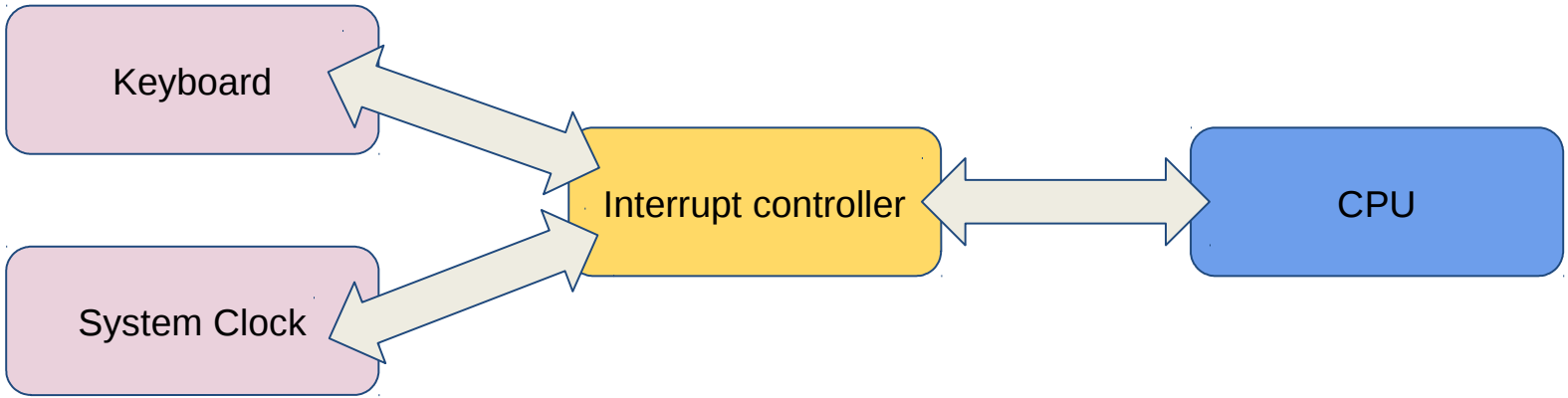
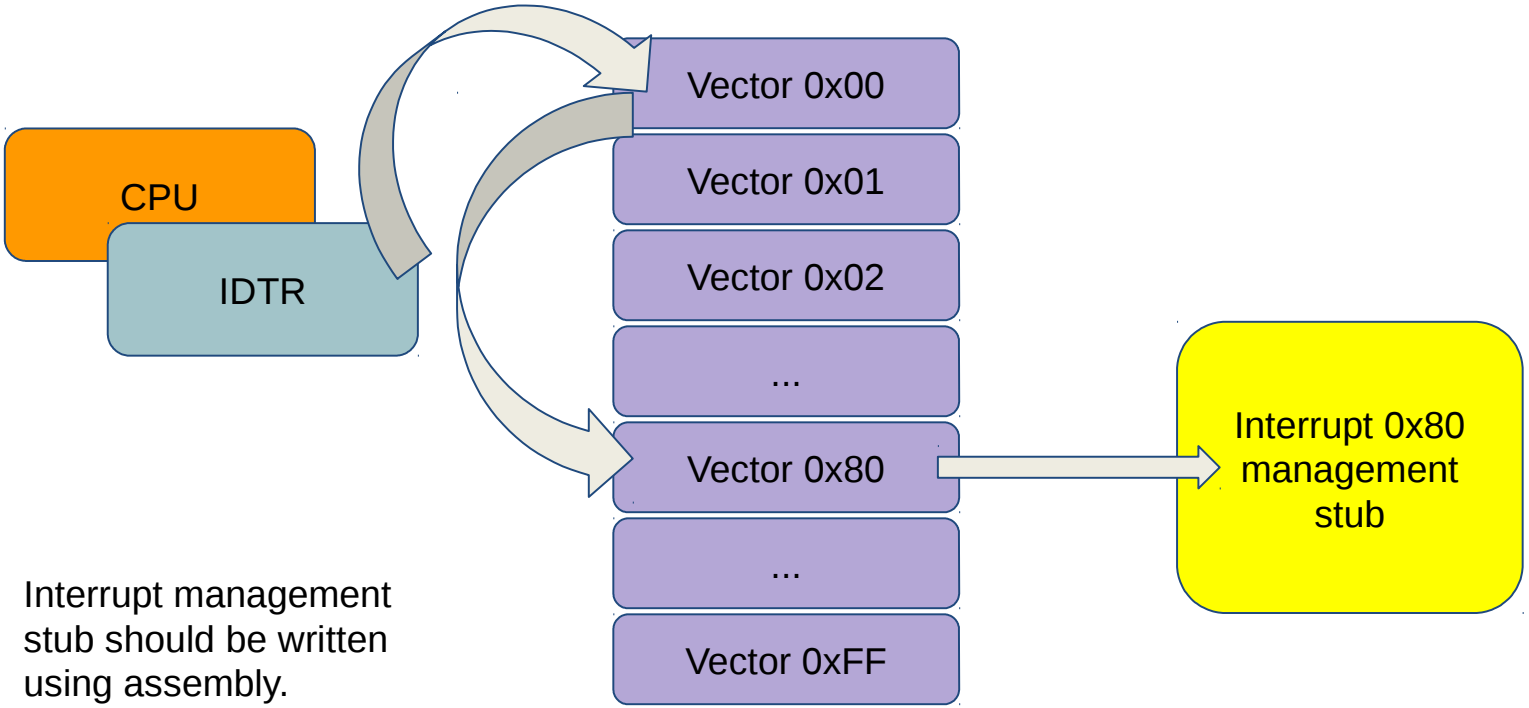
Interrupt Management

We are interrupting this presentation so we can have a quick chat about x86 interrupt management.

High level view of the system

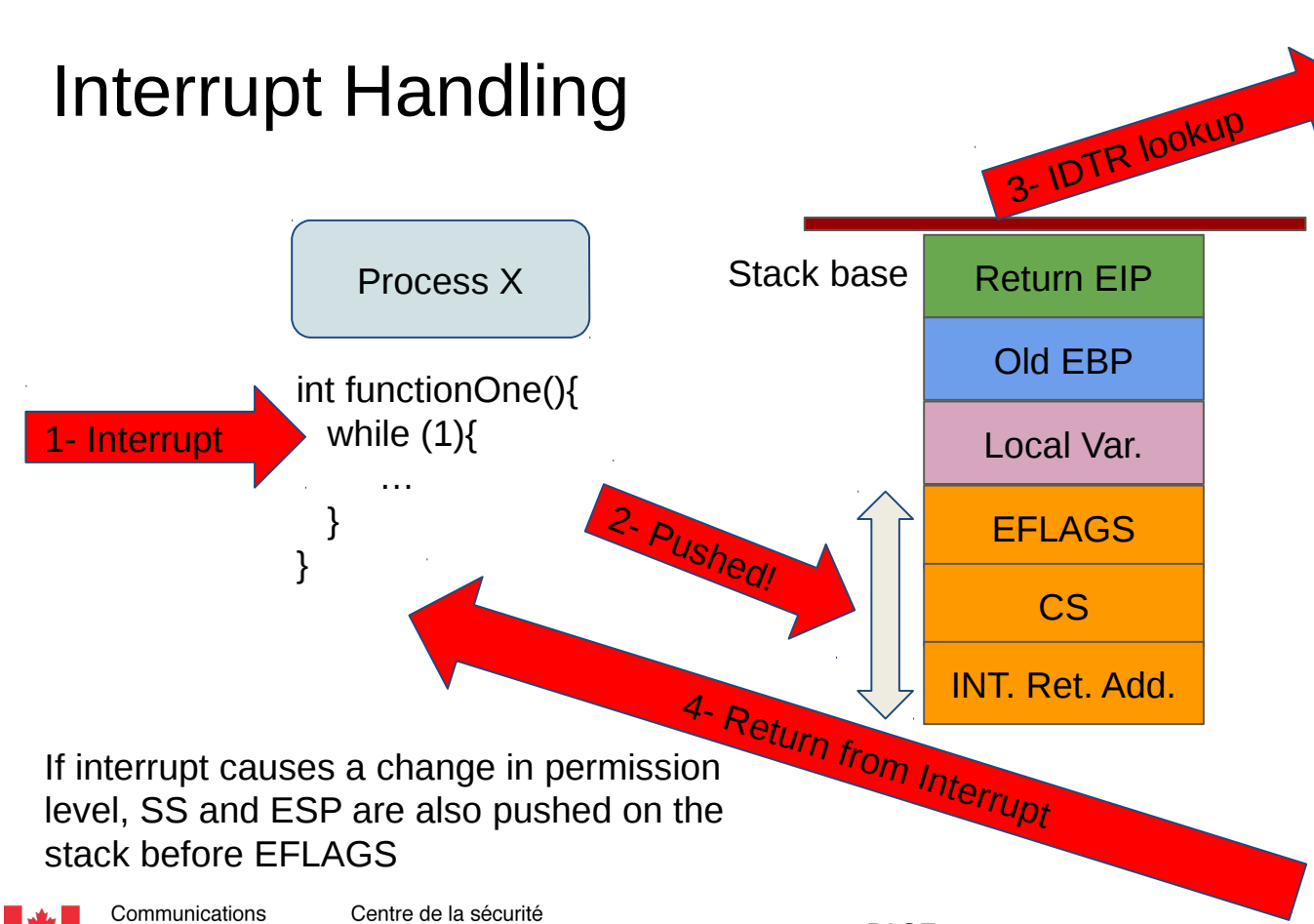


Interrupt delivery to the CPU



Interrupt management stub should be written using assembly.

Interrupt Handling




```
_handlerKeyboard:
extern _readKeyboard
extern khandleKeyboard
    pusha
    mov al, 0x20 ; Interrupt ACK
    out 0x20, al; ack irq
    out 0xA0, al; ack irq

    xor eax, eax
    call _readKeyboard
    test eax, eax
    je handlerK_NoData

    push eax
    call khandleKeyboard
    add esp, 0x04
handlerK_NoData:
    popa
    iret
```

If interrupt causes a change in permission level, SS and ESP are also pushed on the stack before EFLAGS

x86 Exceptions

Exception	Description
0	Divide error
6	Invalid opcode
12	Stack segment fault
14	Page fault
17	Alignement check
19	SIMD floating point exception
...	...

On Linux

traps.c

```
83 gate_desc idt_table[NR_VECTORS] __page_aligned_bss;
```

segment.h

```
217 #define IDT_ENTRIES                     256
218 #define NUM_EXCEPTION_VECTORS           32
```

head_32.s

```
391         movl $idt_table,%edi
392         movl $early_idt_handler_array,%eax
393         movl $NUM_EXCEPTION_VECTORS,%ecx
394 1:
395         movl %eax,(%edi)
396         movl %eax,4(%edi)
397         /* interrupt gate, dpl=0, present */
398         movl $(0x8E000000 + __KERNEL_CS),2(%edi)
399         addl $EARLY_IDT_HANDLER_SIZE,%eax
400         addl $8,%edi
401         loop 1b
402
403         movl $256 - NUM_EXCEPTION_VECTORS,%ecx
404         movl $ignore_int,%edx
405         movl $(__KERNEL_CS << 16),%eax
406         movw %dx,%ax            /* selector = 0x0
407         movw $0x8E00,%dx        /* interrupt gate
408 2:
409         movl %eax,(%edi)
410         movl %edx,4(%edi)
411         addl $8,%edi
412         loop 2b
```

head_32.s

```
641 idt_descr:
642         .word IDT_ENTRIES*8-1
643         .long idt_table
```

```
349         lgdt early_gdt_descr
350         lidt idt_descr
351         ljmp $(__KERNEL_CS),$1f
352 1:      movl $(__KERNEL_DS),%eax
353         movl %eax,%ss
```

On Linux

irq_vectors.h

```
49 #define IA32_SYSCALL_VECTOR             0x80
```

desc.h

```
506 /*
507  * This routine sets up an interrupt gate at directory privilege level
508  */
509 static inline void set_system_intr_gate(unsigned int n, void *addr)
510 {
511         BUG_ON((unsigned)n > 0xFF);
512         _set_gate(n, GATE_INTERRUPT, addr, 0x3, 0, __KERNEL_CS);
513 }
```

traps.c

```
969 void __init trap_init(void)
970 {
971         int i;
972
973 #ifdef CONFIG_EISA
974         void __iomem *p = early_ioremap(0x0FFFD9, 4);
975
976         if (readl(p) == 'E' + ('I'<<8) + ('S'<<16) + ('A'<<24))
977                 EISA_bus = 1;
978         early_iounmap(p, 4);
979 #endif
980
981         set_intr_gate(X86_TRAP_DE, divide_error);
982         set_intr_gate_ist(X86_TRAP_NMI, &nmi, NMI_STACK);
983         /* int4 can be called from all */
984         set_system_intr_gate(X86_TRAP_OF, &overflow);
985         set_intr_gate(X86_TRAP_BR, bounds);
986         set_intr_gate(X86_TRAP_UD, invalid_op);
987         set_intr_gate(X86_TRAP_NM, device_not_available);
1015 #ifdef CONFIG_X86_32
1016         set_system_intr_gate(IA32_SYSCALL_VECTOR, entry_INT80_32);
1017         set_bit(IA32_SYSCALL_VECTOR, used_vectors);
1018 #endif
```

entry_32.s

```
519 ENTRY(entry_INT80_32)
520         ASM_CLAC
521         pushl   %eax
522         SAVE_ALL pt_regs_ax=$-ENOSYS
523
524         /*
525          * User mode is traced as though
526          * turned them off.
527          */
528         TRACE_IRQS_OFF
529
530         movl    %esp, %eax
531         call    do_int80_syscall_32
```

On Linux

common.c

```
335 /* Handles int $0x80 */
336 __visible void do_int80_syscall_32(struct pt_regs *regs)
337 {
338         enter_from_user_mode();
339         local_irq_enable();
340         do_syscall_32_irqs_on(regs);
341 }
```

You are free to continue your exploration from here :)
The point is simply to make you understand a little bit of what happens when a syscall is made on Linux before moving one to understanding how to do a syscall using assembly programming.