

PPL 2/19

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- I really need to fucking pay attention today so I locked myself out of my phone. Here goes nothing.

LL Parsing

- It's been two minutes and I already opened GroupMe on my laptop send help

Removing Left Recursion (important to remove it for some reason)

- This has left recursion:

$$\begin{array}{l} \text{id-list} \rightarrow \text{id} \\ \quad \quad | \quad \text{id-list, id} \end{array}$$

- This fixes the left recursion

$$\begin{array}{l} \text{id-list} \rightarrow \text{id id-list-tail} \\ \text{id-list-tail} \rightarrow , \text{id id-list-tail} \\ \quad \quad | \quad \text{epsilon} \end{array}$$

- You can solve the common prefixes issue by “left-factoring” which is demonstrated as shown:

Example: $\text{stmt} \rightarrow \text{id} := \text{id} (\text{arg-list})$

$$\begin{array}{l} \text{Fixed: } \text{stmt} \rightarrow \text{id id-stmt-trail} \\ \text{id-stmt-trail} \rightarrow := \text{expr} \\ \quad \quad | \quad (\text{arg-list}) \end{array}$$

- There's also an issue with the “dangling else” that prevents grammars from being LL(1)
 - Dangling else basically means when there's an else that could go to two different if statements and you don't know which one it belongs to
- The example in the slides changes to ambiguous, dangling else grammar into a less natural grammar that can be parsed from bottom up but not top down, that's a lot of LaTeX to write so I'll try to remember to do it later
- You can employ explicit end-markers which is a better solution and fixes the issues with both of the previous grammars
- Usually you can use the ambiguous grammar with a *disambiguating rule* that says the else should go with the closest then, or the first two possible productions is the one to predict/reduce
- I still ended up getting distracted sorry friends