

# Dealing with shellcode

*Tooling ourselves to open that code*

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# What is a shellcode?



# Shellcode

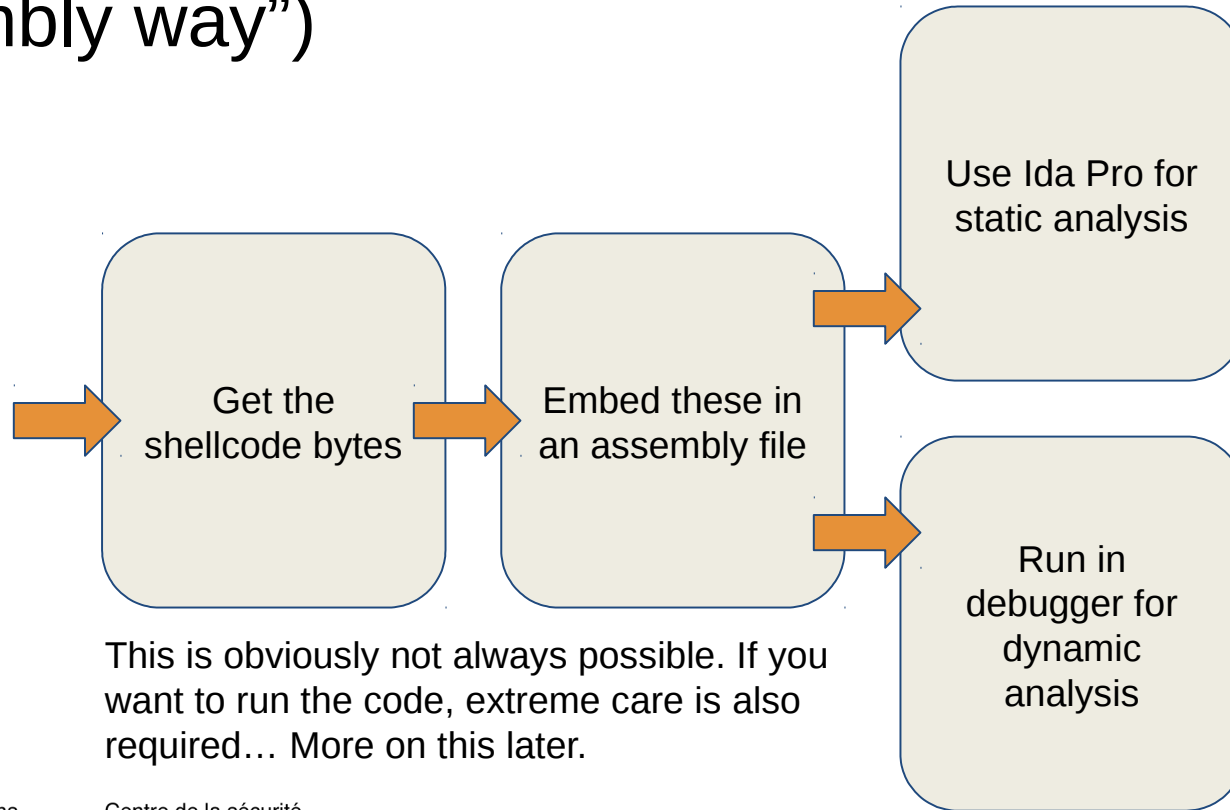
- A shellcode is a piece of software that is usually present as part of an attack
  - As part of an exploit, the shellcode will generally be the code that will be launched after the attacker gained code execution on a machine
- Shellcodes tend to be small
- A great deal of care is put in crafting these
  - A bad shellcode would mean a failure for an attacker
  - Often need to remove all NULL chars from a shellcode
  - Need to be position independent code (PIC)
- Shellcode analysis can be a bit daunting if you never did it
  - This is why we're going to do a small break in the programming aspect of the class and will be doing shellcode analysis for a little while! :)



All shellcode analysed here are part of  
the open source Metasploit framework



# Shellcode analysis process (easy “hey I know assembly way”)



# The dynamic way

Let's pretend the following bytes were found and are suspected part of a shell code:

```
6a 0b 58 99 52 66 68 2d 63 89 e7 68 2f 73 68
00 68 2f 62 69 6e 89 e3 52 e8 09 00 00 00 65
63 68 6f 20 6d 30 30 00 57 53 89 e1 cd 80
```

If you are familiar with x86 opcode, you might have noticed the ending:  
0xCD 0x80

This generally suggest a system call since this directly translate to:  
int 0x80

And therefore suggest a possible Linux shellcode

Original code used in this example can be found at please play fairgame and don't check the original code before the end of demo::  
[https://github.com/rapid7/metasploit-framework/blob/master/external/source/shellcode/linux/ia32/single\\_exec.asm](https://github.com/rapid7/metasploit-framework/blob/master/external/source/shellcode/linux/ia32/single_exec.asm)



# The dynamic way

Why do the unconditional jump to main in a case like this one?

Packaging the shellcode like this has many advantages...

Also when you actually are ready to run the code, you can change “jmp main” to “jmp shellcode” and you will be able to run the shellcode in a debugger.

```
2 section .text
3 global main
4
5 main:
6     jmp main
7
8 shellcode:
9     db 0x6a, 0x0b, 0x58, 0x99, 0x52, 0x66, 0x68
10    db 0x2d, 0x63, 0x89, 0xe7, 0x68, 0x2f, 0x73
11    db 0x68, 0x00, 0x68, 0x2f, 0x62, 0x69, 0x6e
12    db 0x89, 0xe3, 0x52, 0xe8, 0x09, 0x00, 0x00
13    db 0x00, 0x65, 0x63, 0x68, 0x6f, 0x20, 0x6d
14    db 0x30, 0x30, 0x00, 0x57, 0x53, 0x89, 0xe1
15    db 0xcd, 0x80
```

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[https://github.com/rapid7/metasploit-framework/blob/master/external/source/shellcode/linux/ia32/single\\_exec.asm](https://github.com/rapid7/metasploit-framework/blob/master/external/source/shellcode/linux/ia32/single_exec.asm)



Let's do this one together and then you'll  
do one on your own.

