Introduction to Operating Systems

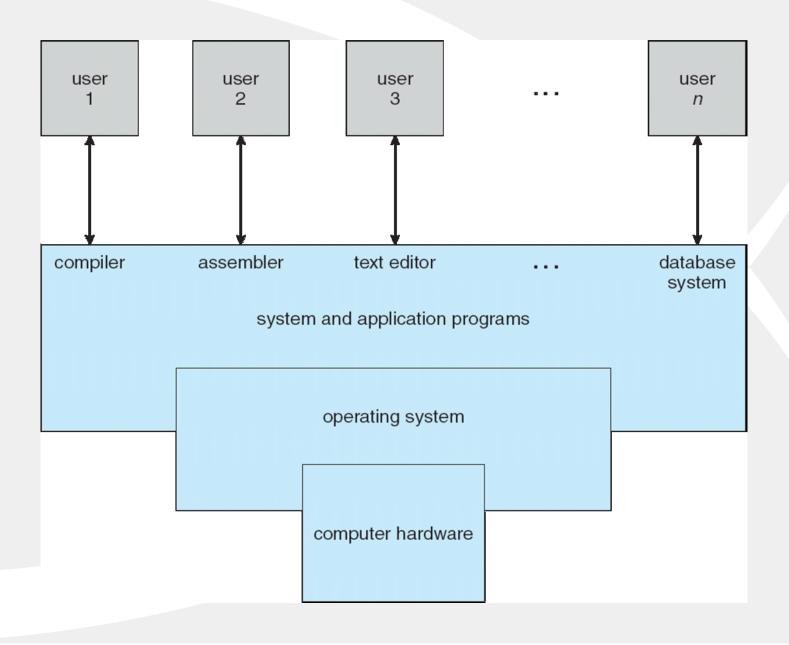
What is an Operating System?

- A program that acts as an intermediary between a user of a computer and the computer hardware
- Operating system goals:
 - Execute user programs and make solving user problems easier
 - Make the computer system convenient to use
 - Use the computer hardware in an efficient manner

Computer System Structure

- Computer system can be divided into four components:
 - Hardware provides basic computing resources
 - CPU, memory, I/O devices
 - Operating system
 - Controls and coordinates use of hardware among various applications and users
 - Application programs define the ways in which the system resources are used to solve the computing problems of the users
 - Word processors, compilers, web browsers, database systems, video games
 - Users
 - People, machines, other computers

Four Components of a Computer System



What Operating Systems Do?

- Users want convenience, ease of use and good performance
 - Don't care about resource utilization
- But shared computer such as mainframe or minicomputer must keep all users happy
- Users of dedicate systems such as workstations have dedicated resources but frequently use shared resources from servers
- Handheld computers are resource poor, optimized for usability and battery life
- Some computers have little or no user interface, such as embedded computers in devices and automobiles could have real-time requirements

So What are the Major OS Functions?

- Control access and provide interfaces
 - To the OS and devices attached to the system
 - Provide interfaces for human-machine and machine-machine transactions
- Manage resources (Resource Allocator)
 - Mediate resource usage among different tasks
 - Implement policies

So What are the Major OS Functions?

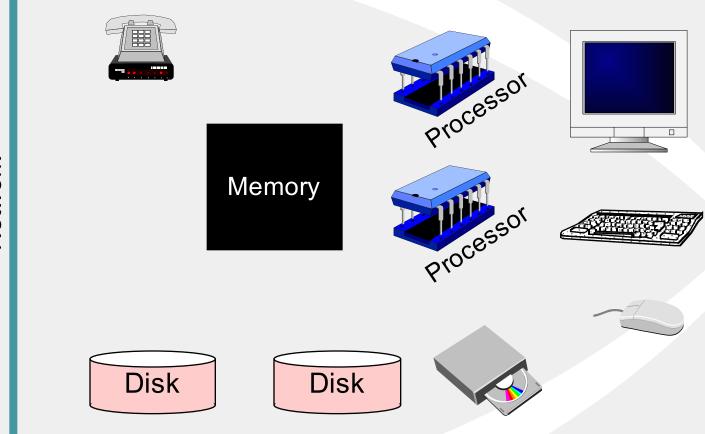
Provide abstractions

- Hide the peculiarities of the hardware.
- Example: device independent I/O

Consume resources*

OS runs on the system that is being managed... so it is going to consume resources!

Hardware



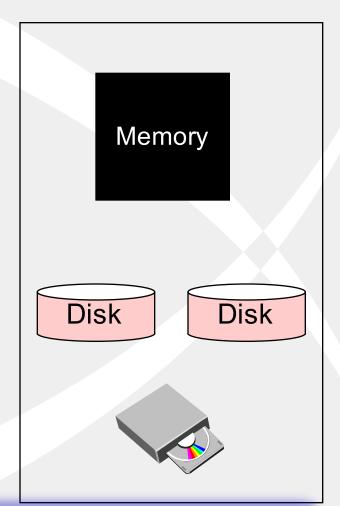
Operating System needs to give an easy-to-use abstraction over a complicated collection of hardware components

Network

Files

How does OS manage persistent storage (e.g., disks)?



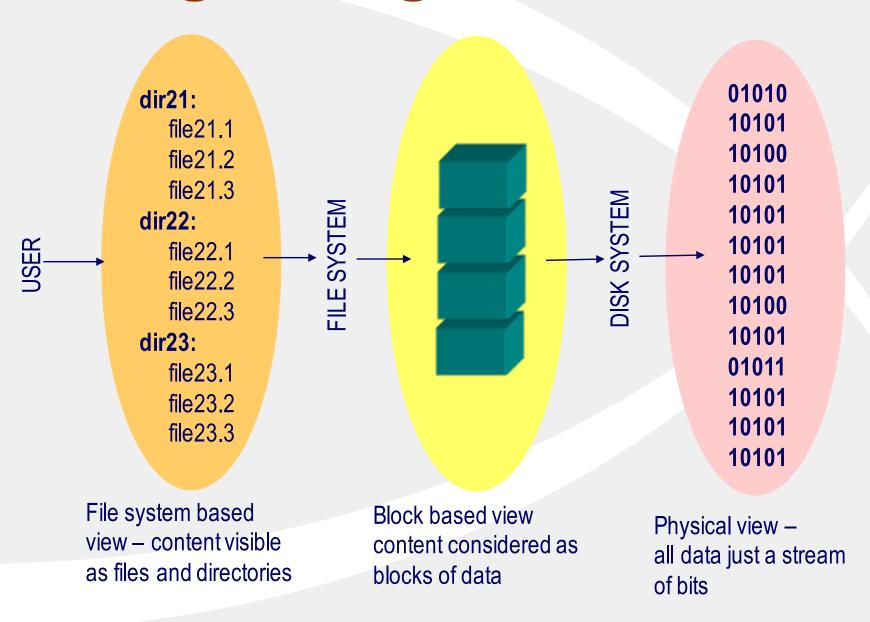


OSes use the File concept to give a familiar way of managing the persistent storage

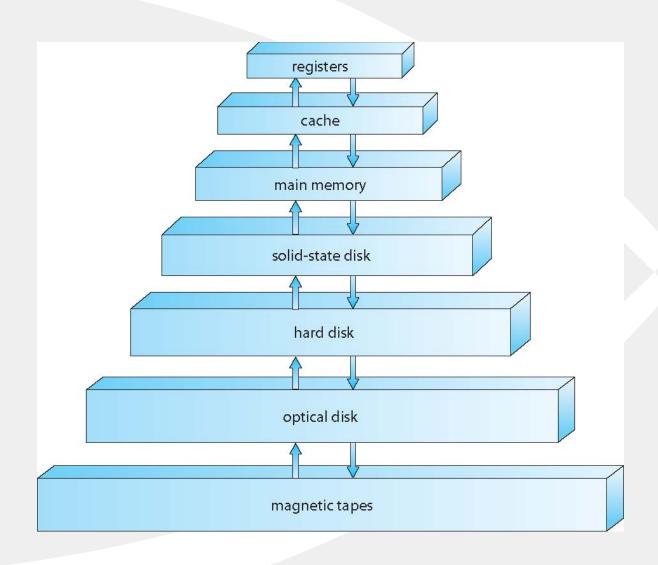
Storage Management: In Words

- OS provides uniform, logical view of information storage
 - Abstracts physical properties to logical storage unit file
 - Each medium is controlled by device (i.e., disk drive, tape drive)
 - Varying properties include access speed, capacity, datatransfer rate, access method (sequential or random)
- File-System management
 - Files usually organized into directories
 - Access control on most systems to determine who can access what
 - OS activities include
 - Creating and deleting files and directories
 - Primitives to manipulate files and directories
 - Mapping files onto secondary storage
 - Backup files onto stable (non-volatile) storage media

Storage Management: In Pictures



Storage-Device Hierarchy



Storage Hierarchy

- Storage systems organized in hierarchy
 - Speed
 - Cost
 - Volatility
- Caching copying information into faster storage system; main memory can be viewed as a cache for secondary storage
- Device Driver for each device controller to manage I/O
 - Provides uniform interface between controller and kernel

Storage Structure

- Main memory only large storage media that the CPU can access directly
 - Random access
 - Typically volatile
- Secondary storage extension of main memory that provides large nonvolatile storage capacity
- Hard disks rigid metal or glass platters covered with magnetic recording material
 - Disk surface is logically divided into tracks, which are subdivided into sectors
 - The disk controller determines the logical interaction between the device and the computer
- Solid-state disks faster than hard disks, nonvolatile
 - Various technologies
 - Becoming more popular

Programs

Memory How to load a program? Data Disk Disk Code

Programs are made up of Code & Data Need to load both into memory and make the Code run

Process Management

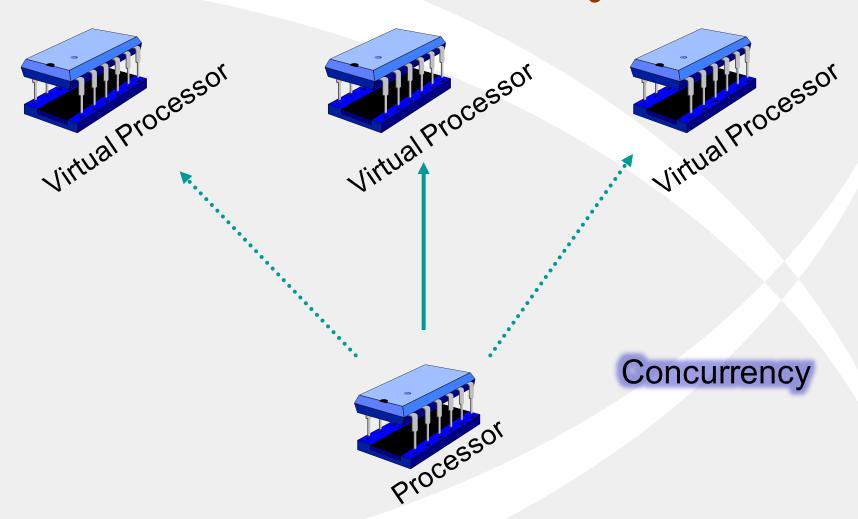
- A process is a program in execution. It is a unit of work within the system. Program is a passive entity, process is an active entity.
- Process needs resources to accomplish its task
 - CPU, memory, I/O, files
 - Initialization data
- Process termination requires reclaim of any reusable resources
- Single-threaded process has one program counter specifying location of next instruction to execute
 - Process executes instructions sequentially, one at a time, until completion
- Multi-threaded process has one program counter per thread
- Typically system has many processes, some user, some operating system running concurrently on one or more CPUs
 - Concurrency by multiplexing the CPUs among the processes / threads

Process Management Activities

The operating system is responsible for the following activities in connection with process management:

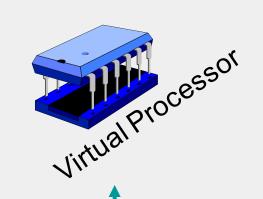
- Creating and deleting both user and system processes
- Suspending and resuming processes
- Providing mechanisms for process synchronization
- Providing mechanisms for process communication
- Providing mechanisms for deadlock handling

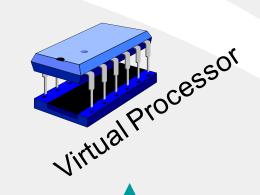
Concurrency

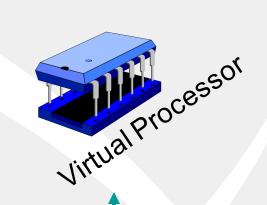


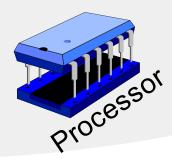
With multiple processes, we can think of each process getting a virtual processor

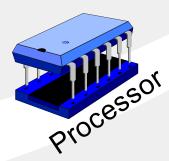
Parallelism

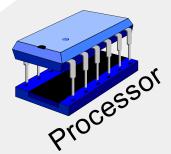












Parallelism

Memory Sharing

Program 1

Program 2

Program 3

Operating System

Memory

OS needs to share the memory among all the processes that are active so that a process does not corrupt memory belonging to another process

Memory Management

- To execute a program all (or part) of the instructions must be in memory
- All (or part) of the data that is needed by the program must be in memory.
- Memory management determines what is in memory and when
 - Optimizing CPU utilization and computer response to users
- Memory management activities
 - Keeping track of which parts of memory are currently being used and by whom
 - Deciding which processes (or parts thereof) and data to move into and out of memory
 - Allocating and deallocating memory space as needed

Operating System Definition

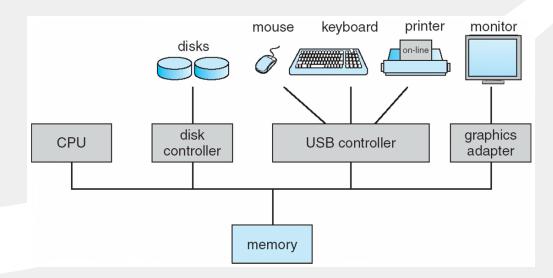
- No universally accepted definition
- "Everything a vendor ships when you order an operating system" is a good approximation
 - But varies wildly
- "The one program running at all times on the computer" is the kernel.
- Everything else is either
 - a system program (ships with the operating system), or
 - an application program.

Computer Startup

- bootstrap program is loaded at power-up or reboot
 - Typically stored in ROM or EPROM, generally known as firmware
 - Initializes all aspects of system
 - Loads operating system kernel and starts execution

Computer System Org.

- Computer-system operation
 - One or more CPUs, device controllers connect through common bus providing access to shared memory
 - Concurrent execution of CPUs and devices competing for memory cycles



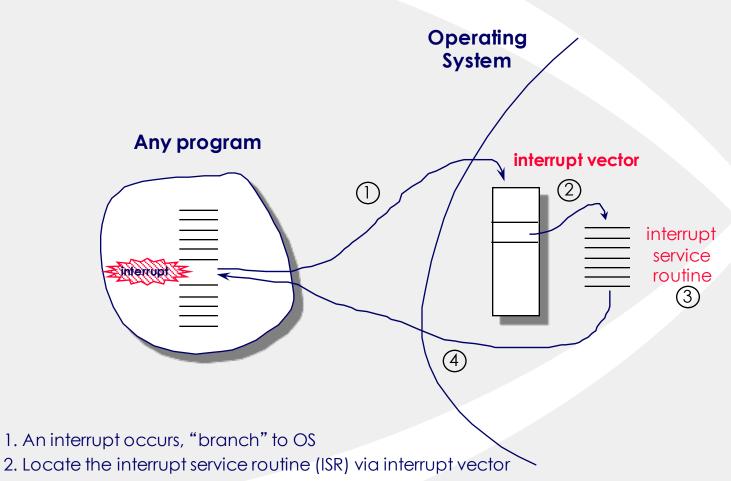
Computer-System Operation

- I/O devices and the CPU can execute concurrently
- Each device controller is in charge of a particular device type
- Each device controller has a local buffer
- CPU moves data from/to main memory to/from local buffers
- I/O is from the device to local buffer of controller
- Device controller informs CPU that it has finished its operation by causing an interrupt

Common Functions of Interrupts

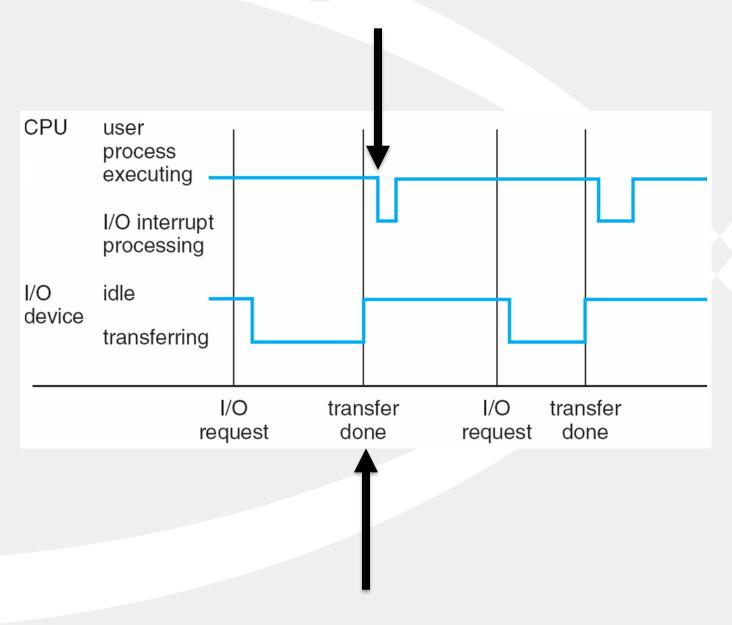
- Interrupt transfers control to the interrupt service routine generally, through the interrupt vector, which contains the addresses of all the service routines
- Interrupt architecture must save the address of the interrupted instruction
- A trap or exception is a software-generated interrupt caused either by an error or a user request
- An operating system is interrupt driven

Interrupt Servicing: In Picture

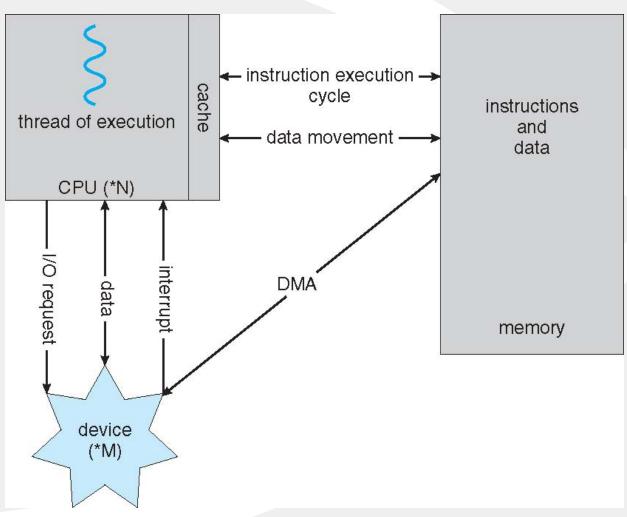


- 3. Execute the ISR
- 4. Return to interrupted program

Interrupt Timeline



How a Modern Computer Works

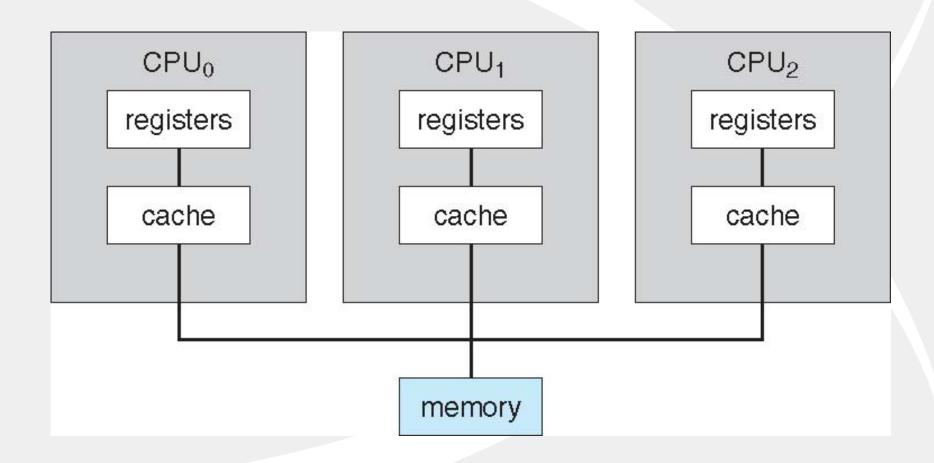


A von Neumann architecture

Computer-System Arch.

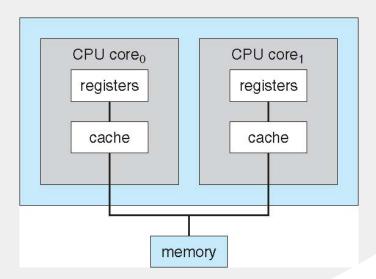
- Started migrating from single processors to multiprocessors (multi-cores are common)
- Special purpose processors are also heavily used (e.g., GPUs)
 - Also known as parallel systems, tightly-coupled systems
 - Advantages include:
 - 1. Increased throughput
 - 2. Economy of scale
 - 3. Increased reliability graceful degradation or fault tolerance
 - Two types:
 - Asymmetric Multiprocessing each processor is assigned a specie task.
 - Symmetric Multiprocessing each processor performs all tasks

Symmetric Multiprocessing Architecture



A Dual-Core Design

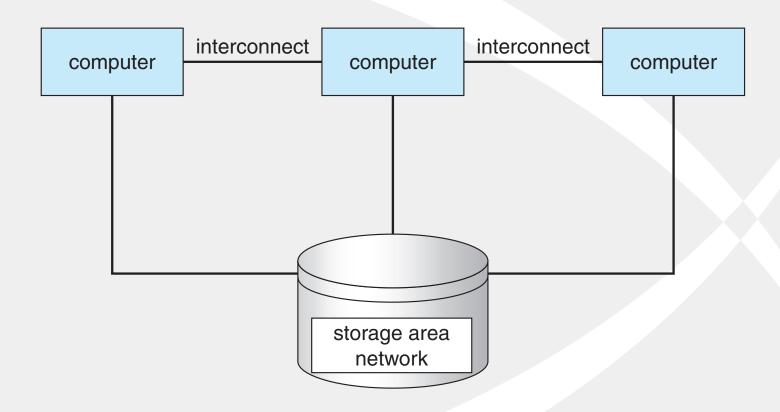
- Multi-chip and multicore
- Systems containing all chips
 - Chassis containing multiple separate systems



Clustered Systems

- Like multiprocessor systems, but multiple systems working together
 - Usually sharing storage via a storage-area network (SAN)
 - Provides a high-availability service which survives failures
 - Asymmetric clustering has one machine in hotstandby mode
 - Symmetric clustering has multiple nodes running applications, monitoring each other
 - Some clusters are for high-performance computing (HPC)
 - Applications must be written to use parallelization
 - Some have distributed lock manager (DLM) to avoid conflicting operations

Clustered Systems



Operating System Structure

- Multiprogramming (Batch system) needed for efficiency
 - Single user cannot keep CPU and I/O devices busy at all times
 - Multiprogramming organizes jobs (code and data) so CPU always has one to execute
 - A subset of total jobs in system is kept in memory
 - One job selected and run via job scheduling
 - When it has to wait (for I/O for example), OS switches to another job

Operating System Structure

- Timesharing (multitasking) is logical extension in which CPU switches jobs so frequently that users can interact with each job while it is running, creating interactive computing
 - Response time should be < 1 second</p>
 - Each user has at least one program executing in memory ⇒ process
 - ♦ If several jobs ready to run at the same time ⇒
 CPU scheduling
 - If processes don't fit in memory, swapping moves them in and out to run
 - Virtual memory allows execution of processes not completely in memory

OS Operations

- Interrupt driven (hardware and software)
 - Hardware interrupt by one of the devices
 - Software interrupt (exception or trap):
 - Software error (e.g., division by zero)
 - Request for operating system service
 - Other process problems include infinite loop, processes modifying each other or the operating system

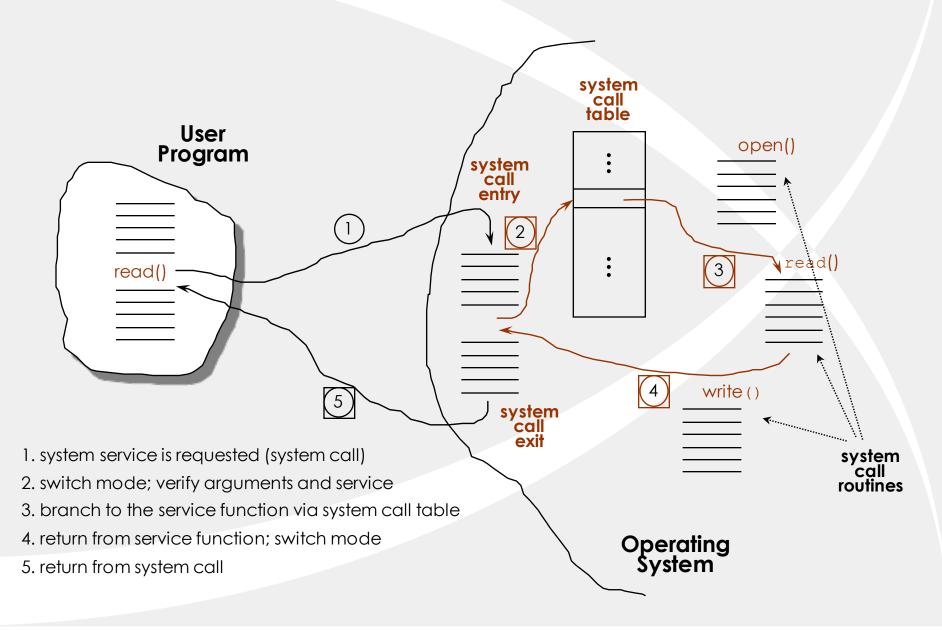
OS Operations (cont.)

- Dual-mode operation allows OS to protect itself and other system components
 - User mode and kernel mode
 - Mode bit provided by hardware
 - Provides ability to distinguish when system is running user code or kernel code
 - Some instructions designated as privileged, only executable in kernel mode
 - System call changes mode to kernel, return from call resets it to user
- Increasingly CPUs support multi-mode operations
 - i.e. virtual machine manager (VMM) mode for guest VMs

User to Kernel Mode Transition

- Timer to prevent infinite loop / process hogging resources
 - Timer is set to interrupt the computer after some time period
 - Keep a counter that is decremented by the physical clock.
 - Operating system set the counter (privileged instruction)
 - When counter zero generate an interrupt
 - Set up before scheduling process to regain control or terminate program that exceeds allotted time

How system call is processed?



Protection and Security

- Protection any mechanism for controlling access of processes or users to resources defined by the OS
- Security defense of the system against internal and external attacks
 - Huge range, including denial-of-service, worms, viruses, identity theft, theft of service
- Systems generally first distinguish among users, to determine who can do what
 - User identities (user IDs, security IDs) include name and associated number, one per user
 - User ID then associated with all files, processes of that user to determine access control
 - Group identifier (group ID) allows set of users to be defined and controls managed, then also associated with each process, file
 - Privilege escalation allows user to change to effective ID with more rights

Computing Environments - Traditional

- Stand-alone general purpose machines
- But blurred as most systems interconnect with others (i.e., the Internet)
- Portals provide web access to internal systems
- Network computers (thin clients) are like Web terminals
- Mobile computers interconnect via wireless networks
- Networking becoming ubiquitous even home systems use firewalls to protect home computers from Internet attacks

Computing Environments – Distributed

- Distributed computing
 - Collection of separate, possibly heterogeneous, systems networked together
 - Network is a communications path, TCP/IP most common
 - Local Area Network (LAN)
 - Wide Area Network (WAN)
 - Metropolitan Area Network (MAN)
 - Personal Area Network (PAN)
 - Network Operating System provides features between systems across network
 - Communication scheme allows systems to exchange messages
 - Illusion of a single system

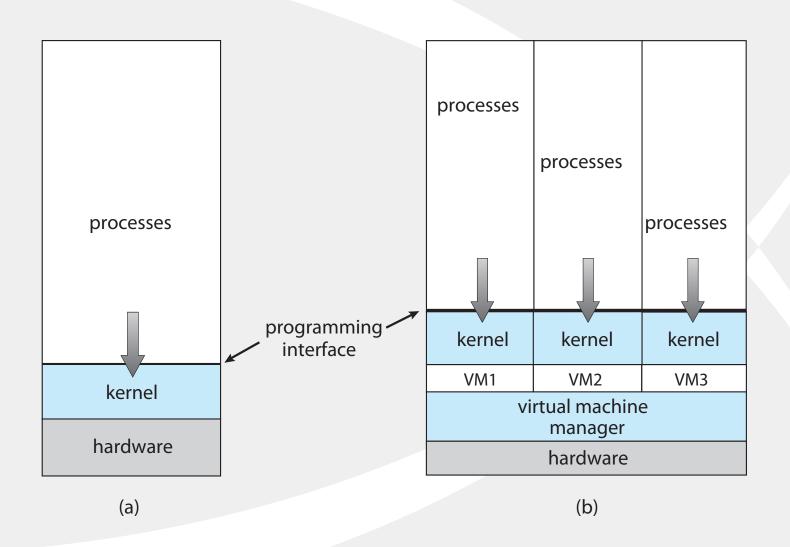
Computing Environments - Virtualization

- Allows operating systems to run applications within other OSes
 - Vast and growing industry
- **Emulation** used when source CPU type different from target type (i.e. PowerPC to Intel x86)
 - Generally slowest method
 - When computer language not compiled to native code – Interpretation
- Virtualization OS natively compiled for CPU, running guest OSes also natively compiled
 - Consider VMware running WinXP guests, each running applications, all on native WinXP host OS
 - VMM (virtual machine Manager) provides virtualization services

Computing Environments - Virtualization

- Use cases involve laptops and desktops running multiple OSes for exploration or compatibility
 - Apple laptop running Mac OS X host, Windows as a guest
 - Developing apps for multiple OSes without having multiple systems
 - QA testing applications without having multiple systems
 - Executing and managing compute environments within data centers
- VMM can run natively, in which case they are also the host
 - There is no general purpose host then (VMware ESX and Citrix XenServer)

Computing Environments - Virtualization



Computing Environments – Cloud Computing

- Delivers computing, storage, even apps as a service across a network
- Logical extension of virtualization because it uses virtualization as the base for it functionality.
 - Amazon EC2 has thousands of servers, millions of virtual machines, petabytes of storage available across the Internet, pay based on usage

Many types

- Public cloud available via Internet to anyone willing to pay
- Private cloud run by a company for the company's own use
- Hybrid cloud includes both public and private cloud components
- Software as a Service (SaaS) one or more applications available via the Internet (i.e., word processor)
- Platform as a Service (PaaS) software stack ready for application use via the Internet (i.e., a database server)
- Infrastructure as a Service (laaS) servers or storage available over Internet (i.e., storage available for backup use)

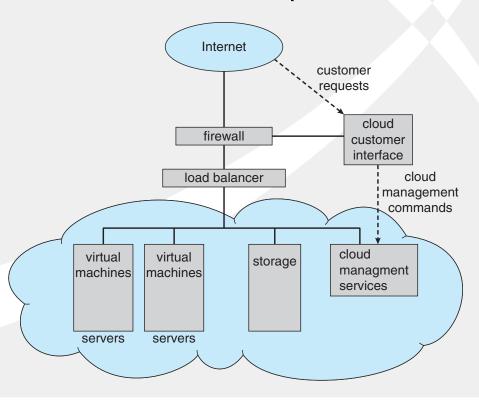
Computing Environments – Cloud Computing

Cloud computing environments composed of traditional OSes, plus VMMs, plus cloud management tools

Internet connectivity requires security like firewalls

Load balancers spread traffic across multiple

applications



Computing Environments – Real-Time Embedded Systems

- Real-time embedded systems most prevalent form of computers
 - Vary considerable, special purpose, limited purpose OS, real-time OS
 - Use expanding
- Many other special computing environments as well
 - Some have OSes, some perform tasks without an OS
- Real-time OS has well-defined fixed time constraints
 - Processing must be done within constraint
 - Correct operation only if constraints met