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# Numerical Analysis for Machine Learning

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These are unreviewed notes and may contain errors.  
Made by Roberto Benatuil Valera**



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## Chapter 1

# Numerical Linear Algebra tools

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### 1.1 Introduction: Recap of Linear Algebra

In this section we will review some basic concepts of Linear Algebra that will be useful for the rest of the course.

#### 1.1.1 Matrix-vector multiplication

Given a matrix  $A \in \mathbb{R}^{m \times n}$  and a vector  $x \in \mathbb{R}^n$ , the matrix-vector multiplication  $y = Ax$  is defined as:

$$y_i = \sum_{j=1}^n A_{ij}x_j \quad (1.1)$$

A matrix-vector multiplication can be considered as a linear combination of the columns of the matrix  $A$ . Lets see an example:

$$\begin{bmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix} x_1 + \begin{bmatrix} 4 \\ 5 \\ 6 \end{bmatrix} x_2 \quad (1.2)$$

#### 1.1.2 Column space of a matrix

The column space of a matrix  $A \in \mathbb{R}^{m \times n}$  is the subspace of  $\mathbb{R}^m$  spanned by the columns of  $A$ . In other words, it is the set of all possible linear combinations of the columns of  $A$ . The column space of a matrix is denoted as  $C(A)$ .

If the columns of  $A$  are linearly independent, then the column space of  $A$  is the entire  $\mathbb{R}^m$ . If the columns of  $A$  are linearly dependent, then the column space of  $A$  is a subspace of  $\mathbb{R}^m$  with dimension equal to the rank of  $A$ .

The rank of a matrix  $A$  is the size of the largest set of linearly independent columns of  $A$ . It is denoted as  $rank(A)$ . Note that  $rank(A) = rank(A^T)$ .

### 1.1.3 System of linear equations

A system of linear equations is a set of  $m$  equations with  $n$  unknowns of the form:

$$\begin{aligned} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n &= b_1 \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n &= b_2 \\ &\vdots \\ a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n &= b_m \end{aligned} \tag{1.3}$$

This system can be written in matrix form as  $Ax = b$ , where  $A \in \mathbb{R}^{m \times n}$ ,  $x \in \mathbb{R}^n$  and  $b \in \mathbb{R}^m$ .

The system  $Ax = b$  has a solution if and only if  $b \in C(A)$ . If  $b \in C(A)$ , then the system has a unique solution if and only if  $\text{rank}(A) = n$ . If  $\text{rank}(A) < n$ , then the system has infinitely many solutions.

### 1.1.4 CR factorization

The CR factorization of a matrix  $A \in \mathbb{R}^{m \times n}$ , with  $m \geq n$ , is a factorization of  $A$  as  $A = CR$ , where  $C \in \mathbb{R}^{m \times r}$  is a matrix with the linearly independent columns of  $A$  and  $R \in \mathbb{R}^{r \times n}$  is obtained by determining the coefficients of the linear combination of the columns of  $C$  that give the columns of  $A$ . In this factorization,  $r = \text{rank}(A)$ .

Lets see an example:

$$A = \begin{bmatrix} 1 & 4 & 7 \\ 2 & 5 & 8 \\ 3 & 6 & 9 \end{bmatrix} = \begin{bmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{bmatrix} \begin{bmatrix} 1 & 0 & -1 \\ 0 & 1 & 2 \end{bmatrix} = CR \tag{1.4}$$

The matrix  $C$  is also called the Row Reduced Echelon Form of  $A$ .

### 1.1.5 Matrix-matrix multiplication

Given two matrices  $A \in \mathbb{R}^{m \times n}$  and  $B \in \mathbb{R}^{n \times p}$ , the matrix-matrix multiplication  $C = AB$  is defined as:

$$C_{ij} = \sum_{k=1}^n A_{ik}B_{kj} \tag{1.5}$$

A matrix-matrix multiplication can be considered as the outer product of the columns of  $A$  and the rows of  $B$ . Lets see an example:

$$\begin{bmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix} \begin{bmatrix} 1 & 2 \end{bmatrix} + \begin{bmatrix} 4 \\ 5 \\ 6 \end{bmatrix} \begin{bmatrix} 3 & 4 \end{bmatrix} \tag{1.6}$$

Note that each outer product generates a matrix of the same size as the result matrix, but always with rank 1. So the matrix-matrix multiplication can be considered as a sum of rank 1 matrices, obtained by the outer products of the columns of  $A$  and the rows of  $B$ .



### 1.1.6 Null space of a matrix

The null space of a matrix  $A \in \mathbb{R}^{m \times n}$  is the set of all vectors  $x \in \mathbb{R}^n$  such that  $Ax = 0$ . The null space of a matrix is denoted as  $N(A)$ . It is also called the kernel of  $A$ , denoted as  $\ker(A)$ .

Formally, we have that:

$$N(A) = \{x \in \mathbb{R}^n : Ax = 0\} \quad (1.7)$$

The null space of a matrix is a subspace of  $\mathbb{R}^n$ . The dimension of the null space of a matrix is called the nullity of the matrix.

### 1.1.7 Fundamental subspaces of a matrix

Given a matrix  $A \in \mathbb{R}^{m \times n}$ , we can define four fundamental subspaces:

- The column space of  $A$ , denoted as  $C(A)$
- The row space of  $A$ , denoted as  $C(A^T)$
- The null space of  $A$ , denoted as  $N(A)$
- The left null space of  $A$ , denoted as  $N(A^T)$

These subspaces are related by the following properties:

$$\begin{aligned} C(A) &\perp N(A^T) \\ C(A^T) &\perp N(A) \end{aligned} \quad (1.8)$$

They also satisfy the following dimensions properties:

$$\begin{aligned} \dim(C(A)) + \dim(N(A)) &= n \\ \dim(C(A^T)) + \dim(N(A^T)) &= m \end{aligned} \quad (1.9)$$

This is known as the Rank-Nullity Theorem.

### 1.1.8 Orthogonal matrices

An orthogonal matrix is a square matrix  $Q \in \mathbb{R}^{n \times n}$  such that  $Q^T Q = I$ , where  $I$  is the identity matrix. This implies that  $Q^T = Q^{-1}$ .

Now, consider that  $Q$  is an orthogonal matrix, and set  $w = Q^T x$ . Then we have that:

$$\begin{aligned} \|w\|^2 &= w^T w = x^T Q Q^T x \\ &= x^T x = \|x\|^2 \end{aligned} \quad (1.10)$$

This means that the norm of a vector is preserved under an orthogonal transformation. This is called an isometry. It is a useful property for numerical algorithms, as it helps to avoid numerical instability.

There are two main types of orthogonal transformations that we are interested:

## Rotation matrices

A rotation matrix is an orthogonal matrix that represents a rotation in  $\mathbb{R}^2$  or  $\mathbb{R}^3$ . In  $\mathbb{R}^2$ , a rotation matrix is of the form:

$$Q(\theta) = \begin{bmatrix} \cos(\theta) & -\sin(\theta) \\ \sin(\theta) & \cos(\theta) \end{bmatrix} \quad (1.11)$$

## Reflection matrices

A reflection matrix is an orthogonal matrix that represents a reflection with respect to a hyperplane. If  $n$  denotes the unit normal vector to the hyperplane, then the reflection matrix is of the form:

$$Q = I - 2nn^T \quad (1.12)$$

Note that the inverse of this matrix is itself, as  $Q^T = Q^{-1}$  and in this case,  $Q$  is symmetric ( $Q = Q^T$ ).

### 1.1.9 QR factorization

The QR factorization of a matrix  $A \in \mathbb{R}^{m \times n}$ , with  $m \geq n$ , is a factorization of  $A$  as  $A = QR$ , where  $Q \in \mathbb{R}^{m \times n}$  is an orthogonal matrix and  $R \in \mathbb{R}^{n \times n}$  is an upper triangular matrix.

## Gram-Schmidt process

The Gram-Schmidt process is a method to compute the QR factorization of a matrix. Given a matrix  $A \in \mathbb{R}^{m \times n}$ , the Gram-Schmidt process computes an orthonormal basis for the column space of  $A$ , as follows:

$$\begin{aligned} q_1 &= \frac{a_1}{\|a_1\|} \\ q_i &= a_i - \sum_{j=1}^{i-1} (q_j^T a_i) q_j \quad \forall i = 2, \dots, n \end{aligned} \quad (1.13)$$

where  $a_i$  denotes the  $i$ -th column of  $A$ . The matrix  $Q$  is obtained by stacking the vectors  $q_i$  as columns. The matrix  $R$  is obtained by computing the coefficients of the linear combination of the columns of  $Q$  that give the columns of  $A$ .

### 1.1.10 Eigenvalues and eigenvectors

Given a square matrix  $A \in \mathbb{R}^{n \times n}$ , a scalar  $\lambda$  is called an eigenvalue of  $A$  if there exists a vector  $v \in \mathbb{R}^n$  such that:

$$Av = \lambda v \quad (1.14)$$

The vector  $v$  is called an eigenvector of  $A$  associated with the eigenvalue  $\lambda$ .

Let  $P$  be the matrix whose columns are the eigenvectors of  $A$ , and  $\Lambda$  be the diagonal matrix whose diagonal elements are the eigenvalues of  $A$ . Then we have that:

$$A = P\Lambda P^{-1} \quad (1.15)$$

This is called the eigendecomposition of  $A$ .

The eigenvalues of a matrix are the roots of the characteristic polynomial of  $A$ , which is defined as:

$$\det(A - \lambda I) = 0 \quad (1.16)$$

### 1.1.11 Similar matrices

Two square matrices  $A$  and  $B$  are called similar if there exists a non-singular matrix  $M$  such that:

$$B = M^{-1}AM \quad (1.17)$$

Similar matrices have the same eigenvalues, but not necessarily the same eigenvectors. Let  $(\lambda, y)$  be an eigenpair of  $B$ , then we have:

$$By = \lambda y \Rightarrow M^{-1}AMy = \lambda y \Rightarrow A(My) = \lambda(My) \quad (1.18)$$

This means that  $My$  is an eigenvector of  $A$  associated with the eigenvalue  $\lambda$ . So, to obtain the eigenvectors of  $A$  from the eigenvectors of  $B$ , we need to multiply the eigenvectors of  $B$  by  $M$ .

This property can be useful. For example, if we want to compute the eigenvalues of a matrix  $A$ , we can find some transformation  $M$  such that  $M^{-1}AM = B$  is a simpler matrix to work with, usually a lower triangular matrix. Then we can compute the eigenvalues of  $B$  and obtain the eigenvalues of  $A$ .  $M$  is obtained by the permutation matrices to get from  $A$  to  $B$ . The Givens and Householder transformations are examples of such method.

## 1.2 Power method

The power method is an iterative algorithm to compute the dominant eigenvalue of a matrix (i.e., the eigenvalue with the largest magnitude). The algorithm is as follows:

---

### Algorithm 1 Power method

---

- 1: Choose a random vector  $x^{(0)}$ , s.t.  $\|x^{(0)}\| = 1$
  - 2: **for**  $k = 1, 2, \dots$  **do**
  - 3:    $y^{(k)} = Ax^{(k-1)}$
  - 4:    $x^{(k)} = \frac{y^{(k)}}{\|y^{(k)}\|}$
  - 5:    $\lambda^{(k)} = x^{(k)T}Ax^{(k)}$
  - 6: **end for**
-

The convergence rate is determined by the ratio of the largest eigenvalue to the second largest eigenvalue.

### 1.2.1 Rayleigh quotient

The Rayleigh quotient is the base of the power method algorithm. Given a matrix  $A \in \mathbb{R}^{n \times n}$  and a vector  $x \in \mathbb{R}^n$ , the Rayleigh quotient is defined as:

$$R(x) = \frac{x^T A x}{x^T x} \quad (1.19)$$

For every eigenpair  $(\lambda, v)$  of  $A$ , we have that:

$$R(v) = \frac{v^T A v}{v^T v} = \frac{v^T \lambda v}{v^T v} = \lambda \quad (1.20)$$

This means that the Rayleigh quotient is equal to the eigenvalue associated with the eigenvector  $v$ . This property is used in the power method to compute the dominant eigenvalue of a matrix.

### 1.2.2 Proof of convergence for the power method

The power method converges to the dominant eigenvalue of a matrix. The sketch proof is as follows:

Let  $x^{(0)}$  be our initial vector, and let  $\{v_1, \dots, v_n\}$  be the eigenvectors of  $A$ . We can write  $x^{(0)}$  as a linear combination of the eigenvectors of  $A$  (since the eigenvectors of  $A$  form a basis of  $\mathbb{R}^n$ ):

$$x^{(0)} = \sum_{i=1}^n \alpha_i v_i \quad (1.21)$$

Then we have that:

$$A x^{(0)} = \sum_{i=1}^n \alpha_i A v_i = \sum_{i=1}^n \alpha_i \lambda_i v_i \quad (1.22)$$

Since in every iteration  $k$  of the power method we apply the matrix  $A$  to the vector  $x^{(k-1)}$ , we have that:

$$A x^{(k-1)} = A^k x^{(0)} = \sum_{i=1}^n \alpha_i \lambda_i^k v_i \quad (1.23)$$

Now, let us factorize the previous equation by the dominant eigenvalue  $(\lambda_1)^k$ :

$$A^k x^{(0)} = (\lambda_1)^k \left( \alpha_1 v_1 + \sum_{i=2}^n \alpha_i \left( \frac{\lambda_i}{\lambda_1} \right)^k v_i \right) \quad (1.24)$$

Note that the term inside the parenthesis converges to zero as  $k \rightarrow \infty$ , since the ratio of the other eigenvalues to the dominant eigenvalue is less than 1. This means that  $x^{(k)}$  converges to the direction of the dominant eigenvector  $v_1$ . When it is normalized, its Rayleigh quotient converges to the dominant eigenvalue  $\lambda_1$ .

### 1.2.3 Inverse power method

The inverse power method is an iterative algorithm to compute the eigenvalue with the smallest magnitude of a matrix. Note that the smallest eigenvalue of a matrix is the largest eigenvalue of its inverse, since:

$$Ax = \lambda x \Rightarrow A^{-1}x = \frac{1}{\lambda}x \quad (1.25)$$

The algorithm is similar to the power method, but instead of applying the matrix  $A$  to the vector  $x$ , we apply the inverse of the matrix  $A$ :

---

**Algorithm 2** Inverse power method

---

```
1: Choose a random vector  $x^{(0)}$ , s.t.  $\|x^{(0)}\| = 1$ 
2: for  $k = 1, 2, \dots$  do
3:    $y^{(k)} = A^{-1}x^{(k-1)}$ 
4:    $x^{(k)} = \frac{y^{(k)}}{\|y^{(k)}\|}$ 
5:    $\lambda^{(k)} = x^{(k)T}Ax^{(k)}$ 
6: end for
```

---

In practice, we normally don't compute the inverse of the matrix  $A$ , but instead solve the linear system  $Ax = y^{(k)}$  in each iteration.

### 1.2.4 Shifted inverse power method

The shifted inverse power method is an iterative algorithm to compute the eigenvalue with the smallest magnitude of a matrix, but with a shift  $\mu$  added to the matrix  $A$ . The algorithm is as follows:

---

**Algorithm 3** Shifted inverse power method

---

```
1: Choose a random vector  $x^{(0)}$ , s.t.  $\|x^{(0)}\| = 1$ 
2: for  $k = 1, 2, \dots$  do
3:    $y^{(k)} = (A - \mu I)^{-1}x^{(k-1)}$ 
4:    $x^{(k)} = \frac{y^{(k)}}{\|y^{(k)}\|}$ 
5:    $\lambda^{(k)} = x^{(k)T}Ax^{(k)}$ 
6: end for
```

---

Note that with this algorithm, we are computing the eigenvalue of  $A$  that is closest to the shift  $\mu$ . This can be useful to compute the eigenvalues of a matrix that are close to a given value.

### 1.2.5 Applications: Page Rank

The Page Rank algorithm is an algorithm that is used to rank the importance of web pages. The Page Rank algorithm is based on the idea that the importance of a web page is determined by the number of links that point to the page. The Page Rank algorithm is used by search engines to rank web pages in search results.

The Page Rank algorithm is based on the idea of random walks on the web graph. The web graph is a directed graph where the nodes represent web pages and the edges represent links between web pages. The Page Rank algorithm computes the importance of a web page by simulating a random walk on the web graph.

Suppose that we have a web graph with  $n$  web pages. Let  $A \in \{0,1\}^{n \times n}$  be the adjacency matrix of the web graph, where  $A_{ij} = 1$  if there is a link from web page  $j$  to web page  $i$ , and  $A_{ij} = 0$  otherwise. Now, suppose that from each web page  $j$ , it is equally likely to follow any of the links that point to another web page  $i$ . Then, the probability of following a link from web page  $j$  to web page  $i$  is given by:

$$P_{ij} = \frac{A_{ij}}{\sum_{k=1}^n A_{kj}} \quad (1.26)$$

The matrix  $P \in \mathbb{R}^{n \times n}$  is called the transition matrix of the web graph. Now, suppose that we are on a certain page as an initial state, and we want to know the probability of being on each page after a certain number of steps. This can be computed by the following equation:

$$x^{(t+1)} = Px^{(t)} \quad (1.27)$$

where  $x^{(t)} \in \mathbb{R}^n$  is the probability distribution of being on each page at time  $t$ . Notice that  $x^{(0)}$  is the initial state, with  $x_i^{(0)} = 1$  if we start on page  $i$ , and  $x_i^{(0)} = 0$  otherwise.

Now, we define the steady state probability distribution as the probability distribution of being on each page after an infinite number of steps. We call this vector  $\pi \in \mathbb{R}^n$ . Then, we have that:

$$P\pi = \pi \quad (1.28)$$

Meaning that, the steady state probability distribution is the eigenvector of the transition matrix  $P$  with eigenvalue 1. But notice that it is not always true that the steady state probability distribution is unique. In fact, the transition matrix  $P$  is a stochastic matrix, and it is guaranteed to have an eigenvector with eigenvalue 1. However, this eigenvector may not be unique. It can even happen that this distribution does not exist.

We formulate the Perron-Frobenius theorem, which states that a stochastic matrix has a unique eigenvector with eigenvalue 1 if the matrix is in its irreducible form. Formally, a matrix is irreducible if it is not in the form of:

$$P = \begin{bmatrix} P_{11} & P_{12} \\ 0 & P_{22} \end{bmatrix} \quad (1.29)$$

The theorem even states that the dominant eigenvalue is 1, and the other eigenvalues are less than 1.

When we have an irreducible form, we can just use the power iteration method to find the eigenvector with eigenvalue 1. However, when  $P$  is reducible, we need to introduce the following convex combination:

$$\hat{P} = \alpha P + (1 - \alpha) \frac{1}{n} \mathbb{1} \mathbb{1}^T \quad (1.30)$$

where  $\alpha$  is a damping factor, and  $\mathbb{1}$  is a vector of ones. Notice that now,  $\hat{P}$  is irreducible. We can think of it as adding a teleportation to the random walk, meaning that with probability  $\alpha$ , we follow a link, and with probability  $1 - \alpha$ , we teleport to a random page.

The Page Rank algorithm is based on the idea of finding the steady state probability distribution of the web graph. Naturally, the page with the highest probability is the most important page.

### 1.3 Symmetric matrices

A matrix  $A \in \mathbb{R}^{n \times n}$  is called symmetric if  $A = A^T$ . Symmetric matrices have important properties:

- The eigenvectors of a symmetric matrix form an orthonormal basis of  $\mathbb{R}^n$ .
- All the eigenvalues of a symmetric matrix are real.

Let us prove the first property:

Let  $A$  be a symmetric matrix, and let  $\lambda_1, \dots, \lambda_n$  be its eigenvalues. Let  $v_1, \dots, v_n$  be the eigenvectors associated with the eigenvalues  $\lambda_1, \dots, \lambda_n$ . Let us take two eigenvectors  $v_i$  and  $v_j$ , such that  $i \neq j$ . Then we have that:

$$Av_i = \lambda_i v_i \quad \text{and} \quad Av_j = \lambda_j v_j \quad (1.31)$$

Then we have that:

$$\begin{aligned} (A - \lambda_i I)v_i &= 0 \quad \text{and} \quad (A - \lambda_i I)v_j = (\lambda_j - \lambda_i)v_j \\ \Rightarrow v_i &\in N(A - \lambda_i I) \quad \text{and} \quad v_j \in C(A - \lambda_i I) \end{aligned} \quad (1.32)$$

Since  $A$  is symmetric, we have that  $A - \lambda_i I$  is also symmetric. Then we have that:

$$N(A - \lambda_i I) = N((A - \lambda_i I)^T) \perp C(A - \lambda_i I) \quad (1.33)$$

Concluding that  $v_i$  and  $v_j$  are orthogonal. Since this holds for all pairs of eigenvectors, we have that the eigenvectors of a symmetric matrix form an orthonormal basis of  $\mathbb{R}^n$ .

Now, let us prove the second property:

Since  $A$  is symmetric, we have that  $A = A^T$ . Then we have that:

$$\begin{aligned} Ax &= \lambda x \\ A\bar{x} &= \bar{\lambda} \bar{x} \end{aligned} \quad (1.34)$$

Then we have that:

$$\begin{aligned}\bar{x}^T Ax &= \lambda x^T x = \lambda \|x\|^2 \\ x^T A\bar{x} &= \bar{\lambda} x^T \bar{x} = \bar{\lambda} \|x\|^2\end{aligned}\tag{1.35}$$

Since  $A = A^T$ , we have that:

$$\bar{x}^T Ax = (Ax)^T \bar{x} = x^T A^T \bar{x} = x^T A\bar{x}\tag{1.36}$$

Then we have that:

$$\lambda \|x\|^2 = \bar{\lambda} \|x\|^2 \Rightarrow \lambda = \bar{\lambda}\tag{1.37}$$

This means that the eigenvalues of a symmetric matrix are real.

### 1.3.1 Symmetric positive definite matrices

A matrix  $A \in \mathbb{R}^{n \times n}$  is called symmetric positive definite if it is symmetric and if for every vector  $x \in \mathbb{R}^n$  we have that:

$$x^T Ax > 0 \quad \forall x \neq 0\tag{1.38}$$

Symmetric positive definite matrices have important properties:

- All the eigenvalues of a symmetric positive definite matrix are positive.
- The Cholesky factorization of a symmetric positive definite matrix exists and is unique:

$$A = LL^T\tag{1.39}$$

In fact, a symmetric matrix is positive definite if and only if all its eigenvalues are positive, so:

$$x^T Ax > 0 \quad \forall x \neq 0 \quad \Leftrightarrow \quad \lambda_i > 0 \quad \forall i\tag{1.40}$$

Note that the quantity  $x^T Ax$  is called the energy of the vector  $x$  with respect to the matrix  $A$ . This quantity is always positive for a symmetric positive definite matrix.

## 1.4 Singular Value Decomposition (SVD)

The Singular Value Decomposition (SVD) is a factorization of a matrix  $A \in \mathbb{R}^{m \times n}$  as:

$$A = U\Sigma V^T\tag{1.41}$$

where  $U \in \mathbb{R}^{m \times m}$  and  $V \in \mathbb{R}^{n \times n}$  are orthogonal matrices, and  $\Sigma \in \mathbb{R}^{m \times n}$  is a quasi-diagonal matrix with the singular values of  $A$ . The singular values of  $A$  are the square roots of the eigenvalues of  $A^T A$ .

Note that if  $\text{rank}(A) = r$ , then the matrix  $\Sigma$  has  $r$  non-zero singular values, and the remaining singular values are zero. Assume that the singular values of  $A$  are  $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_r > 0$ . Then we have that:



$$\begin{aligned} Av_i &= \sigma_i u_i \quad \forall i = 1, \dots, r \\ Av_i &= 0 \quad \forall i = r + 1, \dots, n \end{aligned} \tag{1.42}$$

where  $u_i$  and  $v_i$  are the columns of  $U$  and  $V$ , respectively. The vectors  $u_i$  and  $v_i$  are called the left and right singular vectors of  $A$ , respectively.

The SVD can also be written as:

$$A = \sum_{i=1}^r \sigma_i u_i v_i^T \tag{1.43}$$

#### 1.4.1 Economy SVD

The economy SVD is a factorization of a matrix  $A \in \mathbb{R}^{m \times n}$  as:

$$A = U_r \Sigma_r V_r^T \tag{1.44}$$

where  $U \in \mathbb{R}^{m \times r}$ ,  $V \in \mathbb{R}^{n \times r}$  and  $\Sigma \in \mathbb{R}^{r \times r}$ , with  $r = \text{rank}(A)$ .

The economy SVD is useful when we are only interested in the first  $r$  singular values of  $A$ , which are the non-zero singular values.

#### 1.4.2 Low-rank approximation

The SVD can be used to compute a low-rank approximation of a matrix  $A \in \mathbb{R}^{m \times n}$ . Given a rank  $k$ , with  $k < r = \text{rank}(A)$ , the low-rank approximation of  $A$  is given by:

$$A_k = U_k \Sigma_k V_k^T = \sum_{i=1}^k \sigma_i u_i v_i^T \tag{1.45}$$

where  $U_k \in \mathbb{R}^{m \times k}$ ,  $V_k \in \mathbb{R}^{n \times k}$  and  $\Sigma_k \in \mathbb{R}^{k \times k}$ , with  $k < r = \text{rank}(A)$ .

Because the singular values of  $A$  are sorted in decreasing order, the low-rank approximation only considers the first  $k$  singular values of  $A$ , as they are the components of  $A$  with the largest contribution.

The low-rank approximation of a matrix is useful for data compression, as it allows to represent a matrix with a smaller number of parameters. Note that the low-rank approximation of a matrix is the best rank- $k$  approximation of the matrix in the Frobenius norm. This is called the Eckart-Young theorem.

#### Optimal thresholding for low-rank approximation

There are several methods to determine the optimal rank  $k$  for the low-rank approximation of a matrix. A common heuristic approach is to graph the singular values of the matrix and choose the

rank  $k$  as the point where the singular values start to decay rapidly.

We can see the idea of this method in the following figure:

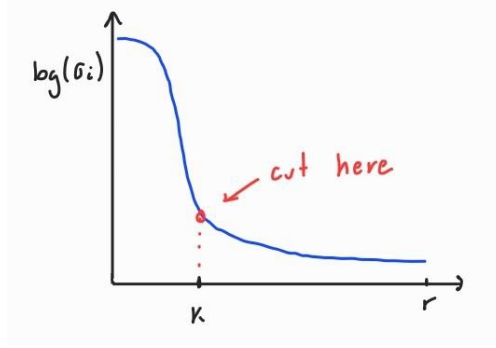


Figure 1.1: Heuristic method for choosing  $k$

Another common method is to use the cumulative energy of the singular values. The cumulative energy is defined as:

$$\frac{\sum_{i=1}^k \sigma_i^2}{\sum_{i=1}^r \sigma_i^2} \quad (1.46)$$

The cumulative energy measures the proportion of the total energy of the matrix that is captured by the first  $k$  singular values. A common heuristic is to choose the rank  $k$  as the smallest value such that the cumulative energy is above a certain threshold, e.g., 90%.

However, the optimal rank  $k$  for the low-rank approximation of a matrix is a complex problem that depends on the specific application. For example, we can use the following model:

$$A = A_{TRUE} + \gamma A_{NOISE} \quad (1.47)$$

where  $A_{TRUE}$  is the true low-rank matrix that we want to approximate,  $A_{NOISE}$  is the noise matrix, and  $\gamma$  is the noise level. In this case, the optimal rank  $k$  is the one that minimizes the error of the low-rank approximation of  $A_{TRUE}$ .

For this case, depending of the knowledge of the noise level, we can use the following methods to compute a threshold.

1. If the noise level is known, and  $A \in \mathbb{R}^{n \times n}$  is a square matrix, then we can use the following threshold:

$$\tau = \frac{4}{\sqrt{3}} \sqrt{n} \gamma \quad (1.48)$$

2. If  $n \ll m$ :

$$\tau = \lambda \frac{n}{m} \sqrt{n} \quad (1.49)$$

Where  $\lambda(\cdot)$  is a computable function (specific for this model, we don't care about it).

3.  $\gamma$  is unknown:

$$\tau = \omega\left(\frac{n}{m}\right) \sigma_{med} \quad (1.50)$$

Where  $\sigma_{med}$  is the median of the singular values of  $A$ , and  $\omega(\cdot)$  is a computable function (specific for this model, we don't care about it).

The idea is to select the singular values that are above the threshold.

### 1.4.3 Existence of the SVD

The SVD of a matrix  $A \in \mathbb{R}^{m \times n}$  always exists. The proof is as follows:

Let  $A \in \mathbb{R}^{m \times n}$  be a matrix. Then we have that  $A^T A \in \mathbb{R}^{n \times n}$ . We have that  $A^T A$  is symmetric and positive semidefinite. In fact:

- $A^T A$  is symmetric:  $(A^T A)^T = A^T (A^T)^T = A^T A$
- $A^T A$  is positive semidefinite:  $x^T A^T A x = (Ax)^T Ax = \|Ax\|^2 \geq 0$

Therefore,  $A^T A$  has an eigendecomposition:

$$A^T A = V \Lambda V^T \quad \text{s.t. } V^T V = I$$

Let us show now that  $\text{rank}(A) = \text{rank}(A^T A)$ . We can prove this by showing that the null space of  $A$  is the same as the null space of  $A^T A$ . Let  $x \in N(A)$ , then we have that:

$$Ax = 0 \Rightarrow A^T Ax = 0$$

This means that  $N(A) \subseteq N(A^T A)$ . Now, let  $x \in N(A^T A)$ , then we have that:

$$A^T Ax = 0 \Rightarrow x^T A^T Ax = 0 \Rightarrow \|Ax\|^2 = 0 \Rightarrow Ax = 0$$

This means that  $N(A^T A) \subseteq N(A)$ . Therefore,  $N(A) = N(A^T A)$ , and we have that  $\text{rank}(A) = \text{rank}(A^T A)$ .

Now, let us consider the eigenpairs  $(\lambda_i, v_i)$  of  $A^T A$ . Then let us define  $u_i$  as:

$$u_i = \frac{1}{\sqrt{\lambda_i}} A v_i$$

We will prove that each  $u_i$  is a unitary vector, and that the vectors  $u_i$  are mutually orthogonal. Let us first show that each  $u_i$  is a unitary vector:

$$\|u_i\|^2 = \left\| \frac{1}{\sqrt{\lambda_i}} Av_i \right\|^2 = \frac{1}{\lambda_i} v_i^T A^T A v_i = \frac{1}{\lambda_i} \lambda_i = 1$$

Now, let us show that the vectors  $u_i$  are mutually orthogonal. Let  $i \neq j$ , then we have that:

$$\begin{aligned} u_i^T u_j &= \frac{1}{\sqrt{\lambda_i}} v_i^T A^T \frac{1}{\sqrt{\lambda_j}} A v_j = \frac{1}{\sqrt{\lambda_i \lambda_j}} v_i^T A^T A v_j \\ &= \frac{1}{\sqrt{\lambda_i \lambda_j}} \lambda_j v_i^T v_j = 0 \end{aligned}$$

Now, we can show that the vectors  $u_i$  are the eigenvectors of  $AA^T$ , with eigenvalues  $\lambda_i$ :

$$\begin{aligned} AA^T u_i &= AA^T \frac{1}{\sqrt{\lambda_i}} A v_i = \\ \frac{1}{\sqrt{\lambda_i}} AA^T A v_i &= \frac{1}{\sqrt{\lambda_i}} A \lambda_i v_i = \lambda_i u_i \end{aligned}$$

Therefore, we have that  $AA^T = U\Lambda U^T$ , where  $U$  is the matrix whose columns are the vectors  $u_i$ . Now, let us define  $\Sigma$  as the diagonal matrix whose diagonal elements are the square roots of the eigenvalues of  $A^T A$ :

$$\Sigma = \begin{bmatrix} \sqrt{\lambda_1} & 0 & \dots & 0 \\ 0 & \sqrt{\lambda_2} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \sqrt{\lambda_n} \end{bmatrix}$$

Then we have that:

$$A = U\Sigma V^T$$

where  $V$  is the matrix whose columns are the eigenvectors of  $A^T A$ . Therefore, the SVD of  $A$  exists.

#### 1.4.4 Computation of the SVD

The SVD of a matrix can be computed using the following steps:

1. Compute the eigendecomposition of  $A^T A$ .
2. Compute the singular values of  $A$  as the square roots of the eigenvalues of  $A^T A$ .
3. Compute the right singular vectors of  $A$  as the eigenvectors of  $A^T A$ .
4. Compute the left singular vectors of  $A$  as  $u_i = \frac{1}{\sqrt{\lambda_i}} A v_i$ .

### 1.4.5 Geometrical interpretation of the SVD

The SVD of a matrix  $A \in \mathbb{R}^{m \times n}$  can be interpreted geometrically as a composition of three transformations:

- The matrix  $V$  represents a rotation in  $\mathbb{R}^n$ .
- The matrix  $\Sigma$  represents a scaling in  $\mathbb{R}^{m \times n}$ .
- The matrix  $U$  represents a rotation in  $\mathbb{R}^m$ .

In  $\mathbb{R}^2$ , the SVD of a matrix  $A$  can be interpreted as a rotation of the unit circle, followed by a scaling along the axes, followed by another rotation. So this transformation can be represented with 4 parameters: two angles and two scaling factors. In  $\mathbb{R}^3$ , it's the same idea, but with 6 parameters: three angles and three scaling factors.

### 1.4.6 Polar decomposition

Let  $A \in \mathbb{R}^{n \times n}$  be a symmetric matrix. The polar decomposition of  $A$  is a factorization of  $A$  as:

$$A = QS \quad (1.51)$$

where  $Q$  is an orthogonal matrix and  $S$  is a symmetric positive definite matrix. The polar decomposition can be derived from the SVD of  $A$  as:

$$A = U\Sigma V^T = (UV^T)(V\Sigma V^T) = QS \quad (1.52)$$

where  $Q = UV^T$  and  $S = V\Sigma V^T$ . So, the polar decomposition of a symmetric matrix always exists.

### 1.4.7 Eckart-Young theorem

The Eckart-Young theorem states that the best rank- $k$  approximation of a matrix  $A \in \mathbb{R}^{m \times n}$  in the Frobenius norm is given by the low-rank approximation of  $A$ , denoted as  $A_k$ :

$$\min_{\text{rank}(B)=k} \|A - B\|_F = \|A - A_k\|_F \quad (1.53)$$

This also applies to the spectral norm  $\|\cdot\|_2$ . To prove this, let us define the mentioned norms:

#### Frobenius norm

The Frobenius norm of a matrix  $A \in \mathbb{R}^{m \times n}$  is defined as:

$$\|A\|_F = \sqrt{\sum_{i=1}^m \sum_{j=1}^n A_{ij}^2} \quad (1.54)$$

The Frobenius norm is the Euclidean norm of the vector obtained by stacking the columns of  $A$ . It is also equal to:

$$\|A\|_F = \sqrt{\text{tr}(A^T A)} \quad (1.55)$$

It has 3 important properties:

1.  $\|A\|_F = \|A^T\|_F$
2. It is invariant under orthogonal transformations:  $\|QA\|_F = \|A\|_F$
3. It is equal to the square root of the sum of the squared singular values of  $A$ :

$$\|A\|_F = \sqrt{\sum_{i=1}^r \sigma_i^2} \quad (1.56)$$

### Spectral norm

The spectral norm of a matrix  $A \in \mathbb{R}^{n \times n}$  is defined as:

$$\|A\|_2 = \sup_{\|x\|_2=1} \|Ax\|_2 \quad (1.57)$$

It is equal to the largest singular value of  $A$ . This is denoted as:

$$\|A\|_2 = \sigma_1 \quad (1.58)$$

### Proof of the Eckart-Young theorem

For both norms  $\|\cdot\|_F$  and  $\|\cdot\|_2$ , we will prove that:

$$\|A - A_k\| \leq \|A - B\| \quad \forall B \in \mathbb{R}^{m \times n} \quad \text{s.t.} \quad \text{rank}(B) = k$$

Let  $B \in \mathbb{R}^{m \times n}$  be a matrix such that  $\text{rank}(B) = k$ . Then we have that:

$$\dim(N(B)) = n - k$$

Let us consider the matrix  $V_{k+1}$  whose columns are the vectors  $v_1, \dots, v_{k+1}$  from the SVD of  $A$ . Then we have that:

$$\begin{aligned} \text{rank}(V_{k+1}) = k + 1 &\Rightarrow \dim(C(V_{k+1})) = k + 1 \\ \Rightarrow \dim(N(B)) + \dim(C(V_{k+1})) &= n - k + k + 1 = n + 1 \end{aligned}$$

This means that  $N(B) \cap C(V_{k+1}) \neq \emptyset$ . Let  $w \in N(B) \cap C(V_{k+1})$ , such that  $\|w\| = 1$ . Then we have that, for coefficients  $\alpha_i$ :

$$w = \sum_{i=1}^k \alpha_i v_i \quad \wedge \quad Bw = 0$$

Since  $\|w\| = 1$ , then  $\sum_{i=1}^k \alpha_i^2 = 1$ . Now, we have that, for the spectral norm:

$$\|A - B\|_2^2 \geq \|(A - B)w\|_2^2 = \|Aw\|_2^2 = w^T A^T A w$$

$$\begin{aligned}
&= w^T V \Sigma^2 V^T w = \sum_{i=1}^k \sigma_i^2 \alpha_i^2 \geq \\
&\geq \sigma_{k+1}^2 \sum_{i=1}^k \alpha_i^2 = \sigma_{k+1}^2 \\
&= \|A - A_k\|_2^2
\end{aligned}$$

Therefore:

$$\|A - A_k\|_2 \leq \|A - B\|_2 \quad \forall B \in \mathbb{R}^{m \times n} \quad \text{s.t.} \quad \text{rank}(B) = k$$

For the Frobenius norm, we need to use the Weyl's inequality:

$$\sigma_{i+j-1}(X + Y) \leq \sigma_i(X) + \sigma_j(Y) \quad (1.59)$$

where  $\sigma_i(X)$  denotes the  $i$ -th singular value of a matrix  $X$ . Using this inequality, let us define  $X = A - B$  and  $Y = B$ . Then we have that:

$$\sigma_{i+k}(A) \leq \sigma_i(A - B) + \sigma_{k+1}(B)$$

Note that  $\sigma_{k+1}(B) = 0$ , since  $B$  has rank  $k$ . Therefore:

$$\sigma_{i+k}(A) \leq \sigma_i(A - B)$$

Then, we have that:

$$\begin{aligned}
\|A - A_k\|_F^2 &= \sum_{i=k+1}^r \sigma_i^2(A) = \sum_{i=1}^{r-k} \sigma_{i+k}^2(A) \leq \\
&\leq \sum_{i=1}^{r-k} \sigma_i^2(A - B) \leq \sum_{i=1}^n \sigma_i^2(A - B) = \|A - B\|_F^2
\end{aligned}$$

Therefore:

$$\|A - A_k\|_F \leq \|A - B\|_F \quad \forall B \in \mathbb{R}^{m \times n} \quad \text{s.t.} \quad \text{rank}(B) = k$$

This completes the proof.

### 1.4.8 Principal Component Analysis (PCA)

Principal Component Analysis (PCA) is a technique to reduce the dimensionality of a dataset. Given a dataset  $X \in \mathbb{R}^{m \times n}$ , PCA computes the SVD of the centered data matrix  $\bar{B}$ , where  $\bar{B}$  is obtained by subtracting the mean of each column of  $X$ . Then, the principal components of the dataset are the right singular vectors of  $\bar{B}$ . The principal components are the directions of the dataset with the largest variance.

In detail, let  $X \in \mathbb{R}^{m \times n}$  the dataset, with  $m$  samples and  $n$  features. Let  $\bar{x}$  be mean vector of the data samples. Then, we define  $\bar{X}$  as:

$$\bar{X} = \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix} \bar{x}^T \quad (1.60)$$

Then, we define the centered data matrix  $\bar{B}$  as:

$$\bar{B} = X - \bar{X} \quad (1.61)$$

The covariance matrix is then defined as:

$$C = \frac{1}{m-1} \bar{B}^T \bar{B} \quad (1.62)$$

We observe that the eigendecomposition of the covariance matrix  $C$  is as follows:

$$C = V \Lambda V^T \quad (1.63)$$

We can obtain the principal components of the dataset as the columns of  $V$ . The principal components are the directions of the dataset with the largest variance. The first principal component is the direction with the largest variance, the second principal component is the direction with the second largest variance, and so on.

By doing the SVD of the centered data matrix  $\bar{B}$ , we can obtain the principal components of the dataset, such that:

$$\bar{B} = U \Sigma V^T \Rightarrow V = \text{principal components} \quad \wedge \quad \Lambda = \frac{1}{m-1} \Sigma^2 \quad (1.64)$$

#### 1.4.9 Data pre-processing: sensitivity to transformations

The SVD is specially sensitive to scaling, rotating and translating the data. This is because the SVD is a geometrical transformation of the data, and these transformations change the geometry of the data. For that reason, if we want to apply the SVD to a dataset, we need to pre-process the data to remove or somehow filter the effects of these transformations.

For example, if we have a dataset that consist of face images, and we want to apply the SVD to this dataset, we need to pre-process the images to remove the effects of scaling, rotating and translating the faces. This can be done by aligning the faces to a common reference frame, and by normalizing the faces to have the same size and orientation.

In general, the pre-processing of the data is an important step in the application of the SVD to a dataset. The pre-processing step can have a significant impact on the results of the SVD, and it is important to carefully choose the pre-processing steps to obtain meaningful results.

#### 1.4.10 Randomized SVD

The SVD of a matrix can be computed using the randomized SVD algorithm. The randomized SVD is an iterative algorithm that approximates the SVD of a matrix using random projections. The algorithm is as follows:



---

**Algorithm 4** Randomized SVD

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- 1: Choose a random matrix  $Y \in \mathbb{R}^{n \times k}$ , with  $k \ll n$ .
  - 2: Compute the matrix  $Z = AY$ .
  - 3: Compute the QR factorization of  $Z = QR$ .
  - 4: Compute the matrix  $B = Q^T A$ .
  - 5: Compute the SVD of  $B = \hat{U}\Sigma V^T$ .
  - 6: Obtain  $U$  as  $U = Q\hat{U}$ .
  - 7: The approximate SVD of  $A$  is given by  $A \approx U\Sigma V^T$ .
- 

The randomized SVD algorithm is a fast and efficient way to compute the SVD of a matrix. The algorithm is based on the fact that the SVD of a matrix can be approximated using random projections. The algorithm is especially useful when the matrix is large and when only an approximate SVD is needed.

The idea behind this algorithm is a sampling of the column space of the matrix  $A$  using the matrix  $Y$ .

This method is most widely used in the context of large-scale data analysis, where the matrix  $A$  is too large to be stored in memory. In this case, the randomized SVD algorithm can be used to compute an approximate SVD of the matrix using only a small amount of memory.

## 1.5 Matrix completion

Let  $X \in \mathbb{R}^{n \times p}$  be a matrix that is partially observed. That is, we only know  $X_{ij}$  for  $(i, j) \in \Omega$ , where  $\Omega$  is the set of observed entries. Assume that  $\text{rank}(X) = r \ll \min(n, p)$ . Our goal is to estimate  $X_{ij}$  for every  $(i, j) \notin \Omega$ .

Formally, we express this as:

$$\hat{X} = \operatorname{argmin}_Z \{\text{rank}(Z)\} \quad \text{s.t.} \quad Z_{ij} = X_{ij} \quad \forall (i, j) \in \Omega \quad (1.65)$$

However, this formulation is a non-convex optimization problem, and it is not solvable in practical cases. Instead, we can use the following convex optimization problem:

$$\hat{X} = \operatorname{argmin}_Z \{\|Z\|_*\} \quad \text{s.t.} \quad Z_{ij} = X_{ij} \quad \forall (i, j) \in \Omega \quad (1.66)$$

where  $\|Z\|_*$  is the nuclear norm of  $Z$ , defined as:

$$\|Z\|_* = \sum_{i=1}^r \sigma_i(Z) \quad (1.67)$$

It is possible to prove that minimizing the nuclear norm of  $Z$  is highly probable to be equivalent to minimizing the rank of  $Z$ . This is because the nuclear norm is the convex relaxation of the rank function.

There are many ways of solving this optimization problem. One of the most common methods is the Singular Value Thresholding (SVT) algorithm. The SVT algorithm is an iterative algorithm that approximates the solution of the optimization problem. The algorithm is as follows:

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**Algorithm 5** Singular Value Thresholding (SVT)

---

```

1: Initialize  $Z_0 = X$ .
2: for  $k = 1, 2, \dots$  do
3:   Compute the SVD of  $Z_{k-1} = U\Sigma V^T$ .
4:   Compute the singular value thresholding of  $\Sigma$  as  $\Sigma_\lambda = \max(\Sigma - \lambda, 0)$ .
5:   Compute the matrix  $Z_k = U\Sigma_\lambda V^T$ .
6:   Make  $Z_k$  satisfy the constraints  $Z_k = X$  for  $(i, j) \in \Omega$ .
7:   If  $\|Z_k - Z_{k-1}\| < \epsilon$ , then stop.
8: end for

```

---

The SVT algorithm is a fast and efficient way to solve the matrix completion problem. The algorithm is based on the fact that the nuclear norm of a matrix is the sum of its singular values. The algorithm iteratively computes the singular value thresholding of the matrix, and then makes the matrix satisfy the constraints. The algorithm converges to the solution of the optimization problem, and it is guaranteed to converge to the global minimum.

## 1.6 Regression Methods

Regression methods are a class of machine learning algorithms that are used to predict the value of a continuous variable based on the value of one or more input variables. Regression methods are widely used in data analysis, and they are used to model the relationship between the input variables and the output variable.

### 1.6.1 Least Squares

The least squares problem is a problem to find the vector  $x$  that minimizes the residual of a linear system  $Ax = b$ . The least squares solution is given by:

$$x = (A^T A)^{-1} A^T b \quad (1.68)$$

Here, we are assuming that  $A \in \mathbb{R}^{m \times n}$ , with  $m > n$  and  $\text{rank}(A) = n$ . The least squares solution is the solution that minimizes the residual  $\|Ax - b\|_2$ . We can prove this in 2 different ways:

- **Geometrical interpretation:**

The least squares solution is the solution that minimizes the residual  $\|Ax - b\|_2$ . This means that the residual is orthogonal to the column space of  $A$ . Let  $r = Ax - b$  be the residual, then we have that:

$$\begin{aligned}
r \perp C(A) &\Rightarrow r \in N(A^T) \Rightarrow A^T r = 0 \\
&\Rightarrow A^T (Ax - b) = 0 \Rightarrow A^T Ax = A^T b
\end{aligned}$$

Note that  $A$  has full column rank, so  $A^T A$  is invertible (in fact, it is s.d.p). Therefore, we have that:

$$x = (A^T A)^{-1} A^T b$$

• **Derivation:**

The least squares solution is the solution that minimizes the residual  $\|Ax - b\|_2$ . This means that the least squares solution is the solution that minimizes the function:

$$\mathcal{L}(x) = \|Ax - b\|_2^2 = (Ax - b)^T (Ax - b) = x^T A^T A x - 2b^T A x + b^T b$$

To find the minimum of this function, we need to find the critical points of the function. Let us take the derivative of the function with respect to  $x$ :

$$\begin{aligned} \frac{\partial \mathcal{L}(x)}{\partial x} &= 2A^T A x - 2A^T b = 0 \\ \Rightarrow A^T A x &= A^T b \end{aligned}$$

By the same argument as before, we have that:

$$x = (A^T A)^{-1} A^T b$$

In the context of data analysis, we often have a matrix  $X \in \mathbb{R}^{n \times p}$  that represents  $n$  samples of  $p$  features. In this case, the least squares solution is used to find the coefficients  $w$  of a linear model that best fits the data. The system of equations is given by:

$$Xw = y \tag{1.69}$$

where  $X$  is the matrix of features,  $w$  is the vector of coefficients, and  $y$  is the vector of labels. Note that an exact solution may not exist, so we need to find the least squares solution. The least squares solution, as we deduced before, would be given by:

$$w = (X^T X)^{-1} X^T y \tag{1.70}$$

Note that the matrix  $(X^T X)^{-1} X^T$  is called the pseudo-inverse of  $X$ , and it is denoted as  $X^\dagger$ .

### Least squares and SVD

Computing the inverse of the matrix  $X^T X$  can be computationally expensive, especially when the matrix  $X$  is large. In this case, the least squares solution can be computed using the SVD of the matrix  $X$ . In fact, we have that:

$$w = V \Sigma^\dagger U^T y \tag{1.71}$$

where  $U$ ,  $\Sigma$  and  $V$  are the SVD of  $X$ , and  $\Sigma^\dagger$  is the pseudo-inverse of  $\Sigma$ . The pseudo-inverse of  $\Sigma$  is obtained by taking the reciprocal of the non-zero singular values of  $\Sigma$ , and then taking

the transpose of the resulting matrix. The pseudo-inverse of  $\Sigma$  is denoted as  $\Sigma^\dagger$ . Visually, the pseudo-inverse of  $\Sigma$  is as follows:

$$\Sigma^\dagger = \begin{bmatrix} \frac{1}{\sigma_1} & 0 & \dots & 0 & \dots & 0 \\ 0 & \frac{1}{\sigma_2} & \dots & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots & \dots & \vdots \\ 0 & 0 & \dots & \frac{1}{\sigma_p} & \dots & 0 \end{bmatrix} \quad (1.72)$$

### Underdetermined system case (wide matrix)

In the previous case, we assumed that  $n > p$ , and  $\text{rank}(X) = p$ . However, in some cases, we have that  $n < p$ , and  $\text{rank}(X) = n$ . In this case, the system of equations  $Xw = y$  is underdetermined, and there are infinitely many solutions. Still, the least square solution is still useful, as it gives the solution that minimizes the norm of the coefficients.

We can prove this as follows. Let  $w$  be a solution of the system of equations  $Xw = y$  and let  $w_0$  be the least squares solution. Then we have that:

$$\begin{aligned} \|w\|^2 &= \|(w - w_0) + w_0\| = [(w - w_0) + w_0]^T [(w - w_0) + w_0] \\ &= \|w - w_0\|^2 - 2(w - w_0)^T w_0 + \|w_0\|^2 \end{aligned}$$

Note that

$$\begin{aligned} (w - w_0)^T w_0 &= (w - w_0)^T X^T (XX^T)^{-1} y \\ &= (X(w - w_0))^T (XX^T)^{-1} y \end{aligned}$$

And since  $Xw = y$  and  $Xw_0 = y$ , we have that  $X(w - w_0) = 0$ . Therefore, we have that:

$$\|w\|^2 = \|w - w_0\|^2 + \|w_0\|^2 \geq \|w_0\|^2$$

This means that the least squares solution is the solution that minimizes the norm of the coefficients. We often use this in practice to avoid error amplification in our calculations.

### 1.6.2 Ridge regression

Let  $w_{LS} = (X^T X)^{-1} X^T y$  be the least squares solution of the system of equations  $Xw = y$ . Then, let us assume the following model:

$$y = Xw^* + \varepsilon \quad (1.73)$$

where  $w^*$  is the true coefficients of the model, and  $\varepsilon$  is the noise. Then:

$$w_{LS} = (X^T X)^{-1} X^T (Xw^* + \varepsilon) = w^* + (X^T X)^{-1} X^T \varepsilon \quad (1.74)$$

Note that  $X = U\Sigma V^T$ , so we have that  $(X^T X)^{-1} X^T = V\Sigma^\dagger U^T$ . Then:

$$w_{LS} = w^* + V\Sigma^\dagger U^T \varepsilon \quad (1.75)$$

The least squares solution is the true coefficients plus a term that depends on the noise. Let us suppose that for some  $p$ , we have that  $\sigma_p \approx 0$ . Then we have that  $\Sigma_{pp}^\dagger = 1/\sigma_p \gg 0$ , and this will amplify the noise.

To avoid this, we can use the ridge regression. The ridge regression is a technique that adds a regularization term to the least squares solution. The optimization problem of the ridge regression is as follows:

$$w_R = \operatorname{argmin} \{ \|Xw - y\|_2^2 + \lambda \|w\|_2^2 \} \quad (1.76)$$

where  $\lambda$  is the regularization parameter. The solution of the ridge regression is given by:

$$w_R = (X^T X + \lambda I)^{-1} X^T y \quad (1.77)$$

Let us analyse this solution with the SVD of  $X$ :

$$\begin{aligned} w_R &= (V \Sigma^T U^T U \Sigma V^T + \lambda V V^T)^{-1} V \Sigma U^T y \\ &= (V (\Sigma^T \Sigma + \lambda I) V^T)^{-1} V \Sigma U^T y \\ &= V (\Sigma^T \Sigma + \lambda I)^{-1} \Sigma U^T y \end{aligned}$$

Note that, if  $S = (\Sigma^T \Sigma + \lambda I)^{-1} \Sigma^T$ , then we have that:

$$S_{pp} = \frac{\sigma_p}{\sigma_p + \lambda}$$

So we have 2 cases:

- If  $\sigma_p \gg 0$ , then  $S_{pp} \approx 1/\sigma_p \approx 0$ .
- If  $0 < \sigma_p \ll 1$ , then  $S_{pp} \approx 0$ .

In both cases, the ridge regression avoids the error amplification of the noise. However, notice that the ridge regression introduces a bias in the coefficients. This bias is controlled by the regularization parameter  $\lambda$ . The larger the value of  $\lambda$ , the larger the bias in the coefficients.

In general, when we know that  $\varepsilon = 0$ , the least squares solution is the best solution. However, when we have noise in the data, the ridge regression is a better solution, as it avoids the error amplification of the noise.

The ridge regression is a special case of the Tikhonov regularization, which is a general technique to regularize ill-posed problems. The Tikhonov regularization adds a regularization term to the least squares solution, which is controlled by the regularization parameter  $\lambda$ .

### 1.6.3 LASSO regression

The LASSO regression is a technique that adds a regularization term to the least squares solution. The optimization problem of the Lasso regression is as follows:

$$w_L = \operatorname{argmin} \{ \|Xw - y\|_2^2 + \lambda \|w\|_1 \} \quad (1.78)$$

where  $\lambda$  is the regularization parameter. Notice that the Lasso regression adds a regularization term that is the  $L_1$  norm of the coefficients. This approach is useful when we want to find a sparse solution, i.e., a solution with few non-zero coefficients.

LASSO stands for Least Absolute Shrinkage and Selection Operator. In general, we have that:

- When  $\lambda$  is large, we tend to have a higher sparsity in the solution.
- When  $\lambda$  is small, we tend to have a solution that is closer to the least squares solution.

We can mix both LASSO and Ridge regression in the Elastic Net regression. The Elastic Net regression is a technique that adds both  $L_1$  and  $L_2$  regularization terms to the least squares solution. The optimization problem of the Elastic Net regression is as follows:

$$w_E = \operatorname{argmin} \{ \|Xw - y\|_2^2 + \lambda_1 \|w\|_1 + \lambda_2 \|w\|_2^2 \} \quad (1.79)$$

where  $\lambda_1$  and  $\lambda_2$  are the regularization parameters. The Elastic Net regression is useful when we want to find a solution that is both sparse and close to the least squares solution.

### 1.6.4 Kernel methods

Let us formulate the Ridge Regression solution in terms of the SVD of  $X$ . We have that:

$$\begin{aligned} w_R &= V \Sigma^T U^T U (\Sigma \Sigma^T + \lambda I)^{-1} U^T y \\ &= X^T \beta \end{aligned}$$

where  $\beta = U(\Sigma \Sigma^T + \lambda I)^{-1} \Sigma U^T y = (X X^T + \lambda I)^{-1} y$ . We call  $\beta$  the dual coefficients of the Ridge Regression.

In general, we can have a linear separator in the input space, given by:

$$y = f(x) = w^T x = x_{i1} w_1 + \dots + x_{ip} w_p$$

But we could also have a linear separator in a new feature space  $\phi(\cdot)$ , given by:

$$y = f(x) = w^T \phi(x) = \phi(x)_{i1} w_1 + \dots + \phi(x)_{id} w_d$$

where  $\phi(x)$  is a non-linear transformation of the input space. This is the idea behind the kernel methods. Notice that the dimension of the feature space can be larger than the dimension of the input space. For example:

$$\phi(x) = (x_1, x_2, x_1^2, x_2^2, x_1 x_2)$$

In this case, the dimension of the feature space is 5, larger than the dimension of the input space, which is 2. Note that the feature space adds non-linear features to the input space.

In this sense, we can create the following matrix:

$$\Phi = \begin{bmatrix} \phi(x_1)^T \\ \vdots \\ \phi(x_n)^T \end{bmatrix}$$

Then, to find the coefficients  $w$  of the linear separator in the feature space, we can use the Ridge Regression solution:

$$w = \Phi^T \beta$$

where  $\beta = (\Phi\Phi^T + \lambda I)^{-1}y$ . To avoid the computation of the feature space, we can use kernel functions. A kernel function is a function that computes the inner product of two vectors in the feature space without computing the feature space. Formally:

$$K(x, x') = \phi(x)^T \phi(x') \quad (1.80)$$

The most common kernel functions are:

- Linear kernel:  $K(x, x') = x^T x'$
- Polynomial kernel:  $K(x, x') = (x^T x' + c)^d$
- Gaussian kernel:  $K(x, x') = \exp(-\frac{\|x-x'\|^2}{2\sigma^2})$

In the previous case for  $\phi(x) = (x_1, x_2, x_1^2, x_2^2, x_1x_2)$ , we have that:

$$\begin{aligned} K(x, x') &= (x^T x')^2 \\ \Rightarrow K(x, x') &= (x_1x'_1 + x_2x'_2 + x_1^2x'^2_1 + x_2^2x'^2_2 + x_1x_2x'_1x'_2) = \phi(x)^T \phi(x') \end{aligned}$$

Now, notice that:

$$[\Phi\Phi^T]_{ij} = \phi(x_i)^T \phi(x_j) = K(x_i, x_j)$$

Therefore, we have that:

$$\beta = (K + \lambda I)^{-1}y \quad (1.81)$$

where  $K$  is the kernel matrix, defined as  $K_{ij} = K(x_i, x_j)$ . Then, the coefficients  $w$  of the linear separator in the feature space are given by:

$$w = \Phi^T \beta \quad (1.82)$$

Suppose now that we have a new sample  $x$  and we want to predict its label  $y$ . Then, we have that:

$$y = w_R^T \cdot \phi(x) = \beta^T \phi(x) = \sum_{i=1}^n \beta_i K(x_i, x)$$

Note that the magnitude of the coefficients  $\beta_i$  is related to the importance of the sample  $x_i$  in the prediction of the label  $y$ . The larger the magnitude of  $\beta_i$ , the more important the sample  $x_i$  is in the prediction of  $y$ .

### 1.6.5 Support Vector Regression (SVR)

Until now, we have used the following loss function to obtain the coefficients of the linear separator:

$$\mathcal{L}(w) = \|Xw - y\|_2^2 + \lambda \|w\|_2^2 \quad (1.83)$$

With  $\lambda = 0$ , we have the least squares solution. With  $\lambda > 0$ , we have the Ridge Regression solution. Now, let us propose a new loss function:

$$\mathcal{L}(w) = \frac{1}{n} \sum_{i=1}^n \max(0, |y_i - w^T x_i| - \varepsilon) + \lambda \|w\|_2^2 \quad (1.84)$$

We can even introduce a feature map  $\phi(\cdot)$ , as in the kernel methods:

$$\mathcal{L}(w) = \frac{1}{n} \sum_{i=1}^n \max(0, |y_i - w^T \phi(x_i)| - \varepsilon) + \lambda \|w\|_2^2 \quad (1.85)$$

This loss function is the loss function of the Support Vector Regression (SVR). The SVR is a technique that is used to predict the value of a continuous variable based on the value of one or more input variables. The SVR is similar to the Support Vector Machine (SVM), but it is used for regression problems instead of classification problems.

We can reformulate the SVR in terms of the dual coefficients  $\beta$ :

$$\mathcal{L}(\beta) = \frac{1}{2} \beta^T K \beta - \beta^T y + \varepsilon \|\beta\|_1 \quad \text{s.t.} \quad |\beta_i| \leq \frac{1}{2n\lambda} \quad (1.86)$$

We say that the support vectors are the samples  $x_i$  such that  $\beta_i \neq 0$ . With a bigger value of  $\varepsilon$ , we have a higher sparsity in the solution, meaning that we have fewer support vectors. Notice that when we don't have a feature map, the matrix  $K$  is equivalent to  $XX^T$ .



## Chapter 2

# Automatic differentiation

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### 2.1 Introduction: differentiation methods

Suppose that we have a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  that we want to differentiate. We have many methods to do so:

1. **Manual differentiation:** we can differentiate the function by hand.
  - Pros: it is exact, fast, and good for theory
  - Cons: it is error-prone, time-consuming, and difficult for complex functions
2. **Numerical differentiation:** we can approximate the derivative by finite differences.
  - Pros: it is easy to implement
  - Cons: it is slow, inaccurate (round-off and truncation errors), and not suitable for complex functions

Let us look at the example of finite differences in 1 dimension. We can formulate 3 ways to approximate the derivative of a function  $f : \mathbb{R} \rightarrow \mathbb{R}$  at a point  $x$ :

- Forward difference:  $f'(x) \approx \frac{f(x+h)-f(x)}{h}$
- Backward difference:  $f'(x) \approx \frac{f(x)-f(x-h)}{h}$
- Central difference:  $f'(x) \approx \frac{f(x+h)-f(x-h)}{2h}$

The problem with the Finite Differences method (FDM) is that it presents two types of errors:

- **Truncation error:** it is the error that arises due to the approximation of the derivative. It is a consequence of the fact that we are using a Taylor series to approximate the derivative. This error decreases as  $h$  decreases. For forward and backward differences, the error is  $O(h)$ , while for central differences, the error is  $O(h^2)$ .
- **Round-off error:** it is the error that arises due to the finite precision of the computer. It is a consequence of adding two numbers of different magnitudes, or subtracting two numbers that are close to each other. This error increases as  $h$  decreases.

The following graph represents the trade-off between the truncation and round-off errors:

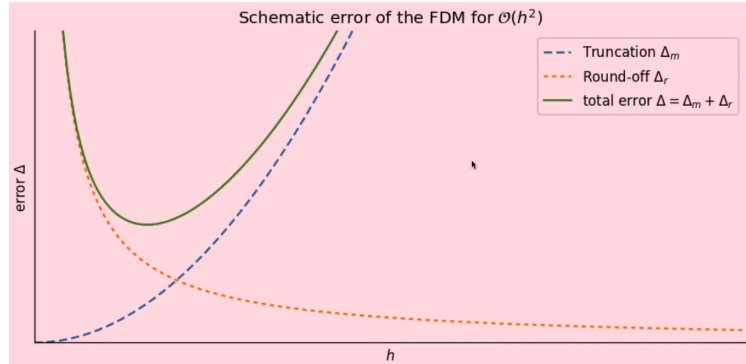


Figure 2.1: Trade-off between truncation and round-off errors

And we can actually see this in an example:

```

1 import numpy as np
2
3 def forward_diff(f, x, h):
4     return (f(x + h) - f(x)) / h
5
6 def centered_diff(f, x, h):
7     return (f(x + h) - f(x - h)) / (2 * h)
8
9 h = np.logspace(-16, 0, num=500, endpoint=True)
10 x = 1.0
11 D1f = forward_diff(np.sin, x, h)
12 D2f = centered_diff(np.sin, x, h)
13 err1 = np.abs(D1f - np.cos(x))
14 err2 = np.abs(D2f - np.cos(x))

```

Then, the plot of the errors is:

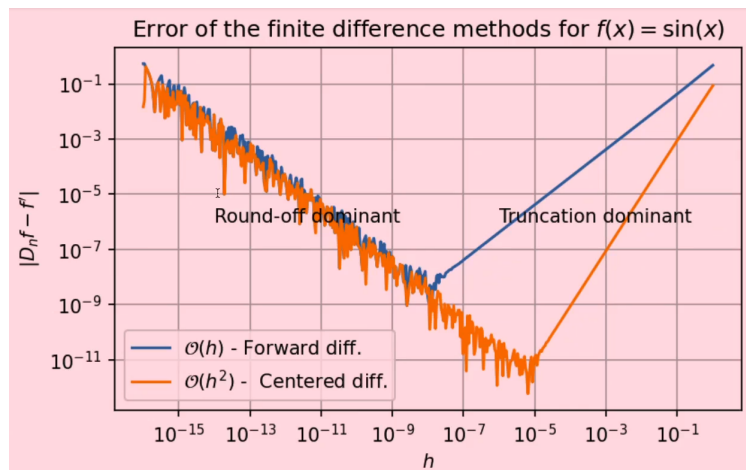


Figure 2.2: Errors of the finite differences method

3. **Symbolic differentiation:** we can use a computer algebra system (CAS) to compute the derivative symbolically.

- Pros: it is exact, good for theory.
- Cons: it is slow, memory-intensive, and can cause expression swell for complex functions.

4. **Automatic differentiation:** it is a technique that computes the derivative of a function by applying the chain rule.

- Pros: it is exact, fast, and general (good for complex functions).
- Cons: it is difficult to implement, and it requires a good understanding of the chain rule.

In this chapter, we will focus on automatic differentiation (AD). One of the main methods of AD is the backward propagation algorithm, which is used in deep learning to compute the gradients of the loss function with respect to the parameters of the model.

Automatic differentiation (AD) is a technique that it is based on the idea of splitting the computation of the derivative of a function into elementary operations:

- Unitary operations
- Exponentials, logarithms, trigonometric functions, etc.

## 2.2 Automatic differentiation: Wengert list

Automatic differentiation (AD) is a technique that computes the derivative of a function by applying the chain rule. It is based on the idea of splitting the computation of the derivative of a function into elementary operations. The main idea is to represent the function as a composition of elementary functions and then apply the chain rule to compute the derivative.

To implement AD, we take advantage of the evaluation trace of the function, also called the forward primal trace. This trace is a sequence of elementary operations that are applied to the input variables to compute the output of the function. It keeps track of the intermediate values of the variables and the operations that are applied to them. These variables are called primal variables, and are typically denoted by  $v_i$  for functions  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$ , following these rules:

- **Input variables:**  $v_{i-n} = x_i, \quad i = 1, \dots, n$
- **Intermediate variables:**  $v_i = g_i(v_{j_1}, \dots, v_{j_k}), \quad i = 1, \dots, \ell, \quad j_1, \dots, j_k \subseteq [1 - n, \ell]$
- **Output variables:**  $y_{m-i} = v_{\ell-i}, \quad i = m - 1, \dots, 0$

For example, consider the function  $f : \mathbb{R}^2 \rightarrow \mathbb{R}$  defined as  $f(x, y) = \sin((x_1 + x_2)x_2^2)$ . The evaluation trace of this function is:

- $v_{-1} = x_1, v_0 = x_2$
- $v_1 = v_{-1} + v_0$

- $v_2 = v_0^2$
- $v_3 = v_1 \cdot v_2$
- $v_4 = \sin(v_3)$
- $y_1 = v_4$

This trace is also called the Wengert list (as it was introduced by R. E. Wengert in 1964). We can represent the Wengert list as a directed acyclic graph (DAG), where nodes represent variables and edges describe the computational hierarchy of input to output transformations. This graph is also commonly referred to as the computational graph. For the previous example, we get:

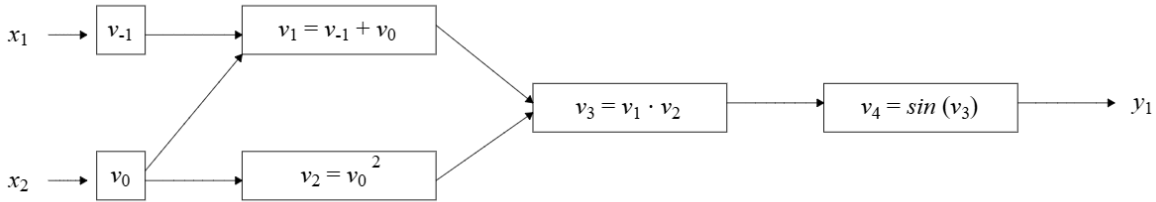


Figure 2.3: Comp. graph of the function  $f(x, y) = \sin((x_1 + x_2)x_2^2)$

These tools are essential for the implementation of AD, in particular, its two modes:

- **Forward mode**
- **Backward mode**

## 2.3 Forward mode of AD

The forward mode of automatic differentiation (AD) adopts the idea of the evaluation trace, but introduces a new concept: the tangent variables, usually denoted by  $\dot{v}_i$ . These tangents carry the derivative information of the primal variables with respect to a single input variable of interest. For example, if we are interested in computing the derivative of a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  with respect to the variable  $x_i$ , we write:

$$\dot{v}_j = \frac{\partial v_j}{\partial x_i} \quad j = 1 - n, \dots, \ell$$

The forward mode of AD computes the derivative of a function by applying the chain rule to the evaluation trace. It is based on the idea of propagating the derivative information from the input variables to the output variables. For this, we calculate what is called the forward tangent trace, by simply applying derivation rules to the evaluation trace. For example, consider the same function  $f(x, y) = \sin((x_1 + x_2)x_2^2)$ . The forward tangent trace of this function is:

- $\dot{v}_{-1} = \dot{x}_1, \dot{v}_0 = \dot{x}_2$
- $\dot{v}_1 = \dot{v}_{-1} + \dot{v}_0$

- $\dot{v}_2 = 2v_0\dot{v}_0$
- $\dot{v}_3 = v_2\dot{v}_1 + v_1\dot{v}_2$
- $\dot{v}_4 = \cos(v_3)\dot{v}_3$
- $\dot{y}_1 = \dot{v}_4$

In general, by following the computational graph, we can propose the following formula for the tangent variables:

$$\dot{v}_i = \sum_{j \in \{\text{predecessors of } i\}} \frac{\partial v_i}{\partial v_j} \dot{v}_j \quad (2.1)$$

In the case of the input variables (that don't have predecessors), we notice that the tangent variables, by definition, are always either 0 or 1. This is because the input variables are independent of each other, so their derivatives are 0, except for the variable of interest, which has a derivative of 1, as it is the variable with respect to which we are computing the derivative.

In fact, by setting the tangent variables of the input variables to 0 or 1, we can compute the derivative of the function with respect to any input variable, following the forward tangent trace. This is the main idea of the forward mode of AD.

For example, let us calculate the derivative of the previous function in the point  $(x, y) = (1, 2)$  with respect to the variable  $x_1$ . We can do this by setting the tangent variables of the input variables to 0 or 1, as well as the primal input variables to their values at the point  $(1, 2)$ . Then, we can compute the forward primal and tangent traces to get the derivative of the function with respect to  $x_1$ :

Forward Primal Trace	Output	Forward Tangent Trace	Output
$v_{-1} = x_1$	1	$\dot{v}_{-1} = \dot{x}_1$	1
$v_0 = x_2$	2	$\dot{v}_0 = \dot{x}_2$	0
$v_1 = v_{-1} + v_0$	$1 + 2 = 3$	$\dot{v}_1 = \dot{v}_{-1} + \dot{v}_0$	$1 + 0 = 1$
$v_2 = v_0^2$	$2^2 = 4$	$\dot{v}_2 = 2v_0 \cdot \dot{v}_0$	$2 \cdot 2 \cdot 0 = 0$
$v_3 = v_1 \cdot v_2$	$3 \cdot 4 = 12$	$\dot{v}_3 = v_2\dot{v}_1 + \dot{v}_2 \cdot v_1$	$4 \cdot 1 + 0 \cdot 3 = 4$
$v_4 = \sin v_3$	$\sin(12) = -0.54$	$\dot{v}_4 = \cos(v_3) \cdot \dot{v}_3$	$\cos(12) \cdot 4 = 3.38$
$y_1 = v_4$	-0.54	$\dot{y}_1 = \dot{v}_4$	3.38

Table 2.1: Example of forward mode of AD

The same can be done, setting the tangent variable  $\dot{v}_0 = 1$ , to compute the derivative of the function with respect to  $x_2$ . The result is  $\dot{y}_1 = -13.501$ . So in the end, we have computed the derivative of the function with respect to  $x_1$  and  $x_2$  at the point  $(1, 2)$ , which is 3.38 and  $-13.501$ , respectively.

This process is the essence of forward mode AD. At every elementary operation for a given function, compute intermediate variables (primals) by applying basic arithmetic operations, and in

synchrony, compute their derivatives (tangents).

We can generalize this procedure to compute Jacobians of functions  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  evaluated at a point  $x_0 \in \mathbb{R}^n$ . Remember that the Jacobian is a matrix that contains the partial derivatives of the function with respect to each input variable:

$$J_f(x_0) = \begin{bmatrix} \frac{\partial f_1}{\partial x_1} & \cdots & \frac{\partial f_1}{\partial x_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial f_m}{\partial x_1} & \cdots & \frac{\partial f_m}{\partial x_n} \end{bmatrix} \quad (2.2)$$

To compute the Jacobian of a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  evaluated at a point  $x_0 \in \mathbb{R}^n$ , we can use the forward mode of AD. We set the input variables to  $x_0$ , and the tangent variables equal to the canonical basis vectors  $e_i$  of  $\mathbb{R}^n$ , for  $i = 1, \dots, n$ . Every pass through the forward mode of AD for each  $e_i$  will give us the  $i$ -th column of the Jacobian.

Because the function  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  has  $n$  input variables, we need to perform  $n$  passes through the forward mode of AD to compute the Jacobian. So this procedure has a complexity of  $O(n)$ . This is the main drawback of the forward mode of AD: it is not efficient for functions with a large number of input variables, with respect to the output variables, i.e.,  $n \gg m$ .

### 2.3.1 Note: Forward mode and Jacobian-vector product (JVP)

In practice, we are often interested in computing the product of the Jacobian of a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  evaluated at a point  $x_0 \in \mathbb{R}^n$  with a vector  $r \in \mathbb{R}^n$ . This operation is called the Jacobian-vector product, and it is denoted by  $J_f(x_0) \cdot r$ .

The Jacobian-vector product can be computed using the forward mode of AD. We set the input variables to  $x_0$ , and the tangent variables equal to  $r$ . Then, a single pass through the forward mode of AD will give us the product of the Jacobian with the vector  $r$ . This is a more efficient way to compute the Jacobian-vector product, and sometimes it is more convenient than computing the full Jacobian.

## 2.4 Backward mode of AD

Backward mode of automatic differentiation (AD), also known as reverse mode of AD, is a technique that computes the derivative of a function by applying the chain rule in reverse order. It is based on the idea of propagating the derivative information from the output variables to the input variables.

This method introduces a new concept: the adjoint variables, usually denoted by  $\bar{v}_i$ . These adjoints carry the derivative information of a single output variable with respect to a variable (intermediate or input) of interest. For example, if we are interested in computing the derivative of the output variable  $y_j$  with respect to the variable  $v_i$ , we write:

$$\bar{v}_i = \frac{\partial y_j}{\partial v_i} \quad i = 1 - n, \dots, \ell$$

As in forward mode, in reverse mode we also perform a pass through the evaluation trace of the function. But this time, the adjoint variables are not computed alongside the primal variables. Instead, we store any dependency information that is needed to compute the adjoints in the computational graph (Wengert list). Then we perform a backward pass through the computational graph to compute the adjoints.

Given the computational graph, the adjoint variables can be computed by the following formula:

$$\bar{v}_i = \sum_{j \in \{\text{successors of } i\}} \frac{\partial v_j}{\partial v_i} \bar{v}_j \quad (2.3)$$

In the case of the output variables (that don't have successors), we notice that the adjoint variables are always either 0 or 1. This is because, by definition, the adjoint of the output variables are calculated with respect to the output variables, so their derivatives are 0, except for the variable of interest, which has a derivative of 1, as it is the variable with respect to which we are computing the derivative.

Notice that when we reach to the input variables, we have computed the derivative of the function with respect to all the output variables. This is the main idea of the backward mode of AD. At every elementary operation for a given function, compute intermediate variables (primals) by applying basic arithmetic operations, and store any dependency information that is needed to compute the adjoints. Then, perform a backward pass through the computational graph to compute the adjoints.

For example, consider the same function  $f(x, y) = \sin((x_1 + x_2)x_2^2)$ . The computational graph of this function is shown in Figure 2.3. The adjoint trace of this function is:

- $\bar{y}_1 = 1$
- $\bar{v}_4 = 1 \cdot \bar{y}_1$
- $\bar{v}_3 = \cos(v_3) \cdot \bar{v}_4$
- $\bar{v}_2 = \bar{v}_3 \cdot v_1$
- $\bar{v}_1 = \bar{v}_3 \cdot v_2$
- $\bar{v}_0 = 2v_0 \cdot \bar{v}_2 + 1 \cdot \bar{v}_1$
- $\bar{v}_{-1} = 1 \cdot \bar{v}_1$

For the same example, let us calculate the derivative of the function in the point  $(x, y) = (1, 2)$ :

Forward Primal Trace	Output	Backward Adjoint Trace	Output
$v_{-1} = x_1$	1	$\bar{v}_{-1} = 1 \cdot \bar{v}_1$	3.375
$v_0 = x_2$	2	$\bar{v}_0 = 2v_0 \cdot \bar{v}_2 + 1 \cdot \bar{v}_1$	$4 \cdot 2.351 + 3.375 = 13.501$
$v_1 = v_{-1} + v_0$	$1 + 2 = 3$	$\bar{v}_1 = \bar{v}_3 \cdot v_2$	$0.843 \cdot 4 = 3.375$
$v_2 = v_0^2$	$2^2 = 4$	$\bar{v}_2 = \bar{v}_3 \cdot v_1$	$0.843 \cdot 3 = 2.531$
$v_3 = v_1 \cdot v_2$	$3 \cdot 4 = 12$	$\bar{v}_3 = \cos(v_3) \cdot \bar{v}_4$	$\cos(12) \cdot 1 = 0.843$
$v_4 = \sin v_3$	$\sin(12) = -0.54$	$\bar{v}_4 = 1 \cdot \bar{y}_1$	1
$y_1 = v_4$	-0.54	$\bar{y}_1 = 1$	1

Table 2.2: Example of backward mode of AD

Notice that the result of the derivative of the function with respect to  $x_1$  and  $x_2$  at the point  $(1, 2)$  is the same as in the forward mode of AD. But in this case, we only needed to perform a single pass through the computational graph to compute the derivatives.

As well as in the forward mode of AD, we can generalize this procedure to compute Jacobians of functions  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  evaluated at a point  $x_0 \in \mathbb{R}^n$ . We just need to set the input variables to  $x_0$  for the evaluation trace, and the adjoint output variables to the canonical basis vectors  $e_i$  of  $\mathbb{R}^m$ , for  $j = 1, \dots, m$ . Then, perform a single pass through the computational graph to compute the Jacobian. For each  $e_i$ , we will get the  $i$ -th row of the Jacobian.

Because the function  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  has  $m$  output variables, we need to perform  $m$  passes through the backward mode of AD to compute the Jacobian. So this procedure has a complexity of  $O(m)$ . This is the main advantage of the backward mode of AD: it is more efficient for functions with a large number of output variables, with respect to the input variables, i.e.,  $n \gg m$ , which is usually the case in deep learning.

#### 2.4.1 Note: Backward mode and vector-Jacobian product (VJP)

In practice, we are often interested in computing the product of a vector  $r \in \mathbb{R}^m$  with the Jacobian of a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$  evaluated at a point  $x_0 \in \mathbb{R}^n$ . This operation is called the vector-Jacobian product, and it is denoted by  $r^T \cdot J_f(x_0) = J_f(x_0)^T \cdot r$ .

In the same way as in the forward mode of AD, the vector-Jacobian product can be computed using the backward mode of AD. We set the input variables to  $x_0$  for the evaluation trace, and the adjoint output variables to  $r$ . Then, a single pass through the computational graph will give us the product of the vector  $r^T$  with the Jacobian. This is a more efficient way to compute the vector-Jacobian product, and sometimes it is more convenient than computing the full Jacobian.

## 2.5 Dual numbers

We define a dual number as a pair of real numbers  $(a, b)$ , where  $a$  is the real part, and  $b$  is the dual part. It is expressed as follows:



$$z = a + b\varepsilon \quad (2.4)$$

where  $\varepsilon$  is a number such that  $\varepsilon \neq 0 \wedge \varepsilon^2 = 0$ . We can define the operations of addition and multiplication as follows:

- Addition:

$$(a + b\varepsilon) + (c + d\varepsilon) = (a + c) + (b + d)\varepsilon$$

- Multiplication:

$$(a + b\varepsilon) \cdot (c + d\varepsilon) = ac + (ad + bc)\varepsilon$$

Now, let us consider a generic function  $f(\cdot)$ . We will evaluate the function in a dual number  $z = x + \varepsilon$ . We can expand the function in a Taylor series around  $x$ :

$$f(z) = f(x + \varepsilon) = f(x) + f'(x)\varepsilon + O(\varepsilon^2) \quad (2.5)$$

Because  $\varepsilon^2 = 0$ , we have that:

$$f(z) = f(x) + f'(x)\varepsilon \quad (2.6)$$

This means that the dual part of the function evaluated in a dual number is the derivative of the function evaluated at the real part of the dual number. This is the key idea behind the dual numbers: we can use them to compute the derivative of a function by evaluating the function in a dual number.

Let us expand the idea to a general dual number  $a + b\varepsilon$ . We can evaluate the function  $f(\cdot)$  in the dual number as follows:

$$f(a + b\varepsilon) = f(a) + f'(a)b\varepsilon \quad (2.7)$$

With this property, we can also compute the derivative of a composite function. Suppose that we have two functions  $f(\cdot)$  and  $g(\cdot)$ . We can compute the derivative of the composite function  $h(\cdot) = f(g(\cdot))$  as follows:

$$h(x + \varepsilon) = f(g(x + \varepsilon)) = f(g(x) + g'(x)\varepsilon) = f(g(x)) + f'(g(x))g'(x)\varepsilon \quad (2.8)$$

This means that we can compute the derivative of a composite function by evaluating the functions in dual numbers.

We can also redefine the dual numbers to obtain the second derivative of a function. We just need to define  $\varepsilon^2 \neq 0$  and  $\varepsilon^3 = 0$ .



## Chapter 3

# Numerical Optimization and Training of Neural Networks

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### 3.1 Introduction: the Perceptron

The first model that introduced the idea of a neural network was the perceptron. The perceptron is a simple model that takes a set of inputs and produces an output. The model is based on the idea of a neuron in the brain, which receives signals from other neurons and produces an output signal.

The perceptron is a linear model that takes a set of binary inputs  $x = (x_1, x_2, \dots, x_n)$  and produces an output  $y$  in the following way:

$$y = \begin{cases} 1 & \text{if } w_1x_1 + w_2x_2 + \dots + w_nx_n + b > 0 \\ 0 & \text{otherwise} \end{cases} \quad (3.1)$$

where  $w = (w_1, w_2, \dots, w_n)$  are the weights of the model and  $b$  is the bias. Note that this can be represented as a vector product:

$$y = \begin{cases} 1 & \text{if } w \cdot x + b > 0 \\ 0 & \text{otherwise} \end{cases} \quad (3.2)$$

The perceptron is a simple model that can be used to classify data into two classes. The model is trained by adjusting the weights and bias to minimize the error on the training data.

Note that the perceptron is a linear model, which means that it can only learn linearly separable functions. This means that the model can only learn functions that can be separated by a hyperplane. If the data is not linearly separable, the perceptron will not be able to learn the function.

This problem is illustrated by the XOR function, which is not linearly separable.

#### 3.1.1 The XOR problem

The XOR logical function is a simple function that takes two binary inputs and produces an output, which is 1 if the inputs are different and 0 if the inputs are the same. The XOR function is defined

as follows:

$x_1$	$x_2$	$y$
0	0	0
0	1	1
1	0	1
1	1	0

Table 3.1: XOR function

Note that the XOR function is not linearly separable, i.e., it is not possible to separate the data points with a hyperplane. For instance, the following plot shows the XOR function in the input space:

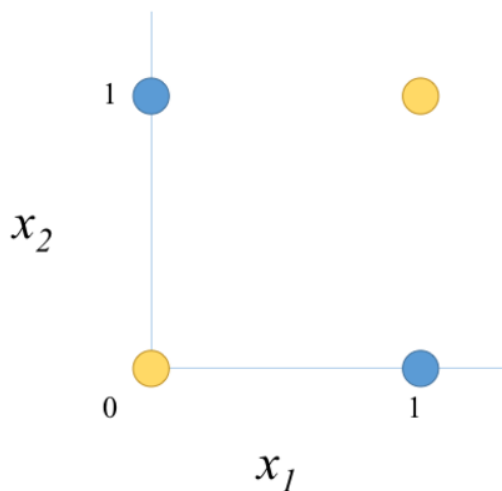


Figure 3.1: XOR function in the input space

As we can see, the data points are not linearly separable, which means that the perceptron alone cannot learn the XOR function. This is a limitation of the perceptron model. However, let us take a look at how we can overcome this limitation by using a more complex model: the multi-layer perceptron.

### 3.1.2 NAND gate

The NAND gate is a simple logical function that takes two binary inputs and produces an output, which is 0 if both inputs are 1 and 1 otherwise. The NAND gate is defined as follows:

$x_1$	$x_2$	$y$
0	0	1
0	1	1
1	0	1
1	1	0

Table 3.2: NAND gate

The NAND gate is an example of a function that is linearly separable, i.e., it can be separated by a hyperplane. This means that the perceptron can learn the NAND gate function. Notice that the NAND gate is an universal gate, which means that it can be used to implement any logical function. This is an important property of the NAND gate, as it allows us to build more complex functions using simple building blocks.

Because the NAND gate can represent any logical function by combining multiple gates, and the perceptron can learn the NAND gate, we can use the perceptron to learn any logical function, just by combining multiple perceptrons. This is the raw idea behind the multi-layer perceptron.

### 3.1.3 The problem of the step function

Notice that the perceptron uses a step function to produce the output. It takes the value 1 if the input is greater than 0 and 0 otherwise. This has some problems, as it implies that possible small changes in the input and weights can lead to a discontinuous change in the output, i.e., big jumps in the output. This can make the optimization process difficult.

To overcome this problem, we can use a different activation functions, such as the sigmoid function, which is a smooth function that takes values between 0 and 1. The idea of using a continuous activation function, combined with multiple layers of perceptrons, is the basis of the multi-layer perceptron.

## 3.2 The Multi-Layer Perceptron

The multi-layer perceptron (MLP) is a neural network model that consists of multiple layers of perceptrons. On the output layer, we use a different activation function than the step function, such as the sigmoid function. This allows the model to learn more complex functions, as it can approximate any continuous function.

We can represent the MLP model using a graph, where each node represents a perceptron and each edge represents a weight. The graph is organized in layers, where each layer is connected to the next layer. The input layer represents the input data, the output layer represents the output of the model, and the hidden layers represent the intermediate layers of the model.

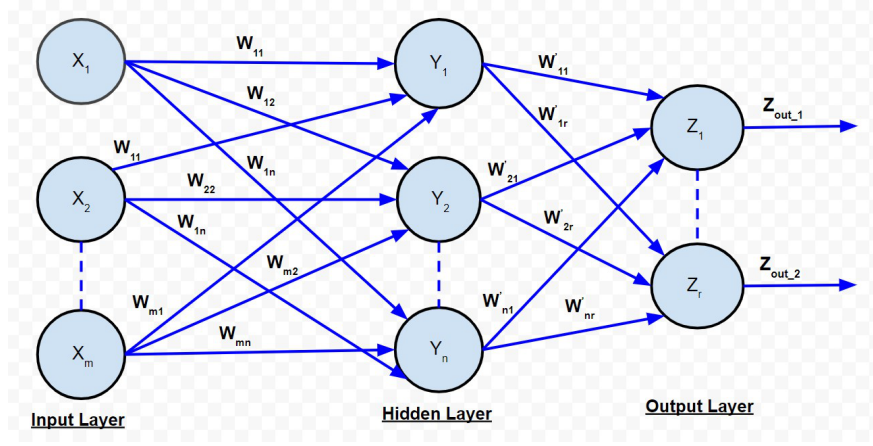


Figure 3.2: Graph representation of the MLP model

For a neuron  $j$  in layer  $\ell$ , we have the following:

$$z_j^\ell = \sum_k w_{jk}^\ell a_k^{\ell-1} + b_j^\ell \quad (3.3)$$

where  $z_j^\ell$  is the weighted sum of the inputs of neuron  $j$  in layer  $\ell$ ,  $w_{jk}^\ell$  is the weight of the connection between neuron  $k$  in layer  $\ell - 1$  and neuron  $j$  in layer  $\ell$ ,  $a_k^{\ell-1}$  is the output after activation of neuron  $k$  in layer  $\ell - 1$ , and  $b_j^\ell$  is the bias of neuron  $j$  in layer  $\ell$ .

Notice that in each output layer, we apply an activation function to the weighted sum of the inputs. This activation function is a non-linear function that allows the model to learn complex functions. This is represented by:

$$a_j^\ell = \sigma(z_j^\ell) \quad (3.4)$$

where  $\sigma$  is the activation function. In general, we can represent the weights as a matrix  $W^\ell$  and the biases as a vector  $b^\ell$ . This allows us to write the equation that represents the output of a particular layer as follows:

$$a^\ell = \sigma(W^\ell a^{\ell-1} + b^\ell) \quad (3.5)$$

where  $a^\ell$  is the output of layer  $\ell$ ,  $a^{\ell-1}$  is the output of layer  $\ell - 1$ , and  $\sigma$  is the activation function. This equation represents the forward pass of the model, i.e., the process of computing the output of the model given an input.

Now, on the output layer, we need a function that can represent the distance between the output of the model and the true output. This function is called the loss function, and it measures the error of the model. The main goal of training a neural network is to minimize the loss function.

### 3.3 Training a Neural Network

To train a neural network, we need to define a loss function that measures the error of the model. This function takes the output of the model and the true output and produces a value that represents how well the model is performing. In general, a loss function  $J$  should have the following properties:

1.  $J$  can be written as an average of the loss over the training data.

$$J = \frac{1}{N} \sum_{i=1}^N J_s$$

where  $N$  is the number of samples in the training data and  $J_s$  is the loss for sample  $s$ .

2.  $J$  is a function only of the output of the model and the true output.

Now, to train the model, we need to minimize the loss function. This is done by adjusting the weights and biases of the model to reduce the error. This process is called optimization. There are many optimization algorithms that can be used to train a neural network, but in general, all of them follow the same idea: compute the gradient of the loss function with respect to the weights and biases and update the weights and biases in the opposite direction of the gradient. This is called gradient descent.

Therefore, our aim is to compute the derivative of the loss function with respect to the weights and biases of the model, i.e., we need to compute:

$$\frac{\partial J}{\partial w_{jk}^\ell}, \frac{\partial J}{\partial b_j^\ell} \quad \forall \ell, j, k \quad (3.6)$$

where  $w_{jk}^\ell$  is the weight of the connection between neuron  $k$  in layer  $\ell - 1$  and neuron  $j$  in layer  $\ell$ , and  $b_j^\ell$  is the bias of neuron  $j$  in layer  $\ell$ .

For this, let us define the vector of errors (or sensitivities)  $\delta^\ell$  for layer  $\ell$  as follows:

$$\delta_j^\ell = \frac{\partial J}{\partial z_j^\ell} \quad \forall j \quad (3.7)$$

where  $z_j^\ell$  is the weighted sum of the inputs of neuron  $j$  in layer  $\ell$ . The vector of errors  $\delta^\ell$  represents how much the error changes with respect to the weighted sum of the inputs of neuron  $j$  in layer  $\ell$ .

Now, we will deduce 4 fundamental equations that will allow us to compute the gradient of the loss function with respect to the weights and biases of the model. These equations are called the backpropagation equations, and they are the key to training a neural network.

### 3.3.1 Backpropagation equations

1. Compute the error of the output layer  $\delta^L$ :

$$\delta_j^L = \frac{\partial J}{\partial z_j^L} = \sum_k \frac{\partial J}{\partial a_k^L} \cdot \frac{\partial a_k^L}{\partial z_j^L}$$

Because  $a_k^L$  only depends on  $z_k^L$  if  $k = j$ , we have:

$$\delta_j^L = \frac{\partial J}{\partial a_j^L} \cdot \frac{\partial a_j^L}{\partial z_j^L}$$

Notice that:

$$\frac{\partial a_j^L}{\partial z_j^L} = \sigma'(z_j^L)$$

where  $\sigma'$  is the derivative of the activation function. Then, we have:

$$\delta_j^L = \frac{\partial J}{\partial a_j^L} \cdot \sigma'(z_j^L)$$

Notice that the value of  $z_j^L$  has already been computed in the forward pass. Rewriting the equation in vector form, we have:

$$\delta^L = \nabla_{a^L} J \odot \sigma'(z^L) \quad (3.8)$$

where  $\nabla_a J$  is the gradient of the loss function with respect to the output of the model and  $\odot$  is the element-wise product.

2. Compute the error of the hidden layers  $\delta^\ell$  with respect to  $\delta^{\ell+1}$ :

$$\delta_j^\ell = \frac{\partial J}{\partial z_j^\ell} = \sum_k \frac{\partial J}{\partial z_k^{\ell+1}} \cdot \frac{\partial z_k^{\ell+1}}{\partial z_j^\ell}$$

Notice that:

$$\frac{\partial J}{\partial z_k^{\ell+1}} = \delta_k^{\ell+1}$$

Then, we have:

$$\delta_j^\ell = \sum_k \delta_k^{\ell+1} \cdot \frac{\partial z_k^{\ell+1}}{\partial z_j^\ell}$$

Now, we need to compute the derivative of  $z_k^{\ell+1}$  with respect to  $z_j^\ell$ . Notice that:



$$z_k^{\ell+1} = \sum_m w_{km}^{\ell+1} a_m^\ell + b_k^{\ell+1} = \sum_m w_{km}^{\ell+1} \sigma(z_m^\ell) + b_k^{\ell+1}$$

Then, we have:

$$\frac{\partial z_k^{\ell+1}}{\partial z_j^\ell} = w_{kj}^{\ell+1} \cdot \sigma'(z_j^\ell)$$

Finally, we have:

$$\delta_j^\ell = \sum_k \delta_k^{\ell+1} \cdot w_{kj}^{\ell+1} \cdot \sigma'(z_j^\ell)$$

Rewriting the equation in vector form, we have:

$$\delta^\ell = ((W^{\ell+1})^T \delta^{\ell+1}) \odot \sigma'(z^\ell) \quad (3.9)$$

3. Compute the gradient of the loss function with respect to  $b^\ell$ :

$$\frac{\partial J}{\partial b_j^\ell} = \frac{\partial J}{\partial z_j^\ell} \cdot \frac{\partial z_j^\ell}{\partial b_j^\ell}$$

Notice that:

$$z_j^\ell = \sum_k w_{jk}^\ell a_k^{\ell-1} + b_j^\ell$$

Then, we have:

$$\frac{\partial z_j^\ell}{\partial b_j^\ell} = 1$$

Finally, we have:

$$\frac{\partial J}{\partial b_j^\ell} = \delta_j^\ell$$

Rewriting the equation in vector form, we have:

$$\nabla_{b^\ell} J = \delta^\ell \quad (3.10)$$

4. Compute the gradient of the loss function with respect to  $W^\ell$ :

$$\frac{\partial J}{\partial w_{jk}^\ell} = \frac{\partial J}{\partial z_j^\ell} \cdot \frac{\partial z_j^\ell}{\partial w_{jk}^\ell}$$

Notice that:

$$z_j^\ell = \sum_k w_{jk}^\ell a_k^{\ell-1} + b_j^\ell$$

Then, we have:

$$\frac{\partial z_j^\ell}{\partial w_{jk}^\ell} = a_k^{\ell-1}$$

Finally, we have:

$$\frac{\partial J}{\partial w_{jk}^\ell} = \delta_j^\ell \cdot a_k^{\ell-1}$$

Rewriting the equation in matrix form, we have:

$$\nabla_{W^\ell} J = \delta^\ell (a^{\ell-1})^T \quad (3.11)$$

These are the backpropagation equations, which allow us to compute the gradient of the loss function with respect to the weights and biases of the model. This is the key to training a neural network. By computing the gradient of the loss function and updating the weights and biases in the opposite direction of the gradient, we can minimize the error of the model and learn complex functions.

Summarizing the equations:

1. Compute the error of the output layer:

$$\delta^L = \nabla_{a^L} J \odot \sigma'(z^L)$$

2. Compute the error of the hidden layers:

$$\delta^\ell = ((W^{\ell+1})^T \delta^{\ell+1}) \odot \sigma'(z^\ell)$$

3. Compute the gradient of the loss function with respect to  $b^\ell$ :

$$\nabla_{b^\ell} J = \delta^\ell$$

4. Compute the gradient of the loss function with respect to  $W^\ell$ :

$$\nabla_{W^\ell} J = \delta^\ell (a^{\ell-1})^T$$

Then, the backpropagation algorithm can be summarized as follows:

1.  $x$  is the input to the model. Compute  $a^1$
2. For each layer  $\ell$  from 2 to  $L$ :
  - (a) Compute  $z^\ell = W^\ell a^{\ell-1} + b^\ell$
  - (b) Compute  $a^\ell = \sigma(z^\ell)$
3. Compute the loss function  $J$
4. Compute the error of the output layer  $\delta^L$
5. For each layer  $\ell$  from  $L - 1$  to 1:
  - (a) Compute the error of the hidden layers  $\delta^\ell$
  - (b) Compute the gradient of the loss function with respect to  $b^\ell$
  - (c) Compute the gradient of the loss function with respect to  $W^\ell$
  - (d) Update the weights and biases:
    - $W^\ell = W^\ell - \alpha \nabla_{W^\ell} J$
    - $b^\ell = b^\ell - \alpha \nabla_{b^\ell} J$

where  $\alpha$  is the learning rate.

### 3.4 Activation functions

The activation function of a neuron is a non-linear function that takes the weighted sum of the inputs of the neuron and produces an output. The activation function is a key component of a neural network, as it allows the model to learn complex functions. There are many activation functions that can be used in a neural network, but some of the most common are:

1. Sigmoid function:

$$\sigma(z) = \frac{1}{1 + e^{-z}} \quad (3.12)$$

The sigmoid function takes values between 0 and 1, which makes it useful for binary classification problems. However, the sigmoid function has some problems, such as the vanishing gradient problem, which can make training difficult.

2. Hyperbolic tangent function:

$$\tanh(z) = \frac{e^z - e^{-z}}{e^z + e^{-z}} \quad (3.13)$$

The hyperbolic tangent function is similar to the sigmoid function, but it takes values between -1 and 1. This can make training easier, as the output is centered around 0. However, the hyperbolic tangent function also has the vanishing gradient problem.

3. Rectified Linear Unit (ReLU) function:

$$\text{ReLU}(z) = \max(0, z) \quad (3.14)$$

The ReLU function is a simple function that takes the maximum between 0 and the input. The ReLU function is widely used in practice, as it is simple and efficient. However, the ReLU function has some problems, such as the dying ReLU problem, which can make some neurons inactive.

4. Leaky ReLU function:

$$\text{Leaky ReLU}(z) = \begin{cases} z & \text{if } z > 0 \\ \alpha z & \text{otherwise} \end{cases} \quad (3.15)$$

The Leaky ReLU function is a variant of the ReLU function that allows a small gradient when the input is negative. This can help to overcome the dying ReLU problem.

5. Swish function (Google):

$$\text{swish}(z) = \frac{z}{1 + e^{-z}} \quad (3.16)$$

The swish function is a new activation function that has been proposed recently. The swish function is similar to the sigmoid function, but it has a different shape that can make training easier. The swish function has been shown to perform well on a variety of tasks.

6. Softmax function:

$$\text{softmax}(z)_i = \frac{e^{z_i}}{\sum_j e^{z_j}} \quad (3.17)$$

The softmax function is a generalization of the sigmoid function that takes a vector of inputs and produces a vector of outputs that sum to 1. The softmax function is often used in the output layer of a neural network to produce probabilities for a classification problem. The softmax function is useful when the output of the model needs to be interpreted as a probability distribution.

## 3.5 Loss functions

The loss function of a neural network is a function that measures the error of the model. The loss function takes the output of the model and the true output and produces a value that represents how well the model is performing. There are many loss functions that can be used in a neural network, but some of the most common are:

1. Mean Squared Error (MSE):

$$J = \frac{1}{2N} \sum_{i=1}^N (y_i - \hat{y}_i)^2 \quad (3.18)$$

The mean squared error (MSE) is a simple loss function that measures the squared difference between the output of the model and the true output. The MSE is often used in regression problems, where the output of the model is a continuous value.

2. Mean Absolute Error (MAE):

$$J = \frac{1}{N} \sum_{i=1}^N |y_i - \hat{y}_i| \quad (3.19)$$

The mean absolute error (MAE) is a loss function that measures the absolute difference between the output of the model and the true output. The MAE is often used in regression problems, where the output of the model is a continuous value.

An issue of the MSE is that it can be sensitive to outliers. The MAE is more robust to outliers, but it can be harder to optimize. In practice, the choice of loss function depends on the problem at hand.

Another big issue of using the MSE is that it can lead to the vanishing gradient problem. This is because when we derive the MSE with respect to the weights and biases, we get an expression that depends on the derivative of the activation function. If the activation function is the sigmoid function, the derivative of the activation function can be very small, which can make the optimization process difficult.

To overcome this problem, we can use the cross-entropy loss function.

### 3.5.1 Cross-entropy loss function

The cross-entropy loss function is a loss function that measures the error of the model by comparing the output of the model with the true output. The cross-entropy loss function is often used in classification problems, where the output of the model is a probability distribution. The cross-entropy loss function is defined as follows:

$$J = -\frac{1}{N} \sum_{i=1}^N \sum_j y_{ij} \log(\hat{y}_{ij}) \quad (3.20)$$

where  $y_{ij}$  is the true output of class  $j$  for sample  $i$  and  $\hat{y}_{ij}$  is the predicted output of class  $j$  for sample  $i$ .

Let us see why the cross-entropy loss function is useful for avoiding the vanishing gradient problem. For this, let us take a simple binary classification problem, with the following loss function:

$$J = -\frac{1}{N} \sum_{i=1}^N y_i \log(\hat{y}_i) + (1 - y_i) \log(1 - \hat{y}_i) \quad (3.21)$$

where  $y_i$  is the true output for sample  $i$  and  $\hat{y}_i$  is the predicted output for sample  $i$ .

Now, let us compute the derivative of the loss function with respect to one of the weights of the model. For this, we have:

$$\frac{\partial J}{\partial w} = \frac{\partial J}{\partial \hat{y}} \cdot \frac{\partial \hat{y}}{\partial z} \cdot \frac{\partial z}{\partial w}$$

where  $z$  is the weighted sum of the inputs of the neuron and  $\hat{y}$  is the output of the neuron after activation. So we have:

$$\frac{\partial J}{\partial w} = \frac{\partial J}{\partial \hat{y}} \cdot \sigma'(z) \cdot x$$

where  $x$  is the input of the neuron. Now, let us compute the derivative of the loss function with respect to the output of the neuron:

$$\frac{\partial J}{\partial \hat{y}} = -\frac{1}{N} \sum_{i=1}^N \left( \frac{y_i}{\hat{y}_i} - \frac{1-y_i}{1-\hat{y}_i} \right)$$

Notice that we can rewrite  $\hat{y}_i$  as  $\sigma(z_i)$ , where  $z_i$  is the weighted sum of the inputs of the neuron. Then, simplifying, we have:

$$\frac{\partial J}{\partial \hat{y}} = -\frac{1}{N} \sum_{i=1}^N \frac{\sigma(z_i) - y_i}{\sigma(z_i)(1 - \sigma(z_i))}$$

Combining the terms, and using the fact that  $\sigma'(z) = \sigma(z)(1 - \sigma(z))$  for the sigmoid, we have:

$$\frac{\partial J}{\partial w} = -\frac{1}{N} \sum_{i=1}^N (\sigma(z_i) - y_i) \cdot x$$

This is the gradient of the loss function with respect to the weights of the model. Notice that the gradient depends on the difference between the output of the model and the true output. This is the key to avoiding the vanishing gradient problem.

The cross-entropy loss function is a generalization of the binary classification loss function to multi-class classification problems. It avoids the vanishing gradient problem when using the softmax activation function in the output layer, which is the generalization of the sigmoid function to multi-class classification.

How do we obtain the cross-entropy loss function? Let us derive it.

### 3.5.2 Derivation of the cross-entropy loss function

The idea is to avoid the vanishing gradient problem by letting the derivative of the loss function to not depend on the derivative of the activation function. In other terms, we want that these terms:

$$\frac{\partial J}{\partial w} = (a - y)\sigma'(z)x, \quad \frac{\partial J}{\partial b} = (a - y)\sigma'(z)$$

do not depend on  $\sigma'(z)$ , i.e., we want that:

$$\frac{\partial J}{\partial w} = (a - y)x, \quad \frac{\partial J}{\partial b} = (a - y) \tag{3.22}$$

where  $a$  is the output of the neuron after activation,  $y$  is the true output,  $z$  is the weighted sum of the inputs of the neuron,  $w$  is the weight of the connection between the neuron and the previous

layer, and  $b$  is the bias of the neuron.

Let us take the derivative with respect to the bias  $b$ , and notice the following:

$$\begin{aligned}\frac{\partial J}{\partial b} &= \frac{\partial J}{\partial a} \cdot \frac{\partial a}{\partial z} \\ &= \frac{\partial J}{\partial a} \cdot \sigma'(z) = \frac{\partial J}{\partial a} \cdot a(1-a)\end{aligned}$$

Now, we want that this term does not depend on  $\sigma'(z)$ , and for that, we use the equation (3.22):

$$\frac{\partial J}{\partial a} a(1-a) = a - y$$

Then, we have:

$$\frac{\partial J}{\partial a} = \frac{a - y}{a(1-a)}$$

Integrating, we obtain:

$$J = -(y \log(a) + (1-y) \log(1-a)) + C$$

which is the cross-entropy loss function.

## 3.6 Regularization

Regularization is a technique used to prevent overfitting in a neural network. Overfitting occurs when the model learns the training data too well and fails to generalize to new data. Regularization is a way to prevent overfitting by adding a penalty term to the loss function. This penalty term discourages the model from learning complex functions that fit the training data too well.

There are many regularization techniques that can be used in a neural network, but some of the most common are:

### 1. L2 regularization:

$$J = J_0 + \lambda \frac{1}{2N} \sum_{\ell} \sum_{j,k} (w_{jk}^{\ell})^2 \quad (3.23)$$

L2 regularization is a technique that adds a penalty term to the loss function that is proportional to the square of the weights of the model. L2 regularization encourages the model to learn small weights, which can help to prevent overfitting by reducing the complexity of the model.

Notice that, when computing the update of the weights, we have:

$$w^{(k+1)} = w^{(k)} - \alpha \nabla_{w^{(k)}} J = w^{(k)} - \alpha \nabla_{w^{(k)}} J_0 - \alpha \frac{\lambda}{N} w^{(k)}$$

which is equivalent to:

$$w^{(k+1)} = (1 - \alpha \frac{\lambda}{N})w^{(k)} - \alpha \nabla_{w^{(k)}} J_0$$

When choosing appropriate values for  $\lambda$ , notice that the term  $\alpha \frac{\lambda}{N} < 1$ . This causes the weights to decay over time, which can help to prevent overfitting. This is called weight decaying.

## 2. L1 regularization:

$$J = J_0 + \lambda \frac{1}{N} \sum_{\ell} \sum_j |w_j^{\ell}| \quad (3.24)$$

L1 regularization is a technique that adds a penalty term to the loss function that is proportional to the absolute value of the weights of the model. L1 regularization encourages the model to learn sparse weights, i.e., weights that are close to zero. This can help to prevent overfitting by reducing the complexity of the model.

Notice that, when computing the update of the weights, we have:

$$w^{(k+1)} = w^{(k)} - \alpha \nabla_{w^{(k)}} J = w^{(k)} - \alpha \nabla_{w^{(k)}} J_0 - \alpha \frac{\lambda}{N} \text{sign}(w^{(k)})$$

This means that the weights are updated by subtracting (or adding) a constant value from the weights. This can help to prevent overfitting by encouraging the model to learn sparse weights.

## 3. Dropout regularization:

Dropout regularization is a technique that randomly sets a fraction of the neurons in the model to zero during training. This can help to prevent overfitting by reducing the complexity of the model. Dropout regularization is a simple and effective technique that can be used in a variety of models.

The idea is, for some layer  $\ell$ , to set a fraction  $p$  of the neurons to zero. This can be done by multiplying the output of the layer by a mask  $m$  that has values of 0 or 1. The mask  $m$  is generated by sampling from a Bernoulli distribution with probability  $p$ . Then, the output of the layer is given by:

$$a^{\ell} = m \odot \sigma(z^{\ell})$$

where  $\odot$  is the element-wise product. This can help to prevent overfitting by reducing the complexity of the model.



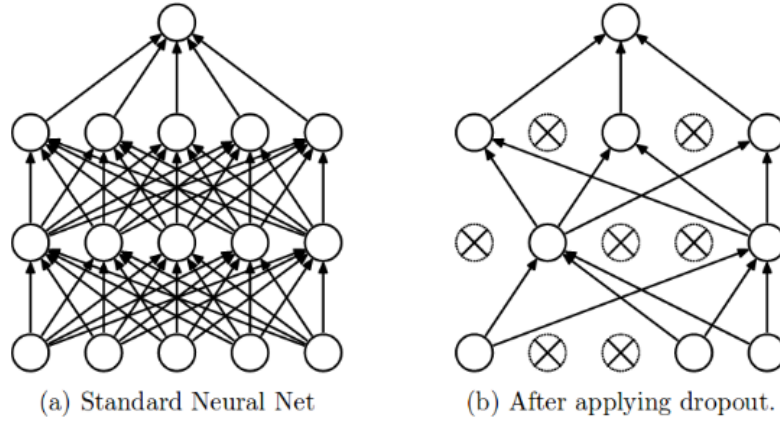


Figure 3.3: Dropout regularization

### 3.7 Optimization methods: Gradient Descent

The process of training a neural network involves solving an optimization problem to minimize the loss function. There are many optimization algorithms that can be used to train a neural network. To study these algorithms, let us first define some basic concepts.

Let us consider a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  convex and differentiable. A convex function has these properties:

1. The domain of  $f$  is a convex set, i.e., for any  $x, y \in \text{dom}(f)$ :

$$\lambda x + (1 - \lambda)y \in \text{dom}(f)$$

2. For any  $x, y \in \text{dom}(f)$  and  $\lambda \in [0, 1]$ :

$$f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y)$$

Now, let us propose the following optimization problem:

$$\min_{x \in \mathbb{R}^n} f(x) \tag{3.25}$$

An idea to solve this problem is to propose an iterative algorithm such that, starting from an initial point  $x^{(0)}$ , we update the point  $x^{(k)}$  to  $x^{(k+1)}$  in such a way that:

$$x^{(k+1)} = \min_{x \in \mathbb{R}^n} \tilde{f}_k(x)$$

where  $\tilde{f}_k$  is a function that approximates  $f$  at the point  $x^{(k)}$ . Some ways to approximate  $f$  are:

#### 1. Linear approximation

A linear approximation of  $f$  at the point  $x^{(k)}$  is given by:

$$\tilde{f}_k(x) = c_k + b_k^T(x - x^{(k)})$$

where  $c_k$  is a scalar and  $b_k$  is a vector. Right away, we notice that this problem is unbounded, and we need to add some constraints to make it feasible.

## 2. Quadratic approximation

A quadratic approximation of  $f$  at the point  $x^{(k)}$  is given by:

$$\tilde{f}_k(x) = c_k + b_k^T(x - x^{(k)}) + \frac{1}{2}(x - x^{(k)})^T A_k (x - x^{(k)})$$

where  $A_k$  is a matrix. Let us use this form of approximation, and set some conditions:

- (a)  $f(x^{(k)}) = \tilde{f}_k(x^{(k)})$
- (b)  $\nabla \tilde{f}_k(x^{(k)}) = \nabla f(x^{(k)})$

From the first condition, we have:

$$c_k = f(x^{(k)})$$

From the second condition, we have:

$$b_k = \nabla f(x^{(k)})$$

Then, let us set  $A_k = \frac{1}{\eta}I$ , where  $\eta$  is a positive scalar. Then, we have:

$$\tilde{f}_k(x) = f(x^{(k)}) + \nabla f(x^{(k)})^T(x - x^{(k)}) + \frac{1}{2\eta}\|x - x^{(k)}\|^2$$

Now, notice that, since the function  $\tilde{f}_k$  is convex, we can minimize it by setting the gradient to zero. Then, we have:

$$\begin{aligned}\nabla \tilde{f}_k(x) &= \nabla f(x^{(k)}) + \frac{1}{\eta}(x - x^{(k)}) = 0 \\ \Rightarrow x^{(k+1)} &= x^{(k)} - \eta \nabla f(x^{(k)})\end{aligned}$$

This is the **gradient descent algorithm**, and we call  $\eta$  the learning rate.

### 3.7.1 Gradient descent algorithm

More formally, the gradient descent algorithm is defined as follows, for a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ :

---

**Algorithm 6** Gradient descent algorithm

---

**Require:**  $x^{(0)}, \eta, \varepsilon, T$

```
1: for  $k = 0, 1, 2, \dots, T$  do  
2:    $x^{(k+1)} = x^{(k)} - \eta \nabla f(x^{(k)})$   
3:   if  $\text{STOP}(\varepsilon)$  then  
4:     break  
5:   end if  
6: end for
```

---

In this case, we can define the stopping criterion as:

$$\text{STOP}(\varepsilon) = \begin{cases} \text{true} & \text{if } \|\nabla f(x^{(k)})\| \leq \varepsilon \\ \text{false} & \text{otherwise} \end{cases} \quad (3.26)$$

Or:

$$\text{STOP}(\varepsilon) = \begin{cases} \text{true} & \text{if } |f(x^{(k+1)}) - f(x^{(k)})| \leq \varepsilon \\ \text{false} & \text{otherwise} \end{cases} \quad (3.27)$$

Or:

$$\text{STOP}(\varepsilon) = \begin{cases} \text{true} & \text{if } \|x^{(k+1)} - x^{(k)}\| \leq \varepsilon \\ \text{false} & \text{otherwise} \end{cases} \quad (3.28)$$

Now, to choose the learning rate  $\eta$ , we have many options. One of them is to use a fixed learning rate. Another option is to use a learning rate that decreases over time. This is called the learning rate schedule, and effectively, just decreases the learning rate depending on the iteration. For example:

$$\eta_k = \frac{\eta_0}{1 + k} \quad (3.29)$$

where  $\eta_0$  is the initial learning rate. Another option is to use the **Backtracking line search**, which is an algorithm that finds an effective learning rate at each iteration. It works as follows:

---

**Algorithm 7** Backtracking line search

---

**Require:**  $x^{(k)}, \eta_0, \alpha, \beta$

```
1:  $\eta = \eta_0$   
2: while  $f(x^{(k)} - \eta \nabla f(x^{(k)})) > f(x^{(k)}) - \alpha \eta \|\nabla f(x^{(k)})\|^2$  do  
3:    $\eta = \beta \eta$   
4: end while
```

---

We usually set  $\alpha = 0.5$  and  $\beta = 0.5$ .

Another way of setting the learning rate is to use the **Exact line search**, which is an algorithm that finds the optimal learning rate at each iteration. It works by solving the following optimization problem at each iteration:

$$\eta^{(k)} = \arg \min_{\eta > 0} f(x^{(k)} - \eta \nabla f(x^{(k)})) \quad (3.30)$$

This is a one-dimensional optimization problem that can be solved using standard optimization, but it can be computationally expensive, as it may not have a closed-form solution. In practice, the exact line search is rarely used.

### Wolfe conditions

In general, a good learning rate is one that satisfies the **Wolfe conditions**. These are a set of conditions that ensure that the learning rate is not too large or too small. The Wolfe conditions are defined as follows:

For  $0 < c_1 < c_2 < 1$ , and a descent direction  $d$ , the Wolfe conditions are:

1. **Sufficient decrease condition:**

$$f(x + \eta d) \leq f(x) + c_1 \eta \nabla f(x)^T d$$

This condition ensures that the learning rate is not too large, i.e., that the function value decreases sufficiently at each iteration.

2. **Curvature condition:**

$$\nabla f(x + \eta d)^T d \geq c_2 \nabla f(x)^T d$$

This condition ensures that the learning rate is not too small, i.e., that the gradient of the function is not too flat at each iteration.

The Wolfe conditions are a good way to choose the learning rate, as they ensure that the learning rate is not too large or too small. In practice, the Wolfe conditions are often used in conjunction with the backtracking line search to find an effective learning rate at each iteration.

### 3.7.2 Convergence of the gradient descent algorithm

Let us state the following theorem:

**Theorem 1.** *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  be a convex function that is differentiable and has a Lipschitz continuous gradient with constant  $L$ . Let  $\{x^{(k)}\}$  be the sequence generated by the gradient descent algorithm with learning rate  $\eta \leq \frac{1}{L}$ . Then, we have:*

$$f(x^{(k)}) - f(x^*) \leq \frac{\|x^{(0)} - x^*\|^2}{2\eta k} \quad (3.31)$$

where  $x^*$  is the optimal solution of the optimization problem.

*Proof.* This proof is different from the one presented in the professor's lecture. Nonetheless, it is a valid proof. It is based on the following lemma:

**Lemma 1** (Descent lemma). *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  be a differentiable function that has a Lipschitz continuous gradient with constant  $L$ . Then, for any  $x, y \in \mathbb{R}^n$ , we have:*

$$f(y) \leq f(x) + \nabla f(x)^T(y - x) + \frac{L}{2}\|y - x\|^2 \quad (3.32)$$

*Proof.* Let us define the function  $g(t) = f((1 - t)x + ty)$ . Then, by the fundamental theorem of calculus, we have:

$$\begin{aligned} g(1) &= g(0) + \int_0^1 g'(t)dt \\ f(y) &= f(x) + \int_0^1 \nabla f((1 - t)x + ty)^T(y - x)dt \\ &= f(x) + \nabla f(x)^T(y - x) + \int_0^1 (\nabla f((1 - t)x + ty) - \nabla f(x))^T(y - x)dt \end{aligned}$$

By the Cauchy-Schwarz inequality (more generally, the Hölder's inequality), we have:

$$(\nabla f((1 - t)x + ty) - \nabla f(x))^T(y - x) \leq \|\nabla f((1 - t)x + ty) - \nabla f(x)\| \|y - x\|$$

And then, by the Lipschitz continuity of the gradient, we have:

$$\begin{aligned} \|\nabla f((1 - t)x + ty) - \nabla f(x)\| &\leq L\|(1 - t)x + ty - x\| \\ &= Lt\|y - x\| \end{aligned}$$

Then, we have:

$$\begin{aligned} f(y) &\leq f(x) + \nabla f(x)^T(y - x) + \int_0^1 Lt\|y - x\|^2 dt \\ &= f(x) + \nabla f(x)^T(y - x) + \frac{L}{2}\|y - x\|^2 \end{aligned}$$

This is the desired result. □

Observe that a direct consequence of the descent lemma is that the gradient descent algorithm is in fact a descent algorithm, i.e., it decreases the value of the function at each iteration. We can propose our second lemma:

**Lemma 2** (Gradient descent lemma). *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  be a function that is differentiable and has a Lipschitz continuous gradient with constant  $L$ . Let  $\{x^{(k)}\}$  be the sequence generated by the gradient descent algorithm with learning rate  $\eta \leq \frac{1}{L}$ . Then, we have:*

$$f(x^{(k+1)}) \leq f(x^{(k)}) - \frac{\eta}{2}\|\nabla f(x^{(k)})\|^2 \quad (3.33)$$

*Proof.* By the descent lemma, we have:

$$\begin{aligned}
f(x^{(k+1)}) &\leq f(x^{(k)}) + \nabla f(x^{(k)})^T (x^{(k+1)} - x^{(k)}) + \frac{L}{2} \|x^{(k+1)} - x^{(k)}\|^2 \\
&= f(x^{(k)}) - \eta \|\nabla f(x^{(k)})\|^2 + \frac{L}{2} \|\eta \nabla f(x^{(k)})\|^2 \\
&= f(x^{(k)}) - \left(1 - \frac{L\eta}{2}\right) \eta \|\nabla f(x^{(k)})\|^2
\end{aligned}$$

Now notice that, since  $\eta \leq \frac{1}{L}$ , we have:

$$1 - \frac{L\eta}{2} \geq \frac{1}{2}$$

Then, we have:

$$f(x^{(k+1)}) \leq f(x^{(k)}) - \frac{\eta}{2} \|\nabla f(x^{(k)})\|^2$$

This is the desired result. □

Now, let us prove the theorem. To do so, let us define the following operator:

$$V_t = \frac{\|x^{(t)} - x^*\|^2}{2}$$

Then, we have:

$$\begin{aligned}
V_{t+1} &= \frac{1}{2} \|x^{(t+1)} - x^*\|^2 = \frac{1}{2} \|x^{(t)} - \eta \nabla f(x^{(t)}) - x^*\|^2 \\
&= \frac{1}{2} \left[ \|x^{(t)} - x^*\|^2 + \eta^2 \|\nabla f(x^{(t)})\|^2 - 2\eta \nabla f(x^{(t)})^T (x^{(t)} - x^*) \right] \\
&= V_t + \frac{\eta^2}{2} \|\nabla f(x^{(t)})\|^2 - \eta \nabla f(x^{(t)})^T (x^{(t)} - x^*)
\end{aligned}$$

Here, we use the fact that  $f$  is convex, meaning that:

$$\begin{aligned}
f(x^*) &\geq f(x^{(t)}) + \nabla f(x^{(t)})^T (x^* - x^{(t)}) \\
\Rightarrow \nabla f(x^{(t)})^T (x^{(t)} - x^*) &\geq f(x^{(t)}) - f(x^*)
\end{aligned}$$

Then, we have:

$$V_{t+1} \leq V_t + \frac{\eta^2}{2} \|\nabla f(x^{(t)})\|^2 - \eta (f(x^{(t)}) - f(x^*))$$

Now, notice that, by the gradient descent lemma, we have:

$$\frac{\eta}{2} \|\nabla f(x^{(t)})\|^2 \leq f(x^{(t)}) - f(x^{(t+1)})$$

Replace this in the previous equation, we have:

$$\begin{aligned}
V_{t+1} &\leq V_t + \frac{\eta^2}{2} \|\nabla f(x^{(t)})\|^2 - \eta(f(x^{(t)}) - f(x^*)) \\
&\leq V_t + \eta(f(x^{(t)}) - f(x^{(t+1)})) - \eta(f(x^{(t)}) - f(x^*)) \\
&= V_t - \eta(f(x^{(t+1)}) - f(x^*))
\end{aligned}$$

We obtain that:

$$f(x^{(t+1)}) - f(x^*) \leq \frac{1}{\eta}(V_t - V_{t+1})$$

Now, using the fact that for each step, the value of the function decreases, we have:

$$\begin{aligned}
T(f(x^{(T)}) - f(x^*)) &\leq \sum_{t=0}^{T-1} (f(x^{(t+1)}) - f(x^*)) \\
&\leq \frac{1}{\eta} \sum_{t=0}^{T-1} (V_t - V_{t+1})
\end{aligned}$$

Notice that the sum on the right-hand side telescopes, and we obtain:

$$T(f(x^{(T)}) - f(x^*)) \leq \frac{1}{\eta}(V_0 - V_T) \leq \frac{1}{\eta}V_0$$

Rearranging, we obtain:

$$f(x^{(T)}) - f(x^*) \leq \frac{\|x^{(0)} - x^*\|^2}{2\eta T}$$

This is the desired result. □

The result of the theorem can be generalized when used the backtracking line search. In this case, we have:

$$f(x^{(k)}) - f(x^*) \leq \frac{1}{2\eta_{\min}k} \|x^{(0)} - x^*\|^2 \quad (3.34)$$

where  $\eta_{\min} = \min\{1, \beta/L\}$ .

With this result, we can calculate the number of iterations needed to reach a certain level of accuracy. For example, if we want to reach an accuracy of  $\varepsilon$ , we have:

$$k \geq \frac{1}{2\eta\varepsilon} \|x^{(0)} - x^*\|^2 \in O\left(\frac{1}{\varepsilon}\right) \quad (3.35)$$

This means that the gradient descent algorithm converges at a rate of  $O(1/\varepsilon)$ , called **sub-linear convergence**.

### 3.7.3 Stochastic gradient descent (SGD)

The gradient descent algorithm is a simple and effective optimization algorithm that can be used to train a neural network. However, the gradient descent algorithm can be slow to converge when the training data is large. This is because the gradient descent algorithm computes the gradient of the loss function with respect to all the training data at each iteration. This can be computationally expensive, especially when the training data is large.

To overcome this problem, we can use the stochastic gradient descent algorithm. The stochastic gradient descent algorithm is a variant of the gradient descent algorithm that computes the gradient of the loss function with respect to a random sample of the training data at each iteration. This can be much faster than computing the gradient with respect to all the training data.

The stochastic gradient descent algorithm works as follows:

---

**Algorithm 8** Stochastic gradient descent algorithm

---

**Require:**  $x^{(0)}, \eta, \varepsilon, T$

```

1: for  $k = 0, 1, 2, \dots, T$  do
2:   Sample a random batch of training data  $\{x_i, y_i\}_{i=1}^m$ 
3:   Compute the gradient of the loss function with respect to the batch
4:   Update the weights of the model using the gradient
5:   if STOP( $\varepsilon$ ) then
6:     break
7:   end if
8: end for
```

---

The sample can be chosen using one of the following strategies:

1. Random sampling **with replacement**:

In this strategy, in every iteration, we sample  $m$  samples from the training data with replacement, i.e., we can sample the same sample more than once. To implement that, for each iteration, we have to choose a new set of indices to sample.

2. Random sampling **without replacement**:

In this strategy, in every iteration, we sample  $m$  samples from the training data without replacement, i.e., we can sample the same sample only once. To implement that, we just shuffle the training data once and then go through the shuffled data in batches of size  $m$ .



In practice, the second strategy is more common, as it can be more efficient. For every epoch of training, we shuffle the training data and then go through the shuffled data in batches of size  $m$ .

Also, note that SGD is originally formulated by taking just one sample at a time, i.e.,  $m = 1$ . When  $m > 1$ , we call it **mini-batch stochastic gradient descent**.

### 3.7.4 Convergence of SGD

To study the convergence of the SGD algorithm, we will make some assumptions:

1.  $|u^T H_J u| \leq L \|u\|^2$ ,  $L > 0$ , where  $H_J$  is the Hessian of the loss function.
2. The loss function  $J$  is  $\mu$ -strongly convex, i.e., for all  $x, y \in \mathbb{R}^n$ :

$$J(y) \geq J(x) + \nabla J(x)^T (y - x) + \frac{\mu}{2} \|y - x\|^2$$

3.  $\|\nabla J_i(x)\| < C$  for all  $i$  and  $x$ , where  $C$  is a constant. Note that  $J(x) = \frac{1}{m} \sum_{i=1}^m J_i(x)$  for  $m$  the number of samples.
4.  $\mathbb{E}[\nabla J_i(x)] = \nabla J(x)$ . This means that  $\nabla J_i(x)$  is an unbiased estimator of the gradient of the loss function.
5.  $0 < \eta \leq \frac{1}{2\mu}$ , where  $\eta$  is the learning rate.

From the first assumption, we obtain that:

$$J(y) \leq J(x) + \nabla J(x)^T (y - x) + \frac{L}{2} \|y - x\|^2 \quad (3.36)$$

And from the second assumption, we have:

$$J(x^*) \geq J(x) - \frac{1}{2\mu} \|\nabla J(x)\|^2 \quad (3.37)$$

Now, we would like to find an upper bound for these two differences:

- $J(w^k) - J(w^*)$
- $\|w^{(k)} - w^*\|^2$

Lets us start with the first one. By using the SGD step and the first assumption, we have:

$$\begin{aligned} J(w^{(k+1)}) &\leq J(w^{(k)}) + \nabla J(w^{(k)})^T (w^{(k+1)} - w^{(k)}) + \frac{L}{2} \|w^{(k+1)} - w^{(k)}\|^2 \\ &= J(w^{(k)}) - \eta \nabla J(w^{(k)})^T \nabla J_i(w^{(k)}) + \frac{L\eta^2}{2} \|\nabla J_i(w^{(k)})\|^2 \\ &\leq J(w^{(k)}) - \eta \nabla J(w^{(k)})^T \nabla J_i(w^{(k)}) + \frac{L\eta^2}{2} C^2 \end{aligned}$$

Now, bt subtracting  $J(w^*)$  from both sides, we have:

$$J(w^{(k+1)}) - J(w^*) \leq J(w^{(k)}) - J(w^*) - \eta \nabla J(w^{(k)})^T \nabla J_i(w^{(k)}) + \frac{L\eta^2}{2} C^2$$

Now, we can take the expectation of both sides:

$$\begin{aligned} \mathbb{E}[J(w^{(k+1)}) - J(w^*)] &\leq J(w^{(k)}) - J(w^*) - \eta \nabla J(w^{(k)})^T \mathbb{E}[\nabla J_i(w^{(k)})] + \frac{L\eta^2}{2} C^2 \\ &= J(w^{(k)}) - J(w^*) - \eta \nabla J(w^{(k)})^T \nabla J(w^{(k)}) + \frac{L\eta^2}{2} C^2 \\ &= J(w^{(k)}) - J(w^*) - \eta \|\nabla J(w^{(k)})\|^2 + \frac{L\eta^2}{2} C^2 \end{aligned}$$

Now, we can use the equation (3.37) to obtain:

$$-\|\nabla J(w^{(k)})\|^2 \leq -2\mu(J(w^{(k)}) - J(w^*))$$

Then, we have:

$$\begin{aligned} \mathbb{E}[J(w^{(k+1)}) - J(w^*)] &\leq J(w^{(k)}) - J(w^*) - 2\mu\eta(J(w^{(k)}) - J(w^*)) + \frac{L\eta^2}{2} C^2 \\ &= (1 - 2\mu\eta)(J(w^{(k)}) - J(w^*)) + \frac{L\eta^2}{2} C^2 \end{aligned}$$

Now, notice that we have the following inequality:

$$\mathbb{E}[J(w^{(k+1)}) - J(w^*)] \leq (1 - 2\mu\eta)\mathbb{E}[J(w^{(k)}) - J(w^*)] + \frac{L\eta^2}{2} C^2$$

If we reapply the expected value to the equation, the left-hand side will remain the same (as the expectation is a constant value), and we will have:

$$\mathbb{E}[J(w^{(k+1)}) - J(w^*)] \leq (1 - 2\mu\eta)\mathbb{E}[J(w^{(k)}) - J(w^*)] + \frac{L\eta^2}{2} C^2$$

This is a recursive equation. We get as solution:

$$\mathbb{E}[J(w^{(k)}) - J(w^*)] \leq (1 - 2\mu\eta)^k (J(w^{(0)}) - J(w^*)) + \frac{L\eta^2}{2} C^2 \sum_{i=0}^{k-1} (1 - 2\mu\eta)^i$$

Now, we can use the geometric series formula to obtain:

$$\sum_{i=0}^{k-1} (1 - 2\mu\eta)^i = \frac{1 - (1 - 2\mu\eta)^k}{2\mu\eta}$$

And notice that, since  $0 < \eta \leq \frac{1}{2\mu}$ , we have:

$$0 < 1 - 2\mu\eta < 1$$

So we have:

$$\frac{1 - (1 - 2\mu\eta)^k}{2\mu\eta} \leq \frac{1}{2\mu\eta}$$

Using this result, we have:

$$\mathbb{E}[J(w^{(k)}) - J(w^*)] \leq (1 - 2\mu\eta)^k (J(w^{(0)}) - J(w^*)) + \frac{L\eta}{4\mu} C^2 \quad (3.38)$$

Notice that the term  $(1 - 2\mu\eta)^k$  is a decreasing function of  $k$ . This means that the expected value of the loss function decreases with the number of iterations, with exception of the term  $\frac{L\eta}{4\mu} C^2$ . This term is a constant value that depends on the learning rate  $\eta$  and the constants  $L$ ,  $\mu$  and  $C$ . This is why we need to choose the learning rate carefully, as it can affect the convergence of the algorithm.

Similarly, we can obtain an upper bound for the expected value of the second term,  $\|w^{(k)} - w^*\|^2$ :

$$\mathbb{E}[\|w^{(k)} - w^*\|^2] \leq (1 - 2\mu\eta)^k \|w^{(0)} - w^*\|^2 + \frac{\eta}{2\mu} C^2 \quad (3.39)$$

This means that the expected value of the distance between the weights and the optimal solution also decreases with the number of iterations, with exception of the term  $\frac{\eta}{2\mu} C^2$ .

This result shows that the SGD algorithm is expected to converge to some vicinity with radius  $\frac{\eta}{2\mu} C^2$  of the optimal solution. This is a good result, as it shows that the SGD algorithm is able to find a good solution to the optimization problem.

However, notice that these results have strong assumptions, that in practice are not always satisfied. For example, the assumption that the loss function is strongly convex is not always true. In practice, the loss function is often non-convex, which can make the optimization problem more difficult.

### 3.8 Improving the SGD: Momentum and Adaptive learning rates

There are many techniques that can be used to improve the performance of the SGD algorithm. Two of the most popular techniques are **momentum** and **adaptive learning rates**. These techniques can help to speed up the convergence of the SGD algorithm and improve the performance of the model.

In this section, we will discuss the following algorithms:

- SGD with momentum
- Nesterov accelerated gradient (NAG)
- AdaGrad
- AdaDelta + RMSProp
- Adam

### 3.8.1 SGD with momentum

The SGD with momentum algorithm is a variant of the SGD algorithm that uses a momentum term to accelerate the convergence of the algorithm. The momentum term is a moving average of the gradients of the loss function. This can help to smooth out the updates to the weights and improve the performance of the model.

It's step is given by:

$$\begin{aligned}v^{(k+1)} &= \gamma v^{(k)} + \eta \nabla J(w^{(k)}) \\w^{(k+1)} &= w^{(k)} - v^{(k+1)}\end{aligned}\tag{3.40}$$

Note that  $v^{(k)}$  represents the momentum term, and  $\gamma$  is the momentum parameter. The momentum parameter is a hyperparameter that controls the influence of the momentum term on the updates to the weights.

The momentum term allows us to take into account the previous gradients when updating the weights. In some way, we are considering the velocity of the descent, and not only the gradient at the current point. This helps us to avoid local minima and saddle points, in which the gradient is close to zero.

### 3.8.2 Nesterov accelerated gradient (NAG)

The Nesterov accelerated gradient (NAG) algorithm is a variant of the SGD with momentum algorithm that uses a different update rule for the weights. The NAG algorithm uses the gradient of the loss function at the point  $w^{(k)} - \gamma v^{(k)}$  to update the weights. The step of the NAG algorithm is given by:

$$\begin{aligned}v^{(k+1)} &= \gamma v^{(k)} + \eta \nabla J(w^{(k)} - \gamma v^{(k)}) \\w^{(k+1)} &= w^{(k)} - v^{(k+1)}\end{aligned}\tag{3.41}$$

In this case, the expression  $\nabla J(w^{(k)} - \gamma v^{(k)})$  is an approximation of the gradient of the loss function at the point  $w^{(k+1)}$ . The term  $w^{(k)} - \gamma v^{(k)}$  is known as a **predictor corrector** step, as we are in some way "looking ahead" to the value of the gradient at the possible next point. Then we use this to calculate the real update to the weights. This can help to improve the performance of the algorithm and speed up the convergence.

### 3.8.3 AdaGrad

The AdaGrad algorithm is a variant of the SGD algorithm that uses an adaptive learning rate to update the weights. The AdaGrad algorithm uses a different learning rate for each weight, based on the history of the gradients of the loss function. The step of the AdaGrad algorithm is given by:

$$\begin{aligned}G^{(k)} &= G^{(k-1)} + (\nabla J(w^{(k)}))^2 \\w^{(k+1)} &= w^{(k)} - \eta_0 (G^{(k)} + \varepsilon \cdot \mathbb{1})^{-1/2} \odot \nabla J(w^{(k)})\end{aligned}\tag{3.42}$$

Here,  $G^{(k)}$  contains the sum of the squares of the gradients of the loss function up to the  $k$ -th iteration. The term  $\varepsilon$  is a small constant that is added to the denominator to avoid division by

zero. The term  $\mathbb{1}$  is a vector of ones with the same size as the weights. The term  $(G^{(k)} + \varepsilon \cdot \mathbb{1})^{-1/2}$  is the element-wise inverse square root of the sum of the squares of the gradients. This term is used to scale the gradients of the loss function before updating the weights.

The AdaGrad algorithm adapts the learning rate for each weight based on the history of the gradients of the loss function. This can help to improve the performance of the algorithm and speed up the convergence.

However, the AdaGrad algorithm has some limitations. As the sum of the squares of the gradients increases, the learning rate decreases. This can cause the learning rate to become very small, which can slow down the convergence of the algorithm. To overcome this problem, we can use the AdaDelta algorithm.

### 3.8.4 AdaDelta (and RMSProp)

This algorithm is a variant of the AdaGrad algorithm that uses an adaptive learning rate to update the weights. But in this case, instead of using the sum of the squares of the gradients up to the  $k$ -th iteration, we use a moving average of the squares of the gradients, i.e., we take the average of the squares of the gradients on the last  $n$  iterations.

Since keeping track of the last  $n$  gradients can be computationally expensive, we can use an exponential moving average to approximate the moving average. The step of the AdaDelta algorithm is given by:

$$\begin{aligned} E[g^2]^{(k)} &= \rho E[g^2]^{(k-1)} + (1 - \rho)(\nabla J(w^{(k)}))^2 \\ \Delta w^{(k+1)} &= -\eta_0 (E[g^2]^{(k)} + \varepsilon \cdot \mathbb{1})^{-1/2} \odot \nabla J(w^{(k)}) \\ w^{(k+1)} &= w^{(k)} + \Delta w^{(k+1)} \end{aligned} \tag{3.43}$$

where  $E[g^2]^{(k)}$  is the notation for the moving average of the squares of the gradients up to the  $k$ -th iteration. The term  $\rho$  is the decay rate of the moving average, that controls the influence of the previous gradients on the moving average. The term  $\varepsilon$  is a small constant that is added to the denominator to avoid division by zero.

The term  $\sqrt{E[g^2]^{(k+1)} + \varepsilon}$  is the element-wise square root of the moving average of the squares of the gradients. It is called the **root mean square** (RMS) of the gradients. This is why the AdaDelta algorithm is also known as the RMSProp algorithm, as we are using the RMS of the gradients to update the weights.

In other words, if we define:

$$RMS[g]^{(k)} = \sqrt{E[g^2]^{(k)} + \varepsilon \cdot \mathbb{1}} \tag{3.44}$$

Then, the step of the AdaDelta algorithm can be written as:

$$\begin{aligned}
E[g^2]^{(k)} &= \rho E[g^2]^{(k-1)} + (1 - \rho)(\nabla J(w^{(k)}))^2 \\
\Delta w^{(k)} &= -\frac{\eta_0}{RMS[g]^{(k)}} \odot \nabla J(w^{(k)}) \\
w^{(k+1)} &= w^{(k)} + \Delta w^{(k)}
\end{aligned} \tag{3.45}$$

The hyperparameter  $\rho$  is usually set to a value close to 1, e.g.,  $\rho = 0.9$ .

The algorithm also uses an unit corrector factor. It is based on the fact that the update term  $\Delta w^{(k)}$  is not proportional to the units of the weights. To correct this, we use the RMS of the update term  $\Delta w^{(k)}$  to scale the update term, instead of a fixed learning rate  $\eta_0$ :

$$\Delta w^{(k)} = -\frac{RMS[\Delta w]^{(k-1)}}{RMS[g]^{(k)}} \odot \nabla J(w^{(k)}) \tag{3.46}$$

So, the AdaDelta algorithm can be written as:

$$\begin{aligned}
E[g^2]^{(k)} &= \rho E[g^2]^{(k-1)} + (1 - \rho)(\nabla J(w^{(k)}))^2 \\
\Delta w^{(k)} &= -\frac{RMS[\Delta w]^{(k-1)}}{RMS[g]^{(k)}} \odot \nabla J(w^{(k)}) \\
E[\Delta w^2]^{(k)} &= \rho E[\Delta w^2]^{(k-1)} + (1 - \rho)(\Delta w^{(k)})^2 \\
w^{(k+1)} &= w^{(k)} + \Delta w^{(k)}
\end{aligned} \tag{3.47}$$

The idea of using the RMS of the gradients to update the weights is that it can help to avoid the problem of the learning rate becoming very small, as in the AdaGrad algorithm. This allows the model to adapt to the changing conditions of the optimization problem and keep learning even when the gradients are small.

### 3.8.5 Adam

Adam is a very famous optimization algorithm that combines the ideas of the SGD with momentum and the AdaDelta algorithm. The Adam algorithm uses a moving average of the gradients of the loss function and the moving average of the squares of the gradients to update the weights. The step of the Adam algorithm is given by:

$$\begin{aligned}
m^{(k)} &= \beta_1 m^{(k-1)} + (1 - \beta_1) \nabla J(w^{(k)}) \\
v^{(k)} &= \beta_2 v^{(k-1)} + (1 - \beta_2) (\nabla J(w^{(k)}))^2 \\
\hat{m}^{(k)} &= \frac{m^{(k)}}{1 - \beta_1^k} \\
\hat{v}^{(k)} &= \frac{v^{(k)}}{1 - \beta_2^k} \\
w^{(k+1)} &= w^{(k)} - \frac{\eta_0}{\sqrt{\hat{v}^{(k)} + \varepsilon}} \odot \hat{m}^{(k)}
\end{aligned} \tag{3.48}$$

Here,  $m^{(k)}$  is the moving average of the gradients of the loss function,  $v^{(k)}$  is the moving average of the squares of the gradients,  $\beta_1$  and  $\beta_2$  are the decay rates of the moving averages, and  $\varepsilon$  is a

small constant that is added to the denominator to avoid division by zero. The terms  $\hat{m}^{(k)}$  and  $\hat{v}^{(k)}$  are the bias-corrected moving averages, which are used to correct the bias of the moving averages.

The Adam algorithm uses the moving average of the gradients to update the weights, and the moving average of the squares of the gradients to scale the updates. This can help to improve the performance of the algorithm and speed up the convergence. The hyperparameters  $\beta_1$  and  $\beta_2$  control the influence of the moving averages on the updates to the weights. They are usually set to values close to 1, e.g.,  $\beta_1 = 0.9$  and  $\beta_2 = 0.999$ .

The Adam algorithm is a very powerful optimization algorithm that is widely used in practice. It can help to improve the performance of the model and speed up the convergence of the algorithm. It is recommended to use the Adam algorithm when training a neural network, as it can help to achieve better results in less time.

### 3.9 2nd order optimization methods: Newton's method and BFGS

The gradient descent algorithm (and its variants) is a first-order optimization algorithm, as it uses the first derivative of the loss function to update the weights. However, there are also second-order optimization algorithms that use the second derivative of the loss function to update the weights. These algorithms can converge faster than first-order optimization algorithms, as they take into account the curvature of the loss function.

In this section, we will discuss two second-order optimization algorithms: Newton's method and the Broyden-Fletcher-Goldfarb-Shanno (BFGS) algorithm.

#### 3.9.1 Newton's method

The Newton's method is derived from a quadratic approximation of the loss function. The idea is to approximate the loss function with a quadratic function and then find the minimum of the quadratic function. Let us consider the following quadratic approximation of the loss function  $J$  at the point  $w^{(k)}$ :

$$J(w) \approx \tilde{J}(w) = J(w^{(k)}) + \nabla J(w^{(k)})^T (w - w^{(k)}) + \frac{1}{2} (w - w^{(k)})^T H_J(w^{(k)}) (w - w^{(k)}) \quad (3.49)$$

Here,  $H_J(w^{(k)})$  is the Hessian of the loss function at the point  $w^{(k)}$ . Now, we can find the minimum of the quadratic function by setting the gradient of the quadratic function to zero:

$$\nabla \tilde{J}(w) = \nabla J(w^{(k)}) + H_J(w^{(k)}) (w - w^{(k)}) = 0$$

Solving this equation, we obtain:

$$w = w^{(k)} - H_J(w^{(k)})^{-1} \nabla J(w^{(k)}) \quad (3.50)$$

This is the update rule of the Newton's method. In some way, this method uses the inverse of the Hessian to adapt the direction of the gradient, according to the curvature of the loss function. Note that, in practice, we do not compute the inverse of the Hessian, and instead, we solve the

linear system  $H_J(w^{(k)})\Delta w = -\nabla J(w^{(k)})$  to find the update  $\Delta w$ .

We can also add a learning rate to the update rule to control the step size of the update:

$$w^{(k+1)} = w^{(k)} - \eta H_J(w^{(k)})^{-1} \nabla J(w^{(k)}) \quad (3.51)$$

Similar to what happens in the gradient descent algorithm, the learning rate  $\eta$  can be a fixed value or can be adapted during training, and should follow the Wolfe conditions.

The Newton's method is a very powerful optimization algorithm that can converge very fast. However, it has some limitations. One of the main limitations of the Newton's method is that it requires the computation of the inverse of the Hessian of the loss function, which can be computationally expensive, especially when the number of parameters is large. In practice, the inverse of the Hessian is often approximated using Quasi-Newton methods, such as the BFGS algorithm.

### 3.9.2 Secant condition

The methods for approximating the Hessian are based on the secant condition. To illustrate this, let us consider the case of minimizing a one-dimensional function,  $f : \mathbb{R} \rightarrow \mathbb{R}$ . At the  $k+1$ -th iteration, let us state the quadratic approximation of the function at the point  $x^{(k+1)}$ :

$$\tilde{f}^{(k+1)}(x) = f(x^{(k+1)}) + f'(x^{(k+1)})(x - x^{(k+1)}) + \frac{1}{2}h(x^{(k+1)})(x - x^{(k+1)})^2$$

Where  $h(x^{(k+1)})$  is the approximation of the second derivative of the function at the point  $x^{(k+1)}$ . On this approximation, we want to enforce the following condition:

$$\tilde{f}^{(k+1)'}(x^{(k+1)}) = f'(x^{(k)})$$

This will allow us to determine an expression for  $h(x^{(k+1)})$ . By solving this equation, we obtain:

$$h(x^{(k+1)}) = \frac{f'(x^{(k+1)}) - f'(x^{(k)})}{x^{(k+1)} - x^{(k)}}$$

By using this expression on the update rule of the Newton's method, we obtain:

$$x^{(k+1)} = x^{(k)} - \frac{f'(x^{(k)})}{h(x^{(k)})} = x^{(k)} - \frac{(x^{(k+1)} - x^{(k)})}{f'(x^{(k+1)}) - f'(x^{(k)})} f'(x^{(k)}) \quad (3.52)$$

This is called the secant method, for 1D functions. For multi-dimensional functions, we have the same idea, but now we have to consider the Hessian matrix. Let us again propose a quadratic approximation of the function at the point  $x^{(k+1)}$ :

$$\tilde{f}^{(k+1)}(x) = f(x^{(k+1)}) + \nabla f(x^{(k+1)})^T (x - x^{(k+1)}) + \frac{1}{2}(x - x^{(k+1)})^T \tilde{H}_f(x^{(k+1)})(x - x^{(k+1)})$$

Where  $\tilde{H}_f(x^{(k+1)})$  is the approximation of the Hessian of the function at the point  $x^{(k+1)}$ . By enforcing the following condition:

$$\nabla \tilde{f}^{(k+1)}(x^{(k+1)}) = \nabla f(x^{(k)})$$



We can determine that the approximation of the Hessian should obey the following expression:

$$\tilde{H}_f(x^{(k+1)})(x^{(k+1)} - x^{(k)}) = \nabla f(x^{(k+1)}) - \nabla f(x^{(k)}) \quad (3.53)$$

This is the secant condition for multi-dimensional functions. But notice now that in this expression, the unknown is the Hessian matrix, and not the second derivative of the function. As the Hessian is symmetric, this means that for a  $n \times n$  matrix, we have  $n(n+1)/2$  unknowns. The equation propose a total of  $n$  conditions (one for each dimension). This means that we still need to impose a total of:

$$\frac{n(n+1)}{2} - n = \frac{n(n-1)}{2}$$

additional conditions to solve the system. By formulating the rest of the conditions, we can determine the approximation of the Hessian matrix. This is the idea behind the BFGS algorithm.

Before deducing the BFGS algorithm, let us repropose the secant condition in a more general form:

$$H_{k+1}s_k = y_k \quad (3.54)$$

where:

$$s_k = x^{(k+1)} - x^{(k)}, \quad y_k = \nabla f(x^{(k+1)}) - \nabla f(x^{(k)})$$

Also notice that we are reformulating the notation of the Hessian matrix, as  $H_{k+1}$ , to indicate that on each iteration, we have a different approximation of the Hessian, that should fullfil the previous secant condition.

### 3.9.3 BFGS algorithm

This algorithm is named after Broyden, Fletcher, Goldfarb, and Shanno, who proposed the algorithm in the 1970s. The BFGS algorithm is a Quasi-Newton method that uses the secant condition to approximate the Hessian of the loss function. The BFGS algorithm enforces also the following conditions:

- $H_{k+1}$  is symmetric
- $H_{k+1}$  is positive definite
- $H_{k+1}$  is a good approximation of the previous Hessian matrix  $H_k$

In other words, the BFGS algorithm proposes the following problem to solve:

$$\begin{aligned} \min_{H_{k+1}} & \|H_{k+1} - H_k\|_F^2 \\ \text{s.t.} & H_{k+1}s_k = y_k \\ & H_{k+1} \text{ is symmetric} \end{aligned}$$

Where  $\|\cdot\|_F$  is the Frobenius norm. Now, since what we really want is to find the inverse of the Hessian, we can define the matrix  $B_{k+1} = H_{k+1}^{-1}$ , and then solve the following problem:

$$\begin{aligned} \min_{B_{k+1}} & \|B_{k+1} - B_k\|_F^2 \\ \text{s.t. } & B_{k+1}y_k = s_k \\ & B_{k+1} \text{ is symmetric} \end{aligned}$$

To solve this problem, the algorithm proposes the following update rule to the Hessian:

$$H_{k+1} = H_k + auu^T + bvv^T$$

where  $u, v$  are linear independent vectors and  $a, b$  are scalars. This is called a rank-2 update to the Hessian. By imposing the secant condition, we obtain:

$$u = y_k, \quad v = H_k s_k, \quad a = \frac{1}{y_k^T s_k}, \quad b = -\frac{1}{s_k^T H_k s_k}$$

So, the update rule to the Hessian is given by:

$$H_{k+1} = H_k + \frac{y_k y_k^T}{y_k^T s_k} - \frac{H_k s_k s_k^T H_k}{s_k^T H_k s_k} \quad (3.55)$$

But notice that we are interested in the inverse of the Hessian, not the Hessian itself. So, we can compute  $B_{k+1} = H_{k+1}^{-1}$  by using the **Sherman-Morrison formula**:

$$(A + UCV)^{-1} = A^{-1} - A^{-1}U(C^{-1} + VA^{-1}U)^{-1}VA^{-1} \quad (3.56)$$

We can use this formula by defining:

$$A = H_k, \quad U = \begin{bmatrix} H_k s_k & y_k \end{bmatrix}, \quad C = \begin{bmatrix} \frac{-1}{s_k^T H_k s_k} & 0 \\ 0 & \frac{1}{y_k^T s_k} \end{bmatrix}, \quad V = \begin{bmatrix} H_k s_k \\ y_k \end{bmatrix}$$

By using the Sherman-Morrison formula, we can obtain the update rule to the inverse of the Hessian:

$$B_{k+1} = (I - \rho s_k y_k^T) B_k (I - \rho y_k s_k^T) + \rho s_k s_k^T \quad (3.57)$$

where:

$$\rho = \frac{1}{y_k^T s_k}$$

This is the final update rule of the BFGS algorithm. But we still need to ensure that the matrix  $B_{k+1}$  is positive definite. Let us consider the generic vector  $z$ :

$$z^T B_{k+1} z = (z - \rho s_k^T z y_k)^T B_k (z - \rho s_k^T z y_k) + \rho (s_k^T z)^2$$

Now, as long as  $B_k$  is positive definite, the term  $(z - \rho s_k^T z y_k)^T B_k (z - \rho s_k^T z y_k)$  is positive. The term  $\rho (s_k^T z)^2$  is also positive, if and only if  $\rho > 0$ . This means that the matrix  $B_{k+1}$  is positive

definite, as long as we start with a positive definite matrix  $B_0$  and  $\rho > 0$ .

The condition  $\rho > 0$  (also formulated as  $y_k^T s_k > 0$ ) is called the **curvature condition**. This condition ensures that the update to the inverse of the Hessian is positive definite. For strongly convex functions, this condition is always satisfied. For other functions, we need to impose this condition in some way, for example, by finding a point  $x_{k+1}$  that satisfies Wolfe conditions, which entail that the curvature condition is satisfied, using line search.

To initialize the BFGS algorithm, we need to start with a positive definite matrix  $B_0$ . A common choice is to start with the identity matrix,  $B_0 = I$ . This is a good choice, as the identity matrix is always positive definite and obviously symmetric.

### Low memory BFGS (L-BFGS)

The BFGS algorithm requires the storage of the Hessian matrix, which can be computationally expensive when the number of parameters is large. To overcome this problem, the low memory BFGS (L-BFGS) algorithm was proposed.

Instead of storing the complete Hessian, L-BFGS stores only a few vectors that represent the approximation implicitly. Due to its resulting linear memory requirement, the L-BFGS method is particularly well suited for optimization problems with many variables.

Instead of the inverse Hessian  $B_k$ , L-BFGS maintains a history of the past  $m$  updates of the position  $x$  and gradient  $\nabla f(x)$ , where generally the history size  $m$  can be small (often  $m < 10$ ). These updates are used to implicitly do operations requiring the  $B_k$ -vector product.

### 3.9.4 Comparison between Newton's method and BFGS

The Newton's method and the BFGS algorithm are both second-order optimization algorithms that use the second derivative of the loss function to update the weights. However, the main difference between the two algorithms is the time complexity of computing the update rule.

Let us assume a problem of  $n$  dimensions. The Newton's method requires the computation of the inverse of the Hessian matrix, which has a time complexity of  $O(n^3)$ . This means that the Newton's method is computationally expensive when the number of parameters is large.

On the other hand, the BFGS algorithm uses an approximation of the Hessian matrix, which can be computed as we saw before. The time complexity of these matrix-vector products is  $O(n^2)$ . This means that the BFGS algorithm is more efficient than the Newton's method when the number of parameters is large.

In summary, the complexities are:

- Newton's method:  $O(n^3)$
- BFGS algorithm:  $O(n^2)$

But on the other hand, the Newton's method is more accurate than the BFGS algorithm, as it uses the exact Hessian matrix to update the weights. This means that the Newton's method can converge faster than the BFGS algorithm.

## Chapter 4

# Approximation properties of neural networks

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In this chapter, we will study the approximation properties of neural networks. We will show that neural networks can approximate any continuous function on a compact set to arbitrary precision. This result is known as the universal approximation theorem. We will also discuss the limitations of neural networks in approximating functions.

## 4.1 Universal approximation theorem

In this section, we will state and prove the universal approximation theorem. First, we need some concepts:

### 4.1.1 Preliminary notions

**Definition 1.** A function  $\sigma : \mathbb{R} \rightarrow \mathbb{R}$  is called a **sigmoidal function** if:

$$\lim_{x \rightarrow -\infty} \sigma(x) = 0 \quad \text{and} \quad \lim_{x \rightarrow \infty} \sigma(x) = 1$$

**Definition 2.** Let  $n$  be a natural number and  $I_n = [0, 1]^n$ . We say that an activation function  $f : \mathbb{R} \rightarrow \mathbb{R}$  is  **$n$ -discriminatory** if the only signed Borel measure  $\mu$  such that:

$$\int_{I_n} f(y \cdot x + \theta) d\mu(w) = 0 \quad \forall y \in \mathbb{R}^n, \theta \in \mathbb{R}$$

is the zero measure.

**Definition 3.** We say an activation function  $f : \mathbb{R} \rightarrow \mathbb{R}$  is **discriminatory** if it is  $n$ -discriminatory for all  $n \in \mathbb{N}$ .

**Remark 1.** A discriminatory function  $\sigma$  is **volumetrically non-destructive** when it acts on linear transformations of input.

**Definition 4.** A subspace  $U$  of  $X$  is called **dense** in  $X$  if  $\overline{U} = X$ . In other words,  $U$  is dense in  $X$  if there are elements  $u \in U$  arbitrarily close to any element  $x \in X$ . Alternatively:

1.  $\forall x \in X$ , there is a sequence  $\{u_n\} \subset U$  s.t.  $u_n \rightarrow x$ .
2.  $\forall x \in X, \forall \epsilon > 0, \exists u \in U$  s.t.  $\|x - u\| < \epsilon$ .

The fact that the subspace  $U$  is **not dense** in  $X$  can be described as:

1. There are elements  $x_0 \in X$  such that no element  $u \in U$  is arbitrarily close to  $x_0$ .
2. There is a  $\delta > 0$  such that  $\forall u \in U$  we have  $\|x_0 - u\| \geq \delta$ .

**Definition 5.** We say that a neural network is an **universal approximator** for the space  $(S, d)$  if the space of outcomes  $U$  is  $d$ -dense in  $S$ , i.e.:

$$\forall f \in S, \forall \epsilon > 0, \exists g \in U \text{ s.t. } d(f, g) < \epsilon$$

In practice it means that for any function  $f$  in the space  $S$ , there are functions in  $U$  in any proximity to  $f$ .

Let us also introduce some notations:

- Let  $K$  denote a compact set in  $\mathbb{R}^n$ , and denote  $C(K)$  the set of real-valued continuous functions on  $K$ .
- Let  $M(I_n)$  denote the set of signed regular Borel measures on  $I_n$ . By regular, we mean that, while different measures may define different sizes for a single set, they all ideally convey some idea of how much space that set takes up relative to the larger space it is contained in.

**Theorem 2** (Representation of linear bounded functionals). *Let  $K$  be a compact set in  $\mathbb{R}^n$ , and let  $F$  be a linear bounded functional on  $C(K)$ . Then there exists a unique signed regular Borel measure  $\mu$  on  $K$  such that:*

$$F(f) = \int_K f(x) d\mu(x) \quad \forall f \in C(K)$$

Moreover, the norm of  $F$  is equal to the total variation of  $\mu$ , i.e.:

$$\|F\| = |\mu|(K)$$

**Theorem 3** (Hahn-Banach). *Let  $X$  be a linear real vector space,  $X_0$  a linear subspace,  $p$  a linear convex functional on  $X$ , and  $f$  a linear functional on  $X_0$  such that:*

$$f(x) \leq p(x) \quad \forall x \in X_0$$

Then there exists a linear functional  $F$  on  $X$  such that:

$$(i) \quad F(x) = f(x) \quad \forall x \in X_0$$

$$(ii) \quad F(x) \leq p(x) \quad \forall x \in X$$

**Remark 2.** The HB theorem tells us that we can always extend a linear functional defined on a subspace to obtain a linear functional on the whole space that behaves in the same way. This is useful because it allows us to study the behavior of linear functionals on a larger space, which can provide more information about the structure of the original subspace.

From the HB theorem, we have the following 2 lemmas:

**Lemma 3.** Let  $U$  be a linear subspace of a normed linear space  $X$  and consider  $x_0 \in X$  such that:

$$d(x_0, U) \geq \delta$$

for some  $\delta > 0$ , i.e.:

$$\|x_0 - u\| \geq \delta \quad \forall u \in U$$

Then there exists a bounded linear functional  $F$  on  $X$  such that:

$$(i) \quad \|F\| \leq 1$$

$$(ii) \quad F(u) = 0 \quad \forall u \in U$$

$$(iii) \quad F(x_0) = \delta$$

**Lemma 4.** Let  $U$  be a linear, non-dense subspace of a normed linear space  $X$ . Then there is a bounded linear functional  $F$  on  $X$  such that:

$$(i) \quad F \neq 0$$

$$(ii) \quad F(u) = 0 \quad \forall u \in U$$

**Lemma 5.** Let  $U$  be a linear, non-dense subspace of  $C(I_n)$ . Then, there is a measure  $\mu \in M(I_n)$  such that:

$$\int_{I_n} h \, d\mu = 0 \quad \forall h \in U$$

*Proof.* Considering  $X = C(I_n)$  in the previous lemma, there is a bounded linear functional  $F : C(I_n) \rightarrow \mathbb{R}$  such that  $F \neq 0$  and  $F(h) = 0 \quad \forall h \in U$ . Applying the representation theorem, there is a measure  $\mu \in M(I_n)$  such that:

$$F(f) = \int_{I_n} f \, d\mu \quad \forall f \in C(I_n)$$

In particular, for any  $h \in U$ , we have  $F(h) = 0$ , which implies:

$$\int_{I_n} h \, d\mu = 0 \quad \forall h \in U$$

This completes the proof. □

**Remark 3.** Note that  $F \neq 0$  implies that  $\mu$  is not the zero measure.

#### 4.1.2 Important results

**Proposition 1.** Let  $\sigma$  be any continuous discriminatory function. Then, the finite sums of the form:

$$G(x) = \sum_{j=1}^N \alpha_j \sigma(w_j^T \cdot x + \theta_j), \quad w_j \in \mathbb{R}^n, \alpha_j, \theta_j \in \mathbb{R}$$

are dense in  $C(I_n)$ .

*Proof.* Since  $\sigma$  is continuous, it follows that:

$$U = \left\{ G : G(x) = \sum_{j=1}^N \alpha_j \sigma(w_j^T \cdot x + \theta_j) \right\}$$

is a linear subspace of  $C(I_n)$ . We will show that  $U$  is dense in  $C(I_n)$ . Assume by contradiction that  $U$  is not dense in  $C(I_n)$ . Then, by HB theorem, there is a bounded linear functional  $F$  on  $C(I_n)$  such that:

$$(i) \quad F(G) = 0 \quad \forall G \in U$$

$$(ii) \quad F \neq 0$$

By the representation theorem, there is a measure  $\mu \in M(I_n)$  such that:

$$F(f) = \int_{I_n} f d\mu \quad \forall f \in C(I_n)$$

In particular, since  $\sigma(w^T \cdot x + \theta) \in U$ , we have:

$$F(\sigma(w^T \cdot x + \theta)) = \int_{I_n} \sigma(w^T \cdot x + \theta) d\mu(x) = 0$$

However, we have assumed that  $\sigma$  is discriminatory, so this implies that  $\mu$  is the zero measure. This contradicts the fact that  $F \neq 0$ , so  $U$  must be dense in  $C(I_n)$ . □

Let us propose now some extra definitions:

**Definition 6.** *Let:*

- $P_{w,\theta} = \{x : w^T x + \theta = 0\}$  be a hyperplane in  $\mathbb{R}^n$  with normal vector  $w$  and intercept  $\theta$ .
- $H_{w,\theta}^+ = H_{w,\theta} = \{x : w^T x + \theta > 0\}$  be the positive half-space defined by the hyperplane.
- $H_{w,\theta}^- = \{x : w^T x + \theta < 0\}$  be the negative half-space.

**Lemma 6.** *Let  $\mu \in M(I_n)$ . If  $\mu$  vanishes on all hyperplanes and open half-spaces in  $\mathbb{R}^n$ , then  $\mu$  is the zero measure. More precisely, if:*

$$\mu(P_{w,\theta}) = 0 \quad \text{and} \quad \mu(H_{w,\theta}) = 0 \quad \forall w, \theta$$

*then  $\mu = 0$ .*

Now, let us prove some more propositions:

**Proposition 2.** *Any continuous sigmoidal function is discriminatory for all measures  $\mu \in M(I_n)$ .*



*Proof.* Let  $\mu \in M(I_n)$  be a fixed measure. Choose a continuous sigmoidal function  $\sigma$ , that satisfies:

$$\int_{I_n} \sigma(w^T \cdot x + \theta) d\mu(x) = 0 \quad \forall w \in \mathbb{R}^n, \theta \quad (\star)$$

We need to show that  $\mu = 0$ . First, construct the continuous function:

$$\sigma_\lambda(x) = \sigma(\lambda(w^T x + \theta) + \phi)$$

for given  $w, \theta, \phi$  and use the definition of a sigmoidal to note that:

$$\lim_{\lambda \rightarrow \infty} \sigma_\lambda(x) = \begin{cases} 1, & \text{if } w^T x + \theta > 0 \\ 0, & \text{if } w^T x + \theta < 0 \\ \sigma(\phi), & \text{if } w^T x + \theta = 0 \end{cases}$$

Define the bounded function:

$$\gamma(x) = \begin{cases} 1, & \text{if } x \in H_{w,\theta}^+ \\ 0, & \text{if } x \in H_{w,\theta}^- \\ \sigma(\phi), & \text{if } x \in P_{w,\theta} \end{cases}$$

and notice that  $\sigma_\lambda(x) \rightarrow \gamma(x)$  pointwise on  $\mathbb{R}$ , as  $\lambda \rightarrow \infty$ . The dominated convergence theorem allows switching the limit and the integral, so:

$$\begin{aligned} \lim_{\lambda \rightarrow \infty} \int_{I_n} \sigma_\lambda(x) d\mu(x) &= \int_{I_n} \gamma(x) d\mu(x) = \\ &= \int_{H_{w,\theta}^+} \gamma(x) d\mu(x) + \int_{H_{w,\theta}^-} \gamma(x) d\mu(x) + \int_{P_{w,\theta}} \gamma(x) d\mu(x) = \\ &= \mu(H_{w,\theta}^+) + \sigma(\phi)\mu(P_{w,\theta}) \end{aligned}$$

Equation  $(\star)$  implies that the limit of the left-hand side is zero, so we must have:

$$\mu(H_{w,\theta}^+) + \sigma(\phi)\mu(P_{w,\theta}) = 0$$

Since this relationship holds for any value of  $\phi$ , taking  $\phi \rightarrow \infty$  and using properties of the sigmoidal function, we get:

$$\mu(H_{w,\theta}^+) + \mu(P_{w,\theta}) = 0 \quad \forall w, \theta$$

Similarly, taking  $\phi \rightarrow -\infty$  gives:

$$\mu(H_{w,\theta}^-) + \mu(P_{w,\theta}) = 0 \quad \forall w, \theta \quad (\star\star)$$

therefore:

$$\mu(P_{w,\theta}) = 0 \quad \forall w, \theta$$

Note that, since  $H_{w,\theta}^+ = H_{-w,-\theta}^-$ , the relation  $(\star\star)$  implies that the measure  $\mu$  vanishes on all hyperplanes and open half-spaces in  $\mathbb{R}^n$ . By the previous lemma, this implies that  $\mu = 0$ . Therefore, the function  $\sigma$  is discriminatory. □

**Proposition 3.** *The ReLU function is 1-discriminatory.*

*Proof.* Let  $\mu$  be a signed Borel measure, and assume the following holds for all  $y \in \mathbb{R}$  and  $\theta \in \mathbb{R}$ :

$$\int_{I_1} \text{ReLU}(yx + \theta) d\mu(x) = 0$$

We want to show that  $\mu$  is the zero measure. For that, we will construct a sigmoid bounded, continuous (and therefore Borel measurable) function from subtracting two ReLU functions with different parameters. In particular, consider the function:

$$f(x) = \begin{cases} 0, & \text{if } x < 0 \\ x, & \text{if } x \in [0, 1] \\ 1, & \text{if } x > 1 \end{cases}$$

Then any function of the form  $g(x) = f(yx + \theta)$  with  $y \neq 0$  can be described as:

$$g(x) = \text{ReLU}(yx + \theta_1) - \text{ReLU}(yx + \theta_2)$$

by setting  $\theta_1 = -\theta/y$  and  $\theta_2 = (1 - \theta)/y$ . If  $y = 0$ , then instead set:

$$g(x) = f(\theta) = \begin{cases} \text{ReLU}(f(\theta)), & \text{if } f(\theta) \geq 0 \\ -\text{ReLU}(-f(\theta)), & \text{if } f(\theta) < 0 \end{cases}$$

which means that for any  $y \in \mathbb{R}, \theta \in \mathbb{R}$ :

$$\begin{aligned} \int f(yx + \theta) d\mu(x) &= \int (\text{ReLU}(yx + \theta_1) - \text{ReLU}(yx + \theta_2)) d\mu(x) = \\ &= \int \text{ReLU}(yx + \theta_1) d\mu(x) - \int \text{ReLU}(yx + \theta_2) d\mu(x) = \\ &= 0 - 0 = 0 \end{aligned}$$

By the previous lemma,  $f$  is discriminatory, so  $\mu$  must be the zero measure. Therefore, the ReLU function is 1-discriminatory. □

Finally, to complement the previous results, we have the following proposition:

**Definition 7.** *For  $f : \mathbb{R} \rightarrow \mathbb{R}$  an activation function, we define:*

$$\Sigma_n(f) = \text{span}\{f(y \cdot x + \theta) : y \in \mathbb{R}^n, \theta \in \mathbb{R}\}$$

**Proposition 4.** *If  $\Sigma_1(f)$  is dense in  $C(I_1)$ , then  $\Sigma_n(f)$  is dense in  $C(I_n)$ .*

### 4.1.3 Universal approximation theorem

Now, we are ready to state the universal approximation theorem:

**Theorem 4** (Universal approximation theorem). *Let  $\sigma$  be a continuous sigmoidal function. Then, for any compact set  $K \subset \mathbb{R}^n$  and any function  $f \in C(K)$ , there exist a neural network  $G$  with a single hidden layer and a finite number of neurons such that:*

$$\|f - G\|_K < \epsilon$$

*for any  $\epsilon > 0$ . This also holds for ReLU activation functions.*

*Proof.* Note that a single hidden layer neural network can be written as:

$$G(x) = \sum_{j=1}^N \alpha_j \sigma(w_j^T \cdot x + \theta_j)$$

where  $N$  is the number of neurons in the hidden layer, and  $\sigma$  is the activation function, which is either a continuous sigmoidal function or the ReLU function.

We have shown that the finite sums of the form  $G(x)$  are dense in  $C(I_n)$ , when  $\sigma$  is a discriminatory function. We also proved that any continuous sigmoidal function is discriminatory, so the result holds for sigmoidal functions.

For the ReLU function, we have shown that it is 1-discriminatory, so the linear combinations of  $n$  ReLU functions are dense in  $C(I_n)$ . Therefore, the universal approximation theorem holds for ReLU functions as well. □

### 4.1.4 Complexity of the neural network

Let us introduce some definitions:

**Definition 8.** *Let:*

- $W_m^n$  be the class of  $n$ -variable functions with partial derivatives up to order  $m$ .
- $W_m^{n,2} \subset W_m^n$  be the compositional subclass following binary tree structure.

Then, we have:

**Theorem 5.** *Let  $\sigma : \mathbb{R} \rightarrow \mathbb{R}$  be infinitely differentiable, and not polynomial. For  $f \in W_m^n$ , the complexity of shallow networks that provide an accuracy of  $\epsilon$  is:*

$$O(\epsilon^{-n/m})$$

*and it is the best possible.*

**Theorem 6.** *For  $f \in W_m^{n,2}$ , consider a deep network with the same compositional architecture and with an activation function  $\sigma : \mathbb{R} \rightarrow \mathbb{R}$  which is infinitely differentiable and not polynomial. Then, the complexity of the network that provides an accuracy of  $\epsilon$  is:*

$$O((n-1)\epsilon^{-2/m})$$

**Theorem 7.** *Let  $f$  be a  $L$ -Lipschitz continuous function of  $n$  variables. Then, the complexity of a network which is a linear combination of ReLU providing an approximation with accuracy at least  $\epsilon$  is:*

$$O\left(\left(\frac{\epsilon}{L}\right)^{-n}\right)$$

*whereas that of a deep compositional architecture is:*

$$O\left((n-1)\left(\frac{\epsilon}{L}\right)^{-2}\right)$$

## Chapter 5

# Other topics

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### 5.1 Toeplitz matrices

A Toeplitz matrix is a matrix in which each descending diagonal from left to right is constant. In other words, a matrix  $T$  is Toeplitz if  $T_{i,j} = T_{i+1,j+1}$  for all  $i, j$ . A Toeplitz matrix is completely determined by its first row and first column, i.e., for a matrix  $T$  of size  $n \times n$ , the set:

$$\{t_i \in \mathbb{R} : -(n-1) \leq i \leq (n-1)\}$$

determines the matrix  $T$  in the following way:

$$T = \begin{bmatrix} t_0 & t_{-1} & t_{-2} & \cdots & t_{-(n-1)} \\ t_1 & t_0 & t_{-1} & \cdots & t_{-(n-2)} \\ t_2 & t_1 & t_0 & \cdots & t_{-(n-3)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ t_{n-1} & t_{n-2} & t_{n-3} & \cdots & t_0 \end{bmatrix}$$

### 5.2 Circulant matrices

A circulant matrix is a special kind of Toeplitz matrix in which each row is a cyclic shift of the previous row. In other words, a matrix  $C$  is circulant if  $C_{i,j} = C_{i+1,j+1}$  for all  $i, j$ , where the indices are taken modulo  $n$  (the size of the matrix). A circulant matrix is completely determined by its first row, i.e., for a matrix  $C$  of size  $n \times n$ , the set:

$$\{c_i \in \mathbb{R} : 0 \leq i \leq n-1\}$$

determines the matrix  $C$  in the following way:

$$C = \begin{bmatrix} c_0 & c_{n-1} & c_{n-2} & \cdots & c_1 \\ c_1 & c_0 & c_{n-1} & \cdots & c_2 \\ c_2 & c_1 & c_0 & \cdots & c_3 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ c_{n-1} & c_{n-2} & c_{n-3} & \cdots & c_0 \end{bmatrix}$$

Notice that a circulant matrix can be represented using the following permutation matrix (example with  $n = 4$ ):

$$P = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

and the first row of the circulant matrix  $\vec{c} = [c_0, c_1, c_2, c_3]$  in the following way:

$$C = \begin{bmatrix} c_0 & c_1 & c_2 & c_3 \\ c_3 & c_0 & c_1 & c_2 \\ c_2 & c_3 & c_0 & c_1 \\ c_1 & c_2 & c_3 & c_0 \end{bmatrix} = (c_0I + c_1P + c_2P^2 + c_3P^3)$$

Also, note that  $P \in \mathbb{R}^{n \times n}$  (the permutation matrix in  $\mathbb{R}^{n \times n}$ ) is a circulant matrix, and has the following property:

$$P^m = P^{(m \bmod n)}$$

with  $P^0 = I$ . This is a very useful property when working with circulant matrices.

For example, assume that  $C, D \in \mathbb{R}^{4 \times 4}$  are circulant matrices with first rows  $\vec{c} = [c_0, c_1, c_2, c_3]$  and  $\vec{d} = [d_0, d_1, d_2, d_3]$ , respectively. Then, the product  $CD$  is also a circulant matrix with first row  $\vec{e} = [e_0, e_1, e_2, e_3]$ .

In fact, we can show that the multiplication of these matrices is a multiplication of polynomials in the following way:

$$CD = (c_0I + c_1P + c_2P^2 + c_3P^3)(d_0I + d_1P + d_2P^2 + d_3P^3)$$

and by the cycling property of  $P$ , we have that the result will be also a polynomial in  $P$  of degree at  $n - 1 = 3$ :

$$CD = e_0I + e_1P + e_2P^2 + e_3P^3$$

where the coefficients  $e_i$  are given by:

$$\vec{e} = \vec{c} \circledast \vec{d}$$

where  $\circledast$  denotes the cyclic convolution of the vectors  $\vec{c}$  and  $\vec{d}$ , defined as:

$$e_i = \sum_{j=0}^{n-1} c_j d_{i-j}$$

### 5.3 Eigenvalues and eigenvectors of circulant matrices

The eigenvalues and eigenvectors of circulant matrices have a very nice structure. In particular, the eigenvalues of a circulant matrix  $C$  are given by:

$$\lambda_k = c_0 + c_1\omega^k + c_2\omega^{2k} + \cdots + c_{n-1}\omega^{(n-1)k}$$

where  $\omega = e^{2\pi i/n}$  is a primitive  $n$ -th root of unity. The eigenvectors of  $C$  are given by:

$$v_k = \frac{1}{\sqrt{n}}[1, \omega^k, \omega^{2k}, \dots, \omega^{(n-1)k}]^T$$

where  $k = 0, 1, \dots, n-1$ . Notice that the eigenvectors of a circulant matrix are the columns of the  $n \times n$  DFT matrix (Fourier matrix), which we will denote as  $F_n$ . The DFT matrix is defined as:

$$F_n = \frac{1}{\sqrt{n}} \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \omega & \omega^2 & \dots & \omega^{n-1} \\ 1 & \omega^2 & \omega^4 & \dots & \omega^{2(n-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{n-1} & \omega^{2(n-1)} & \dots & \omega^{(n-1)(n-1)} \end{bmatrix}$$

For example, for  $n = 4$ , and the circulant matrix  $C$  with first row  $\vec{c} = [c_0, c_1, c_2, c_3]$ , the eigenvalues and eigenvectors are given by:

$$\begin{aligned} \lambda_0 &= c_0 + c_1 + c_2 + c_3 \\ \lambda_1 &= c_0 + c_1 i + c_2 i^2 + c_3 i^3 \\ \lambda_2 &= c_0 + c_1 i^2 + c_2 i^4 + c_3 i^6 \\ \lambda_3 &= c_0 + c_1 i^3 + c_2 i^6 + c_3 i^9 \end{aligned}$$

as  $\omega = e^{2\pi i/4} = i$ . The eigenvectors are:

$$\begin{aligned} v_0 &= \frac{1}{2}[1, 1, 1, 1]^T \\ v_1 &= \frac{1}{2}[1, i, -1, -i]^T \\ v_2 &= \frac{1}{2}[1, -1, 1, -1]^T \\ v_3 &= \frac{1}{2}[1, -i, -1, i]^T \end{aligned}$$

which are the columns of the DFT matrix  $F_4$ :

$$F_4 = \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & i & i^2 & i^3 \\ 1 & i^2 & i^4 & i^6 \\ 1 & i^3 & i^6 & i^9 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{bmatrix}$$

Note that we can calculate the eigenvalues and eigenvectors of a circulant matrix using the DFT matrix  $F_n$  in the following way:

$$\begin{aligned} \vec{\lambda}(C) &= F_n \vec{c} \\ V(C) &= F_n \end{aligned}$$

where  $\vec{\lambda}(C)$  is the vector of eigenvalues of  $C$ , and  $V(C)$  is the matrix of eigenvectors of  $C$ .

## 5.4 Convolution rule

Note that, from the properties of before, we can write the product of two circulant matrices as a new circulant matrix, with the coefficients given by the cyclic convolution of the coefficients of the original matrices. In other words, if  $C, D \in \mathbb{R}^{n \times n}$  are circulant matrices with first rows  $\vec{c} = [c_0, c_1, \dots, c_{n-1}]$  and  $\vec{d} = [d_0, d_1, \dots, d_{n-1}]$ , respectively, then the product  $CD$  is also a circulant matrix with first row  $\vec{e} = [e_0, e_1, \dots, e_{n-1}]$ , where

$$\vec{e} = \vec{c} \circledast \vec{d}$$

Now, the eigenvalues of the product  $CD$  are given by:

$$\vec{\lambda}(CD) = F_n \vec{e} = F_n (\vec{c} \circledast \vec{d})$$

Now, notice that:

$$CDx = C(Dx) = C(\lambda(D)x) = \lambda(D)(Cx) = \lambda(D)(\lambda(C)x) = \lambda(C)\lambda(D)x$$

where  $\lambda(C)$  and  $\lambda(D)$  are the eigenvalues of  $C$  and  $D$ , respectively. Therefore, the eigenvalues of the product  $CD$  are given by the element-wise product of the eigenvalues of  $C$  and  $D$ :

$$\vec{\lambda}(CD) = \vec{\lambda}(C) \odot \vec{\lambda}(D)$$

where  $\odot$  denotes the element-wise product of two vectors. But we now that:

$$\vec{\lambda}(C) = F_n \vec{c}, \quad \vec{\lambda}(D) = F_n \vec{d}$$

So, combining these results, we have that:

$$\vec{\lambda}(CD) = F_n \vec{c} \odot F_n \vec{d} = F_n (\vec{c} \circledast \vec{d})$$

Formally, we denote this relation as the convolution rule for circulant matrices (also known as the circular convolution theorem):

**Theorem 8.** For  $\vec{c} \in \mathbb{R}^n$  and  $\vec{d} \in \mathbb{R}^n$ , we have that:

$$F_n (\vec{c} \circledast \vec{d}) = F_n \vec{c} \odot F_n \vec{d}$$

where  $\circledast$  denotes the cyclic convolution of the vectors  $\vec{c}$  and  $\vec{d}$ , and  $\odot$  denotes the element-wise product of two vectors.

Notice that this is a very useful property. Indeed, notice that performing a circular convolution between 2 vectors takes  $O(n^2)$  operations, and then computing the DFT of the result takes  $O(n \log n)$  operations, resulting in a total complexity of  $O(n \log n + n^2) = O(n^2)$  operations. On the other hand, computing the DFT of the two vectors takes  $O(n \log n)$  operations, and then performing the element-wise product takes  $O(n)$  operations, resulting in a total complexity of  $O(n \log n + n) = O(n \log n)$  operations.

Therefore, the convolution rule allows us to compute the left-hand side of the equation in less operations.