

Assignment 1 - CS 4071 - Spring 2018

Due: 2018-01-26

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1. Exercise 1.1

Problem: Trace the action of the left-to-right binary method to compute:

a. x^{123}

$$\begin{aligned} 123 &= 1111011_2 \\ x &\rightarrow x \times (x)^2 = x^3 \rightarrow x \times (x^3)^2 = x^7 \rightarrow x \times (x^7)^2 \\ &= x^{15} \rightarrow (x^{15})^2 = x^{30} \rightarrow x \times (x^{30})^2 = x^{61} \rightarrow x \times (x^{61})^2 = x^{123} \end{aligned}$$

b. x^{64}

$$\begin{aligned} 64 &= 1000000_2 \\ x &\rightarrow (x)^2 = x^2 \rightarrow (x^2)^2 = x^4 \rightarrow (x^4)^2 \\ &= x^8 \rightarrow (x^8)^2 = x^{16} \rightarrow (x^{16})^2 = x^{32} \rightarrow (x^{32})^2 = x^{64} \end{aligned}$$

c. x^{65}

$$\begin{aligned} 65 &= 1000001_2 \\ x &\rightarrow (x)^2 = x^2 \rightarrow (x^2)^2 = x^4 \rightarrow (x^4)^2 \\ &= x^8 \rightarrow (x^8)^2 = x^{16} \rightarrow (x^{16})^2 = x^{32} \rightarrow x \times (x^{32})^2 = x^{65} \end{aligned}$$

d. x^{711}

$$\begin{aligned} 711 &= 1011000111_2 \\ x &\rightarrow (x)^2 = x^2 \rightarrow x \times (x^2)^2 = x^5 \rightarrow x \times (x^5)^2 = x^{11} \rightarrow (x^{11})^2 = x^{22} \rightarrow (x^{22})^2 \\ &= x^{44} \rightarrow (x^{44})^2 = x^{88} \rightarrow x \times (x^{88})^2 = x^{177} \rightarrow x \times (x^{177})^2 = x^{355} \rightarrow x \times (x^{355})^2 = x^{711} \end{aligned}$$

2-old. Exercise 1.7

Problem: Trace the action of the algorithm *GCD* for the following input pairs:

a. (24, 108)

$$\begin{aligned} & \text{NaiveGCD}(24, 108) \\ &= \text{gcd}(24, 84) = \text{gcd}(24, 60) = \text{gcd}(24, 36) \\ &= \text{gcd}(24, 12) = \text{gcd}(12, 12) = 12 \end{aligned}$$

b. (23, 108)

$$\begin{aligned} & \text{NaiveGCD}(23, 108) \\ &= \text{gcd}(23, 85) = \text{gcd}(23, 62) = \text{gcd}(23, 39) \\ &= \text{gcd}(23, 16) = \text{gcd}(7, 16) = \text{gcd}(7, 9) = \text{gcd}(7, 2) \\ &= \text{gcd}(5, 2) = \text{gcd}(3, 2) = \text{gcd}(1, 2) = \text{gcd}(1, 1) = 1 \end{aligned}$$

c. (89, 144)

$$\begin{aligned} & \text{NaiveGCD}(89, 144) \\ &= \text{gcd}(89, 55) = \text{gcd}(34, 55) = \text{gcd}(34, 21) \\ &= \text{gcd}(13, 21) = \text{gcd}(13, 8) = \text{gcd}(5, 8) = \text{gcd}(5, 3) \\ &= \text{gcd}(2, 3) = \text{gcd}(2, 1) = \text{gcd}(1, 1) = 1 \end{aligned}$$

d. (1953, 1937)

$$\begin{aligned} & \text{NaiveGCD}(1953, 1937) \\ &= \text{gcd}(16, 1937) = \text{gcd}(16, 1921) = \text{gcd}(16, 1905) = \text{gcd}(16, 1889) \\ &= \dots = \text{gcd}(16, 49) = \text{gcd}(16, 33) = \text{gcd}(16, 17) = \text{gcd}(16, 1) \\ &= \text{gcd}(15, 1) = \text{gcd}(14, 1) = \dots = \text{gcd}(3, 1) = \text{gcd}(2, 1) = \text{gcd}(1, 1) = 1 \end{aligned}$$

2-new. Exercise 1.8

Problem: Trace the action of the algorithm *EuclidGCD* for the following input pairs:

a. (24, 108)

$$\begin{aligned} & \text{EuclidGCD}(24, 108) \\ &= \text{gcd}(108, 24) = \text{gcd}(24, 12) = \text{gcd}(12, 0) = 12 \end{aligned}$$

b. (23, 108)

$$\begin{aligned}
& \text{EuclidGCD}(23, 108) \\
&= \gcd(108, 23) = \gcd(23, 16) = \gcd(16, 7) \\
&= \gcd(7, 2) = \gcd(2, 1) = \gcd(1, 0) = 1
\end{aligned}$$

c. (89, 144)

$$\begin{aligned}
& \text{EuclidGCD}(89, 144) \\
&= \gcd(144, 89) = \gcd(89, 55) = \gcd(55, 34) = \gcd(34, 21) \\
&= \gcd(21, 13) = \gcd(13, 8) = \gcd(8, 5) = \gcd(5, 3) \\
&= \gcd(3, 2) = \gcd(2, 1) = \gcd(1, 0) = 1
\end{aligned}$$

d. (1953, 1937)

$$\begin{aligned}
& \text{EuclidGCD}(1953, 1937) \\
&= \gcd(1937, 16) = \gcd(16, 1) = \gcd(1, 0) = 1
\end{aligned}$$

3. Exercise 1.9

Problem: Give a formula for $\text{lcm}(a, b)$ in terms of $\gcd(a, b)$.

$$\text{lcm}(a, b) = a \times b \div \gcd(a, b)$$

4. Exercise 1.17

Problem: Trace the action of Horner's rule for the polynomial $7x^5 - 3x^3 + 2x^2 + x - 5$.

Horner's rule for 5th degree polynomial:

$$((((a_5 \times x + a_4) \times x + a_3) \times x + a_2) \times x + a_1) \times x + a_0$$

For the polynomial $7x^5 - 3x^3 + 2x^2 + x - 5$ the coefficients are:

$$a_5 = 7, \quad a_4 = 0, \quad a_3 = -3, \quad a_2 = 2, \quad a_1 = 1, \quad a_0 = -5$$

So Horner's rule is:

$$((((7 \times x + 0) \times x + -3) \times x + 2) \times x + 1) \times x + -5$$

For $x = 7$:

$$\begin{aligned}
& (((((7 \times 7 + 0) \times 7 + -3) \times 7 + 2) \times 7 + 1) \times 7 + -5 \\
& (((((49) \times 7 + -3) \times 7 + 2) \times 7 + 1) \times 7 + -5 \\
& (((((340) \times 7 + 2) \times 7 + 1) \times 7 + -5 \\
& (((((2,382) \times 7 + 1) \times 7 + -5 \\
& (((((16,675) \times 7 + -5 \\
& 116,720
\end{aligned}$$

5. Exercise 2.13

Problem: Design an algorithm that tests whether or not two input lists of size n have at least one element in common. Give formulas for $B(n)$ and $W(n)$ for your algorithm.

function CheckShareElements($a[0 : n - 1], b[0 : n - 1]$)

Input: $a[0 : n - 1], b[0 : n - 1]$ (two arrays of real numbers)

Output: `true` if a and b have at least 1 element in common, else `false`

```

1  for (i = 0; i < n; i++)
2      for (j = 0; j < n; j++)
3          if (i == j) return true
4  return false

```

If $a[0] = b[0]$, then it'll return true after the very first comparison. Thus, $B(n) = 1$.

If $a \cap b = \emptyset$ (or if the only common element is the last element of a and the last element of b), then it'll have to compare every element of a to every element of b . Thus, $W(n) = n^2$.

6. Exercise 2.16

Part a.

Problem: Write a procedure that finds both the largest and second-largest elements in a list $L[0 : n - 1]$ of size n .

function FindTwoLargest($L[0 : n - 1]$)

Input: $L[0 : n - 1]$ (an array of real numbers)

Output: a and b (the largest two numbers from L)

```

1  # start by assuming the first two elements are the largest
2  a = L[0]
3  b = L[1]
4  if (b > a) # make sure a is bigger than b
5      temp = a

```

```

6      a = b
7      b = temp
8
9      # search the list for bigger numbers, and update a and b as needed
10     for (i = 2; i < n; i++)
11         if (L[i] > a)
12             b = a
13             a = L[i]
14         elif (L[i] > b)
15             b = L[i]
16
17     return a, b

```

Part b.

Problem: Determine $B(n)$ and $W(n)$ for the algorithm in a.

Our basic operation is comparisons. We do one comparison at the beginning, before the loop. Then on each iteration of the loop, we do one comparison if $L[i] > a$; if not, we end up doing two comparisons. Thus, our best case would be if the list sorted from smallest to largest (meaning $L[i] > a$ will be true on every iteration). Our worst case would any list where none of the other elements are larger than $L[0]$ or $L[1]$ (meaning $L[i] > a$ will be false on every iteration). Thus:

$$B(n) = 1 + (n - 2) = n - 1$$

$$W(n) = 1 + 2(n - 2) = 2n - 3$$

7. Exercise 2.30

Part a.

Problem: Design a recursive version of *InsertionSort*.

Part b.

Problem: Design a recursive linked list version of *InsertionSort*.

8. Exercise 1.19

Problem: Describe a modification of *HornerEval* that solves this particular evaluation problem using only $n + 1$ multiplications and $n + 1$ additions.

function **ModifiedHornerEval**($a[0 : n], v$)

Input: $a[0 : n]$ (an array of real numbers), v (a real number)

Output: the values of polynomial $P(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$ at $x = v$ and $x = -v$

```
1  v2 = v*v
2  if n is even:
3      EvenSum <= a[n]
4      OddSum <= 0
5  else:
6      EvenSum <= 0
7      OddSum <= a[n]
8  endif
9  for i <= n-1 downto 0 do:
10     if i is even:
11         EvenSum <= EvenSum * v2 + a[i]
12     else:
13         OddSum <= OddSum * v2 + a[i]
14     endif
15 endfor
16 OddSum = OddSum*v
17 return(P(v) is EvenSum+OddSum and P(-v) is EvenSum-OddSum)
```