

Cryptography 2/28

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- $\Pi(n) = \#\{p \leq n \mid p \in \text{prime}\}$
- $\Pi(n) \approx \frac{n}{\log(n)}$
- Factoring in \mathbb{Z} is hard but factoring in $\mathbb{F}_2[x]$ is easy
- In \mathbb{F}_{p^n} , the number of multiplicative generators is:

$$\frac{\phi(p^n - 1)}{p^n - 1}$$

Block Cipher

- Random substitution, while very safe, makes your codebook too big
- You can do a permutation cipher to make the codebook smaller but it's not all that safe tbh
 - Your keyspace is $n!$ and your ciphertext space is $|\Sigma|^n$
 - You can create a permutation matrix, where the inverse of the matrix is the decryption of the message
 - * **I'll try to remember to upload a picture of this at some point**
- You also have a linear cipher/hill cipher
 - $\vec{c} = A\vec{m}$
 - This is encryption, A is invertible $\iff \det(A)$ is a unit
 - plaintext and ciphertext spaces are Σ^n , the keyspace is...? Σ^{n^2} ?