

Patterns for Efficient Software

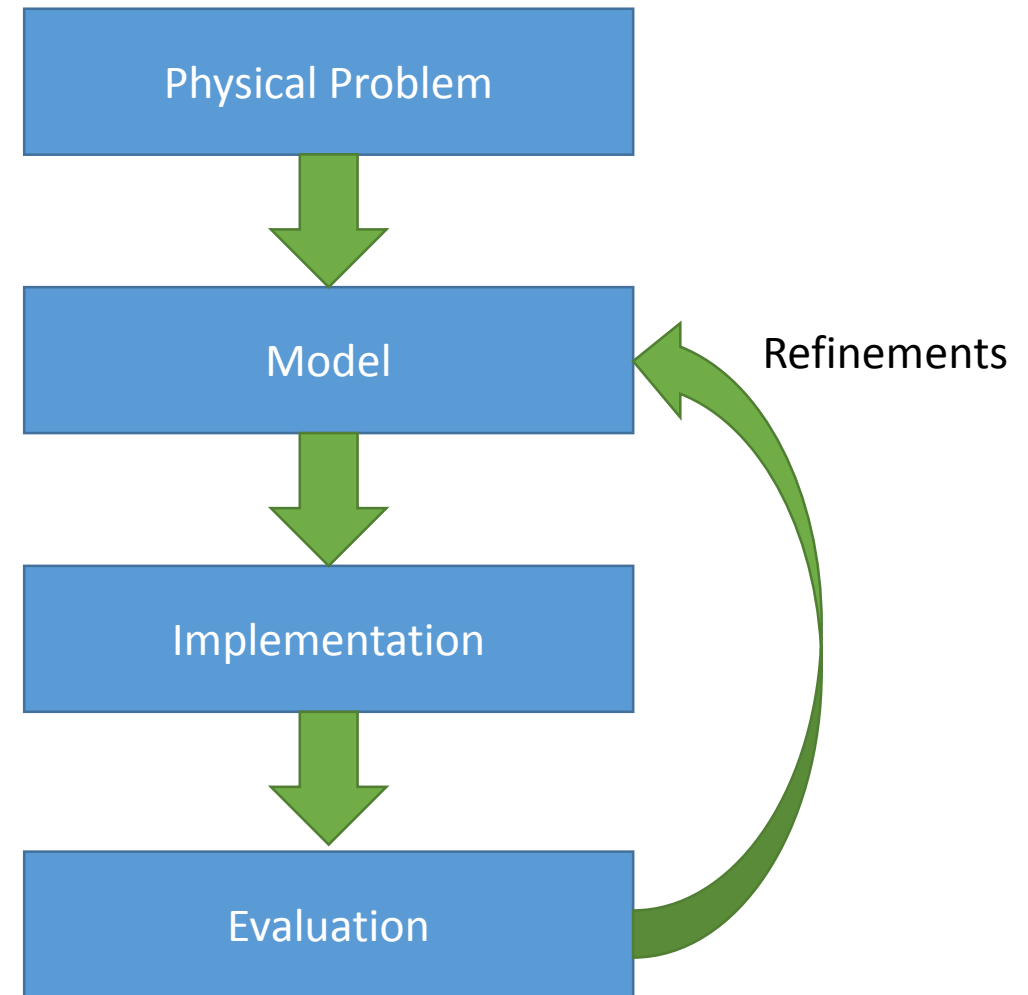
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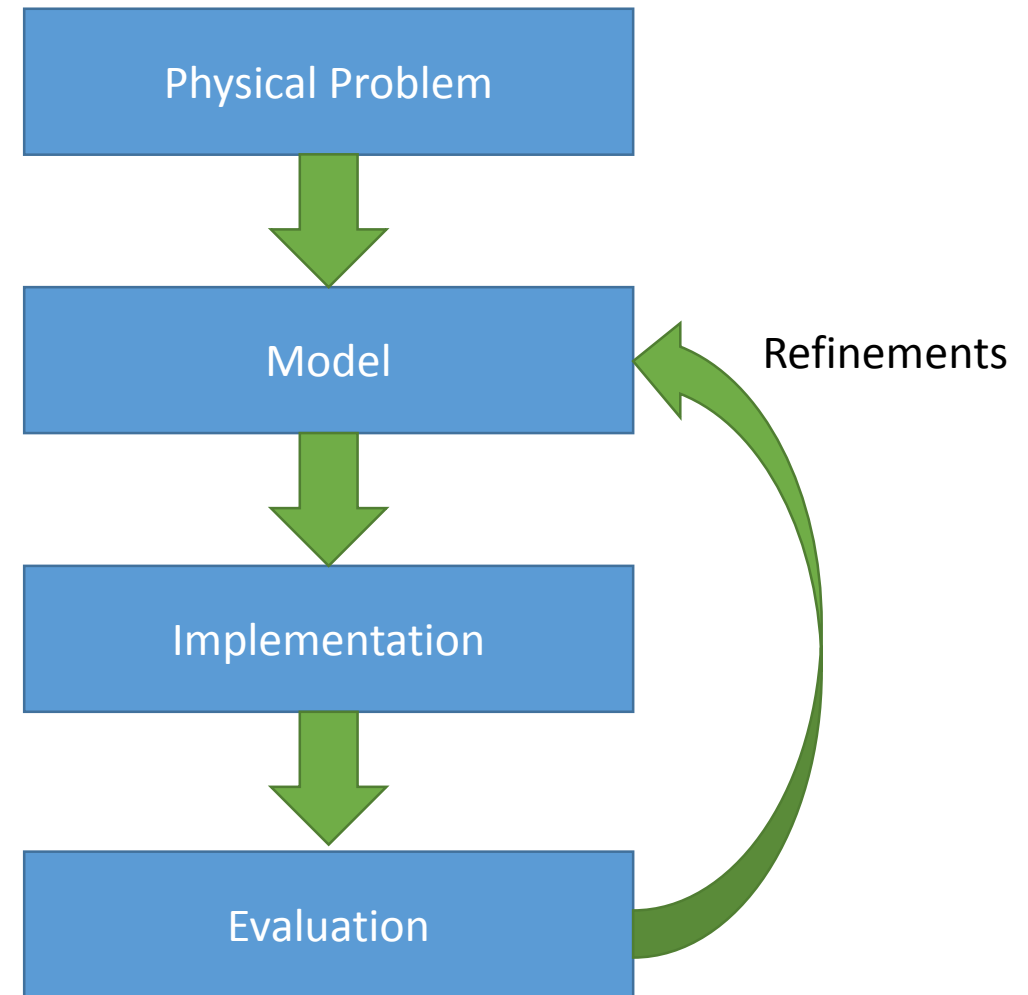
What writing scientific codes looks like...

- Many of us write programs to solve specific problems in science
- We create and use models to describe our problems
- These models are implemented as code and produce results
- Evaluating these results allows us to validate our models and improve them



What writing scientific codes looks like...

- Performance, Maintainability, Portability of those programs is not our main concern
- What matters is that the results we compute are **correct!**
 - First working solution often ends up in production code
 - Sub-optimal solutions are chosen, which only get noticed later in the process (e.g. when working on bigger data sets)



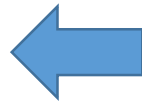
There is more than one way to Rome...

- There is usually more than one way of solving your problem
- It's very likely that someone other than you has already solved at least parts of your problem (and in a better way than you could have ever imagined)
- So the key approach to this is:
 - Know what's out there
 - Figure out a way to choose the best solution



What this talk is about

- Choosing algorithms
 - How to compare algorithms
 - Example: Sorting
- Choosing data structures
 - Performance Characteristics
 - Types



also covered:

Generic Programming,
Examples: C++ STL

Algorithms

“Informally, an **algorithm** is any well-defined computational procedure that takes some value, or set of values, as **input** and produces some value, or set of values, as **output**.

We can also view an algorithm as a tool for solving a well-specified **computational problem**. The statement of the problem specifies in general terms the desired input/output relationship. The algorithm describes a specific computational procedure for achieving that input/output relationship.”

from *Introduction to Algorithms* (3. Edition), by Thomas H. Cormen et. al.

Comparing Algorithms

Theoretically:

- **by runtime complexity**
- **by memory usage**

Practically:

- **by measurement (profiling)**

Sorting

Input Sequence

10	2	7	5	1	4	9	3	6	8
----	---	---	---	---	---	---	---	---	---

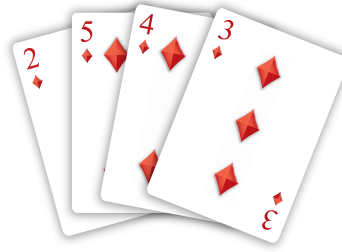
Algorithm



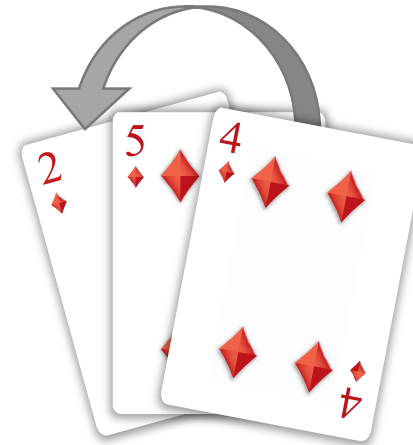
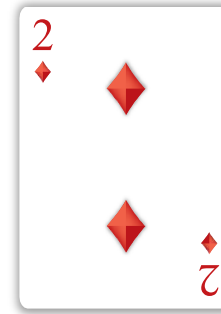
Output Sequence

1	2	3	4	5	6	7	8	9	10
---	---	---	---	---	---	---	---	---	----

Insertion Sort

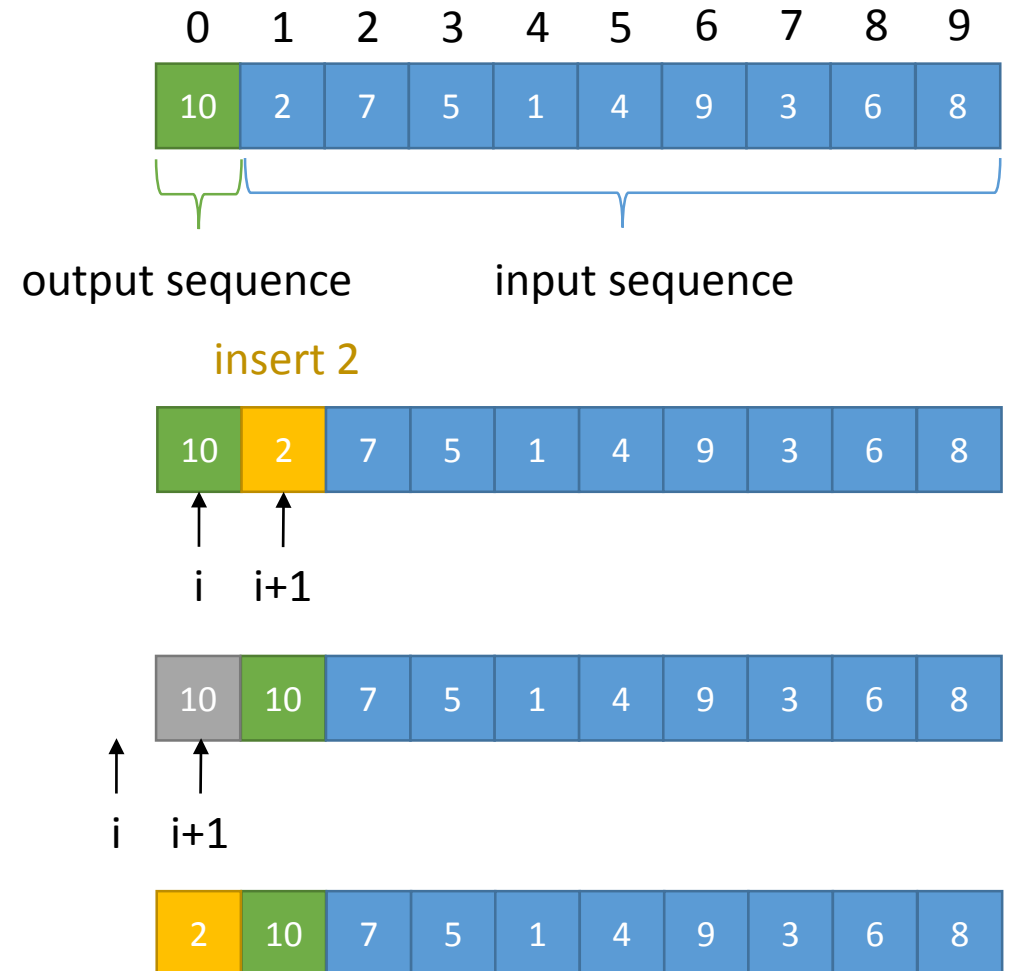


- Start out with an unordered input sequence of elements and an empty output sequence
- You then remove one element from the input sequence and insert it into the output sequence from right to left at the correct location
- Repeat until the input sequence is empty



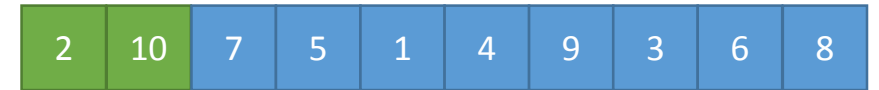
Insertion Sort – Implementation (C++)

```
void insertion_sort(std::vector<int> & v) {  
    for(int j = 1; j < v.size(); j++) {  
        int key = v[j];  
        int i = j - 1;  
        while(i >= 0 && v[i] > key) {  
            v[i+1] = v[i];  
            i--;  
        }  
        v[i+1] = key;  
    }  
}
```



Insertion Sort – Implementation (C++)

```
void insertion_sort(std::vector<int> & v) {  
    for(int j = 1; j < v.size(); j++) {  
        int key = v[j];  
        int i = j - 1;  
        while(i >= 0 && v[i] > key) {  
            v[i+1] = v[i];  
            i--;  
        }  
        v[i+1] = key;  
    }  
}
```

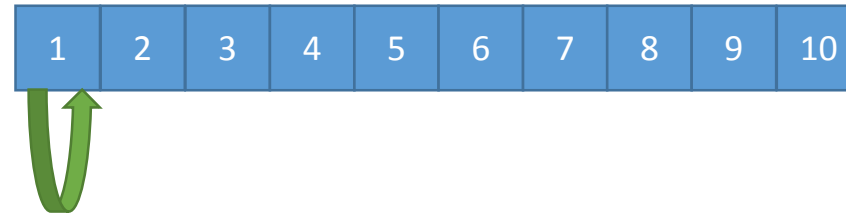


insert 7

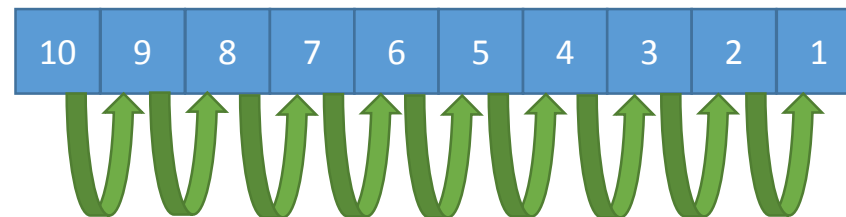


Insertion Sort

Best case: Sequence already in sorted order



Worst case: Sequence in reverse order



Insertion Sort

```

void insertion_sort(std::vector<int> & v) {
    for(int j = 1; j < v.size(); j++) {
        int key = v[j];
        int i = j - 1;

        while(i >= 0 && v[i] > key) {

            v[i+1] = v[i];

            i--;
        }
        v[i+1] = key;
    }
}

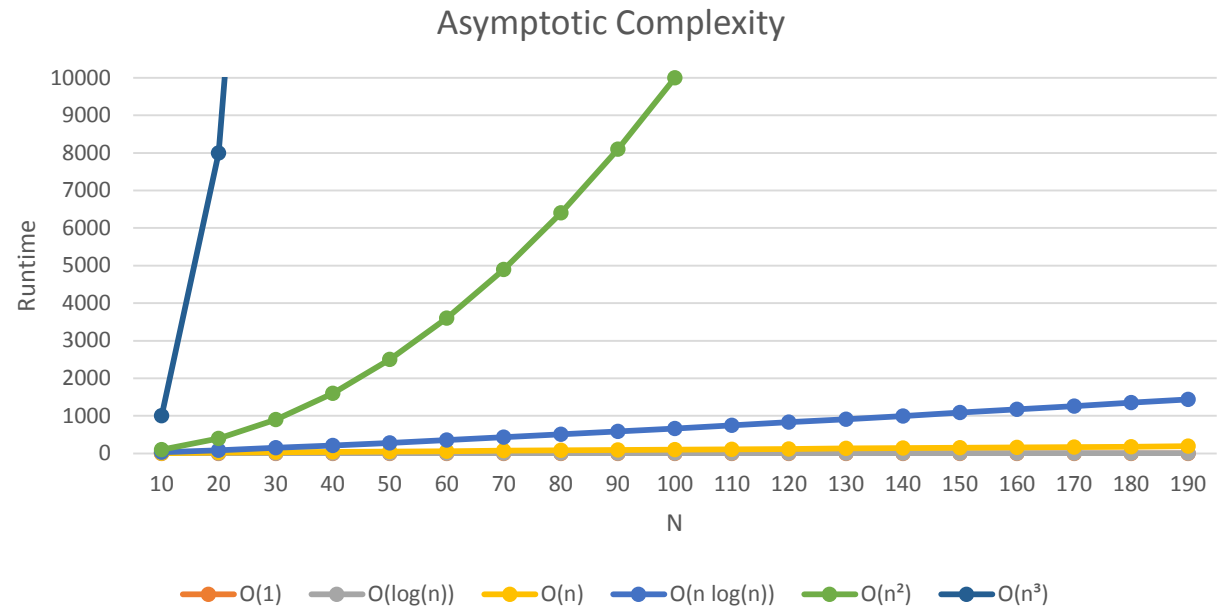
```

times	Best case ($t_j = 1$)	Worst case ($t_j = j$)
n	n	n
n-1	n-1	n-1
n-1	n-1	n-1
$\sum_{j=1}^{n-1} t_j$	n-1	$\frac{n(n+1)}{2} - 1$
$\sum_{j=1}^{n-1} (t_j - 1)$	0	$\frac{n(n-1)}{2}$
$\sum_{j=1}^{n-1} (t_j - 1)$	0	$\frac{n(n-1)}{2}$
n-1	n-1	n-1
	O(n)	O(n²)

Asymptotic Complexity & O-Notation

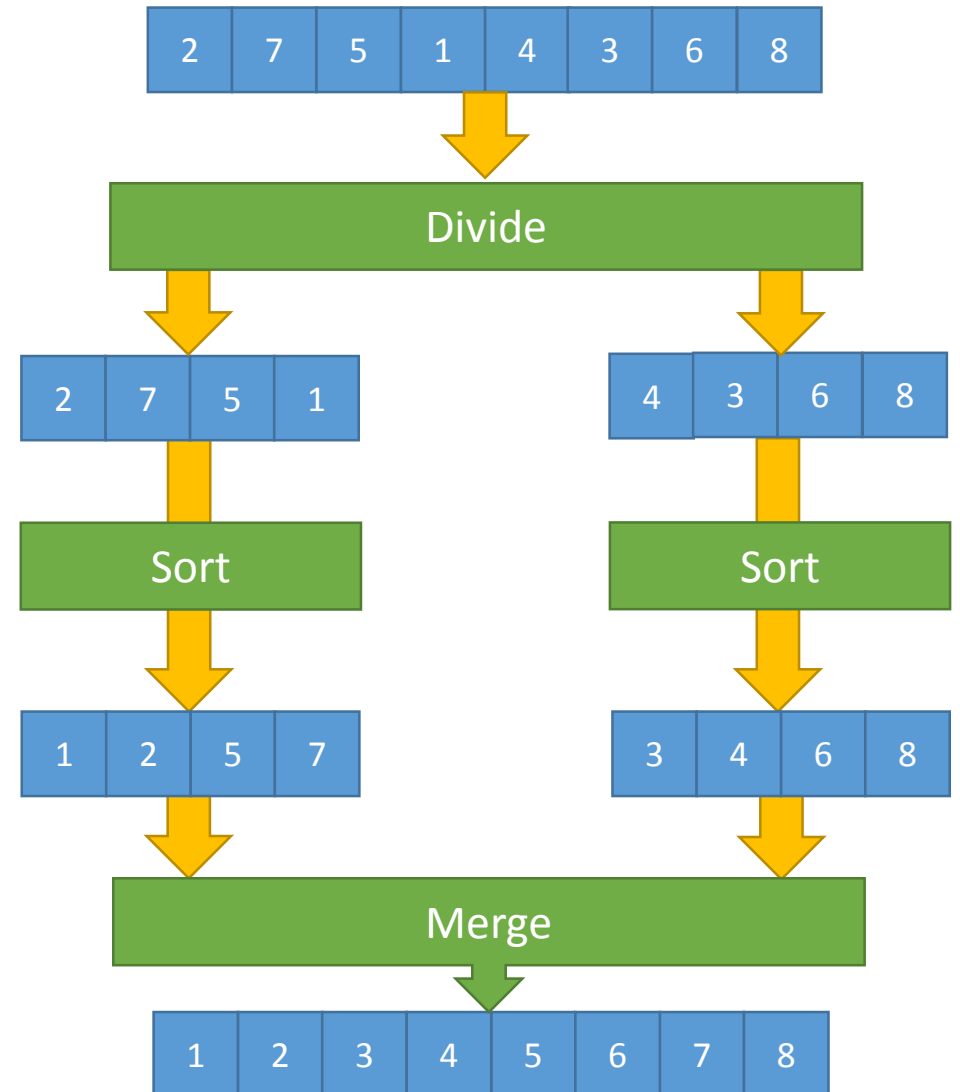
- $O(1)$... constant
- $O(\log n)$... logarithmic
- $O(n)$... linear
- $O(n \log n)$... quasi-linear
- $O(n^2)$... quadratic
- $O(n^3)$... cubic
- $O(c^n)$... exponential

n	$O(1)$	$O(\log(n))$	$O(n)$	$O(n \log(n))$	$O(n^2)$	$O(n^3)$
10	const	3	10	33	100	1000
1000	const	10	1000	9966	1E+06	1E+09
100000	const	17	100000	1660964	1E+10	1E+15
1000000	const	20	1000000	19931569	1E+12	1E+18



Merge Sort

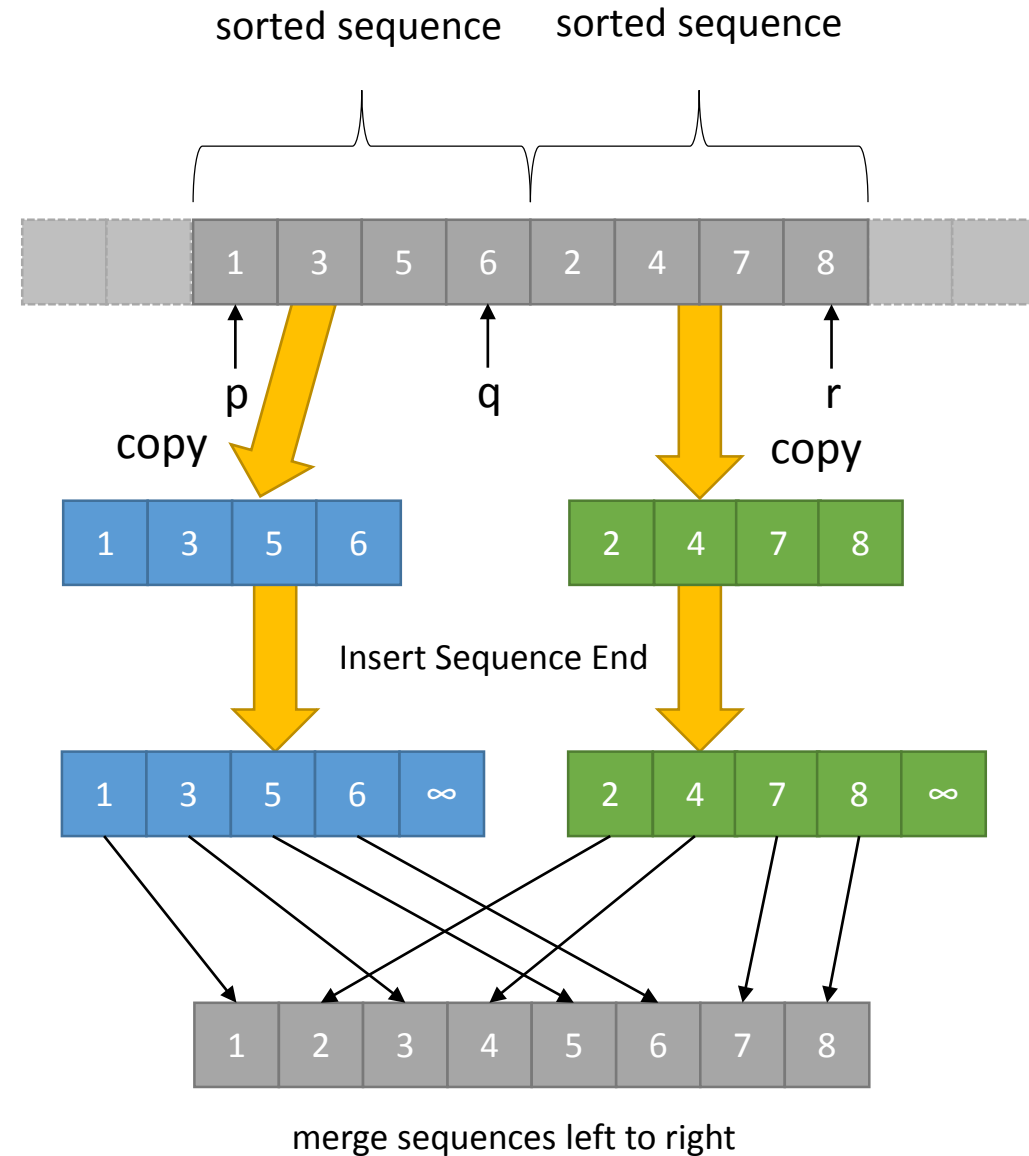
- Divide-And-Conquer Approach
- Subdivide sequence into smaller subsequences
- Sort each subsequence
- Merge ordered subsequences to final sequence



Merge Sort

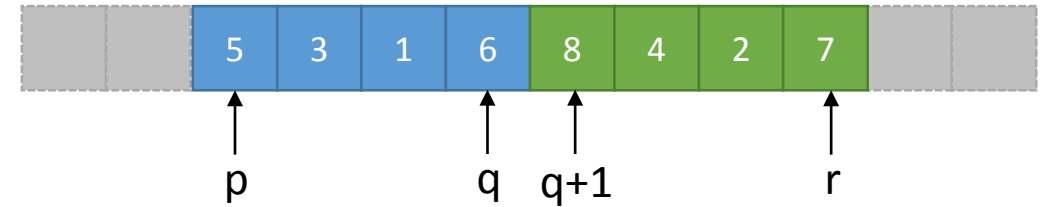
```
void merge(vector<int> & v, int p, int q, int r) {  
    vector<int> left(v.begin()+p, v.begin()+q+1);  
    vector<int> right(v.begin()+q+1, v.begin()+r+1);  
  
    left.push_back(/* infinity */);  
    right.push_back(/* infinity */);  
    int i = 0;  
    int j = 0;  
  
    for(int k = p; k <= r; k++) {  
        if(left[i] <= right[j])  
            v[k] = left[i++];  
        else  
            v[k] = right[j++];  
    }  
}
```

$O(n)$



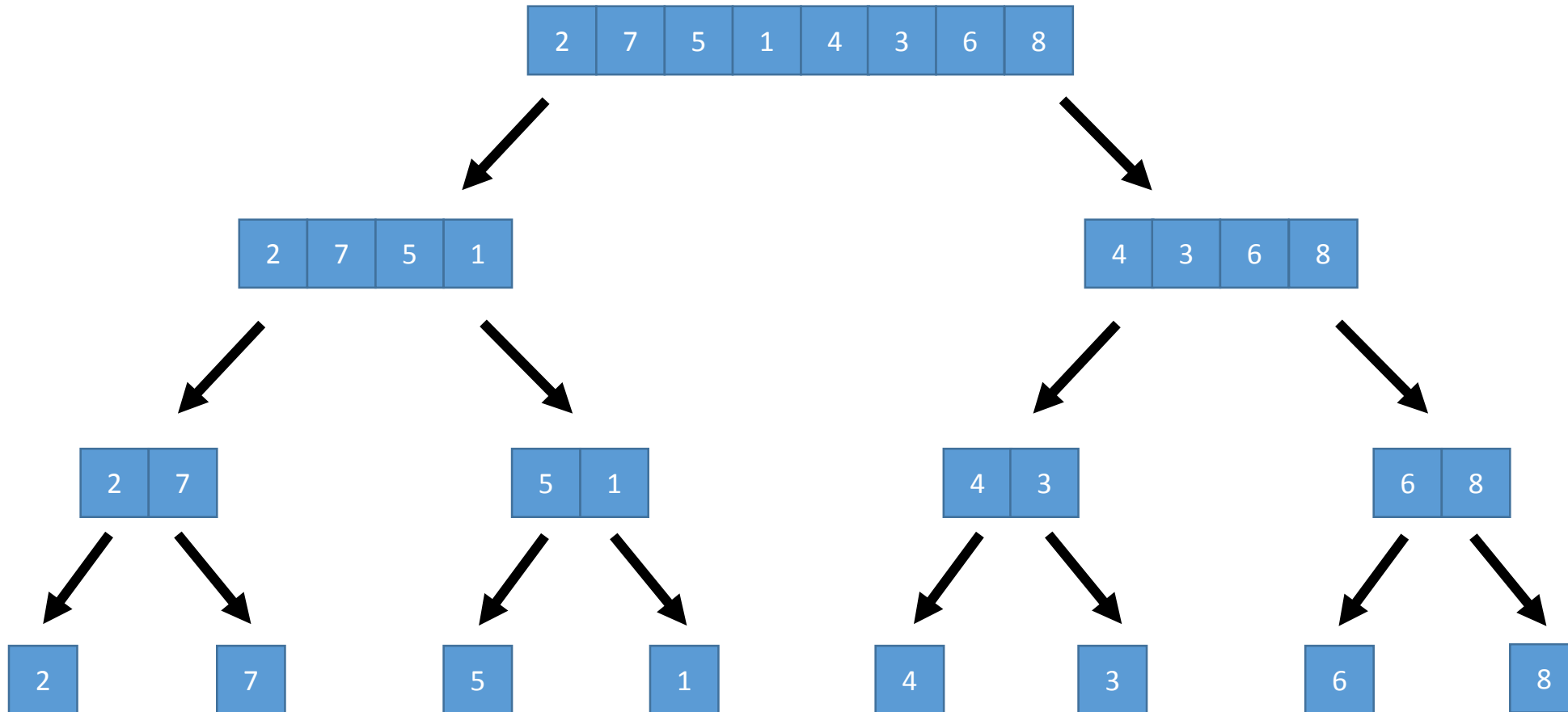
Merge Sort

```
void merge_sort(std::vector<int> & v, int p, int r) {  
    if(p < r) {  
        int q = (p+r)/2;  
        merge_sort(v, p, q);  
        merge_sort(v, q+1, r);  
        merge(v, p, q, r);  
    }  
}
```

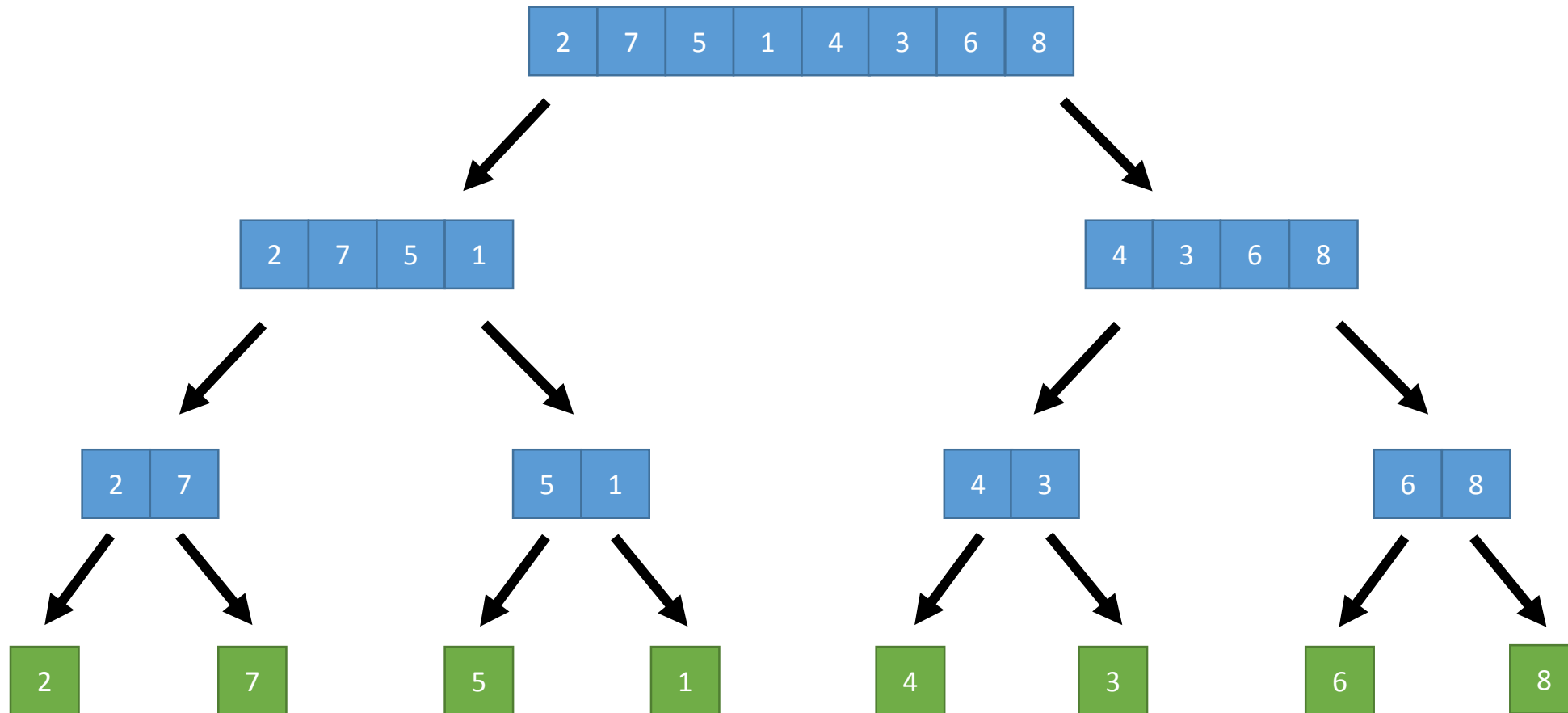


```
void merge_sort(std::vector<int> v) {  
    merge_sort(v, 0, v.size());  
}
```

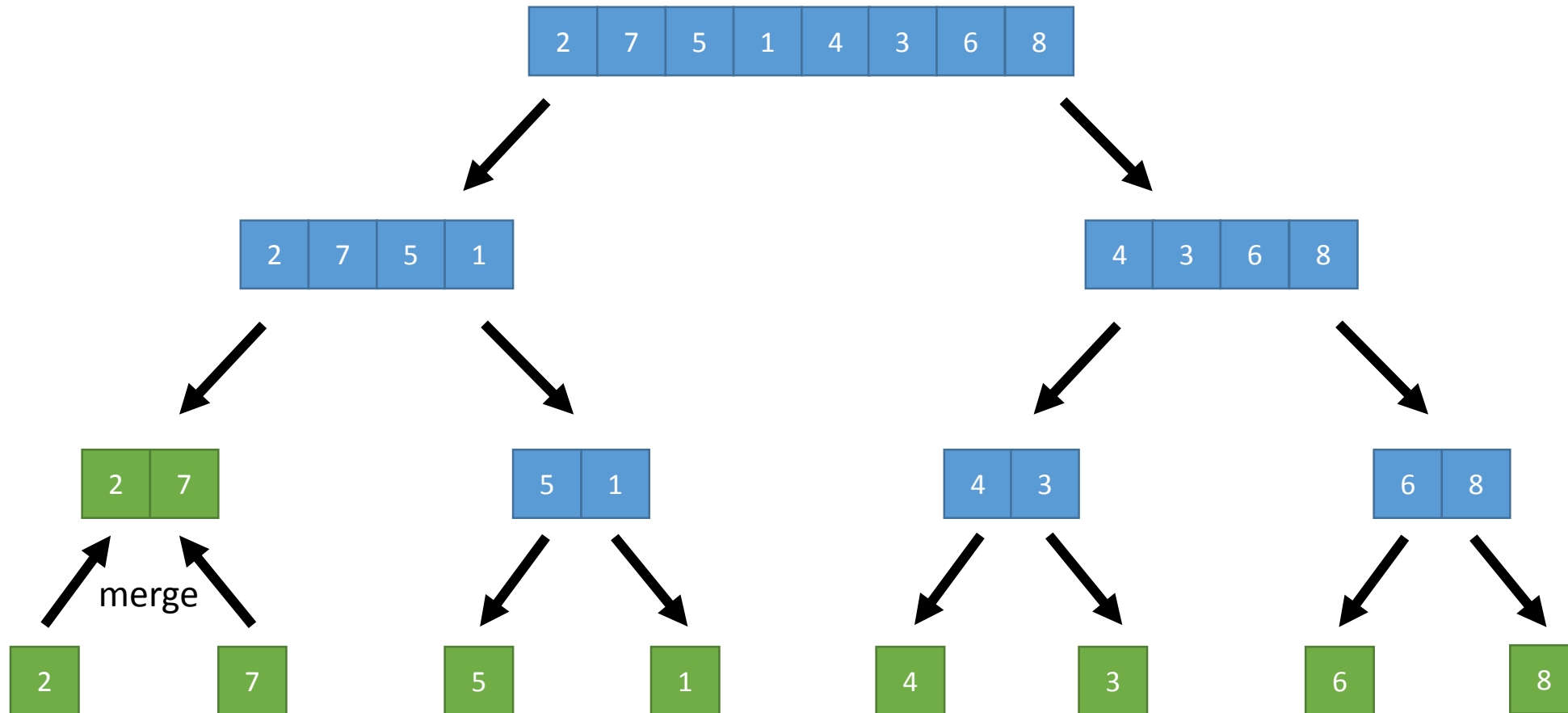
Merge Sort



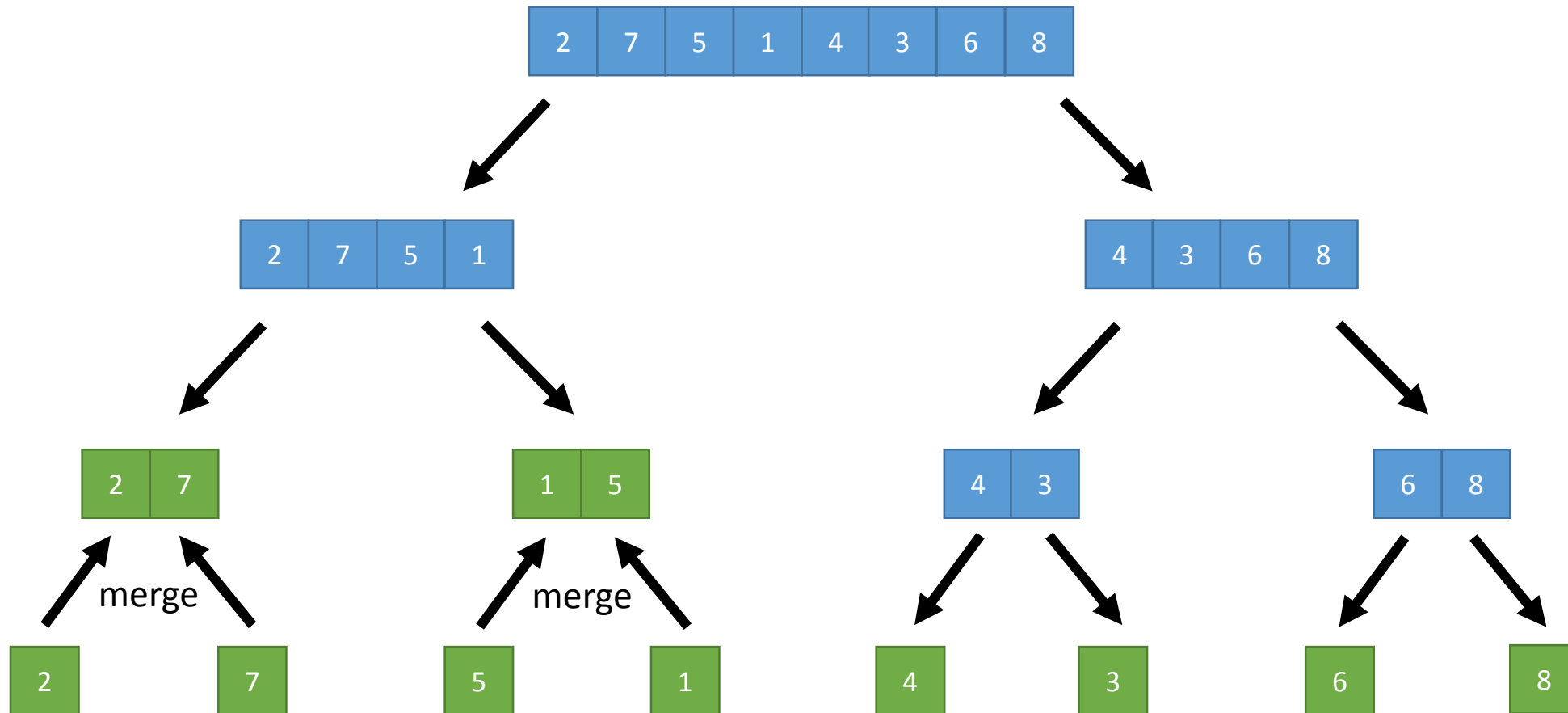
Merge Sort



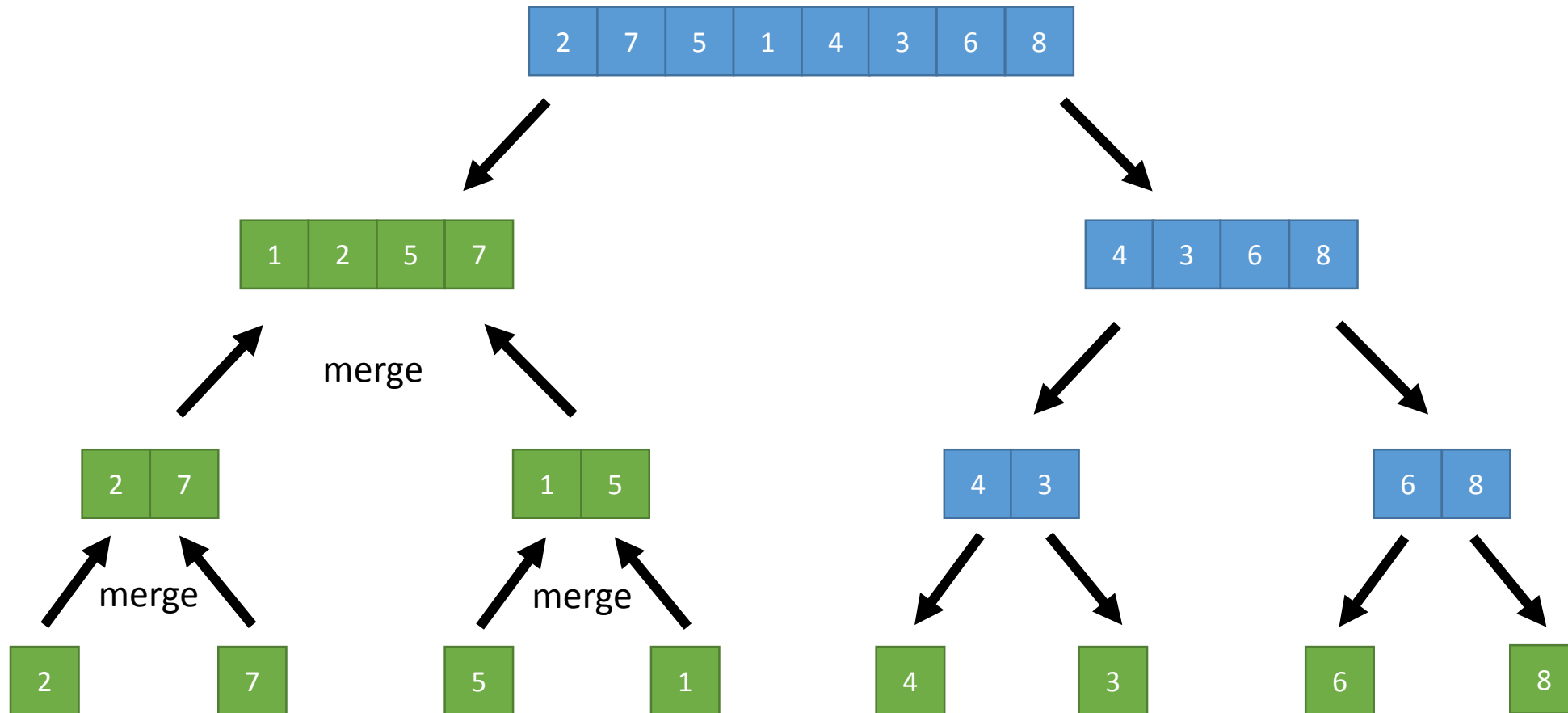
Merge Sort



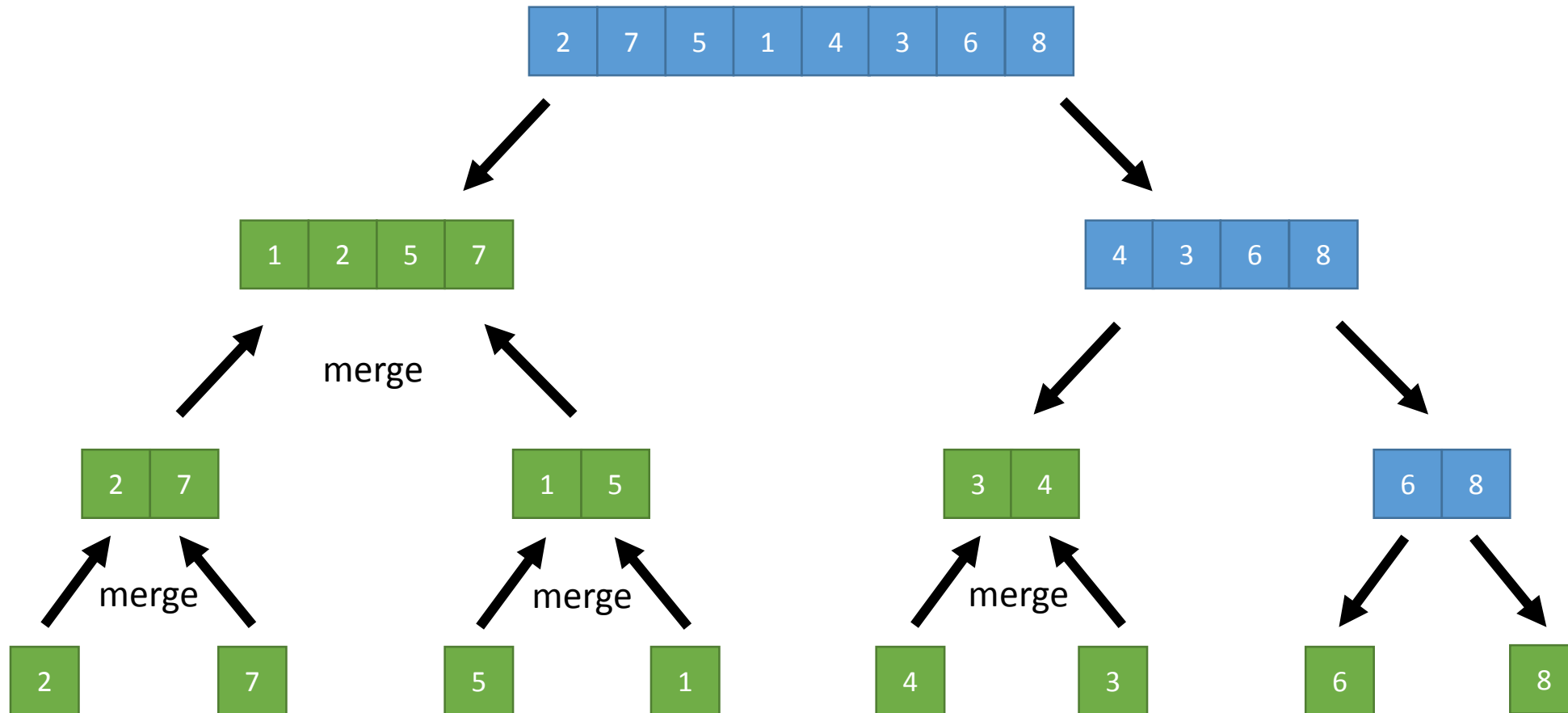
Merge Sort



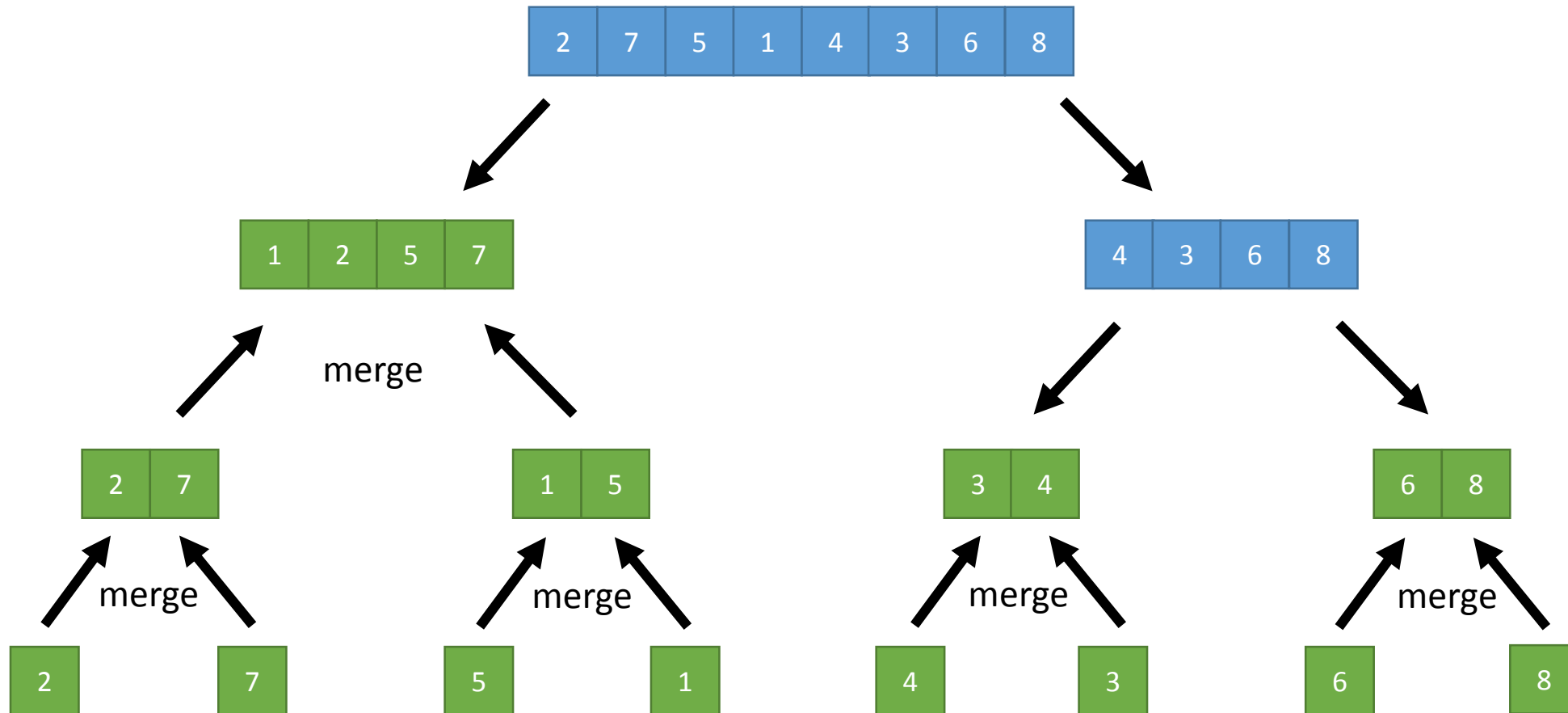
Merge Sort



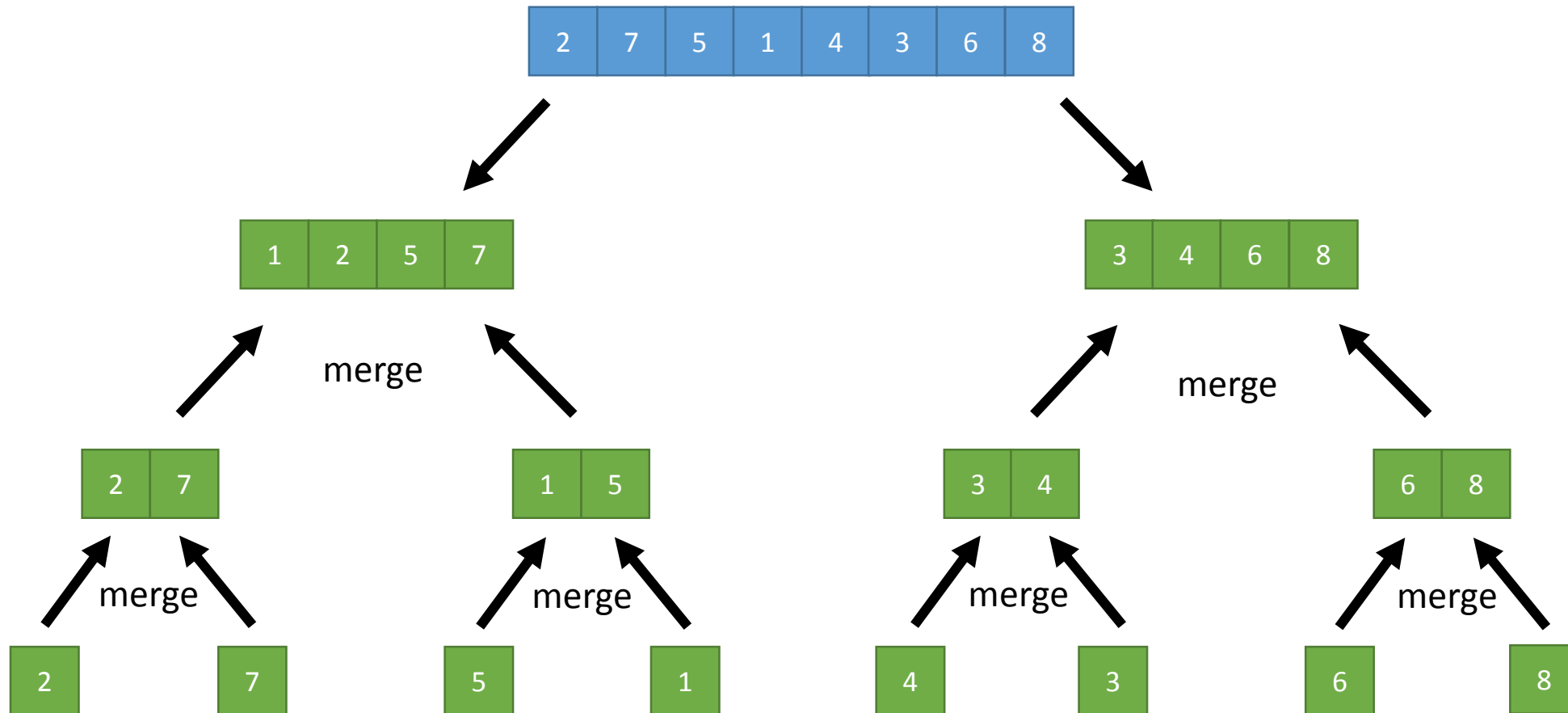
Merge Sort



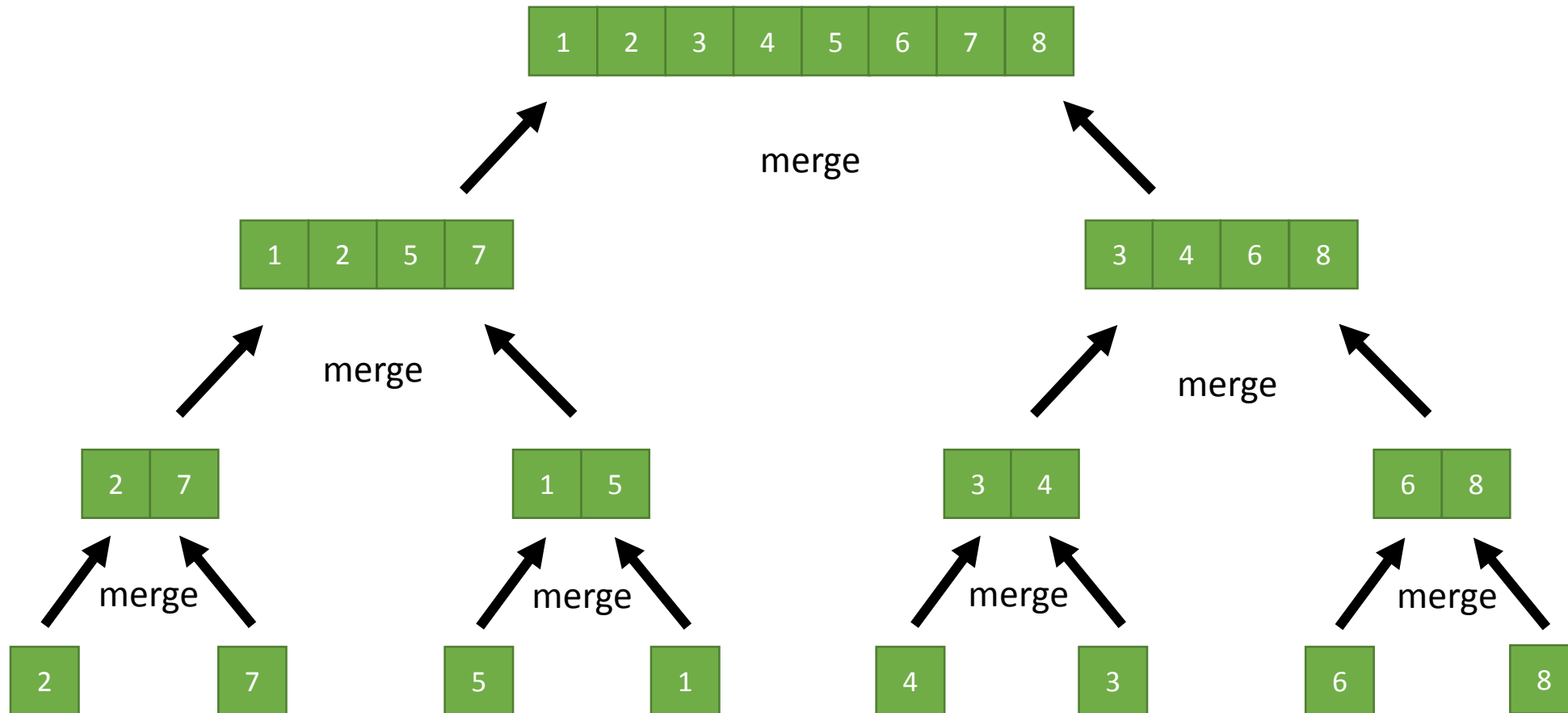
Merge Sort



Merge Sort

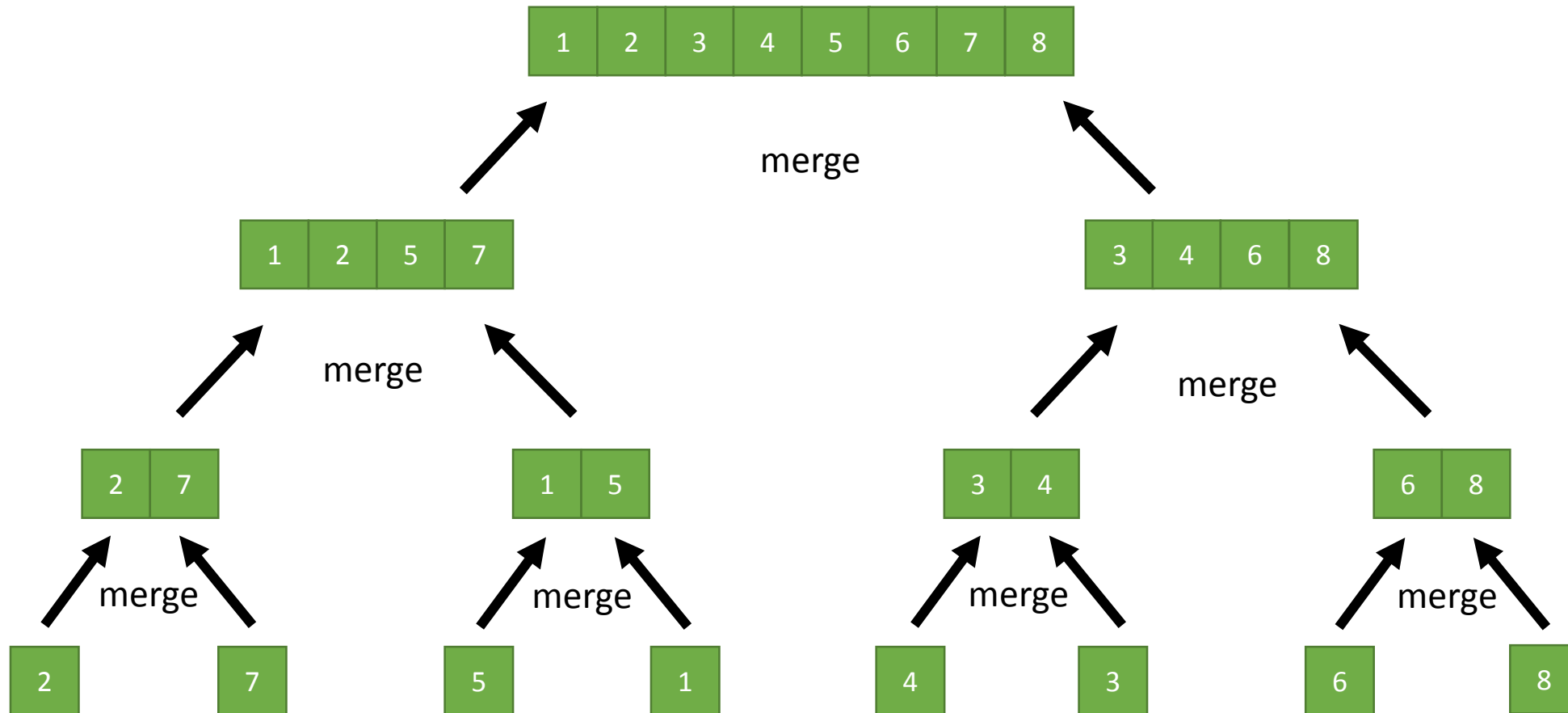


Merge Sort



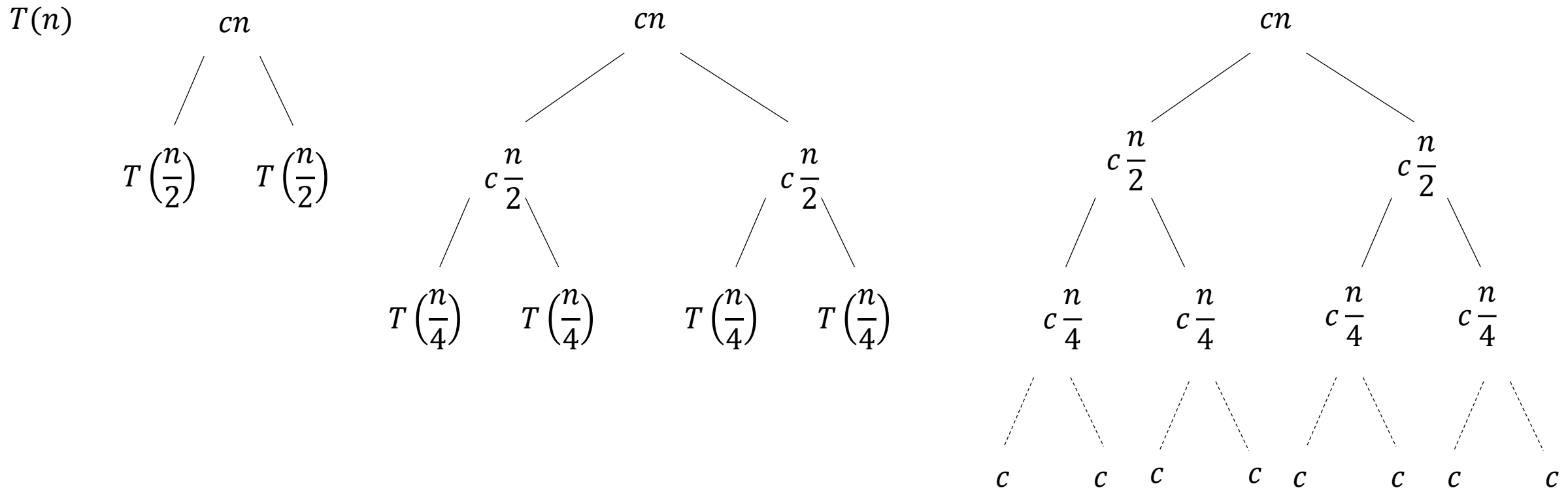
Merge Sort

$$T(n) = \begin{cases} c & \text{if } n = 1 \\ 2T\left(\frac{n}{2}\right) + cn & \text{if } n > 1 \end{cases}$$



Merge Sort – Recursive Tree

$$T(n) = \begin{cases} c & \text{if } n = 1 \\ 2T\left(\frac{n}{2}\right) + cn & \text{if } n > 1 \end{cases}$$



Merge Sort - Complexity

$$T(n) = \begin{cases} c & \text{if } n = 1 \\ 2T\left(\frac{n}{2}\right) + cn & \text{if } n > 1 \end{cases}$$

$$T(n) = cn(\log_2 n + 1)$$

$$O(n \log n)$$

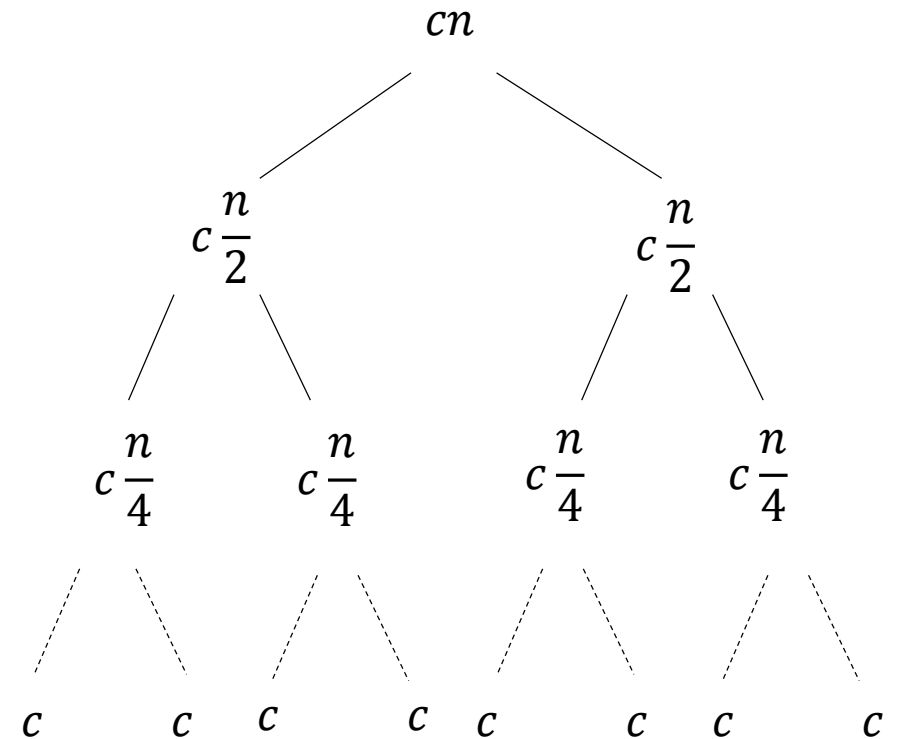
$$\log_2 n + 1$$

cn

cn

cn

cn



Comparing Algorithms

Theoretically:

- by **runtime complexity**
 - Insertion Sort $O(n^2)$ slower than Merge Sort $O(n \log n)$
- by **memory usage**
 - Insertion Sort can be implemented as in-place algorithm (work on same memory block)
 - Merge Sort creates copies of sub-sequences → 2X maximum memory usage

Practically:

- by **measurement (profiling)**

Sorting

- Quick Sort:
 - Worst case: $O(n^2)$
 - Best case: $O(n \log n)$
 - Average case: $O(n \log n)$
- Heap Sort:
 - Worst case: $O(n \log n)$
 - Best case: $O(n \log n)$
 - Average case: $O(n \log n)$

Visualization of Sorting

- 15 Sorting Algorithms in 6 Minutes
<http://youtu.be/kPRA0W1kECg>

Keep this in mind when optimizing

Finding a **better** algorithm $>$ Optimizing an algorithm

$O(n)$ beats $O(n^2)$

$O(1)$ beats $O(n)$

...

Generic Programming

- Defining functions and classes which work with many different types

C++: strict type system

```
void sort(int* ar, int length);  
void sort(float* ar, int length);  
void sort(double* ar, int length);
```

C++ templates:

allow you to write generalized algorithms that can work with any type

```
template<typename T>  
void sort(T* ar, int length);
```

```
template<typename T>  
class SomeClass {  
    T memberVariable;  
}
```

Python: duck typing

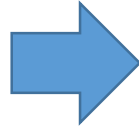
```
def sort(ar, length):  
    ...
```

Python doesn't care what types come in as parameters. It will try to use them. Type checking is done at runtime.

"If it quacks like a duck, looks like a duck, then it is a duck."

Generic Algorithms

```
void insertion_sort(std::vector<int> & v) {  
    for(int j = 1; j < v.size(); j++) {  
        int key = v[j];  
        int i = j - 1;  
        while(i >= 0 && v[i] > key) {  
            v[i+1] = v[i];  
            i--;  
        }  
        v[i+1] = key;  
    }  
}
```



```
template<typename T>  
void insertion_sort(std::vector<T> & v) {  
    for(int j = 1; j < v.size(); j++) {  
        T key = v[j];  
        int i = j - 1;  
        while(i >= 0 && v[i] > key) {  
            v[i+1] = v[i];  
            i--;  
        }  
        v[i+1] = key;  
    }  
}
```

Usage:

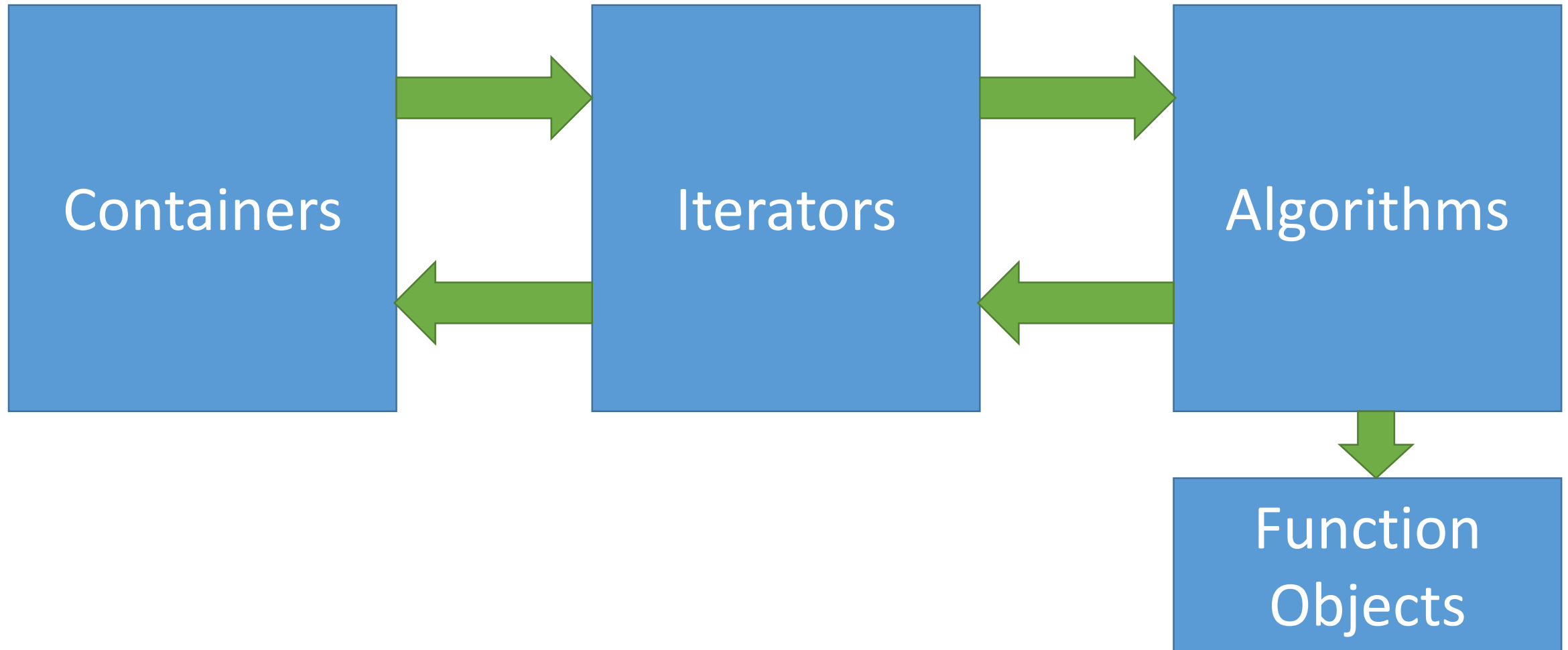
```
std::vector<int> i_list;  
std::vector<double> fp_list;  
insertion_sort(i_list);  
insertion_sort(fp_list);
```

C++ Standard Template Library

- Common Algorithms and Data Structures
- Optimized for General Computing
- Uses C++ template mechanisms extensively
- No inheritance, no virtual calls

C++ Standard Template Library (STL)

Part of the C++ Standard Library



Iterators

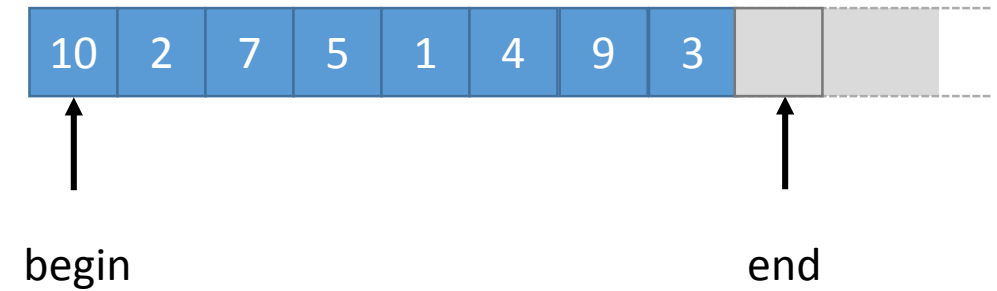
„Generalized Pointers“

```
int myArray[100];

int * b = &myArray[0];
int * e = &myArray[0] + 100;

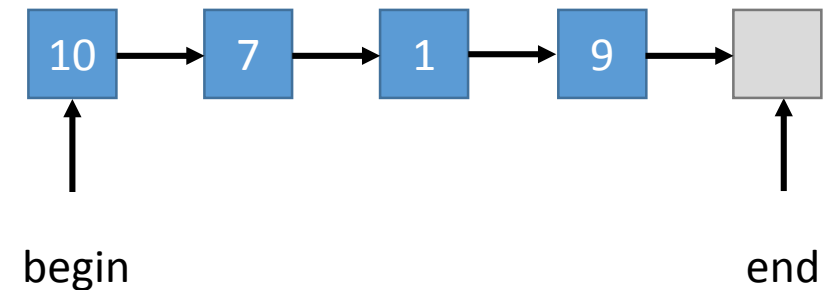
for(int i = 0; i < 100; ++i) {
    myArray[i] = /* calculation */;
}

for(int * a = b; a != e; ++a) {
    *a = /* calculation */;
}
```



```
vector<int> v(100, 0);

for(vector<int>::iterator it = v.begin(); it != v.end(); ++it)
    *it = /* calculation */;
}
```



Function Objects

- Python

- any function can be used as an object
- lambda functions: anonymous functions (functions without a name)

```
def less_than(a, b):  
    return a < b
```

```
# store function in a variable  
a_variable = less_than
```

```
# call function through variable  
a_variable(a, b)
```

- C++

- Something like C-callbacks, but it can store its state
- Definition of a class which implements special operator() member function
- C++11 lambda functions

```
struct less_than {  
    bool operator()(int a, int b) {  
        return a < b;  
    }  
}
```

```
less_than a_object;  
// use that object like a function  
bool result = a_object(10, 20);
```

std::sort and std::stable_sort

stable_sort ensures the order of objects which are equal is not changed after sorting.

```
#include <vector>
#include <list>
#include <algorithm> // sort
#include <functional> // function objects
```

```
std::vector<int> myVector;
```

```
// sort in ascending order
```

```
std::sort(myVector.begin(), myVector.end());
```

Function Object

```
// equivalent
```

```
std::sort(myVector.begin(), myVector.end(), std::less_than())
```

```
// descending order
```

```
std::sort(myVector.begin(), myVector.end(), std::greater_than())
```


STL Algorithms

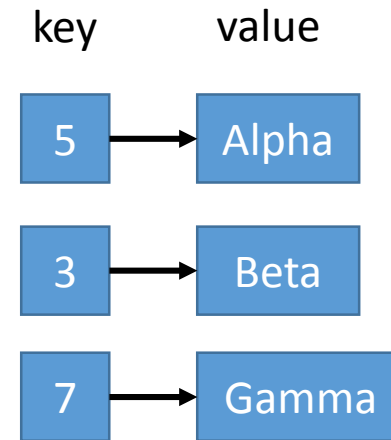
- `sort`
- `stable_sort`
- `swap`
- `find_if`
- `any`
- `rotate`
- `lower_bound`

Data Structures



Sequence

e.g., C-Arrays, `std::vector`, `std::deque`, `std::list`

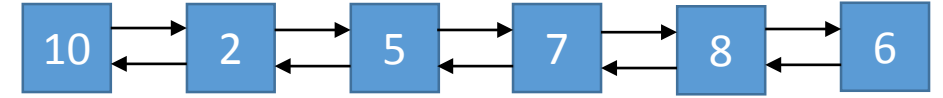


Associative

e.g., C-Arrays, `std::map`, `std::set`, `std::unordered_map`

Data Structures

- Operations:
 - Insertion
 - Searching
 - Deletion
- Variants:
 - Ordered
 - Unordered



Sequential Containers

Arrays, Lists, Queues, Stacks

STL Containers

- Sequence Containers
 - vector (flexible sequence)
 - deque (double-ended queue)
 - list (double linked list)
 - array (fixed sequence, C++11)
 - forward_list (single linked list, C++11)

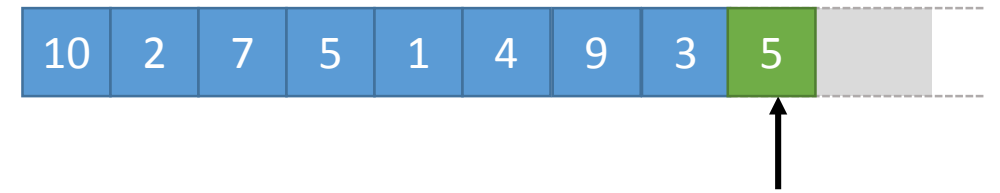
C-Arrays

- Simplest Sequence Data Structure
- Data stored in range $[0, \text{numElements})$
- Fixed Size, Wasteful
- Consecutive Memory (efficient access)

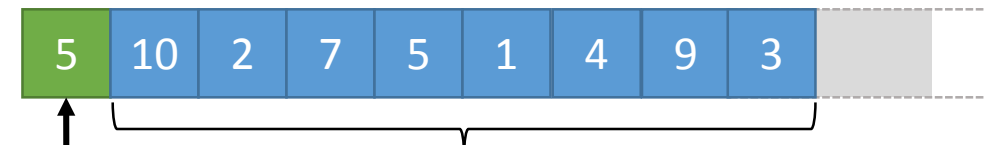
```
int a[10000];  
int numElements = 0;
```

```
// insertion at end  $O(1)$   
a[numElements++] = new_value;
```

```
// insertion at beginning  $O(n)$   
for(int i = numElements; i > 0; i--) a[i] = a[i-1];  
a[0] = new_value;  
numElements++;
```

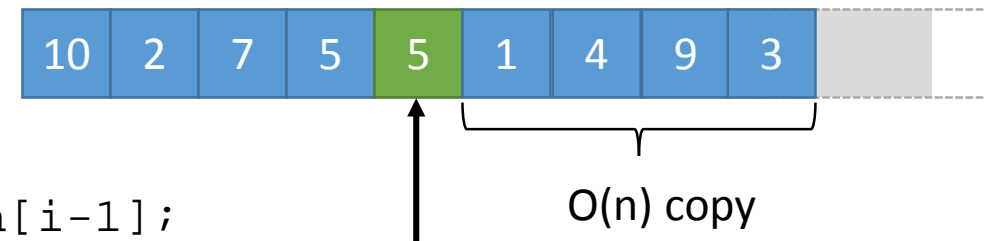


Inserting at end is $O(1)$



$O(n)$ copy of previous values to new location

Therefore inserting at beginning is $O(n)$



inserting in the middle is $O(n)$

std::vector

```
#include <iostream>
#include <vector>

using namespace std;

// empty construction
vector<int> a;
// sized construction
vector<int> a(10);
// sized construction with initial value
vector<int> a(100, -1);
// C++ 11 initializer lists
vector<int> a { 3, 5, 7, 9, 11 };

// insertion at end
a.push_back(3);
a.push_back(5);
a.push_back(7);

// delete at end
a.pop_back();

// insertion at beginning
a.insert(a.begin(), new_value);
```

```
// accessing elements just like arrays
for(int i = 0; i < a.size(); i++) {
    cout << a[i] << endl;
}

// using iterators
for(auto i = a.begin(); i != a.end(); ++i) {
    cout << *i << endl;
}

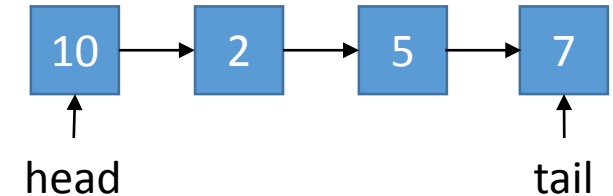
// C++11 for each
for(auto element : a) {
    cout << element << endl;
}
```

Linked-List

- List Elements connected through pointers
- First Element (head) and last element (tail) are always known
- Insertion/Deletion at **both** ends in $O(1)$
- Insertion in the middle is also cheaper
 - Finding insertion location is $O(n)$ compared to $O(1)$ with C-Arrays
 - But insertion itself happens in $O(1)$ instead of $O(n)$ copies
- Dynamic Size
- Distributed in memory

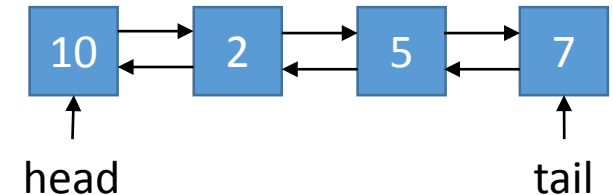
Single Linked-List:

only pointer of next element



Double Linked-List:

pointer of previous and next element



```
struct Node {  
    Node * prev;  
    Node * next;  
    int data;  
}
```


std::list

```
#include <iostream>
#include <list>

using namespace std;

// empty construction
list<int> a;
// sized construction
list<int> a(10);
// sized construction with initial value
list<int> a(100, -1);
// C++ 11 initializer lists
list<int> a { 3, 5, 7, 9, 11 };
```

```
// insertion at beginning
a.push_front(3);
```

```
// insertion at end
a.push_back(3);
```

```
// delete at beginning
a.pop_front();
```

```
// delete at end
a.pop_back();
```

```
// access front element
int first = a.front();
```

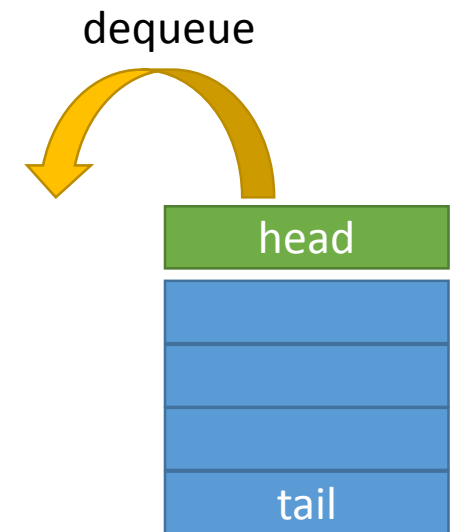
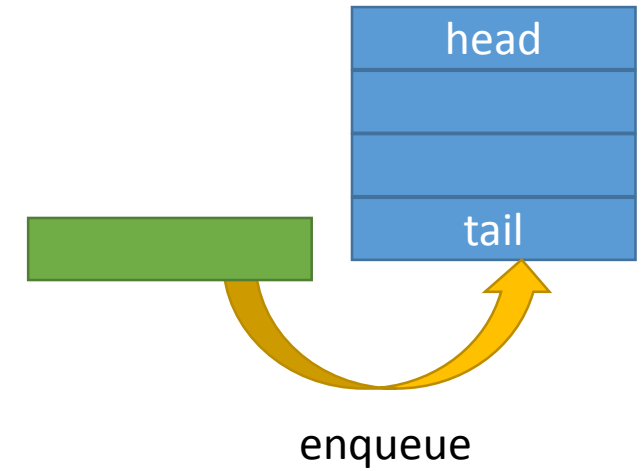
```
// access last element
int last = a.back();
```

```
// using iterators
for(auto i = a.begin(); i != a.end(); ++i) {
    cout << *i << endl;
}
```

```
// C++11 for each
for(auto element : a) {
    cout << element << endl;
}
```

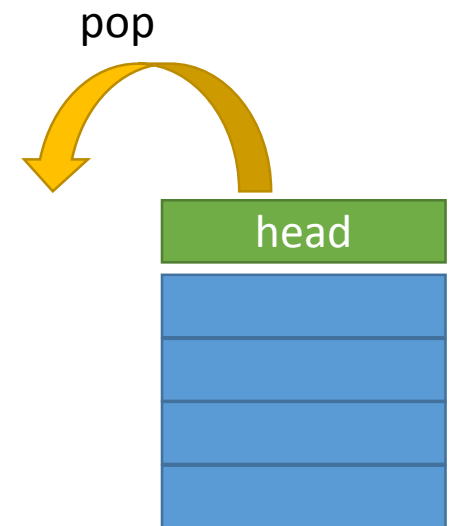
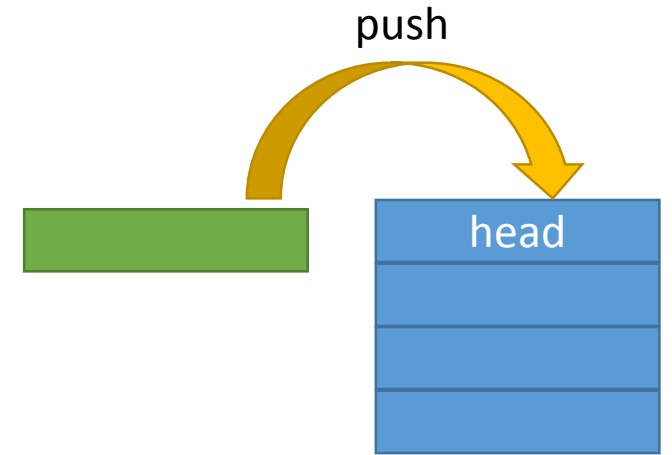
Queue

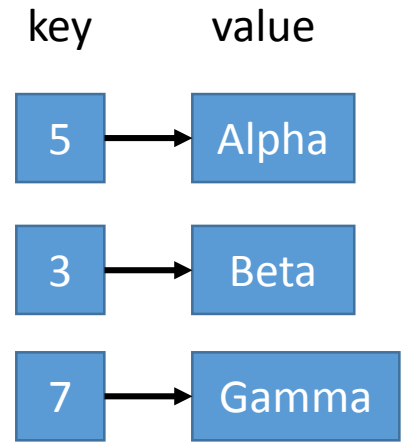
- First-In-First-Out (FIFO) data structure
- Implementations:
 - Double-Linked-List
- Operations:
 - **enqueue**: put element in queue (insert at tail)
 - **dequeue**: get first element in queue (remove head)



Stack

- Last-In-First-Out (LIFO) data structure
- Implementations:
 - C-Array
 - Single-Linked-List
- Operations:
 - **push**: put element on stack (insert as first element)
 - **pop**: get first element on stack (remove head)





Associative Containers

Dictionaries, Maps, Sets

Associative Containers

- Map a key to a value
- Searching for a specific element in unsorted sequential containers takes **linear** time $O(n)$
- Getting a specific element from an associative container can be as fast as **constant** time $O(1)$

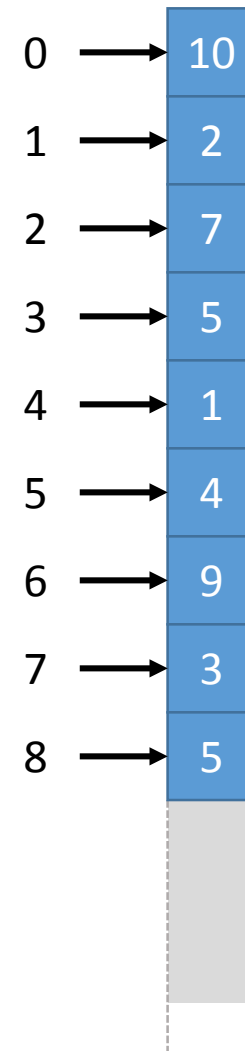
STL Containers

- Associative Containers
 - map
 - set
 - multimap
 - multiset
- unordered_map (C++11)
- unordered_set (C++11)
- unordered_multimap (C++11)
- unordered_multiset (C++11)

C-Array as Associative Container

- Simplest associative data structure
- maps **integer number** to data
 - 0 -> a[0]
 - 1 -> a[1]
 - ...
- **efficient access in $O(1)$**
- **inefficient storage**
- **limited to positive integer numbers as keys**

```
int a[10000];
```

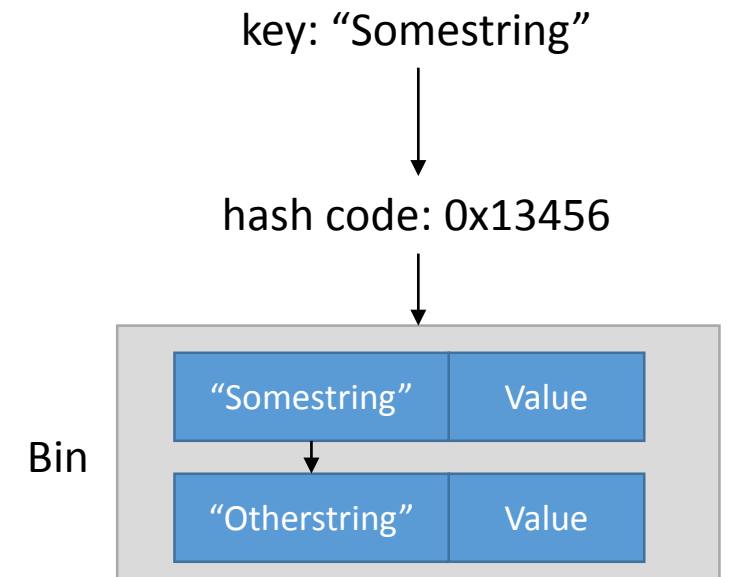


Unordered maps / Hash maps

- Maps **arbitrary keys** (objects, basic types) to **arbitrary values** (objects, basic types)
- **On average accessing a hash map through keys takes $O(1)$**
- In general unordered structure - you can't get out objects in the same order you inserted them.
- a number, called a **hash code**, is generated using a **hash function** based on key in $O(1)$
- Each hash code can be mapped to a location called a bin
- A bin stores nodes with keys which map to the same hash code
- Lookup therefore consists of:
 - Determining the hash code of the key $O(1)$
 - Selecting the correct node inside the bin is in the worst case $O(n)$

On average lookup times are $O(1)$. But this is only true if there are only few hash collisions.

Hash maps require a **good hashing function**, which reduces the amount of hash collisions.



Ordered maps

- Maps **arbitrary keys** (objects, basic types) to **arbitrary values** (objects, basic types)
- Basic idea: if keys are sortable, we can store nodes in a data structure sorted by its keys. Sorted data structures can be searched more quickly, e.g. with binary search in $O(\log(n))$
- Elements ordered by key
- **Worst case lookup time is $O(\log(n))$**

std::map

```
#include <iostream>
#include <map>
#include <string>

using namespace std;

map<string, string> capitals;

// setting value for key
capitals["Austria"] = "Vienna";
capitals["France"] = "Paris";
capitals["Italy"] = "Rome";

// getting value from key
cout << "Capital of Austria: " << capitals["Austria"] << endl;
string & capital_of_france = capitals["France"];
cout << "Capital of France: " << capitals << endl;

// check if key is set
if (capitals.find("Spain") != capitals.end()) {
    cout << "Capital of Spain is " << capitals["Spain"] << endl;
} else {
    cout << "Capital of Spain not found!" << endl;
}
```

std::map

```
// iterate over all elements
for (map<string, string>::iterator it = capitals.begin(); it != capitals.end(); ++it) {
    string & key = it->first;
    string & value = it->second;
    cout << "The capitol of " << key << " is " << value << endl;
}
```

```
// C++11: iterate over all elements
for (auto it = capitals.begin(); it != capitals.end(); ++it) {
    string & key = it->first;
    string & value = it->second;
    cout << "The capitol of " << key << " is " << value << endl;
}
```

```
// C++11: iterate over all elements
for (auto & kv : capitals) {
    string & key = kv.first;
    string & value = kv.second;
    cout << "The capitol of " << key << " is " << value << endl;
}
```

C++ Resources

- (Unofficial) C++ STL online reference:
<http://www.cplusplus.com/reference/stl/>
- C++11 working draft:
<http://www.open-std.org/jtc1/sc22/wg21/docs/papers/2011/n3242.pdf>