

Accessing Parameters

Working the stack!

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Quick reminder about C

```
#include <stdio.h>

int main(int argc, char* argv[]){
    printf("%d arguments\n", argc);
    printf("%s\n", argv[0]);

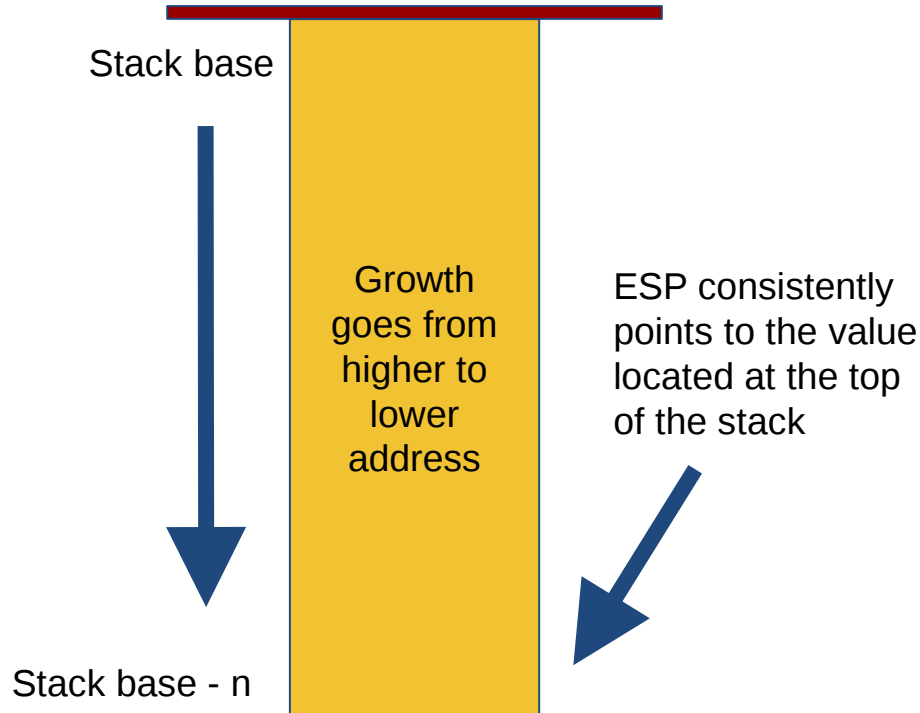
    if (argc > 1){
        printf("%s\n", argv[1]);
    }
}
```

main receives 2 parameters

- **argc** is the **number of** command line **parameters**
- **argv** is an array of char* (strings) each containing **separate** command line **arguments**

On IA-32, parameters are passed on the stack in reverse order of appearance in a function declaration.

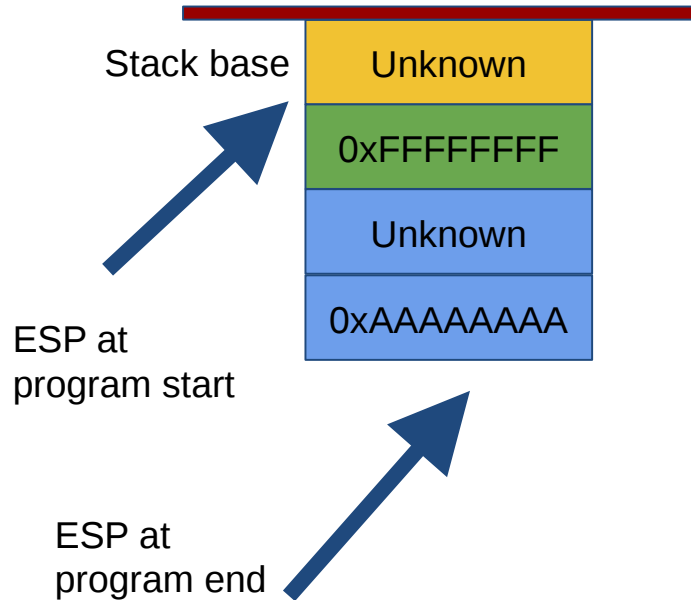
The stack



Multiple instructions can change the stack pointer location and add or remove elements from the stack among the most popular:

- *push value*
- *pop register*
- *call label*
- *ret*
- *add* destination, source (addition)
- *sub* destination, source (subtraction)
- etc...

Example



The following instructions changes the stack to the presented state:

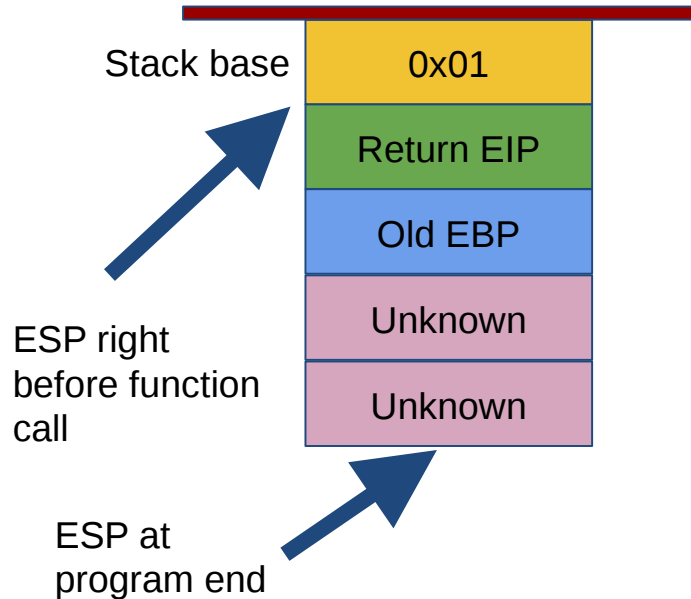
- `push 0xFFFFFFFF`
- `sub esp, 0x08`
- `mov [esp], 0xAAAAAAAA`

What would be the result of instruction:

- `add esp, 0xC`

What are the consequence of using that instruction?

More complicated example

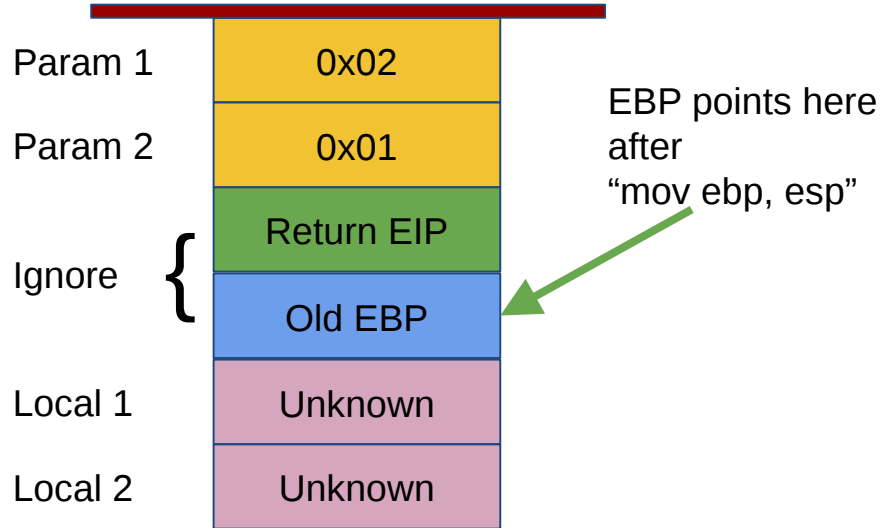


- *push 0x01*
- *call randomFunction*
- ***** Code for randomFunction*
- *push ebp*
- *mov ebp, esp*
- *sub esp, 0x08*

For now, ignore the return EIP and the old EBP parts. We will be coming back to this later.

We will call everything from the red line to the esp after the sub a "Stack frame"

Stack Frame



- *push 0x02*
- *push 0x01*
- *call randomFunction*
- **** Code for randomFunction
- *push ebp*
- *mov ebp, esp*
- *sub esp, 0x08*

For now, ignore the return EIP and the old EBP parts. We will be coming back to this later.

Param 1 can be accessed as $[ebp + 0xC]$

Param 2 can be accessed as $[ebp + 0x8]$

Why?

We will call everything from the red line to the esp after the sub a "Stack frame"



The compiler monster!



Different compilers, different flavors of code!

```
(gdb) disass $eip
Dump of assembler code for function main:
   0x08048410 <+0>:      push    ebp
   0x08048411 <+1>:      mov     ebp,esp
   0x08048413 <+3>:      sub     esp,0x18
=> 0x08048416 <+6>:      mov     eax,DWORD PTR [ebp+0xc]
   0x08048419 <+9>:      mov     ecx,DWORD PTR [ebp+0x8]
   0x0804841c <+12>:     mov     DWORD PTR [ebp-0x4],0x0
   0x08048423 <+19>:     mov     DWORD PTR [ebp-0x8],ecx
   0x08048426 <+22>:     mov     DWORD PTR [ebp-0xc],eax
   0x08048429 <+25>:     mov     DWORD PTR [ebp-0x10],0x0
   0x08048430 <+32>:     mov     DWORD PTR [ebp-0x14],0x0
   0x08048437 <+39>:     mov     eax,DWORD PTR [ebp-0xc]
   0x0804843a <+42>:     mov     eax,DWORD PTR [eax]
   0x0804843c <+44>:     mov     ecx,esp
   0x0804843e <+46>:     mov     DWORD PTR [ecx],eax
   0x08048440 <+48>:     call    0x80482e0 <strlen@plt>
   0x08048445 <+53>:     mov     DWORD PTR [ebp-0x10],eax
   0x08048448 <+56>:     cmp     DWORD PTR [ebp-0x8],0x1
   0x0804844c <+60>:     jle     0x8048464 <main+84>
   0x08048452 <+66>:     mov     eax,DWORD PTR [ebp-0xc]
   0x08048455 <+69>:     mov     eax,DWORD PTR [eax+0x4]
   0x08048458 <+72>:     mov     ecx,esp
   0x0804845a <+74>:     mov     DWORD PTR [ecx],eax
   0x0804845c <+76>:     call    0x80482e0 <strlen@plt>
   0x08048461 <+81>:     mov     DWORD PTR [ebp-0x14],eax
   0x08048464 <+84>:     mov     eax,DWORD PTR [ebp-0x8]
   0x08048467 <+87>:     add     esp,0x18
   0x0804846a <+90>:     pop     ebp
   0x0804846b <+91>:     ret
End of assembler dump.
```

clang is presented at the left while gcc is presented at the right. The compiler has a direct impact on the code being generated! Visual C++ is closer to clang. The important thing is the end code. All of these are relatively equivalent.

```
Dump of assembler code for function main:
   0x0804840b <+0>:      lea     ecx,[esp+0x4]
   0x0804840f <+4>:      and     esp,0xffffffff
   0x08048412 <+7>:      push    DWORD PTR [ecx-0x4]
   0x08048415 <+10>:     push    ebp
   0x08048416 <+11>:     mov     ebp,esp
   0x08048418 <+13>:     push    ebx
   0x08048419 <+14>:     push    ecx
=> 0x0804841a <+15>:     sub     esp,0x10
   0x0804841d <+18>:     mov     ebx,ecx
   0x0804841f <+20>:     mov     DWORD PTR [ebp-0x10],0x0
   0x08048426 <+27>:     mov     DWORD PTR [ebp-0xc],0x0
   0x0804842d <+34>:     mov     eax,DWORD PTR [ebx+0x4]
   0x08048430 <+37>:     mov     eax,DWORD PTR [eax]
   0x08048432 <+39>:     sub     esp,0xc
   0x08048435 <+42>:     push    eax
   0x08048436 <+43>:     call    0x80482e0 <strlen@plt>
   0x0804843b <+48>:     add     esp,0x10
   0x0804843e <+51>:     mov     DWORD PTR [ebp-0x10],eax
   0x08048441 <+54>:     cmp     DWORD PTR [ebx],0x1
   0x08048444 <+57>:     jle     0x804845d <main+82>
   0x08048446 <+59>:     mov     eax,DWORD PTR [ebx+0x4]
   0x08048449 <+62>:     add     eax,0x4
   0x0804844c <+65>:     mov     eax,DWORD PTR [eax]
   0x0804844e <+67>:     sub     esp,0xc
   0x08048451 <+70>:     push    eax
   0x08048452 <+71>:     call    0x80482e0 <strlen@plt>
   0x08048457 <+76>:     add     esp,0x10
   0x0804845a <+79>:     mov     DWORD PTR [ebp-0xc],eax
   0x0804845d <+82>:     mov     eax,DWORD PTR [ebx]
   0x0804845f <+84>:     lea     esp,[ebp-0x8]
   0x08048462 <+87>:     pop     ecx
   0x08048463 <+88>:     pop     ebx
   0x08048464 <+89>:     pop     ebp
   0x08048465 <+90>:     lea     esp,[ecx-0x4]
   0x08048468 <+93>:     ret
```


Different level of optimization also change the code

```
(gdb) disass main
Dump of assembler code for function main:
   0x080483db <+0>:    mov     eax,DWORD PTR [esp+0x4]
   0x080483df <+4>:    ret
End of assembler dump.
```

GCC used with -O3
optimization

Why was all the code removed?



Understanding stack frames is central to assembly programming and software reverse engineering.



Review

- Let's say a function has 4 parameters, how can we access parameter #3?
- And if we wanted to access parameter #4?
- And the first parameter?
- EBP is normally used to access both local variables and function parameters. How can we quickly know if an access (relative to EBP) is done to function parameters or to the local variables?
- How about writing some code...

