

# An Open-Source Tool for Generating Domain-Specific Accelerators for Resource-Constrained Computing

David T Kerns<sup>1</sup> and Tosiron Adegbija<sup>1</sup>

<sup>1</sup> Department of Electrical & Computer Engineering, University of Arizona, Tucson, Arizona, USA

DOI: [10.xxxxxx/draft](https://doi.org/10.xxxxxx/draft)

## Software

- [Review](#) ↗
- [Repository](#) ↗
- [Archive](#) ↗

Editor: ↗

Submitted: 01 December 2023

Published: unpublished

## License

Authors of papers retain copyright and release the work under a Creative Commons Attribution 4.0 International License ([CC BY 4.0](#)).

## Summary

Domain-specific accelerators (DSAs) (Hennessy & Patterson, 2019) are crucial in modern computer architecture as they enable highly efficient processing for specialized tasks, significantly improving performance and energy efficiency over general-purpose computing systems. Previous work (Limaye & Adegbija, 2021) introduced the *Superblock* (SB) as a new granularity for designing DSAs. The SB provides a middle ground to the existing extreme granularities of creating DSAs only from Basic Blocks (BBs) or whole functions. This paper describes an Open Source Software (OSS), called **D2**, that implements the SB approach and will enable researchers and developers in the DSA community to easily experiment with this new granularity towards finding an optimal solution to the DSA definition at hand. Our results show that the SB should be taken seriously as a candidate for creating DSAs.

## Statement of need

The DSA field currently offers few, if any, open-source automation tools. An important challenge in the design of DSAs is identifying what portions of a set of domain programs should be implemented in hardware (i.e., accelerated) to maximize the performance and energy benefits of the DSA. (Hennessy & Patterson, 2019) even argue for a new domain-specific language. The SB as a new DSA granularity offers real benefits for efficiently addressing this challenge. However, manually identifying these SBs can be prohibitively time-consuming and error-prone. This ground-up re-write as open source from conception hopes to bring the SB construct to the community at large. The **D2** tool provides an exhaustive list of SBs that can be simulated to determine the optimal size.

## Key aspects of D2

This is a follow-on work of (Limaye & Adegbija, 2021), re-engineered from the ground up, with special emphasis on making it open source. Additionally, the novel parts are:

- user-controlled constraints
- normalization of BBs
- ranking at the BB level and then mapping the BB ranking onto SBs
- maintaining a link back to the source so that the accelerators can be generated directly from the C source code rather than the LLVM intermediate representation (IR) files

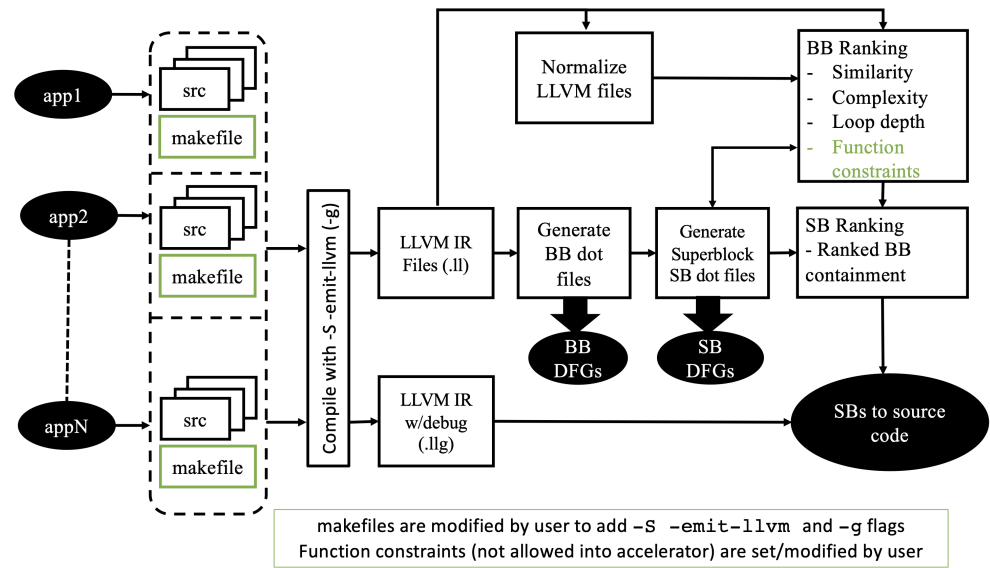


Figure 1: D2 Flow

## D2 Overview

**D2**, as depicted in Figure 1, is a tool that accepts a system's tree of source code. The user makes minor modifications to the makefiles to generate the required LLVM IR files. **D2** then evaluates the IR files to create a set of metadata files. It then identifies and ranks BBs and SBs as candidates for FPGA acceleration to produce an optimal set of accelerators for the system. The identified source lines are then given to a synthesis tool, such as Xilinx Vivato HLS (Xilinx, 2021) to produce an FPGA accelerator that is integrated back into the source code. The resulting system runs faster and more efficiently, potentially making the system even viable on a resource-constrained target.

## Target system model

The target is assumed to be a resource-constrained computing system in which the generated DSA can be integrated. For example, this could be a CPU+FPGA system, wherein the DSA is implemented on the FPGA, while the unaccelerated portions of the target workloads are run on the CPU. In this scenario, the FPGA should be addressable from the CPU such that the identified code that would normally be run by the CPU can be implemented on the FPGA and offloaded from the CPU without a significant bottleneck. This model is commonly found in resource-constrained computing devices where power constraints are a major concern, like human-implanted medical devices (Karageorgos et al., 2020), for example.

## D2 implementation details

**D2** is a shell script that calls the individual programs written in both Python and C++ that process the LLVM files generated by the compiler.

### The search for superblocks

The control flow graph (CFG) is easily generated from the LLVM IR (LLVM Language Reference Manual, n.d.-a) files. We use Graphviz (Gansner et al., 1993) .dot file format to represent the CFG. The heart of the overall **D2** package is the sb application (sb.cc) that reads the .dot file

59 and iteratively identifies every SB in the CFG file. Using the C++ Standard Template Library  
60 (STL) (Plauger et al., 2000), a map is constructed of each BB as a node with a vector of  
61 connected nodes. Once the graph is parsed into the map/tree, it recursively walks the tree  
62 identifying SBs by adding one adjacent node at a time and checking if the “One In, One Out”  
63 condition of an SB is satisfied. Once identified, the SB is added to a vector of SBs that is  
64 output and then ranked in a later stage.

#### 65 Identification of accelerator candidates

66 The key to achieving maximum acceleration is to choose areas of the software that are  
67 compute-bound. Further, if an accelerator can be used by multiple applications, the number of  
68 accelerators can be reduced. **D2** makes a concerted effort to identify common and frequently  
69 used code segments for acceleration.

#### 70 Constraints

71 One piece of metadata that is captured in the CFG file is all functions that the BB calls. These  
72 are labeled as constraints; in that if the BB is to be realized by an accelerator, any function  
73 called by the BB must also be incorporated by the accelerator. Thus, a BB that calls I/O  
74 functions, for example, is not a candidate for acceleration. During the **D2** processing of the  
75 source tree, the user is presented the list of constraints or functions called, if any, and given  
76 the option to remove any functions that could be realized by an accelerator.

#### 77 Normalization of basic blocks

78 **D2** finds common accelerators through a process we call *normalization*. The normalization  
79 process strips the code of all data, both variables and constants, and replaces them with named  
80 registers. This process is greatly simplified via LLVM's existing Static Single Assignment (SSA)  
81 (LLVM Language Reference Manual, n.d.-b) strategy. Once normalized, BBs that perform  
82 the same set of instructions on a given set of data inputs are identified and consolidated to  
83 minimize redundancy.

#### 84 Ranking and selection

85 Ranking is accomplished by keeping metrics on each BB. Often, several SBs are subsets  
86 of a larger SB. While the largest SB often offers the best candidate for acceleration for a  
87 specific workload, a smaller SB that is common to multiple workloads may offer better overall  
88 acceleration to the entire system. **D2** makes it easier to iterate through the many possibilities  
89 to find the optimal solution for the target workloads.

#### 90 Back to the source

91 **D2** tracks the original C source files and lines of code that comprise each basic block. This  
92 allows us to use generic high-level synthesis tools like Xilinx Vivado (Xilinx, 2021) to produce  
93 the hardware description code (e.g., Verilog, SystemVerilog, VHDL) from the C source. This  
94 also simplifies modification of the original source.

#### 95 Conclusion

96 We believe there is a future in right-sizing the DSA and that the **D2** tool can provide valuable  
97 input to that end. Because there are very few OSS tools geared towards the automation of  
98 DSA identification, we hope that the **D2** tool will be utilized and expanded upon within the  
99 computer architecture community to become a valuable resource and additionally, make the  
100 concept of the SB more accessible to the community as a whole.

## References

- Gansner, E. R., Koutsofios, E., North, S. C., & Vo, K.-P. (1993). A technique for drawing directed graphs. *IEEE Trans. Software Eng.*, 19(3), 214–230. <https://doi.org/10.1109/32.221135>
- Hennessy, J., & Patterson, D. (2019). A new golden age for computer architecture. In *Communications of the ACM* (Vol. 62, pp. 48–60). ACM. <https://doi.org/10.1145/3282307>
- Karageorgos, I., Sriram, K., Veselý, J., Wu, M., Powell, M., Borton, D., Manohar, R., & Bhattacharjee, A. (2020). Hardware-software co-design for brain-computer interfaces. *2020 ACM/IEEE 47th Annual International Symposium on Computer Architecture (ISCA)*, 391–404. <https://doi.org/10.1109/ISCA45697.2020.00041>
- Limaye, A., & Adegbija, T. (2021). DOSAGE: Generating domain-specific accelerators for resource-constrained computing. *2021 IEEE/ACM International Symposium on Low Power Electronics and Design (ISLPED)*, 1–6. <https://doi.org/10.1109/ISLPED52811.2021.9502501>
- LLVM language reference manual*. (n.d.-b). <https://llvm.org/docs/LangRef.html#abstract>
- LLVM language reference manual*. (n.d.-a). <https://llvm.org/docs/LangRef.html#introduction>
- Plauger, P. J., Lee, M., Musser, D., & Stepanov, A. A. (2000). *C++ standard template library* (1st ed.). Prentice Hall PTR. ISBN: 0134376331
- Xilinx, I. (2021). *Introduction to FPGA design with vivado high-level synthesis*. Xilinx. <https://docs.xilinx.com/v/u/en-US/ug998-vivado-intro-fpga-design-hls>