ICPC Team Notebook (2018-19)

Contents

1	Con	nbinatorial optimization	
	1.1	Dense max-flow	
	1.2	Min-cost max-flow	
	1.3	Min-cost matching	
	1.4	Sparse max-flow	
	1.5	Global min-cut	
	1.6	Push-relabel max-flow	
	1.7	Max bipartite matching	
	1.8	Graph cut inference	
	1.9	General Matching	
	1.10	Edmond Blossom Matching	
	1.11	Min Edge Cover	
	1.12	Stable Marriage Problem	
2		metry	(
	2.1	Lines	
	2.2	Circles	
	2.3	Triangles	
	2.4	Polygon	
	2.5	Convex hull	1
	2.6	3D Convex Hull	1
	2.7	Slow Delaunay triangulation	1
	2.8	Closest Pair	1
	2.9	Rotating Caliphers	1
3	Nun	nerical algorithms	1
o	3.1	Fast exponentiation	1
	3.2	Prime numbers	1
	3.3	Miller-Rabin Primality Test (C)	1
	3.4	Number theory (modular, Chinese remainder, linear Diophantine)	1
	3.5	Systems of linear equations, matrix inverse, determinant	1
	3.6	Reduced row echelon form, matrix rank	1
	3.6 3.7		1
	3.7	Simplex algorithm	1
	3.7 3.8	Simplex algorithm	1
	3.7 3.8 3.9 3.10	Simplex algorithm	1 1 1 1
4	3.7 3.8 3.9 3.10	Simplex algorithm	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
4	3.7 3.8 3.9 3.10 Gra j	Simplex algorithm	1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1
4	3.7 3.8 3.9 3.10 Gra] 4.1 4.2	Simplex algorithm	1 1 1 1 1 1 1
4	3.7 3.8 3.9 3.10 Gra] 4.1 4.2 4.3	Simplex algorithm	1,11,11,11,11,11,11,11,11,11,11,11,11,1
4	3.7 3.8 3.9 3.10 Gra j 4.1 4.2 4.3 4.4	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number Ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
4	3.7 3.8 3.9 3.10 Gra] 4.1 4.2 4.3 4.4 4.5	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees	1 1 1 1 1 1 1 1 1 1
4	3.7 3.8 3.9 3.10 Gra j 4.1 4.2 4.3 4.4 4.5 4.6	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number Ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor	1; 1; 1; 1; 1; 1; 1; 1; 1; 1; 1; 1;
4	3.7 3.8 3.9 3.10 Gra] 4.1 4.2 4.3 4.4 4.5	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees	1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra j 4.1 4.2 4.3 4.4 4.5 4.6 4.7	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points	1 1 1 1 1 1 1 1 1 1 1 1
4 5	3.7 3.8 3.9 3.10 Gra y 4.1 4.2 4.3 4.4 4.5 4.6 4.7	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra j 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Data	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra y 4.1 4.2 4.3 4.4 4.5 4.6 4.7	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Dat: 5.1 5.2 5.3	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree Union-find	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra j 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Data 5.1 5.2	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree Union-find KD-tree	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Data 5.1 5.2 5.3 5.4 5.5	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree Union-find KD-tree Splay tree	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Dat: 5.1 5.2 5.3 5.4 5.5 5.6	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree Union-find KD-tree Splay tree Splay tree Splay Link Cut Trees	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Data 5.1 5.2 5.3 5.4 5.5	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree Union-find KD-tree Splay tree Splay Link Cut Trees	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	3.7 3.8 3.9 3.10 Gra 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Data 5.1 5.2 5.3 5.4 5.5 5.6 5.7	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree Union-find KD-tree Splay tree Splay tree Splay tree Splay Link Cut Trees Sparse Table	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 2 2 2
	3.7 3.8 3.9 3.10 Gray 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Data 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 String	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree Union-find KD-tree Splay tree Splay tree Splay Link Cut Trees Sparse Table Lazy Segment Tree	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 2 2 2 2
5	3.7 3.8 3.9 3.10 Gray 4.1 4.2 4.3 4.4 4.5 4.6 4.7 Data 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8	Simplex algorithm Fast Fourier transform (C++) Pollard Rho Algorithm Big Number ph algorithms Dijkstra's algorithm Strongly connected components Eulerian path Kruskal's algorithm Minimum spanning trees Lowest common ancestor Bridge and Articulation Points a structures Binary Indexed Tree 2D Binary Indexed Tree Union-find KD-tree Splay tree Splay tree Splay tree Sparse Table Lazy Segment Tree	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 2 2 2

22
2
2
2
2
2
2
23
23
2
23

1 Combinatorial optimization

1.1 Dense max-flow

```
// Adjacency matrix implementation of Dinic's blocking flow algorithm.
// Running time:
       0(|V|^4)
// INPUT:
       - graph, constructed using AddEdge()
// OUTPUT:
       - maximum flow value
       - To obtain the actual flow, look at positive values only.
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
struct MaxFlow {
  int N;
  VVI cap, flow;
  VI dad, Q;
   N(N), cap(N, VI(N)), flow(N, VI(N)), dad(N), Q(N) {}
  void AddEdge(int from, int to, int cap) {
   this->cap[from][to] += cap;
  int BlockingFlow(int s, int t) {
    fill(dad.begin(), dad.end(), -1);
    dad[s] = -2;
    int head = 0, tail = 0;
    Q[tail++] = s;
    while (head < tail) {
      int x = Q[head++];
for (int i = 0; i < N; i++) {</pre>
        if (dad[i] == -1 && cap[x][i] - flow[x][i] > 0) {
           dad[i] = x;
           Q[tail++] = i;
    if (dad[t] == -1) return 0;
    int totflow = 0;
    for (int i = 0; i < N; i++) {
      if (dad[i] == -1) continue;
      int amt = cap[i][t] - flow[i][t];
      for (int j = i; amt && j != s; j = dad[j])
  amt = min(amt, cap[dad[j]][j] - flow[dad[j]][j]);
      if (amt == 0) continue;
      il (amt -- 0, County)
flow[i][t] += amt;
flow[t][i] -= amt;
for (int j = i; j!= s; j = dad[j]) {
    flow[dad[j]][j] += amt;
}
         flow[j][dad[j]] -= amt;
      totflow += amt;
```

```
return totflow;
 int GetMaxFlow(int source, int sink) {
   int totflow = 0;
    while (int flow = BlockingFlow(source, sink))
     totflow += flow;
    return totflow;
int main() {
 mf.AddEdge(0, 1, 3);
 mf.AddEdge(0, 2, 4);
mf.AddEdge(0, 3, 5);
 mf.AddEdge(0, 4, 5);
 mf.AddEdge(1, 2, 2);
 mf.AddEdge(2, 3, 4);
 mf.AddEdge(2, 4, 1);
 mf.AddEdge(3, 4, 10);
  // should print out "15"
 cout << mf.GetMaxFlow(0, 4) << endl;
// The following code solves SPOJ problem #203: Potholers (POTHOLE)
#ifdef COMMENT
int main() {
 int t:
 cin >> t;
 for (int i = 0; i < t; i++) {
   int n;
    cin >> n;
    MaxFlow mf(n);
    for (int j = 0; j < n-1; j++) {
     int m;
      cin >> m;
      for (int k = 0; k < m; k++) {
       int p;
       cin >> p;
        int cap = (j == 0 || p == n-1) ? 1 : INF;
        mf.AddEdge(j, p, cap);
   cout << mf.GetMaxFlow(0, n-1) << endl;
 return 0;
#endif
// END CUT
```

1.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
                      O(|V|^3) augmentations
     max flow.
      min cost max flow: O(|V|^4 * MAX_EDGE_COST) augmentations
// TNPIIT ·
      - graph, constructed using AddEdge()
      - source
      - sink
// OUTPUT:
      - (maximum flow value, minimum cost value)
      - To obtain the actual flow, look at positive values only.
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
```

```
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow
 int N;
  VVL cap, flow, cost;
 VI found;
 VL dist, pi, width;
 VPII dad;
 MinCostMaxFlow(int N) :
   N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
   this->cap[from][to] = cap;
    this->cost[from][to] = cost;
  void Relax(int s, int k, L cap, L cost, int dir) {
    L \text{ val} = \text{dist}[s] + \text{pi}[s] - \text{pi}[k] + \text{cost};
    if (cap && val < dist[k]) {
     dist[k] = val;
      dad[k] = make_pair(s, dir);
     width[k] = min(cap, width[s]);
  L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
     int best = -1;
      found[s] = true;
      for (int k = 0; k < N; k++) {
       if (found[k]) continue;
       Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
        Relax(s, k, flow[k][s], -cost[k][s], -1);
       if (best == -1 || dist[k] < dist[best]) best = k;</pre>
     s = best:
    for (int k = 0: k < N: k++)
     pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
   L totflow = 0, totcost = 0;
   while (L amt = Dijkstra(s, t)) {
  totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
          totcost += amt * cost[dad[x].first][x];
        else (
          flow[x][dad[x].first] -= amt;
          totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
};
// BEGIN CUT
// The following code solves UVA problem #10594: Data Flow
int main() {
 int N, M;
  while (scanf("%d%d", &N, &M) == 2) {
   VVL v(M, VL(3));
    for (int i = 0; i < M; i++)
     scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
    scanf("%Ld%Ld", &D, &K);
    MinCostMaxFlow mcmf(N+1);
    for (int i = 0; i < M; i++) {
     mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i][2]);
     mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i][2]);
    mcmf.AddEdge(0, 1, D, 0);
```

```
pair<L, L> res = mcmf.GetMaxFlow(0, N);
if (res.first == D) {
    printf("%Ld\n", res.second);
} else {
    printf("Impossible.\n");
}
return 0;
}
```

1.3 Min-cost matching

```
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
     cost[i][j] = cost for pairing left node i with right node j
     Lmate[i] = index of right node that left node i pairs with
     Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
  int n = int(cost.size());
  // construct dual feasible solution
  VD v(n):
  for (int i = 0; i < n; i++) {
    u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
  for (int j = 0; j < n; j++) {
    v[j] = cost[0][j] - u[0];
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
  // construct primal solution satisfying complementary slackness
  Lmate = VI(n, -1);
  Rmate = VI(n, -1);
  int mated = 0;
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {</pre>
      if (Rmate[j] != -1) continue;
      if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {
        Lmate[i] = j;
        Rmate[j] = i;
        mated++;
        break:
  VD dist(n):
  VI dad(n):
  VI seen(n);
  // repeat until primal solution is feasible
  while (mated < n) {
    // find an unmatched left node
    int s = 0:
    while (Lmate[s] !=-1) s++;
    // initialize Dijkstra
    fill(dad.begin(), dad.end(), -1);
    fill(seen.begin(), seen.end(), 0);
    for (int k = 0; k < n; k++)
     dist[k] = cost[s][k] - u[s] - v[k];
```

```
int j = 0;
  while (true) {
    // find closest
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      if (j == -1 || dist[k] < dist[j]) j = k;</pre>
    seen[j] = 1;
    // termination condition
    if (Rmate[j] == -1) break;
    // relax neighbors
    const int i = Rmate[j];
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
      if (dist[k] > new_dist) {
        dist[k] = new_dist;
        dad[k] = j;
  // update dual variables
  for (int k = 0; k < n; k++)
   if (k == j || !seen[k]) continue;
const int i = Rmate[k];
    v[k] += dist[k] - dist[j];
    u[i] -= dist[k] - dist[j];
 u[s] += dist[j];
  // augment along path
  while (dad[j] >= 0) {
    const int d = dad[j];
    Rmate[i] = Rmate[d];
    Lmate[Rmate[i]] = i;
    j = d;
  Rmate[i] = s;
  Lmate[s] = j;
 mated++:
double value = 0;
for (int i = 0; i < n; i++)
 value += cost[i][Lmate[i]];
return value;
```

1.4 Sparse max-flow

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
// Running time:
      O(IVI^2 IEI)
// INPUT:
      - graph, constructed using AddEdge()
      - source and sink
// OUTPUT:
      - maximum flow value
      - To obtain actual flow values, look at edges with capacity > 0
        (zero capacity edges are residual edges).
typedef long long LL;
struct Edge {
 int u, v;
 LL cap, flow;
 Edge(int u, int v, LL cap): u(u), v(v), cap(cap), flow(0) {}
};
struct Dinic {
 int N;
 vector<Edge> E;
 vector<vector<int>> q;
```

```
vector<int> d, pt;
 Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}
  void AddEdge(int u, int v, LL cap) {
     E.emplace_back(u, v, cap);
      g[u].emplace_back(E.size() - 1);
      E.emplace_back(v, u, 0);
     g[v].emplace_back(E.size() - 1);
 bool BFS(int S, int T) {
   queue<int> q({S});
    fill(d.begin(), d.end(), N + 1);
    while(!q.empty()) {
     int u = q.front(); q.pop();
      if (u == T) break;
      for (int k: g[u]) {
        Edge &e = E[k];
        if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
          d[e.v] = d[e.u] + 1;
          q.emplace(e.v);
    return d[T] != N + 1;
  LL DFS (int u, int T, LL flow = -1) {
   if (u == T || flow == 0) return flow;
    for (int &i = pt[u]; i < g[u].size(); ++i) {</pre>
     Edge &e = E[q[u][i]];
     Edge &oe = E[g[u][i]^1];
if (d[e.v] == d[e.u] + 1) {
       LL amt = e.cap - e.flow;
        if (flow !=-1 && amt > flow) amt = flow;
        if (LL pushed = DFS(e.v, T, amt)) {
         e.flow += pushed;
          oe.flow -= pushed;
          return pushed;
    return 0:
  LL MaxFlow(int S, int T) {
   I.I. total = 0:
    while (BFS(S, T)) {
     fill(pt.begin(), pt.end(), 0);
      while (LL flow = DFS(S, T))
       total += flow;
    return total:
};
// BEGIN CUT
// The following code solves SPOJ problem #4110: Fast Maximum Flow (
     FASTFIOW)
int main()
 int N. E:
  scanf("%d%d", &N, &E);
 Dinic dinic(N);
  for(int i = 0; i < E; i++)</pre>
   int u, v;
   LL cap;
scanf("%d%d%lld", &u, &v, &cap);
    dinic.AddEdge(u - 1, v - 1, cap);
    dinic.AddEdge(v - 1, u - 1, cap);
  printf("%11d\n", dinic.MaxFlow(0, N - 1));
  return 0;
// END CUT
```

1.5 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// Running time:
      0(|V|^3)
// INPUT:
       - graph, constructed using AddEdge()
// OUTPUT:
       - (min cut value, nodes in half of min cut)
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
  int N = weights.size();
  VI used(N), cut, best_cut;
  int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
    VI w = weights[0];
    VI added = used;
    int prev, last = 0;
    for (int i = 0; i < phase; i++) {
      prev = last;
      last = -1;
      for (int j = 1; j < N; j++)
       if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][</pre>
        for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j</pre>
        used[last] = true;
        cut.push_back(last);
        if (best_weight == -1 || w[last] < best_weight) {</pre>
          best_cut = cut;
          best weight = w[last];
      } else {
        for (int j = 0; j < N; j++)
         w[j] += weights[last][j];
        added[last] = true;
  return make pair (best weight, best cut);
// BEGIN CUT
// The following code solves UVA problem #10989: Bomb, Divide and
     Conquer
int main() {
 int N;
  cin >> N:
  for (int i = 0; i < N; i++) {
   int n. m:
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
      int a, b, c;
      cin >> a >> b >> c;
      weights[a-1][b-1] = weights[b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
    cout << "Case #" << i+1 << ": " << res.first << endl;
// END CUT
```

1.6 Push-relabel max-flow

```
// Adjacency list implementation of FIFO push relabel maximum flow
// with the gap relabeling heuristic. This implementation is
// significantly faster than straight Ford-Fulkerson. It solves
// random problems with 10000 vertices and 1000000 edges in a few
// seconds, though it is possible to construct test cases that
// achieve the worst-case.
//
// Running time:
// O(|V|^3)
//
// INPUT:
// - graph, constructed using AddEdge()
```

```
- source
      - sink
// OUTPUT:
       - maximum flow value
       - To obtain the actual flow values, look at all edges with
        capacity > 0 (zero capacity edges are residual edges).
typedef long long LL;
 int from, to, cap, flow, index;
 Edge (int from, int to, int cap, int flow, int index) :
   from (from), to (to), cap (cap), flow (flow), index (index) {}
struct PushRelabel {
 int N;
  vector<vector<Edge> > G;
  vector<LL> excess;
  vector<int> dist, active, count;
  queue<int> 0:
 PushRelabel(int N) : N(N), G(N), excess(N), dist(N), active(N),
       count (2*N) {}
  void AddEdge(int from, int to, int cap) {
   G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
   G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
  void Enqueue(int v) {
   if (!active[v] && excess[v] > 0) { active[v] = true; Q.push(v); }
  void Push (Edge &e) {
   int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
   if (dist[e.from] <= dist[e.to] || amt == 0) return;</pre>
   e.flow += amt;
   G[e.to][e.index].flow -= amt;
   excess[e.to] += amt;
   excess[e.from] -= amt;
   Enqueue (e.to);
 void Gap(int k) {
   for (int v = 0; v < N; v++) {
     if (dist[v] < k) continue;</pre>
      count[dist[v]]--;
      dist[v] = max(dist[v], N+1);
      count[dist[v]]++;
      Enqueue (v);
  void Relabel(int v) {
   count[dist[v]]--:
    dist[v] = 2*N;
   for (int i = 0; i < G[v].size(); i++)</pre>
     if (G[v][i].cap - G[v][i].flow > 0)
       dist[v] = min(dist[v], dist[G[v][i].to] + 1);
   count[dist[v]]++;
   Enqueue (v);
  void Discharge(int v) {
   for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[v][i</pre>
   if (excess[v] > 0) {
     if (count[dist[v]] == 1)
       Gap(dist[v]);
      else
       Relabel(v);
 LL GetMaxFlow(int s, int t) {
   count[0] = N-1;
   count[N] = 1;
    dist[s] = N;
    active[s] = active[t] = true;
    for (int i = 0; i < G[s].size(); i++) {</pre>
      excess[s] += G[s][i].cap;
      Push(G[s][i]);
    while (!Q.empty()) {
     int v = Q.front();
      Q.pop();
```

```
active[v] = false;
     Discharge(v);
   LL totflow = 0;
   for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;</pre>
   return totflow;
// The following code solves SPOJ problem #4110: Fast Maximum Flow (
 int n, m;
 scanf("%d%d", &n, &m);
 PushRelabel pr(n);
 for (int i = 0; i < m; i++) {
   scanf("%d%d%d", &a, &b, &c);
   if (a == b) continue;
   pr.AddEdge(a-1, b-1, c);
   pr.AddEdge(b-1, a-1, c);
 printf("%Ld\n", pr.GetMaxFlow(0, n-1));
 return 0;
// END CUT
```

1.7 Max bipartite matching

```
// This code performs maximum bipartite matching.
// Running time: O(|E| |V|) -- often much faster in practice
    INPUT: w[i][j] = edge between row node i and column node j
    OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
             mc[j] = assignment for column node j, -1 if unassigned
             function returns number of matches made
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
 for (int j = 0; j < w[i].size(); j++) {
   if (w[i][j] && !seen[j]) {
    seen[j] = true;
     if (mc[j] < 0 || FindMatch(mc[j], w, mr, mc, seen)) {</pre>
       mr[i] = j;
       mc[i] = i;
        return true:
  return false;
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
 mr = VI(w.size(), -1);
 mc = VI(w[0].size(), -1);
 int ct = 0:
 for (int i = 0; i < w.size(); i++) {</pre>
    VI seen(w[0].size());
    if (FindMatch(i, w, mr, mc, seen)) ct++;
 return ct:
1. max matching should have been found first
2. change each edge in matching into a directed edge from right to
     left
3. change each edge not used in matching into a directed edge from
     left to right
4. compute T: set of vertices reachable from unmatched vertices on the
       left (including themselves)
5. \mbox{MVC} = vertices cover consists of all vertices on the right that are
      in T, and all vertices on the left that are not in \tilde{T}
// OUTPUT: <row/col idx, 0:row/1:col>
const int LEFT = 0;
```

```
const int RIGHT = 1;
vii MVC(vvi &w, vi &mr, vi &mc) {
    set<ii>> T;
    queue<ii> q;
    for(int i=0;i<mr.size();i++){</pre>
        if (mr[i]==-1) {
            q.push({i, LEFT});
            T.insert({i, LEFT});
    while(!q.empty()) {
        ii curr = q.front(); q.pop();
         int u = curr.first;
        int type = curr.second;
        if(type == LEFT) {
            for (int v=0; v<mc.size(); v++) {</pre>
                 ii next = {v, RIGHT};
                 if (w[u][v] && mr[u]!=v && !T.count(next)) {
                     T.insert (next);
                     q.push(next);
        } else {
             // RIGHT
            for (int v=0; v<mr.size(); v++) {</pre>
                 ii next = {v, LEFT};
                 if (w[v][u] && mr[v] == u && !T.count(next))
                    T.insert(next);
                     q.push(next);
    for(int i=0;i<mr.size();i++) if(!T.count({i, LEFT})) mvc.push_back</pre>
          ({i, LEFT});
    for(int i=0;i<mc.size();i++) if(T.count({i, RIGHT})) mvc.push_back</pre>
          ({i, RIGHT});
    return mvc;
```

1.8 Graph cut inference

```
// Special-purpose {0,1} combinatorial optimization solver for
// problems of the following by a reduction to graph cuts:
                            sum i psi i(x[i])
   x[1]...x[n] in \{0,1\} + sum_{\{i < j\}} phi_{ij}(x[i], x[j])
// where
        psi_i : {0, 1} --> R
     phi_{ij} : \{0, 1\} \times \{0, 1\} \longrightarrow R
// such that
    phi_{ij}(0,0) + phi_{ij}(1,1) \le phi_{ij}(0,1) + phi_{ij}(1,0)
// This can also be used to solve maximization problems where the
// direction of the inequality in (*) is reversed.
// INPUT: phi -- a matrix such that phi[i][j][u][v] = phi_{ij}(u, v) // psi -- a matrix such that psi[i][u] = psi_i(u)
           x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution
// To use this code, create a GraphCutInference object, and call the
// DoInference() method. To perform maximization instead of
      minimization.
// ensure that #define MAXIMIZATION is enabled.
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef vector<VVI> VVVI;
typedef vector<VVVI> VVVVI;
const int TNF = 10000000000:
// comment out following line for minimization
#define MAXIMIZATION
struct GraphCutInference {
  int N;
```

```
VVI cap, flow;
  VI reached;
  int Augment(int s, int t, int a) {
    reached[s] = 1;
    if (s == t) return a;
    for (int k = 0; k < N; k++) {
      if (reached[k]) continue;
      if (int aa = min(a, cap[s][k] - flow[s][k])) {
        if (int b = Augment(k, t, aa)) {
          flow[s][k] += b;
           flow[k][s] = b;
           return b;
    return 0;
  int GetMaxFlow(int s, int t) {
    N = cap.size();
    flow = VVI(N, VI(N));
    reached = VI(N);
    int totflow = 0;
    while (int amt = Augment(s, t, INF)) {
      fill(reached.begin(), reached.end(), 0);
    return totflow;
  int DoInference (const VVVVI &phi, const VVI &psi, VI &x) {
    int M = phi.size();
    cap = VVI(M+2, VI(M+2));
    VI b(M);
    int c = 0;
    for (int i = 0; i < M; i++) {
      b[i] += psi[i][1] - psi[i][0];
      c += psi[i][0];
      for (int j = 0; j < i; j++)</pre>
      for (int j = 0; j < 1; j++;
b[i] += phi[i][j][j][1] - phi[i][j][0][1];
for (int j = i+1; j < M; j++) {
    cap[i][j] = phi[i][j][0][1] + phi[i][j][1][0] - phi[i][j]
        ][0][0] - phi[i][j][1][1];
b[i] += phi[i][j][1][0] - phi[i][j][0][0];
        c += phi[i][j][0][0];
#ifdef MAXIMIZATION
    for (int i = 0; i < M; i++) {
      for (int j = i+1; j < M; j++)
        cap[i][j] *= -1;
      b[i] *= -1;
    c *= -1:
#endif
    for (int i = 0; i < M; i++) {
      if (b[i] >= 0) {
        cap[M][i] = b[i];
      else
        cap[i][M+1] = -b[i];
        c += b[i];
    int score = GetMaxFlow(M, M+1);
    fill(reached.begin(), reached.end(), 0);
    Augment (M, M+1, INF);
    x = VI(M);
    for (int i = 0; i < M; i++) x[i] = reached[i] ? 0 : 1;</pre>
    score += c:
#ifdef MAXIMIZATION
    score \star = -1;
#endif
    return score;
int main() {
  // solver for "Cat vs. Dog" from NWERC 2008
  int numcases;
```

```
cin >> numcases;
for (int caseno = 0; caseno < numcases; caseno++) {
 int c, d, v;
 cin >> c >> d >> v;
 VVVVI phi(c+d, VVVI(c+d, VVI(2, VI(2))));
 VVI psi(c+d, VI(2));
 for (int i = 0; i < v; i++) {
   char p, q;
   int u, v;
   cin >> p >> u >> q >> v;
   if (p == 'C') {
     phi[u][c+v][0][0]++;
      phi[c+v][u][0][0]++;
     phi[v][c+u][1][1]++;
      phi[c+u][v][1][1]++;
 GraphCutInference graph;
 cout << graph.DoInference(phi, psi, x) << endl;</pre>
return 0:
```

1.9 General Matching

```
// Unweighted general matching.
// Vertices are numbered from 1 to V.
// G is an adjlist.
// G[x][0] contains the number of neighbours of x.
// The neigbours are then stored in G[x][1] .. G[x][G[x][0]].
// Mate[x] will contain the matching node for x.
// V and E are the number of edges and vertices.
// Slow Version (2x on random graphs) of Gabow's implementation
// of Edmonds' algorithm (O(V^3)).
const int MAXV = 250;
int G[MAXV][MAXV];
int VLabel[MAXV];
int Queue[MAXV];
int Mate[MAXV];
      Save[MAXV];
int
      Used[MAXV];
int
int
       Up, Down;
int
void ReMatch(int x, int y)
  int m = Mate[x]; Mate[x] = y;
 if (Mate[m] == x)
      if (VLabel[x] <= V)</pre>
          Mate[m] = VIabel[x]:
          ReMatch(VLabel[x], m);
      else
          int a = 1 + (VLabel[x] - V - 1) / V;
          int b = 1 + (VLabel[x] - V - 1) % V;
          ReMatch(a, b); ReMatch(b, a);
void Traverse(int x)
  for (int i = 1; i <= V; i++) Save[i] = Mate[i];</pre>
  ReMatch(x, x);
  for (int i = 1; i <= V; i++)</pre>
      if (Mate[i] != Save[i]) Used[i]++;
      Mate[i] = Save[i];
void ReLabel(int x, int y)
  for (int i = 1; i <= V; i++) Used[i] = 0;</pre>
 Traverse(x); Traverse(y);
```

```
for (int i = 1; i <= V; i++)
      if (Used[i] == 1 && VLabel[i] < 0)</pre>
          VLabel[i] = V + x + (y - 1) * V;
          Queue[Up++] = i;
// Call this after constructing G
void Solve()
  for (int i = 1; i <= V; i++)</pre>
    if (Mate[i] == 0)
        for (int j = 1; j <= V; j++) VLabel[j] = -1;
        VLabel[i] = 0; Down = 1; Up = 1; Queue[Up++] = i;
        while (Down != Up)
            int x = Queue[Down++];
            for (int p = 1; p <= G[x][0]; p++)
                int y = G[x][p];
                if (Mate[y] == 0 && i != y)
                    Mate[y] = x; ReMatch(x, y);
                    Down = Up; break;
                if (VLabel[y] >= 0)
                    ReLabel(x, y);
                    continue;
                if (VLabel[Mate[y]] < 0)</pre>
                    VLabel[Mate[y]] = x;
                    Queue[Up++] = Mate[y];
// Call this after Solve(). Returns number of edges in matching (half
      the number of matched vertices)
int Size()
 int Count = 0;
 for (int i = 1; i <= V; i++)
   if (Mate[i] > i) Count++;
  return Count;
```

1.10 Edmond Blossom Matching

```
#include <algorithm>
#include <iostream>
#include <queue>
#include <vector>
class edmond blossom {
private:
    std::vector<std::vector<int>> adi list;
    std::vector<int> match:
    std::vector<int> parent;
    std..vector<int> blossom root:
    std::vector<bool> in_queue;
    std..vector<bool> in blossom:
    std::queue<int> process_queue;
    int num v:
    /// Find the lowest common ancestor between u and v.
    /// The specified root represents an upper bound.
    int get_lca(int root, int u, int v) {
        std::vector<bool> in_path(num_v, false);
        for (u = blossom_root[u]; ; u = blossom_root[parent[match[u
            in_path[u] = true;
            if (u == root) {
               break:
```

```
for (v = blossom_root[v]; ; v = blossom_root[parent[match[v
        if (in_path[v]) {
           return v;
/// Mark the vertices between u and the specified lowest
/// common ancestor for contraction where necessary.
void mark_blossom(int lca, int u) {
   while (blossom_root[u] != lca) {
       int v = match[u];
        in_blossom[blossom_root[u]] = true;
        in_blossom[blossom_root[v]] = true;
       if (blossom_root[u] != lca) {
           parent[u] = v;
/// Contract the blossom that is formed after processing
void contract_blossom(int source, int u, int v) {
   int lca = get_lca(source, u, v);
   std::fill(in_blossom.begin(), in_blossom.end(), false);
   mark_blossom(lca, u);
   mark_blossom(lca, v);
   if (blossom_root[u] != lca) {
       parent[u] = v;
   if (blossom_root[v] != lca) {
       parent[v] = u;
   for (int i = 0; i < num v; ++i) {</pre>
       if (in blossom[blossom root[i]]) {
            blossom_root[i] = lca;
           if (!in_queue[i]) {
               process_queue.push(i);
               in_queue[i] = true;
/// Return the vertex at the end of an augmenting path
/// starting at the specified source, or -1 if none exist.
int find_augmenting_path(int source) {
   for (int i = 0; i < num_v; ++i) {
   in_queue[i] = false;</pre>
        parent[i] = -1;
        blossom_root[i] = i;
   // Empty the gueue
   process_queue = std::queue<int>();
   process_queue.push(source);
   in_queue[source] = true;
   while (!process_queue.empty()) {
       int u = process_queue.front();
       process_queue.pop();
        for (int v : adj_list[u]) {
            if (blossom_root[u] != blossom_root[v] && match[u] !=
                // Process if
                // + u-v is not an edge in the matching
                // && u and v are not in the same blossom (yet)
                if (v == source || (match[v] != -1 && parent[match
                      [v]] != -1)) {
                    // Contract a blossom if
                    // + v is the source
                    // || v is matched and v's match has a parent.
                    // The fact that parents are assigned to
                          vertices
                    // with odd distances from the source is used
                    // check if a cycle is odd or even. u is
                          always an
                    // even distance away from the source, so if v
```

```
// match is assigned a parent, you have an odd
                               cycle.
                        contract_blossom(source, u, v);
                    } else if (parent[v] == -1) {
                        parent[v] = u;
                        if (match[v] == -1) {
                            // v is unmatched; augmenting path found
                            return v;
                            // Enqueue v's match.
                            int w = match[v];
                            if (!in_queue[w]) {
                                process_queue.push(w);
                                in_queue[w] = true;
        return -1:
    /// Augment the path that ends with the specified vertex
    /// using the parent and match fields. Returns the increase
    /// in the number of matchings. (i.e. 1 if the path is valid,
    /// 0 otherwise)
    int augment_path(int end) {
        int u = end;
        while (u != -1) {
            // Currently w===v---u
            int v = parent[u];
            int w = match[v];
            // Change to w---v===u
           match[v] = u;
           match[u] = v;
           u = w;
        // Return 1 if the augmenting path is valid
        return end == -1 ? 0 : 1;
public:
    edmond blossom(int v) :
        adj list(v),
        match(v, -1),
        parent(v),
        blossom root(v),
        in queue(v).
        in blossom(v).
        num v(v) {}
    /// Add a bidirectional edge from u to v.
    void add_edge(int u, int v) {
        adj_list[u].push_back(v);
        adj_list[v].push_back(u);
    /// Returns the maximum cardinality matching
    int get_max_matching() {
        int ans = 0:
        // Reset
        std::fill(match.begin(), match.end(), -1);
        for (int u = 0; u < num_v; ++u) {</pre>
           if (match[u] == -1) {
                int v = find_augmenting_path(u);
                if (v != −1) {
                    // An augmenting path exists
                    ans += augment_path(v);
        return ans;
    /// Constructs the maximum cardinality matching
    std::vector<std::pair<int, int>> construct_matching() {
        std::vector<std::pair<int, int>> output;
```

```
std::vector<bool> is_processed(num_v, false);
        for (int u = 0; u < num_v; ++u) {
           if (!is_processed[u] && match[u] != -1) {
                output.emplace_back(u, match[u]);
                is_processed[u] = true;
                is_processed[match[u]] = true;
        return output;
int main() {
    10 18
   0 2
   1 2
   std::ios_base::sync_with_stdio(false);
   std::cin.tie(nullptr);
   int num vertices;
   int num edges;
   std::cin >> num vertices >> num edges;
    edmond blossom eb(num vertices):
   for (int i = 0; i < num_edges; ++i) {</pre>
       int 11. v:
       std::cin >> u >> v:
       eb.add edge(u, v);
    std::cout << "Maximum Cardinality: " << eb.get_max_matching() << "
   std::vector<std::pair<int, int>> matching = eb.construct matching
    for (auto@ match : matching) {
       std::cout << match.first << " " << match.second << "\n";
```

1.11 Min Edge Cover

/* If a minimum edge cover contains C edges, and a maximum matching contains M edges, then C + M = jV j. To obtain the edge cover, start with a maximum matching, and then, for every vertex not matched, just select some edge incident upon it and add it to the edge set */

1.12 Stable Marriage Problem

```
// Gale-Shapley algorithm for the stable marriage problem.
// madj[i][j] is the jth highest ranked woman for man i.
// fpref[i][j] is the rank woman i assigns to man j.
// Returns a pair of vectors (mpart, fpart), where mpart[i] gives the partner of man i, and fpart is analogous
```

```
pair<vector<int>, vector<int> > stable_marriage(vector<vector<int> >&
      madj, vector<vector<int> >& fpref) {
       int n = madj.size();
       vector<int> mpart(n, -1), fpart(n, -1);
        vector<int> midx(n);
        queue<int> mfree;
        for (int i = 0; i < n; i++) {
               mfree.push(i);
        while (!mfree.empty()) {
                int m = mfree.front(); mfree.pop();
                int f = madj[m][midx[m]++];
                if (fpart[f] == -1) {
                       mpart[m] = f; fpart[f] = m;
                } else if (fpref[f][m] < fpref[f][fpart[f]]) {</pre>
                       mpart[fpart[f]] = -1; mfree.push(fpart[f]);
                        mpart[m] = f; fpart[f] = m;
                        mfree.push(m);
        return make_pair(mpart, fpart);
```

2 Geometry

2.1 Lines

```
#include <algorithm>
#include <cstdio>
#include <cmath>
using namespace std;
#define INF 1e9
#define EPS 1e-9
#define PI acos(-1.0) // important constant; alternative #define PI
      (2.0 * acos(0.0))
double DEG_to_RAD(double d) { return d * PI / 180.0; }
double RAD to DEG(double r) { return r * 180.0 / PI; }
// struct point_i { int x, y; }; // basic raw form, minimalist mode
struct point_i { int x, y; // whenever possible, work with point_i
                                                 // default constructor
 point i() \{ x = y = 0; \}
 point_i(int _x, int _y) : x(_x), y(_y) {} };
                                                         // user-defined
\textbf{struct} \text{ point } \{ \text{ } \textbf{double} \text{ } \textbf{x, } \textbf{y;} \text{ } \textit{// only used if more precision is needed}
                                                 // default constructor
 point() { x = y = 0.0; }
  point(double _x, double _y) : x(_x), y(_y) {}
                                                         // user-defined
 bool operator < (point other) const { // override less than operator</pre>
                                                   // useful for sorting
   if (fabs(x - other.x) > EPS)
                                    // first criteria , by x-coordinate
     return x < other.x;
  return y < other.y; } // second criteria, by y-coord y use EPS (le-9) when testing equality of two floating points
                                    // second criteria, by y-coordinate
  bool operator == (point other) const {
   return (fabs(x - other.x) < EPS && (fabs(y - other.y) < EPS)); } };
double dist(point p1, point p2) {
                                                   // Euclidean distance
                       // hypot(dx, dy) returns sqrt(dx * dx + dy * dy)
  return hypot (p1.x - p2.x, p1.y - p2.y); }
                                                        // return double
// rotate p by theta degrees CCW w.r.t origin (0, 0)
point rotate(point p, double theta) {
 double rad = DEG_to_RAD(theta);  // multiply theta with PI / 180.0
  return point (p.x * cos(rad) - p.y * sin(rad),
               p.x * sin(rad) + p.y * cos(rad)); }
struct line { double a, b, c; };
                                            // a wav to represent a line
// the answer is stored in the third parameter (pass by reference)
void pointsToLine(point p1, point p2, line &1) {
 if (fabs(p1.x - p2.x) < EPS) {</pre>
                                               // vertical line is fine
    1.a = 1.0; 1.b = 0.0; 1.c = -p1.x;
                                                        // default values
 else (
   1.a = -(double) (p1.y - p2.y) / (p1.x - p2.x);
                            // IMPORTANT: we fix the value of b to 1.0
   1.b = 1.0;
   1.c = -(double) (1.a * p1.x) - p1.y;
```

```
// not needed since we will use the more robust form: ax + by + c = 0
      (see above)
struct line2 { double m, c; };
                                 // another way to represent a line
int pointsToLine2(point p1, point p2, line2 &1) {
if (abs(p1.x - p2.x) < EPS) {</pre>
                                      // special case: vertical line
                               // 1 contains m = INF and c = x_value
  l.m = INF;
  1.c = p1.x;
                              // to denote vertical line x = x_value
   return 0; // we need this return variable to differentiate result
  1.m = (double) (p1.y - p2.y) / (p1.x - p2.x);
   1.c = p1.y - 1.m * p1.x;
   return 1; // 1 contains m and c of the line equation y = mx + c
bool areParallel(line 11, line 12) {
                                       // check coefficients a & b
  return (fabs(11.a-12.a) < EPS) && (fabs(11.b-12.b) < EPS); }
bool areSame(line 11, line 12) {
                                        // also check coefficient c
 return areParallel(11 ,12) && (fabs(11.c - 12.c) < EPS); }
// returns true (+ intersection point) if two lines are intersect
bool areIntersect(line 11, line 12, point &p) {
 if (areParallel(11, 12)) return false;
  // solve system of 2 linear algebraic equations with 2 unknowns
 p.x = (12.b * 11.c - 11.b * 12.c) / (12.a * 11.b - 11.a * 12.b);
  // special case: test for vertical line to avoid division by zero
  if (fabs(11.b) > EPS) p.y = -(11.a * p.x + 11.c);
                     p.y = -(12.a * p.x + 12.c);
 else
 return true: }
struct vec { double x, y; // name: 'vec' is different from STL vector
 vec(double _x, double _y) : x(_x), y(_y) {} };
vec toVec(point a, point b) {
                                  // convert 2 points to vector a->b
 return vec(b.x - a.x, b.y - a.y); }
vec scale(vec v, double s) {
                                  // nonnegative s = \{<1 ... 1 ... >1\}
 return vec(v.x * s, v.v * s); }
                                              // shorter.same.longer
point translate(point p, vec v) {
                                       // translate p according to v
 return point(p.x + v.x , p.y + v.y); }
// convert point and gradient/slope to line
void pointSlopeToLine(point p, double m, line &1) {
                                                        // always -m
 1.a = -m:
                                                         // always 1
 1.b = 1:
 1.c = -((1.a * p.x) + (1.b * p.y)); }
                                                      // compute this
void closestPoint(line 1, point p, point &ans) {
                          // perpendicular to 1 and pass through p
  line perpendicular:
 if (fabs(1.b) < EPS) {
                                    // special case 1: vertical line
   ans.x = -(1.c); ans.y = p.y;
                                      return; }
 if (fabs(l.a) < EPS) {
                                  // special case 2: horizontal line
   ans.x = p.x; ans.y = -(1.c); return; }
 pointSlopeToLine(p, 1 / 1.a, perpendicular);
                                                       // normal line
  // intersect line 1 with this perpendicular line
  // the intersection point is the closest point
 areIntersect(l, perpendicular, ans); }
// returns the reflection of point on a line
void reflectionPoint(line 1, point p, point &ans) {
 point b;
  closestPoint(1, p, b);
                                            // similar to distToLine
  vec v = toVec(p, b);
                                                 // create a vector
 ans = translate(translate(p, v), v); }
                                                // translate p twice
double dot(vec a, vec b) { return (a.x * b.x + a.y * b.y); }
double norm_sq(vec v) { return v.x * v.x + v.y * v.y; }
// returns the distance from p to the line defined by
// two points a and b (a and b must be different)
// the closest point is stored in the 4th parameter (byref)
double distToLine(point p, point a, point b, point &c) {
 // formula: c = a + u * ab
  vec ap = toVec(a, p), ab = toVec(a, b);
 double u = dot(ap, ab) / norm_sq(ab);
  c = translate(a, scale(ab, u));
                                                 // translate a to c
 return dist(p, c); }
                                // Euclidean distance between p and c
// returns the distance from p to the line segment ab defined by
// two points a and b (still OK if a == b)
// the closest point is stored in the 4th parameter (byref)
double distToLineSegment (point p, point a, point b, point &c) {
  vec ap = toVec(a, p), ab = toVec(a, b);
```

```
double u = dot(ap, ab) / norm_sq(ab);
 if (u < 0.0) { c = point(a.x, a.y);</pre>
                                                     // closer to a
   return dist(p, a); } // Euclidean distance between p and a
  return dist(p, b); } // Euclidean distance between p and b
  return distToLine(p, a, b, c); }
                                         // run distToline as above
double angle (point a, point o, point b) { // returns angle aob in rad
 vec oa = toVec(o, a), ob = toVec(o, b);
  return acos(dot(oa, ob) / sqrt(norm_sq(oa) * norm_sq(ob))); }
double cross(vec a, vec b) { return a.x * b.y - a.y * b.x; }
//// another variant
//int area2(point p, point q, point r) { // returns 'twice' the area
     of this triangle A-B-c
   return p.x * q.y - p.y * q.x +
          q.x * r.y - q.y * r.x +
          r.x * p.y - r.y * p.x;
1/3
// note: to accept collinear points, we have to change the '> 0'
// returns true if point r is on the left side of line pg
bool ccw(point p, point q, point r) {
 return cross(toVec(p, q), toVec(p, r)) > 0; }
// returns true if point r is on the same line as the line pg
bool collinear (point p, point q, point r) {
 return fabs(cross(toVec(p, q), toVec(p, r))) < EPS; }</pre>
int main() {
 point P1, P2, P3(0, 1); // note that both P1 and P2 are (0.00, 0.00)
 printf("%d\n", P1 == P2);
                                                            // true
 printf("%d\n", P1 == P3);
  vector<point> P:
 P.push_back(point(2, 2));
  P.push_back(point(4, 3));
 P.push back(point(2, 4));
 P.push back(point(6, 6));
 P.push back(point(2, 6));
 P.push back(point(6, 5));
  // sorting points demo
 sort (P.begin(), P.end());
 for (int i = 0; i < (int)P.size(); i++)</pre>
   printf("(%.21f, %.21f)\n", P[i].x, P[i].y);
 // rearrange the points as shown in the diagram below
 P.clear();
 P.push_back(point(2, 2));
 P.push back(point(4, 3));
 P.push back(point(2, 4));
 P.push back(point(6, 6));
 P.push back(point(2, 6));
 P.push back(point(6, 5));
 P.push back(point(8, 6));
 // the positions of these 7 points (0-based indexing)
  6 P4
          P3 P6
             P5
     P2
        P1
    P0
 0 1 2 3 4 5 6 7 8
  double d = dist(P[0], P[5]);
 printf("Euclidean distance between P[0] and P[5] = %.21f\n", d); //
        should be 5 000
 // line equations
 line 11, 12, 13, 14;
 pointsToLine(P[0], P[1], 11);
 printf("%.21f * x + %.21f * y + %.21f = 0.00\n", 11.a, 11.b, 11.c);
        // should be -0.50 * x + 1.00 * y - 1.00 = 0.00
 pointsToLine (P[0],\ P[2],\ 12);\ //\ a\ vertical\ line,\ not\ a\ problem\ in\ "
       ax + by + c = 0" representation
 printf("%.21f * x + %.21f * y + %.21f = 0.00\n", 12.a, 12.b, 12.c);
        // should be 1.00 * x + 0.00 * y - 2.00 = 0.00
 // parallel, same, and line intersection tests
 pointsToLine(P[2], P[3], 13);
 printf("11 & 12 are parallel? %d\n", areParallel(11, 12)); // no
 printf("11 & 13 are parallel? %d\n", areParallel(11, 13)); // yes,
```

11 (P[0]-P[1]) and 13 (P[2]-P[3]) are parallel

```
pointsToLine(P[2], P[4], 14);
printf("11 & 12 are the same? %d\n", areSame(11, 12)); // no
printf("12 & 14 are the same? %d\n", areSame(12, 14)); // yes, 12 (P
      [0]-P[2]) and 14 (P[2]-P[4]) are the same line (note, they are
       two different line segments, but same line)
bool res = areIntersect(11, 12, p12); // yes, 11 (P[0]-P[1]) and 12
      (P[0]-P[2]) are intersect at (2.0, 2.0)
printf("11 & 12 are intersect? %d, at (%.21f, %.21f)\n", res, p12.x,
// other distances
point ans;
d = distToLine(P[0], P[2], P[3], ans);
printf("Closest point from P[0] to line
                                               (P[21-P[31): (%.21f.
       %.21f), dist = %.21f\n", ans.x, ans.y, d);
closestPoint(13, P[0], ans);
printf("Closest point from P[0] to line V2
                                               (P[2]-P[3]): (%.21f,
       .21f, dist = .21fn", ans.x, ans.y, dist(P[0], ans));
d = distToLineSegment(P[0], P[2], P[3], ans);
printf("Closest point from P[0] to line SEGMENT (P[2]-P[3]): (%.21f,
       %.21f), dist = %.21f\n", ans.x, ans.y, d); // closer to A (or
       P[2]) = (2.00, 4.00)
d = distToLineSegment(P[1], P[2], P[3], ans);
printf("Closest point from P[1] to line SEGMENT (P[2]-P[3]): (%.21f,
      %.21f), dist = %.21f\n", ans.x, ans.y, d); // closer to
     midway between AB = (3.20, 4.60)
d = distToLineSegment(P[6], P[2], P[3], ans);
printf("Closest point from P[6] to line SEGMENT (P[2]-P[3]): (%.21f,
       %.21f), dist = %.21f\n", ans.x, ans.y, d); // closer to B (or
       P[3]) = (6.00, 6.00)
reflectionPoint(14, P[1], ans);
printf("Reflection point from P[1] to line
                                            (P[2]-P[4]): (%.21f,
       %.21f)\n", ans.x, ans.y); // should be (0.00, 3.00)
printf("Angle P[0]-P[4]-P[3] = %.21f\n", RAD to DEG(angle(P[0], P
     [4], P[3]))); // 90 degrees
printf("Angle P[0]-P[2]-P[1] = %.21f\n", RAD to DEG(angle(P[0], P
     [2], P[1]))); // 63.43 degrees
printf("Angle P[4]-P[3]-P[6] = %.21f\n", RAD_to_DEG(angle(P[4], P
     [3], P[6]))); // 180 degrees
printf("P[0], P[2], P[3] form A left turn? %d\n", ccw(P[0], P[2], P
     [31)): // no
printf("P[0], P[3], P[2] form A left turn? %d\n", ccw(P[0], P[3], P
     [2])); // yes
printf("P[0], P[2], P[3] are collinear? %d\n", collinear(P[0], P[2],
      P[3])); // no
printf("P[0], P[2], P[4] are collinear? %d\n", collinear(P[0], P[2], P[4])
      P[4])); // yes
point p(3, 7), q(11, 13), r(35, 30); // collinear if r(35, 31)
printf("r is on the %s of line p-r\n", ccw(p, q, r) ? "left" : "
     right"); // right
// the positions of these 6 points
  E<-- 4
               B D<--
        2 A C
-4-3-2-1 0 1 2 3 4 5 6
       - 1
       -2
F<-- -3
// translation
point A(2.0, 2.0);
point B(4.0, 3.0);
vec v = toVec(A, B); // imagine there is an arrow from A to B (see
     the diagram above)
point C(3.0, 2.0);
point D = translate(C, v); // D will be located in coordinate (3.0 +
      2.0, 2.0 + 1.0) = (5.0, 3.0)
printf("D = (%.21f, %.21f)\n", D.x, D.y);
point E = translate(C, scale(v, 0.5)); // E will be located in
     coordinate (3.0 + 1/2 * 2.0, 2.0 + 1/2 * 1.0) = (4.0, 2.5)
printf("E = (%.21f, %.21f)\n", E.x, E.y);
printf("B = (%.21f, %.21f)\n", B.x, B.y); // B = (4.0, 3.0)
point F = rotate(B, 90); // rotate B by 90 degrees COUNTER clockwise
     F = (-3.0, 4.0)
```

2.2 Circles

```
#include <cstdio>
#include <cmath>
using namespace std:
#define TNF 1e9
#define EPS 1e-9
#define PT acos (-1.0)
double DEG to RAD(double d) { return d * PI / 180.0; }
double RAD to DEG(double r) { return r * 180.0 / PI; }
point_i() { x = y = 0; }
                                              // default constructor
 point_i(int _x, int _y) : x(_x), y(_y) {} };
                                                      // constructor
struct point { double x, y; // only used if more precision is needed
 point() { x = y = 0.0; }
                                       // default constructor
  point(double _x, double _y) : x(_x), y(_y) {} }; // constructor
int insideCircle(point_i p, point_i c, int r) { // all integer version
 int dx = p.x - c.x, dy = p.y - c.y;
int Euc = dx * dx + dy * dy, rSq = r * r;
                                                      // all integer
 return Euc < rSq ? 0 : Euc == rSq ? 1 : 2; } //inside/border/outside
bool circle2PtsRad(point p1, point p2, double r, point &c) {
  double d2 = (p1.x - p2.x) * (p1.x - p2.x) +
            (p1.y - p2.y) * (p1.y - p2.y);
  double det = r * r / d2 - 0.25;
 if (det < 0.0) return false;</pre>
  double h = sqrt(det);
 c.x = (p1.x + p2.x) * 0.5 + (p1.y - p2.y) * h;
  c.y = (p1.y + p2.y) * 0.5 + (p2.x - p1.x) * h;
                    // to get the other center, reverse p1 and p2
 return true: }
 // circle equation, inside, border, outside
  point i pt(2, 2);
 int r = 7;
 point i inside(8, 2);
 printf("%d\n", insideCircle(inside, pt, r));
                                                          // 0-inside
 point i border (9, 2);
 printf("%d\n", insideCircle(border, pt, r));
                                                       // 1-at border
 point_i outside(10, 2);
 printf("%d\n", insideCircle(outside, pt, r));
                                                         // 2-outside
  double d = 2 * r;
 printf("Diameter = %.21f\n", d);
  double c = PT * d:
 printf("Circumference (Perimeter) = %.21f\n", c);
 double A = PI * r * r;
 printf("Area of circle = %.21f\n", A);
 printf("Length of arc (central angle = 60 degrees) = %.21f\n",
       60.0 / 360.0 * c);
  printf("Length of chord (central angle = 60 degrees) = %.21f\n",
       sgrt((2 * r * r) * (1 - cos(DEG_to_RAD(60.0))));
  printf("Area of sector (central angle = 60 degrees) = %.21f\n",
       60.0 / 360.0 * A);
 point p1:
  point p2(0.0, -1.0);
  point ans:
  circle2PtsRad(p1, p2, 2.0, ans);
  printf("One of the center is (%.21f, %.21f)\n", ans.x, ans.y);
 circle2PtsRad(p2, p1, 2.0, ans); // we simply reverse p1 with p2 printf("The other center is (%.21f, %.21f) \n", ans.x, ans.v);
 return 0:
```

2.3 Triangles

```
#include <cstdio>
#include <cmath>
using namespace std:
#define EPS 1e-9
#define PI acos (-1.0)
double DEG_to_RAD(double d) { return d * PI / 180.0; }
double RAD_to_DEG(double r) { return r * 180.0 / PI; }
// default constructor
  point i() \{ x = y = 0; \}
  point_i(int _x, int _y) : x(_x), y(_y) {} };
                                                     // constructor
struct point { double x, y;  // only used if more precision is needed
  point() { x = y = 0.0; }
                                              // default constructor
  point (double _x, double _y) : x(_x), y(_y) {} }; // constructor
double dist(point p1, point p2) {
  return hypot(p1.x - p2.x, p1.y - p2.y); }
double perimeter(double ab, double bc, double ca) {
  return ab + bc + ca; }
double perimeter(point a, point b, point c) {
  return dist(a, b) + dist(b, c) + dist(c, a); }
double area(double ab, double bc, double ca) {
  // Heron's formula, split sqrt(a * b) into sqrt(a) * sqrt(b); in
       implementation
  double s = 0.5 * perimeter(ab, bc, ca);
  return sqrt(s) * sqrt(s - ab) * sqrt(s - bc) * sqrt(s - ca); }
double area(point a, point b, point c) {
  return area(dist(a, b), dist(b, c), dist(c, a)); }
// from ch7_01_points_lines
struct line { double a, b, c; }; // a way to represent a line
// the answer is stored in the third parameter (pass by reference)
void pointsToLine(point p1, point p2, line &1) {
  if (fabs(p1.x - p2.x) < EPS) {</pre>
                                          // vertical line is fine
    1.a = 1.0; 1.b = 0.0; 1.c = -p1.x;
                                                  // default values
   1.a = -(double) (p1.y - p2.y) / (p1.x - p2.x);
   1.b = 1.0; // IMPORTANT: we fix the value of b to 1.0
   1.c = -(double) (1.a * p1.x) - p1.y;
bool areParallel(line 11, line 12) {
                                         // check coefficient a + b
  return (fabs(11.a-12.a) < EPS) && (fabs(11.b-12.b) < EPS); }
// returns true (+ intersection point) if two lines are intersect
bool areIntersect(line 11, line 12, point &p) {
  if (areParallel(11, 12)) return false;
                                                  // no intersection
  // solve system of 2 linear algebraic equations with 2 unknowns
  p.x = (12.b * 11.c - 11.b * 12.c) / (12.a * 11.b - 11.a * 12.b);
  // special case: test for vertical line to avoid division by zero
  if (fabs(11.b) > EPS) p.y = -(11.a * p.x + 11.c);
                     p.y = -(12.a * p.x + 12.c);
  else
  return true: }
struct vec { double x, y; // name: 'vec' is different from STL vector
  vec(double _x, double _y) : x(_x), y(_y) {} };
vec toVec(point a, point b) {
                                  // convert 2 points to vector a->b
  return vec(b.x - a.x, b.y - a.y); }
vec scale (vec v, double s) {
                                 // nonnegative s = \{<1 ... 1 ... > 1\}
  return vec(v.x * s, v.v * s); }
                                              // shorter.same.longer
point translate(point p, vec v) {
                                        // translate p according to v
  return point(p.x + v.x , p.y + v.y); }
double rInCircle(double ab, double bc, double ca) {
  return area(ab, bc, ca) / (0.5 * perimeter(ab, bc, ca)); }
double rInCircle(point a, point b, point c) {
  return rInCircle(dist(a, b), dist(b, c), dist(c, a)); }
```

```
// assumption: the required points/lines functions have been written
// returns 1 if there is an inCircle center, returns 0 otherwise
// if this function returns 1, ctr will be the inCircle center
// and r is the same as rInCircle
int inCircle(point p1, point p2, point p3, point &ctr, double &r) {
  r = rInCircle(p1, p2, p3);
  if (fabs(r) < EPS) return 0;</pre>
                                                                            // no inCircle center
  line 11, 12;
                                                     // compute these two angle bisectors
  double ratio = dist(p1, p2) / dist(p1, p3);
  point p = translate(p2, scale(toVec(p2, p3), ratio / (1 + ratio)));
   pointsToLine(p1, p, 11);
   ratio = dist(p2, p1) / dist(p2, p3);
  p = translate(p1, scale(toVec(p1, p3), ratio / (1 + ratio)));
   pointsToLine(p2, p, 12);
  areIntersect(11, 12, ctr);
                                                        // get their intersection point
  return 1: }
double rCircumCircle(double ab, double bc, double ca) {
  return ab * bc * ca / (4.0 * area(ab, bc, ca)); }
double rCircumCircle(point a, point b, point c) {
  return rCircumCircle(dist(a, b), dist(b, c), dist(c, a)); }
// assumption: the required points/lines functions have been written
// returns 1 if there is a circumCenter center, returns 0 otherwise
// if this function returns 1, ctr will be the circumCircle center
// and r is the same as rCircumCircle
int circumCircle(point p1, point p2, point p3, point &ctr, double &r) {
  double a = p2.x - p1.x, b = p2.y - p1.y;
   double c = p3.x - p1.x, d = p3.y - p1.y;
   double e = a * (p1.x + p2.x) + b * (p1.y + p2.y);
   double f = c * (p1.x + p3.x) + d * (p1.y + p3.y);
   double g = 2.0 * (a * (p3.y - p2.y) - b * (p3.x - p2.x));
   if (fabs(g) < EPS) return 0;</pre>
   ctr.x = (d*e - b*f) / q;
  ctr.v = (a*f - c*e) / q;
   r = dist(p1, ctr); // r = distance from center to 1 of the 3 points
   return 1; }
// returns true if point d is inside the circumCircle defined by a,b,c
int inCircumCircle(point a, point b, point c, point d) {
  return (a.x - d.x) * (b.y - d.y) * ((c.x - d.x) * (c.x - d.x) + (c.y)
              - d.y) * (c.y - d.y)) +
              (a.y - d.y) * ((b.x - d.x) * (b.x - d.x) + (b.y - d.y) * (b.y
                      - d.v) + (c.x - d.x) +
              ((a.x - d.x) * (a.x - d.x) + (a.y - d.y) * (a.y - d.y)) * (b.
                     x - d.x) * (c.y - d.y) -
              ((a.x - d.x) * (a.x - d.x) + (a.y - d.y) * (a.y - d.y)) * (b.
                     y - d.y) * (c.x - d.x) -
              (a.y - d.y) * (b.x - d.x) * ((c.x - d.x) * (c.x - d.x) + (c.y)
              (a.x - d.y) * (c.y - d.y) - (a.x - d.x) * ((b.x - d.x)) * ((b.y - d.y)) * (b.y - d.y) * (b.y - d.y
                        -d.y)) * (c.y - d.y) > 0 ? 1 : 0;
bool canFormTriangle(double a, double b, double c) {
  return (a + b > c) && (a + c > b) && (b + c > a): }
int main() {
  double base = 4.0, h = 3.0;
  double A = 0.5 + base + b:
  printf("Area = %.21f\n", A);
   point a:
                                                                                // a right triangle
  point b(4.0, 0.0);
   point c(4.0, 3.0);
   double p = perimeter(a, b, c);
  double s = 0.5 * p;
  A = area(a, b, c);
  printf("Area = %.21f\n", A);
                                                                 // must be the same as above
   double r = rInCircle(a, b, c);
  printf("R1 (radius of incircle) = %.21f\n", r);
                                                                                                  // 1.00
   point ctr;
   int res = inCircle(a, b, c, ctr, r);
  // same, 1.00
   printf("R2 (radius of circumcircle) = %.21f\n", rCircumCircle(a, b,
          c)); // 2.50
   res = circumCircle(a, b, c, ctr, r);
  printf("R2 (radius of circumcircle) = %.21f\n", r); // same, 2.50
  printf("Center = (%.21f, %.21f)\n", ctr.x, ctr.y); // (2.00, 1.50)
```

```
// inside triangle and circumCircle
point d(2.0, 1.0);
printf("d inside circumCircle (a, b, c) ? %d\n", inCircumCircle(a, b
      , c, d));
point e(2.0, 3.9); // outside the triangle but inside circumCircle
printf("e inside circumCircle (a, b, c) ? %d\n", inCircumCircle(a, b
point f(2.0, -1.1);
                                                   // slightly outside
printf("f inside circumCircle (a, b, c) ? %d\n", inCircumCircle(a, b
     , c, f));
// Law of Cosines
double ab = dist(a, b);
double bc = dist(b, c);
double ca = dist(c, a);
double alpha = RAD_to_DEG(acos((ca * ca + ab * ab - bc * bc) / (2.0
      * ca * ab)));
printf("alpha = %.21f\n", alpha);
double beta = RAD_to_DEG(acos((ab * ab + bc * bc - ca * ca) / (2.0
      \star ab \star bc)));
printf("beta = %.21f\n", beta);
double gamma = RAD_to_DEG(acos((bc * bc + ca * ca - ab * ab) / (2.0
      * bc * ca)));
printf("gamma = %.21f\n", gamma);
// Law of Sines
printf("%.21f == %.21f == %.21f\n", bc / sin(DEG_to_RAD(alpha)), ca
      / sin(DEG_to_RAD(beta)), ab / sin(DEG_to_RAD(gamma)));
// Phytagorean Theorem
printf("%.21f^2 == %.21f^2 + %.21f^2\n", ca, ab, bc);
// Triangle Inequality
printf("(%d, %d, %d) => can form triangle? %d\n", 3, 4, 5,
      canFormTriangle(3, 4, 5)); // yes
printf("(%d, %d, %d) => can form triangle? %d\n", 3, 4, 7,
canFormTriangle(3, 4, 7)); // no, actually straight line
printf("(%d, %d, %d) => can form triangle? %d\n", 3, 4, 8,
      canFormTriangle(3, 4, 8)); // no
return 0;
```

2.4 Polygon

```
#include <algorithm>
#include <cstdio>
#include <cmath>
#include <stack>
#include <vector>
using namespace std;
#define EPS 1e-9
#define PT acos (-1.0)
double DEG to RAD(double d) { return d * PI / 180.0: }
double RAD to DEG(double r) { return r * 180.0 / PI; }
struct point { double x, y; // only used if more precision is needed
 point() { x = y = 0.0; }
                                                // default constructor
  point (double _x, double _y) : x(_x), y(_y) {}
                                                      // user-defined
  bool operator == (point other) const {
  return (fabs(x - other.x) < EPS && (fabs(y - other.y) < EPS)); };</pre>
struct vec { double x, y; // name: 'vec' is different from STL vector
 vec(double _x, double _y) : x(_x), y(_y) {} };
vec toVec(point a, point b) {
                                // convert 2 points to vector a->b
 return vec(b.x - a.x, b.y - a.y); }
double dist(point p1, point p2) {
                                                 // Euclidean distance
 return hypot(p1.x - p2.x, p1.y - p2.y); }
                                                     // return double
// returns the perimeter, which is the sum of Euclidian distances
// of consecutive line segments (polygon edges)
double perimeter(const vector<point> &P) {
 double result = 0.0;
  for (int i = 0; i < (int)P.size()-1; i++) // remember that P[0] = P</pre>
        [n-1]
   result += dist(P[i], P[i+1]);
  return result; }
// returns the area, which is half the determinant
double area(const vector<point> &P) {
```

```
double result = 0.0, x1, y1, x2, y2;
  for (int i = 0; i < (int)P.size()-1; i++) {</pre>
    x1 = P[i].x; x2 = P[i+1].x;
    y1 = P[i].y; y2 = P[i+1].y;
    result += (x1 * y2 - x2 * y1);
  return fabs(result) / 2.0; }
double dot(vec a, vec b) { return (a.x * b.x + a.y * b.y); }
double norm sq(vec v) { return v.x * v.x + v.y * v.y; }
double angle (point a, point o, point b) { // returns angle aob in rad
  vec oa = toVec(o, a), ob = toVec(o, b);
  return acos(dot(oa, ob) / sqrt(norm_sq(oa) * norm_sq(ob))); }
double cross(vec a, vec b) { return a.x * b.y - a.y * b.x; }
// note: to accept collinear points, we have to change the '> 0'
// returns true if point r is on the left side of line pg
bool ccw(point p, point q, point r) {
  return cross(toVec(p, q), toVec(p, r)) > 0; }
// returns true if point r is on the same line as the line pg
bool collinear (point p, point q, point r) {
  return fabs(cross(toVec(p, q), toVec(p, r))) < EPS; }
// returns true if we always make the same turn while examining
// all the edges of the polygon one by one
bool isConvex(const vector<point> &P) {
  int sz = (int)P.size();
  if (sz <= 3) return false; // a point/sz=2 or a line/sz=3 is not</pre>
       convex
  bool isLeft = ccw(P[0], P[1], P[2]);
                                                     // remember one
       result
  for (int i = 1; i < sz-1; i++)
                                            // then compare with the
       others
    if (ccw(P[i], P[i+1], P[(i+2) == sz ? 1 : i+2]) != isLeft)
     return false;
                              // different sign -> this polygon is
           concave
                                                  // this polygon is
  return true; }
        convex
// returns true if point p is in either convex/concave polygon P
bool inPolygon(point pt, const vector<point> &P) {
  if ((int)P.size() == 0) return false;
  double sum = 0; // assume the first vertex is equal to the last
        vertex
  for (int i = 0; i < (int)P.size()-1; i++) {</pre>
   if (ccw(pt, P[i], P[i+1]))
                                                           // left
        sum += angle(P[i], pt, P[i+1]);
             turn/ccw
    else sum -= angle(P[i], pt, P[i+1]); }
                                                           // right
          turn/cw
  return fabs(fabs(sum) - 2*PI) < EPS; }
// line segment p-q intersect with line A-B.
point lineIntersectSeg(point p, point q, point A, point B) {
 double a = B.y - A.y;
  double b = A.x - B.x:
  double c = B.x * A.y - A.x * B.y;
  double u = fabs(a \star p.x + b \star p.y + c);
  double v = fabs(a * q.x + b * q.y + c);
  return point((p.x * v + q.x * u) / (u+v), (p.y * v + q.y * u) / (u+v
// cuts polygon Q along the line formed by point a -> point b
// (note: the last point must be the same as the first point)
vector<point> cutPolygon(point a, point b, const vector<point> &Q) {
  vector<point> P:
  for (int i = 0; i < (int)Q.size(); i++) {</pre>
    double left1 = cross(toVec(a, b), toVec(a, Q[i])), left2 = 0;
    if (i != (int)Q.size()-1) left2 = cross(toVec(a, b), toVec(a, Q[i
          +1]));
    if (left1 > -EPS) P.push_back(Q[i]);
                                               // O[il is on the left
          of ab
                                     // edge (Q[i], Q[i+1]) crosses
    if (left1 * left2 < -EPS)</pre>
          line ab
     P.push_back(lineIntersectSeg(Q[i], Q[i+1], a, b));
  if (!P.empty() && !(P.back() == P.front()))
    P.push_back(P.front()); // make P's first point = P's last
          point
  return P; }
point pivot;
bool angleCmp(point a, point b) {
                                                  // angle-sorting
      function
  if (collinear(pivot, a, b))
                                                            // special
```

```
closer
 double d1x = a.x - pivot.x, d1y = a.y - pivot.y;
  double d2x = b.x - pivot.x, d2y = b.y - pivot.y;
  return (atan2(d1y, d1x) - atan2(d2y, d2x)) < 0; } // compare two
vector<point> CH(vector<point> P) { // the content of P may be
      reshuffled
  int i, j, n = (int)P.size();
 if (n <= 3) {
   if (!(P[0] == P[n-1])) P.push_back(P[0]); // safeguard from corner
    return P;
                                        // special case, the CH is P
  // first, find PO = point with lowest Y and if tie: rightmost X
  int P0 = 0;
  for (i = 1; i < n; i++)
   if (P[i].y < P[P0].y || (P[i].y == P[P0].y && P[i].x > P[P0].x))
  point temp = P[0]; P[0] = P[P0]; P[P0] = temp; // swap P[P0] with
  // second, sort points by angle w.r.t. pivot PO
  pivot = P[0];
                                  // use this global variable as
       reference
  sort (++P.begin(), P.end(), angleCmp);
                                                      // we do not sort
        P [ 0 ]
  // third, the ccw tests
  vector<point> S;
  S.push_back(P[n-1]); S.push_back(P[0]); S.push_back(P[1]); //
       initial S
  i = 2;
                                                 // then, we check the
        rest
  while (i < n) {
                            // note: N must be >= 3 for this method to
        work
    j = (int) S.size()-1;
   if (ccw(S[j-1], S[j], P[i])) S.push_back(P[i++]); // left turn,
         accept
    else S.pop_back(); } // or pop the top of S until we have a left
          turn
  return S: }
                                                       // return the
       result
int main() {
 // 6 points, entered in counter clockwise order, 0-based indexing
  vector<point> P:
  P.push back(point(1, 1));
  P.push back(point(3, 3));
  P.push back(point(9, 1));
  P.push back(point(12, 4));
 P.push back(point(9, 7));
  P.push back(point(1, 7));
 P.push_back(P[0]); // loop back
  printf("Perimeter of polygon = %.21f\n", perimeter(P)); // 31.64
 printf("Area of polygon = %.21f\n", area(P)); // 49.00
printf("Is convex = %d\n", isConvex(P)); // false (P1 is the culprit
  //// the positions of P6 and P7 w.r.t the polygon
  //7 P5-----P4
  1/6 1
  //5
  //4 I P7
                            P3
  //3 1
       P1
  //2 | / P6
  //1 PO
  //0 1 2 3 4 5 6 7 8 9 101112
  point P6(3, 2); // outside this (concave) polygon
 printf("Point P6 is inside this polygon = %d\n", inPolygon(P6, P));
       // false
  point P7(3, 4); // inside this (concave) polygon
  printf("Point P7 is inside this polygon = %d\n", inPolygon(P7, P));
       // true
  // cutting the original polygon based on line P[2] \rightarrow P[4] (get the
       left side)
  //7 P5-----
  1/6 1
  //5 /
  1/4 1
```

return dist(pivot, a) < dist(pivot, b); // check which one is

```
//0 1 2 3 4 5 6 7 8 9 101112
// new polygon (notice the index are different now):
//7 P4----
1/6 1
1/5 1
//4
//2 / /
P = cutPolygon(P[2], P[4], P);
printf("Perimeter of polygon = %.21f\n", perimeter(P)); // smaller
printf("Area of polygon = %.21f\n", area(P)); // 40.00
// running convex hull of the resulting polygon (index changes again
//7 P3----P2
//5
1/4 1
//3
//0 1 2 3 4 5 6 7 8 9
P = CH(P); // now this is a rectangle
printf("Perimeter of polygon = %.21f\n", perimeter(P)); // precisely
       28.00
printf("Area of polygon = %.21f\n", area(P)); // precisely 48.00
printf("Is convex = %d\n", isConvex(P)); // true
printf("Point P6 is inside this polygon = %d\n", inPolygon(P6, P));
      // true
printf("Point P7 is inside this polygon = %d\n", inPolygon(P7, P));
      // true
return 0;
```

2.5 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone
     chain
// algorithm. Eliminate redundant points from the hull if
      REMOVE REDUNDANT is
  #defined.
// Running time: O(n log n)
    INPUT: a vector of input points, unordered.
    OUTPUT: a vector of points in the convex hull, counterclockwise,
      starting
              with bottommost/leftmost point
#define REMOVE REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
 T x, y;
 PT(T x, T y) : x(x), y(y) {}
 bool operator<(const PT &rhs) const { return make_pair(v,x) <</pre>
       make_pair(rhs.y,rhs.x); }
 bool operator == (const PT &rhs) const { return make pair(v,x) ==
        make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a
     ); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
 return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 && (a.y
        -b.y) * (c.y-b.y) <= 0);
#endif
void ConvexHull (vector<PT> &pts) {
 sort(pts.begin(), pts.end());
 pts.erase(unique(pts.begin(), pts.end()), pts.end());
```

```
vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i])
          >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i])
          <= 0) dn.pop_back();
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push back(up[i]);
#ifdef REMOVE REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {
    if (between (dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
#endif
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP
int main() {
  int t:
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {
   int n;
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
    vector<PT> h(v);
    map<PT, int> index;
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h);
    double len = 0:
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {
      if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

2.6 3D Convex Hull

```
const double eps = le-8;
int mark[1005][1005];
Point info[1005];
int n, cnt;
double mix(const Point &a, const Point &b, const Point &c) {
    return a.dot(b.cross(c));}
double area(int a, int b, int c) {
    return ((info[b] - info[a]).cross(info[c] - info[a])).length();}
double volume(int a, int b, int c, int d) {
    return mix(info[b] - info[a], info[c] - info[a], info[d] - info[a
    ]);}
struct Face {
    int a, b, c;
    Face() {
    Face(int a, int b, int c): a(a), b(b), c(c) {
        int &operator [](int k) { return k==0?a:k==1?b:c; }
    };
```

```
vector <Face> face;
inline void insert(int a, int b, int c) { face.push_back(Face(a, b, c)
void add(int v) {
    vector <Face> tmp;
    int a, b, c;
    for (int i = 0; i < SIZE(face); i++) {</pre>
        a = face[i][0]; b = face[i][1]; c = face[i][2];
        if (Sign(volume(v, a, b, c)) < 0)</pre>
            mark[a][b] = mark[b][a] = mark[b][c] = mark[c][b] = mark[c
                  ][a] = mark[a][c] = cnt;
        else tmp.push_back(face[i]);
    for (int i = 0; i < SIZE(tmp); i++) {</pre>
        a = face[i][0]; b = face[i][1]; c = face[i][2];
        if (mark[a][b] == cnt) insert(b, a, v);
        if (mark[b][c] == cnt) insert(c, b, v);
        if (mark[c][a] == cnt) insert(a, c, v);
    for (int i = 2; i < n; i++) {
        Point ndir = (info[0] - info[i]).cross(info[1] - info[i]);
        if (ndir == Point()) continue;
        swap(info[i], info[2]);
        for (int j = i + 1; j < n; j++)
   if (Sign(volume(0, 1, 2, j)) != 0) {</pre>
                swap(info[j], info[3]);
                insert(0, 1, 2); insert(0, 2, 1);
                return 1;
   return 0:
int main() {
    for (; scanf("%d", &n) == 1; ) {
       for (int i = 0; i < n; i++)
            info[i].Input();
        sort(info, info + n);
        n = unique(info, info + n) - info;
        face.clear();
        random shuffle(info, info + n);
        if (Find()) {
            memset(mark, 0, sizeof(mark));
            cnt = 0:
            for (int i = 3; i < n; i++) add(i);</pre>
            vector<Point> Ndir;
            for (int i = 0; i < SIZE(face); ++i) {</pre>
                Point p = (info[face[i][0]] - info[face[i][1]]).cross(
                      info[face[i][2]] - info[face[i][1]]);
                p = p / p.length();
                Ndir.push_back(p);
            sort (Ndir.begin(), Ndir.end());
            int ans = unique(Ndir.begin(), Ndir.end()) - Ndir.begin();
            printf("%d\n", ans);
        else (
            printf("1\n");
```

2.7 Slow Delaunay triangulation

```
// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
//
// Running time: O(n^4)
//
// INPUT: x[] = x-coordinates
//
// OUTPUT: triples = a vector containing m triples of indices
corresponding to triangle vertices
typedef double T;

struct triple {
   int i, j, k;
   triple() {
      triple(int i, int j, int k) : i(i), j(j), k(k) {}
   }
};
```

```
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
        int n = x.size();
        vector<T> z(n);
        vector<triple> ret;
        for (int i = 0; i < n; i++)</pre>
            z[i] = x[i] * x[i] + y[i] * y[i];
        for (int i = 0; i < n-2; i++) {
            for (int j = i+1; j < n; j++) {
   for (int k = i+1; k < n; k++) {</pre>
                     if (j == k) continue;
                      double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])
                            *(z[j]-z[i]);
                      double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])
                            *(z[k]-z[i]);
                      double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])
                            *(y[j]-y[i]);
                     bool flag = zn < 0;
                     for (int m = 0; flag && m < n; m++)</pre>
                         flag = flag && ((x[m]-x[i])*xn +
                                           (y[m]-y[i])*yn +
                                           (z[m]-z[i])*zn <= 0);
                     if (flag) ret.push_back(triple(i, j, k));
        return ret:
int main()
   T xs[]={0, 0, 1, 0.9};
T ys[]={0, 1, 0, 0.9};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);
    //expected: 0 1 3
    int i;
    for(i = 0; i < tri.size(); i++)</pre>
       printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
    return 0;
```

2.8 Closest Pair

```
// Source: e-maxx.ru
#define upd_ans(x, y) {}
double mindist = 1e20; // will be the result
void rec(int 1, int r, Point a[]) {
   if (r - 1 <= 3) {
        for (int i=1; i<=r; ++i)
           for (int j=i+1; j<=r; ++j)
                   upd_ans(a[i], a[j]);
        sort (a+1, a+r+1); // compare by y
   int m = (1 + r) >> 1;
   int midx = a[m].x;
   rec(1, m, a), rec(m+1, r, a);
   static Point t[MAXN];
   merge(a+1, a+m+1, a+m+1, a+r+1, t); // compare by y
   copy(t, t+r-1+1, a+1);
   int tsz = 0;
   for (int i=1; i<=r; ++i)</pre>
       if (fabs(a[i].x - midx) < mindist) {</pre>
            for (int j=tsz-1; j>=0 && a[i].y - t[j].y < mindist; --j)</pre>
               upd_ans(a[i], t[j]);
            t[tsz++] = a[i];
```

2.9 Rotating Caliphers

```
// Rotating calipers
```

```
double convex_diameter(Polygon pt) {
    const int n = pt.size();
    int is = 0, js = 0;
for (int i = 1; i < n; ++i) {</pre>
        if (pt[i].y > pt[is].y) is = i;
        if (pt[i].y < pt[js].y) js = i;
    double maxd = (pt[is]-pt[js]).norm();
    int i, maxi, j, maxj;
    i = maxi = is;
    j = maxj = js;
         int jj = j+1; if (jj == n) jj = 0;
        if ((pt[i] - pt[jj]).norm() > (pt[i] - pt[j]).norm()) j = (j
         else i = (i+1) % n;
        if ((pt[i]-pt[j]).norm() > maxd) {
            maxd = (pt[i]-pt[j]).norm();
            \max i = i; \max j = j;
    } while (i != is || j != js);
    return maxd; /* farthest pair is (maxi, maxj). */
```

3 Numerical algorithms

3.1 Fast exponentiation

double n = 2.37;

```
Uses powers of two to exponentiate numbers and matrices. Calculates
n^k in O(\log(k)) time when n is a number. If A is an n x n matrix,
calculates A^k in O(n^3*log(k)) time.
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T power(T x, int k) {
  T ret = 1;
    if(k & 1) ret *= x;
    k >>= 1; x *= x;
  return ret;
VVT multiply(VVT& A, VVT& B) {
  int n = A.size(), m = A[0].size(), k = B[0].size();
  VVT C(n, VT(k, 0));
  for(int i = 0; i < n; i++)</pre>
    for(int j = 0; j < k; j++)
      for(int 1 = 0; 1 < m; 1++)
        C[i][j] += A[i][1] * B[1][j];
  return C;
VVT power(VVT& A, int k) {
  int n = A.size();
  VVT ret(n, VT(n)), B = A;
  for(int i = 0; i < n; i++) ret[i][i]=1;</pre>
    if(k & 1) ret = multiply(ret, B);
    k >>= 1; B = multiply(B, B);
  return ret:
int main()
  /* Expected Output:
     2.37^48 = 9.72569e+17
     376 264 285 220 265
550 376 529 285 484
     484 265 376 264 285
     285 220 265 156 264
     529 285 484 265 376 */
```

3.2 Prime numbers

```
typedef unsigned long long 11;
typedef vector<11> v11;
typedef vector<int> vi;
ll _sieve_size;
bitset<10000010> bs;
vll primes;
void sieve(ll upper) {
    _sieve_size = upper + 1;
     bs.set(); // set all to one
     bs[0] = bs[1] = 0;
    for(ll i = 2; i < _sieve_size; i++) if (bs[i]) {
    for(ll j = i*i; j < _sieve_size; j+= i) {</pre>
                  bs[j] = 0;
              primes.push_back((int) i);
bool isPrime(ll n) {
    if (n <= _sieve_size) return bs[n];</pre>
    for(int i = 0; i < (int) primes.size(); i++) {
   if (n % primes[i] == 0) return false;</pre>
         if (primes[i] * primes[i] > n) return true;
    return true:
bool isPrime_slow(ll n) {
    if(n < 2) return false;</pre>
    if(n == 2 | | n == 3) return true;
if(n % 2 == 0 | | n % 3 == 0) return false;
    int limit = sqrt(n);
    for(int i = 5; i <= limit; i += 6) {
   if(n % i == 0 || n % (i+2) == 0)</pre>
              return false:
    return true:
vi primeFactors(ll N) {
     vi factors:
     11 PF_index = 0; 11 PF = primes[PF_index];
    while (PF*PF <= N) {
         while (N%PF == 0) {
              N /= PF; factors.push_back(PF);
         PF = primes[++PF_index];
    if(N != 1) factors.push_back(N);
    return factors;
// Primes less than 1000:
```

```
5 7 11 13 17 19
                                          23
        43
     41
              47 53 59
                            61
                                 67
                                      71
                                          73
     97 101 103 107 109
                          113 127 131 137
        163 167 173 179 181 191
                                    193 197
    227 229 233 239 241
                           251
                                257
                                    263 269
        293 307 311 313 317 331
                                    337
                                         347
        373
            379 383 389
                                         419
      433
    439 443 449 457
                                          487
      503
        521
                                         571
     593
    599
        601 607 613 617
                           619
                                631
                                    641
                                         643
      659
    661 673 677 683
                      691
                           701
                                709
                                    719
                                         727
     743
    751
        757
            761 769 773
                           787
                                797
                                    809
                                         811
      827
    829
       839
             853 857
                      859
                           863
                                877
                                    881
                                          883
     911
        929
            937 941 947
                           953 967 971 977 983 991
// Other primes: largest prime smaller than X is Y
   10 is 7.
    100 is 97.
    1000 is 997.
    10000 is 9973.
    100000 is 99991.
    10000000 is 999983.
    10000000 is 9999991.
    100000000 is 99999989.
    1000000000 is 999999937.
    10000000000 is 9999999967.
    100000000000 is 99999999977.
    10000000000000 is 999999999999999.
    100000000000000 is 9999999999971
    1000000000000000 is 9999999999973.
    100000000000000000 is 999999999999937.
    1000000000000000000 is 9999999999999997.
```

29

79

647

733

887

821 823

139 149

271 277

3.3 Miller-Rabin Primality Test (C)

```
// Randomized Primality Test (Miller-Rabin):
    Error rate: 2^(-TRIAL)
    Almost constant time. srand is needed
#include <stdlib.h>
#define EPS 1e-7
typedef long long LL;
LL ModularMultiplication(LL a, LL b, LL m)
        LL ret=0. c=a:
        while (b)
                if (b&1) ret=(ret+c) %m;
                b>>=1: c=(c+c)%m:
        return ret:
LL Modular Exponentiation (LL a, LL n, LL m)
        LL ret=1, c=a;
        while (n)
                if(n&1) ret=ModularMultiplication(ret, c, m);
                n>>=1; c=ModularMultiplication(c, c, m);
        return ret:
bool Witness(LL a, LL n)
        T_{n}T_{n-1} = n-1:
  int t=0;
```

```
while (! (u&1)) {u>>=1; t++;}
        LL x0=ModularExponentiation(a, u, n), x1;
        for(int i=1;i<=t;i++)</pre>
                x1=ModularMultiplication(x0, x0, n);
                if (x1==1 && x0!=1 && x0!=n-1) return true;
        if(x0!=1) return true;
        return false;
LL Random(LL n)
  LL ret=rand(); ret*=32768;
        ret+=rand(); ret*=32768;
        ret+=rand(); ret*=32768;
        ret+=rand();
  return ret%n;
bool IsPrimeFast (LL n, int TRIAL)
  while (TRIAL--)
    LL a=Random(n-2)+1;
    if(Witness(a, n)) return false;
  return true:
```

3.4 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1:
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
               b >>= 1:
        return ret:
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
       int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
               int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
```

```
int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                x = mod(x*(b / g), n);
                for (int i = 0; i < g; i++)</pre>
                        ret.push_back(mod(x + i*(n / g), n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
        if (q > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = 1 cm (m1, m2)
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
        int s, t;
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g)
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make_pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {</pre>
                ret = chinese remainder theorem(ret.second, ret.first,
                      m[i], r[i]);
                if (ret.second == -1) break;
        return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
       if (!a && !b)
                if (c) return false;
                x = 0; y = 0;
                return true:
        if (!a)
                if (c % b) return false:
                x = 0; v = c / b;
                return true;
        if (!h)
                if (c % a) return false;
                x = c / a; y = 0;
                return true:
        int g = gcd(a, b);
        if (c % g) return false;
        x = c / g * mod_inverse(a / g, b / g);
        v = (c - a*x) / b;
        return true:
int main() {
        // expected: 2
        cout << gcd(14, 30) << endl;
        // expected: 2 -2 1
        int x, y;
       int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
        // expected: 95 451
        VI sols = modular_linear_equation_solver(14, 30, 100);
        for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
        // expected: 8
        cout << mod inverse(8, 9) << endl;
```

```
// expected: 23 105
// 11 12
PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
cout << ret.first << " " << ret.second << endl;
ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << endl;
// expected: 5 -15
if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;
return 0;</pre>
```

3.5 Systems of linear equations, matrix inverse, determinant

// Gauss-Jordan elimination with full pivoting.

```
// Uses:
     (1) solving systems of linear equations (AX=B)
      (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
// INPUT:
                a[][] = an nxn matrix
                b[][] = an nxm matrix
// OUTPUT:
                         = an nxm matrix (stored in b[][])
                A^{-1} = an nxn matrix (stored in a[][])
                returns determinant of a[][]
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT &a, VVT &b) {
 const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
 T \det = 1;
  for (int i = 0; i < n; i++) {
    int pj = -1, pk = -1;
int pj = -1, pk = -1;
for (int j = 0; j < n; j++) if (!ipiv[j])
    for (int k = 0; k < n; k++) if (!ipiv[k])
    if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk
     \begin{tabular}{ll} \textbf{if} & (fabs(a[pj][pk]) < EPS) & (cerr << "Matrix is singular." << endl \\ \end{tabular} 
           ; exit(0); }
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c; for (int p = 0; p < m; p++) b[pk][p] *= c; for (int p = 0; p < n; p++) if (p != pk) {
      c = a[p][pk];
       for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
       for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
  return det;
int main() {
```

```
const int n = 4;
const int m = 2;
double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
for (int i = 0; i < n; i++) {
 a[i] = VT(A[i], A[i] + n);
  b[i] = VT(B[i], B[i] + m);
double det = GaussJordan(a, b);
cout << "Determinant: " << det << endl;
// expected: -0.233333 0.166667 0.133333 0.0666667
             0.166667 0.166667 0.333333 -0.333333
             0.233333 0.833333 -0.133333 -0.0666667
             0.05 -0.75 -0.1 0.2
cout << "Inverse: " << endl;
for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)
   cout << a[i][j] << ' ';
// expected: 1.63333 1.3
             -0.166667 0.5
             -1.85 -1.35
cout << "Solution: " << endl;
for (int i = 0; i < n; i++) {
 for (int j = 0; j < m; j++)
   cout << b[i][j] << ' ';
  cout << endl;
```

3.6 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT: a[][] = an nxm matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
             returns rank of a[][]
const. double EPSILON = 1e-10:
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
  int n = a.size();
  int m = a[0].size();
  int r = 0:
  for (int c = 0; c < m && r < n; c++) {
    int i = r:
    for (int i = r + 1; i < n; i++)
      if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
for (int i = 0; i < n; i++) if (i != r) {</pre>
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];</pre>
    r++;
  return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
```

```
{ 5, 11, 10, 8},
  { 9, 7, 6, 12},
  { 4, 14, 15, 1},
  {13, 21, 21, 13}};
VVT a(n);
for (int i = 0; i < n; i++)</pre>
 a[i] = VT(A[i], A[i] + m);
int rank = rref(a);
// expected: 3
cout << "Rank: " << rank << endl;
// expected: 1 0 0 1
              0 0 0 3.10862e-15
             0 0 0 2.22045e-15
cout << "rref: " << endl;
for (int i = 0; i < 5; i++)
 for (int j = 0; j < 4; j++)
  cout << a[i][j] << ' ';</pre>
  cout << endl;
```

3.7 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear programs of the form
         subject to Ax <= b
                          x >= 0
// INPUT: A -- an m x n matrix
            b -- an m-dimensional vector
             c -- an n-dimensional vector
             x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
             above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
  VI B, N;
   \begin{array}{l} LPSolver(\textbf{const} \ VVD \ \&A, \ \textbf{const} \ VD \ \&b, \ \textbf{const} \ VD \ \&c): \\ m(b.size()), \ n(c.size()), \ N(n+1), \ B(m), \ D(m+2, \ VD(n+2)) \ \{ \\ \textbf{for (int} \ i = 0; \ i < m; \ i++) \ \textbf{for (int} \ j = 0; \ j < n; \ j++) \ D[i][j] = \\ \end{array} 
     for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + 1]
             11 = b[i]; }
    for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
N[n] = -1; D[m + 1][n] = 1;</pre>
  void Pivot(int r, int s) {
    double inv = 1.0 / D[r][s];
for (int i = 0; i < m + 2; i++) if (i != r)</pre>
       for (int j = 0; j < n + 2; j++) if (j != s)
    D[i][j] -= D[r][j] * D[i][s] * inv;
for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;
    D[r][s] = inv;
     swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
     while (true) {
       int s = -1;
        for (int j = 0; j <= n; j++) {
```

```
if (phase == 2 && N[j] == -1) continue;
        if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j]
               < N[s]) s = j;
      if (D[x][s] > -EPS) return true;
      for (int i = 0; i < m; i++) {
        if (D[i][s] < EPS) continue;</pre>
        if (r == -1 \mid | D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s]
           (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s]) && B[i] <
      if (r == -1) return false;
      Pivot(r, s);
  DOUBLE Solve(VD &x) {
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -</pre>
            numeric_limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1;
        for (int j = 0; j <= n; j++)
          if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s] && N[
                j] < N[s]) s = j;
        Pivot(i, s);
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];</pre>
    return D[m][n + 1];
};
int main() {
  const int m = 4;
  const int n = 3;
  DOUBLE A[m][n] =
    \{6, -1, 0\},
    \{-1, -5, 0\},
    { 1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE _b[m] = { 10, -4, 5, -5 };
 DOUBLE _{c[n]} = \{ 1, -1, 0 \};
  VVD A(m);
  VD b(_b, _b + m);
  VD c(\_c, \_c + n);
  for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
  LPSolver solver(A, b, c);
  VD x:
 DOUBLE value = solver Solve(x):
 cerr << "VALUE: " << value << endl; // VALUE: 1.29032 cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
  cerr << endl:
 return 0:
```

3.8 Fast Fourier transform (C++)

```
// Convolution using the fast Fourier transform (FFT).
// INPIIT ·
      a[1...n]
      b[1...m]
// OUTPUT:
      c[1...n+m-1] such that c[k] = sum_{i=0}^k a[i] b[k-i]
// Alternatively, you can use the DFT() routine directly, which will
// zero-pad your input to the next largest power of 2 and compute the
// DFT or inverse DFT.
typedef long double DOUBLE;
```

```
typedef complex<DOUBLE> COMPLEX;
typedef vector<DOUBLE> VD;
typedef vector<COMPLEX> VC;
struct FFT {
  int n, L;
  int ReverseBits(int k) {
    for (int i = 0; i < L; i++) {
      ret = (ret << 1) | (k & 1);
    return ret;
  void BitReverseCopy(VC a) {
    for (n = 1, L = 0; n < a.size(); n <<= 1, L++);
    for (int k = 0; k < n; k++)</pre>
      A[ReverseBits(k)] = a[k];
  VC DFT(VC a, bool inverse) {
    BitReverseCopy(a);
    for (int s = 1; s <= L; s++) {
      int m = 1 << s;</pre>
      COMPLEX wm = \exp(COMPLEX(0, 2.0 * M_PI / m));
      if (inverse) wm = COMPLEX(1, 0) / wm;
      for (int k = 0; k < n; k += m) {
        COMPLEX w = 1;
         for (int j = 0; j < m/2; j++) {
          COMPLEX t = w * A[k + j + m/2];
           COMPLEX u = A[k + j];
          A[k + j] = u + t;

A[k + j + m/2] = u - t;
           w = w \star wm;
    if (inverse) for (int i = 0; i < n; i++) A[i] /= n;</pre>
    return A;
  // c[k] = sum_{i=0}^k a[i] b[k-i]
  VD Convolution(VD a, VD b) {
    int T_i = 1:
    while ((1 << L) < a.size()) L++;</pre>
    while ((1 << L) < b.size()) L++;
    int n = 1 << (L+1):
    for (size_t i = 0; i < n; i++) aa.push_back(i < a.size() ? COMPLEX
      \text{ for } \substack{ (\text{a[i], 0) : 0);} \\ \text{for } (\text{size\_t i = 0; i < n; i++) bb.push\_back(i < b.size() ? COMPLEX} } \\ \textbf{3.10} \quad \textbf{Big Number} 
           (b[i], 0) : 0);
    VC AA = DFT(aa, false);
    VC BB = DFT(bb, false);
    VC CC:
    for (size_t i = 0; i < AA.size(); i++) CC.push_back(AA[i] * BB[i])</pre>
    VC cc = DFT(CC, true);
    for (int i = 0; i < a.size() + b.size() - 1; i++) c.push_back(cc[i</pre>
          ].real());
    return c;
int main() {
  double a[] = \{1, 3, 4, 5, 7\};
  double b[] = {2, 4, 6};
  VD c = fft.Convolution(VD(a, a + 5), VD(b, b + 3));
  // expected output: 2 10 26 44 58 58 42
  for (int i = 0; i < c.size(); i++) cerr << c[i] << " ";</pre>
  cerr << endl;
  return 0;
```

3.9 Pollard Rho Algorithm

```
#include <cstdio>
using namespace std;
#define abs_val(a) (((a)>=0)?(a):-(a))
typedef long long 11;
ll mulmod(ll a, ll b, ll c) { // returns (a \star b) \$ c, and minimize
  11 x = 0, y = a % c;
  while (b > 0) {
   if (b % 2 == 1) x = (x + y) % c;
   y = (y * 2) % c;
 return x % c;
11 gcd(11 a, 11 b) { return !b ? a : gcd(b, a % b); }
     standard gcd
11 pollard rho(11 n) {
 int i = 0, k = 2;
 11 x = 3, y = 3;
                                 // random seed = 3, other values
       possible
  while (1) {
   i++;
   x = (mulmod(x, x, n) + n - 1) % n;
                                                     // generating
         function
   11 d = gcd(abs_val(y - x), n);
                                                         // the key
          insight
   if (d != 1 && d != n) return d;
                                           // found one non-trivial
         factor
   if (i == k) y = x, k \neq 2;
int main() {
 11 n = 2063512844981574047LL;
                                  // we assume that n is not a large
      prime
 ll ans = pollard rho(n);
                                   // break n into two non trivial
      factors
 if (ans > n / ans) ans = n / ans;
                                            // make ans the smaller
       factor
 printf("%lld %lld\n", ans, n / ans); // should be: 1112041493
       1855607779
} // return 0;
```

```
// Depending on your application it's pretty unlikely that you'll have
      to type out the entirety of this struct. For example, in most
      cases you won't need division, and you can leave out most of the
      operators too.
// NB: The are fairly terrible implementations. Multiplication is
      about twice as slow as Python, and division is about 20 times
      slower. Use only as a last resort when for some reason you can't
      use Java's native bignums.
struct bignum
       typedef unsigned int uint;
       vector<uint> digits:
       static const uint RADIX = 1000000000;
       bignum(): digits(1, 0) {}
       bignum (const bignum& x): digits (x.digits) {}
       bignum (unsigned long long x) {
                *this = x:
       bignum(const char* x) {
                *this = x;
       bignum (const string& s) {
                *this = s;
```

bignum& operator=(const bignum& y) {

```
digits = y.digits; return *this;
bignum& operator=(unsigned long long x) {
        digits.assign(1, x%RADIX);
        if (x >= RADIX) {
                digits.push_back(x/RADIX);
        return *this;
bignum& operator=(const char* s) {
        int slen=strlen(s),i,l;
        digits.resize((slen+8)/9);
        for (1=0; slen>0; 1++,slen-=9) {
                digits[1]=0;
                for (i=slen>9?slen-9:0; i<slen; i++) {
                        digits[1]=10*digits[1]+s[i]-'0';
        while (digits.size() > 1 && !digits.back()) digits.
bignum& operator=(const string& s) {
        return *this = s.c str();
void add(const bignum& x) {
        int l = max(digits.size(), x.digits.size());
        digits.resize(1+1);
        for (int d=0, carry=0; d<=1; d++) {
                uint sum=carry;
                if (d<digits.size()) sum+=digits[d];</pre>
                if (d<x.digits.size()) sum+=x.digits[d];</pre>
                digits[d]=sum;
                if (digits[d]>=RADIX) {
                        digits[d]-=RADIX; carry=1;
                } else {
        if (!digits.back()) digits.pop_back();
void sub(const bignum& x) {
        // if ((*this)<x) throw; //negative numbers not yet
              supported
        for (int d=0, borrow=0; d<digits.size(); d++) {
                digits[d] -= borrow;
                if (d<x.digits.size()) digits[d]-=x.digits[d];
if (digits[d]>>31) { digits[d]+=RADIX; borrow
                       =1; } else borrow=0;
        while (digits.size() > 1 && !digits.back()) digits.
              pop_back();
void mult (const bignum& x) {
        vector<uint> res(digits.size() + x.digits.size());
        unsigned long long y, z;
        for (int i=0; i<digits.size(); i++) {
                for (int j=0; j<x.digits.size(); j++) {</pre>
                        unsigned long long y=digits[i]; y*=x.
                              digits[j];
                         unsigned long long z=y/RADIX;
                         res[i+j+1]+=z; res[i+j]+=y-RADIX*z; //
                               mod is slow
                        if (res[i+j] >= RADIX) { res[i+j] -=
                               RADIX; res[i+j+1]++; }
                         for (int k = i+j+1; res[k] >= RADIX;
                               res[k] -= RADIX, res[++k]++);
        digits = res;
        while (digits.size() > 1 && !digits.back()) digits.
              pop_back();
// returns the remainder
bignum div(const bignum& x)
        bignum dividend(*this);
        bignum divisor(x);
        fill(digits.begin(), digits.end(), 0);
        // shift divisor up
        int pwr = dividend.digits.size() - divisor.digits.size
        if (pwr > 0) {
```

```
divisor.digits.insert(divisor.digits.begin(),
        while (pwr >= 0) {
                if (dividend.digits.size() > divisor.digits.
                        unsigned long long q = dividend.digits
                              .back();
                        q *= RADIX; q += dividend.digits[
                              dividend.digits.size()-2];
                        q /= 1+divisor.digits.back();
                        dividend -= divisor*q; digits[pwr] = q
                        if (dividend >= divisor) { digits[pwr
                              ]++; dividend -= divisor; }
                        assert(dividend.digits.size() <=
                              divisor.digits.size()); continue
                while (dividend.digits.size() == divisor.
                      digits.size()) {
                        uint q = dividend.digits.back() / (1+
                              divisor.digits.back());
                        if (q == 0) break;
                        digits[pwr] += q; dividend -= divisor*
                if (dividend >= divisor) { dividend -= divisor
                     ; digits[pwr]++; }
                pwr--; divisor.digits.erase(divisor.digits.
                      begin());
        while (digits.size() > 1 && !digits.back()) digits.
              pop_back();
        return dividend;
string to_string() const {
        ostringstream oss;
        oss << digits.back();
        for (int i = digits.size() - 2; i >= 0; i--) {
                oss << setfill('0') << setw(9) << digits[i];
        return oss.str();
bignum operator+(const bignum& v) const
        bignum res(*this); res.add(y); return res;
bignum operator-(const bignum@ y) const {
        bignum res(*this); res.sub(y); return res;
bignum operator*(const bignum& v) const {
        bignum res(*this); res.mult(v); return res;
bignum operator/(const bignum& v) const {
        bignum res(*this); res.div(v); return res;
bignum operator% (const bignum& v) const {
        bignum res(*this); return res.div(y);
bignum& operator+=(const bignum& y) {
        add(v): return *this:
bignum& operator-=(const bignum& y) {
        sub(y); return *this;
bignum& operator *= (const bignum& y) {
        mult(v); return *this;
bignum& operator/=(const bignum& y) {
        div(y); return *this;
bignum& operator%=(const bignum& y) {
        *this = div(y);
bool operator == (const bignum& y) const {
        return digits == y.digits;
```

```
bool operator < (const bignum& y) const {
                 if (digits.size() < y.digits.size()) return true;</pre>
                 if (digits.size() > y.digits.size()) return false;
                 for (int i = digits.size()-1; i >= 0; i--) {
                         if (digits[i] < y.digits[i]) {</pre>
                                  return true;
                           else if (digits[i] > y.digits[i]) {
                                  return false;
                 return false;
        bool operator > (const bignum& y) const {
                 return v<*this;
        bool operator<=(const bignum& y) const {</pre>
                 return ! (y<*this);
        bool operator>=(const bignum& y) const {
                 return ! (*this<y);</pre>
};
```

4 Graph algorithms

4.1 Dijkstra's algorithm

```
int V, E, s, u, v, w;
vector<vii> AdjList;
int main() {
    cin>>V>>E>>s;
    AdjList.assign(V, vii());
    for(int i = 0; i < E; i++) {
        cin>>u>>v>>w;
        AdjList[u].push_back(ii(v, w));
    ,
// Dijkstra routine
    vi dist(V, INF); dist[s] = 0;
    // distance, node
    priority_queue<ii, vector<ii>, greater<ii> > pq;
   pq.push(ii(0,s));
    while(!pq.empty()) {
       ii front = pq.top(); pq.pop();
        int d = front.first, u = front.second;
        if (d > dist[u]) continue; // handle duplicates
        for(int j = 0; j < (int) AdjList[u].size(); j++) {</pre>
            ii v = AdjList[u][j];
            if (dist[u] + v.second < dist[v.first]) {</pre>
                dist[v.first] = dist[u] + v.second;
                pq.push(ii(dist[v.first], v.first));
    // SPFA (Faster Bellman Ford)
   queue<int> q; q.push(S);
   vi in queue(n, 0); in queue(S) = 1;
    while(!q.emptv()) {
       int u = q.front(); q.pop(); in_queue[u] = 0;
        for(auto v : AdjList[u]) {
            if (dist[u] + v.second < dist[v.first]) {
                dist[v.first] = dist[u] + v.second;
                if (!in_queue[v.first]) {
                    q.push(v.first); // add to queue only if it's not
                          in queue
                    in_queue[v.first] = 1;
```

4.2 Strongly connected components

```
bool adjList[55][55];
vi dfs_num, dfs_low, S, visited;
int dfs_counter = 0, numSCC = 0;
int m, n; // nLocations, nEdges
void tarjanSCC(int u) {
    dfs_low[u] = dfs_num[u] = dfs_counter++;
    S.push_back(u);
    visited[u] = 1;
    for (int v = 0; v < m; v++) {
        if (adjList[u][v]) {
            if(dfs_num[v] == -1) tarjanSCC(v);
             if(visited[v] == 1)
                 dfs_low[u] = min(dfs_low[v], dfs_low[u]);
    if(dfs_low[u] == dfs_num[u]) { //root
        printf("SCC %d\n", ++numSCC);
         while(1) {
            int v = S.back(); S.pop_back(); visited[v] = 0;
             printf("%d ", v);
             if(u == v) break;
        printf("\n");
int main() {
    m=8,n=9; // edges and vertices
    adjList[0][1] = true; adjList[1][3] = true; adjList[3][4] = true;
    adjList[4][5] = true; adjList[5][7] = true; adjList[6][6] = true; adjList[6][4] = true; adjList[3][2] = true; adjList[2][1] = true;
    dfs_num = vi(m, -1); dfs_low = vi(m, 0); visited = vi(m, 0);
    dfs_counter = numSCC = 0;
    for(int i = 0; i < m; i++) {</pre>
        if (dfs num[i] == -1) {
            tarjanSCC(i);
    // 1: 6 7 4 5
    // 2: 2 3 1
    // 3: 0
```

4.3 Eulerian path

```
struct Edge:
typedef list<Edge>::iterator iter;
struct Edge
         int next vertex:
        iter reverse edge:
         Edge(int next vertex)
                 :next_vertex(next_vertex)
};
const int max_vertices = ;
int num vertices:
                                            // adjacency list
list<Edge> adj[max_vertices];
vector<int> path:
void find_path(int v)
         while (adj[v].size() > 0)
                 int vn = adj[v].front().next_vertex;
adj[vn].erase(adj[v].front().reverse_edge);
                  adj[v].pop_front();
                  find_path(vn);
         path.push_back(v);
```

```
void add_edge(int a, int b)
{
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
}
```

4.4 Kruskal's algorithm

```
Uses Kruskal's Algorithm to calculate the weight of the minimum
      spanning
forest (union of minimum spanning trees of each connected component)
a possibly disjoint graph, given in the form of a matrix of edge
      weights
(-1 if no edge exists). Returns the weight of the minimum spanning
forest (also calculates the actual edges - stored in T). Note: uses a
disjoint-set data structure with amortized (effectively) constant time
union/find. Runs in O(E*log(E)) time.
typedef int T;
struct edge
  int u. v:
 T d;
};
struct edgeCmp
  int operator()(const edge& a, const edge& b) { return a.d > b.d; }
int find(vector \langle int \rangle \& C, int x) { return (C[x] == x) ? x : C[x] =
      find(C, C[x]); }
T Kruskal (vector <vector <T> >& w)
  int n = w.size();
 T weight = 0;
  vector <int> C(n), R(n);
  for(int i=0; i<n; i++) { C[i] = i; R[i] = 0; }</pre>
  vector <edge> T;
 priority_queue <edge, vector <edge>, edgeCmp> E;
  for(int i=0; i<n; i++)
    for(int j=i+1; j<n; j++)</pre>
     if(w[i][j] >= 0)
        e.u = i; e.v = j; e.d = w[i][j];
       E.push(e);
  while (T.size() < n-1 && !E.empty())
    edge cur = E.top(); E.pop();
    int uc = find(C, cur.u), vc = find(C, cur.v);
    if (nc != vc)
     T.push_back(cur); weight += cur.d;
      if(R[uc] > R[vc]) C[vc] = uc;
      else if(R[vc] > R[uc]) C[uc] = vc;
      else { C[vc] = uc; R[uc]++; }
  return weight;
int main()
 int wa[6][6] = {
```

4.5 Minimum spanning trees

```
// This function runs Prim's algorithm for constructing minimum
// weight spanning trees.
// Running time: O(|V|^2)
     INPUT: w[i][j] = cost of edge from i to j
               NOTE: Make sure that w[i][j] is nonnegative and
               symmetric. Missing edges should be given -1
     OUTPUT: edges = list of pair<int,int> in minimum spanning tree
               return total weight of tree
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
T Prim (const VVT &w, VPII &edges) {
 int n = w.size();
  VI found (n);
  VI prev (n, -1);
 VT dist (n, 1000000000);
 int here = 0;
 dist[here] = 0;
  while (here != -1) {
   found[here] = true;
    int best = -1:
   for (int k = 0; k < n; k++) if (!found[k]) {
   if (w[here][k] != -1 && dist[k] > w[here][k]) {
        dist[k] = w[here][k];
        prev[k] = here;
      if (best == -1 || dist[k] < dist[best]) best = k;</pre>
   here = best;
 T tot_weight = 0;
for (int i = 0; i < n; i++) if (prev[i] != -1) {</pre>
   edges.push_back (make_pair (prev[i], i));
   tot_weight += w[prev[i]][i];
 return tot_weight;
int main() {
 int ww[5][5] = {
   {0, 400, 400, 300, 600},
    \{400, 0, 3, -1, 7\},\
    {400, 3, 0, 2, 0},
    {300, -1, 2, 0, 5},
   {600, 7, 0, 5, 0}
  VVT w(5, VT(5));
  for (int i = 0; i < 5; i++)
   for (int j = 0; j < 5; j++)
w[i][j] = ww[i][j];</pre>
 // expected: 305
```

```
// 2 1
// 3 2
// 0 3
// 0 3
// 2 4

VPII edges;
cout << Prim (w, edges) << endl;
for (int i = 0; i < edges.size(); i++)
cout << edges[i].first << " " << edges[i].second << endl;
}
```

4.6 Lowest common ancestor

const int max_nodes, log_max_nodes;

```
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                            // children[i] contains the
      children of node i
int A[max_nodes][log_max_nodes+1];
                                            // A[i][j] is the 2^j-th
      ancestor of node i, or -1 if that ancestor does not exist
                                            // L[i] is the distance
int L[max_nodes];
      between node i and the root
// floor of the binary logarithm of n
int lb(unsigned int n)
    if(n==0)
        return -1:
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16;
    if (n >= 1<< 8) { n >>= 8; p += 8;
    if (n >= 1<< 4) { n >>= 4; p += 4; }
    if (n >= 1<< 2) { n >>= 2; p += 2; }
    if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    for(int j = 0; j < children[i].size(); j++)
    DFS(children[i][j], l+1);</pre>
int LCA(int p, int q)
    // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node p situated on the same
           level as q
    for (int i = log_num_nodes; i >= 0; i--)
   if (L[p] - (1<<i) >= L[q])
            p = A[p][i];
    if(p == q)
        return p;
    // "binary search" for the LCA
    for (int i = log_num_nodes; i >= 0; i--)
   if (A[p][i] != -1 && A[p][i] != A[q][i])
             p = A[p][i];
             q = A[q][i];
    return A[p][0];
int main(int argc,char* argv[])
    // read num nodes, the total number of nodes
    log num nodes=1b(num nodes);
    for(int i = 0; i < num_nodes; i++)</pre>
        // \bar{\text{read}} p, the parent of node i or -1 if node i is the root
        A[i][0] = p;
if(p != -1)
            children[p].push_back(i);
        else
             root = i;
```

4.7 Bridge and Articulation Points

```
#include <bits/stdc++.h>
using namespace std;
#define FOR(x,n) for(int x = 0; x < n; ++x)
typedef unsigned long long 11;
typedef pair<int, int> ii;
typedef vector<int> vi;
typedef vector<ii> vii;
typedef vector<string> vs;
typedef vector<vi> vvi;
#define UNVISITED 0
int dfsCounter, rootChildren, dfsRoot;
vi dfs_num, dfs_low, dfs_parent, art_vertex;
void artPointAndBridge(int u) {
        dfs_low[u] = dfs_num[u] = dfsCounter++;
        for(auto &v : AdjList[u]) {
                 if (dfs_num[v] == UNVISITED) {
                          dfs_parent[v] = u;
                          if(u == dfsRoot) rootChildren++;
                          artPointAndBridge(v);
                         if(dfs_low[v] >= dfs_num[v]) // art point
    art vertex[v] = true;
                          if(dfs_low[v] > dfs_num[u]) {
                                  // (u,v) is bridge
                                  printf("(%d,%d) is a bridge\n", u, v);
                          //update dfs_low
                 dfs_low[u] = min(dfs_low[u], dfs_low[v]);
} else if(v != dfs_parent[u]) {
                         // back edge and not direct cycle
dfs_low[u] = min(dfs_low[u], dfs_low[v]);
int main() {
        int v = 6:
        AdiList.assign(v, vi());
        AdjList[0].push_back(1); AdjList[1].push_back(0);
        AdjList[1].push_back(2); AdjList[2].push_back(1);
        AdjList[1].push_back(3); AdjList[3].push_back(1);
        AdjList[1].push_back(4); AdjList[4].push_back(1);
        AdjList[4].push_back(5); AdjList[5].push_back(4);
        AdjList[1].push_back(5); AdjList[5].push_back(1);
        dfsCounter = 0, dfs_num.assign(v, UNVISITED), dfs_low.assign(v
               , 0);
        dfs_parent.assign(v, 0), art_vertex.assign(v, 0);
        printf("Bridges\n"); // (0,1), (1,2), (1,3)
        for(int i = 0; i < v; i++) {
    if(dfs_num[i] == UNVISITED) {</pre>
                         dfsRoot = 1, rootChildren = 0,
                                artPointAndBridge(i);
                          art_vertex[dfsRoot] = (rootChildren > 1); //
                                special case
```

5 Data structures

5.1 Binary Indexed Tree

```
// BIT with range updates, inspired by Petr Mitrichev
struct BIT (
   int n:
    vector<int> slope:
    vector<int> intercept;
    // BIT can be thought of as having entries f[1], ..., f[n]
    // which are 0-initialized
   BIT(int n): n(n), slope(n+1), intercept(n+1) {}
    // returns f[1] + ... + f[idx-1]
    // precondition idx <= n+1
   int query(int idx) {
       int m = 0, b = 0;
        for (int i = idx-1; i > 0; i -= i&-i) {
            m += slope[i];
           b += intercept[i];
       return m*idx + b;
    ^{\prime\prime} // adds amt to f[i] for i in [idx1, idx2)
    // precondition 1 <= idx1 <= idx2 <= n+1 (you can't update element
    void update(int idx1, int idx2, int amt) {
        for (int i = idx1; i <= n; i += i&-i) {
            slope[i] += amt;
            intercept[i] -= idx1*amt;
        for (int i = idx2; i <= n; i += i&-i) {
            slope[i] -= amt;
            intercept[i] += idx2*amt;
// BIT with range updates, inspired by Petr Mitrichev
class FenwickTree
private: vi ft1, ft2;
   int query(vi &ft, int b) {
        int sum = 0; for (; b; b -= LSOne(b)) sum += ft[b];
    void adjust(vi &ft, int k, int v) {
       for (; k < (int)ft.size(); k += LSOne(k)) ft[k] += v; }</pre>
   FenwickTree() {}
    FenwickTree(int n) { ft1.assign(n + 1, 0); ft2.assign(n+1, 0);}
    int query(int a) { return a * query(ft1, a) - query(ft2, a); }
    int query (int a, int b) { return query (b) - (a == 1 ? 0 : query (a
    void adjust (int a, int b, int value) {
       adjust (ft1, a, value);
        adjust (ft1, b+1, -value);
       adjust(ft2, a, value * (a-1));
       adjust (ft2, b+1, -1 * value * b);
    int get(int n) {
       return query(n) - query(n-1);}
```

5.2 2D Binary Indexed Tree

```
// WARNING NOT FIELD TESTED YET
class FenwickTree {
private:
   vi ft;
```

```
public:
  FenwickTree() {}
  // initialization: n + 1 zeroes, ignore index 0
  FenwickTree(int n) { ft.assign(n + 1, 0); }
  int rsq(int b) {
                                                            // returns RSQ
        (1, b)
    int sum = 0; for (; b; b -= LSOne(b)) sum += ft[b];
    return sum;
  int rsq(int a, int b) {
                                                            // returns RSQ(
        a, b)
    return rsq(b) - (a == 1 ? 0 : rsq(a - 1)); }
  // adjusts value of the k-th element by v (v can be +ve/inc or -ve/
  void adjust (int k, int v) {
                                                     // note: n = ft.size
    for (; k < (int)ft.size(); k += LSOne(k)) ft[k] += v; }</pre>
class FenwickTree2D {
private:
  vector<FenwickTree> ft2d;
  FenwickTree2D() {}
  FenwickTree2D(int n) { ft2d.assign(n+1, FenwickTree(n)); }
  int rsq(int r, int c) {
    int sum = 0;
    for(; r; r \rightarrow LSOne(r)) sum += ft2d[r].rsq(c);
    return sum:
  // top left, bottom right
  int rsq(int r1, int c1, int r2, int c2) {
    return rsq(r2, c2) - rsq(r2, c1-1) - rsq(r1-1, c2) + rsq(r1-1, c2
           -1);
  void adjust(int r, int c, int v) {
    for (; r < (int)ft2d.size(); r += LSOne(r)) ft2d[r].adjust(c, v);</pre>
1:
int main() {
 FenwickTree2D ft2d(4);
  ft2d.adjust(1, 1, 1);
  ft2d.adjust(2, 2, 1);
  ft2d.adjust(3, 3, 1);
 ft2d.adjust(4, 4, 1);

printf("%d\n", ft2d.rsq(1,1));

printf("%d\n", ft2d.rsq(2,2));

printf("%d\n", ft2d.rsq(3,3));
                                         // 2
  printf("%d\n", ft2d.rsq(2,2,3,3)); // 2
  return 0:
```

5.3 Union-find

```
struct UnionFind {
    vector<int> C;
    UnionFind(int n) : C(n) { for (int i = 0; i < n; i++) C[i] = i; }
    int find(int x) { return (C[x] == x) ? x : C[x] = find(C[x]); }
    void merge(int x, int y) { C[find(x)] = find(y); }
};
int main()
{
    int n = 5;
    UnionFind uf(n);
    uf.merge(0, 2);
    uf.merge(1, 0);
    uf.merge(1, 0);
    uf.merge(3, 4);
    for (int i = 0; i < n; i++) cout << i << " " << uf.find(i) << endl
    return 0;</pre>
```

5.4 KD-tree

```
// A straightforward, but probably sub-optimal KD-tree implmentation
// that's probably good enough for most things (current it's a
// 2D-tree)
// - constructs from n points in O(n 1g^2 n) time
// - handles nearest-neighbor query in O(lg n) if points are well
    distributed
// - worst case for nearest-neighbor may be linear in pathological
    case
#include <iostream>
#include <vector>
#include inits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
    point (ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
bool operator == (const point &a, const point &b)
    return a.x == b.x && a.y == b.y;
// sorts points on x-coordinate
bool on_x (const point &a, const point &b)
    return a.x < b.x;
// sorts points on y-coordinate
bool on_y (const point &a, const point &b)
    return a.v < b.v;
// squared distance between points
ntype pdist2(const point &a, const point &b)
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dv*dv:
// bounding box for a set of points
struct bbox
    ntype x0, x1, y0, y1;
    bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
    // computes bounding box from a bunch of points
    void compute (const vector<point> &v) {
       for (int i = 0; i < v.size(); ++i) {
            x0 = min(x0, v[i].x); x1 = max(x1, v[i].x);
           y0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
    // squared distance between a point and this bbox, 0 if inside
    ntype distance (const point &p) {
        if (p.x < x0) {
                               return pdist2(point(x0, y0), p);
            if (p.y < y0)
            else if (p.y > y1)
                               return pdist2(point(x0, y1), p);
            else
                               return pdist2(point(x0, p.y), p);
        else if (p.x > x1) {
           if (p.y < y0)
                               return pdist2(point(x1, y0), p);
            else if (p.y > y1) return pdist2(point(x1, y1), p);
            else
                               return pdist2(point(x1, p.y), p);
        else
            if (p.y < y0)
                               return pdist2(point(p.x, y0), p);
            else if (p.y > y1) return pdist2(point(p.x, y1), p);
            else
                               return 0;
```

```
};
// stores a single node of the kd-tree, either internal or leaf
struct kdnode
   bool leaf:
                    // true if this is a leaf node (has one point)
   point pt;
                    // the single point of this is a leaf
    bbox bound;
                    // bounding box for set of points in children
    kdnode *first, *second; // two children of this kd-node
    kdnode() : leaf(false), first(0), second(0) {}
    "kdnode() { if (first) delete first; if (second) delete second; }
    // intersect a point with this node (returns squared distance)
    ntype intersect (const point &p) {
        return bound.distance(p);
    // recursively builds a kd-tree from a given cloud of points
    void construct (vector<point> &vp)
        // compute bounding box for points at this node
        bound.compute(vp);
        // if we're down to one point, then we're a leaf node
        if (vp.size() == 1) {
           pt = vp[0];
        else
            // split on x if the bbox is wider than high (not best
                 heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
            // otherwise split on y-coordinate
           else
                sort(vp.begin(), vp.end(), on v);
            // divide by taking half the array for each child
            // (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half);
            vector<point> vr(vp.begin()+half, vp.end());
           first = new kdnode(); first->construct(vl);
            second = new kdnode(); second->construct(vr);
};
// simple kd-tree class to hold the tree and handle queries
struct kdtree
   kdnode *root;
    // constructs a kd-tree from a points (copied here, as it sorts
         them)
   kdtree(const vector<point> &vp) {
       vector<point> v(vp.begin(), vp.end());
        root = new kdnode();
       root->construct(v):
    "kdtree() { delete root; }
    // recursive search method returns squared distance to nearest
          point
    ntype search(kdnode *node, const point &p)
       if (node->leaf) {
            // commented special case tells a point not to find itself
             if (p == node->pt) return sentry;
              else
               return pdist2(p, node->pt);
       ntype bfirst = node->first->intersect(p);
       ntype bsecond = node->second->intersect(p);
        // choose the side with the closest bounding box to search
        // (note that the other side is also searched if needed)
        if (bfirst < bsecond) {
            ntype best = search(node->first, p);
            if (bsecond < best)</pre>
               best = min(best, search(node->second, p));
            return best;
```

```
else {
           ntype best = search(node->second, p);
           if (bfirst < best)</pre>
               best = min(best, search(node->first, p));
                                                                               inline void rotate(Node *x, int c)
                                                                                 Node *y = x->pre;
                                                                                 x->pre = y->pre;
                                                                                 if(y->pre != null)
   // squared distance to the nearest
                                                                                  y->pre->ch[y == y->pre->ch[1]] = x;
   ntype nearest (const point &p) {
                                                                                 y - ch[!c] = x - ch[c];
       return search (root, p);
                                                                                 if (x->ch[c] != null)
                                                                                   x->ch[c]->pre = y;
                                                                                 x->ch[c] = y, y->pre = x;
                                                                                 update(y);
                                                                                 if(y == root)
                                                                                  root = x;
// some basic test code here
                                                                               void splay(Node *x, Node *p)
int main()
                                                                                 while (x->pre != p)
   // generate some random points for a kd-tree
   for (int i = 0; i < 100000; ++i) {
       vp.push_back(point(rand()%100000, rand()%100000));
   kdtree tree(vp);
   // query some points
   for (int i = 0; i < 10; ++i) {
       point q(rand()%100000, rand()%100000);
       cout << "Closest squared distance to (" << q.x << ", " << q.y
             << " is " << tree.nearest(q) << endl;
   return 0;
                                                                                 update(x);
```

Splay tree

```
Node *now = root:
                                                                                  while(1)
const int N MAX = 130010;
const int oo = 0x3f3f3f3f3f;
                                                                                    pushDown (now);
                                                                                    int tmp = now->ch[0]->size + 1;
struct Node
                                                                                    if (tmp == k)
                                                                                     break:
 Node *ch[2], *pre;
                                                                                    else if(tmp < k)
 int val, size;
                                                                                     now = now -> ch[1], k -= tmp;
 bool isTurned:
                                                                                    else
} nodePool[N_MAX], *null, *root;
                                                                                     now = now -> ch[0];
Node *allocNode(int val)
                                                                                  splay(now, fa);
  static int freePos = 0:
 Node *x = &nodePool[freePos ++];
 x->val = val, x->isTurned = false;
                                                                               Node *makeTree(Node *p, int 1, int r)
  x->ch[0] = x->ch[1] = x->pre = null;
                                                                                 if(1 > r)
 x->size = 1;
                                                                                   return null;
 return x:
                                                                                 int mid = (1 + r) / 2;
                                                                                 Node *x = allocNode(mid);
inline void update (Node *x)
                                                                                 x-sch[0] = makeTree(x, 1, mid - 1);
                                                                                 x->ch[1] = makeTree(x, mid + 1, r);
  x->size = x->ch[0]->size + x->ch[1]->size + 1;
                                                                                 update(x);
                                                                                 return x;
inline void makeTurned (Node *x)
                                                                               int main()
  if(x == null)
   return:
  swap(x->ch[0], x->ch[1]);
                                                                                 int n, m;
                                                                                 null = allocNode(0);
 x->isTurned ^= 1;
                                                                                 null->size = 0;
                                                                                 root = allocNode(0);
                                                                                 root->ch[1] = allocNode(oo);
inline void pushDown (Node *x)
                                                                                  root->ch[1]->pre = root;
                                                                                 update (root);
  if(x->isTurned)
   makeTurned(x->ch[0]);
                                                                                 root->ch[1]->ch[0] = makeTree(root->ch[1], 1, n);
   makeTurned(x->ch[1]);
                                                                                 splay(root->ch[1]->ch[0], null);
   x->isTurned ^= 1;
```

rotate(x, x == x->pre->ch[0]);

Node *y = x->pre, *z = y->pre;

rotate(y, 1), rotate(x, 1);

rotate(x, 0), rotate(x, 1);

rotate(y, 0), rotate(x, 0);

rotate(x, 1), rotate(x, 0);

if (y == z->ch[0])

else

if(x == y->ch[0])

if(x == y->ch[1])

void select(int k, Node *fa)

```
while (m --)
 int a, b;
 scanf("%d%d", &a, &b);
 a ++, b ++;
 select(a - 1, null);
 select(b + 1, root);
 makeTurned(root->ch[1]->ch[0]);
for(int i = 1; i <= n; i ++)
 select(i + 1, null);
 printf("%d ", root->val);
```

5.6 Splay Link Cut Trees

```
const int MAXN = 110000;
typedef struct _node{
  _node *1, *r, *p, *pp;
  int size; bool rev;
  _node();
  explicit _node(nullptr_t) {
    1 = r = p = pp = this;
 void push() {
   if (rev) {
      1->rev ^= 1; r->rev ^= 1;
      rev = 0; swap(l,r);
 void update();
}* node;
node None = new _node(nullptr);
node v2n[MAXN];
_node::_node(){
 1 = r = p = pp = None;
 size = 1; rev = false;
void node::update(){
 size = (this != None) + 1->size + r->size;
 1->p = r->p = this;
void rotate (node v) {
 assert (v != None && v->p != None);
 assert(!v->rev); assert(!v->p->rev);
 node 11 = v - > p:
 if (v == u->1)
   u->1 = v->r, v->r = u;
 else
   u->r = v->1, v->1 = u;
  swap(u->p,v->p); swap(v->pp,u->pp);
 if (v->p != None) {
   assert(v->p->l == u || v->p->r == u);
if (v->p->r == u) v->p->r = v;
   else v->p->1 = v;
 u->update(); v->update();
void bigRotate(node v) {
 assert (v->p != None);
 v->p->p->push();
 v->p->push();
  v->push();
 v->pusn();
if (v->p->p != None) {
  if ((v->p->1 == v) ^ (v->p->r == v->p))
      rotate(v->p);
    else
      rotate(v):
 rotate(v);
inline void Splay(node v) {
 while (v->p != None) bigRotate(v);
inline void splitAfter(node v) {
 v->push();
 Splay(v);
 v->r->p = None;
```

```
v->r->pp = v;
 v->r = None;
 v->update();
void expose(int x) {
 node v = v2n[x];
  splitAfter(v);
 while (v->pp != None)
   assert (v->p == None);
   splitAfter(v->pp);
   assert (v->pp->r == None);
   assert (v->pp->p == None);
   assert (!v->pp->rev);
   v->pp->r = v;
   v->pp->update();
   v = v - pp;
   v->r->pp = None;
  assert (v->p == None);
 Splay(v2n[x]);
inline void makeRoot(int x) {
 expose(x);
  assert(v2n[x]->p == None);
 assert (v2n[x]->pp == None);
 assert (v2n[x]->r == None);
 v2n[x]->rev ^= 1;
inline void link(int x, int y) {
 makeRoot(x); v2n[x]->pp = v2n[y];
inline void cut(int x,int y) {
 expose(x);
  Splay(v2n[y]);
 if (v2n[y]->pp != v2n[x]) {
   swap(x,v);
   expose(x);
   Splay(v2n[y]);
   assert (v2n[y]->pp == v2n[x]);
 v2n[y] -> pp = None;
inline int get(int x,int y) {
 if (x == y) return 0;
 makeRoot(x):
 expose(y); expose(x);
 Splay(v2n[y]);
 if (v2n[y]->pp != v2n[x]) return -1;
return v2n[y]->size;
```

5.7 Sparse Table

```
#include <algorithm>
#include <cmath>
#include <cstdio>
using namespace std:
#define MAX N 1000
                                                 // adjust this value as
      needed
#define LOG_TWO_N 10
                                   // 2^10 > 1000, adjust this value as
      needed
class RMO
                                                           // Range Minimum
      Ouerv
private:
  int _A[MAX_N], SpT[MAX_N][LOG_TWO_N];
  RMQ(int n, int A[]) { // constructor as well as pre-processing
        routine
    for (int i = 0; i < n; i++) {</pre>
      A[i] = A[i];
      {\tt SpT[i][0]} = {\tt i;} \ // \ {\tt RMQ} \ {\tt of} \ {\tt sub} \ {\tt array} \ {\tt starting} \ {\tt at} \ {\tt index} \ {\tt i} \ + \ {\tt length}
    // the two nested loops below have overall time complexity = O(n
           log n)
    for (int j = 1; (1<<j) <= n; j++) // for each j s.t. 2^j <= n, 0(
      for (int i = 0; i + (1<<j) - 1 < n; i++) // for each valid i,
              O(n)
        if (_A[SpT[i][j-1]] < _A[SpT[i+(1<<(j-1))][j-1]])</pre>
           SpT[i][j] = SpT[i][j-1]; // start at index i of length
                 2^(j-1)
```

5.8 Lazy Segment Tree

#include <iostream>

#include <algorithm>

```
using namespace std:
const int n=10; // number of elements in the tree should be at most
int arr[1<<(n+1)]; // store values</pre>
int low[1<<(n+1)]; // left of range</pre>
int high[1<<(n+1)]; // right of range</pre>
int lazyadd[1<<(n+1)];</pre>
// should be commutative and associative. most commonly used: a+b, max
      (a,b), min(a,b)
int acc(int a, int b) {
    return max(a,b);
int acc2(int a, int b) {
    // a "acts" on b, the function needs to be distributive over acc,
          commutative + associative
    // common uses: a+max(b,c) = max(a+b, a+c), a*max(b,c)=max(a*b,a*c
          ) (only if a>0)
    // a*(b+c) = (a*b)+(a*c)
    return a+b;
void init() {
    for(int i=0; i<(1<<n); i++) {
        low[i+(1<< n)] = i;
        high[i+(1<< n)] = i:
        arr[i+(1<< n)] = 0; // initial value
    for(int i=(1<<n)-1; i>=0; i--) {
        low[i] = min(low[2*i], low[2*i+1]);
        high[i] = max(high[2*i], high[2*i+1]);
        arr[i] = acc(arr[2*i], arr[2*i+1]);
    for(int i=0; i<(1<<(n+1)); i++) {
        lazyadd[i] = 0; // identity of the acc2 function
int value(int node) { // gives the true value of the node
    arr[node] = acc2(lazyadd[node], arr[node]);
    if(node<(1<<n)) { // not the leaf, propagate downwards</pre>
        lazyadd[2*node] = acc2(lazyadd[node], lazyadd[2*node]); //
              stack values on children
        lazyadd[2*node+1] = acc2(lazyadd[node], lazyadd[2*node+1]); //
               stack values on children
    lazyadd[node] = 0; // reset to identity function
    return arr[node];
void update(int node, int left, int right, int change)
    if(right>=high[node] && left<=low[node]) { // case 1: updated</pre>
          range covers node
        lazyadd[node] = acc2(lazyadd[node],change); // stack the
              change
    } else if(right<low[node] || left>high[node]) { // case 2: empty
          intersection
        return;
    } else { // case 3: need to propagate
```

```
update(2*node, left, right, change);
        update(2*node+1, left, right, change);
       arr[node] = acc(value(node*2), value(node*2+1));
void update(int left, int right, int change) {
    update(1, left, right, change);
int query(int node, int left, int right) {
    value(node); // important to call this!
    if(right>=high[node] && left<=low[node])</pre>
       return arr[node];
     else if(right<low[node] || left>high[node]) {
       return -(1<<30); // identity operator of acc
       return acc(query(node*2, left, right), query(node*2+1, left,
int query(int left, int right) {
   return query(1, left, right);
   cout << query(1,5); //0
   update(2,5,3); // 0 3 3 3 3 -> max=3
   cout << query(1,5); // 3
   update(2,3,2); // 0 5 5 3 3 -> max=5
   cout << query(1,5); // 5
   update(2,4,-4); // 0 1 1 -1 3 -> max=5
   cout << query(1,5); // 5
   update(5,5,-1); // 0 1 1 -1 2 -> max=2
   cout << query(1,5); // 2
   return 0;
```

6 String

6.1 Suffix array

```
// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
             of substring s[i...L-1] in the list of sorted suffixes.
             That is, if we take the inverse of the permutation suffix
             we get the actual suffix array.
struct SuffixArray {
 const int L:
 string s:
  vector<vector<int> > P:
  vector<pair<pair<int,int>,int> > M;
  \texttt{SuffixArray}(\textbf{const} \texttt{ string \&s)} \; : \; \texttt{L(s.length()), s(s), P(1, vector<\textbf{int})}
       >(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {</pre>
     P.push_back(vector<int>(L, 0));
      for (int i = 0; i < L; i++)
        M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[
              level-1][i + skip] : -1000), i);
      sort(M.begin(), M.end());
      for (int i = 0; i < L; i++)
        P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first)
               ? P[level][M[i-1].second] : i;
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and
         s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
   int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
      if (P[k][i] == P[k][j]) {
```

```
i += 1 << k:
        i += 1 << k:
        len += 1 << k;
   return len;
// BEGIN CUT
// The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
 int T;
  for (int caseno = 0; caseno < T; caseno++) {</pre>
   cin >> s;
   SuffixArray array(s);
   vector<int> v = array.GetSuffixArray();
   int bestlen = -1, bestpos = -1, bestcount = 0;
   for (int i = 0; i < s.length(); i++) {</pre>
     int len = 0, count = 0;
      for (int j = i+1; j < s.length(); j++) {</pre>
       int 1 = array.LongestCommonPrefix(i, j);
       if (1 >= len) {
          if (1 > len) count = 2; else count++;
          len = 1;
      if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen
          ) > s.substr(i, len)) {
        bestlen = len;
       bestcount = count;
       bestpos = i;
   if (bestlen == 0) {
     cout << "No repetitions found!" << endl;
     cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>
#else
// END CUT
int main() {
  // bobocel is the O'th suffix
  // obocel is the 5'th suffix
     bocel is the 1'st suffix
       ocel is the 6'th suffix
        cel is the 2'nd suffix
         el is the 3'rd suffix
          1 is the 4'th suffix
  SuffixArray suffix("bobocel");
  vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
 for (int i = 0; i < v.size(); i++) cout << v[i] << " ";
 cout << endl:
 cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
// Yan Hao's Implementation
#include <bits/stdc++.h>
using namespace std;
#define all(o) (o).begin(), (o).end()
#define allr(o) (o).rbegin(), (o).rend()
const int INF = 2147483647;
typedef long long 11;
typedef pair<int, int> ii;
typedef vector<int> vi;
typedef vector<ii> vii;
typedef vector<vi> vvi;
typedef vector<vii> vvii;
template <class T> int size(T &x) { return x.size(); }
// assert or gtfo
struct suffix array {
```

```
struct entry {
        pair<int, int> nr;
        int p;
        bool operator < (const entry &other) const {
            return nr < other.nr;
    string s;
    int n;
    vector<vector<int> > P;
    vector<entry> L;
    vector<int> idx;
    suffix_array(string _s) : s(_s), n(s.size()) {
       L = vector<entry>(n);
        P.push_back(vector<int>(n));
        idx = vector<int>(n);
        for (int i = 0; i < n; i++) {
            P[0][i] = s[i];
        for (int stp = 1, cnt = 1; (cnt >> 1) < n; stp++, cnt <<= 1) {
            P.push_back(vector<int>(n));
            for (int i = 0; i < n; i++) {
                L[i].p = i;
                L[i].nr = make_pair(P[stp - 1][i],
                        i + cnt < n ? P[stp - 1][i + cnt] : -1);
            sort(L.begin(), L.end());
            for (int i = 0; i < n; i++) {
                if (i > 0 && L[i].nr == L[i - 1].nr) {
                   P[stp][L[i].p] = P[stp][L[i - 1].p];
                   P[stp][L[i].p] = i;
        for (int i = 0; i < n; i++) {
            idx[P[P.size() - 1][i]] = i;
    int lcp(int x, int y) {
        int res = 0;
       if (x == y) return n - x;
for (int k = P.size() - 1, k >= 0 && x < n && y < n; k--) {
            if (P[k][x] == P[k][y]) {
               x += 1 << k;
                v += 1 << k;
                res += 1 << k;
        return res:
    int longestRepeatedSubsequence() {
        int ans=0:
        for(int i=1: i<n: i++) {</pre>
            ans = max(ans, lcp(idx[i-1], idx[i]));
        return ans:
int main() {
    while(true) {
        string x,y;
        getline(cin, x);
        getline(cin, y);
        y.push_back('$');
        if(!cin.good()) break;
        suffix_array sa(y);
        //do binary search
        int low = 0;
        int high = y.length();
        while(high-low>1) {
            int mid = (high+low)/2;
            int z = sa.idx[mid];
            for(int i=0; i<x.length() && z+i<y.length(); i++) {</pre>
                if(x[i]<y[z+i]) {
                    high=mid; break;
                } else if(x[i]>y[z+i]) {
                    low=mid; break;
```

};

```
if (high!=mid && low!=mid) {
            high=mid;
   int low2 = 0;
   int high2 = y.length();
    while (high2-low2>1) {
        int mid = (high2+low2)/2;
        int z = sa.idx[mid];
        for(int i=0; i<x.length() && z+i<y.length(); i++) {</pre>
            if(x[i]<y[z+i])
               high2=mid; break;
            } else if(x[i]>y[z+i]) {
                low2=mid; break;
        if (high2!=mid && low2!=mid) {
            low2=mid;
    vector<int> ans;
    for (int i=low+1; i<=low2; i++) {
        ans.push_back(sa.idx[i]);
    sort(ans.begin(), ans.end());
    for(int v: ans) {
       cout << v << " ";
   cout << "\n";
return 0;
```

6.2 Knuth-Morris-Pratt

```
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respecitvely.
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
 pi = VI(p.length());
 int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
   while (k >= -1 & & p[k+1] != p[i])
     k = (k == -1) ? -2 : pi[k];
   pi[i] = ++k;
int KMP(string& t, string& p)
 VI pi;
 buildPi(p, pi);
 int k = -1;
  for(int i = 0; i < t.length(); i++) {</pre>
   while (k >= -1 && p[k+1] != t[i])
     k = (k == -1) ? -2 : pi[k];
   k++;
   if(k == p.length() - 1) {
     // p matches t[i-m+1, ..., i]
cout << "matched at index " << i-k << ": ";</pre>
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
 return 0:
int main()
 string a = "AABAACAADAABAABA", b = "AABA";
 KMP(a, b); // expected matches at: 0, 9, 12
 return 0:
```

7 Miscellaneous

7.1 Binary Search

```
// n is size of array, c is value looking for
// sematically equiv to std::lower_bound and std::upper_bound
int lower_bound(int A[], int n, int c) {
    int l = 0, r = n;
    while (l < r) {
        if (A[m] < c) l = m+1; else r=m;
    }
    return l;
}
int upper_bound(int A[], int n, int c) {
    int l = 0, r = n;
    while (l < r) {
        int m = (r-1)/2+1;
        if (A[m] <= c) l = m+1; else r=m;
    }
    return l;
}</pre>
```

7.2 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
    INPUT: a vector of integers
    OUTPUT: a vector containing the longest increasing subsequence
typedef vector<int> VI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
 VPII best:
 VI dad(v.size(), -1);
 for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASING
   PII item = make_pair(v[i], 0);
VPII::iterator it = lower_bound(best.begin(), best.end(), item);
    item.second = i;
#else
    PII item = make_pair(v[i], i);
    VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
   if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().second);
     best.push_back(item);
    } else {
     dad[i] = it == best.begin() ? -1 : prev(it)->second;
      *it = item;
  for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
  reverse(ret.begin(), ret.end());
 return ret:
```

7.3 Median Max/Min Heap

```
#include <bits/stdc++.h>
using namespace std;
```

```
int main() {
        priority_queue<int> maxPQ;
        priority_queue<int, vector<int>, greater<int> > minPQ;
        while(cin >> s) {
               if (s == "#") {
                       int m = minPQ.top(); minPQ.pop();
                        if (minPQ.size() != maxPQ.size()) {
                                minPQ.push(maxPQ.top());
                                maxPQ.pop();
                        cout << m << endl;
                        int c = stoi(s);
                        if(!minPQ.empty() && c > minPQ.top()) {
                                minPQ.push(c);
                                if (minPQ.size() > maxPQ.size() + 1) {
                                        int d = minPQ.top(); minPQ.pop
                                        maxPQ.push(d);
                                maxPQ.push(c);
                                if (maxPQ.size() > minPQ.size()) {
                                       minPQ.push(maxPQ.top());
                                        maxPQ.pop();
```

// Routines for performing computations on dates. In these routines,

7.4 Dates

```
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
  x = jd + 68569;
 n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
 i = (4000 * (x + 1)) / 1461001;

x = 1461 * i / 4 - 31;
  j = 80 * x / 2447;
 d = x - 2447 * j / 80;
 x = j / 11;

m = j + 2 - 12 * x;
 y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
 return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
 int jd = dateToInt (3, 24, 2004);
 int m, d, y;
 intToDate (jd, m, d, y);
  string day = intToDay (jd);
  // expected output:
       2453089
        3/24/2004
  // Wed
  cout << jd << endl
    << m << "/" << d << "/" << y << endl
    << day << endl;
```

7.5 Latitude/longitude

```
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
struct 11
 double r, lat, lon;
struct rect
 double x, y, z;
};
11 convert (rect& P)
 Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
 Q.lat = 180/M_PI*asin(P.z/Q.r);
 Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
 return 0:
rect convert(ll& Q)
 P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
 P.z = Q.r*sin(Q.lat*M_PI/180);
 return P:
int main()
 rect A:
 11 B;
 A.x = -1.0; A.y = 2.0; A.z = -3.0;
 cout << B.r << " " << B.lat << " " << B.lon << endl;
  cout << A.x << " " << A.y << " " << A.z << endl;
```

7.6 Random STL stuff

```
// Example for using stringstreams and next_permutation
int main (void) {
  vector<int> v;
  v.push_back(1);
  v.push_back(2);
  v.push back(3);
  v.push_back(4);
  // Expected output: 1 2 3 4
                        1243
                        4 3 2 1
  do (
   ostringstream oss;    oss << v[0] << " " << v[1] << " " << v[2] << " " << v[3];
    // for input from a string s,
    // istringstream iss(s);
// iss >> variable;
    cout << oss.str() << endl;
  } while (next_permutation (v.begin(), v.end()));
  v.clear();
```

```
v.push_back(1);
v.push_back(2);
v.push_back(1);
v.push_back(1);
v.push_back(3);

// To use unique, first sort numbers. Then call
// unique to place all the unique elements at the beginning
// of the vector, and then use erase to remove the duplicate
// elements.

sort(v.begin(), v.end());
v.erase(unique(v.begin(), v.end()), v.end());

// Expected output: 1 2 3
for (size_t i = 0; i < v.size(); i++)
    cout << v[i] << " ";
cout << endl;</pre>
```

7.7 Longest common subsequence

```
Calculates the length of the longest common subsequence of two vectors
Backtracks to find a single subsequence or all subsequences. Runs in
O(m*n) time except for finding all longest common subsequences, which
may be slow depending on how many there are.
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
void backtrack(VVI& dp, VT& res, VT& A, VT& B, int i, int j)
  if(!i || !j) return;
 if(A[i-1] == B[j-1]) \{ res.push_back(A[i-1]); backtrack(dp, res, A,
       B, i-1, j-1); }
    if(dp[i][j-1] >= dp[i-1][j]) backtrack(dp, res, A, B, i, j-1);
   else backtrack(dp, res, A, B, i-1, j);
void backtrackall(VVI& dp, set<VT>& res, VT& A, VT& B, int i, int j)
 if(!i || !j) { res.insert(VI()); return; }
if(A[i-1] == B[j-1])
    backtrackall(dp, tempres, A, B, i-1, j-1);
    for(set<VT>::iterator it=tempres.begin(); it!=tempres.end(); it++)
     VT temp = *it;
     temp.push_back(A[i-1]);
     res.insert(temp);
  else
    if(dp[i][j-1] >= dp[i-1][j]) backtrackall(dp, res, A, B, i, j-1);
   if(dp[i][j-1] <= dp[i-1][j]) backtrackall(dp, res, A, B, i-1, j);</pre>
VT LCS(VT& A, VT& B)
 int n = A.size(), m = B.size();
  dp.resize(n+1);
  for(int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
  for(int i=1; i<=n; i++)</pre>
    for(int j=1; j<=m; j++)
     if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
     else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
 backtrack(dp, res, A, B, n, m);
```

```
reverse(res.begin(), res.end());
  return res;
set<VT> LCSall(VT& A, VT& B)
  int n = A.size(), m = B.size();
  for (int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
  for (int i=1; i<=n; i++)</pre>
    for(int j=1; j<=m; j++)
      if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
      else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
  backtrackall(dp, res, A, B, n, m);
  return res;
int main()
  int a[] = { 0, 5, 5, 2, 1, 4, 2, 3 }, b[] = { 5, 2, 4, 3, 2, 1, 2,
       1, 3 };
  VI A = VI(a, a+8), B = VI(b, b+9);
  VI C = LCS(A, B);
  for(int i=0; i<C.size(); i++) cout << C[i] << " ";
  cout << endl << endl;
  set <VI> D = LCSall(A, B);
  for(set<VI>::iterator it = D.begin(); it != D.end(); it++)
    for(int i=0; i<(*it).size(); i++) cout << (*it)[i] << " ";</pre>
    cout << endl:
```

8 Language Stuff

8.1 Nifty Tricks

```
// remove duplicated from vector
#define UNIQUE(x) x.erase(unique(x.begin(), x.end()), x.end())
// filters.erase(unique(filters.begin(), filters.end()), filters.end()
// convert string to int
int myint = stoi("123");
// memset
int res[MAX_V][MAX_V];
memset (res, 0, sizeof res);
fill (myvector.begin(), myvector.begin()+4,5);
int myint1 = stoi(str1); // convert string to int
// Convert int to binary string
cout << bitset<32>(val).to_string() << endl;</pre>
// Generate all permutations
sort(nodes.begin(), nodes.end());
do (
  int sum = 0:
  for(int i = 1; i < nodes.size(); i++)</pre>
   sum += __builtin_popcount(nodes[i] & nodes[i-1]);
  best = min(best, sum);
 } while (next_permutation(nodes.begin(), nodes.end()));
// Generate all set of n elements
unsigned next_set_n(unsigned x) {
  unsigned smallest, ripple, new_smallest, ones;
  if(x==0) return 0;
  smallest = (x & -x);
  ripple = x + smallest;
  new_smallest = (ripple & -ripple);
  ones = ((new_smallest/smallest) >> 1 ) - 1;
  return ripple | ones;
```

8.2 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
#define db(x) cerr << #x << "=" << x << endl
#define db2(x, y) cerr << #x << "=" << x << "," << #y << "=" << y <<
#define db3(x, y, z) cerr << \#x << "=" << x << "," << \#y << "=" << y
      << "," << #z << "=" << z << endl
#define LSOne(S) (S & (-S))
// remove duplicated from vector
\#define UNIQUE(x) x.erase(unique(x.begin(), x.end()), x.end())
// Generate all set of n elements
unsigned next_set_n(unsigned x) {
    unsigned smallest, ripple, new_smallest, ones;
    if(x==0) return 0;
    smallest = (x & -x);
   ripple = x + smallest;
   new_smallest = (ripple & -ripple);
   ones = ((new smallest/smallest) >> 1 ) - 1;
   return ripple | ones;
int main()
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
   cout << 100.0/7.0 << endl:
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl:
   cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
   cout.setf(ios::showpos);
cout << 100 << " " << -100 << endl;</pre>
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal cout << hex << 100 << " " << 1000 << " " << 10000 << endl;
    // Convert int to binary string
    cout << bitset<32>(val).to_string() << endl;</pre>
```