

lecture 21, the formal language GPLI

phil1012 introductory logic

overview

this lecture

- an introduction to the formal language GPLI
- some of the limitations of GPL which motivate the shift to GPLI
- the identity relation
- translations into GPLI
- the semantics of identity

learning outcomes

- after doing the relevant reading for this lecture, listening to the lecture, and attending the relevant tutorial, you will be able to:
 - explain the limitations of GPL and the motivations for GPLI
 - explain what the identity relation is
 - translate propositions from English into GPLI
 - explain the motivation behind the semantics of identity
 - determine whether a GPLI proposition is true or false on a model

required reading

- sections 13.1, 13.2, and 13.3 of chapter 13

expressive limitations of GPL

expressive limitations of GPL

- consider the following argument

P1.	John loves something
P2.	John does not love himself
<hr/>	
C1.	John loves something other than himself

- we have the following glossary:
 - j : John
 - Lxy : x loves y
- and we translate the premises as follows:
 - $\exists x Ljx$
 - $\neg Ljj$
- but we have no way of translating the conclusion
- we have no way of saying that John loves someone other than himself

introducing GPLI

introducing GPLI

- we introduce a new two-place predicate into our logical language: the **identity** predicate
 - I^2
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- in GPL I^2 is not used for a two-place predicate. Although I^1 , I^3 , I^4 , and so on are
 - we are extending GPL to GPLI which includes the symbol I^2 in addition to all the symbols of GPL
 - syntactically, I^2 functions just like a two-place predicate in GPL
 - semantically, I^2 is treated differently
 - its meaning does not vary from model to model
 - it is not given a glossary
 - in every model it expresses the **identity relation**
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- since the two-place predicate I^2 is so special in GPLI we usually abbreviate it as follows:
 - $=$
- and we abbreviate negations involving I^2 using:
 - \neq
- it is important to remember, however, that $=$ is just an abbreviation of a two-place predicate

the identity relation

the identity relation

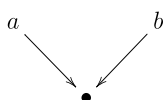
- $=$ expresses the identity relation
 - but what is the relation of identity?
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- B1 and B2 are identical
 - B1 is identical to B2
 - B1 and B2 are *exactly the same in all respects*
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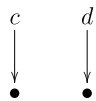
- B1 and B2 are identical
 - B1 is identical to B2
 - B1 and B2 are *one and the same thing*
-

- we have two relations:
 - **identity** (being one and the same thing)
 - **exact similarity** (being the same in all respects)
 - how are they related?
 - Leibniz's Law/Indiscernability of Identicals: if x is identical to y then x is exactly similar to y . (a logical truth)
 - Identity of Indiscernibles: if x is exactly similar to y then x is identical to y . (controversial)
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- a and b are identical (the names ' a ' and ' b ' pick out the same thing):



- c and d are non-identical (the names ' c ' and ' d ' pick out different things):



the predicate =

the predicate =

- consider the following sentences:
- B1 is a banana
- B1 is B2
- the word 'is' has a different meaning in each. in the first case, it is the 'is' of predication. in the second, it is the 'is' of identity.

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- we have the following glossary:
 - a : B1
 - b : B2
 - Bx : x is a banana

-
- we have the following translations:
 - Ba
 - $a = b$

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- recall that I^2 (or $=$) is part of the logical vocabulary
 - we do not put an entry for it in the glossary

translations into GPLI

translations into GPLI

- GPLI allows us to express more than just that a and b are identical
- let's take a look at the expressive power we have gained

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- return to our example:

P1.	John loves something
P2.	John does not love himself
C1.	John loves something other than himself

-
- here's a glossary:
 - j : John
 - Lxy : x loves y

-
- how should we translate the conclusion of the argument?
 - $\exists x (Ljx \wedge \neg I^2 jx)$
 - there is something that John loves and that thing is not John
 - John loves something other than himself!

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- let's look at some other examples . . .
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- here's a glossary:

j : John, m : Mary, s : Sam

Lxy : x loves y , Sx : x is a baker, Px : x is a person

- translate the following into GPLI:

John isn't Mary

$\neg I^2_{jm}$

$j \neq m$

- here's a glossary:

j : John, m : Mary, s : Sam

Lxy : x loves y , Sx : x is a baker, Px : x is a person

- translate the following into GPLI:

John loves someone other than himself

$\exists x (Px \wedge Ljx \wedge \neg I^2_{xj})$

$\exists x (Px \wedge Ljx \wedge x \neq j)$

- here's a glossary:

j : John, m : Mary, s : Sam

Lxy : x loves y , Sx : x is a baker, Px : x is a person

- translate the following into GPLI:

John loves everyone except Sam

$\forall x ((Px \wedge \neg I^2_{xs}) \rightarrow Ljx)$

$\forall x ((Px \wedge x \neq s) \rightarrow Ljx)$

- here's a glossary:

j : John, m : Mary, s : Sam

Lxy : x loves y , Sx : x is a baker, Px : x is a person

- translate the following into GPLI:

Some baker other than John loves Mary

$\exists x (Bx \wedge \neg I^2_{xj} \wedge Lxm)$

$\exists x (Bx \wedge x \neq j \wedge Lxm)$

- here's a glossary:

j : John, m : Mary, s : Sam

Lxy : x loves y , Sx : x is a baker, Px : x is a person

- translate the following into GPLI:

Everyone loves themselves

$$\forall x (Px \rightarrow Lxx)$$

- here's a glossary:

j : John, m : Mary, s : Sam
 Lxy : x loves y , Sx : x is a baker, Px : x is a person

- translate the following into GPLI:

John loves everyone other than himself

$$\forall x ((Px \wedge \neg I^2 xj) \rightarrow Ljx)$$

$$\forall x ((Px \wedge x \neq j) \rightarrow Ljx)$$

- here's a glossary:

j : John, m : Mary, s : Sam
 Lxy : x loves y , Sx : x is a baker, Px : x is a person

- translate the following into GPLI:

John loves everyone other than himself

$$\forall x ((Px \wedge \neg I^2 xj) \rightarrow Ljx) \wedge \neg Ljj$$

$$\forall x ((Px \wedge x \neq j) \rightarrow Ljx) \wedge \neg Ljj$$

- here's a glossary:

j : John, m : Mary, s : Sam
 Lxy : x loves y , Sx : x is a baker, Px : x is a person

- translate the following into GPLI:

John loves everything but does not love himself

$$\forall x Ljx \wedge \neg Ljj$$

- (this is a contradiction. it is false in every model.)

the semantics of $=$

the semantics of $=$

- since I^2 is a two-place predicate, its extension is a set of ordered pairs of members of the domain
 - once we have an extension for I^2 , the semantics for GPLI is just like the semantics of GPL
 - so all we need is an extension for I^2
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- in a model in which the domain is $\{1, 2, 3\}$, the extension of I^2 will be the following set of ordered pairs:
 - $\{\langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 3 \rangle\}$
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- in a model in which the domain is $\{Bill, Ben\}$, the extension of I^2 will be the following set of ordered pairs:
 - $\{\langle Bill, Bill \rangle, \langle Ben, Ben \rangle\}$

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- whatever the domain is, the extension of I^2 will contain exactly one ordered pair for each object in the domain: the pair containing that object in both the first and second place
 - the identity relation is a relation which holds between each object and itself and between no other objects
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- note! The extension of the identity predicate is not exactly the same on all models
- given a domain, however, the extension of the identity predicate is fixed
- this is not true of other predicates
- this is why it is part of the logical vocabulary

propositions of the form I^2ab

- I^2ab is true in a model if and only if the pair consisting of the referent of a and the referent of b is in the extension of I^2
 - Domain: $\{Bill, Ben\}$
 - Referents: $a: Bill, b: Bill$
 - Extension of I^2
 - $\{\langle Bill, Bill \rangle, \langle Ben, Ben \rangle\}$
 - I^2ab is true in a model if and only if a and b have the same referent on that model
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- let's run through a couple of examples
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- here's a model:
 - Domain: $\{Sydney, Canberra, Melbourne\}$
 - Referents: $a: Melbourne, b: Canberra, c: Sydney$
 - Extensions: $N: \{\langle Sydney, Canberra \rangle, \langle Canberra, Sydney \rangle, \langle Melbourne, Canberra \rangle, \langle Canberra, Melbourne \rangle\}$
 - is the following true or false in the model?
 - $\forall x \forall y (Nxy \rightarrow (I^2xb \vee I^2yb))$
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- here's a model:
 - Domain: $\{Sydney, Canberra, Melbourne\}$
 - Referents: $a: Melbourne, b: Canberra, c: Sydney$
 - Extensions: $N: \{\langle Sydney, Canberra \rangle, \langle Canberra, Sydney \rangle, \langle Melbourne, Canberra \rangle, \langle Canberra, Melbourne \rangle\}$
 - is the following true or false in the model?
 - $(I^2ab \rightarrow \forall x \forall y Nxy)$
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wrapping up

this lecture

- some of the limitations of GPL which motivate the shift to GPLI
- introducing GPLI
- the identity relation
- translations into GPLI
- the semantics of identity

next lecture

- lecture 22, trees for GPLI