Cryptography 2/7

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Repeating Squaring Algorithm

• First algorithm: - Simple but effective - Time complexity: $O(log(b)log^2(n))$ * I think he said this is a pessimistic time complexity? * We don't really care about the second half of the complexity Good algorithm - You can remove the last if statement and have the final return take its place modpower(a, b, n) # compute a**b mod n # assume a, n position integer, b non negative integer if b == 0: return 1 if b is even: return modpower(a**2 % n, b/2, n) if b is odd: return a * modpower(a, b - 1, n) % n • Second Algorithm - Assumbe $b = b_k b_{k-1} \cdots b_1$ and $b_k = 1$ - Better than the above algorithm when the base = 2# compute base**b mod n result = base for i in range(k - 1, 0, -1): result == result**2 % n if b_i == 1: result = result * base % n • Lets look at $base^{1010}$ - This means you're calculating $base^{2^3} \times base^{2^2}$ for the first algorithm * You have three squares and one multiplication with the base - This means you're calculating $((base^2)^2base)^2$ for the second algorithm

Chinese Remainder Theorem

- Suppose you have a 6 digit passcode, and one day your "friend" asks you for the remainder when you divide your passcode by 2 (i.e. is it even or odd?) you answer and reveal the last bit of information
 - The "friend" asks again for the remainder by 3, and the next day asks for the remainder by 5
 - "Basically the chinese remainder theorem tells you don't do that okay"

* You have three squares and one multiplication with the base

- You can get a passcode by dividing by 2, 3, 5, 7, 11 I think is what he said?
 - * Ohhhh just prime numbers in general
- You have:

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a \equiv b_1 \mod 2
a \equiv b_2 \mod 3
\vdots
a \equiv b_n \mod p_n
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- $\bullet\,$ It's called the chinese remainder theorem because it was created by a chinese general (General Sun)
 - Not sure if this is legitimate or if it was a set up for a joke but it was a good story that I'll write
 out later if I remember