {
            n = put_tri(board,x,y,n,center,sx,sy);
            n = put_tri(board,bx,y,n,center,bx,sy);
            n = put_tri(board,bx,by,n,center,markx,marky);
            n = put_tri(board,x,by,n,center,sx,by);
            board[sx][by] = n;
            board[sx][sy] = n;
            board[bx][sy] = n;
        }
        else if (markx < bx && marky >= by)//bottom left
        {

```

```

        n = put_tri(board,x,y,n,center,sx,sy);
        n = put_tri(board,bx,y,n,center,bx,sy);
        n = put_tri(board,bx,by,n,center,bx,by);
        n = put_tri(board,x,by,n,center,markx,marky);
        board[sx][sy] = n;
        board[bx][by] = n;
        board[bx][sy] = n;
    }
    return n += 1;
}
}
//print the array to show the result
void print_board(char **board,int side){
    int i, j;
    for (i = 0; i < side; i++)
    {
        for (j = 0; j < side; j++)
        {
            if (board[i][j] != -1)
                printf("%d\t", board[i][j]);
            else
                printf("x\t");
        }
        printf("\n");
    }
    for (i = 0; i < side; i++)
    {
        free(board[i]);
    }
    free(board);
}
int main(int argc,char* argv[])
{
    clock_t start_time, end_time;
    float total_time = 0;
    if(argc == 2)
    {
        int n = atoi(argv[1]);
        start_time = clock();
        int side=pow(2,n);
        struct tri *tri=get_x(side);
        char **board= getborad(tri,side);
        put_tri(board, 0, 0, 1,side, tri->x, tri->y);
        print_board(board, side);
        end_time = clock();
        total_time = (float)(end_time - start_time)/CLOCKS_PER_SEC;
        printf("Board size %d * %d spend %f sec\n",side,side,total_time);
    }
    else
    {
        printf("please enter one argument.\n");
    }
    return 0;
}

```

2.
8*8

```

[us-MacBook-Pro:Desktop roy1u$ ./hw1-1 3
Matrix:8*8
X is in [7, 5]
Get Board
- - - - -
- - - - -
- - - - -
- - - - -
- - - - -
- - - - -
- - - - -
- - - - -
1 1 4 4 16 16 19 19
1 5 5 4 16 20 20 19
2 5 3 3 17 17 20 18
2 2 3 21 21 17 18 18
6 6 9 21 11 11 14 14
6 10 9 9 11 11 15 14
7 10 10 8 12 12 15 13
7 7 8 8 12 x 13 13

```

16*16(if the picture isn't clear, I have upload those picture)

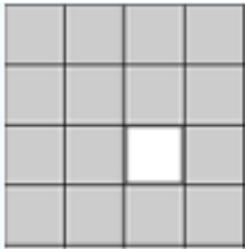
```

[us-MacBook-Pro:Desktop roy1u$ ./hw1-1 4
Matrix:16*16
X is in [0, 7]
Get Board
- - - - - X - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
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- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
- - - - - - - - - - -
1 1 4 4 16 16 19 x 64 64 67 67 79 79 82 82
1 5 5 4 16 20 20 19 64 68 68 67 79 83 83 82
2 5 3 3 17 20 20 18 65 68 66 66 80 80 81 81
2 2 3 21 17 17 18 65 65 66 66 81 81 81 81
6 6 9 21 21 11 14 69 69 72 72 84 84 77 77
6 10 9 9 11 11 15 69 73 73 72 74 74 78 77
7 10 10 8 12 15 15 70 70 73 71 75 78 78 76
7 7 8 8 12 15 15 70 71 71 71 75 75 76 76
22 22 25 25 37 37 40 85 85 43 46 46 58 58 61 61
22 26 26 25 37 41 40 40 43 43 47 46 58 62 62 61
23 26 24 24 38 41 41 39 44 47 47 45 59 59 62 60
25 25 24 42 38 38 39 39 44 44 45 45 63 63 60 60
27 27 30 42 32 35 35 48 48 51 63 63 53 53 56 56
27 31 30 30 32 32 35 35 48 52 51 51 53 53 57 56
28 31 31 29 33 36 36 34 49 52 52 50 54 57 57 55
28 31 29 29 33 33 34 34 49 49 50 50 54 54 55 55
[us-MacBook-Pro:Desktop roy1u$

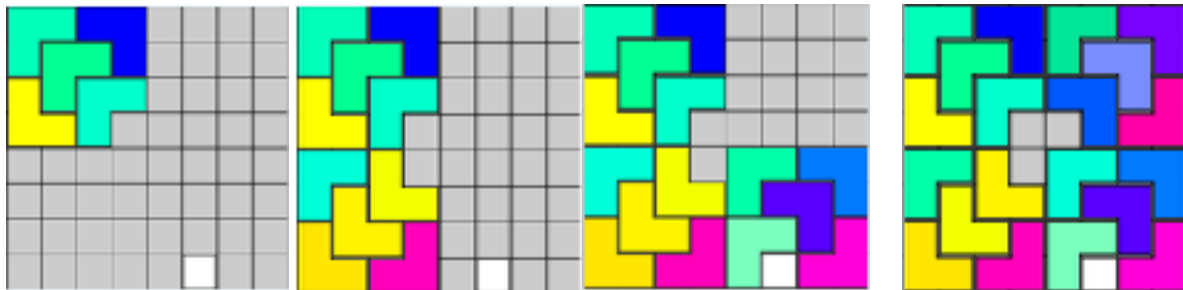
```

3. assume the missing cell is in the [3,2]

Base case: If we have a $2^k \times 2^k$ boardsize, in our algorithm we can divide the board into a lot of 4×4 square like this



Then we can start to fill it with a 2×2 square by the order **top left, bottom left, bottom right, top right and the middle one**. In each step, we will make a cell to be -1 depending on the aspect of the missing cell.



Then we fill each 4×4 square in the same order, and finally we will get a middle empty cell which is facing to the missing cell.

4.

Assume n is side length

$$T(1) = 1$$

$$T(n) = 4T(n/2) + c$$

$$= 4(4T(n/4) + c) + c$$

$$= \dots$$

$$= 4^k T(n/2^k) + c$$

$$= 4^k T(1) + c \text{ where } k = \log_2 n$$

$$= 2^{(2 \log_2 n)} + c$$

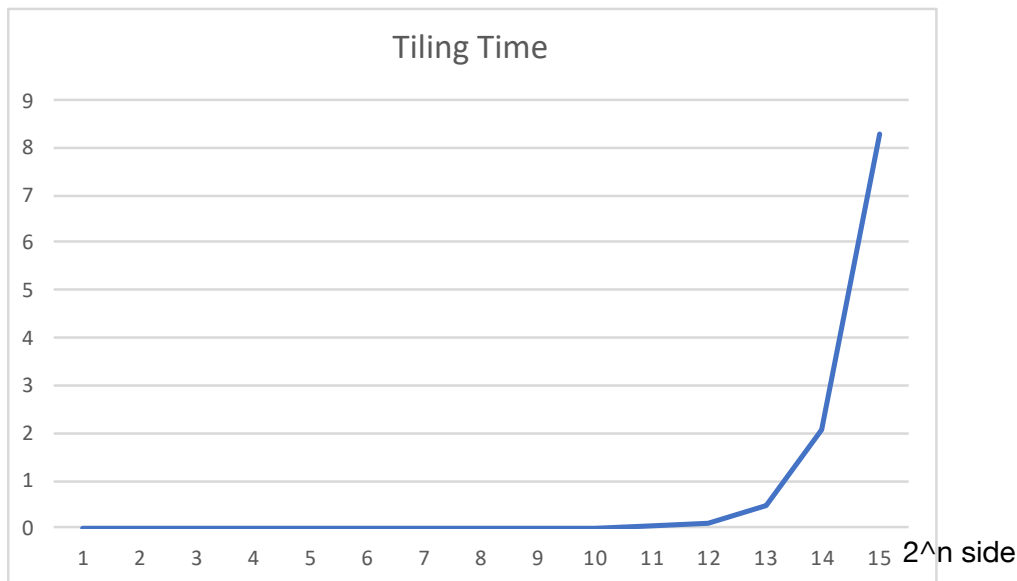
$$= an^2 + c$$

Therefore:

$$T(n) = O(n^2)$$

5.

```
Board size 2 * 2 spend 0.000039 sec
Board size 4 * 4 spend 0.000046 sec
Board size 8 * 8 spend 0.000044 sec
Board size 16 * 16 spend 0.000047 sec
Board size 32 * 32 spend 0.000056 sec
Board size 64 * 64 spend 0.000090 sec
Board size 128 * 128 spend 0.000228 sec
Board size 256 * 256 spend 0.000654 sec
Board size 512 * 512 spend 0.002607 sec
Board size 1024 * 1024 spend 0.009477 sec
Board size 2048 * 2048 spend 0.031669 sec
Board size 4096 * 4096 spend 0.121868 sec
Board size 8192 * 8192 spend 0.485723 sec
Board size 16384 * 16384 spend 2.077413 sec
Board size 32768 * 32768 spend 8.291318 sec
```



We can ignore c because c is very small

From side 2048 to 32768

$$0.031669 = a \cdot 2048 \cdot 2048 \Rightarrow a \approx 7.55 \cdot 10^{-9}$$

$$0.121868 = a \cdot 4096 \cdot 4096 \Rightarrow a \approx 7.26 \cdot 10^{-9}$$

$$0.485723 = a \cdot 8192 \cdot 8192 \Rightarrow a \approx 7.23 \cdot 10^{-9}$$

$$2.077413 = a \cdot 16384 \cdot 16384 \Rightarrow a \approx 7.73 \cdot 10^{-9}$$

To predict the running time of $32768 \cdot 32768$, the I compute the board get the time almost same as predicted time.

$$a \cdot 32768 \cdot 32768 = 8.291318 \text{ sec}$$

Find the ODD Coin

1.

I will divide the coins into three group to compare.

```
#include <stdio.h>
#include <stdlib.h>
#include <time.h>
#include <math.h>
//gcc --std=c99 -o hw4_2 hw4_2.c
//./hw4_2.c

int num=0;
int get_x(int n){
    srand(time(NULL));
    int x=rand()%n;
    return x;
}
void sortcoin(int arr[],int n,int x){
    for(int i=0;i<n;i++){
        if(i==x) arr[i]=0;
        else arr[i]=1;
    }
}
int find_x(int arr[],int s,int n,int weighting){
    int a=((n-s)/3);
    int i,g1=0,g2=0,g3=0;
    for(i=s;i<a+s;i++){
        g1=g1+arr[i];
    }
    for(i=a+s;i<(2*a)+s;i++){
        g2=g2+arr[i];
    }
    for(i=(2*a)+s;i<n;i++){
        g3=g3+arr[i];
    }
    weighting++;
    if((n-s)>3){
        if(g1==g2){
            find_x(arr,s+(2*a),n,weighting);
        }
        else if(g1==g3){
            weighting++;
            find_x(arr,s+a,s+(2*a),weighting);
        }
        else{
            find_x(arr,s,s+a,weighting);
        }
    }
    else{
        if(arr[s]==arr[s+1]) num= n;
        else if(arr[s]==arr[n]) num= s+2;
        else num= s+1;
        printf("the odd coin is the %dth of the sort\n",num);
        printf("the number of weighting %d\n",weighting);
    }
}
```

```

    }
}
int main(int argc, char* argv[])
{
    clock_t start_time, end_time;
    float total_time = 0;
    int n, weighting=0;
    printf("enter a number");
    scanf("%d", &n);
    int m=pow(3,n);
    start_time = clock(); /* mircosecond */
    int arr[m];
    int x=get_x(m);
    sortcoin(arr,m,x);
    for(int i=0;i<m;i++){
        printf("%d ", arr[i]);
    }
    printf("\n");
    find_x(arr,0,m,weighting);
    end_time = clock();
    total_time = (float)(end_time - start_time)/CLOCKS_PER_SEC;
    printf("%f", total_time);
    printf("\n");
    return 0;
}

```

2.

- a. key in a 3^n number
- b. Put it in array and start to find
- c. Divide it into three group and compare the weight of three groups, find the different one and take it to redo step c
- d. If there is only three coin in group, compare those three coin, the different one is odd coin

3.

Number of weightings

For 3 coins, there are three possible solution, so the mean number of weightings is $(1+2+2)/3=5/3$

$W(1)=0$

$W(3)=5/3$

$W(n)=W(n/3)+5/3=W(n/3^k)+5k/3=W(1)+5k/3$ where $k=\log_3 n$

$\Rightarrow W(n)=5\log_3 n/3$

Running time

Assume n is the power of 3^n coin

$T(1)=a$

$T(n)=T(n/3)+n+c=[T(n/9) + n/3 + c] + n + c = [T(n/27) + (n/9 + n/3 + n)]=\dots$

$=T(n/3^k) + n[1-(1/3)^k]/(1-1/3) + kc$

$=T(1) + n[1-(1/3)^k]/(1-1/3) + kc$ where $k=\log_3 n$

$=T(1)+1.5(n-1)+c\log_3 n$

$=O(n)$

4.

I run 9 times of 27 coins

```
1 1 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
the odd coin is the 5th of the sort
the number of weighting 4

1 1 1 1 1 1 1 1 1 1 1 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
the odd coin is the 15th of the sort
the number of weighting 5

1 1 1 1 1 1 1 1 1 1 1 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
the odd coin is the 14th of the sort
the number of weighting 5

1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 1
the odd coin is the 26th of the sort
the number of weighting 3

1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 1 1 1 1 1 1 1 1 1
the odd coin is the 22th of the sort
the number of weighting 4

1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 1 1 1 1 1 1 1 1 1 1 1
the odd coin is the 21th of the sort
the number of weighting 3

1 1 1 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
the odd coin is the 7th of the sort
the number of weighting 3

1 1 1 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
the odd coin is the 6th of the sort
the number of weighting 4

1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1
the odd coin is the 19th of the sort
the number of weighting 3
```

5.(if the picture isn't clear, I have upload those picture)

```

number of weighting 2;3 coins cost 0.000003 sec
number of weighting 1;3 coins cost 0.000003 sec
number of weighting 2;3 coins cost 0.000002 sec
number of weighting 1;3 coins cost 0.000002 sec
number of weighting 2;3 coins cost 0.000003 sec
number of weighting 1;3 coins cost 0.000002 sec
number of weighting 2;3 coins cost 0.000001 sec
number of weighting 2;3 coins cost 0.000001 sec
number of weighting 2;3 coins cost 0.000002 sec
number of weighting 2;3 coins cost 0.000003 sec
mean time:0.000002

number of weighting 3;9 coins cost 0.000003 sec
number of weighting 4;9 coins cost 0.000003 sec
number of weighting 4;9 coins cost 0.000003 sec
number of weighting 3;9 coins cost 0.000003 sec
number of weighting 4;9 coins cost 0.000003 sec
number of weighting 3;9 coins cost 0.000001 sec
number of weighting 3;9 coins cost 0.000001 sec
number of weighting 4;9 coins cost 0.000001 sec
number of weighting 2;9 coins cost 0.000001 sec
number of weighting 3;9 coins cost 0.000001 sec
mean time:0.000002

number of weighting 5;27 coins cost 0.000002 sec
number of weighting 5;27 coins cost 0.000002 sec
number of weighting 5;27 coins cost 0.000003 sec
number of weighting 6;27 coins cost 0.000003 sec
number of weighting 4;27 coins cost 0.000003 sec
number of weighting 6;27 coins cost 0.000003 sec
number of weighting 6;27 coins cost 0.000003 sec
number of weighting 5;27 coins cost 0.000003 sec
number of weighting 5;27 coins cost 0.000003 sec
number of weighting 6;27 coins cost 0.000003 sec
mean time:0.000003

number of weighting 8;81 coins cost 0.000004 sec
number of weighting 7;81 coins cost 0.000004 sec
number of weighting 8;81 coins cost 0.000003 sec
number of weighting 7;81 coins cost 0.000003 sec
number of weighting 7;81 coins cost 0.000004 sec
number of weighting 6;81 coins cost 0.000003 sec
number of weighting 7;81 coins cost 0.000002 sec
number of weighting 7;81 coins cost 0.000003 sec
number of weighting 6;81 coins cost 0.000003 sec
number of weighting 6;81 coins cost 0.000004 sec
mean time:0.000003

```

```

number of weighting 7;243 coins cost 0.000006 sec
number of weighting 7;243 coins cost 0.000006 sec
number of weighting 7;243 coins cost 0.000004 sec
number of weighting 9;243 coins cost 0.000005 sec
number of weighting 9;243 coins cost 0.000006 sec
number of weighting 8;243 coins cost 0.000005 sec
number of weighting 9;243 coins cost 0.000005 sec
number of weighting 7;243 coins cost 0.000006 sec
number of weighting 8;243 coins cost 0.000005 sec
number of weighting 10;243 coins cost 0.000006 sec
mean time:0.000005

number of weighting 10;729 coins cost 0.000012 sec
number of weighting 11;729 coins cost 0.000011 sec
number of weighting 10;729 coins cost 0.000011 sec
number of weighting 9;729 coins cost 0.000012 sec
number of weighting 10;729 coins cost 0.000011 sec
number of weighting 10;729 coins cost 0.000011 sec
number of weighting 11;729 coins cost 0.000010 sec
number of weighting 11;729 coins cost 0.000009 sec
number of weighting 11;729 coins cost 0.000010 sec
number of weighting 10;729 coins cost 0.000011 sec
mean time:0.000011

number of weighting 12;2187 coins cost 0.000027 sec
number of weighting 12;2187 coins cost 0.000028 sec
number of weighting 12;2187 coins cost 0.000028 sec
number of weighting 11;2187 coins cost 0.000027 sec
number of weighting 12;2187 coins cost 0.000032 sec
number of weighting 12;2187 coins cost 0.000031 sec
number of weighting 11;2187 coins cost 0.000029 sec
number of weighting 13;2187 coins cost 0.000041 sec
number of weighting 10;2187 coins cost 0.000028 sec
number of weighting 10;2187 coins cost 0.000028 sec
mean time:0.000038

number of weighting 13;6561 coins cost 0.000100 sec
number of weighting 14;6561 coins cost 0.000078 sec
number of weighting 12;6561 coins cost 0.000081 sec
number of weighting 12;6561 coins cost 0.000087 sec
number of weighting 11;6561 coins cost 0.000077 sec
number of weighting 14;6561 coins cost 0.000077 sec
number of weighting 12;6561 coins cost 0.000076 sec
number of weighting 12;6561 coins cost 0.000079 sec
number of weighting 13;6561 coins cost 0.000079 sec
number of weighting 12;6561 coins cost 0.000078 sec
mean time:0.000081

```

Coin 3 mean weighting:1.7

Coin 9 mean weighting:3.3

Coin 27 mean weighting:5.4

Coin 81 mean weighting:6.9

Coin 243 mean weighting:8.1

Coin 729 mean weighting:10.3

Coin 2187 mean weighting:11.5

Coin 6561 mean weighting:12.5

```

number of weighting 18;19683 coins cost 0.000272 sec
number of weighting 12;19683 coins cost 0.000191 sec
number of weighting 13;19683 coins cost 0.000190 sec
number of weighting 15;19683 coins cost 0.000181 sec
number of weighting 14;19683 coins cost 0.000180 sec
number of weighting 14;19683 coins cost 0.000181 sec
number of weighting 13;19683 coins cost 0.000180 sec
number of weighting 15;19683 coins cost 0.000180 sec
number of weighting 16;19683 coins cost 0.000180 sec
number of weighting 13;19683 coins cost 0.000187 sec
mean time:0.000192

number of weighting 17;59049 coins cost 0.000655 sec
number of weighting 18;59049 coins cost 0.000590 sec
number of weighting 16;59049 coins cost 0.000537 sec
number of weighting 17;59049 coins cost 0.000482 sec
number of weighting 19;59049 coins cost 0.000470 sec
number of weighting 16;59049 coins cost 0.000470 sec
number of weighting 17;59049 coins cost 0.000468 sec
number of weighting 17;59049 coins cost 0.000467 sec
number of weighting 19;59049 coins cost 0.000468 sec
number of weighting 16;59049 coins cost 0.000435 sec
mean time:0.000498

number of weighting 19;177147 coins cost 0.001590 sec
number of weighting 17;177147 coins cost 0.001221 sec
number of weighting 19;177147 coins cost 0.001180 sec
number of weighting 18;177147 coins cost 0.001152 sec
number of weighting 19;177147 coins cost 0.001157 sec
number of weighting 19;177147 coins cost 0.001160 sec
number of weighting 19;177147 coins cost 0.001190 sec
number of weighting 18;177147 coins cost 0.001107 sec
number of weighting 20;177147 coins cost 0.001081 sec
number of weighting 18;177147 coins cost 0.001077 sec
mean time:0.001188

number of weighting 20;531441 coins cost 0.004260 sec
number of weighting 22;531441 coins cost 0.003143 sec
number of weighting 21;531441 coins cost 0.003504 sec
number of weighting 21;531441 coins cost 0.003461 sec
number of weighting 20;531441 coins cost 0.003459 sec
number of weighting 22;531441 coins cost 0.003462 sec
number of weighting 19;531441 coins cost 0.003695 sec
number of weighting 20;531441 coins cost 0.003478 sec
number of weighting 21;531441 coins cost 0.003180 sec
number of weighting 18;531441 coins cost 0.003146 sec
mean time:0.003479

```

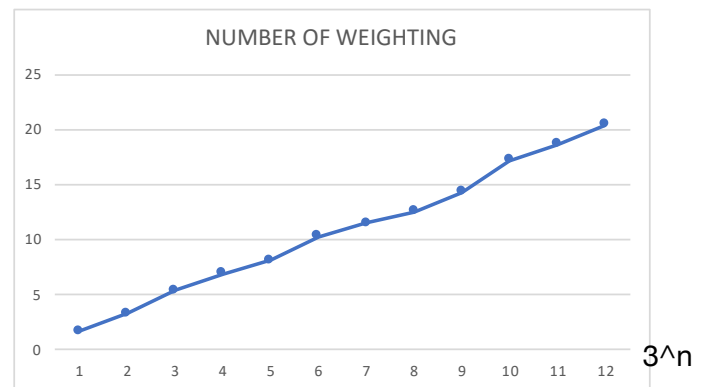
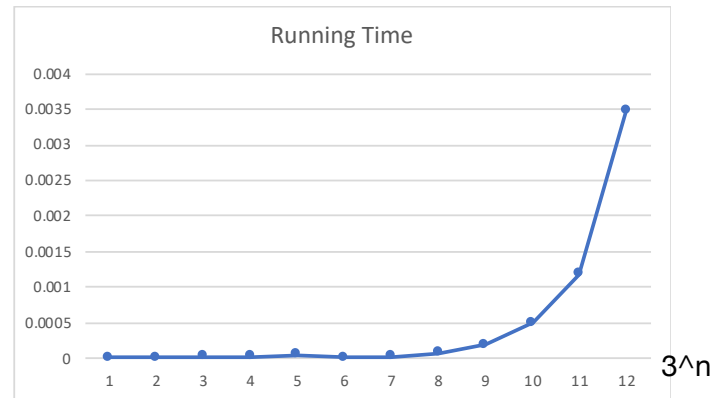
(if the picture isn't clear, I have upload those picture)

Coin 19683 mean weighting:14.3

Coin 59049 mean weighting:17.2

Coin 177147 mean weighting:18.6

Coin 531441 mean weighting:20.4



6.

It is not always take the same time and make the same number of weighings for each problem with n coins, but it is almost same. It is very easy to impact the result because of the limited size.

However, when we run a lot of time and take the mean, it can almost be a specific time for each number of coins.

Running time

The predictions from your difference equations is $T(n)=O(n)$

And we can easily to compare through the plot in Q5, the line is almost straight.

Number of weightings

Take 531441 for example

We predict

$$W(n)=5\log_3 531441/3=5*12/3=20$$

The result of mean weighting we running is 20.4

It is almost the same.