INF4820: Algorithms for Artificial Intelligence and Natural Language Processing

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Common Lisp Fundamentals

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Last Week: What is AI?



- ► Since the 1950s: Chatbots, theorem proving, blocks world, expert and dialogue systems, game playing, . . .
- ► Moving target: Whatever requires 'intelligent' decisions, but seems out of reach, technologically, at the time?
- Recently: Conversational user interfaces, self-driving cars, talking robots, AlphaGo.
- ▶ But also (fuzzily) business intelligence, (big) data analytics, . . .
- ightarrow Toolkit of ('clever') methods for representation and problem solving.





Why Common Lisp?



Eric S. Raymond (2001), How to Become a Hacker.

Lisp is worth learning for the profound enlightenment experience you will have when you finally get it; that experience will make you a better programmer for the rest of your days, even if you never actually use Lisp itself a lot.

- High-level and efficient language with especially strong support for symbolic and functional programming.
- ► Rich language: multitude of built-in data types and operations.
- ► Easy to learn:
 - extremely simple syntax,
 - straightforward semantics.
- ► ANSI-standardized and stable.
- ► Incremental and interactive development.

Lisp



- Conceived in the late 1950s by John McCarthy—one of the founding fathers of AI.
- Originally intended as a mathematical formalism.
- A family of high-level languages.
- Several dialects, e.g. Scheme, Clojure, Emacs Lisp, and Common Lisp.
- Although a multi-paradigm language, functional style prevalent.



Basic Common Lisp in a Couple of Minutes



- ► Testing a few expressions at the REPL;
- ► the read—eval—print loop.
- ► (= the interactive Lisp-environment)
- '?' represents the REPL prompt and'→' what an expression evaluates to.
- Atomic data types like numbers, booleans, and strings are self evaluating.
- Symbols evaluate to whatever value they are bound to.

Examples

- ? "this is a string"
- \rightarrow "this is a string"
- ? 42
- \rightarrow 42
- ? t
- \rightarrow t
- ? nil
- ightarrow nil
- ? pi
- \rightarrow 3.141592653589793d0
 - ? foo
- \rightarrow error; unbound

A Note on Terminology



- ► Lisp manipulates so-called *symbolic expressions*.
- AKA s-expressions or sexps.
- Two fundamental types of sexps;
 - 1. atoms (e.g., nil, t, numbers, strings, symbols)
 - 2. lists containing other sexps.
- ► Sexps are used to represent *both* data and code.

Function Calls



- "Parenthesized prefix notation"
- ► First element (prefix) = operator (i.e. the procedure or function).
- The rest of the list is the operands (i.e. the arguments or parameters).
- Use nesting (of lists) to build compound expressions.
- Expressions can span multiple lines; indentation for readability.

Examples

- ? (+ 1 2)
- \rightarrow 3
- ? (+ 1 2 10 7 5)
- $\rightarrow 25$
- ? (/ (+ 10 20) 2)
- \rightarrow 15
- ? (* (+ 42 58) (- (/ 8 2) 2))
- → 200

The Syntax and Semantics of CL



- ? (expt (- 8 4) 2)
- \rightarrow 16
 - ▶ You now know (almost) all there is to know about (the rules of) CL.
 - ► The first element of a list names a function that is invoked with the values of all remaining elements as its arguments.
 - ► A few exceptions, called special forms, with their own evaluation rules.

Creating our own functions



► The special form defun associates a function definition with a symbol:

General form

 $(defun name (parameter_1 ... parameter_n) body)$

Example

- ? (defun average (x y) (/ (+ x y) 2))
- ? (average 10 20)
- \rightarrow 15

The 'Hello World!' of Functional Programming



- ► Classic example: the factorial function.
- ► A recursive procedure; calls itself, directly or indirectly.
- May seem circular, but is well-defined as long as there's a base case terminating the recursion.
- ► For comparison: a non-recursive implementation (in Python).

```
n! = \begin{cases} 1 & \text{if } n = 0\\ n \times (n-1)! & \text{if } n > 0 \end{cases}
```

```
(defun ! (n)

(if (= n 0)

1

(* n (! (- n 1)))))
```

```
def fac(n):
    r = 1
    while (n > 0):
        r = r * n
        n = n - 1
    return r
```

A Special Case of Recursion: Tail Recursion



- ► A more efficient way to define n! recursively.
- Use a helper procedure with an accumulator variable to collect the product along the way.
- The recursive call is in tail position;

- ▶ no work remains to be done in the calling function.
- ▶ Once we reach the base case, the return value is ready.
- ▶ Most CL compilers do *tail call optimization* (TCO), so that the recursion is executed as an iterative loop.
- ► (The next lecture will cover CL's built-in loop construct.)

Tracing the processes



Recursive

```
(defun ! (n)
   (if (= n 0)
        (* n (! (- n 1)))))
? (! 7)
\Rightarrow (* 7 (! 6))
\Rightarrow (* 7 (* 6 (! 5)))
\Rightarrow (* 7 (* 6 (* 5 (! 4))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (! 3)))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (* 3 (! 2))))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (* 3 (* 2 (! 1)))))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (* 3 (* 2 1))))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (* 3 2)))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 6))))
\Rightarrow (* 7 (* 6 (* 5 24)))
\Rightarrow (* 7 (* 6 120))
\Rightarrow (* 7 720)
\rightarrow 5040
```

Tail-Recursive

```
(defun! (n)
  (!-aux 1 1 n))
(defun !-aux (r i n)
  (if (> i n)
        (!-aux (* r i)
                  (+ i 1)
                  n)))
? (! 7)
\Rightarrow (!-aux 1 1 7)
\Rightarrow (!-aux 1 2 7)
\Rightarrow (!-aux 2 3 7)
\Rightarrow (!-aux 6 4 7)
\Rightarrow (!-aux 24 5 7)
\Rightarrow (!-aux 120 6 7)
\Rightarrow (!-aux 720 7 7)
\Rightarrow (!-aux 5040 8 7)
\rightarrow 5040
```

The quote Operator



- ► A special form making expressions self-evaluating.
- ► The quote operator (or simply ',') suppresses evaluation.
- ? pi→ 3.141592653589793d0
- ? (quote pi) \rightarrow pi
- ? 'pi \rightarrow pi
- ? foobar \rightarrow error; unbound variable
- ? 'foobar \rightarrow foobar
- ? (* 2 pi) → 6.283185307179586d0
- ? $(* 2 pi) \rightarrow (* 2 pi)$
- ?() \rightarrow error; missing procedure
- ? '() \rightarrow ()

Both Code and Data are S-Expressions



- ▶ We've mentioned how sexps are used to represent *both* data and code.
- ► Note the double role of lists:
- ► Lists are function calls:

```
? (* 10 (+ 2 3)) \rightarrow 50
? (bar 1 2) \rightarrow error; function bar undefined
```

► But, lists can also be data:

```
? '(foo bar) \rightarrow (foo bar)
? (list 'foo 'bar) \rightarrow (foo bar)
```

LISP = LISt Processing



► cons builds up new lists; first and rest destructure them.

```
? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)

? (cons 0 '(1 2 3)) \rightarrow (0 1 2 3)

? (first '(1 2 3)) \rightarrow 1

? (rest '(1 2 3)) \rightarrow (2 3)

? (first (rest '(1 2 3))) \rightarrow 2

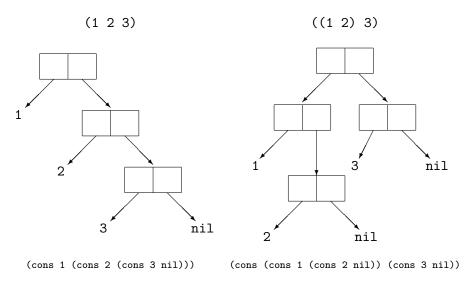
? (rest (rest (rest '(1 2 3)))) \rightarrow nil
```

► Many additional list operations (derivable from the above), e.g.

```
? (list 1 2 3) → (1 2 3)
? (append '(1 2) '(3) '(4 5 6)) → (1 2 3 4 5 6)
? (length '(1 2 3)) → 3
? (reverse '(1 2 3)) → (3 2 1)
? (nth 2 '(1 2 3)) → 3
? (last '(1 2 3)) → (3) Wait, why not 3?
```

Lists are Really Chained 'cons' Cells





Assigning Values: 'Generalized Variables'



► defparameter declares a 'global variable' and assigns a value:

```
? (defparameter *foo* 42) \rightarrow *F00*
```

- ? *foo* \rightarrow 42
- ▶ setf provides a uniform way of assigning values to variables.
- ► General form:

```
(setf place value)
```

... where place can either be a variable named by a symbol or some other storage location:

```
? (setf *foo* (+ *foo* 1))
? *foo* → 43
? (setf *foo* '(2 2 3))
? (setf (first *foo*) 1)
```

Some Other Macros for Assignment



Example	Type of x	Effect
(incf x y)	number	(setf x (+ x y))
(incf x)	number	(incf x 1)
(decf x y)	number	(setf x (-x y))
(decf x)	number	(decf x 1)
(push y x)	list	(setf x (cons y x))
(pop x)	list	<pre>(let ((y (first x))) (setf x (rest x)) y)</pre>
(pushnew y x)	list	<pre>(if (member y x) x (push y x))</pre>

Shall we jointly write our own push and pop?

Local Variables



- ► Sometimes we want to store intermediate results.
- ► let and let* create temporary value bindings for symbols.

▶ Bindings valid only in the body of let.

? baz → error; unbound variable

- ▶ Previously existing bindings are *shadowed* within the lexical scope.
- ▶ let* is like let but binds sequentially.

Predicates



- ► A *predicate* tests some condition.
- ► Evaluates to a boolean truth value:
 - ▶ nil (the empty list) means false.
 - ► Anything non-nil (including t) means *true*.

▶ Plethora of equality tests: eq, eq1, equal, and equalp.

Equality for One and All



- eq tests object identity; not applicable for numbers or characters.
- ▶ eql is like eq, but well-defined on numbers and characters.
- equal tests structural equivalence (recursively for lists and strings).
- equalp is like equal but insensitive to case and numeric type.

```
? (eq (list 1 2 3) '(1 2 3)) \rightarrow nil

? (equal (list 1 2 3) '(1 2 3)) \rightarrow t

? (eq 42 42) \rightarrow ? [implementation-dependent]

? (eq1 42 42) \rightarrow t

? (eq1 42 42.0) \rightarrow nil

? (equalp 42 42.0) \rightarrow t

? (equal "foo" "foo") \rightarrow t

? (equalp "F00" "foo") \rightarrow t
```

► Also many type-specialized tests like =, string=, etc.

Conditional Evaluation



Examples

```
? (defparameter foo 42)
 (if (numberp foo)
      "number"
      "something else")
→ "number"
   (cond ((< foo 3) "less")
         ((> foo 3) "more")
         (else "equal"))
  "more"
```

General Form

```
(if ⟨predicate⟩
    ⟨then clause⟩
    ⟨else clause⟩)

(cond (⟨predicate₁⟩ ⟨clause₁⟩)
    (⟨predicate₂⟩ ⟨clause₂⟩)
    (⟨predicate₁⟩ ⟨clause₁⟩)
    (t ⟨default clause⟩))
```

Rewind: A Note on Symbol Semantics



- ► Symbols can have values as functions and variables at the same time.
- ▶ #' (sharp-quote) gives us the function object bound to a symbol.

```
? (defun foo (x)
          (* x 1000))
```

? (defparameter foo 42) \rightarrow 2

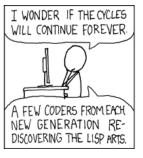
```
? (foo foo) \rightarrow 42000
```

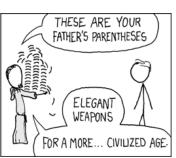
- ? foo $\rightarrow 42$
- ? #'foo \rightarrow #<Interpreted Function FOO>
- ? (funcall #'foo foo) \rightarrow 42000
- #' and funcall (as well as apply) are useful when passing around functions as arguments.

In Conclusion



LISP IS OVER HALF A
CENTURY OLD AND IT
STILL HAS THIS PERFECT,
TIMELESS AIR ABOUT IT.





http://xkcd.com/297/

Next Week



More Common Lisp.

- ► More on argument lists (optional arguments, keywords, defaults).
- ▶ More data types: Hash-tables, a-lists, arrays, sequences, and structures
- ► More higher-order functions.
- ► Iteration (loop) and mapping.