

# Design and Implementation of a Memory Management Simulator

## 1. Introduction

This project implements a Memory Management Simulator that models how an operating system manages memory at runtime. The simulator focuses on dynamic memory allocation, cache behavior, and virtual memory using paging, without implementing a real OS kernel.

The goal is to understand:

- Allocation strategies and fragmentation
  - Performance trade-offs in memory management
  - Interaction between virtual memory, caches, and physical memory
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## 2. System Overview

The simulator consists of four major subsystems:

1. Physical Memory Allocator
2. Buddy Memory Allocation System
3. Multilevel Cache Simulator
4. Virtual Memory System

All components are integrated through a command-line interface (CLI) that allows interactive testing.

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## 3. Physical Memory Model

### 3.1 Memory Representation

- Physical memory is simulated as a contiguous byte-addressable array
- Memory addresses are represented as offsets from base address 0
- Memory blocks are managed using explicit metadata structures

Physical Memory: [0 ..... N-1]

Each block contains:

- Start address
  - Block size
  - Allocation status (free/used)
  - Block ID
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## 4. Dynamic Memory Allocation Strategies

The simulator supports the following allocation strategies:

#### 4.1 First Fit

- Selects the first free block large enough for the request
- Fast but prone to fragmentation

#### 4.2 Best Fit

- Selects the smallest free block that satisfies the request
- Reduces wasted space but increases search time

#### 4.3 Worst Fit

- Selects the largest free block
- Attempts to leave large free regions available

#### 4.4 Allocation Procedure

For each `malloc(size)`:

1. Traverse free block list
2. Select block according to strategy
3. Split block if larger than required
4. Assign block ID and update metadata

#### 4.5 Deallocation Procedure

For each `free(block_id)`:

1. Mark block as free
2. Coalesce with adjacent free blocks
3. Update fragmentation statistics

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### 5. Buddy Memory Allocation System

#### 5.1 Design Rationale

The Buddy system is implemented to:

- Eliminate external fragmentation
- Enable fast coalescing using address arithmetic

#### 5.2 Key Properties

- Total memory size is a power of two
- Allocation sizes are rounded up to nearest power of two
- Minimum block size = 8 bytes
- Free blocks are stored in size-indexed free lists

#### 5.3 Buddy Address Calculation

`buddy_address = address XOR block_size`

## 5.4 Allocation Flow

1. Round request to power of two
2. Find smallest available block
3. Recursively split larger blocks if needed
4. Allocate and record internal fragmentation

## 5.5 Freeing and Coalescing

1. Identify buddy using XOR
2. If buddy is free, merge
3. Repeat recursively until merge not possible

## 5.6 Fragmentation in Buddy System

- Internal fragmentation exists due to rounding
  - External fragmentation = 0% (by design)
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## 6. Fragmentation and Performance Metrics

The simulator tracks:

### 6.1 Internal Fragmentation

Internal Fragmentation = Allocated Block Size – Requested Size

### 6.2 External Fragmentation

External Fragmentation =  $1 - (\text{Largest Free Block} / \text{Total Free Memory})$

### 6.3 Memory Utilization

Utilization =  $\text{Used Memory} / \text{Total Memory}$

### 6.4 Allocation Statistics

- Total allocation requests
  - Successful allocations
  - Failed allocations
  - Allocation success rate
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## 7. Multilevel Cache Simulation

	Cache Size	Block Size	Latency
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L1	128 B	16 B	1 cycle
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L2	512 B	16 B	5 cycles
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## 7.2 Cache Design

- Set-associative cache model
- FIFO replacement policy
- Each access records hit/miss

## 7.3 Cache Access Flow

CPU → L1 → L2 → Main Memory

- Miss penalties propagate to lower levels
  - Hit/miss statistics are tracked per level
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## 8. Virtual Memory Simulation

### 8.1 Address Translation

- Virtual memory implemented using paging
- Page table maps virtual pages to physical frames
- Page size is fixed and documented

### 8.2 Page Replacement Policy

- FIFO replacement
- Page faults trigger eviction
- Disk access latency simulated symbolically

### 8.3 Translation Flow

Virtual Address



Page Table Lookup



Physical Address



Cache Access



Memory

### 8.4 Metrics Tracked

- Page hits
- Page faults
- TLB statistics (optional)

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## 9. Command-Line Interface

Supported commands include:

- `init memory <size>`
- `set allocator <first_fit | best_fit | worst_fit | buddy>`
- `malloc <size>`
- `free <block_id>`
- `dump memory`
- `stats all`
- `access <virtual_address>`

The CLI enables interactive testing and demonstration.

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## 10. Assumptions and Limitations

Assumptions

- Single-process simulation
- No concurrency
- Simplified latency model

Limitations

- No real memory access
  - No write-back cache modeling
  - No multi-process address spaces
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## 11. Conclusion

This simulator successfully models:

- OS-level memory allocation strategies
- Fragmentation behavior
- Cache hierarchy performance
- Virtual memory address translation

The project provides hands-on insight into real operating system memory management design decisions and their performance implications.