# **Anonymity**

- So far: just talked about encryption
- · We leak:
  - o Our IP
  - Timing of our messages
  - Sizes of our message
- Idea: use a relay / VPN
  - A big fan in and fan out
  - o Idea is to make it impossible to figure out which input corresponds to which output
  - Need to create a TLS connection between server and client that is encrypted to ensure relay
- Problems:
  - A compromised relay gets access to everything
  - Performance
  - Timing analysis still works

## **Onion Routing**

- We use a path of relays
- Each relay only knows what sent it to them and where it is going to send it
- Cannot just make a single encrypted connection between each client / relay and relay / relay because then the relay is seeing unencrypted data
- Let's say we have a path of three relays:
  - o C-R1-R2-R3-S
  - C R1 sets up a TLS connection
  - Using this TLS connection, we set up another TLS connection between C and R2 that goes through R1
    - Nested within the C R1 connection
    - We continue nesting connections within each other
- TOR typically uses a circuit of three relays

## Directory

- We want all clients to use the same directory / database of relays
  - An attack could be to give different clients different sets of relays and use that as an attack
- To mitigate this, TOR uses a bunch of directory servers that each sign the directory

#### **TOR Browser**

- Tor browser does more than just use the network
- Trying to limit the amount of unique information leaked
  - Chrome exopses a lot of information like screen size / your graphics card which can build a unique fingerprint of your machine
  - o TOR browser tries to do stuff to make your browser look like everyone else's

## **Bridges**

- Bridge = hidden relays
- Used to combat firewalls that are used for censorship
  - You ask a BridgeDB for a bridge relay
  - o Rate limited
  - Competent firewall operators can try to block this
    - But if you host the BridgeDB on some CDN like Cloudflare then the operator might have to block all of Cloudflare

 Current version of TCP still leaks the website that you are connecting to so good firewall operators (i.e. China) can still block it

### **Anonymous Servers**

- So far we have talked about clients being anonymous
- Instead servers want to be anonymous
  - I.e. whistleblowing or illegal websites
  - Also DOS resistant
- Servers choose some relay point to be a "introduction point" / "mailbox"
  - Builds a circuit to that intro point
- Client chooses some relay as the rendezvous point
  - They then build another circuit to the introduction point and sent the name of the rendezvouz point to the server through it
    - It can also send information like a username to try to convince the server to accept the connection
  - The server, if it chooses to accept the connection, will connect to the Rendezvout point
- · Clever idea:
  - The name of the onion link will just be the public key
  - This is important because:
    - We don't need a common directory to store this because the server can just sign messages with its private key and the client will see that this is legit
    - It also allows you to encrypt messages to the server with the public key because no TLS CA is going to issue a certificate to the server

### **Guard Nodes**

- If we have a path of three nodes but the ends are both compromised (R1 and R3)
  - Can do sophisticated timing attacks by comparing the times
  - Can bypass an honest node
  - As more and more circuits are built by you, the probability increases to 1 that there will be a compromised circuit
- Guard note idea:
  - Choose and keep first node
  - o Only change the second and third node each time
    - You have a chance of getting screwed the first time but your probability won't approach 1 over your lifetime