

CS240, Spring 2022

Assignment 5: Question 1

Q1) Give an algorithm to find the point with the maximum x-value in a 2D kd-tree, and analyze its complexity. For maximum credit, your algorithm must run in $o(n)$ time, where n is the number of points in the kd-tree. For simplicity, you may assume that n is a power of 4. If the run-time $T(n)$ of your algorithm satisfies a recurrence relation that has already been seen in class, you can take for granted the corresponding growth rate for $T(n)$, without giving the proof.

Assume that we already have a KD tree called K which for each split (both in terms of x and y) somewhat equally separates values into two parts.

When our algorithm encounters a split in terms of x , we will always recurse on the right child of K . This is because we are looking for the largest x -value, which should always be larger than our current split. When we reach a split in terms of y we should recurse on both sides as the largest x value can have any y value.

We will assume that we build the KD tree by doing an x split first, and we will use a boolean to represent the type of split. Our algorithm is thus:

Algorithm 1 KD Tree Max X Search

```
findMax(kdTree K, Bool xSplit) {  
  if K is leaf then  
    return k.x  
  end if  
  if xSplit == true then  
    return findMax(k.rightChild, false)  
  end if  
  // if we reach this stage, it means we should do a Y split  
  return max(findMax(k.rightChild, true), findMax(k.leftChild, true))  
}
```

Note that half the values will go under the x split, and the other half will go under the y split and will additionally be split twice, thus our relation is (simplifying using known relations):

$$T(n) = \frac{1}{2}T(n/2) + \frac{1}{2} \times 2T(n/4) + O(1) \quad (1)$$

$$= \frac{1}{2}\sqrt{n} + \frac{1}{2} \times 2T(n/4) + O(1) \quad (2)$$

$$= \frac{1}{2}\sqrt{n} + \frac{1}{2} \log(n) \quad (3)$$

Since we know \sqrt{n} dominates $\log(n)$ we will do our limit test using $\frac{\sqrt{n}}{n}$:

$$\text{limit} = \lim_{n \rightarrow \infty} \frac{\sqrt{n}}{n} \tag{4}$$

$$= \lim_{n \rightarrow \infty} \frac{1}{\sqrt{n}} \tag{5}$$

$$= 0 \tag{6}$$

And so by the limit test we have:

$$\frac{1}{2}\sqrt{n} + \frac{1}{2}\log(n) = o(n)$$