Discussion 0

C, x86

09/03/24

Announcements

Sunday	Monday	Tuesday	Wednesday	Thursday	Friday	Saturday
				Homework 0 Release (8/29)		
	Labor Day	Project 0 Release (9/3)		C Review Session (9/5)	Homework 0 Due Early Drop Deadline (9/6)	
	Midterm Conflict Form Due Homework 1 Release (9/9)					Group Formation Form Due

Important Policies

Contact staff through Ed.

- Use cs162@eecs.berkeley.edu to reach head TAs and professor if Ed is insufficient.
- Only use individual emails for private matters.

3 midterms (check website for date and time).

- No conflicts allowed. Alternates are not guaranteed to happen.

Slip day policy:

- We provide 6 HW slip days and 7 Project slip days which are meant for emergencies
    - Extension form available for DSP and those with extenuating circumstances
    - Form is also accessible from the course website

Discussions 0 and 1 are **unassigned**.

- Feel free to attend different TAs' sections to find a teaching style that suits you best.
- Attendance for discussion 2 and onward are **mandatory**. Preference forms will be available soon.

Follow office hours policies of filling out a **detailed ticket** and being present in the OH room when your ticket is taken.

- All OH will be **in person**.
- Please check the Course calendar to find your OH room.
- Take advantage of empty office hours by **starting assignments early**.

Post your questions on Ed in the appropriate threads.

- No private debugging posts allowed.

C

Resources

- C Review Session
- CS 162 Ladder
  - Overview of C and 61C topics
- CS 61C Resources Page
  - C staff notes, GDB reference card
- Python Tutor C
  - Just like CS 61A's Python Tutor, but for C

Types

C is **statically typed** (i.e. types are known at compile time).

C is **weakly typed** (i.e. can cast between any types).

Primitive types are `char`, `short`, `int`, `long`, `float`.

Arrays are contiguous pieces of memory of a homogenous type

- Denoted with `[]` (e.g. `int []` for an array of integers).
- **String** is an array of `chars` with last element being `null`

Build compound types using `structs`

- Also contiguous in memory

Pointers are references that hold the address of an object in memory.

- Essentially just unsigned integers.
- Prefix a pointer with * to return the value at the memory address that the pointer is holding.
- Prefix a variable with & to return the memory address of the variable.

Little Endian	+3	+2	+1	+0	
0x7FFFFFFC	00	00	00	08	`int a = 8`
0x7FFFFFF8	FF	FF	FF	FF	`int b = -1`
0x7FFFFFF4	7F	FF	FF	FC	`int *p = &a`
0x7FFFFFF0	27	CE	00	'O'	27, CE are garbage
0x7FFFFFEC	'T'	'N'	'I'	'P'	`char s[] = "PINTO"`

0x7FFFFFE8	00	00	00	02	`struct point {`
					`  int x;`
					`  int y;`
					`};`
					
					`struct point pt = {16, 2}`

0x7FFFFFE4	00	00	00	10	

Memory

Typical C program is divided into five segments.

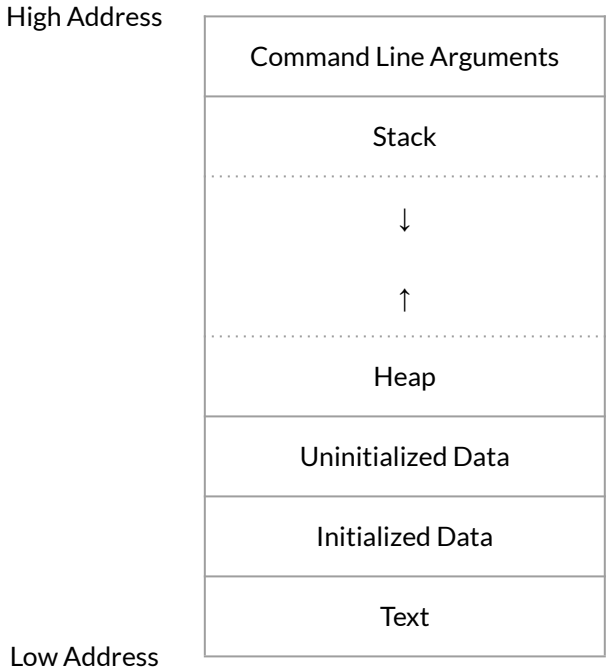
- **Text** contains machine code of the compiled program.
- **(Un)initialized data** contains (un)initialized global/static memory.
- **Heap** contains dynamically allocated memory.
- **Stack** contains local variables and arguments.
- Initialized strings and global constants may be stored in read-only segments (.rodata).

Think of memory as a giant array with elements of one byte.

- Memory addresses = indices of array.

Heap memory needs to be explicitly managed by the user.

- Allocate memory using `malloc`, `calloc`, `realloc` which return a pointer to a chunk of memory.
- Release memory using `free`.



GNU Debugger (GDB)

Need to learn how to use it for 162 even if you skidded by 61C without it.

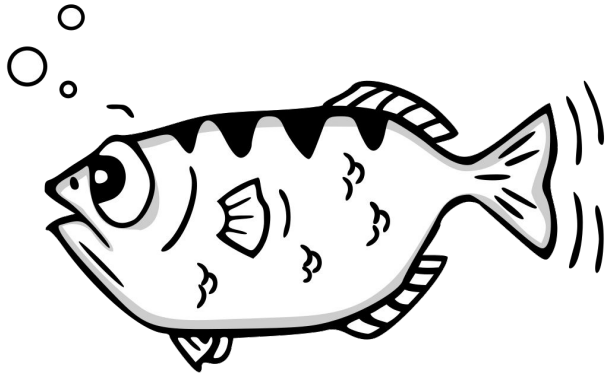
- Staff will not help during OH unless you are able to use GDB.

General workflow of using GDB.

1. Compile program using `-g` flag.
2. Start GDB using `gdb <executable name>`.
3. Set breakpoints using `break <line number>`. Can also break at functions using `break <function name>`.
4. Run program with run. If the program takes in arguments, pass in those after `run` (i.e. `run arg1 arg2 ...`).
5. Once breakpoint is hit, examine variables using `print`. Other commands like `display`, `watch`, and `set` are also useful.

Use GDB frequently to become familiar with the commands.

Check out GDB manual for more details.



Concept Check

1. Consider a valid double pointer `char** dbl_char` in a 32-bit system. What is returned by `sizeof(*dbl_char)`?

   Bonus question: Does sizeof(*dbl_char) error if dbl_char == NULL?

2. Consider strings `char* a = "162 is the best"` and `char b[] = "162 is the best"`. Are a and b different?

3. Consider the following struct declaration:

   ```
   struct point {
     int x;
     int y;
   };
   ```

   Point out a few differences between

   ```
   struct point p;
   printf("%d", p.x = 1);
   ```

   and

   ```
   struct point* p;
   printf("%d", p->x = 1);
   ```

Concept Check

1. **Consider a valid double pointer `char** dbl_char` in a 32-bit system. What is returned by `sizeof(*dbl_char)`?**

   Dereferencing a double pointer gives a single pointer. 32-bit systems have 32-bit = 4 byte memory addresses.

   Bonus question: Does sizeof(*dbl_char) error if dbl_char == NULL?

2. Consider strings `char* a = "162 is the best"` and `char b[] = "162 is the best"`. Are a and b different?

3. Consider the following struct declaration:

```
struct point {
  int x;
  int y;
};
```

   Point out a few differences between

struct point p;	and	struct point* p;
printf("%d", p.x = 1);		printf("%d", p->x = 1);

Concept Check

1. Consider a valid double pointer `char** dbl_char` in a 32-bit system. What is returned by `sizeof(*dbl_char)`?

   Dereferencing a double pointer gives a single pointer. 32-bit systems have 32-bit = 4 byte memory addresses.

   Bonus question: Does sizeof(*dbl_char) error if dbl_char == NULL? **No, since type sizes are known at compile time.**

2. Consider strings `char* a = "162 is the best"` and `char b[] = "162 is the best"`. Are a and b different?

3. Consider the following struct declaration:

```
struct point {
  int x;
  int y;
};
```

   Point out a few differences between

struct point p;	and	struct point* p;
printf("%d", p.x = 1);		printf("%d", p->x = 1);

Concept Check

1. Consider a valid double pointer `char** dbl_char` in a 32-bit system. What is returned by `sizeof(*dbl_char)`?

   Dereferencing a double pointer gives a single pointer. 32-bit systems have 32-bit = 4 byte memory addresses.

   Bonus question: Does sizeof(*dbl_char) error if dbl_char == NULL? No, since type sizes are known at compile time.

2. **Consider strings `char* a = "162 is the best"` and `char b[] = "162 is the best"`. Are a and b different?**

   Yes. a points to a string literal in the read-only segment (.rodata) while b resides on the stack.

3. Consider the following struct declaration:

```
struct point {
  int x;
  int y;
};
```

   Point out a few differences between

struct point p;	and	struct point* p;
printf("%d", p.x = 1);		printf("%d", p->x = 1);

Concept Check

1. Consider a valid double pointer `char** dbl_char` in a 32-bit system. What is returned by `sizeof(*dbl_char)`?

   Dereferencing a double pointer gives a single pointer. 32-bit systems have 32-bit = 4 byte memory addresses.

   Bonus question: Does sizeof(*dbl_char) error if dbl_char == NULL? No, since type sizes are known at compile time.

2. Consider strings `char* a = "162 is the best"` and `char b[] = "162 is the best"`. Are a and b different?

   Yes. a points to a string literal in the read-only segment (.rodata) while b resides on the stack.

3. **Consider the following struct declaration:**

```
struct point {
  int x;
  int y;
};
```

Point out a few differences between

struct point p;	**and**	**struct point* p;**
printf("%d", p.x = 1);		**printf("%d", p->x = 1);**

The first one allocates struct point that is the size of two ints. On the other hand, the second one puts a single pointer to a struct point on the stack, which is the size of a single int. Also, the second one will likely segfault, since the pointer is uninitialized.

Headers

```c
#include <stdio.h>
#include "lib.h"

int main(int argc, char** argv) {
  helper_args_t helper_args;
  helper_args.string = argv[0];
  helper_args.target = '/';
  char* result = helper_func(&helper_args);
  printf("%s\n", result);
  return 0;
}
```

App.c

```c
typedef struct helper_args {
#ifdef ABC
 char* aux;
#endif
 char* string;
 char target;
} helper_args_t;
char* helper_func(helper_args_t* args);
```

lib.h

```c
#include "lib.h"

char* helper_func(helper_args_t* args) {
  int i;
  for (i = 0; args->string[i] != '\0'; i++)
    if (args->string[i] == args->target)
      return &args->string[i + 1];
  return args->string;
}
```

lib.c

You build the program on a 64-bit machine as follows.

> gcc -c app.c -o app.o

> gcc -c lib.c -o lib.o

> gcc app.o lib.o -o app

1. What is the size of a `helper_args_t` struct?
2. Suppose you add a `#define ABC` at the top of `lib.h`. What is the size of a `helper_args_t` struct?

Headers

```c
#include <stdio.h>
#include "lib.h"

int main(int argc, char** argv) {
  helper_args_t helper_args;
  helper_args.string = argv[0];
  helper_args.target = '/';
  char* result = helper_func(&helper_args);
  printf("%s\n", result);
  return 0;
}
```

App.c

```c
typedef struct helper_args {
#ifdef ABC
 char* aux;
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 char* string;
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lib.h

```c
#include "lib.h"

char* helper_func(helper_args_t* args) {
  int i;
  for (i = 0; args->string[i] != '\0'; i++)
    if (args->string[i] == args->target)
      return &args->string[i + 1];
  return args->string;
}
```

lib.c

1. What is the size of a **`helper_args_t`** struct?

16 bytes. Only `char* string` and `char target` meaning 9 bytes but GCC pads structs.

2. Suppose you add a `#define ABC` at the top of `lib.h`. What is the size of a `helper_args_t` struct?

Headers

```c
#include <stdio.h>
#include "lib.h"

int main(int argc, char** argv) {
  helper_args_t helper_args;
  helper_args.string = argv[0];
  helper_args.target = '/';
  char* result = helper_func(&helper_args);
  printf("%s\n", result);
  return 0;
}
```

App.c

```c
typedef struct helper_args {
#ifdef ABC
 char* aux;
#endif
 char* string;
 char target;
} helper_args_t;
char* helper_func(helper_args_t* args);
```

lib.h

```c
#include "lib.h"

char* helper_func(helper_args_t* args) {
  int i;
  for (i = 0; args->string[i] != '\0'; i++)
    if (args->string[i] == args->target)
      return &args->string[i + 1];
  return args->string;
}
```

lib.c

1. What is the size of a `helper_args_t` struct?

16 bytes. Only `char* string` and `char target` meaning 9 bytes but GCC pads structs.

2. **Suppose you add a `#define ABC` at the top of `lib.h`. What is the size of a `helper_args_t` struct?**

24 bytes. Additional 8 bytes from `char* aux` since `ABC` is defined.

Headers

```c
#include <stdio.h>
#include "lib.h"

int main(int argc, char** argv) {
  helper_args_t helper_args;
  helper_args.string = argv[0];
  helper_args.target = '/';
  char* result = helper_func(&helper_args);
  printf("%s\n", result);
  return 0;
}
```

App.c

```c
typedef struct helper_args {
#ifdef ABC
 char* aux;
#endif
 char* string;
 char target;
} helper_args_t;
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#include "lib.h"

char* helper_func(helper_args_t* args) {
  int i;
  for (i = 0; args->string[i] != '\0'; i++)
    if (args->string[i] == args->target)
      return &args->string[i + 1];
  return args->string;
}
```

lib.c

3. Suppose you build the program in a different way with the original files (i.e. none of the changes from previous questions apply).

```
> gcc -DABC -c app.c -o app.o
> gcc -c lib.c -o lib.o
> gcc app.o lib.o -o app
```

The program will now exhibit undefined behavior. Why?

Headers

```c
#include <stdio.h>
#include "lib.h"

int main(int argc, char** argv) {
  helper_args_t helper_args;
  helper_args.string = argv[0];
  helper_args.target = '/';
  char* result = helper_func(&helper_args);
  printf("%s\n", result);
  return 0;
}
```

App.c

```c
typedef struct helper_args {
#ifdef ABC
 char* aux;
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 char* string;
 char target;
} helper_args_t;
char* helper_func(helper_args_t* args);
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lib.h

```c
#include "lib.h"

char* helper_func(helper_args_t* args) {
  int i;
  for (i = 0; args->string[i] != '\0'; i++)
    if (args->string[i] == args->target)
      return &args->string[i + 1];
  return args->string;
}
```

lib.c

3. Suppose you build the program in a different way with the original files (i.e. none of the changes from previous questions apply).

```
> gcc -DABC -c app.c -o app.o
> gcc -c lib.c -o lib.o
> gcc app.o lib.o -o app
```

The program will now exhibit undefined behavior. Why?

`app.c` is compiled with ABC defined but `lib.c` is not

- `main` stored `argv[0]` at address of `helper_args + 8`
- `helper_func` access address of `args` when accessing `args->string`
- First 8 bytes of `args` which `helper_func` is accessing is garbage

Debugging Segmentation Faults

This is a GDB exercise. Here's the Repl if you want to try it out on your own machine:

Debugging Segmentation Faults

```
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

1. We want to debug the program using GDB. How should we compile the program?

Debugging Segmentation Faults

```c
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```c
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

1. We want to debug the program using GDB. How should we compile the program?

```
gcc -g singer.c -o singer
```

Need a -g flag for debugging.

Debugging Segmentation Faults

```
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

2. When running the program without any arguments, what line does the segfault happen? Describe the memory operations happening in that line.

Debugging Segmentation Faults

```c
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```c
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
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  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

2. When running the program without any arguments, what line does the segfault happen? Describe the memory operations happening in that line.

Find segfaulting line by letting the program run until it encounters the fault.

```
> gcc -g singer.c -o singer
> gdb singer
(gdb) run
Starting program: /home/runner/intro/singer
Unsorted: "IU is the best char!"

Program received signal SIGSEGV, Segmentation
fault.
0x00005646308006c8 in swap (a=0x564630800904
"IU is the best singer!", i=1, j=21)
   at singer.c:6
6     a[i] = a[j];
```

Ignore “warning: Error disabling address space randomization: Operation not permitted” if using Replit.

Debugging Segmentation Faults

```c
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```c
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

2. When running the program without any arguments, what line does the segfault happen? Describe the memory operations happening in that line.

Use `backtrace` for more comprehensive breakdown.

```
(gdb) backtrace
#0  0x00005646308006c8 in swap
(a=0x564630800904 "IU is the best singer!",
i=1, j=21)
   at singer.c:6
#1  0x0000564630800773 in partition
(a=0x564630800904 "IU is the best singer!",
l=0, r=21)
   at singer.c:26
#2  0x00005646308007bd in sort
(a=0x564630800904 "IU is the best singer!",
l=0, r=21)
   at singer.c:36
#3  0x0000564630800861 in main (argc=1,
argv=0x7ffd04ac7098) at singer.c:51
```

Debugging Segmentation Faults

```c
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```c
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
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int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
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  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

2. When running the program without any arguments, what line does the segfault happen? Describe the memory operations happening in that line.

Two memory operations

1. Read from `a[j]`.
2. Write to `a[i]`.

Debugging Segmentation Faults

```
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

3. Run the program with and without an argument and observe the memory addresses of a in the segfaulting line. Why are the memory addresses so different?

Debugging Segmentation Faults

```c
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```c
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

3. Run the program with and without an argument and observe the memory addresses of a in the segfaulting line. Why are the memory addresses so different?

Break at line 6 using GDB.

```
> gdb singer
(gdb) break 6
Breakpoint 1 at 0x6ab: file singer.c, line 6.
(gdb) run
Starting program: /home/runner/intro/singer
Unsorted: "IU is the best singer!"

Breakpoint 1, swap (
   a=0x5624e4600904 "IU is the best singer!",
i=1, j=21)
   at singer.c:6
6     a[i] = a[j];
(gdb) print a
$1 = 0x5624e4600904 "IU is the best singer!"
```

Debugging Segmentation Faults

```
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

3. Run the program with and without an argument and observe the memory addresses of a in the segfaulting line. Why are the memory addresses so different?

Break at line 6 using GDB.

```
(gdb) run "Taeyeon is the best singer!"
The program being debugged has been started
already.
Start it from the beginning? (y or n) y
Starting program: /home/runner/intro/asuna
"Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"

Breakpoint 1, swap (
   a=0x7ffcce01bfe5 "Taeyeon is the best
singer!", i=1, j=26)
   at asuna.c:6
(gdb) print a
$2 = 0x7ffcce01bfe5 "Taeyeon is the best
singer!"
```

Debugging Segmentation Faults

```
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

3. Run the program with and without an argument and observe the memory addresses of a in the segfaulting line. Why are the memory addresses so different?

No argument: 0x5624e4600904

- Statically defined strings stored in read-only segment (.rodata).

With argument: 0x7ffcce01bfe5

- Arguments are passed in through the stack.

Debugging Segmentation Faults

```
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

4. How should the code be changed to fix the segfault?

Debugging Segmentation Faults

```
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = "IU is the best singer!";

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

4. How should the code be changed to fix the segfault?

Write to a[i] is the problem since it's in read-only data → need to put a in writable memory.

Allocate memory on the heap and put the default string on there.

- Accomplish with `malloc` followed by `strcpy`.
- Equivalently, use `strdup`.

Replace line 48.

Debugging Segmentation Faults

```c
void swap(char* a, int i, int j) {
  char t = a[i];
  a[i] = a[j];
  a[j] = t;
}

int partition(char* a, int l, int r){
  int pivot = a[l];
  int i = l, j = r+1;

  while (1) {
    do
      ++i;
    while (a[i] <= pivot && i <= r);

    do
      --j;
    while (a[j] > pivot);

    if (i >= j)
      break;

    swap(a, i, j);
  }

  swap(a, l, j);

  return j;
}
```

```c
void sort(char* a, int l, int r){
  if (l < r){
    int j = partition(a, l, r);
    sort(a, l, j-1);
    sort(a, j+1, r);
  }
}

int main(int argc, char** argv){
  char* a = NULL;

  if (argc > 1)
    a = argv[1];
  else
    a = strdup("IU is the best singer!");

  printf("Unsorted: \"%s\"\n", a);
  sort(a, 0, strlen(a) - 1);
  printf("Sorted  : \"%s\"\n", a);
}
```

```
> ./singer "Taeyeon is the best singer!"
Unsorted: "Taeyeon is the best singer!"
Sorted  : "    !Tabeeeeeghiinnorssstty"
> ./singer
Unsorted: "IU is the best singer!"
Segmentation fault (core dumped)
```

4. How should the code be changed to fix the segfault?

Write to a[i] is the problem since it's in read-only data → need to put a in writable memory.

Allocate memory on the heap and put the default string on there.

- Accomplish with `malloc` followed by `strcpy`.
- Equivalently, use `strdup`.

Replace line 48.

x86

Registers

Registers are small storage spaces directly on the processor.

- Allows for fast memory access.

General purpose registers (GPR) store both data and addresses.

- x86 has 8, RISC-V has 32 (x0-x31).
- Started as 16-bits, extend to 32-bit using e prefix (e.g. EAX for AX).
- Access 8-bit LSB by replacing last letter with l (e.g. AL for AX).
- Access 8-bit MSB by replacing last letter with h (e.g. AH for AX) only for AX, BX, CX, DX.

Instruction pointer register holds address of next instruction to execute.

- Called ip which can be extended with e prefix like a GPR.
- Can't be read/modified like a GPR using regular memory instructions.

	Name	Purpose
AX	Accumulator	I/O port access, arithmetic, interrupt calls
BX	Base	Base pointer for memory access
CX	Counter	Loop counting, bit shifts
DX	Data	I/O port access, arithmetic, interrupt calls
SP	Stack Pointer	Top address of stack
BP	Base Pointer	Base address of stack
SI	Source Index	Source for stream operations (e.g. string)
DI	Destination Index	Destination for stream operations (e.g. string)

Syntax

Use **AT&T Syntax** not Intel which is used by GCC.

Prefix registers with **%** (e.g. `%eax`), constants with **$** (e.g. `$4`).

General structure is **`inst src, dest`**.

Address memory with **`offset(base, index, scale)`**.

- `base`, `index` = registers, `offset` = any integer, `scale`= 1, 2, 4, or 8.
- Accesses data at address `base + index * scale + offset`.
- All parameters optional but will see base and offset usually.
- Using `lea` instruction will operate on the address itself not contents.

Suffix instructions to signify operand size.

- Not always necessary but should use them regardless.

`mov 8(%ebx), %eax`	Move contents from address EBX + 8 into EAX
`mov %ecx, -4(%esi, %ebx, 8)`	Move contents in ECX into address ESI + 8 * EBX - 4
`lea 8(%ebx), %eax`	Puts EBX + 8 into EAX

`movb $0, (%esp)`	Zero out a single byte from ESP
`movw $0, (%esp)`	Zero out two bytes from ESP
`movl $0, (%esp)`	Zero out four bytes from ESP

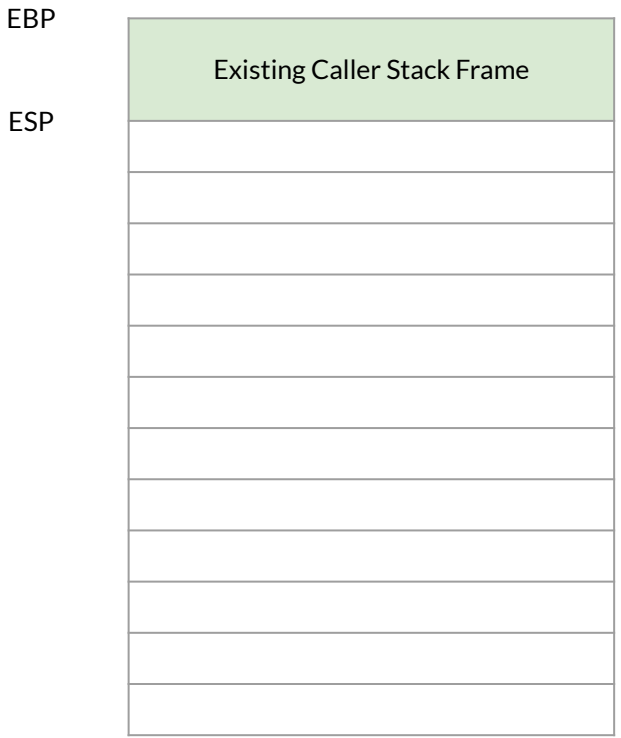
Calling Convention

Calling convention is a procedure for how to call and return from functions.

- Specifies stack management, argument passing, register saving, etc.
- One set of rule *each* for the **caller** and **callee**

Many different calling conventions, will use the one defined in **i386 System V ABI** in this class.

Calling Convention

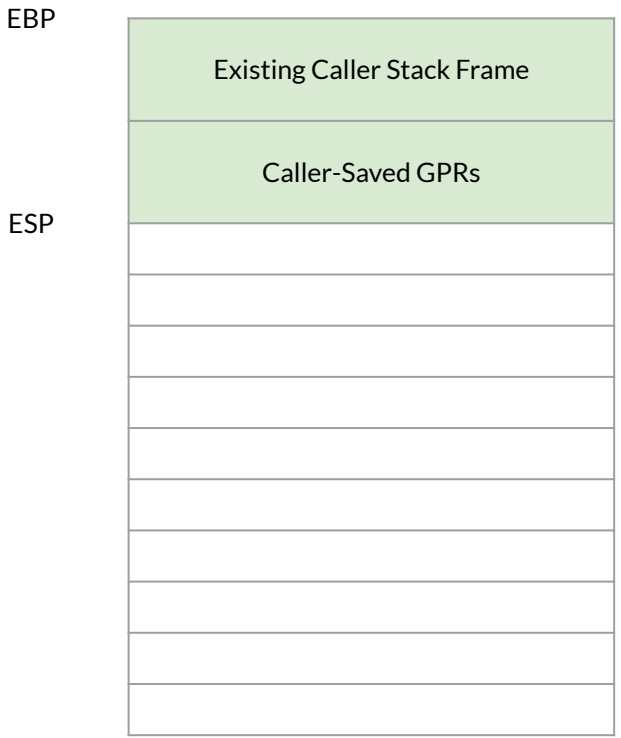


Calling Convention

Caller

Before calling the function (i.e. **prologue**),

1. **Save caller-saved GPRs (EAX, ECX, EDX) onto the stack *if needed after the function call*.**

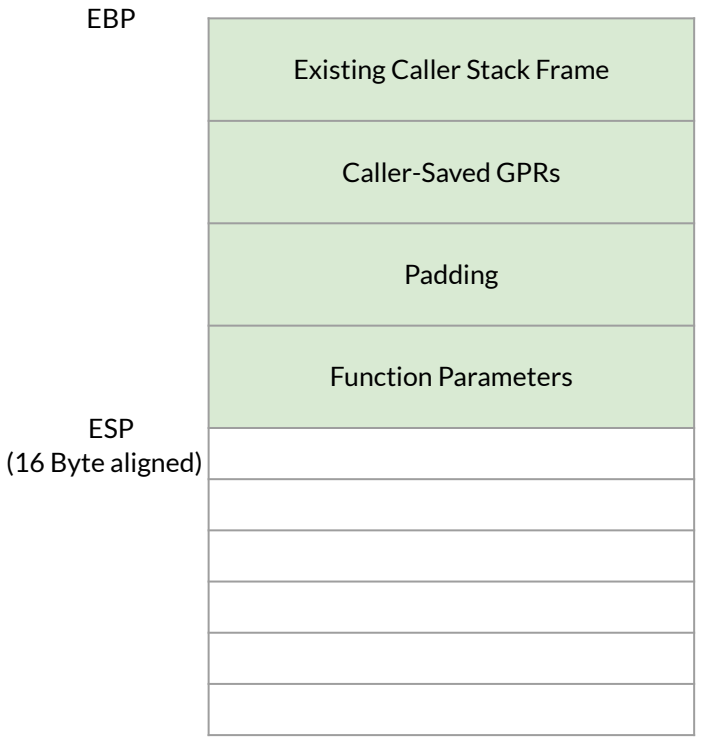


Calling Convention

Caller

Before calling the function (i.e. **prologue**),

1. Save caller-saved GPRs (EAX, ECX, EDX) onto the stack *if needed after the function call*.
2. **Push parameters onto the stack in reverse order. Add necessary padding *before the parameters* to ensure a 16-byte alignment.**



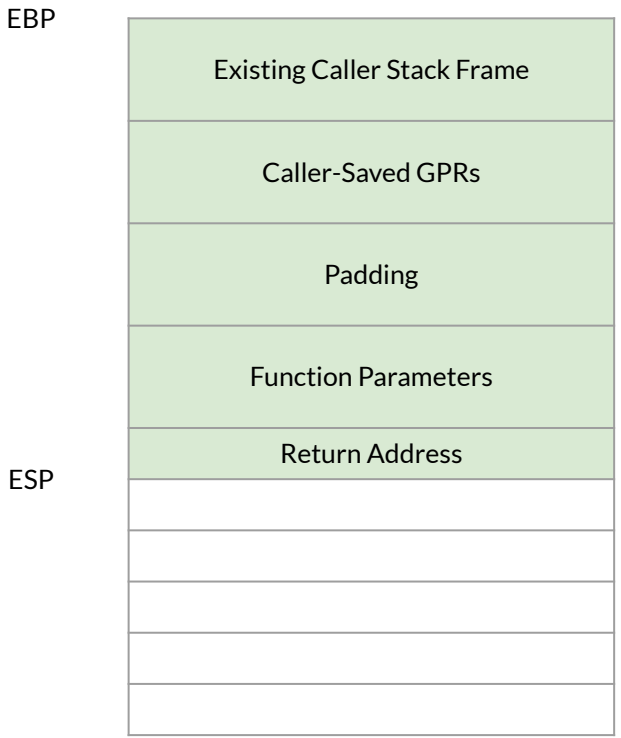
Calling Convention

Caller

Before calling the function (i.e. **prologue**),

1. Save caller-saved GPRs (EAX, ECX, EDX) onto the stack *if needed after the function call*.
2. Push parameters onto the stack in reverse order. Add necessary padding *before the parameters* to ensure a 16-byte alignment.

Call function by pushing the return address onto the stack and jumping to the function.

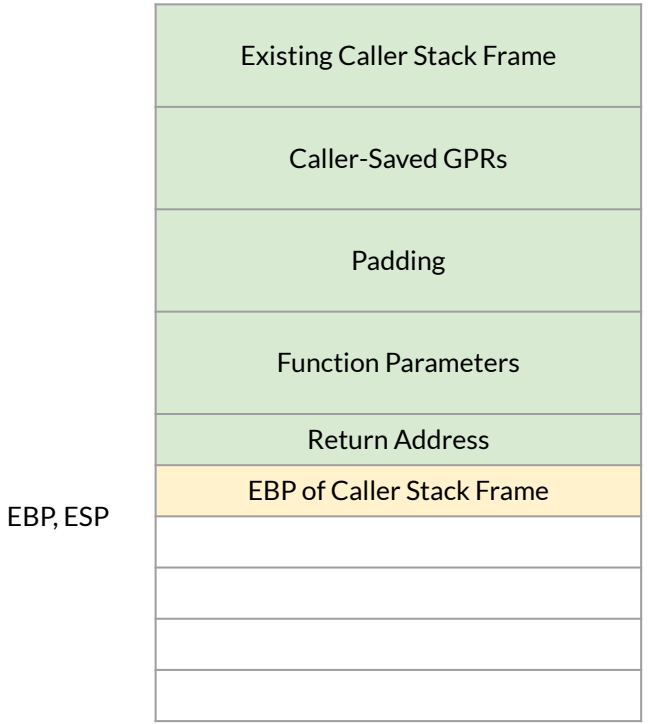


Calling Convention

Callee

Before executing any function logic (i.e. prologue),

1. **Push EBP onto the stack and set EBP to be the new ESP.**

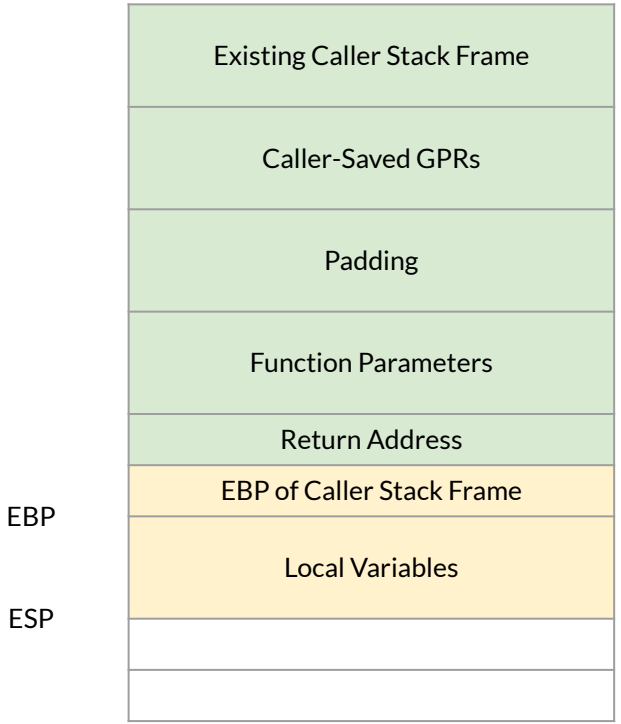


Calling Convention

Callee

Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. **Allocate stack space for local variables.**

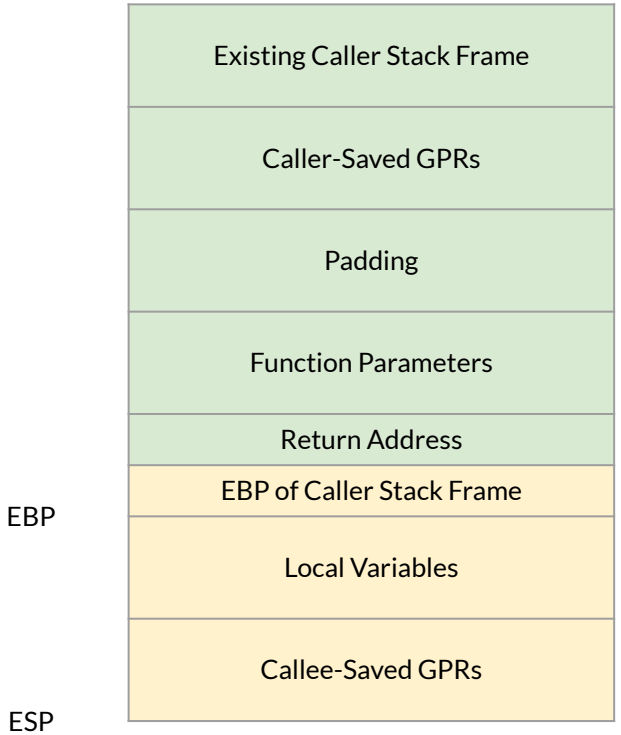


Calling Convention

Callee

Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. Allocate stack space for local variables.
3. **Save callee-saved GPRs (EBX, EDI, ESI) onto the stack *if used during the function logic.***



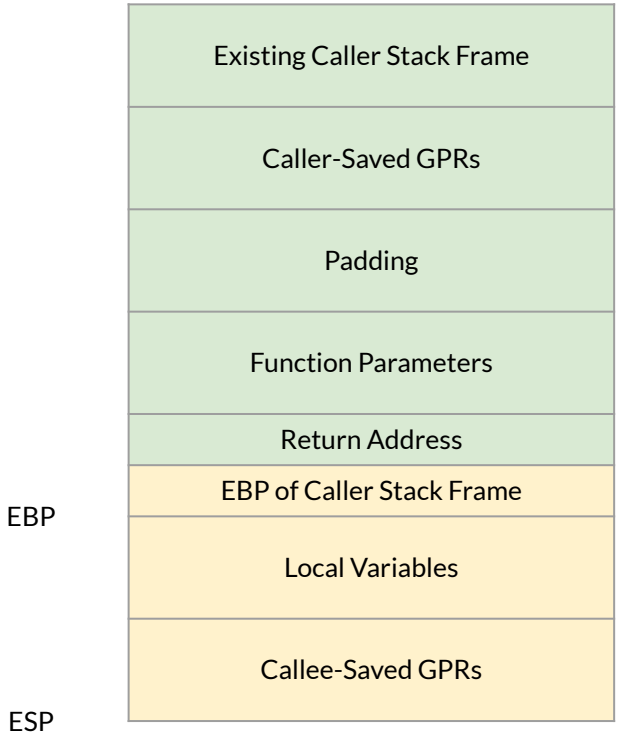
Calling Convention

Callee

Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. Allocate stack space for local variables.
3. Save callee-saved GPRs (EBX, EDI, ESI) onto the stack *if used during the function logic*.

Perform function logic.



Calling Convention

Callee

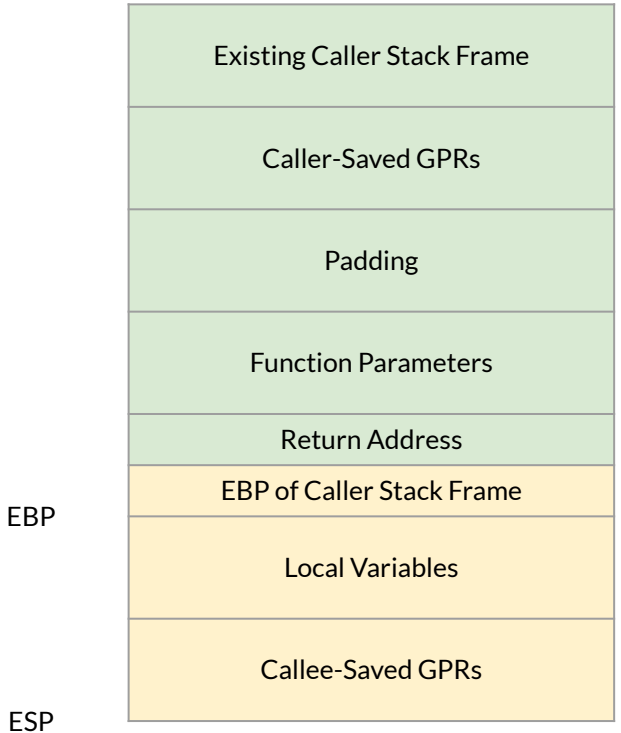
Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. Allocate stack space for local variables.
3. Save callee-saved GPRs (EBX, EDI, ESI) onto the stack *if used during the function logic*.

Perform function logic.

Before returning (i.e. **epilogue**),

1. **Store return value in EAX.**



Calling Convention

Callee

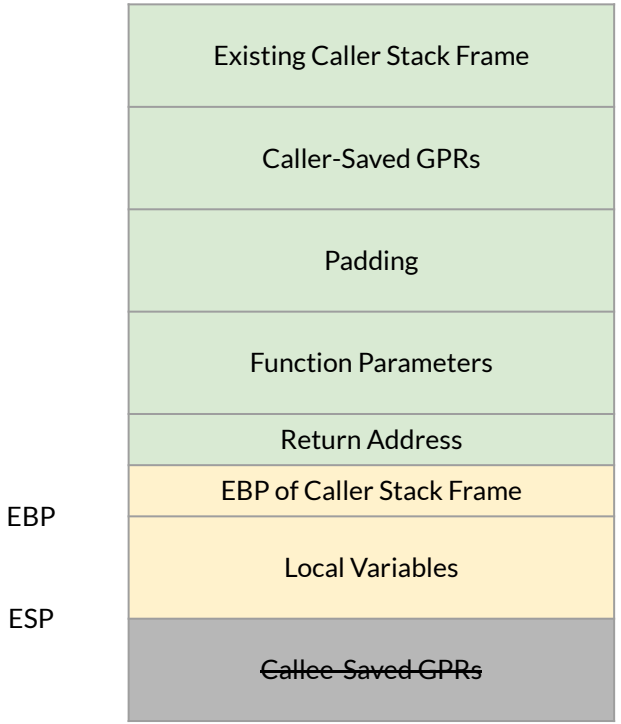
Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. Allocate stack space for local variables.
3. Save callee-saved GPRs (EBX, EDI, ESI) onto the stack *if used during the function logic*.

Perform function logic.

Before returning (i.e. **epilogue**),

1. Store return value in EAX.
2. **Restore callee-saved GPRs if any from the prologue.**



Calling Convention

Callee

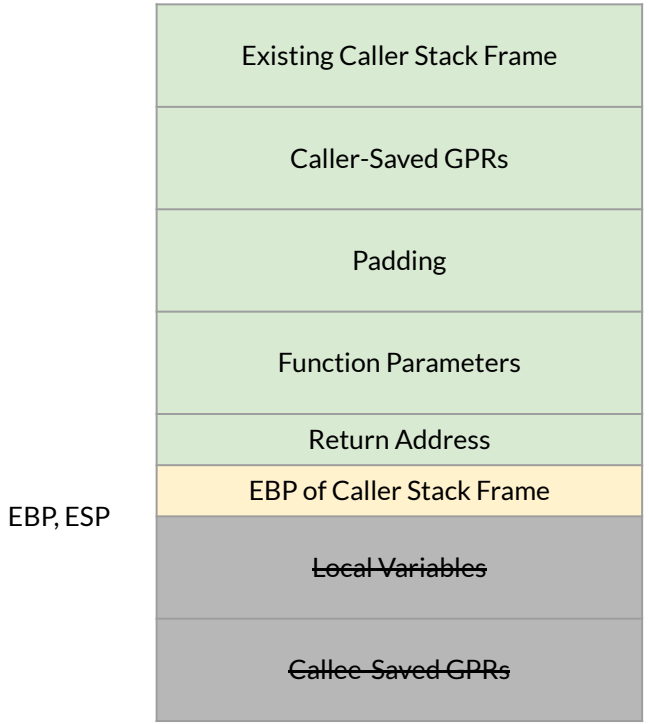
Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. Allocate stack space for local variables.
3. Save callee-saved GPRs (EBX, EDI, ESI) onto the stack *if used during the function logic*.

Perform function logic.

Before returning (i.e. **epilogue**),

1. Store return value in EAX.
2. Restore callee-saved GPRs if any from the prologue.
3. **Deallocate local variables.**



Calling Convention

Callee

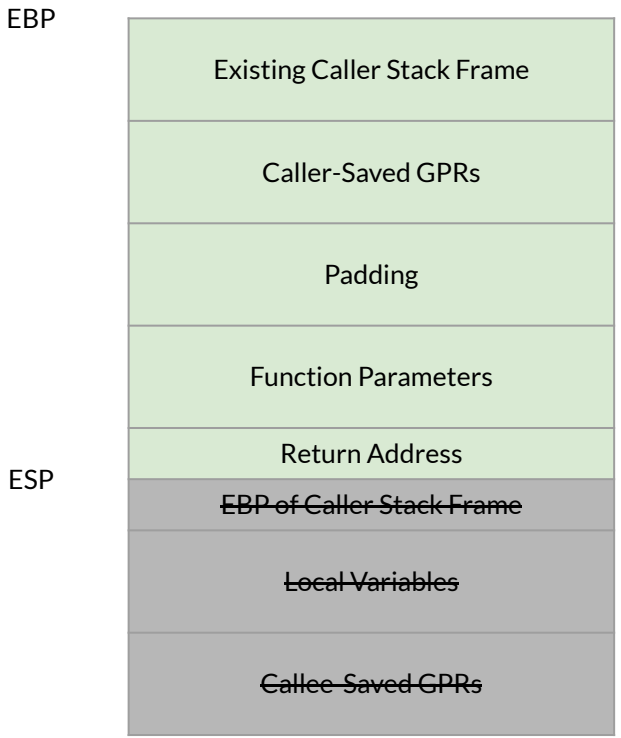
Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. Allocate stack space for local variables.
3. Save callee-saved GPRs (EBX, EDI, ESI) onto the stack *if used during the function logic*.

Perform function logic.

Before returning (i.e. **epilogue**),

1. Store return value in EAX.
2. Restore callee-saved GPRs if any from the prologue.
3. Deallocate local variables.
4. **Restore caller's EBP from stack.**



Calling Convention

Callee

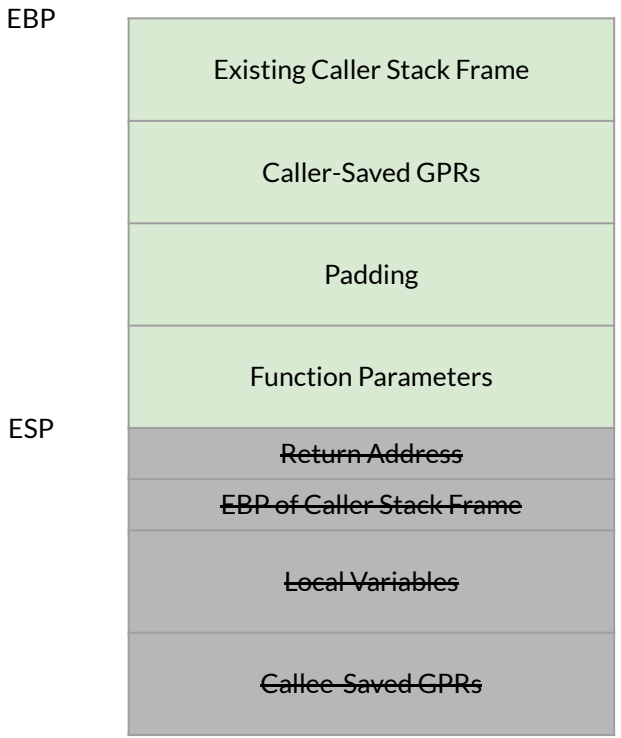
Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. Allocate stack space for local variables.
3. Save callee-saved GPRs (EBX, EDI, ESI) onto the stack *if used during the function logic*.

Perform function logic.

Before returning (i.e. **epilogue**),

1. Store return value in EAX.
2. Restore callee-saved GPRs if any from the prologue.
3. Deallocate local variables.
4. Restore caller's EBP from stack.
5. **Return from function call by popping the return address pushed by the caller in its prologue and jumping to it.**



Calling Convention

Caller

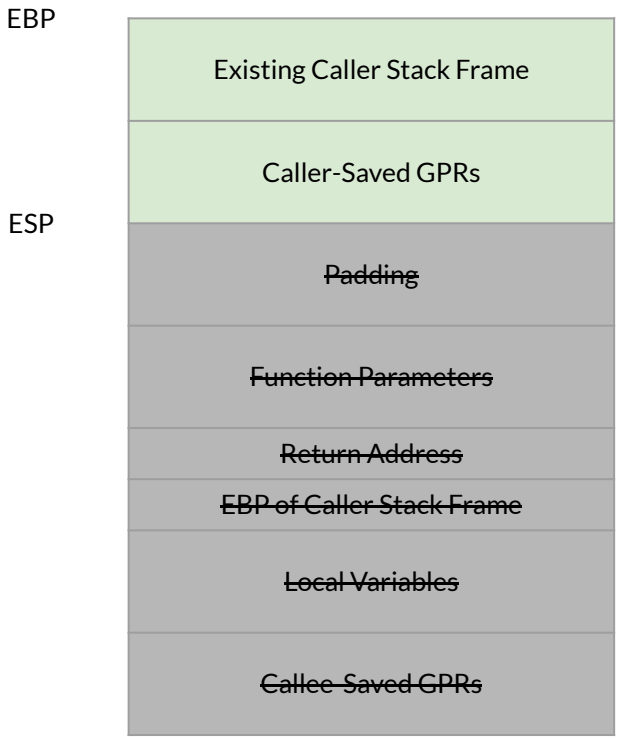
Before calling the function (i.e. **prologue**),

1. Save caller-saved GPRs (EAX, ECX, EDX) onto the stack *if needed after the function call*.
2. Push parameters onto the stack in reverse order. Add necessary padding *before the parameters* to ensure a 16-byte alignment.

Call function by pushing the return address onto the stack and jumping to the function.

Once function call returns (i.e. **epilogue**),

1. **Remove parameters from the stack.**



Calling Convention

Caller

Before calling the function (i.e. **prologue**),

1. Save caller-saved GPRs (EAX, ECX, EDX) onto the stack *if needed after the function call*.
2. Push parameters onto the stack in reverse order. Add necessary padding *before the parameters* to ensure a 16-byte alignment.

Call function by pushing the return address onto the stack and jumping to the function.

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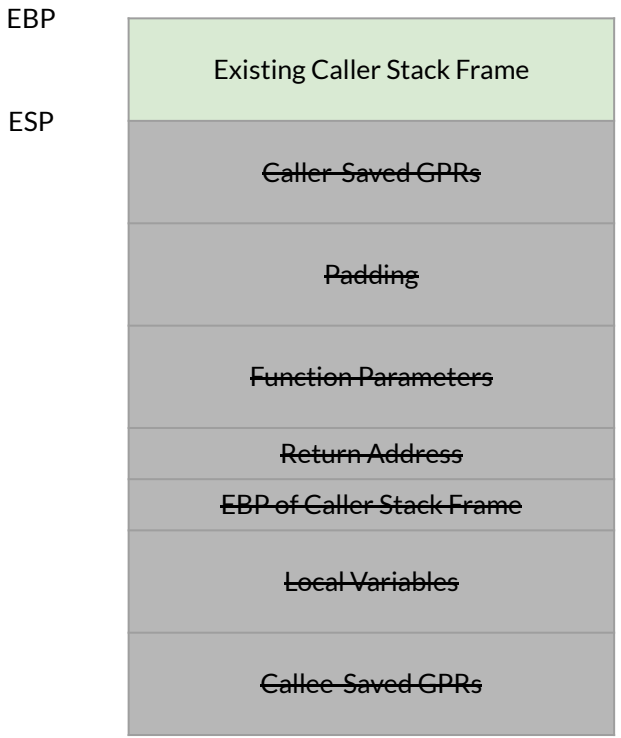

Calling Convention

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Call function by pushing the return address onto the stack and jumping to the function.

Once function call returns (i.e. **epilogue**),

1. Remove parameters from the stack.
2. Restore caller-saved GPRs if any from the prologue.

Callee

Before executing any function logic (i.e. prologue),

1. Push EBP onto the stack and set EBP to be the new ESP.
2. Allocate stack space for local variables.
3. Save callee-saved GPRs (EBX, EDI, ESI) onto the stack *if used during the function logic*.

Perform function logic.

Before returning (i.e. **epilogue**),

1. Store return value in EAX.
2. Restore callee-saved GPRs if any from the prologue.
3. Deallocate local variables.
4. Restore caller's EBP from stack.
5. Return from function call by popping the return address pushed by the caller in its prologue and jumping to it.

Calling Convention

Instruction	Purpose	Effective
`pushl src`	Push `src` onto stack	`subl $4, %esp` `movl src, (%esp)`
`popl dest`	Pop from stack into `dest`	`movl (%esp), dest` `addl $4, %esp`
`call addr`	Push return address onto stack and jump to `addr`	`pushl %eip` `jump addr`
`leave`	Restore EBP and ESP to previous stack frame	`movl %ebp, %esp` `popl %ebp`
`ret`	Pop return address from stack and jump to it	`popl %eip`

Concept Check

1. Between SP and BP, which has a higher memory address?

2. Write three different ways to clear the EAX register (i.e. store a 0).

3. True or False: Right before the caller jumps to the desired function, the stack must be 16-byte aligned.

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   BP. Stack grows downwards → top of the stack (SP) moves towards lower addresses

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movl $0, %eax          subl %eax, %eax          xorl %eax, %eax
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2. Write three different ways to clear the EAX register (i.e. store a 0).

   ```
   movl $0, %eax          subl %eax, %eax          xorl %eax, %eax
   ```

3. **True or False: Right before the caller jumps to the desired function, the stack must be 16-byte aligned.**

   False. Stack needs to be 16-byte aligned *after parameters have been pushed* onto the stack. Return address is pushed right before jumping.

Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
6          subl    $16, %esp
7          movl    8(%ebp), %edx
8          movl    12(%ebp), %eax
9          addl    %edx, %eax
10         subl    16(%ebp), %eax
11         movl    %eax, -4(%ebp)
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25         pushl   $5
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27         pushl   $3
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```

1. Which lines of code correspond to a caller/callee prologue?

Stack Frame

```
                call.s
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2          .zero   4
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34         leave
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```

1. Which lines of code correspond to a caller/callee prologue?

`foo/bar` as caller/callee

24-27 = caller prologue

29 = caller epilogue

4-6 = callee prologue

13-15 = callee epilogue

`foo` as callee

17-20 = callee prologue

33-35 = callee epilogue

Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
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```

2. What does line 19 do in `call.s`? Why is it necessary?

Stack Frame

call.s

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```

2. What does line 19 do in `call.s`? Why is it necessary?

   Saves EBX register since it's callee-saved and `foo` uses it. Doesn't matter that `bar` never uses EBX.

Stack Frame

call.s

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1  p:
2          .zero   4
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```

3. Why is EDX not saved by `foo` before calling `bar` despite the register being overwritten in `bar`?

Stack Frame

call.s

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33         movl    -4(%ebp), %ebx
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```

3. Why is EDX not saved by `foo` before calling `bar` despite the register being overwritten in `bar`?

   EDX does not need to persist after the function call.

Stack Frame

call.s

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4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	
EBX	
EDX	

EBP

ESP

Stack Frame of foo's Caller

Stack Frame

call.s

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EAX	
EBX	
EDX	

EBP

Stack Frame of foo's Caller
Padding
b
a
Return Addr of foo's Caller

ESP

Stack Frame

call.s

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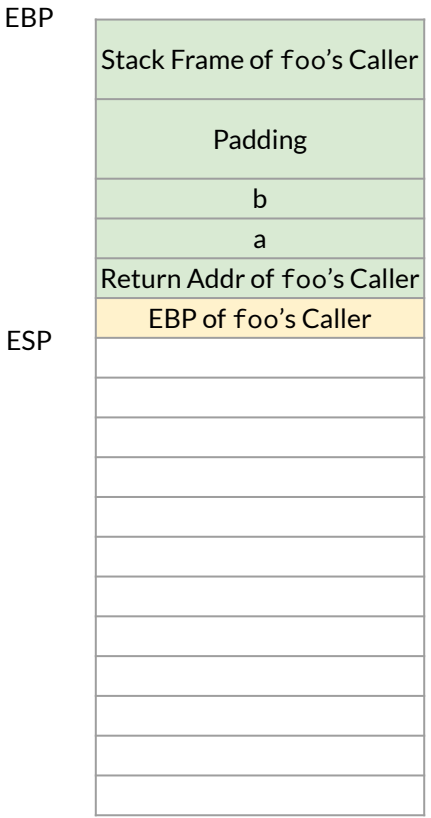
EAX	
EBX	
EDX	



Stack Frame

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EAX	
EBX	
EDX	

EBP, ESP

Stack Frame of foo's Caller
Padding
b
a
Return Addr of foo's Caller
EBP of foo's Caller

Stack Frame

call.s

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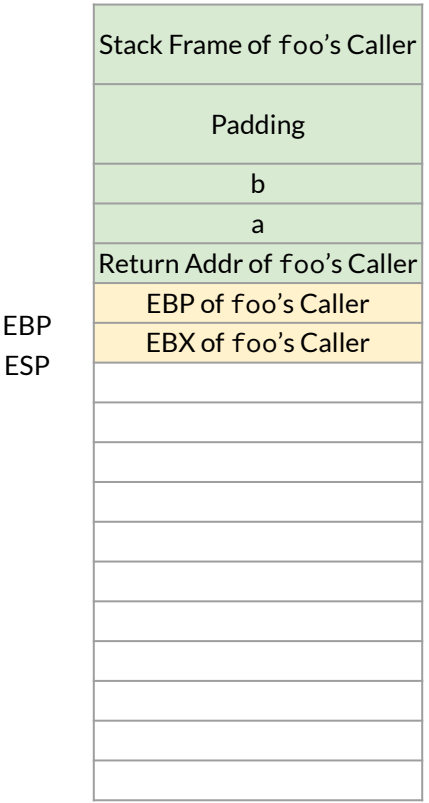
EAX	
EBX	
EDX	



Stack Frame

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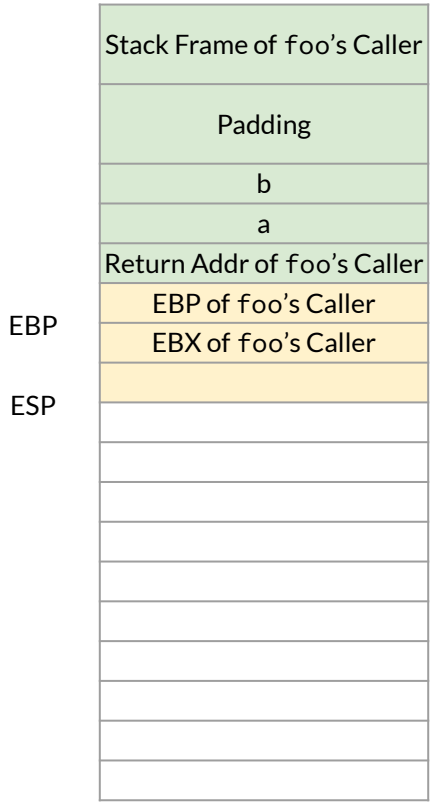
EAX	
EBX	
EDX	



Stack Frame

call.s

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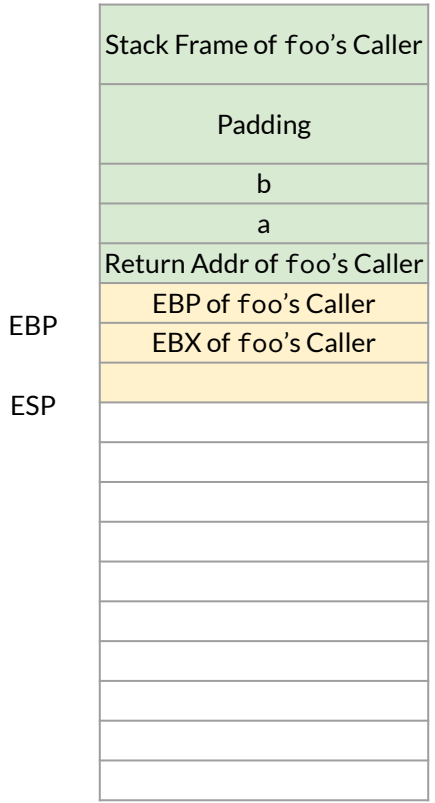
4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	
EBX	
EDX	a



Stack Frame

call.s

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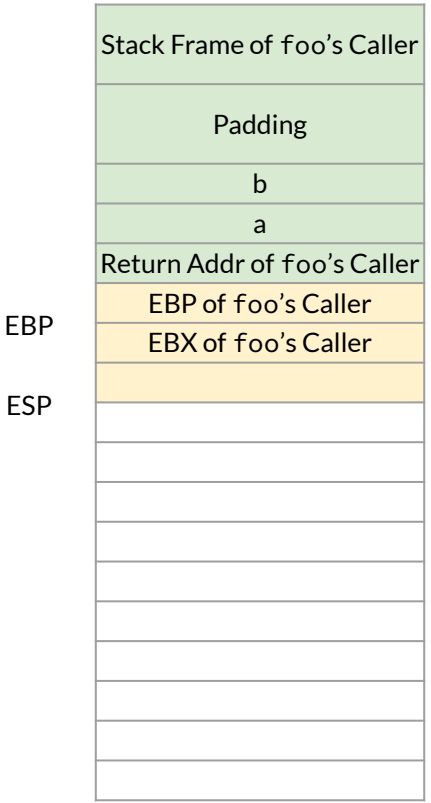
EAX	b
EBX	
EDX	a



Stack Frame

call.s

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2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
6          subl    $16, %esp
7          movl    8(%ebp), %edx
8          movl    12(%ebp), %eax
9          addl    %edx, %eax
10         subl    16(%ebp), %eax
11         movl    %eax, -4(%ebp)
12         movl    -4(%ebp), %eax
13         addl    $1, %eax
14         leave
15         ret
16 foo:
17         pushl   %ebp
18         movl    %esp, %ebp
19         pushl   %ebx
20         subl    $4, %esp
21         movl    8(%ebp), %edx
22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
26         pushl   $4
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31         movl    %eax, p
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35         ret
```
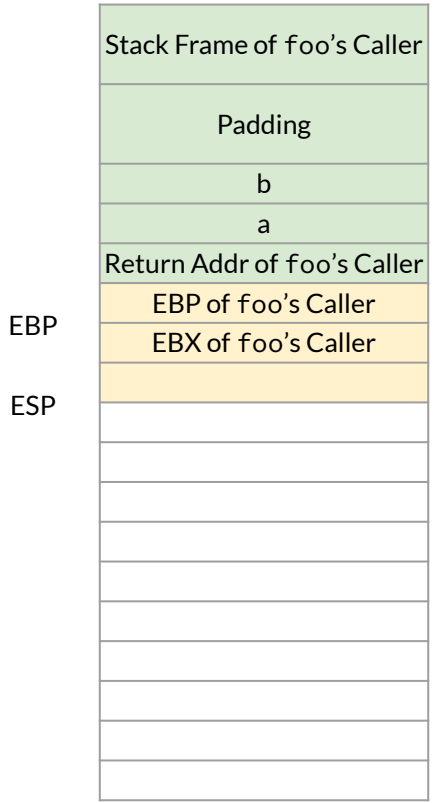

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	b
EBX	a + b
EDX	a



Stack Frame

call.s

```
1  p:
2          .zero   4
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4          pushl   %ebp
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6          subl    $16, %esp
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21         movl    8(%ebp), %edx
22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
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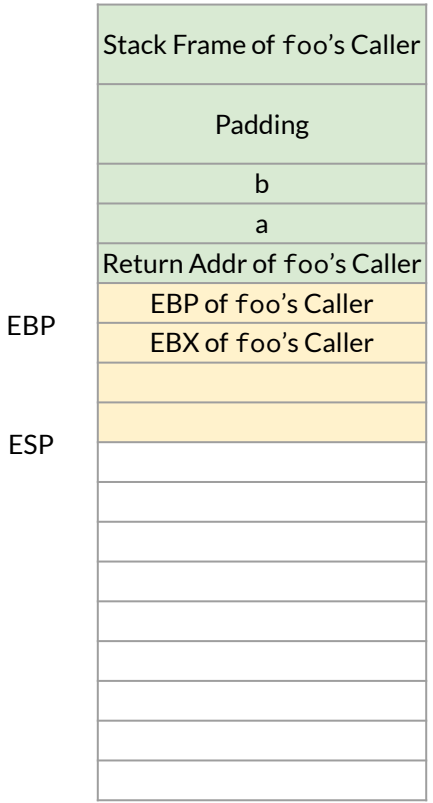

Stack Frame

call.s

```
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4          pushl   %ebp
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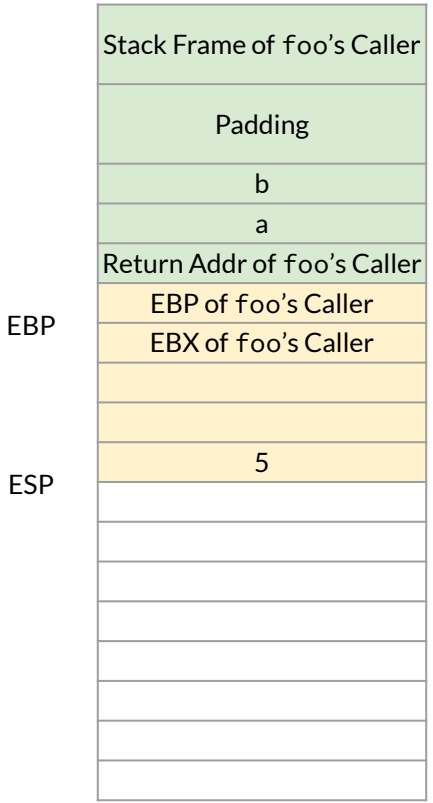

Stack Frame

call.s

```
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4          pushl   %ebp
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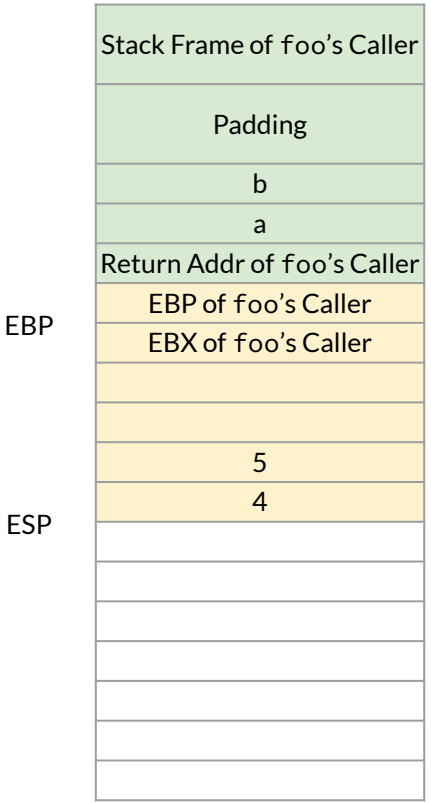

Stack Frame

call.s

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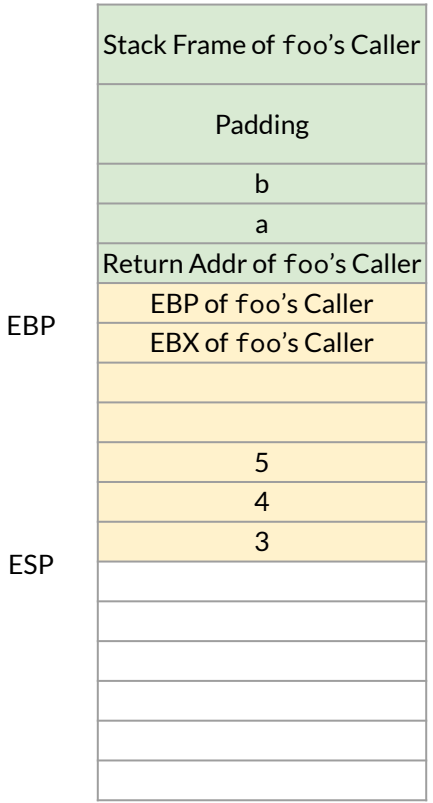

Stack Frame

call.s

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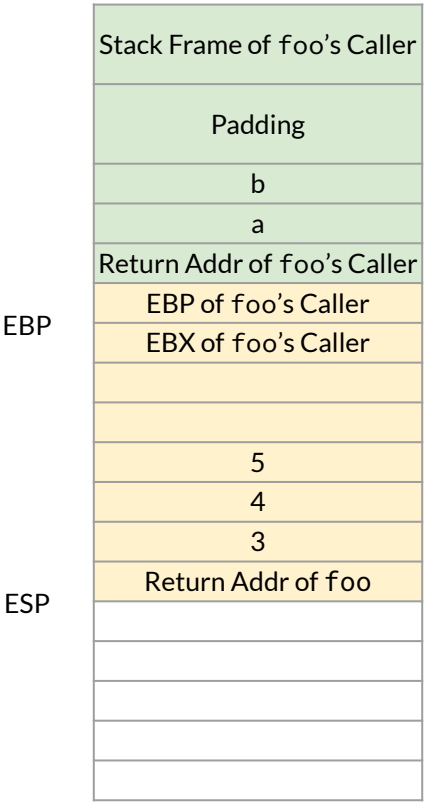

Stack Frame

call.s

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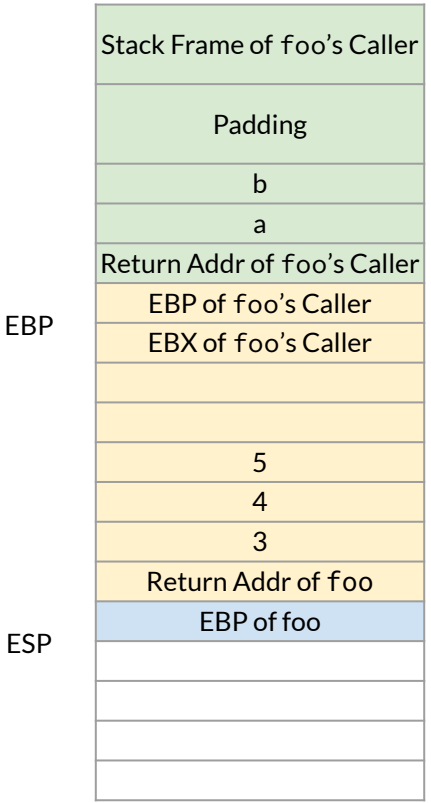

Stack Frame

call.s

```
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4          pushl   %ebp
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```

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

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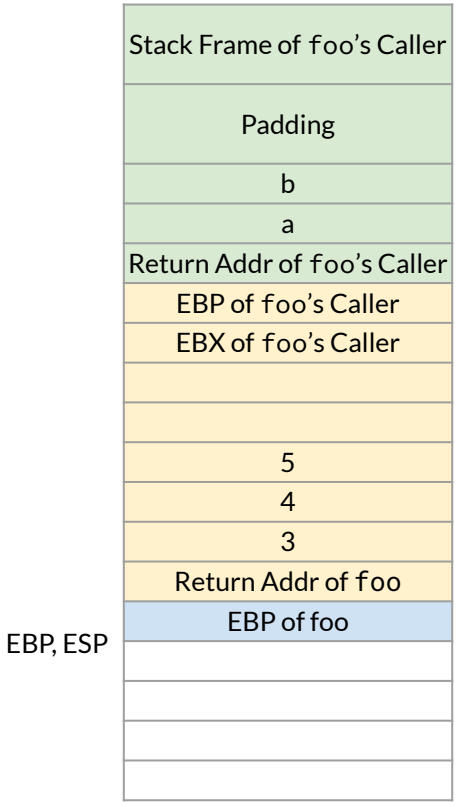

Stack Frame

call.s

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1  p:
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34         leave
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```
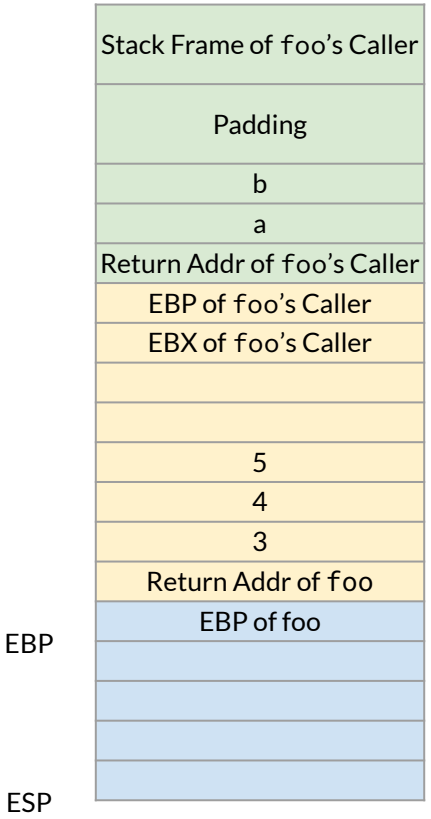

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	b
EBX	a + b
EDX	a



Stack Frame

call.s

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19         pushl   %ebx
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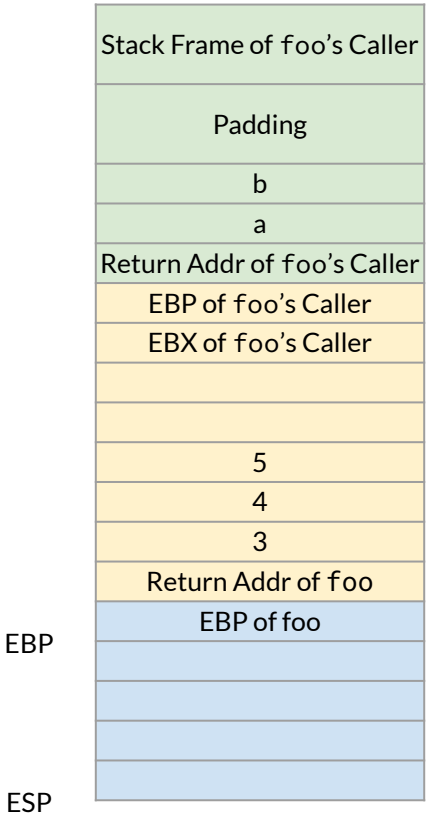
32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	b
EBX	a + b
EDX	3



Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
6          subl    $16, %esp
7          movl    8(%ebp), %edx
8          movl    12(%ebp), %eax
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13         addl    $1, %eax
14         leave
15         ret
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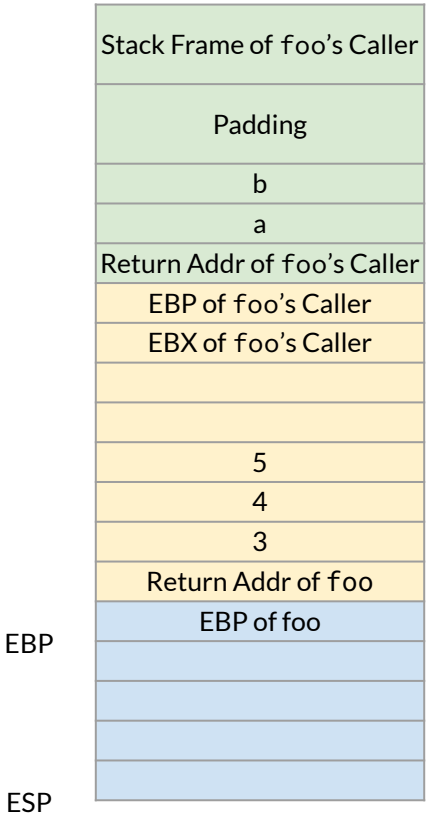
19         pushl   %ebx
20         subl    $4, %esp
21         movl    8(%ebp), %edx
22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
26         pushl   $4
27         pushl   $3
28         call    bar
29         addl    $16, %esp
30         addl    %ebx, %eax
31         movl    %eax, p
32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	4
EBX	a + b
EDX	3



Stack Frame

call.s

```
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19         pushl   %ebx
20         subl    $4, %esp
21         movl    8(%ebp), %edx
22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
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28         call    bar
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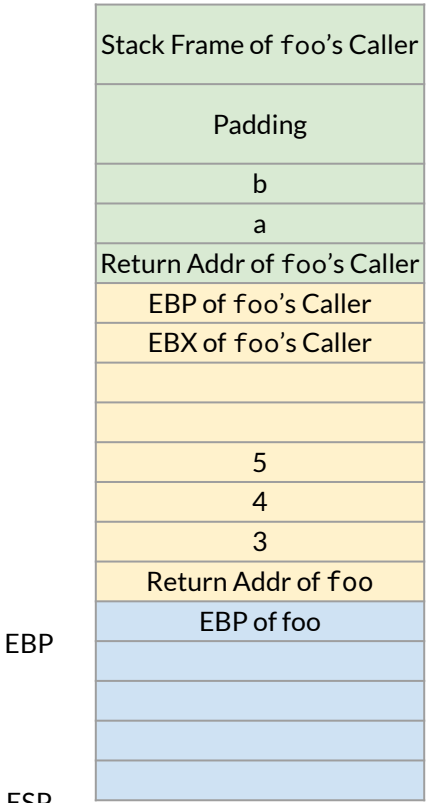
31         movl    %eax, p
32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	7
EBX	a + b
EDX	3



Stack Frame

call.s

```
1  p:
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4          pushl   %ebp
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19         pushl   %ebx
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22         movl    12(%ebp), %eax
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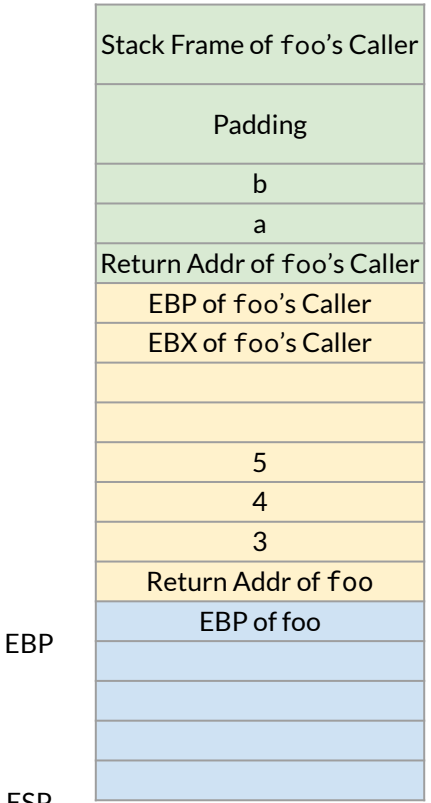
31         movl    %eax, p
32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	2
EBX	a + b
EDX	3



Stack Frame

call.s

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27         pushl   $3
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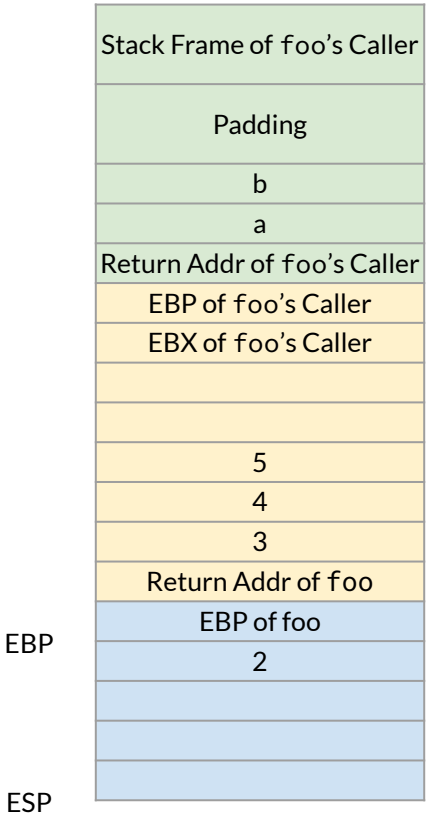

Stack Frame

call.s

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30         addl    %ebx, %eax
31         movl    %eax, p
32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```
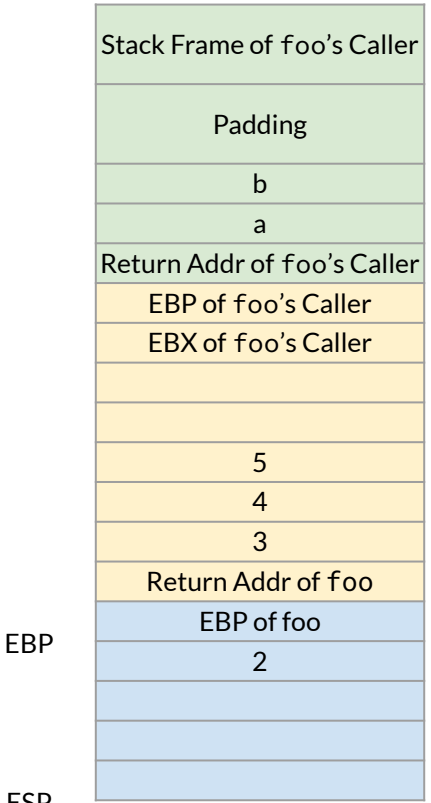

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	2
EBX	a + b
EDX	3



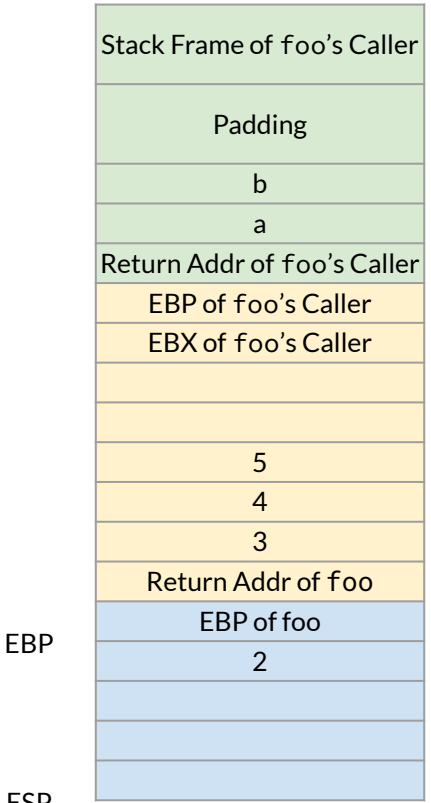
Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
6          subl    $16, %esp
7          movl    8(%ebp), %edx
8          movl    12(%ebp), %eax
9          addl    %edx, %eax
10         subl    16(%ebp), %eax
11         movl    %eax, -4(%ebp)
12         movl    -4(%ebp), %eax
13         addl    $1, %eax
14         leave
15         ret
16 foo:
17         pushl   %ebp
18         movl    %esp, %ebp
19         pushl   %ebx
20         subl    $4, %esp
21         movl    8(%ebp), %edx
22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
26         pushl   $4
27         pushl   $3
28         call    bar
29         addl    $16, %esp
30         addl    %ebx, %eax
31         movl    %eax, p
32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	3
EBX	a + b
EDX	3



Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
6          subl    $16, %esp
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8          movl    12(%ebp), %eax
9          addl    %edx, %eax
10         subl    16(%ebp), %eax
11         movl    %eax, -4(%ebp)
12         movl    -4(%ebp), %eax
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22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
26         pushl   $4
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31         movl    %eax, p
32         nop
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34         leave
35         ret
```
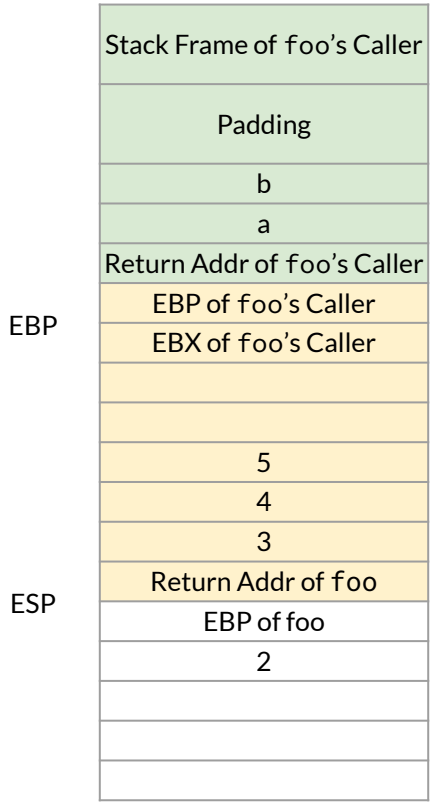

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	3
EBX	a + b
EDX	3



Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
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22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
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27         pushl   $3
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32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```
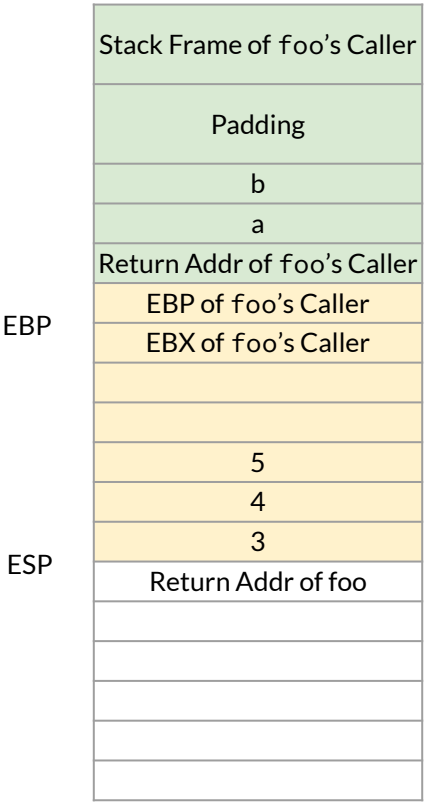

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	3
EBX	a + b
EDX	3



Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
6          subl    $16, %esp
7          movl    8(%ebp), %edx
8          movl    12(%ebp), %eax
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22         movl    12(%ebp), %eax
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24         subl    $4, %esp
25         pushl   $5
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32         nop
33         movl    -4(%ebp), %ebx
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35         ret
```
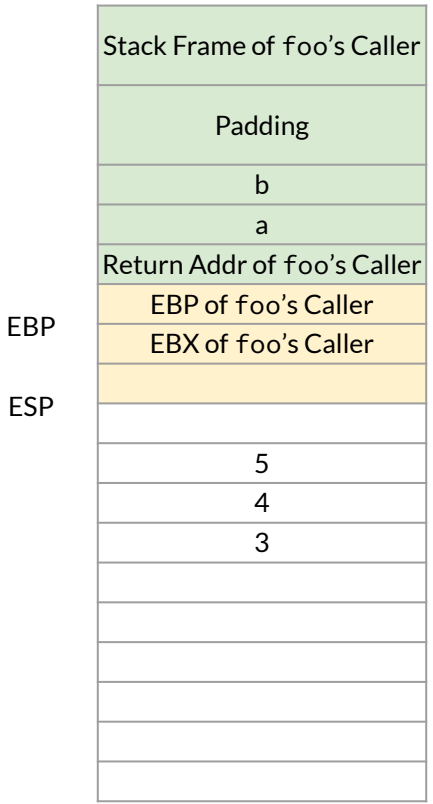

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	3
EBX	a + b
EDX	3



Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
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9          addl    %edx, %eax
10         subl    16(%ebp), %eax
11         movl    %eax, -4(%ebp)
12         movl    -4(%ebp), %eax
13         addl    $1, %eax
14         leave
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18         movl    %esp, %ebp
19         pushl   %ebx
20         subl    $4, %esp
21         movl    8(%ebp), %edx
22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
26         pushl   $4
27         pushl   $3
28         call    bar
29         addl    $16, %esp
30         addl    %ebx, %eax
31         movl    %eax, p
32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```
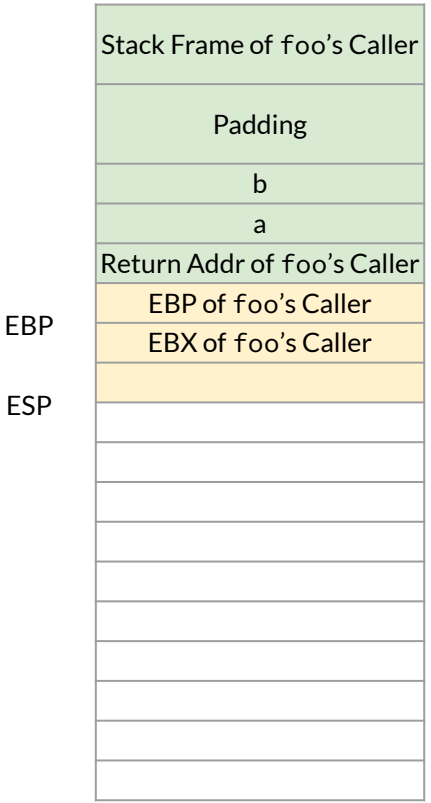

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	3 + a + b
EBX	a + b
EDX	3



Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
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21         movl    8(%ebp), %edx
22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
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25         pushl   $5
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31         movl    %eax, p
32         nop
33         movl    -4(%ebp), %ebx
34         leave
35         ret
```
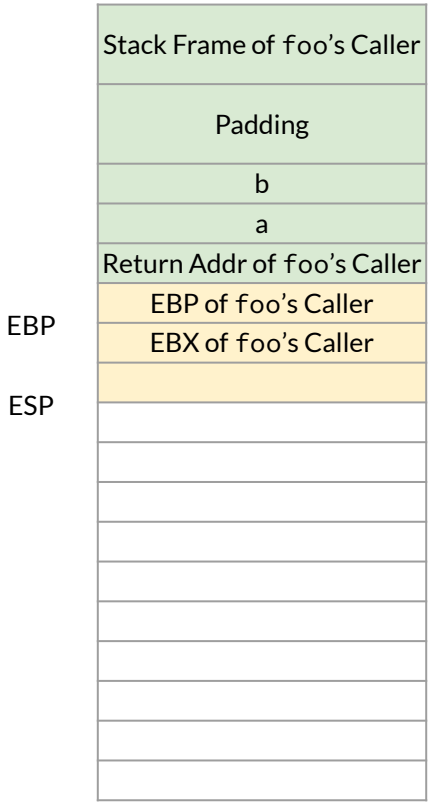

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	3 + a + b
EBX	EBX of foo's Caller
EDX	3



Stack Frame

call.s

```
1  p:
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4          pushl   %ebp
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34         leave
35         ret
```

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	3 + a + b
EBX	EBX of foo's Caller
EDX	3

EBP

ESP

Stack Frame of foo's Caller
Padding
b
a
Return Addr of foo's Caller
EBP of foo's Caller
EBX of foo's Caller

Stack Frame

call.s

```
1  p:
2          .zero   4
3  bar:
4          pushl   %ebp
5          movl    %esp, %ebp
6          subl    $16, %esp
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20         subl    $4, %esp
21         movl    8(%ebp), %edx
22         movl    12(%ebp), %eax
23         leal    (%edx,%eax), %ebx
24         subl    $4, %esp
25         pushl   $5
26         pushl   $4
27         pushl   $3
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```

4. Draw the stack frame with ESP and EBP and contents of EAX, EBX, EDX registers.

EAX	3 + a + b
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