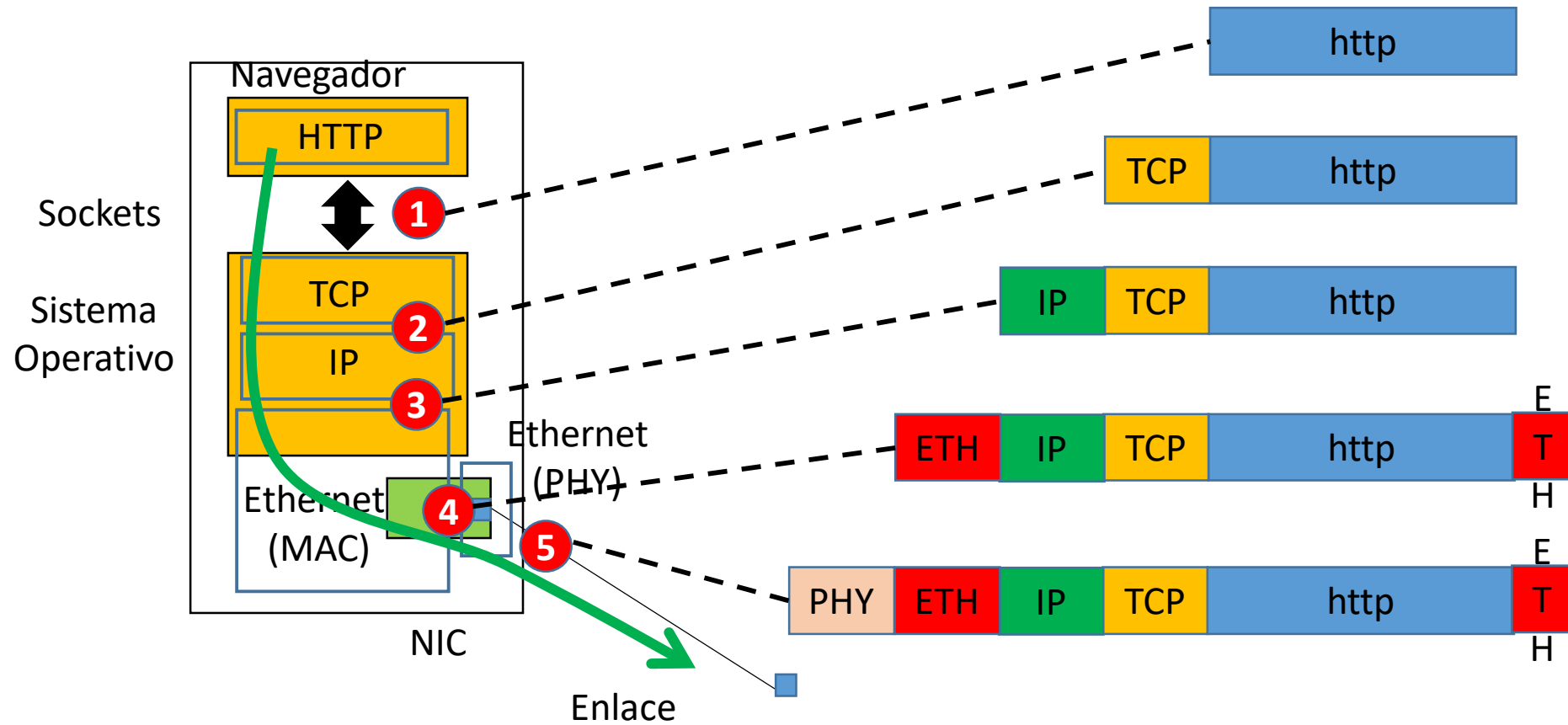
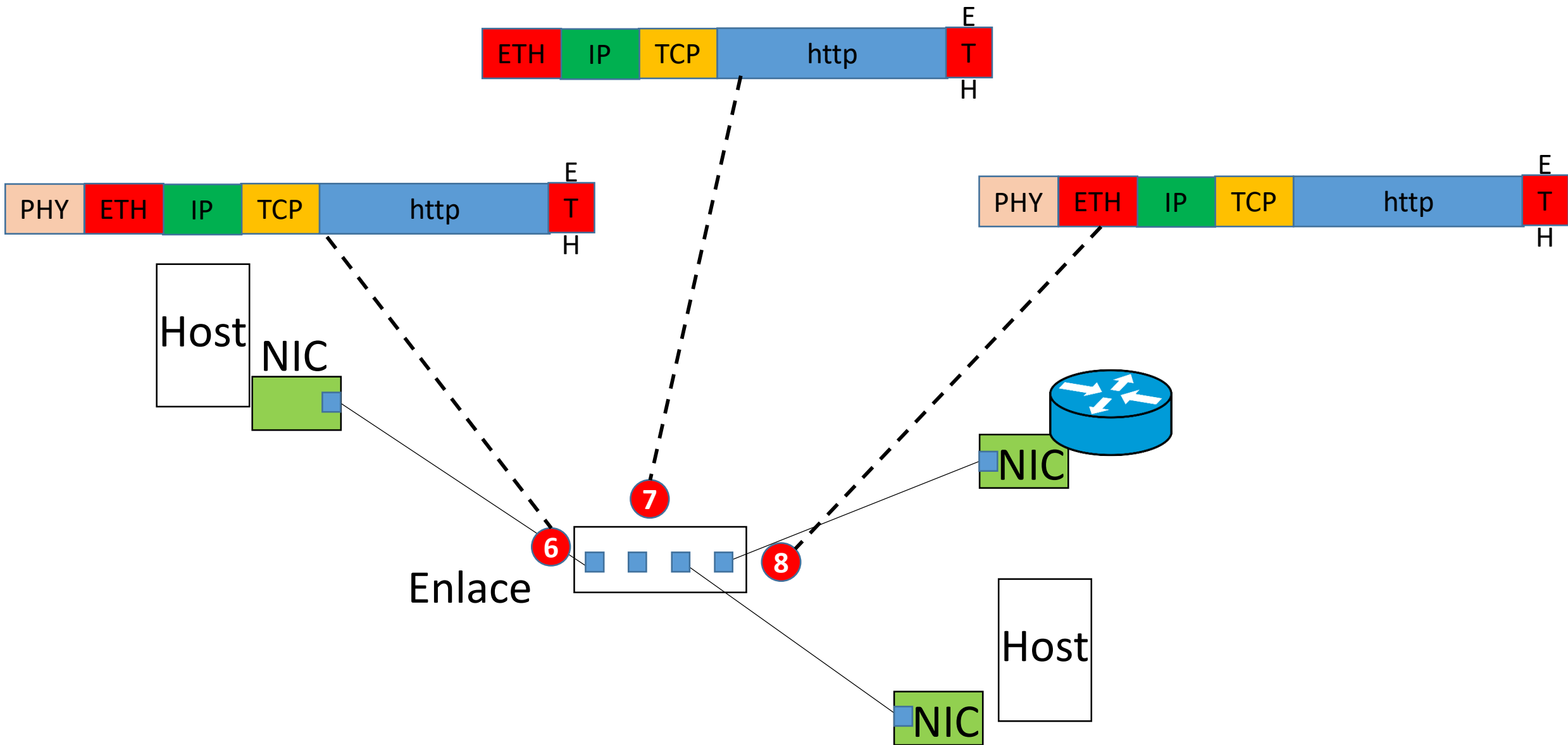
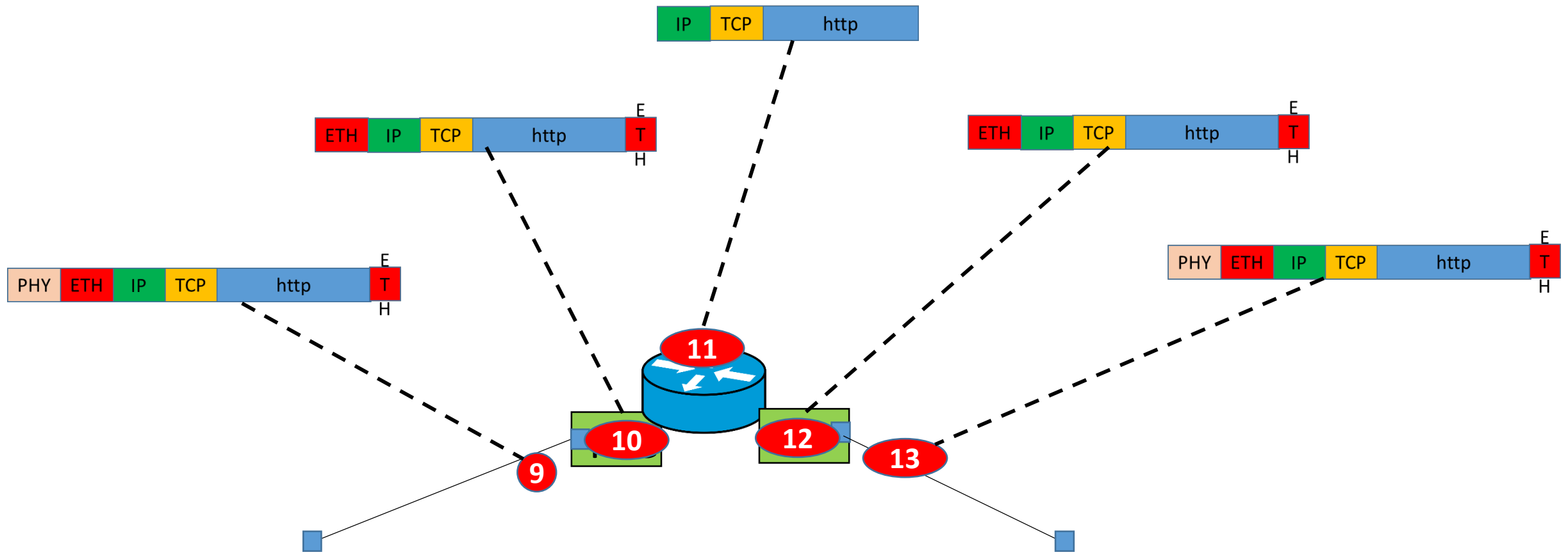


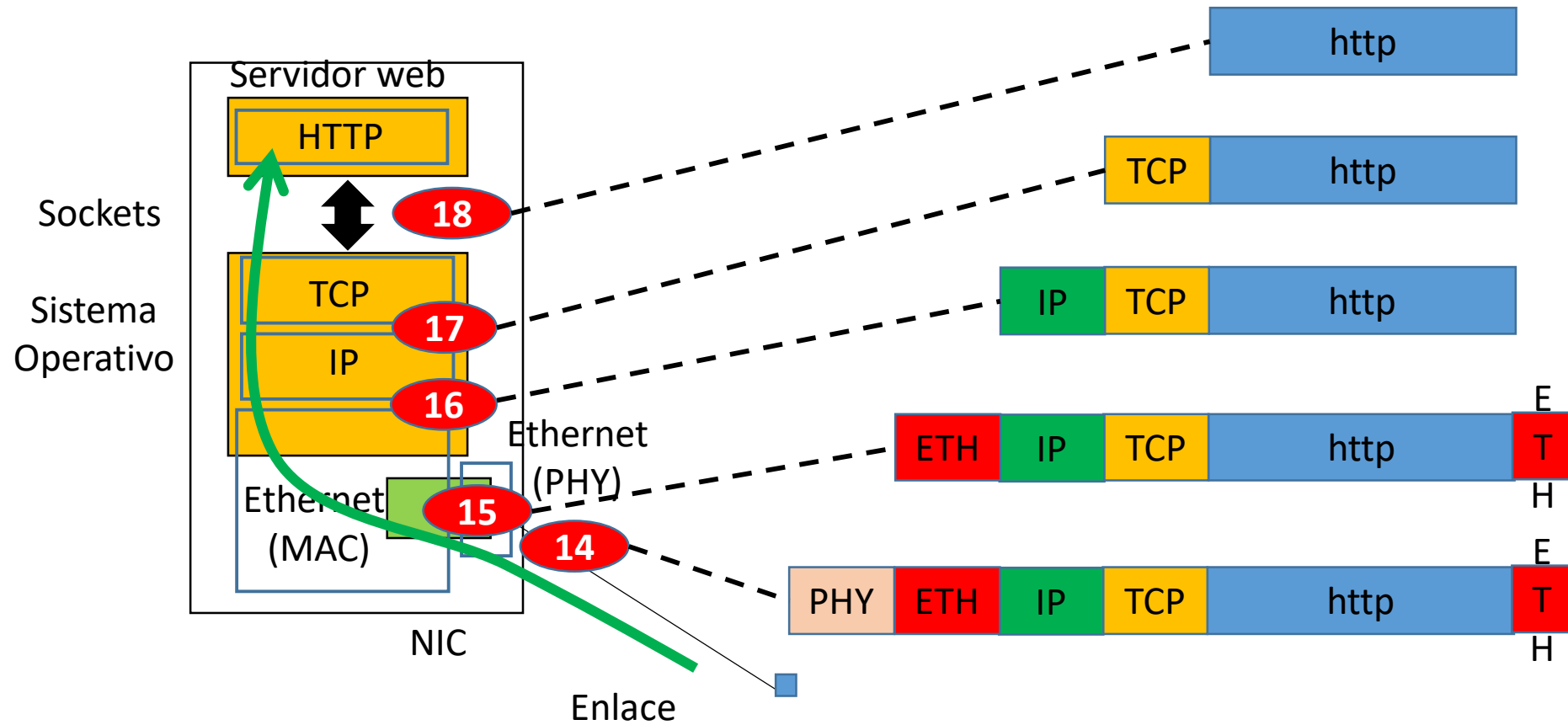
# Generación de paquetes

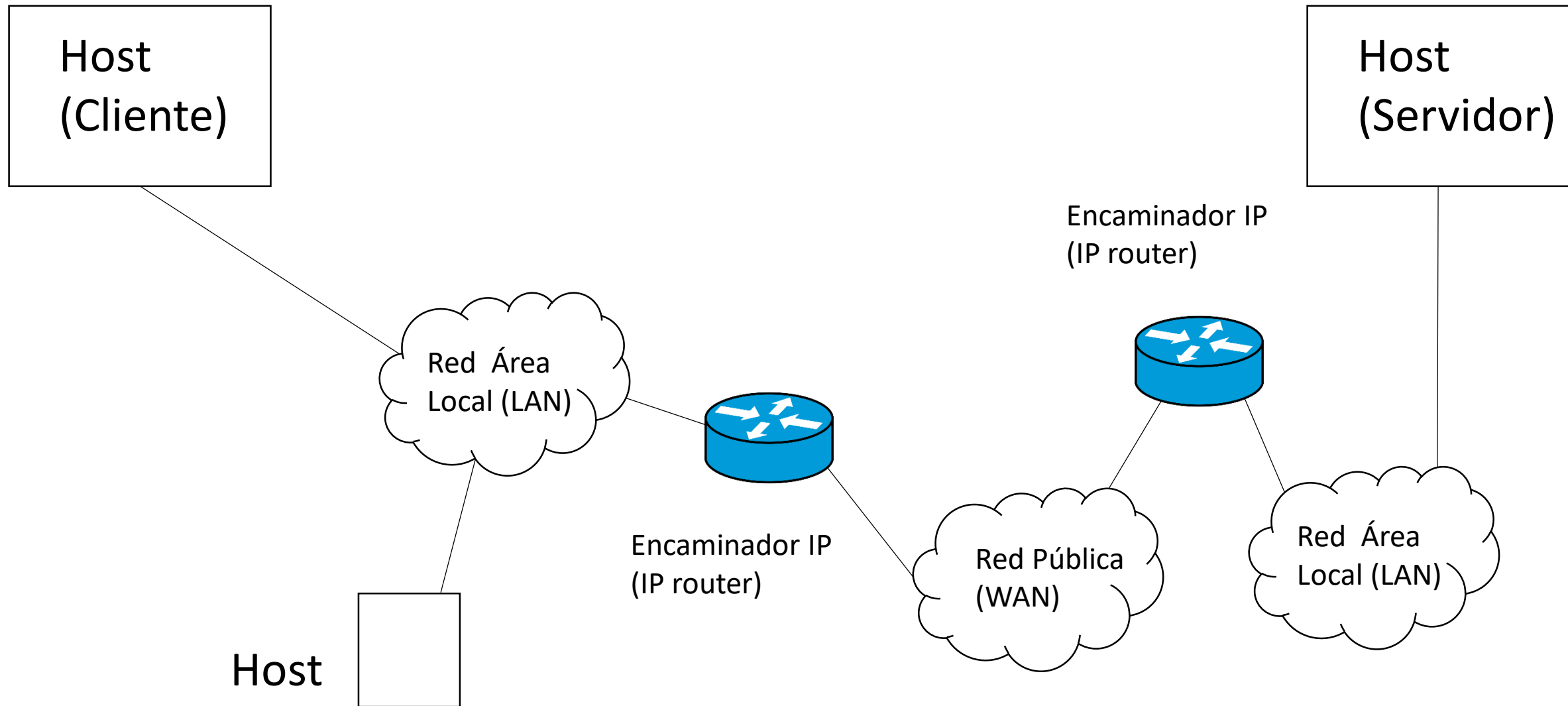


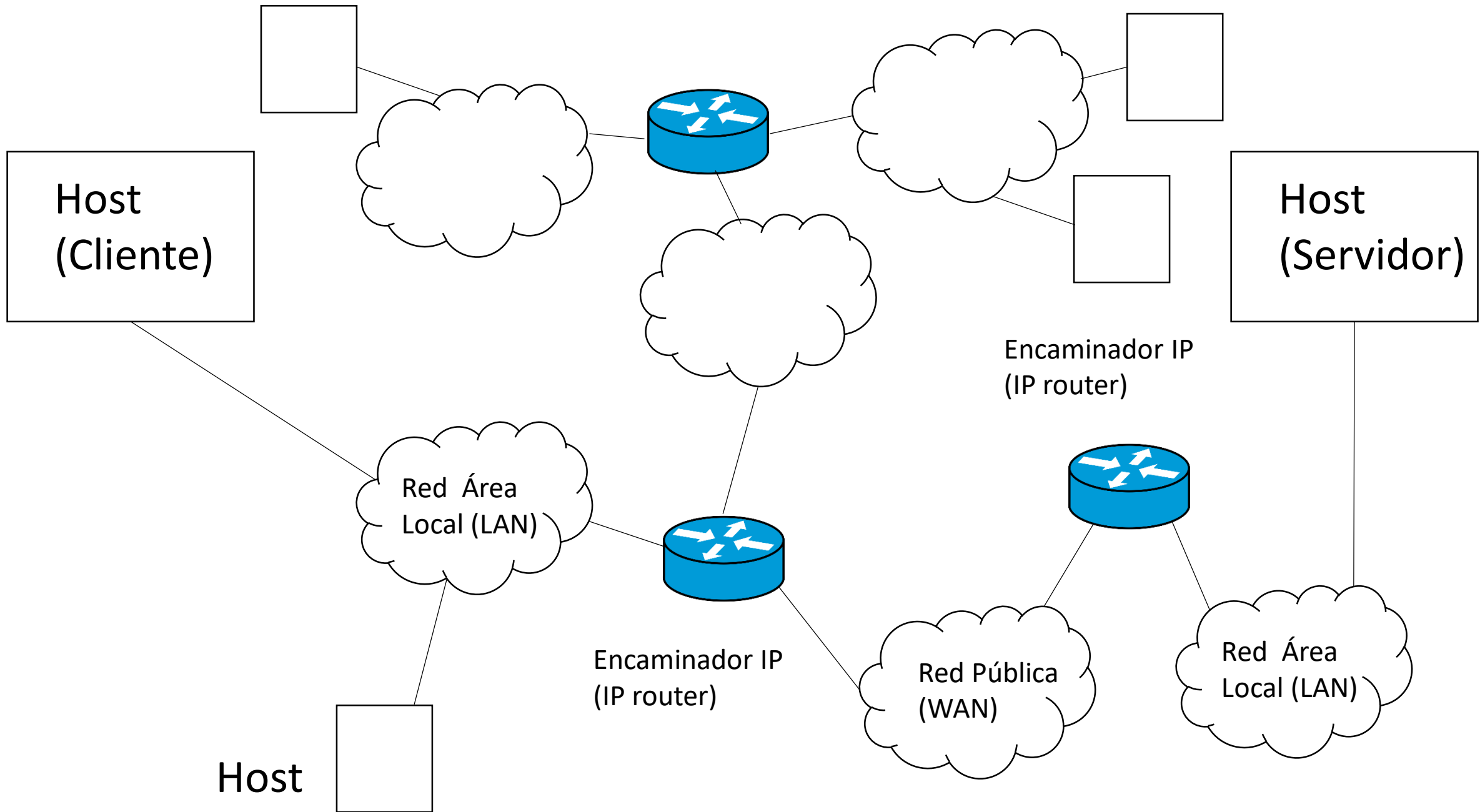




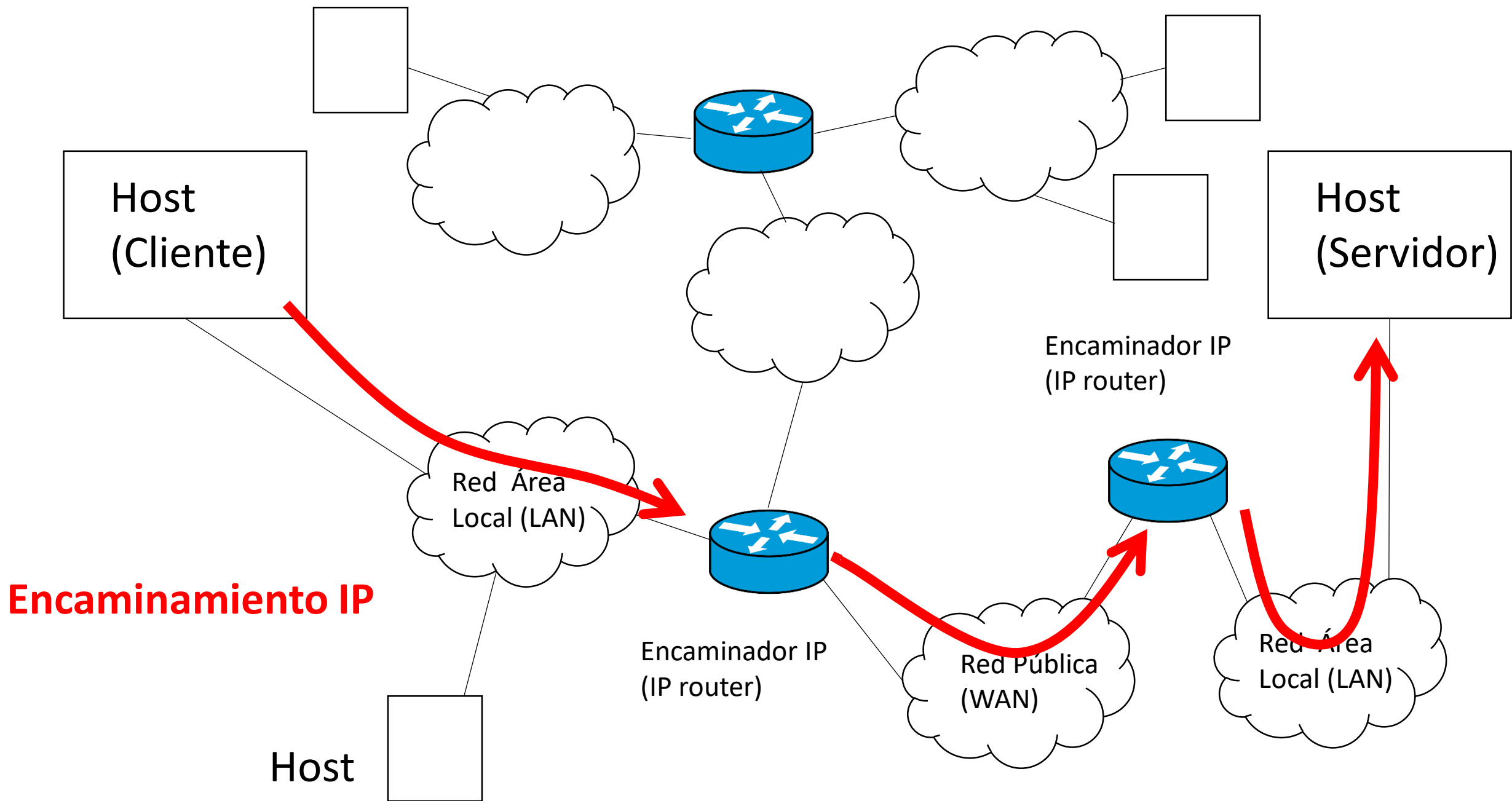
# Proceso de paquetes

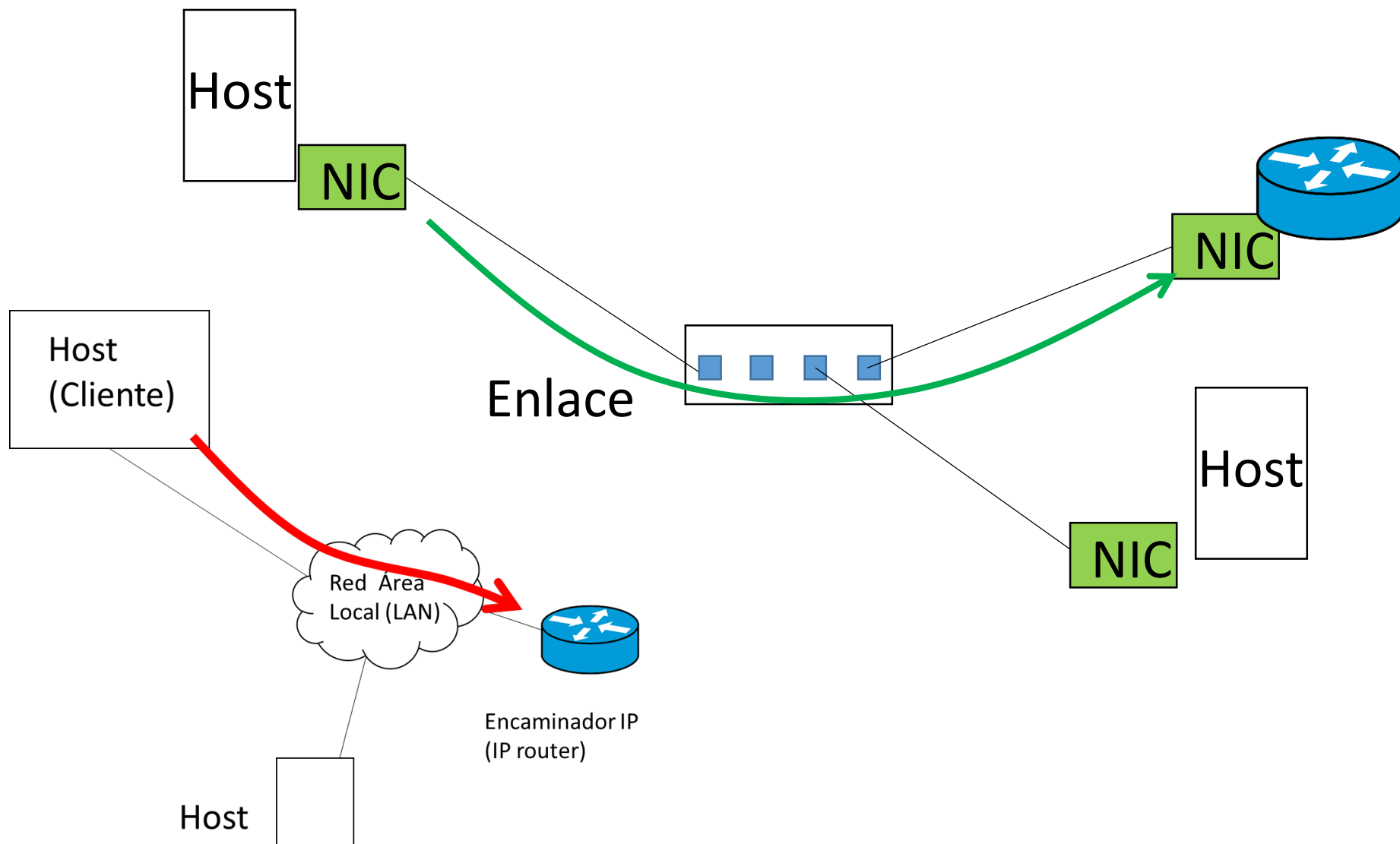


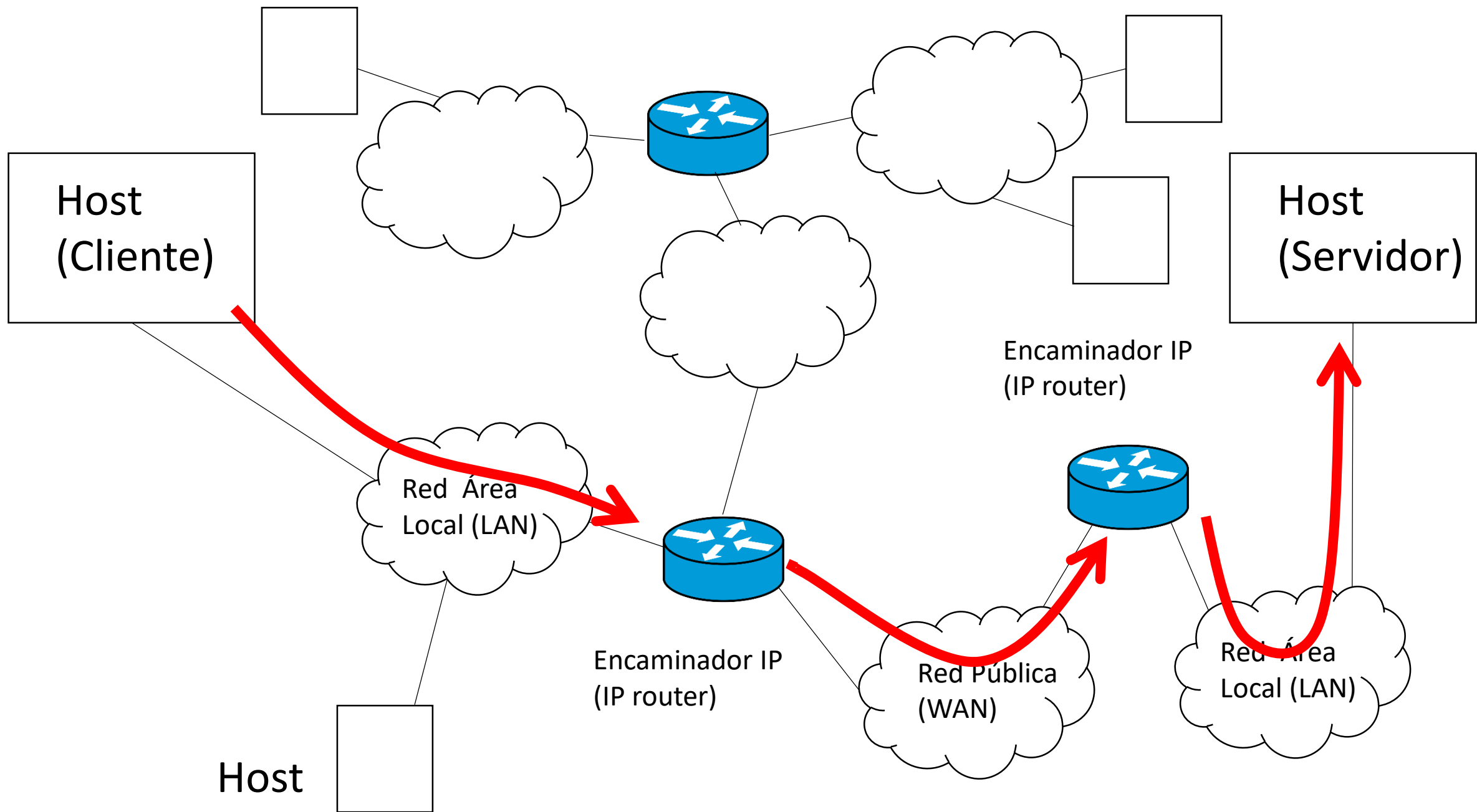












# Etiqueta, dirección, nombre

- **Etiqueta:** secuencia de bits única (ethernet: a2:10:4e:33:30:03)
- **Dirección:** Etiqueta que contiene información de encaminamiento (93 401 6978, direcciones IP: 172.45.67.69)
- **Nombre:** legible por humanos (<http://www.upc.edu>)

$$256 = 2^8$$

1 bit de subnet

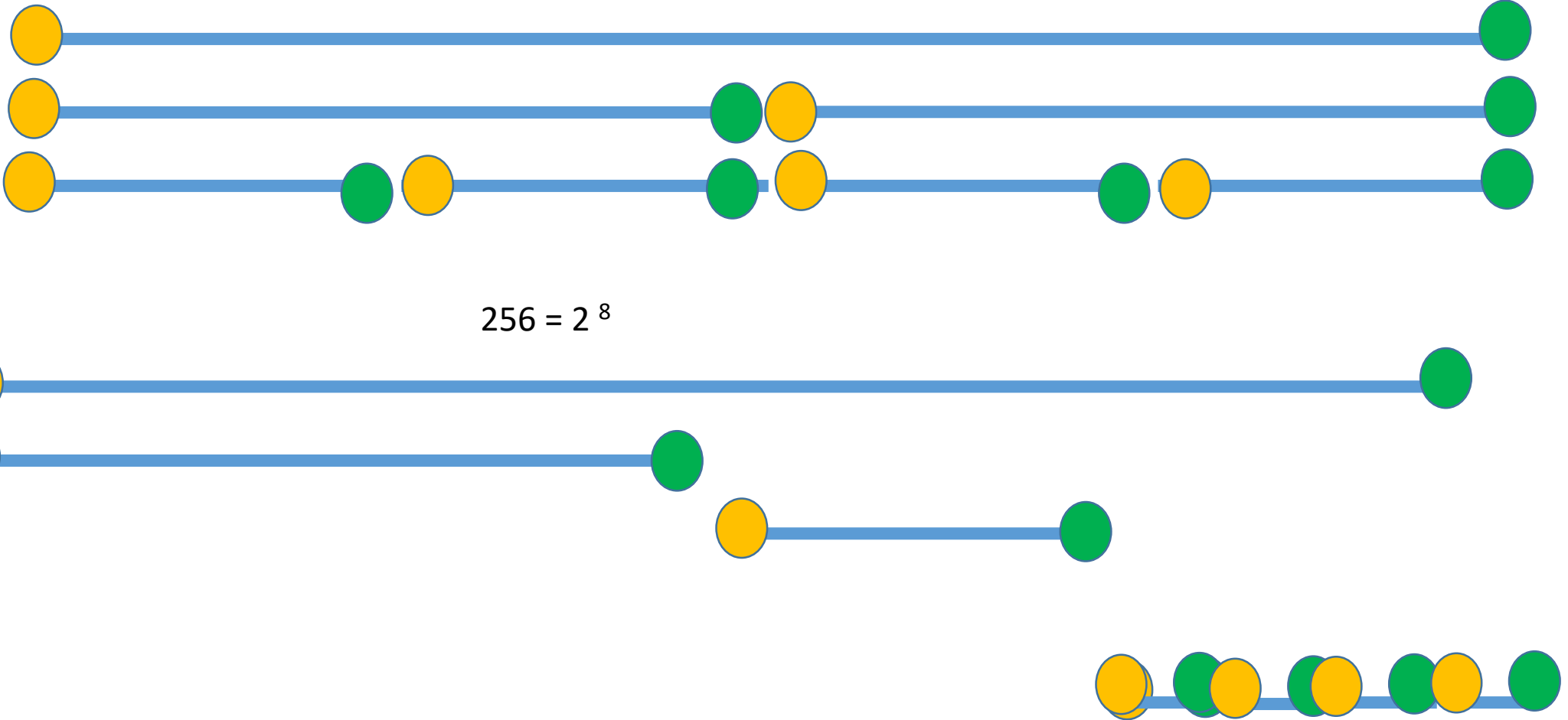
2 bits de subnet

$$256 = 2^8$$

1 bit de subnet

2 bits de subnet

4 bits de subnet



R1

destination	genmask	gateway	interface
200.10.10.0	255.255.255.0 /24	0.0.0.0	eth0
200.20.20.0	255.255.255.0 /24	0.0.0.0	eth1
0.0.0.0	0.0.0.0 /0	200.20.20.1	ppp0

Por el **eth0**, recibimos un paquete IP con @ destino= 200.20.20.10 (al PC2)

Se recorre la table de encaminamiento. Para cada entrada de la tabla se hace lo siguiente

- (1) Aplicamos (AND logico bit a bit) **genmask** a la @destino
- (2) Comparamos con **destination**. Hay match?
- (3) No match: pasamos a la siguiente entrada de la table
- (4) Sí match: **interface**= posible interface de salida; **gateway**: posible next node

Se escoge como ruta la entrada que ha hecho match y tiene más bits de prefijo de genmask (“Longest Prefix Match”)

		destination	genmask	gateway	interface
R1	a	200.10.10.0	255.255.255.0 /24	0.0.0.0	eth0
	b	200.20.20.0	255.255.255.0 /24	0.0.0.0	eth1
	c	0.0.0.0	0.0.0.0 /0	200.30.30.1	ppp0

Por el **eth0**, recibimos un paquete IP con @ destino= 200.20.20.10 (al PC2)

Se recorre la table de encaminamiento. Para cada entrada de la tabla se hace lo siguiente

(1) Aplicamos (AND logico bit a bit) **genmask** a la @destino

(a) 200.20.20.0

(b) 200.20.20.0

(c) 0.0.0.0

(2) Comparamos con **destination**. Hay match?

(3) No match: pasamos a la siguiente entrada de la table

Sí match: **interface**= posible interface de salida; **gateway**: posible next node

(a) NO match

(b) **SÍ match**: posible interface: eth1, posible gateway: directa

(c) **SÍ match**: posible interface: pp0, posible gateway: 200.30.30.1

Escogemos la ruta (b)

		destination	genmask	gateway	interface
R1	a	200.10.10.0	255.255.255.0 /24	0.0.0.0	eth0
	b	200.20.20.0	255.255.255.0 /24	0.0.0.0	eth1
	c	0.0.0.0	0.0.0.0 /0	200.30.30.1	eth2

Por el **eth0**, recibimos un paquete IP con @ destino= 147.82.20.1

Se recorre la table de encaminamiento. Para cada entrada de la tabla se hace lo siguiente

(1) Aplicamos (AND logico bit a bit) **genmask** a la @destino

(a) 147.82.20.0

(b) 147.82.20.0

(c) 0.0.0.0

(2) Comparamos con **destination**. Hay match?

(3) No match: pasamos a la siguiente entrada de la table

Sí match: **interface**= posible interface de salida; **gateway**: posible next node

(a) NO match

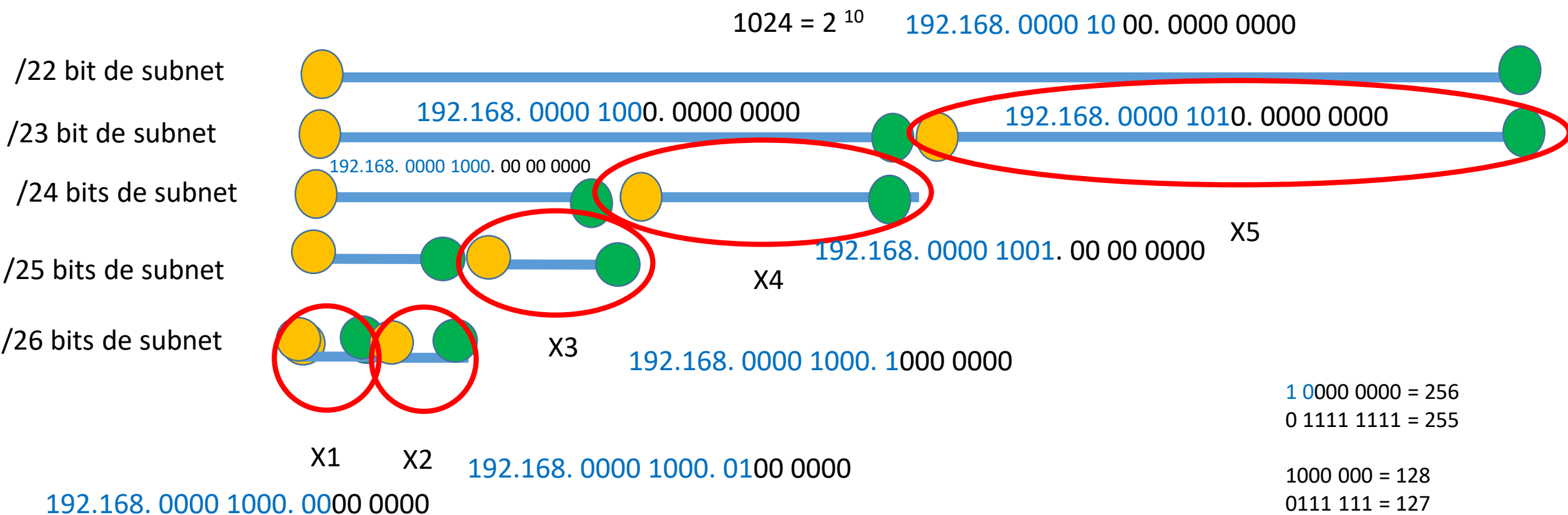
(b) NO match

(c) **Sí match**: posible interface: pp0, posible gateway: 200.30.30.1

Escogemos la ruta (c) (ruta por defecto)



- 192.168.8.0/26
- 192.168.0000 1000.00 00 0000
- 192.168.8.00 00 0001 => 192.168.8.1
- 192.168.8.00 11 1110 => 192.168.8.62
- Rango assignable: 192.168.8.1 a 192.168.8.62
- Dirección de red: 192.168.8.0 /26
- Dirección de broadcast => 192.168.8.63
- 0100 0000 => 64



X3:

# equipos configurables:  $2^7 - 2 = 128 - 2 = 126$

@ de subred: 192.168. 0000 1000. 1000 0000 => 192.168. 8. 128/25

@ broadcast: 192.168.8.255/25

@ router (1a dirección asignable) 192.168.8.129

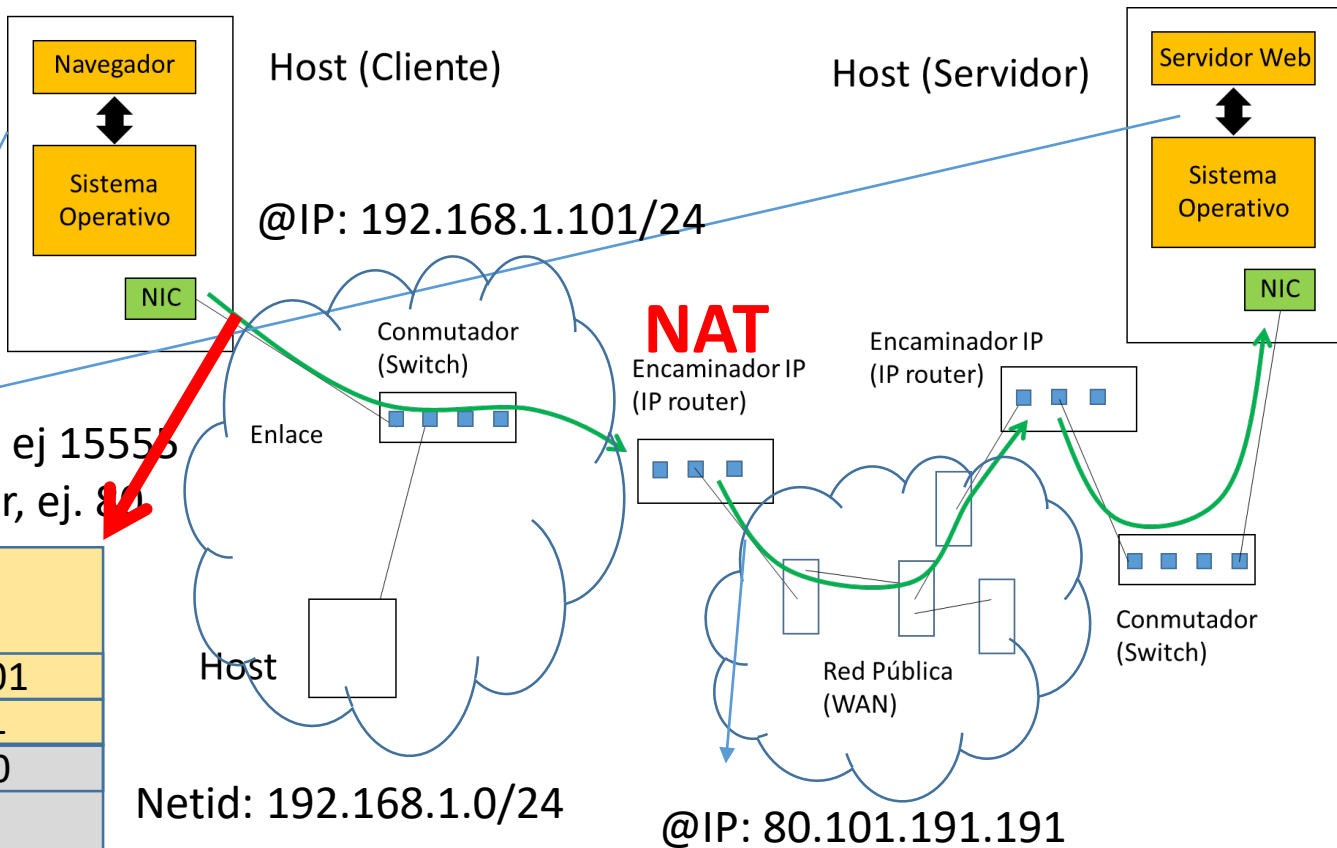
- 80.80.181.178 /18
- Direccion de red: 80.80.128.0/18
- 80.80. 10 11 1111. 255 => @ de broadcast = 80.80.191.255
- 128 + 63= 191
- 128 = 1000 0000
- 64 = 0100 0000
- 181 = 10 11 0101

Cliente

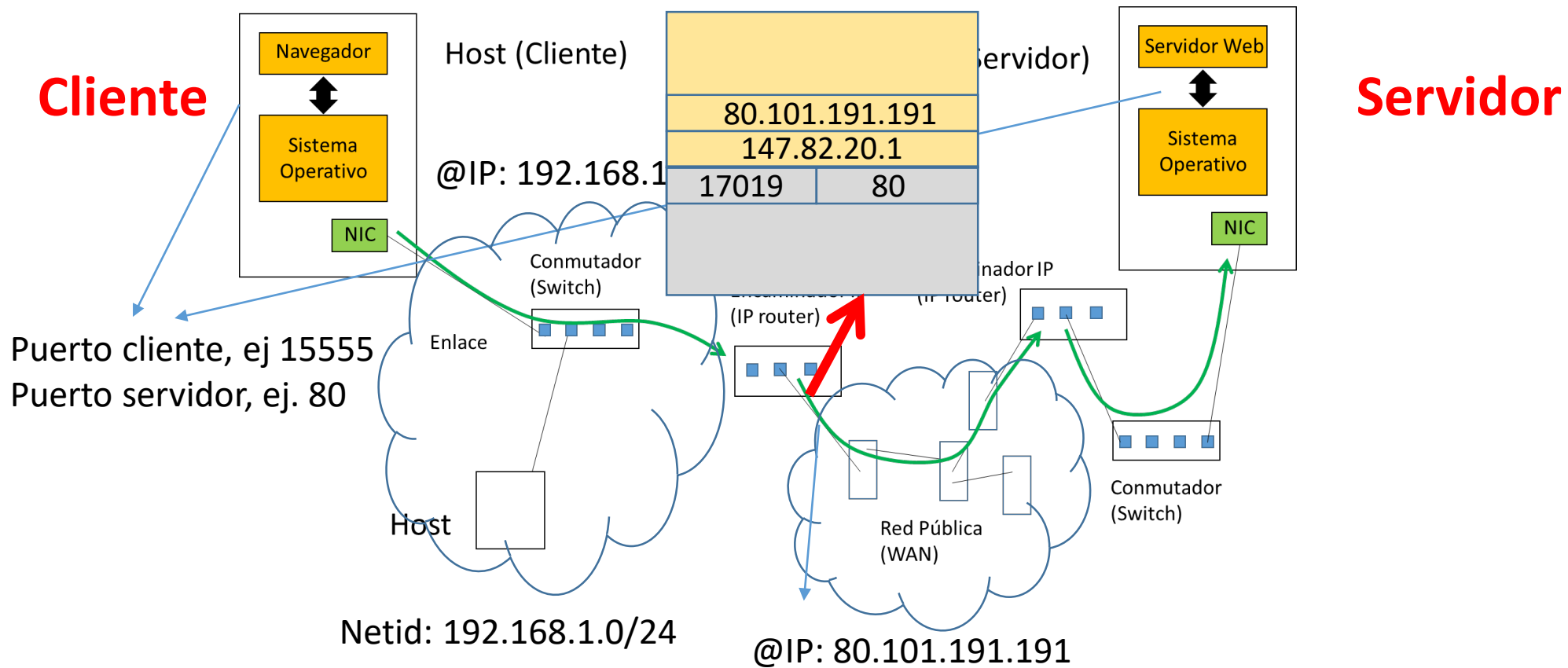
Servidor

192.168.1.101	
147.82.20.1	
15555	80

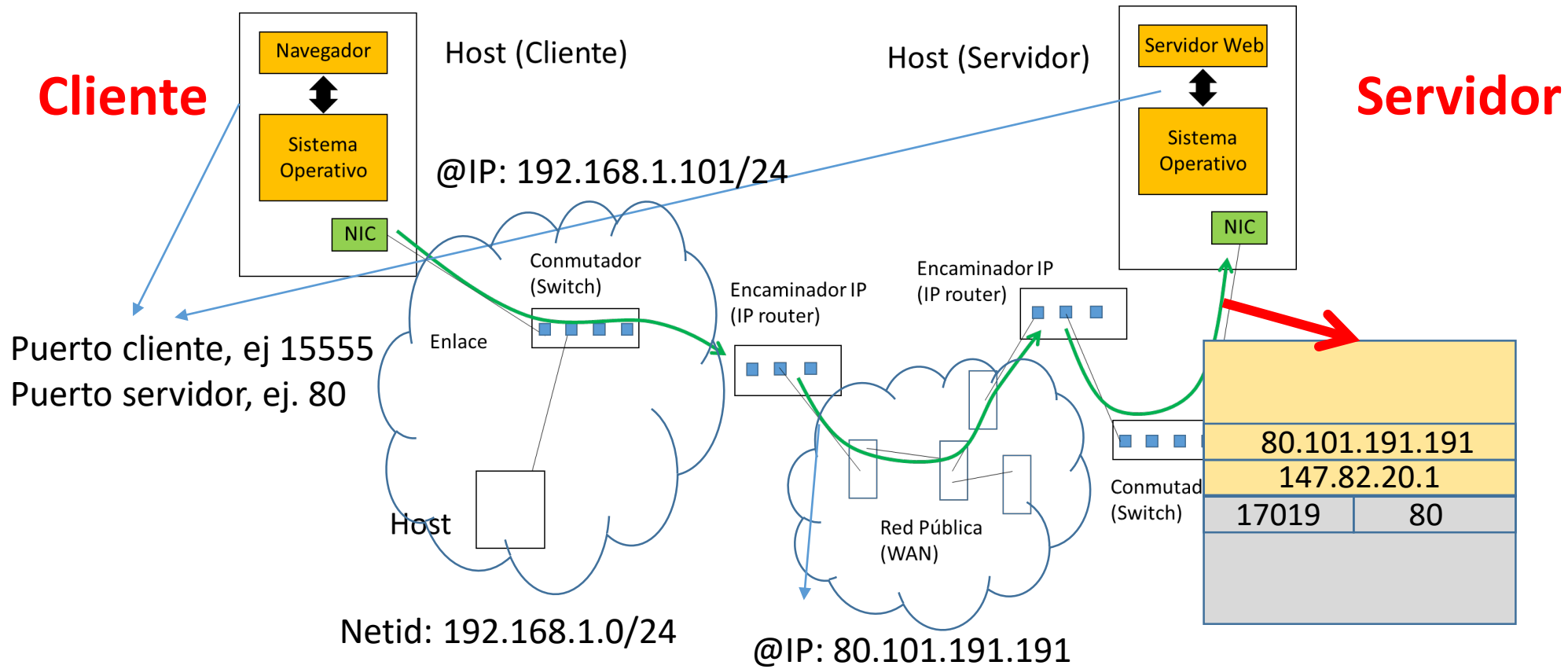
Puerto cliente, ej 15555  
Puerto servidor, ej. 80



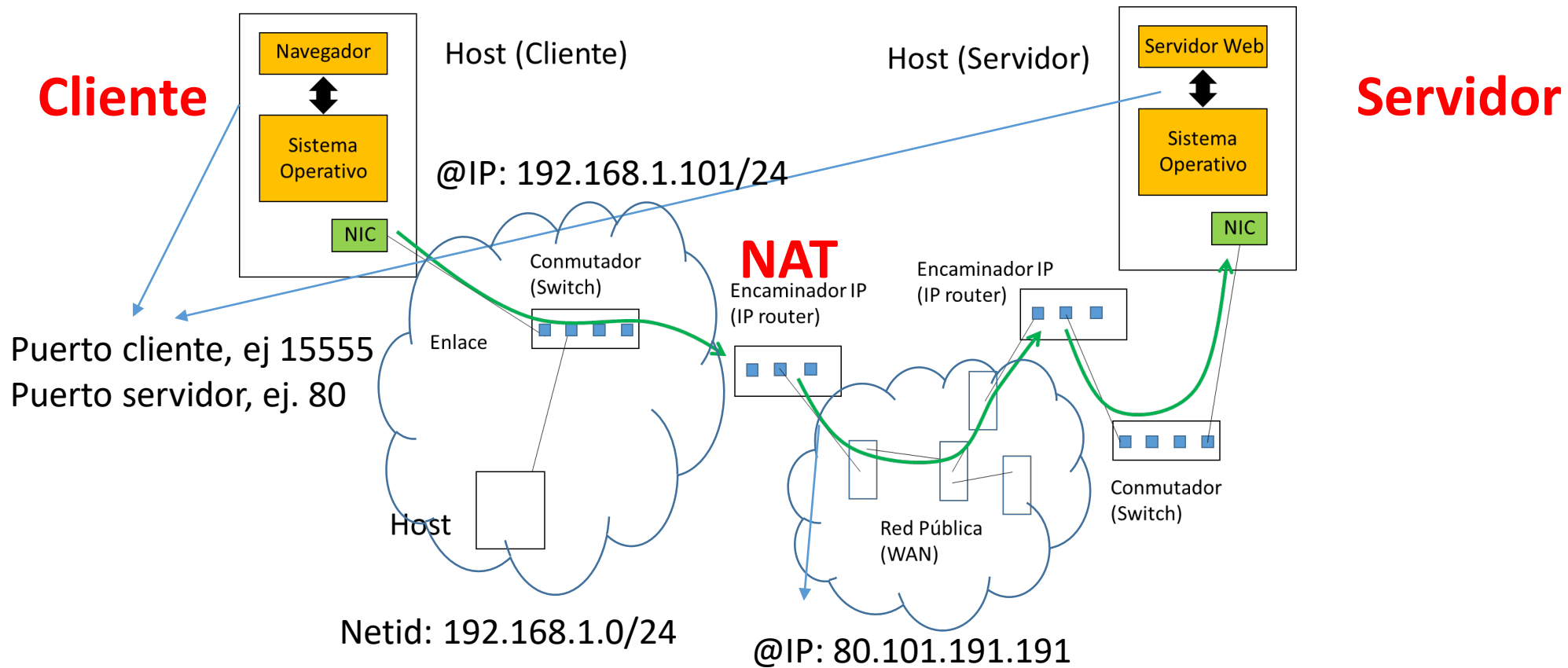
@IP interna	@IP externa	Port interno	Port externo
-------------	-------------	--------------	--------------



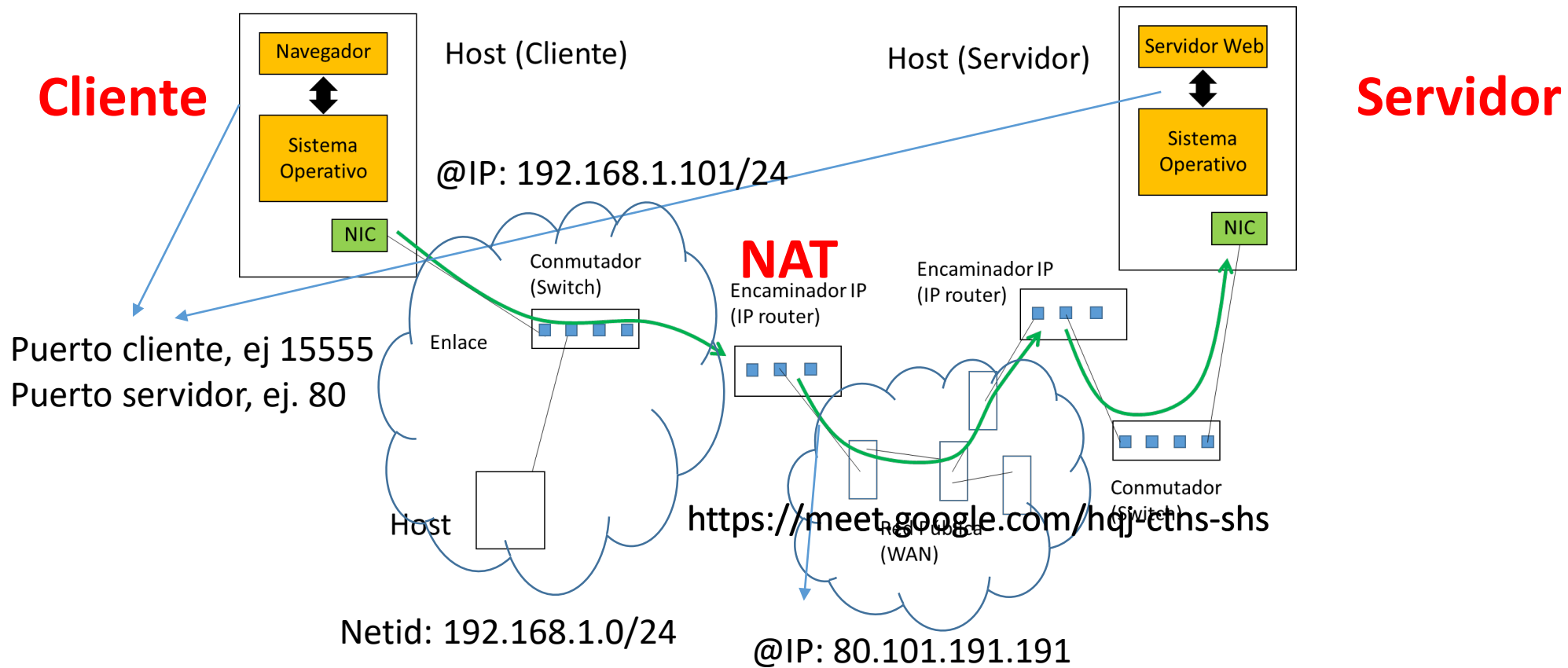
@IP interna	@IP externa	Port interno	Port externo
192.168.1.101	80.101.191.191	15555	17019



@IP interna	@IP externa	Port interno	Port externo
192.168.1.101	80.101.191.191	15555	17019

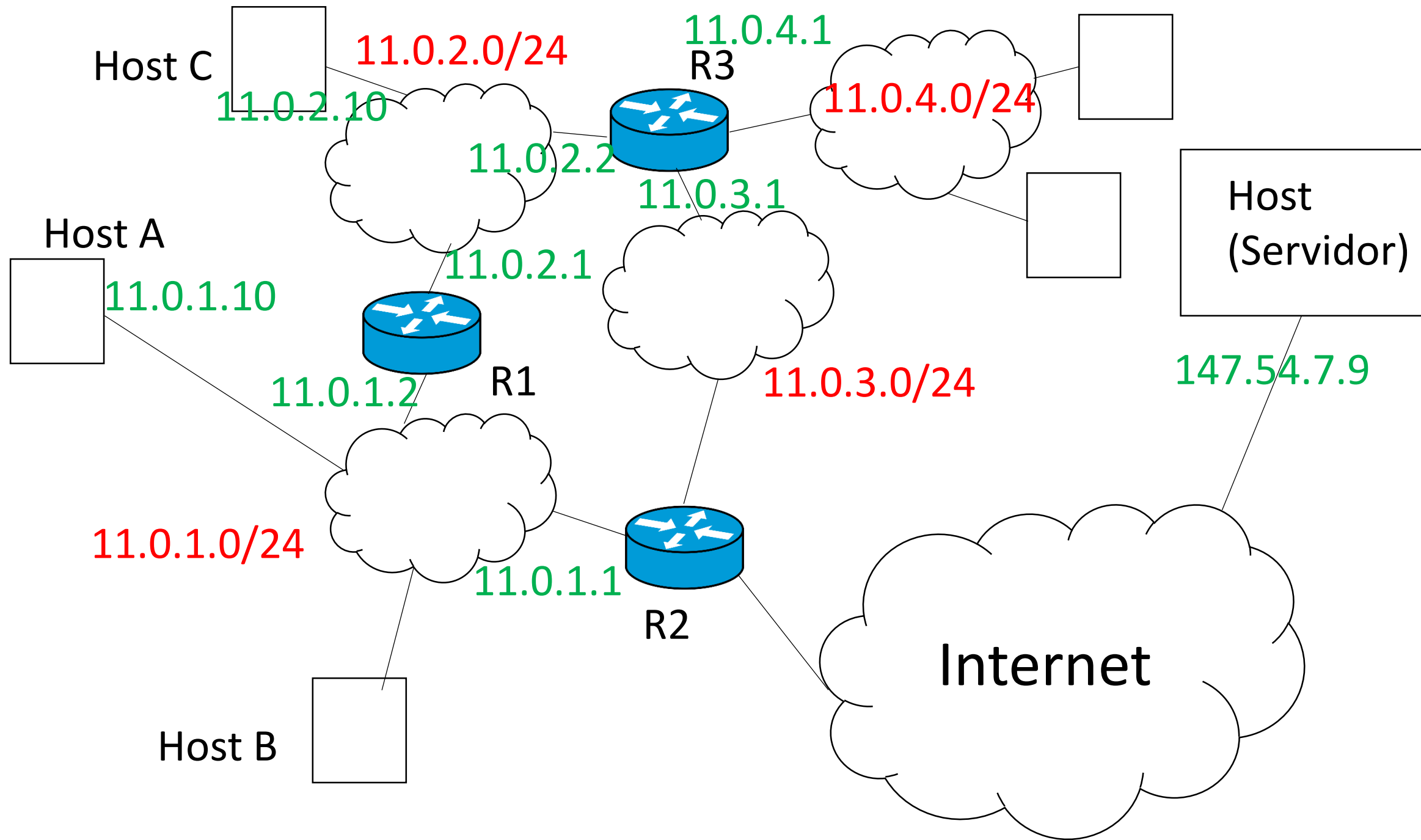


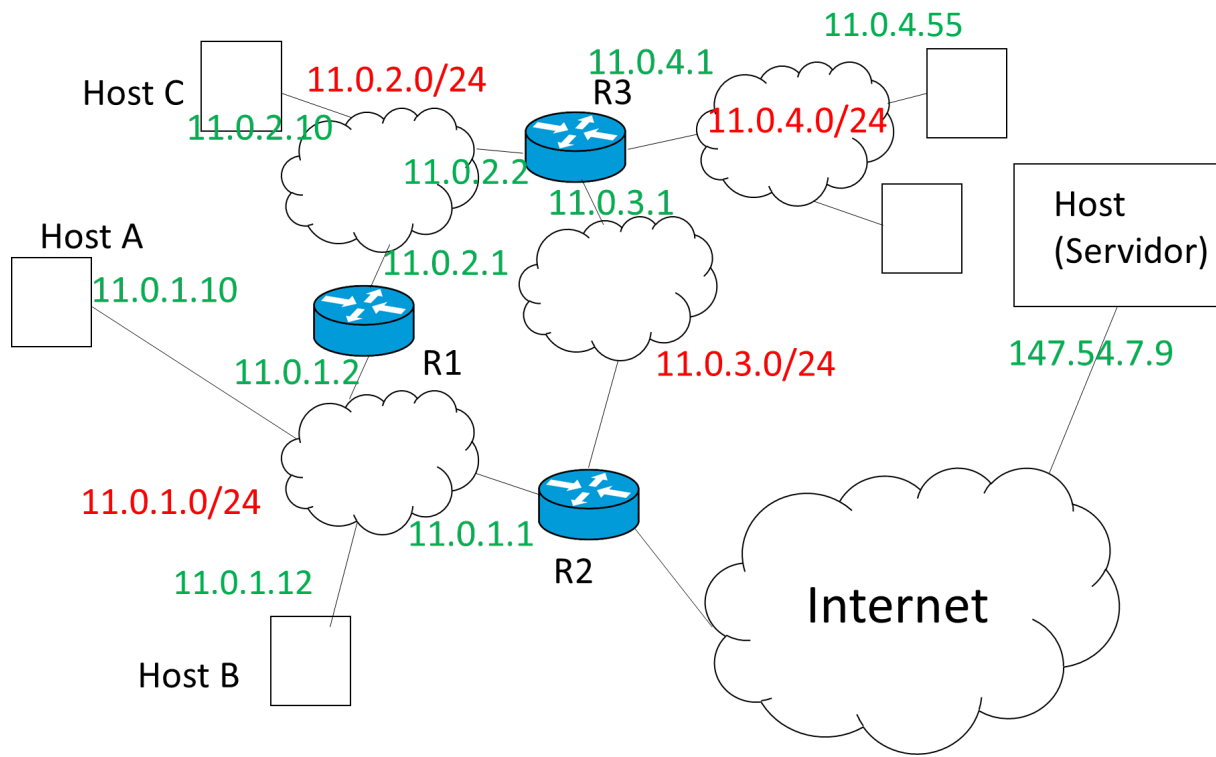
@IP interna	@IP externa	Port interno	Port externo
192.168.1.101	80.101.191.191	15555	17019
192.168.1.102	80.101.191.191	15787	17020



@IP interna	@IP externa	Port interno	Port externo
192.168.1.101	80.101.191.191	15555	17019
192.168.1.102	80.101.191.191	15787	17020
192.168.1.102	80.101.191.191	15555	17021





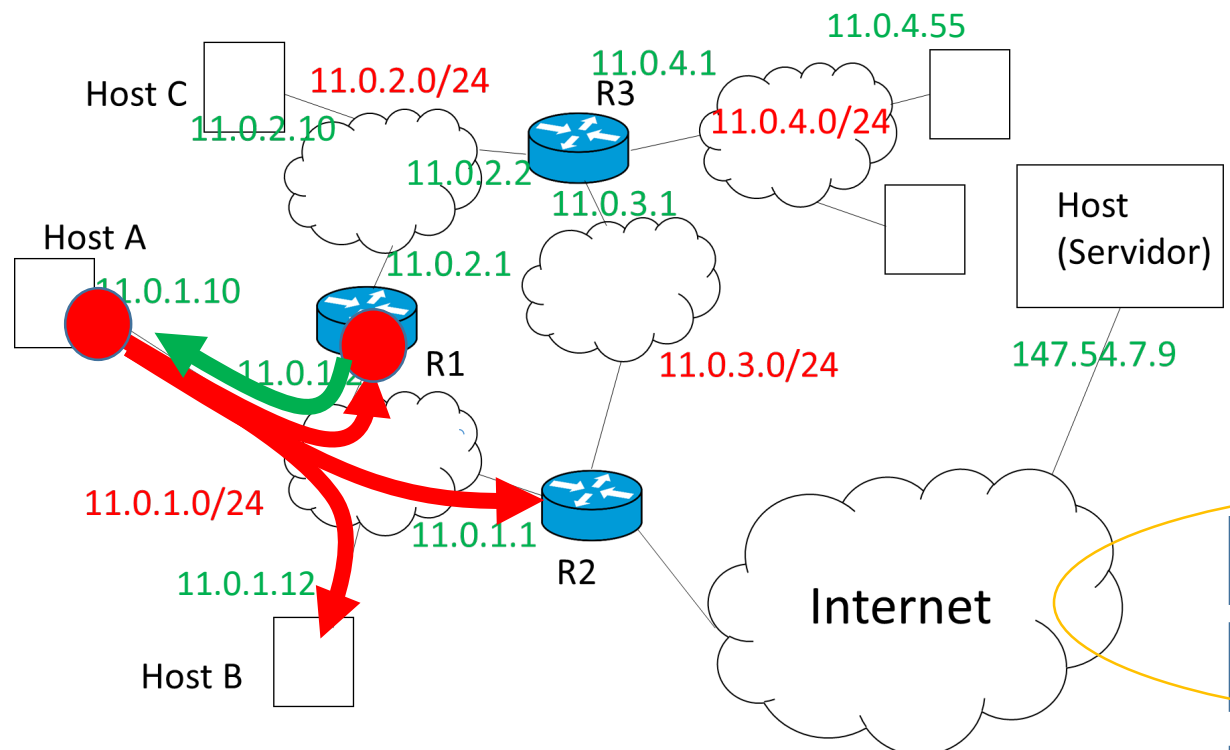


**@IP destino (host B): 11.0.2.10**

**Lookup en la table de encaminatmiento de Host A**  
(1) Aplicamos 255.0.0.0 : 11.0.0.0 **NO MATCH**  
(2) Aplicamos 255.255.255.0: 11.0.2.0 **NO MATCH**  
(3) Aplicamos 255.255.255.0: 11.0.2.0 **MATCH**  
(4) Aplicamos 0.0.0.0: 0.0.0.0 **MATCH**

Host A

destination	genmask	gateway	interface
127.0.0.0	255.0.0.0 (/8)	0.0.0.0 (directa)	loopback
11.0.1.0	255.255.255.0 (/24)	0.0.0.0 (directa)	eth0
11.0.2.0	255.255.255.0 (/24)	11.0.1.2 (R1)	eth0
0.0.0.0	0.0.0.0 (/0)	11.0.1.1 (R2)	eth0



**HostA > ping 11.0.2.10**  
**IP: @org: 11.0.1.10**  
**@dst: 11.0.2.10**  
**ICMP echo request**

Ethernet  
 MAC de A; MAC de R1      payload

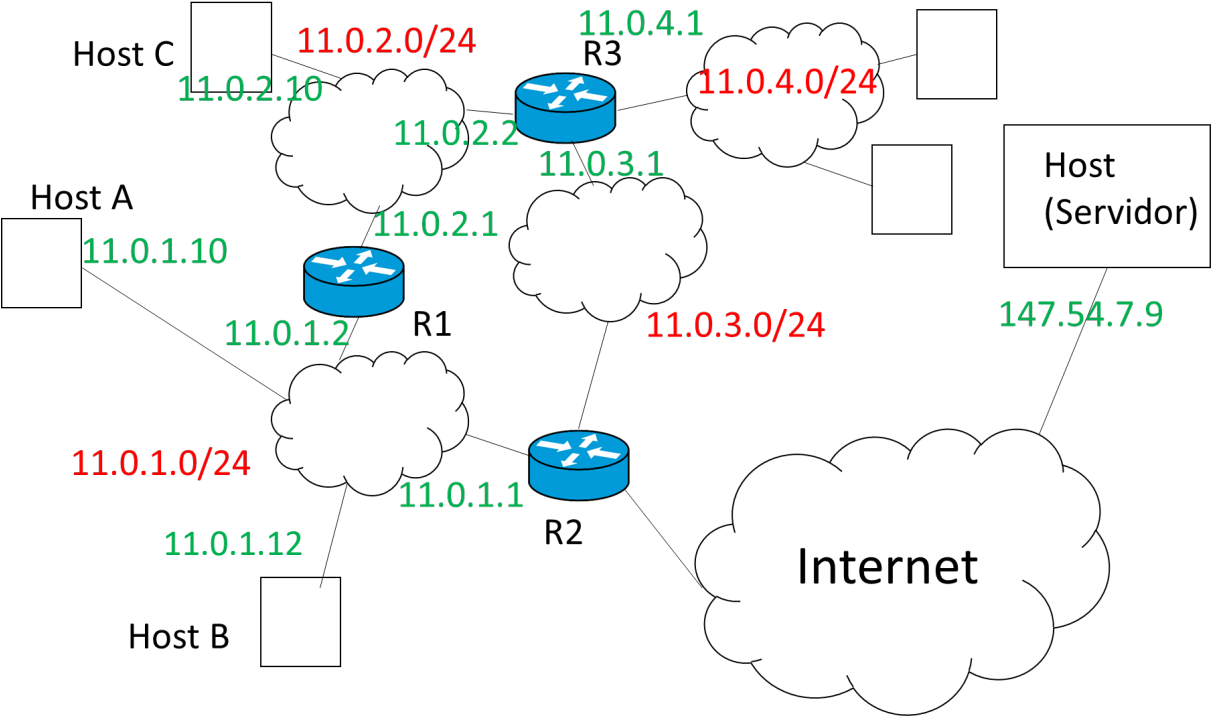
ARP, MAC de A; MAC broadcast  
 Who has 11.0.1.2?      HA -> broadcast

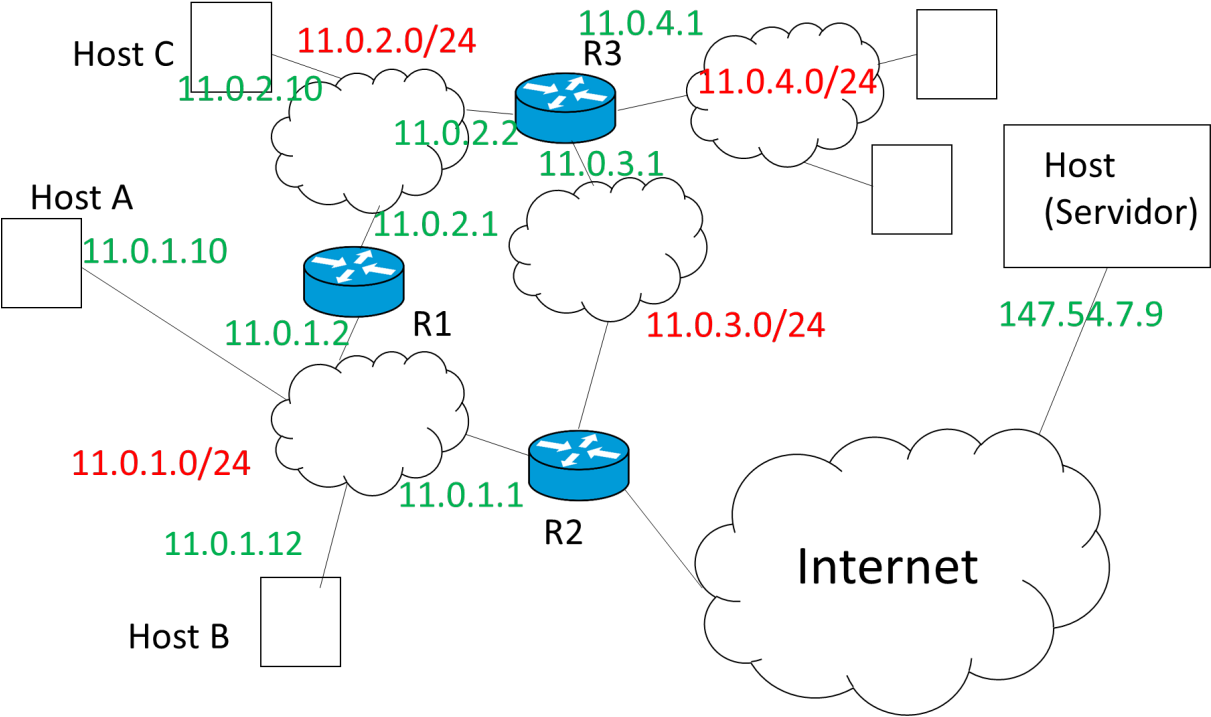
ARP, MAC de R1; MAC A  
 11.0.1.2 -> MAC de R1      R1 -> HA

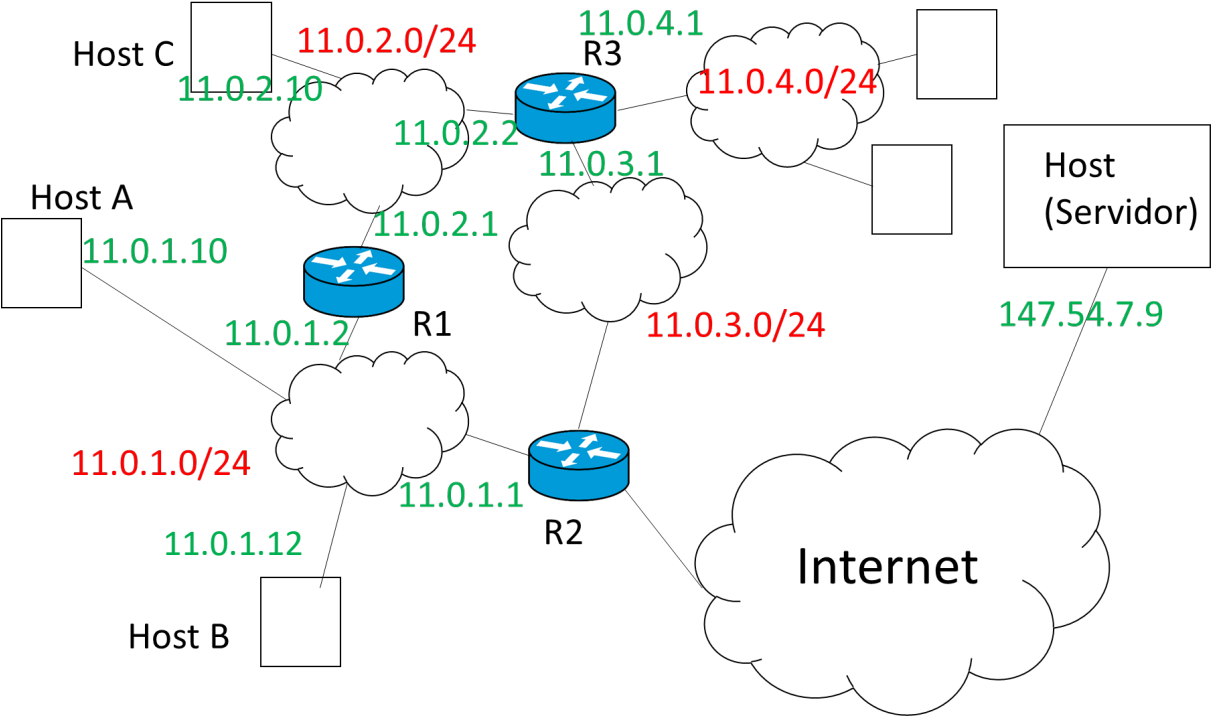
Ethernet  
 MAC de A; MAC de R1      payload      HA -> R1

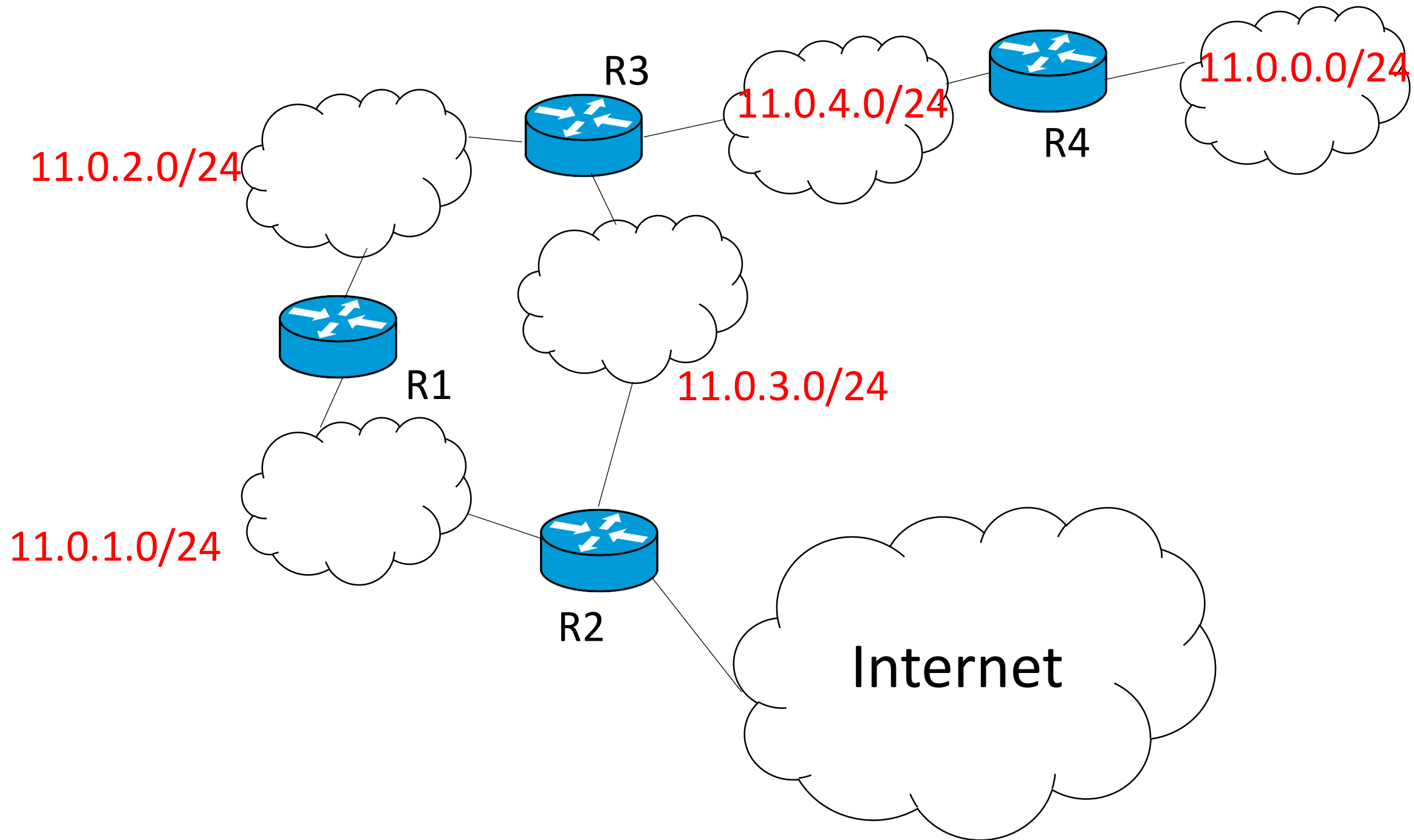
**Host A**

destination	genmask	gateway	interface
127.0.0.0	255.0.0.0 (/8)	0.0.0.0 (directa)	loopback
11.0.1.0	255.255.255.0 (/24)	0.0.0.0 (directa)	eth0
11.0.2.0	255.255.255.0 (/24)	11.0.1.2 (R1)	eth0
0.0.0.0	0.0.0.0 (/0)	11.0.1.1 (R2)	eth0









[\[Docs\]](#) [\[txt|pdf\]](#) [\[draft-ietf-nipv...\]](#) [\[Tracker\]](#) [\[Diff1\]](#) [\[Diff2\]](#) [\[Errata\]](#)

Updated by: [4822](#)

INTERNET STANDARD

[Errata Exist](#)

Network Working Group

G. Malkin

Request for Comments: 2453

Bay Networks

Obsoletes: [1723](#), [1388](#)

November 1998

STD: 56

Category: Standards Track

## RIP Version 2

### Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

### Copyright Notice

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### Abstract

This document specifies an extension of the Routing Information Protocol (RIP), as defined in [\[1\]](#), to expand the amount of useful information carried in RIP messages and to add a measure of security.

A companion document will define the SNMP MIB objects for RIP-2 [\[2\]](#). An additional document will define cryptographic security improvements for RIP-2 [\[3\]](#).

### Acknowledgements

I would like to thank the IETF RIP Working Group for their help in improving the RIP-2 protocol. Much of the text for the background discussions about distance vector protocols and some of the descriptions of the operation of RIP were taken from "Routing Information Protocol" by C. Hedrick [\[1\]](#). Some of the final editing on the document was done by Scott Bradner.

[\[RFC Home\]](#) [\[TEXT|PDF|HTML\]](#) [\[Tracker\]](#) [\[IPR\]](#)

PROPOSED STANDARD

Network Working Group

R. Atkinson

Request for Comments: 4822

Extreme Networks

Obsoletes: [2082](#)

M. Fanto

Updates: [2453](#)

NIST

Category: Standards Track

February 2007

## RIPv2 Cryptographic Authentication

### Status of This Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

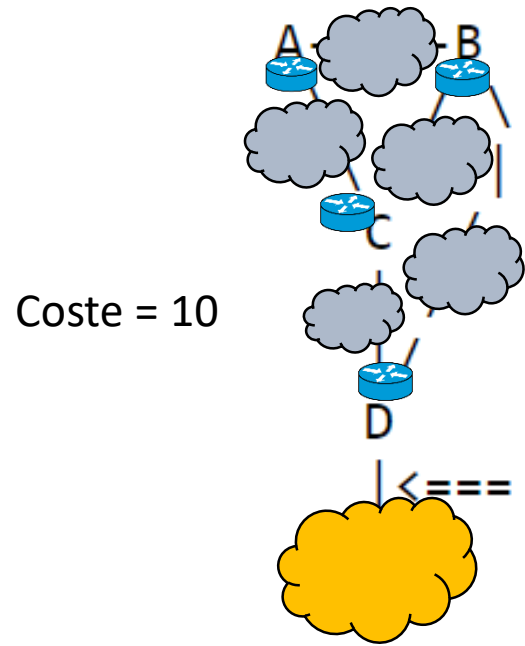
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### IESG Note

In the interests of encouraging rapid migration away from Keyed-MD5 and its known weakness, the IESG has approved this document even though it does not meet the guidelines in [BCP 107](#) ([RFC 4107](#)). However, the IESG stresses that automated key management should be used to establish session keys and urges that the future work on key management described in [Section 5.6](#) of this document should be





all networks have cost 1, except  
for the direct link from C to D, which  
has cost 10

<=== target network

## Rutas a D

D: directly connected, metric 1

*Entrada en la tabla de encamamiento del router D  
que señala a la target network*

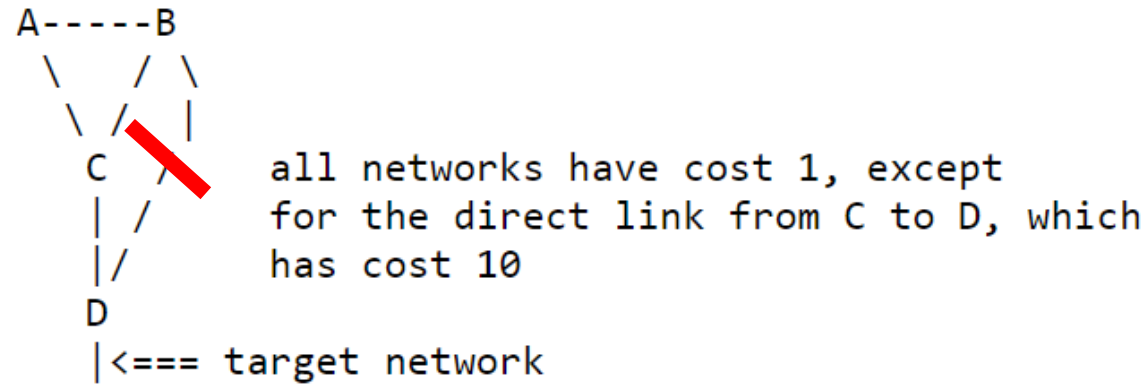
B: route via D, metric 2

C: route via B, metric 3

*Entrada en la tabla de encamamiento del router C  
que señala a la target network*

A: route via B, metric 3

Etc..



Ruta de B a D se rompe (180 sec, 6 updates perdidos,  
timeout)

time ----->

D: dir, 1    dir, 1  
B: unreach C, 4  
C: B, 3    A, 4  
A: B, 3    C, 4

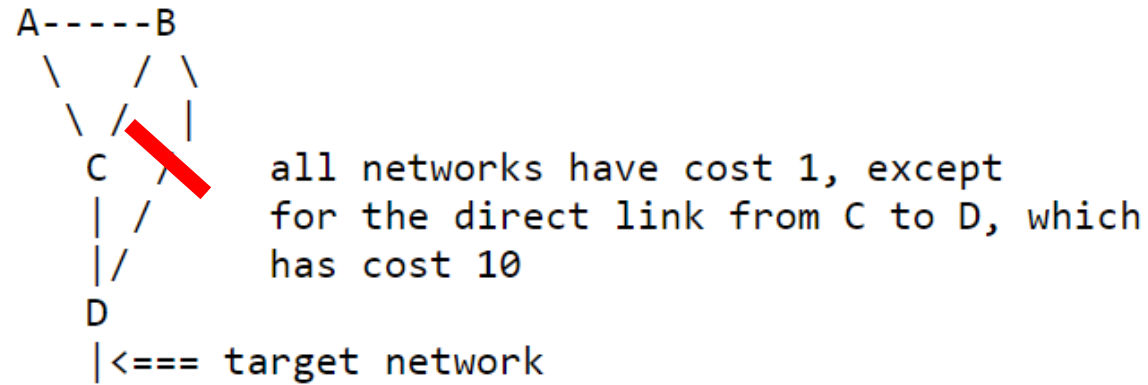
dir = directly connected  
unreach = unreachable

Mensajes RIPv2 recibidos por C:

De B: target network unreachable (coste=16)

De A: target network con coste 3

De D: target network con coste 1



Ruta de B a D se rompe (180 sec, 6 updates perdidos,  
timeout)

time ----->

D:	dir, 1	dir, 1
B:	unreach	C, 4
C:	B, 3	A, 4
A:	B, 3	C, 4

dir = directly connected  
unreach = unreachable

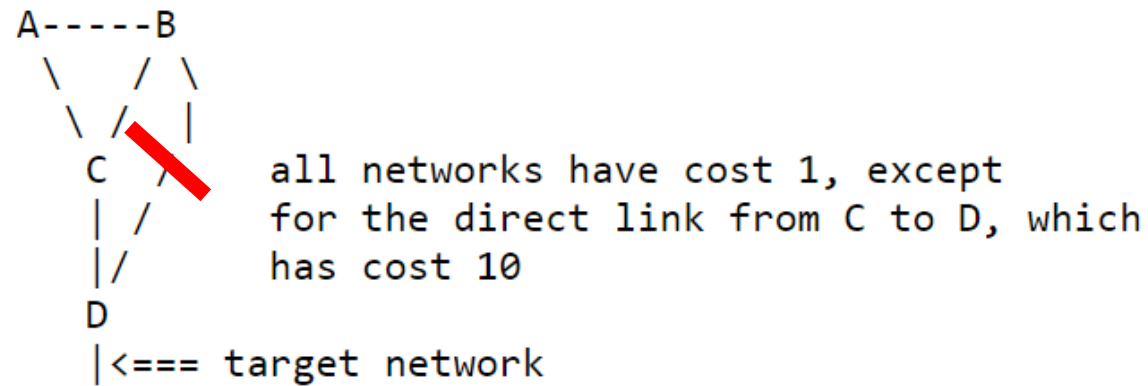
Actualizacion de la tabla de C

**De B: target network unreachable (coste=16)**

De A: target network con coste 3 (coste de C a A es 1) => coste 4

De D: target network con coste 1 (coste de C a D es 10) => coste 11

C actualiza la table a target network Router A, coste 4

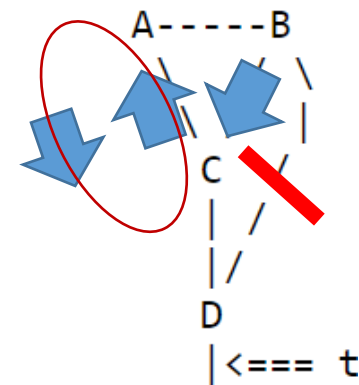


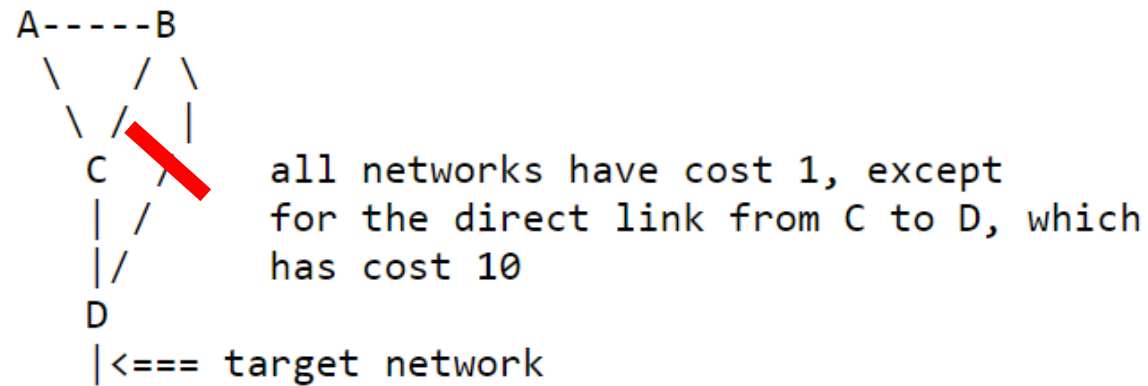
Ruta de B a D se rompe (180 sec, 6 updates perdidos,  
timeout)

time ----->

D: dir, 1      dir, 1  
B: unreach    C,    4  
C: B,    3    A,    4  
A: B,    3    C,    4

dir = directly connected  
unreach = unreachable



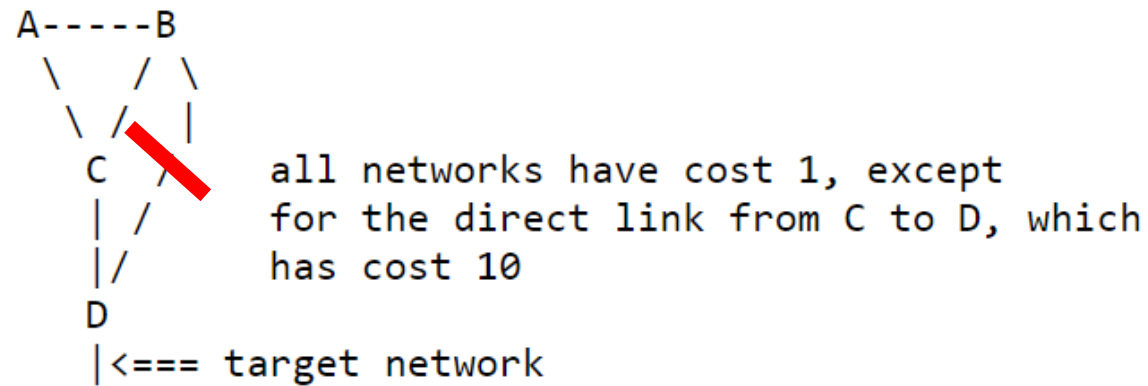


Ruta de B a D se rompe (180 sec, 6 updates perdidos,  
timeout)

time ----->

D: dir, 1	dir, 1	dir, 1	dir, 1
B: unreach	C, 4	C, 5	C, 6
C: B, 3	A, 4	A, 5	A, 6
A: B, 3	C, 4	C, 5	C, 6

dir = directly connected  
unreach = unreachable

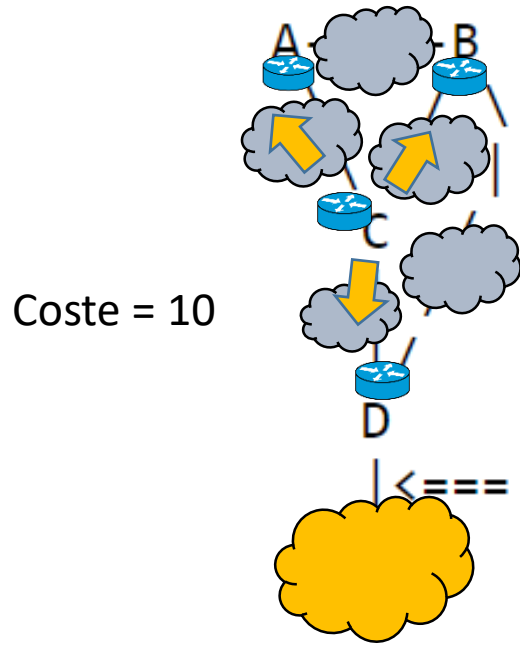


Ruta de B a D se rompe (180 sec, 6 updates perdidos,  
timeout)

time ----->

D: dir, 1	dir, 1	dir, 1	dir, 1	...	dir, 1	dir, 1
B: unreach	C, 4	C, 5	C, 6		C, 11	C, 12
C: B, 3	A, 4	A, 5	A, 6		A, 11	D, 11
A: B, 3	C, 4	C, 5	C, 6		C, 11	C, 12

dir = directly connected  
unreach = unreachable



**Split-Horizon:** Si aprendo una ruta a través de un interface de red, no envío esa ruta en los mensajes RIPv2 que envío por dicho interfaz. Si hago un update de la ruta a través de otro interfaz, entonces sí la incluyo. **Poisoned-reverse:** la incluyo con un coste de 16

all networks have cost 1, except for the direct link from C to D, which has cost 10

<=== target network

## Rutas a D

D: directly connected, metric 1

*Entrada en la tabla de encamamiento del router D que señala a la target network*

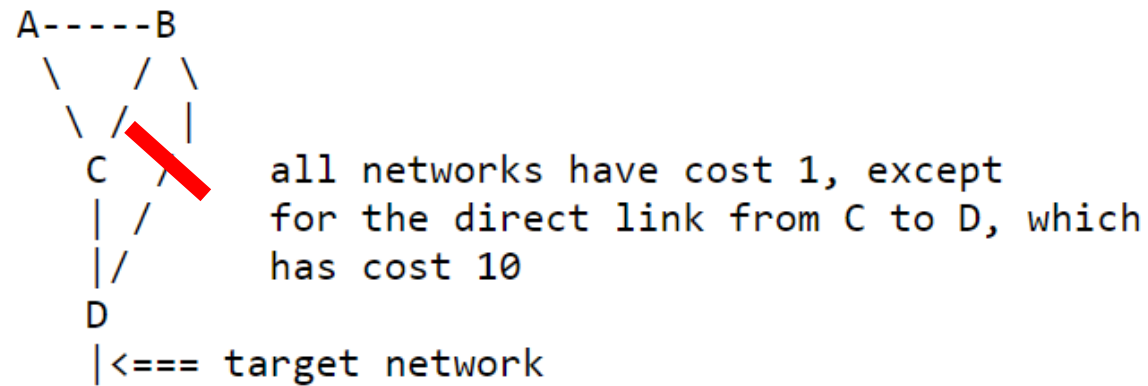
B: route via D, metric 2

C: route via B, metric 3

*Entrada en la tabla de encamamiento del router C que señala a la target network*

A: route via B, metric 3

Etc..



Ruta de B a D se rompe (180 sec, 6 updates perdidos,  
timeout)

time ----->

Triggered updates: no espero 30 sec si hay un cambio en la tabla

30 segundos => pocos segundos

D: dir, 1	dir, 1	dir, 1	dir, 1	...	dir, 1	dir, 1
B: unreach	C, 4	C, 5	C, 6		C, 11	C, 12
C: B, 3	A, 4	A, 5	A, 6		A, 11	D, 11
A: B, 3	C, 4	C, 5	C, 6		C, 11	C, 12

dir = directly connected  
unreach = unreachable