EE270 Large scale matrix computation, optimization and learning

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Randomized Linear Algebra and Optimization Lecture 18: Randomized Kernel Approximations, Effective Dimension and Nystrom Method

Recap: Low-rank matrix approximations

- Singular Value Decomposition (SVD)
- \triangleright $A = U\Sigma V^T$
- ▶ takes $O(nd^2)$ time for $A \in R^{n \times d}$
- ▶ best rank-k approximation is $A_k := U_k \Sigma_k V_k^T = \sum_{i=1}^k \sigma_i u_i v_i^T$
- ► $||A A_k||_2 \le \sigma_{k+1}$

Recap: Randomized low-rank matrix approximations

idea: sample some rows/sketch $A \in \mathbb{R}^{n \times d}$ to get $C \in \mathbb{R}^{n \times m}$

- ightharpoonup C = AS where $S \in \mathbb{R}^{d \times m}$ is a sampling/sketching matrix
- we have $AA^T \approx CC^T$ then consider the best approximation of Ain the range of C = AS

$$\min_{X} \|CX - A\|_{F}$$

- also called CX decomposition
- $ightharpoonup ilde{A}_m := CX^* = CC^{\dagger}A$ is a randomized rank-m approximation

$$(AS)(AS)^{\dagger}_{A} \approx A$$



Recap: Randomized Singular Value Decomposition

CX decomposition provides the approximation

$$(AS)(AS)^{\dagger}A \approx A$$

- ightharpoonup calculate QR decomposition of AS = QR
- ► then $QQ^TA \approx A$, i.e., Q approximates the range space of A, T calculate the SVD $Q^TA = U\Sigma V^T \implies 8.9 A = 9.11 ZV = 0.21$ approximate SVD of A is $A \approx (QU)\Sigma V^T$

Analysis of Randomized Low Rank Approximations

ightharpoonup CX decomposition: form sketch AS, and find the best approximation of A in the range of AS

$$X^* = \arg\min_{X} ||ASX - A||_F^2 = (AS)^{\dagger}A$$

- ▶ approximation $ASX^* = (AS)(AS)^{\dagger}A \approx A$
- ightharpoonup yields randomized SVD : AS = QR and $Q^TA = U\Sigma V^T$
- Let $A = U\Sigma V^T$ and $A_k = \sum_{i=1}^k \sigma_k u_k v_k^T$, i.e., best rank-k approximation of A
- note that

$$||AS(AS)^{\dagger}A - A||_{F}^{2} \leq ||AS(A_{k}S)^{\dagger}A_{k} - A||_{F}^{2} + ||As(SA)^{\dagger}A_{k} - A||_{F}^{2} + ||As(SA)^{\dagger}A_{k} - A^{\dagger}||_{F}^{2}$$

$$= ||A_{k}^{T}(S^{T}A_{k}^{T})^{\dagger}S^{T}A^{T} - A^{T}||_{F}^{2}$$

$$= ||A_{k}^{T}(S^{T}A_{k}^{T})^{\dagger}S^{T}A^{T} - A^{T}||_{F}^{2}$$

Analysis of Randomized Low Rank Approximations

approximation error

$$||AS \underbrace{(AS)^{\dagger}A}_{X^*} - A||_F^2 \le ||AS(A_kS)^{\dagger}A_k - A||_F^2$$

$$= ||A_k^T (S^T A_k^T)^{\dagger} S^T A^T - A^T ||_F^2$$

$$= ||A_k^T \tilde{X} - A^T ||_F^2$$

where

$$\tilde{X} := \arg\min_{\mathbf{v}} \|S^T A_k^T X - S^T A^T\|_F^2$$

Analysis of Randomized Low Rank Approximations

approximation error

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where

$$\tilde{X} := \arg\min_{X} \|S^T A_k^T X - S^T A^T\|_F^2$$

▶ identical to sketching the Generalized Least Squares problem

$$\min_{\mathbf{X}} \|A_k^T \mathbf{X} - A^T\|_F^2$$

Generalized Least Squares Problems

$$\min_{X} \|AX - B\|_F^2$$

► Least Squares problem with multiple right-hand-sides

$$B = [b_1, ..., b_r]$$

$$X = [x_1, ..., x_r]$$

$$\min_{x_1,...,x_r} \sum_{i=1}^r ||Ax_i - b_i||_2^2$$

optimal solution

$$X^* = [x_1^*, ..., x_r^*]$$

= $[A^{\dagger}b_1, ..., A^{\dagger}b_r]$
= $A^{\dagger}B$

Left Sketching Generalized Least Squares Problems

original problem

$$X^* := \arg\min_{X} \|AX - B\|_F^2$$

Form sketches of the data SA and SB, e.g., uniform row sampling, weighted sampling, Gaussian, ± 1 i.i.d, CountSketch, FJLT...

$$\hat{X} := \arg\min_{X} \|SAX - SB\|_F^2$$

$$\hat{X}_i = \arg\min_{x_i} \|SAx_i - Sb_i\|_2^2$$

$$= (SA)^{\dagger}(Sb_i)$$

left-sketch applied to simple Least Squares problem $\min_{x_i} ||Ax_i - b_i||_2^2$



Recall Gaussian Sketch Analysis

Let $A \in \mathbb{R}^{n \times d}$, $S \in \mathbb{R}^{m \times n}$ be i.i.d. Gaussian

$$x^* := \arg\min_{x \in \mathbb{R}^d} \underbrace{\|Ax - b\|_2^2}_{f(x)}$$
 and $\tilde{x} = \arg\min_{x \in \mathbb{R}^d} \|SAx - Sb\|_2^2$

Conditioned on the matrix SA

$$A(\tilde{x} - x^*) \sim N\left(0, \frac{f(x^*)}{m}A(A^TS^TSA)^{-1}A\right)$$

Recall Gaussian Sketch Analysis

Let $A \in \mathbb{R}^{n \times d}$, $S \in \mathbb{R}^{m \times n}$ be i.i.d. Gaussian

$$x^* := \arg\min_{x \in \mathbb{R}^d} \underbrace{\|Ax - b\|_2^2}_{f(x)}$$
 and $\tilde{x} = \arg\min_{x \in \mathbb{R}^d} \|SAx - Sb\|_2^2$

Conditioned on the matrix SA

$$A(\tilde{x} - x^*) \sim N\left(0, \frac{f(x^*)}{m}A(A^TS^TSA)^{-1}A\right)$$

taking expectation over SA, and using $\mathbb{E}[(A^TS^TSA)^{-1}] = (A^TA)^{-1} \frac{m}{m-d-1}$ we get

$$\mathbb{E}||A(\tilde{x} - x^*)||_2^2 = \frac{f(x^*)}{m - d - 1} tr A(A^T A)^{-1} A$$
$$= f(x^*) \frac{\operatorname{rank}(A)}{m - d - 1} = f(x^*) \frac{d}{m - d - 1}$$

Left Sketching Generalized Least Squares Problems

original problem and left-sketch

$$X^* := \arg\min_{X} \|AX - B\|_F^2$$
 and $\hat{X} := \arg\min_{X} \|SAX - SB\|_F^2$

 $\triangleright x_i$: i-th column of \hat{X} satisfies

$$\hat{x}_i = \arg\min_{x_i} \|SAx_i - Sb_i\|_2^2$$

 \triangleright For a Gaussian sketching matrix S we have

$$\mathbb{E}||A(\hat{x}_i - x_i^*)||_2^2 = ||Ax_i^* - b_i||_2^2 \frac{d}{m - d - 1}$$

implies

$$\mathbb{E}||A(\hat{X} - X^*)||_F^2 = \sum_{i=1}^r ||Ax_i^* - b_i||_2^2 \frac{d}{m - d - 1}$$
$$= ||AX^* - B||_F^2 \frac{d}{m - d - 1}$$

Left Sketching Optimality Gap

- ightharpoonup suppose that rank(A) = r
- original problem and left-sketch

$$X^* := \arg\min_{X} \|AX - B\|_F^2$$
 and $\hat{X} := \arg\min_{X} \|SAX - SB\|_F^2$

$$\mathbb{E}||A(\hat{X}-X^*)||_F^2 = ||AX^* - B||_F^2 \frac{r}{m-r-1}$$

$$\mathbb{E}||A\hat{X} - B||_F^2 = \mathbb{E}||AX^* - B + A(\hat{X} - X^*)||_F^2$$

$$= ||AX^* - B||_F^2 + \mathbb{E}||A(\hat{X} - X^*)||_F^2$$

$$= ||AX^* - B||_F^2 \left(1 + \frac{r}{m - r - 1}\right)$$

$$= ||AX^* - B||_F^2 \frac{m - 1}{m - r - 1}$$

Back to Randomized Low Rank Approximations

approximation error

$$\mathbb{E}\|\underbrace{AS}_{X^*}\underbrace{(AS)^{\dagger}A} - A\|_F^2 \leq \mathbb{E}\|AS(A_kS)^{\dagger}A_k - A\|_F^2$$

$$= \|A_k^T(S^TA_k^T)^{\dagger}S^TA^T - A^T\|_F^2$$

$$\mathbb{E}\|AT\tilde{x} - AT\|_2^2$$

$$= \|A_{k}^{T}(S^{T}A_{k}^{T})^{T}S^{T}A^{T} - A^{T}\|_{F}^{2}$$

$$= \mathbb{E}\|A_{k}^{T}\tilde{X} - A^{T}\|_{F}^{2}$$

$$\leq \frac{m-1}{m-k-1}\|A_{k}^{T}(A_{k}^{T})^{\dagger}A^{T} - A^{T}\|_{F}^{2}$$

$$\leq \frac{m-1}{m-k-1}\|A(A_{k}^{\dagger}A_{k}^{T} - I)\|_{F}^{2}$$
provetor into the top k subspace $\frac{m-1}{m-k-1}\|A_{k} - A\|_{F}^{2}$

$$\leq \frac{m-1}{m-k-1}\|A_{k} - A\|_{F}^{2}$$

$$\Rightarrow PF^{OX} \text{ escal } JD,$$

$$\downarrow + C$$

Randomized Low Rank Approximation and Randomized SVD Error Bound

- CX decomposition and randomized SVD
- ► $AS(AS)^{\dagger}A \approx A$
- final Frobenious norm error bound

$$\mathbb{E}||AS(AS)^{\dagger}A - A||_F^2 \le \frac{m-1}{m-k-1}||A_k - A||_F^2$$

▶ valid for any $k \in \{1, ..., rank(A)\}$

Randomized Low Rank Approximation and Randomized SVD Error Bound

- CX decomposition and randomized SVD
- $ightharpoonup AS(AS)^{\dagger}A \approx A$
- final Frobenious norm error bound

$$\sigma_{i}(A^{T}A) = \sigma_{i}(A)$$

$$A^{T}A\cdot A^{T}A = \sigma_{i}^{A}(A)$$

$$\mathbb{E}||AS(AS)^{\dagger}A - A||_F^2 \le \frac{m-1}{m-k-1}||A_k - A||_F^2$$

- ▶ valid for any $k \in \{1, ..., rank(A)\}$
- lacktriangle define the oversampling factor $\ell:=m-k-1$

$$\|AS(AS)^{\dagger}A - A\|_{F}^{2} \leq (1 + \frac{k}{\ell}) \|A_{k} - A\|_{F}^{2}$$

$$\|\sum_{i \in k \in I} \sigma_{i} u_{i} v_{i}^{-1}\|_{F}^{2} = \sum_{i \in k \in I} \sigma_{i}^{-1} u_{i} v_{i}^{-1}\|_{F}^{2}$$

Reducing the Error: Power Iteration

error bounds depend on tail singular values

$$||A_k - A||_F^2 = \sum_{j=k+1}^{\mathsf{rank}(A)} \sigma_j^2$$

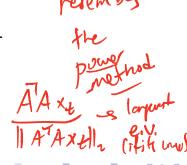
▶ idea: compute the sketch of $(AA^T)^qA$

$$C = (AA^{T})^{q}AS$$

$$S.v = (J; (A))^{21}$$

where q is an integer parameter

- ullet $CC^{\dagger}Approx$ A CC^{\dagger} approximates the range of A better for $q\geq 1$
- ▶ singular values of $(AA^T)^q A$ are $\sigma_i(A)^{2q+1}$ where $\sigma_i(A)$ are the singular values of A



Approximating Large Square Matrices

- Large and square matrices $A \in \mathbb{R}^{n \times n}$
- ► Regularized Least Squares ℓ_2 (Tikhonov) regularization

alternative form

$$= \min_{x} \left\| \left[\begin{array}{c} A \\ \sqrt{\lambda}I \end{array} \right] x - \left[\begin{array}{c} b \\ 0 \end{array} \right] \right\|_{2}^{2}$$

$$\left\| \left(\lambda I \right) \right\|_{2}^{2} = \lambda \left\| \lambda \right\|_{2}^{2}$$

Sketching Regularized Problems

$$\min_{x} \left\| \underbrace{\begin{bmatrix} A \\ \sqrt{\lambda}I \end{bmatrix}}_{\tilde{A}} x - \underbrace{\begin{bmatrix} b \\ 0 \end{bmatrix}}_{\tilde{b}} \right\|_{2}^{2}$$

- (ASSA) ASS b might not be invertible Left sketch min_x $||\tilde{SAx} - \tilde{Sb}||_2^2$ approximates the solution when sketch dimension m > d + 1, e.g., for Gaussian S
- Sketch dimension can be smaller if we use a partial sketch

$$\min_{x} \|SAx - Sb\|_{2}^{2} + \lambda \|x\|_{2}^{2}$$

$$\lim_{x} \|SAx - Sb\|_{2}^{2} + \lambda \|x\|_{2}^{2} + \lambda \|x\|_{2}^{2}$$

$$\lim_{x} \|SAx - Sb\|_{2}^{2} + \lambda \|x\|_{2}^{2} + \lambda \|x\|_{$$

>> (o(m'd)) Joseph ; word; see!

Sketching Regularized Problems

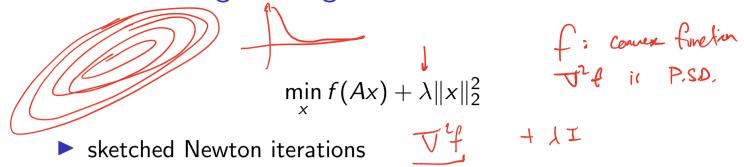


$$x^* = \arg\min_{x} \underbrace{\|Ax - b\|_2^2 + \lambda \|x\|_2^2}_{f(x)}$$

$$\hat{x} = \arg\min_{x} \|SAx - Sb\|_{2}^{2} + \lambda \|x\|_{2}^{2}$$

- ▶ approximation ratio $f(\hat{x}) \leq f(x^*)(1 + \epsilon)$ when $m \geq \text{constant} \times d_e(\lambda)$ for i.i.d. Gaussian, sub-Gaussian and FJLT sketch (ignoring log factors)
- $ightharpoonup d_e(\lambda) = \sum_{i=1}^d \frac{\sigma_i(A)^2}{\sigma_i(A)^2 + \lambda}$ is the effective dimension of A
- ho $d_e(0) = \operatorname{rank}(A)$ our dela) < rank(A)

Hessian Sketching for Regularized Problems



$$x_{t+1} = \arg\min_{x} \frac{1}{2} \|S(\nabla^{2} f(x_{t}))^{1/2} \|_{2}^{2} + (x - x_{t})^{T} \nabla f(x_{t}) + \frac{\lambda}{2} \|x\|_{2}^{2}$$

$$(x_{t})^{\lambda} \|\nabla^{2} f(x_{t})\|^{2} \|\nabla^{2} f(x_{t})\|^{2} \|x\|_{2}^{2}$$

- $(\nabla^2 f(x_t))^{1/2} S^T S(\nabla^2 f(x_t))^{1/2} + \lambda I$ is invertible for all m when $\lambda > 0$
- similar guarantees involving the effective dimension of the Hessian matrix
- $\lambda = 0$ requires m > d for invertibility

▶ Large square matrices $K \in \mathbb{R}^{n \times n}$

Kernel Ridge Regression

K is called the kernel matrix

 κ is the **kernel function**

examples:

the training set are $K\alpha \approx y$

es
$$K \in$$

$$K \in I$$

 $ightharpoonup K = \kappa(x_i, x_j)$ where $x_1, ..., x_n \in \mathbb{R}^d$ are data vectors,

Gaussian kernel $K_{ij} = \kappa(x_i, x_i) = e^{-\frac{1}{\sigma^2} ||x_i - x_j||_2^2}$

▶ prediction at a point x is $\sum_{i=1}^{n} \kappa(x_i, x) \alpha_i$ i.e, predictions on

Kernel matrices typically have low effective dimension, e.g.,

choice of λ provides optimal statistical guarantees

▶ Gaussian kernel has $d_e(\lambda) = O(\sqrt{\log n})$ for $\lambda = \sqrt{\frac{\log n}{n}}$. This

Polynomial kernel $K_{ij} = \kappa(x_i, x_j) = (x_i^T x_j)^r$

$$K \in \mathbb{I}$$

 $\min_{\alpha} ||K\alpha - y||_2^2 + \lambda \alpha^T K\alpha$

Kernel Trick

 $X_i X_j$ line $K_{ij} = (X_i X_j)^2$

Kernel Ridge Regression

$$\min_{\alpha} ||K\alpha - y||_2^2 + \lambda \alpha^T K \alpha$$

example: polynomial kernel (degree 2) $K_{ij} = \kappa(x_i, x_j) = (x_i^T x_j)^2 = (\sum_i x_i x_i^T x_j^T)^2$

maps data to higher dimension ¹

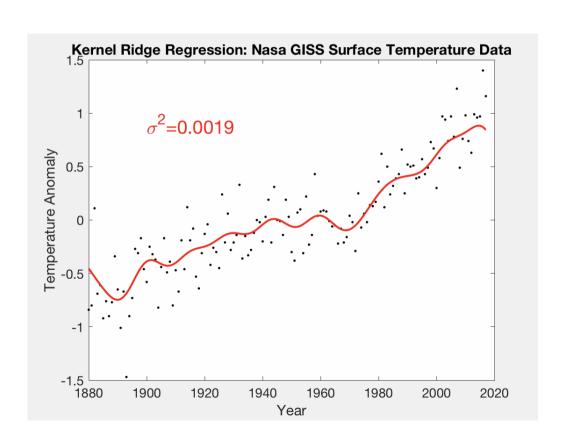
$$A = \begin{bmatrix} x_{11} & \dots & x_{1d} \\ \vdots & & & \\ x_{n1} & \dots & x_{nd} \end{bmatrix} \rightarrow$$

$$\tilde{A} := \begin{bmatrix} x_{11} & \dots & x_{1d} & x_{11}^2 & \dots & x_{1d}^2 \\ \vdots & & & & \\ x_{n1} & \dots & x_{nd}^2 & x_{11}^2 & \dots & x_{nd} \end{bmatrix}$$

Application: Kernel Regression



Gaussian Kernel
$$K_{ij} = e^{-\frac{\|x_i - x_j\|_2^2}{2\sigma^2}}$$



Application: Kernel Classification

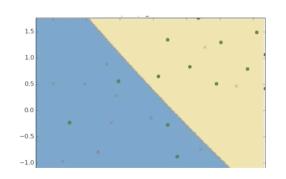
Pight sleeting

$$\begin{array}{c}
(-,y) = (--y)^{n} \\
|K \cdot S + y|^{n} + |A \cdot S | | |K \cdot S + |K \cdot S | |K$$

$$\min_{\alpha} \sum_{i=1}^{n} \ell(K\alpha, y) + \lambda \alpha^{T} K\alpha$$

linear kernel
$$K_{ij} = x_i^T x_j$$

gaussian kernel
$$K_{ij} = e^{-\frac{\|x_i - x_j\|_2^2}{2\sigma^2}}$$





Nystrom Method

KS/KSTK 2K

- - We need a symmetric approximation. CX decomposition is not symmetric.
 - Most kernel matrices are positive semidefinite, i.e., A night he had to compute $K = A^T A$ for some matrix A
 - ▶ Recall the CX decomposition $\tilde{A} = (AS)(AS)^{\dagger}A \approx A$ we used
 - ightharpoonup Consider approximating A^TA via $\tilde{A}^T\tilde{A}$ \prod A \int fill column rank

$$((AS)(AS)^{\dagger}A)^{T}(AS)(AS)^{\dagger}A = A^{T}(AS)(AS)^{\dagger}(AS)(AS)^{\dagger}A$$

$$= A^{T}(AS)(AS)^{\dagger}A$$

$$= A^{T}AS(S^{T}A^{T}AS)^{-1}S^{T}A^{T}A$$

 $\tilde{K} = KS(S^TKS)^{-1}S^TK \approx K$

randomized low rank approximation of K is given by



Nystrom Method: S is uniform column sampling weighted sampling or sketching can also be used

Generalized Nystrom Method

- Nystrom method can be generalized to non symmetric matrices
- ► Consider CX decomposition where C = AS and S is a sketching matrix

$$\min_{X} \|ASX - A\|_{F}$$

 \triangleright Apply another sketching matrix R on the left

$$\min_{X} \|RASX - RA\|_{F}$$

- ▶ solution $X^* = (RAS)^{\dagger}RA$
- ► approximation of A is $AS(RAS)^{\dagger}RA \approx A$
- reduces to the Nystrom method when R = S and $A = A^T$
- ▶ faster than CX and randomized SVD, less accurate



Random Fourier Features

- Random approximations of kernel matrices
- Generate $w \sim N(0, I)$
- ▶ Define features $h(x) := e^{-jw^T x}$ where $j = \sqrt{-1}$ it holds that

$$\mathbb{E}_{w}h(x)h(y)^{*} = \mathbb{E}_{w}e^{-jw^{T}x}e^{+jw^{T}y}$$

$$= \mathbb{E}_{w}e^{-jw^{T}(x-y)}$$

$$= \int p(w)e^{-jw^{T}(x-y)}dw$$

$$= e^{-\frac{1}{2}(x-y)^{T}(x-y)} \mathcal{E}_{\text{faction kernel}}.$$

- ightharpoonup where p(w) is the multivariate Gaussian distribution
- ▶ **Bochner's Theorem:** Fourier transforms of probability distributions correspond to positive semidefinite kernels
- Gaussian distribution corresponds to the Gaussian kernel



Random Fourier Features

- Random approximations of kernel matrices
- ► Generate $w_1, ..., w_m \sim N(0, I)$ i.i.d.
- Define feature vectors

$$h(x) = \frac{1}{\sqrt{m}} \begin{bmatrix} e^{jw_1^T x} \\ e^{jw_2^T x} \\ \dots \\ e^{jw_m^T x} \end{bmatrix} = \underbrace{\downarrow}_{\text{can}} e \times p(S \times \frac{1}{3})$$

then we have

$$\langle h(x), h(y) \rangle = \frac{1}{m} \sum_{i=1}^{m} e^{jw_i^T(x-y)} \approx \mathbb{E}_w e^{jw^T(x-y)} = e^{-\frac{1}{2}(x-y)^T(x-y)}$$

$$= \mathcal{K}(x,y)$$

Rahimi and Recht, Random Features for Large-Scale Kernel Machines, 2007 Keexilds: j) > FXX=K

Random Fourier Features

- ► The embedding is nonlinear $\frac{1}{\sqrt{m}} \exp(iXS)$
- can also be obtained using real valued embeddings
- Generate $w \sim N(0, I)$ i.i.d.
- ► $h(x) = \sqrt{2}\cos(w^Tx + b)$ where $b \sim \text{Uniform}(0, 2\pi)$ also works
- the approximation error

$$\left\| \underbrace{\frac{1}{m} \sum_{i=1}^{m} e^{jw_i^T(x_i - y_i)}} - \mathbb{E} \frac{1}{m} \sum_{i=1}^{m} e^{jw_i^T(x_i - y_i)} \right\|_2 \text{ can be controlled}$$
via matrix concentration bounds