

# Datacenters I: Anatomy and Topologies

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[cs168.io](http://cs168.io)

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Thanks to Ankit Singla and Murphy McCauley for some of the material!

# Recall – where is the Internet?

- Carrier hotel locations.
- Generally for interconnection between networks.
- Some smaller application hosting.
- Where do large applications live?



# A Datacenter



Google datacenter in Belgium - <https://www.google.com/about/datacenters/gallery/>

# Inside a (Google) Datacenter



Server racks in a Google datacenter - <https://www.google.com/about/datacenters/gallery/>

# Infrastructure in a Google Datacenter

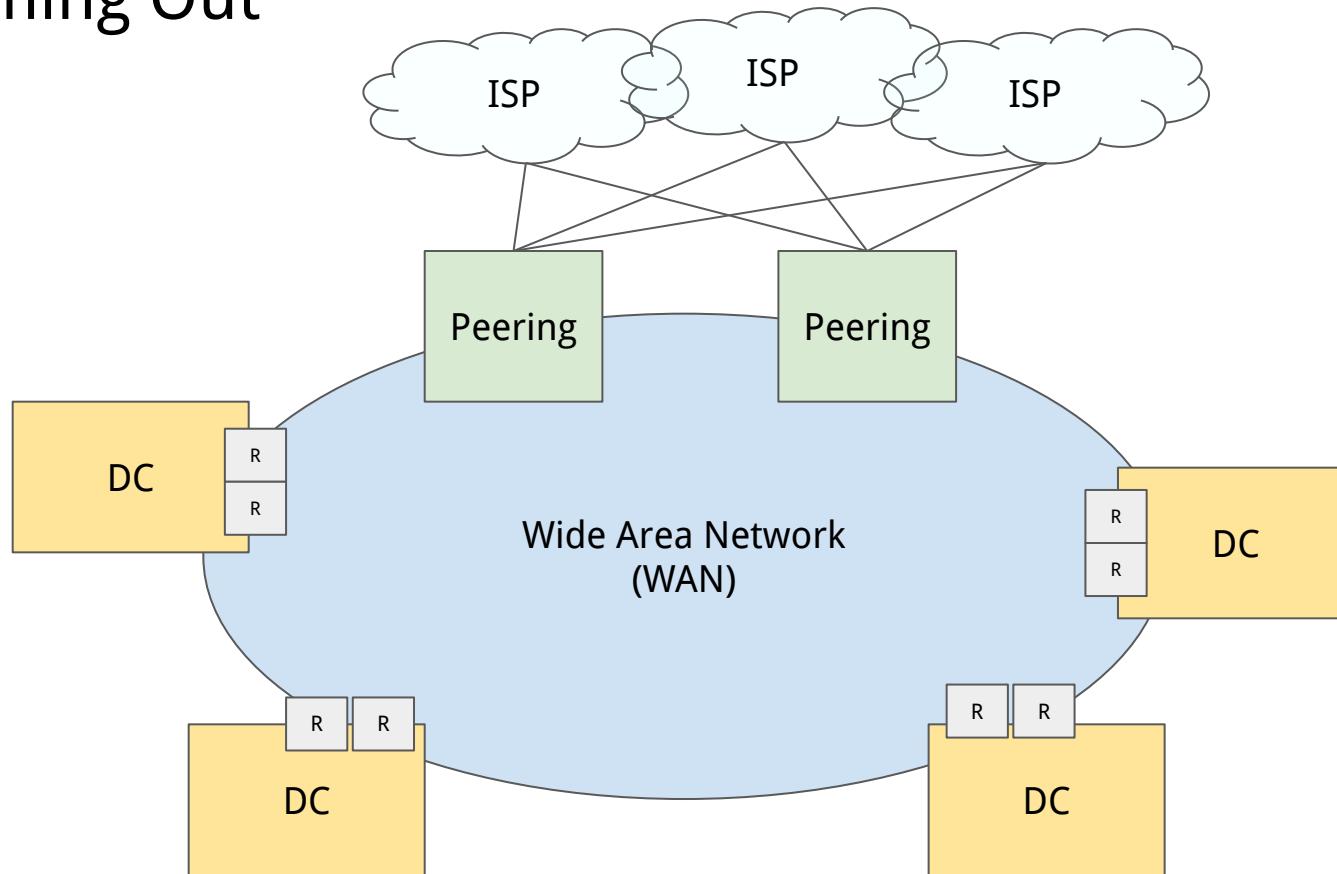


Cooling infrastructure in a Google datacenter - <https://www.google.com/about/datacenters/gallery/>

# Datacenters

- Computing infrastructure, located in one physical location.
- Owned by one organisation.
- But used by multiple users and applications.
- Our focus: modern hyperscale datacenters.
  - Google, Facebook, Microsoft, Meta...
  - Concept scales down.

# Zooming Out



# Anatomy of an Application/Cloud Provider

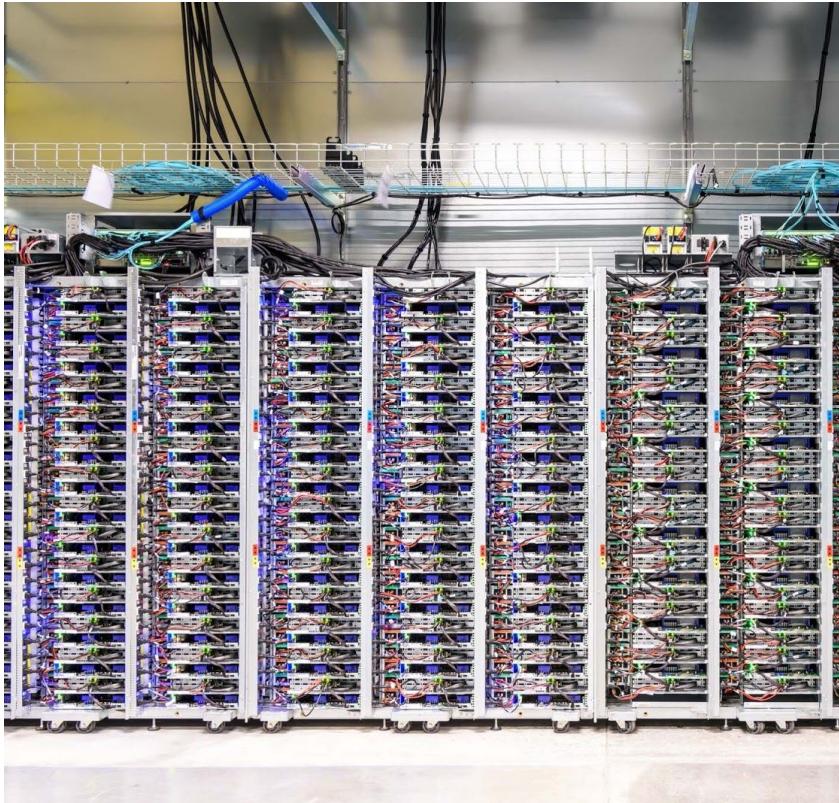
- Data center locations – host servers and application infrastructure.
  - Often huge power requirements.
  - Does not need to be near other networks.
- Peering locations – host network interconnection infrastructure.
  - Typically mostly routers.
  - Needs to be near other networks.
- **Wide Area Network** - connects the different locations together.
- Datacenter network – within a particular DC facility.

# Our focuses

- What does a datacenter network look like?
- What makes a datacenter different to the wide area networks we have discussed thus far?
- Components of a datacenter network.
- Specific solutions for datacenter networking [next time].
  - Routing in datacenters.

Questions?

# Anatomy of a Datacenter

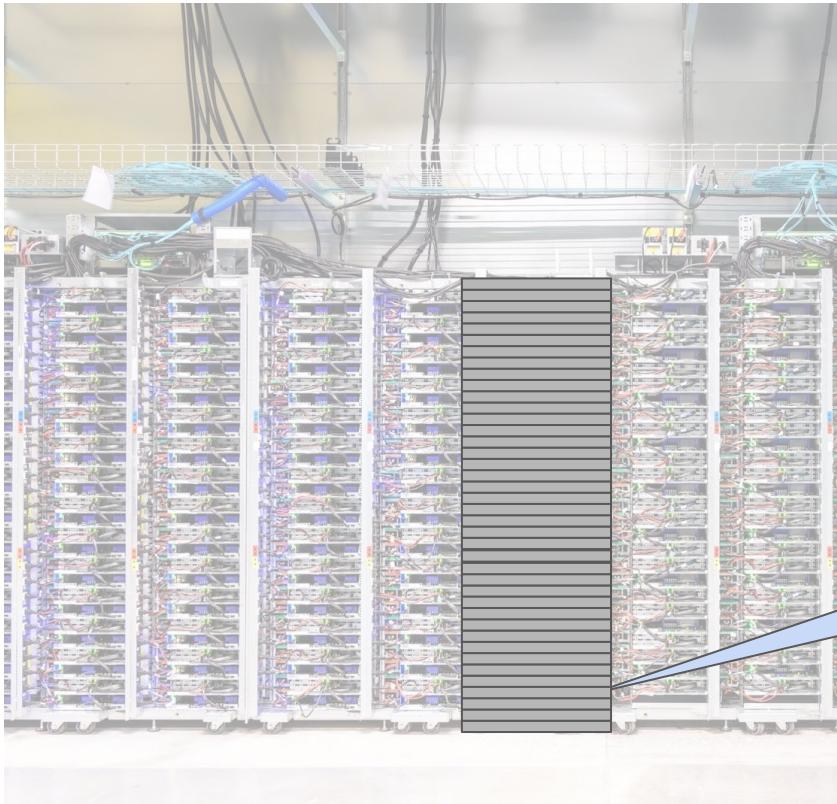


# Anatomy of a Datacenter



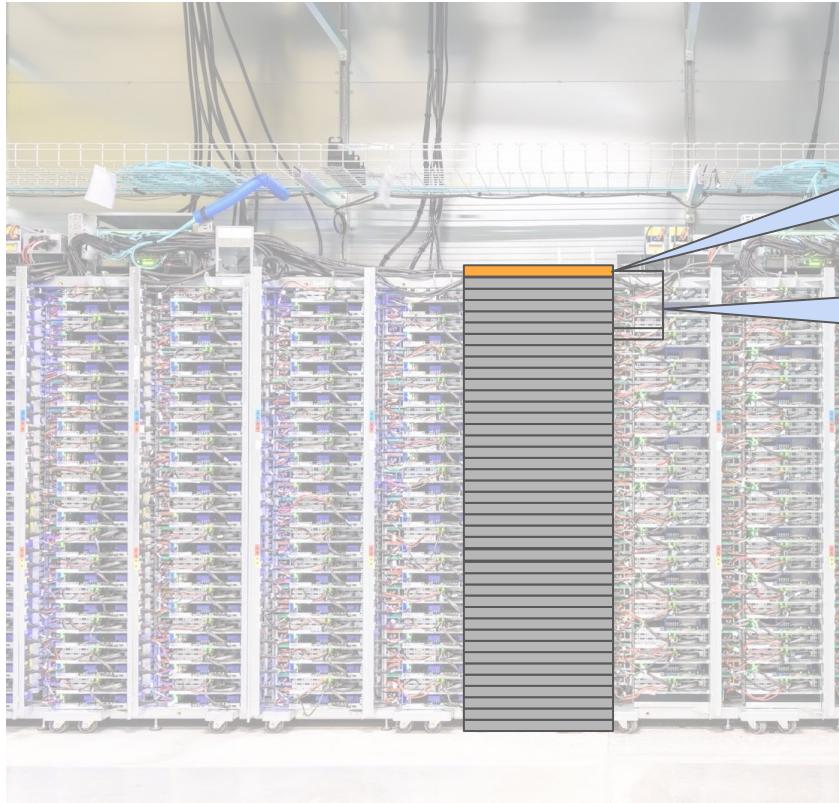
1-2 servers per "U" [0]

# Anatomy of a Datacenter



≈40 "U" per rack.

# Anatomy of a Datacenter



Top of Rack (TOR)  
switch

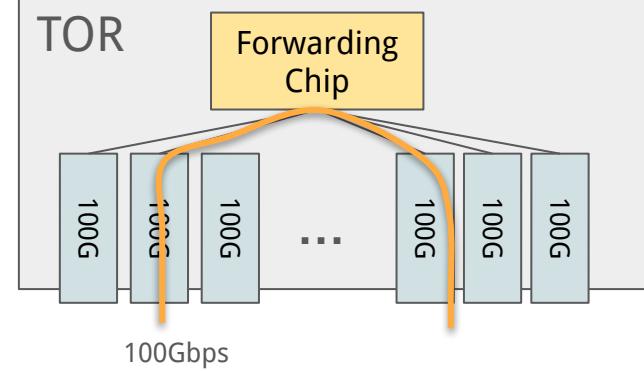
Server “access links” or “uplinks”

# Anatomy of a Datacenter

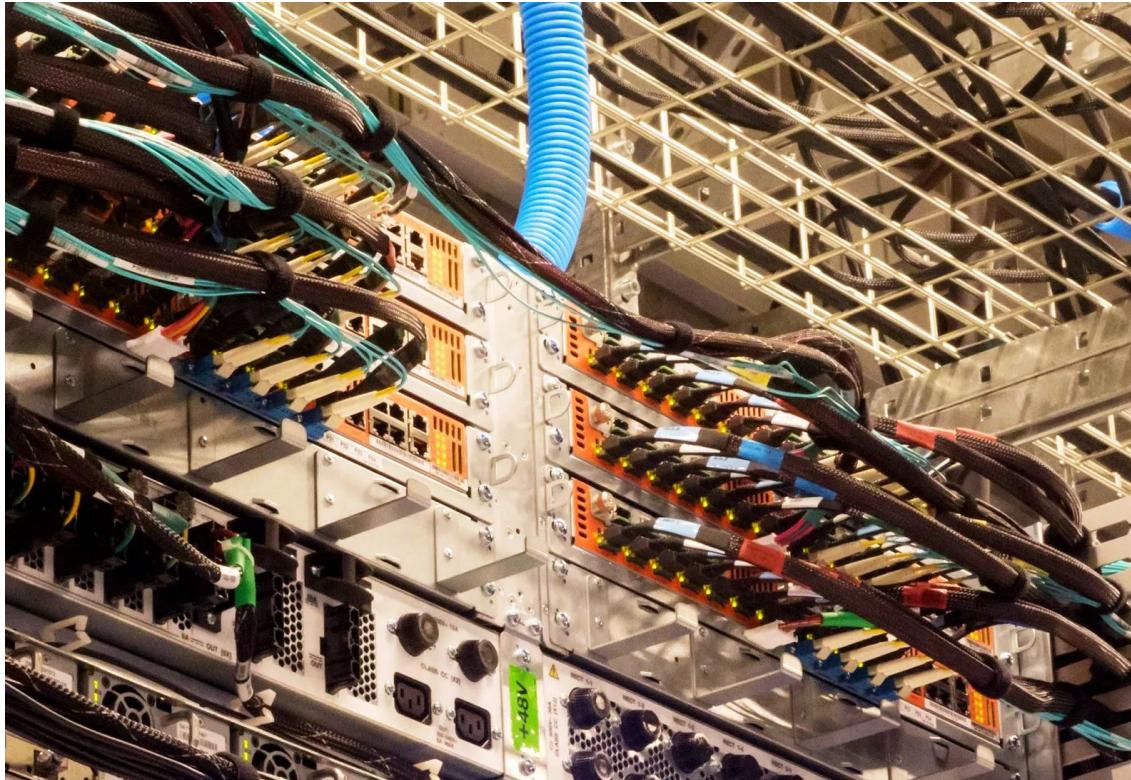


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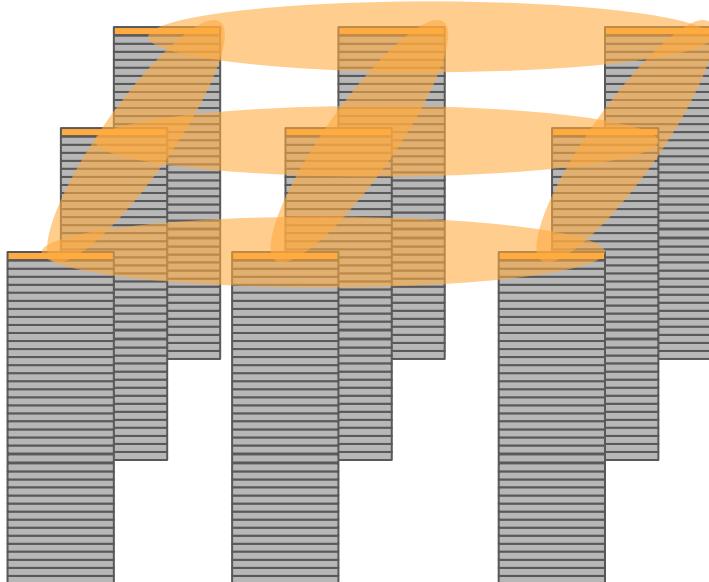
Server “access links” or “uplinks”



# Top-of-Rack Switch



# Anatomy of a Datacenter



- 40-80 servers per rack.
- 100Gbps per server.
- Many racks per datacenter!
- How do we connect racks together?

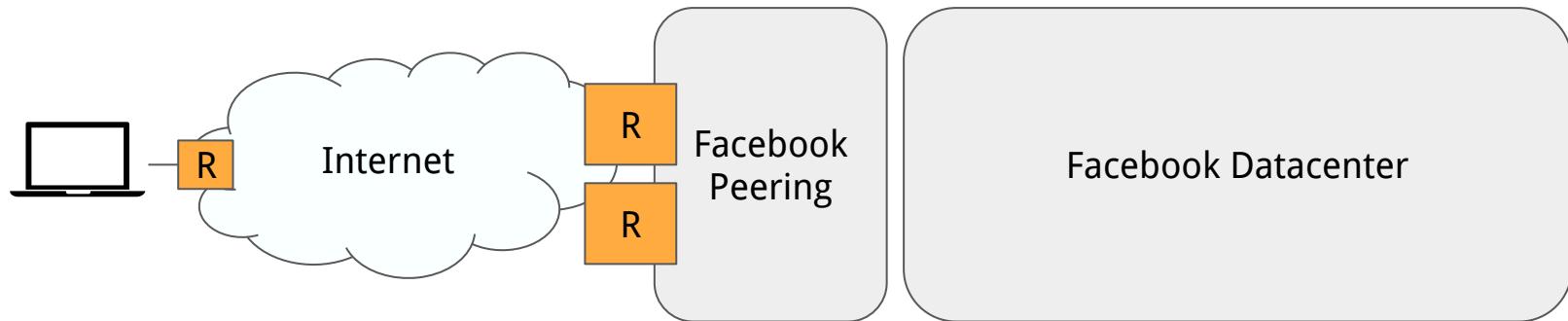
# Why is the datacenter different?

- We have generally been thinking about Wide Area Networks.
- These WANs interconnect to make up the Internet.
- Why might datacenter networks be different?

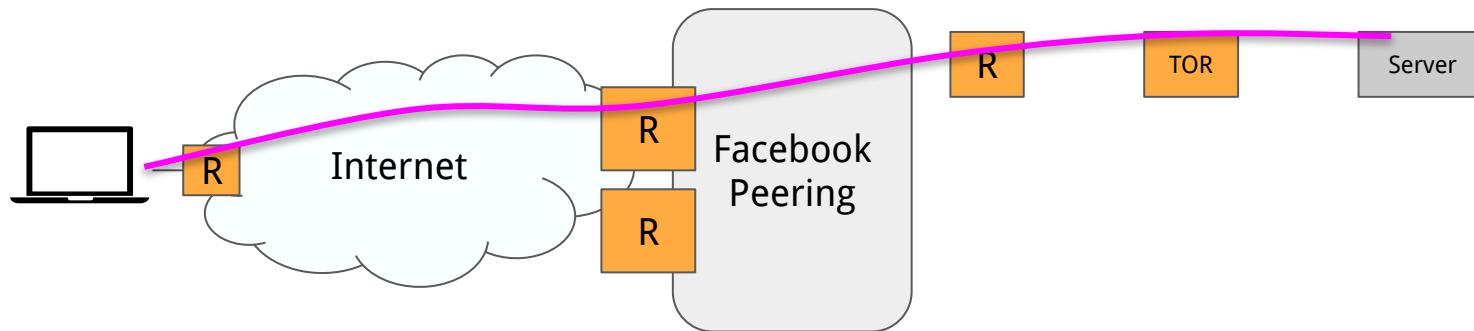
# Why is the datacenter different?

- We have generally been thinking about Wide Area Networks.
- These WANs interconnect to make up the Internet.
- Why might datacenter networks be different?
  - Run by a single organisation
  - Exist in a single physical location
  - High scale (in that single location!)
  - More control over network and hosts (to some degree)
  - Homogeneous
  - Performance, performance, performance!

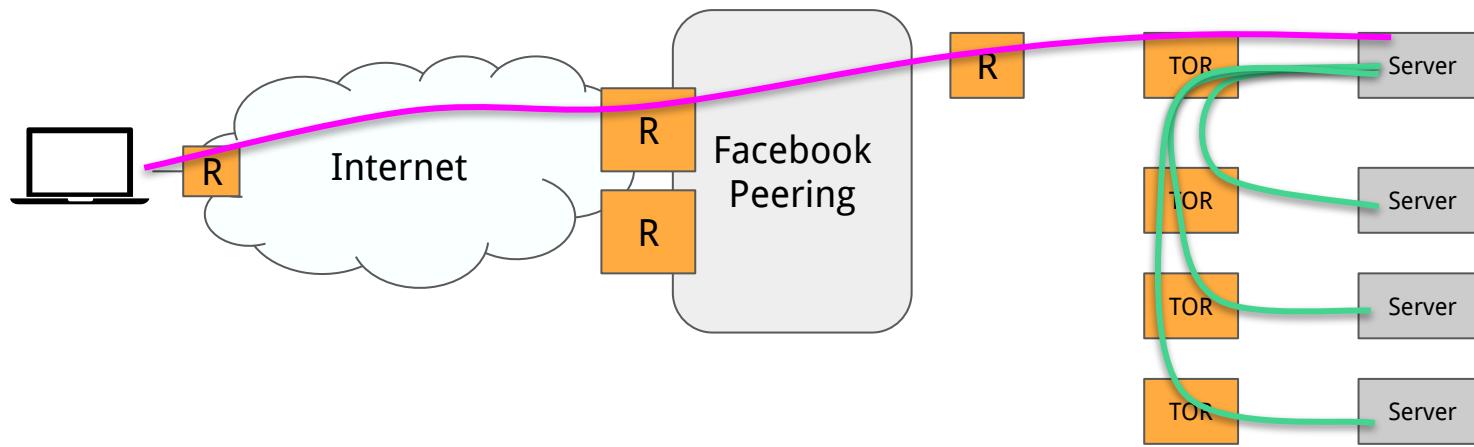
# Accessing an Application



# Accessing an Application



# Accessing an Application



# Accessing an Application

USENIX NSDI, 2013

## Scaling Memcache at Facebook

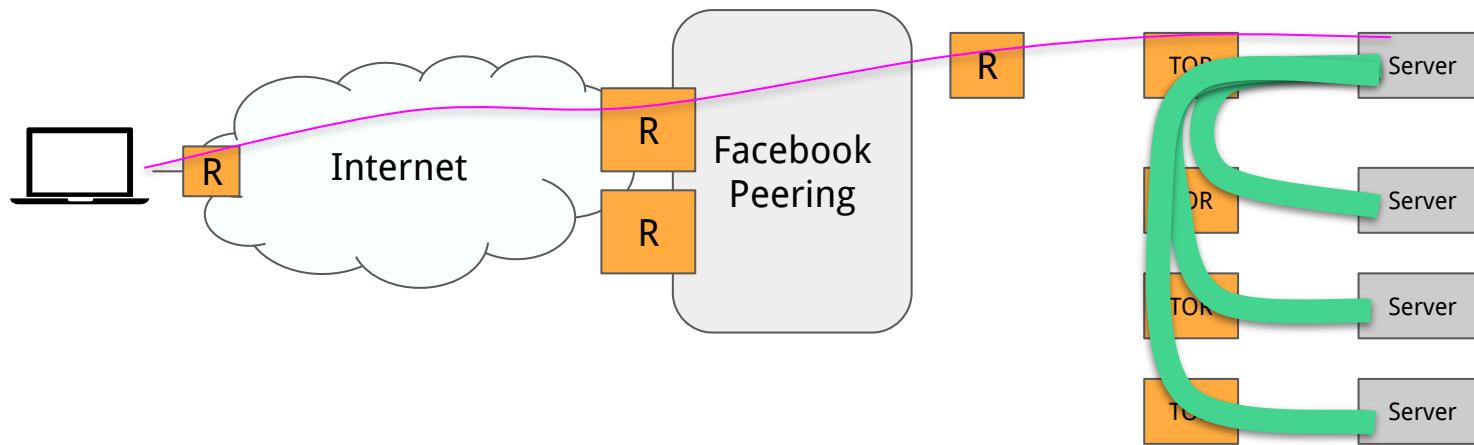
Rajesh Nishtala, Hans Fugal, Steven Grimm, Marc Kwiatkowski, Herman Lee, Harry C. Li,  
Ryan McElroy, Mike Paleczny, Daniel Peek, Paul Saab, David Stafford, Tony Tung,  
Venkateshwaran Venkataramani

{rajeshn,hans}@fb.com, {sgrimm, marc}@facebook.com, {herman, hcli, rm, mpal, dpeek, ps, dstaff, ttung, veeve}@fb.com

Facebook Inc.

1 popular page loaded = **521** distinct memcache loads  
(95th percentile = 1740!)

# Accessing an Application



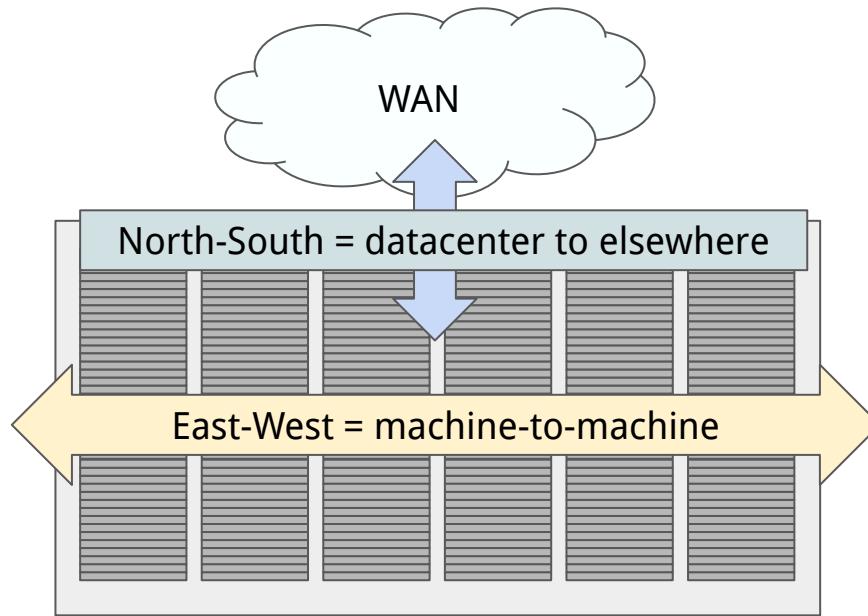
Significantly more inter-machine traffic than “user” to “machine”.

# Other Applications

- Big data analytics
  - e.g., mapreduce
- Significantly more traffic between machines - maybe *no* user-facing traffic.
- We'll come back to some different application communication patterns.

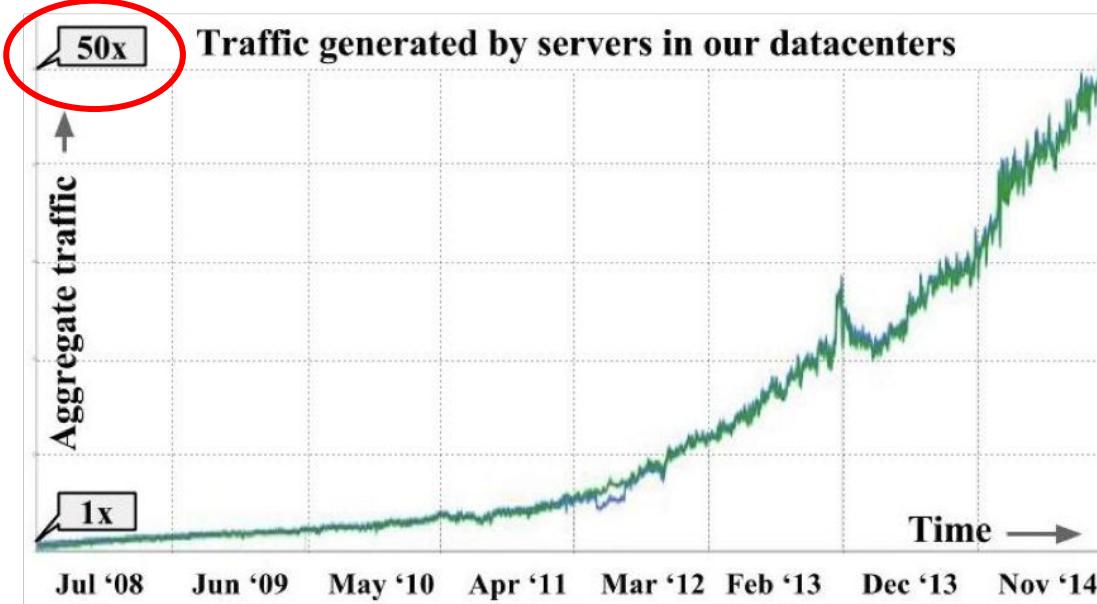


# Datacenter Traffic Patterns



East-West traffic is several orders of magnitude larger than North-South.

# East-West Traffic Volume



"Jupiter Rising: A Decade of Clos Topologies and Centralized Control in Google's Datacenter Network", Arjun Singh et al. @ Google, ACM SIGCOMM'15

Questions?

# How do we support East-West bandwidth?

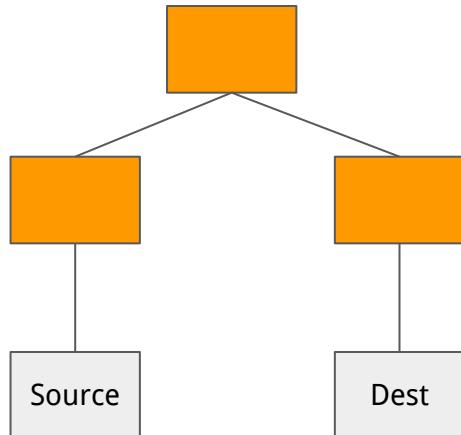
- Ideally any server can talk to any server at line rate.
- We want a network with high **bisection bandwidth**.

# Bisection Bandwidth

- Pick the number of links we must cut in order to partition a network into two halves.
- Bisection bandwidth is the sum of those bandwidths.

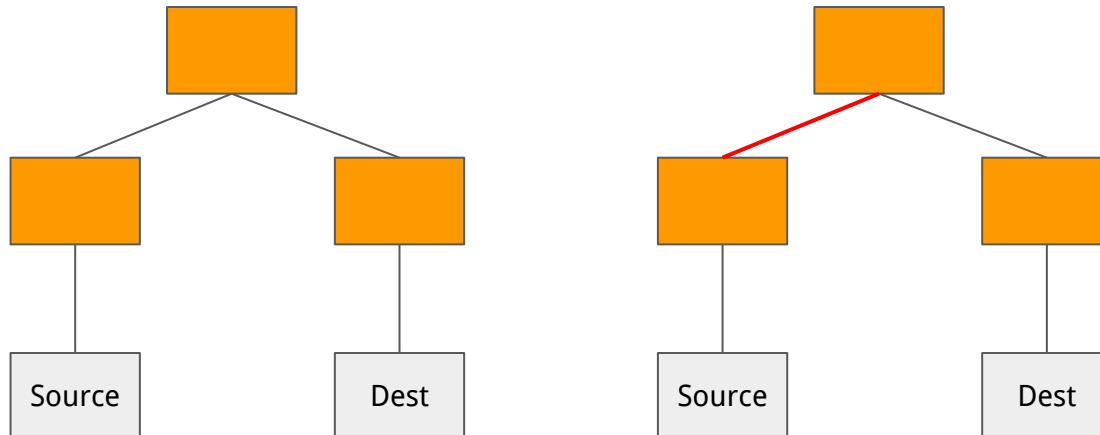
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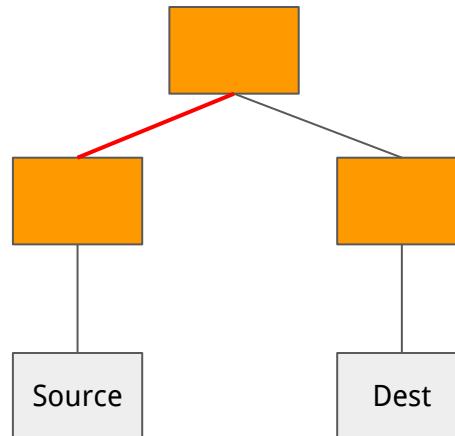
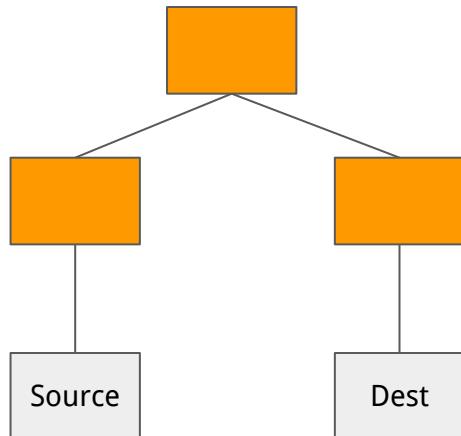
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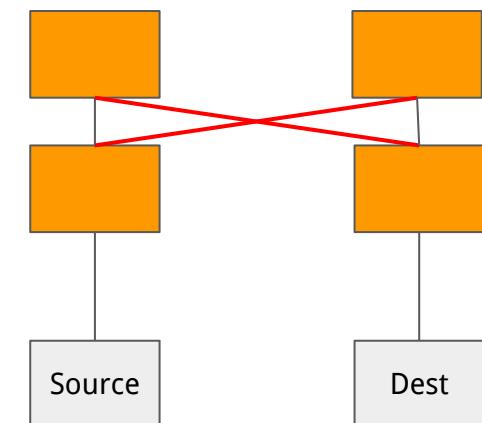
100G

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100G

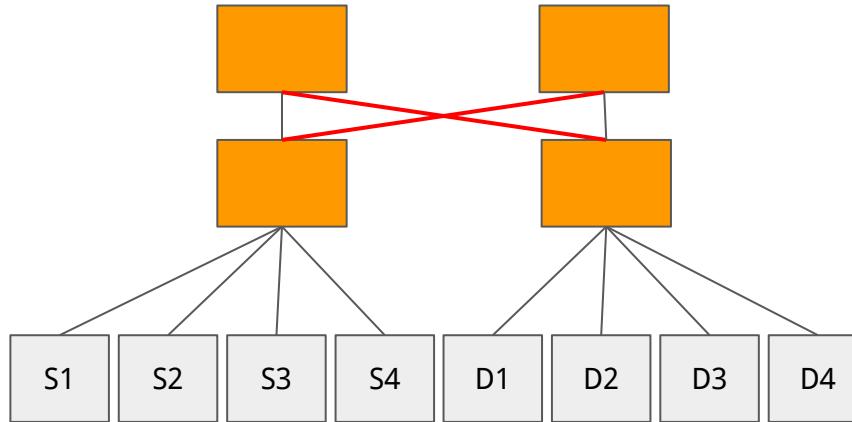


200G

# Bisection Bandwidth

- Pick the number of links we must cut in order to partition a network into two halves.
- Bisection bandwidth is the sum of those bandwidths.
- **Full** bisection bandwidth: Nodes in one partition can communicate simultaneously with nodes in the other partition at full rate.
  - Given  $N$  nodes, each with access link capacity  $R$ , bisection bandwidth =  $N/2 \times R$
- Oversubscription, informally, how far from the full bisection bandwidth we are.
  - Formally: ratio of worst-case achievable bandwidth to full bisection bandwidth.

# Bisection Bandwidth



Bisection Bandwidth: 200G

Full Bisection Bandwidth:  $(8/2) * 100G = 400G$

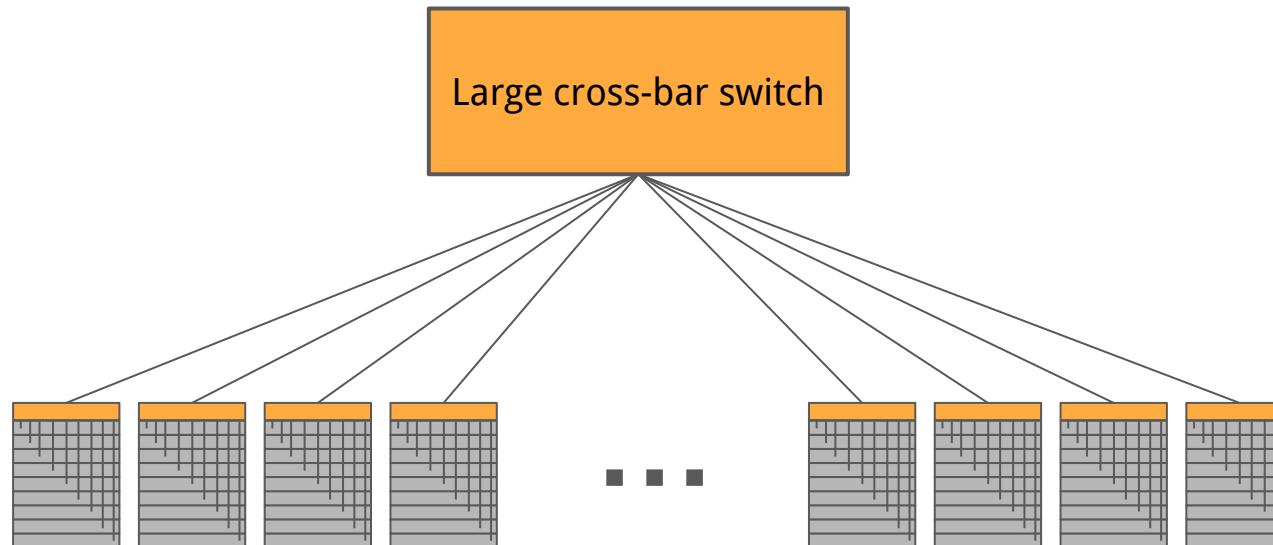
Oversubscription:  $200/400 = 0.5$  (2x)

Questions?

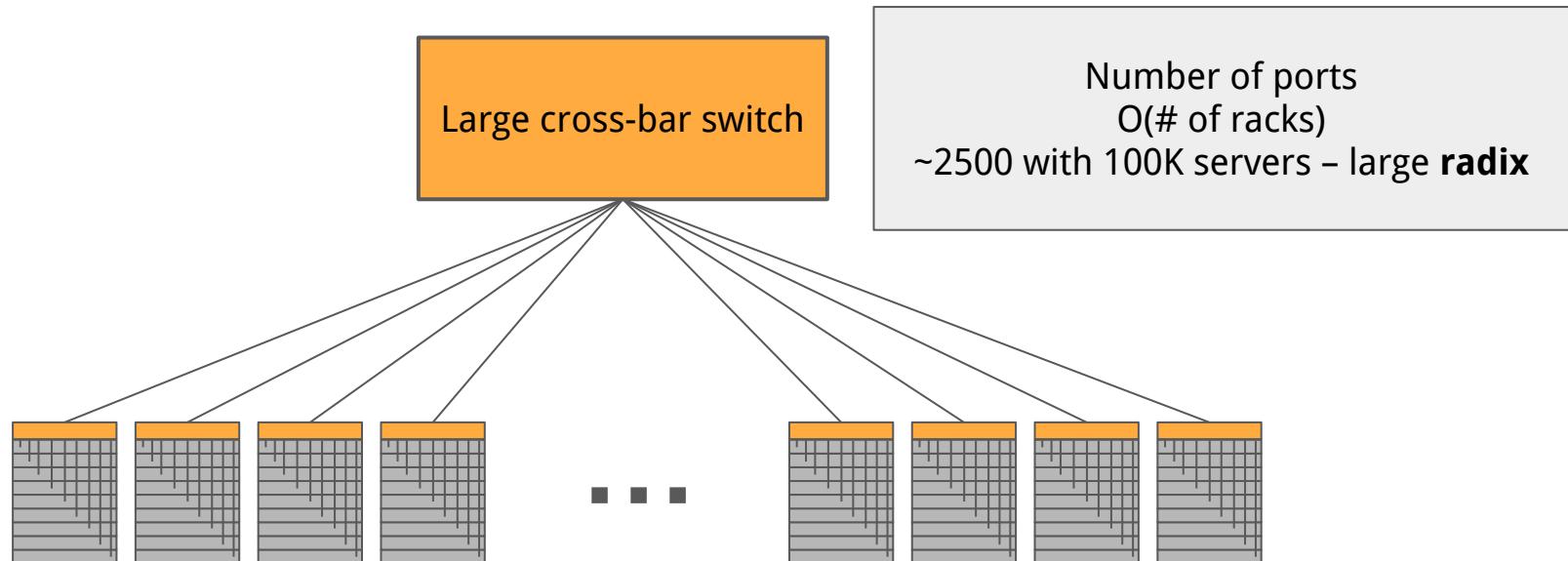
# Maximising Bisection Bandwidth

- As we've seen, bisection bandwidth is a function of the topology of the network.
- In the datacenter we can choose our topology relatively easily.
  - Run more cables (fibre, electrical)
- What topology do we build?

# “Big Switch” Approach for DC Networking



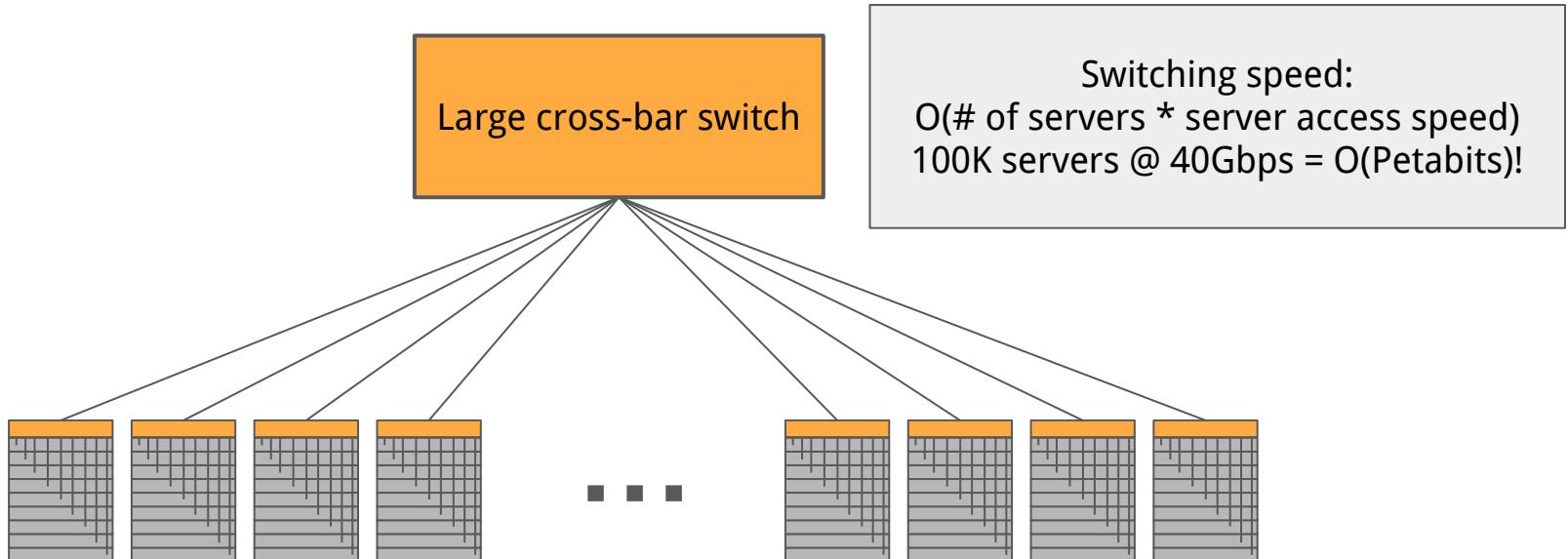
# “Big Switch” Approach for DC Networking



# Switch Radix

- Radix is used to describe the maximum degree a switch can have to other nodes in the network.
- Simply can be thought of as the number of ports the switch can have.
- Switches have constrained radix:
  - Number of line cards they can support.
  - Physical constraints of building fabric cards (remember router architecture).
  - Physical size!

# “Big Switch” Approach for DC Networking



**Does not scale** (and if it did, would be \$\$\$)

# We tried to do this!

10K Gig-E Switch

## 10K Gigabit Ethernet Switch

### Request for Proposal

Google Part 900190  
version 1.0

But what we needed was a 10,000-port switch that cost \$100/port. So, almost exactly 20 years ago, we sent this five-page RFP to four different switch vendors (IIRC: Cisco, Force10, HP, and Quanta) and tried to interest them in building such a switch. They politely declined because “nobody is asking for such a product except for you”, and they anticipated margins to be low.

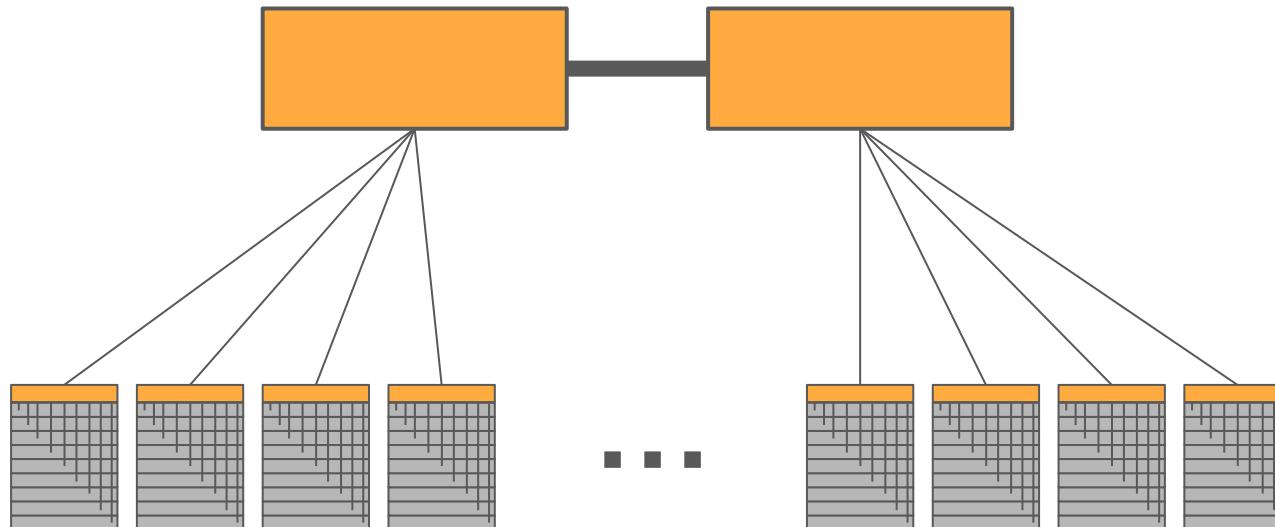
## 6 Implementation Ideas

This section attempts to explain why we believe it is possible to build a 10,000 port non-blocking switch for \$100/port. This section does not imply any requirements for a specific design and should be thought of as one or more potential paths to a solution.

The [Broadcom BCM5670 and BCM5671](#) provide what appear to be ideal solutions for our problem. The BCM5670 can be configured with 2 BCM5690 chips to create a 20 Gig-E port switch with 2 10G uplinks and 2 10G cross links to another box with 20 Gig-E ports. The BCM5670 can be configured with 4 BCM5690 chips to create a 40 Gig-E port switch with 4 10G uplinks. The total chip cost is \$1338. Then we just need to aggregate the uplinks into a non-blocking mesh. We can build such a mesh by configuring 1125 BCM5670s into a CLOS network at a chip cost of \$416250. This leads to a total chip cost of \$75/port. This leaves \$25/port for the phy layers, uplinks, circuit boards, power supplies, CPU or configuration system, etc.

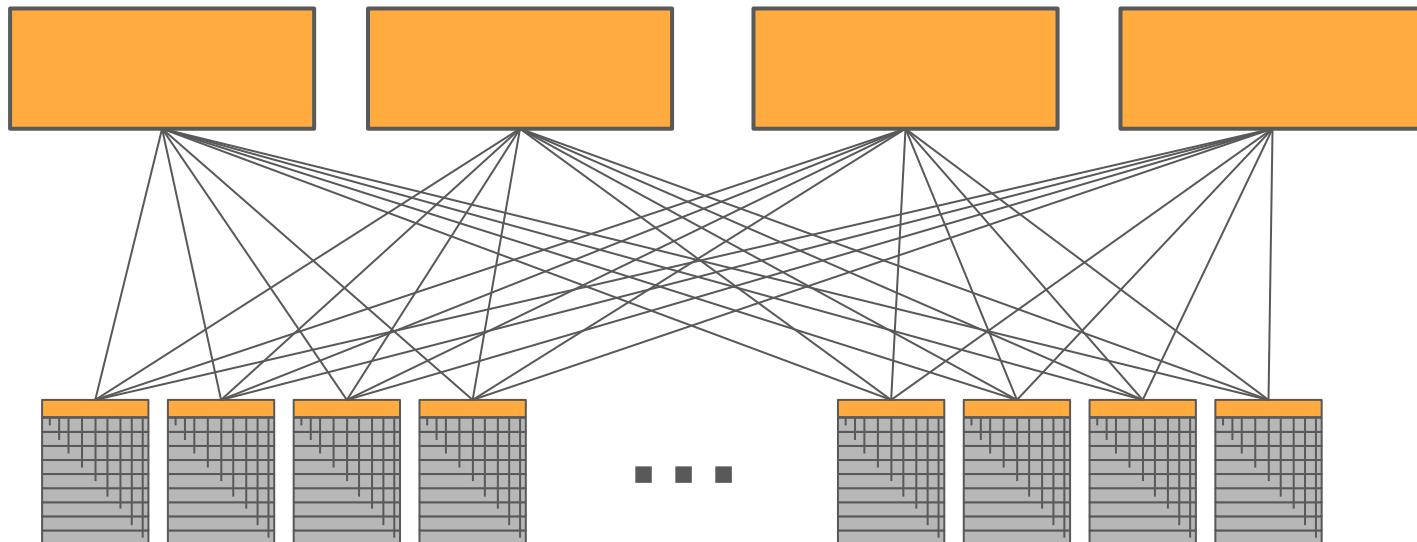
Urs Hözle (Google) on [LinkedIn](#)

# Avoiding a “Big Switch”



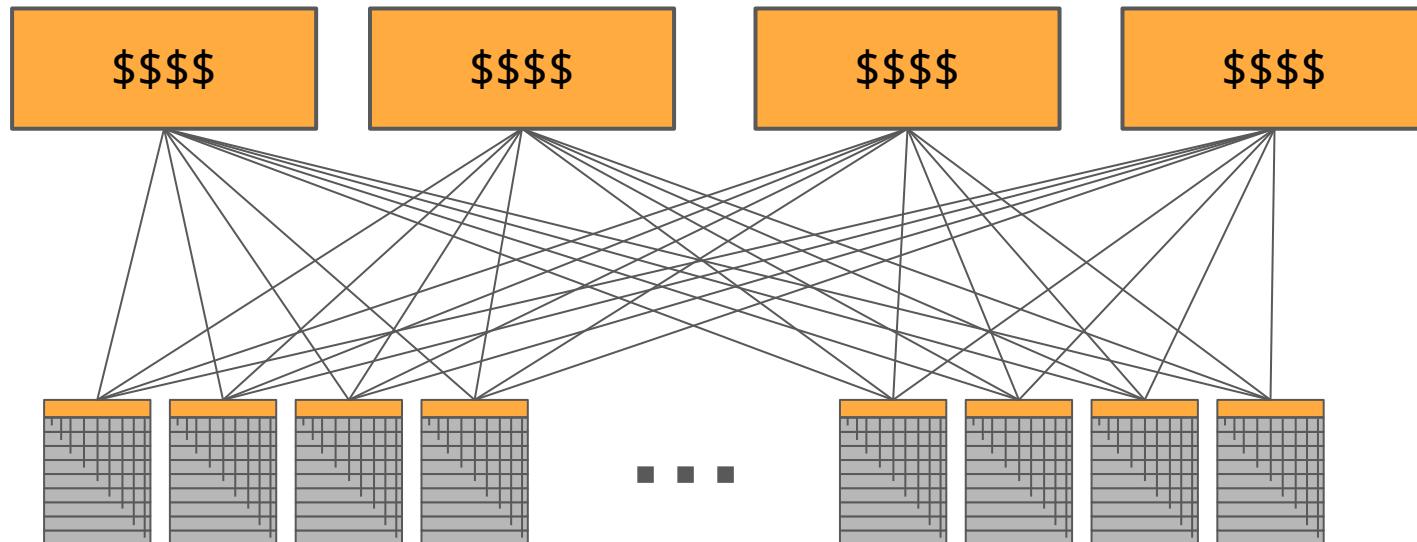
Reduced radix and bandwidth *if we don't care about failures*

# Avoiding a “Big Switch”



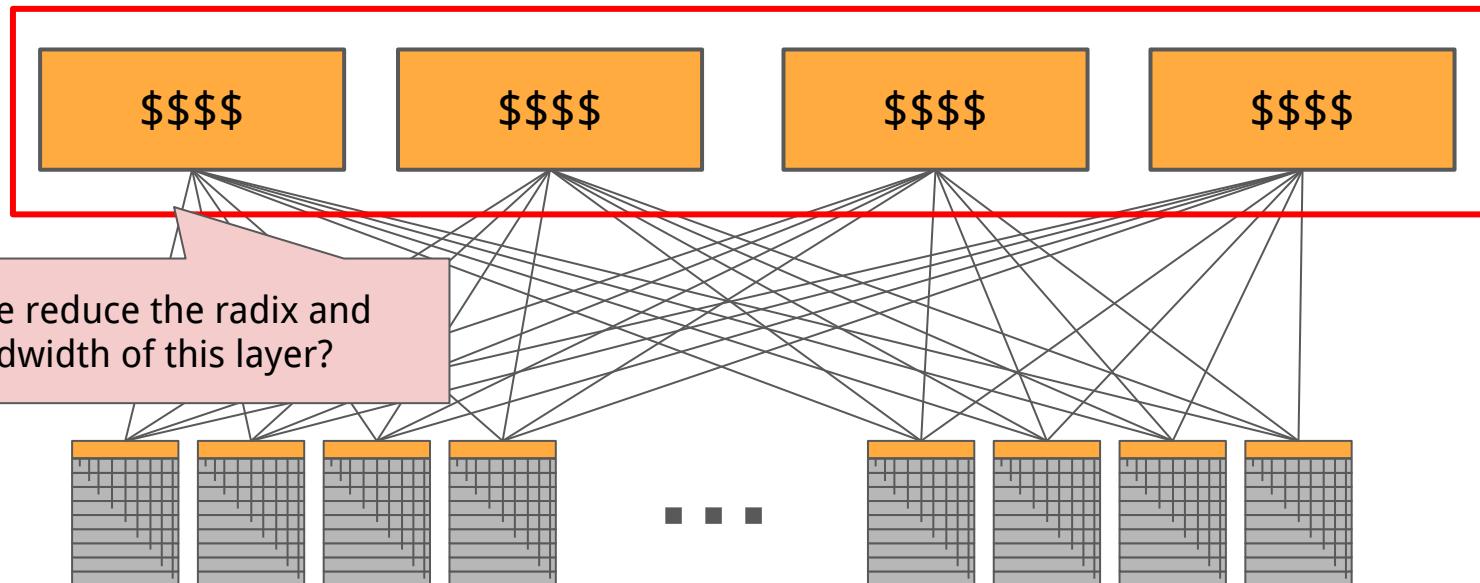
Reduced radix and bandwidth per switch - *if we can use multiple paths*

# Building a DC network



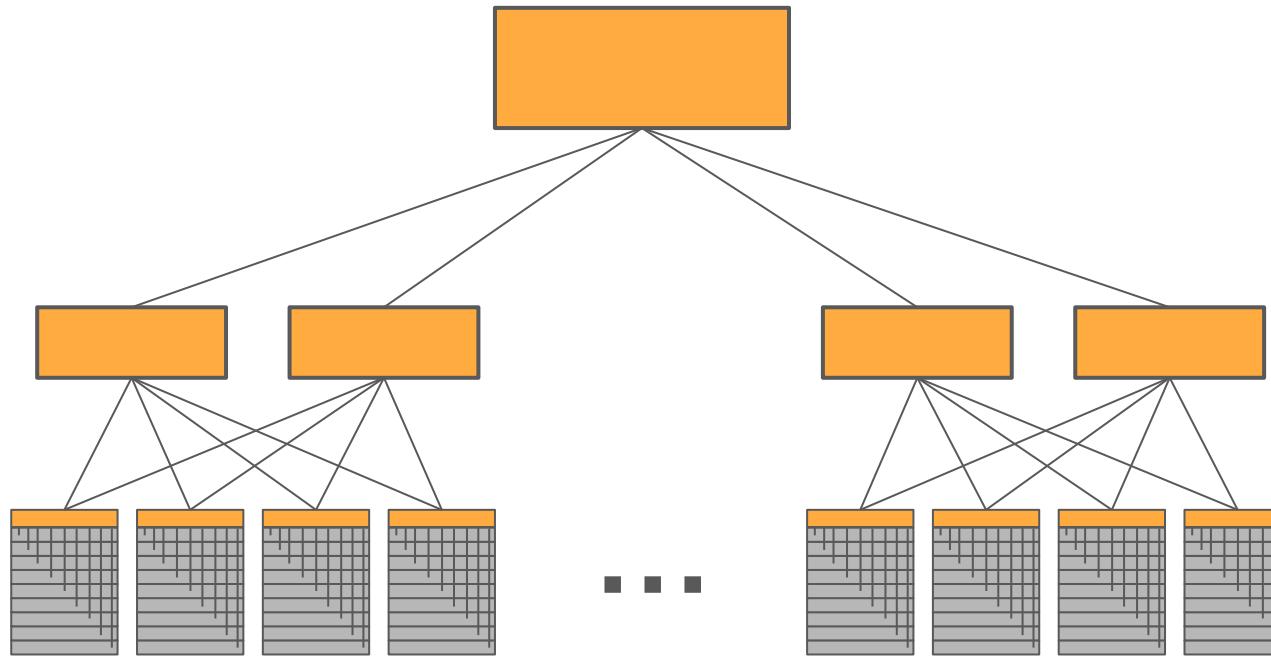
This topology works (and has been used).

# Building a DC network

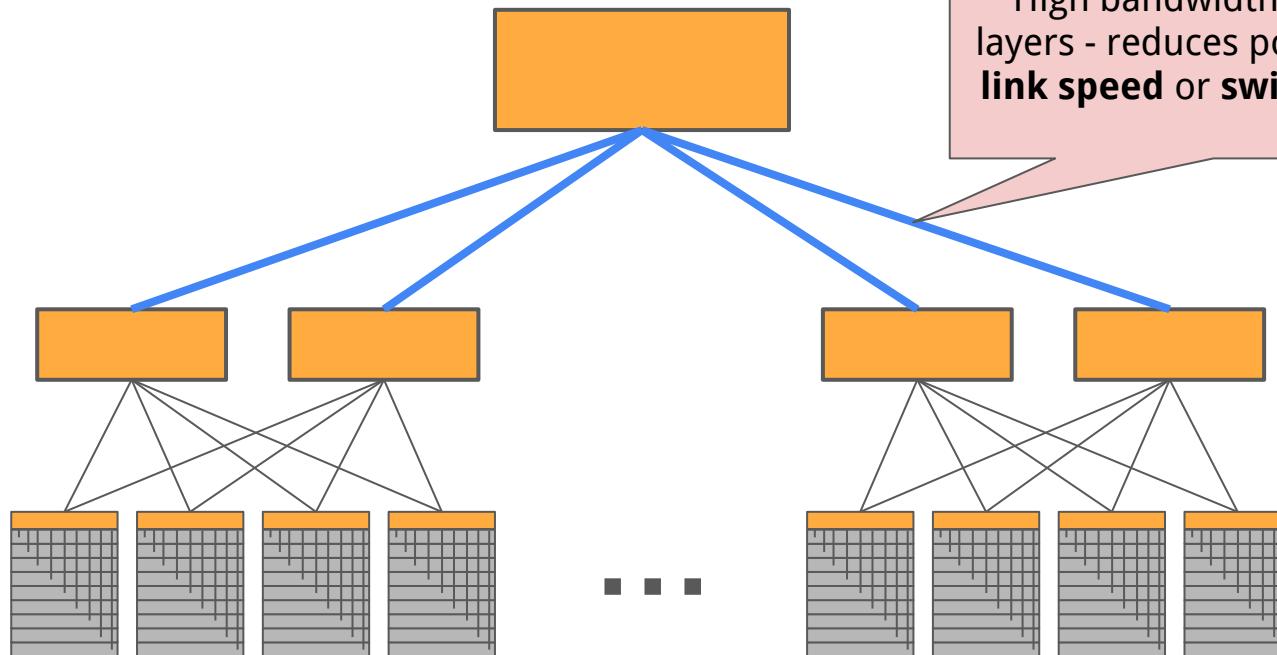


This topology works (and has been used).

# A Tree

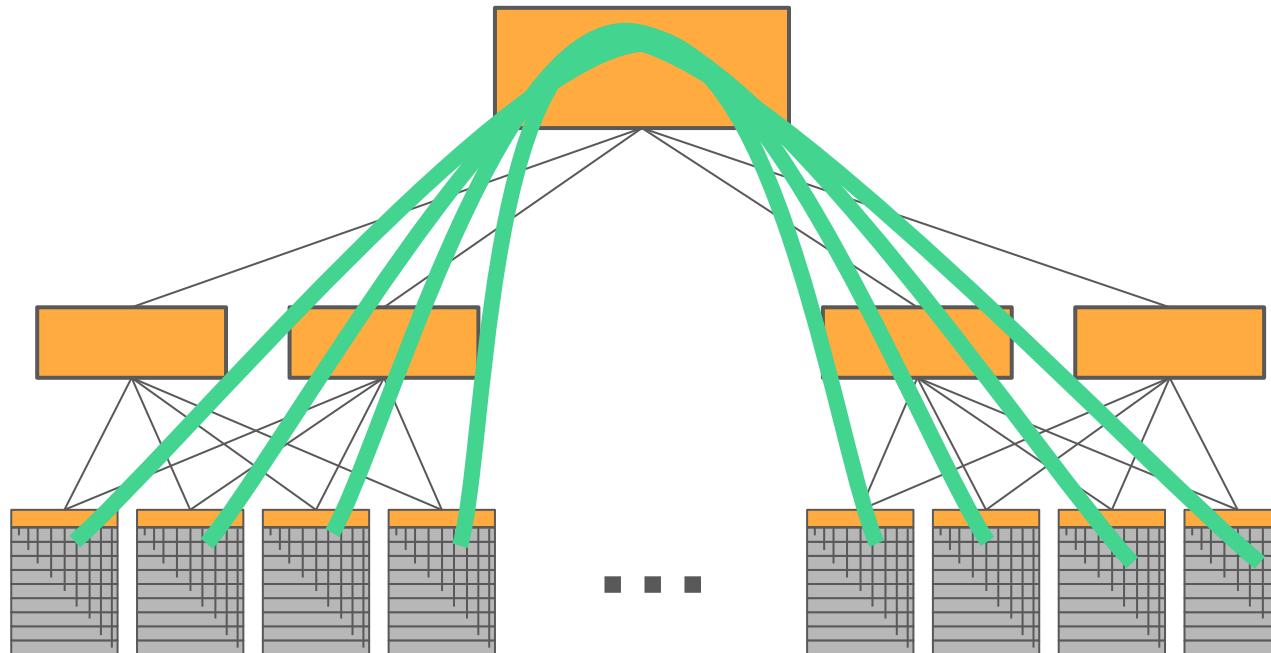


# A Fat Tree



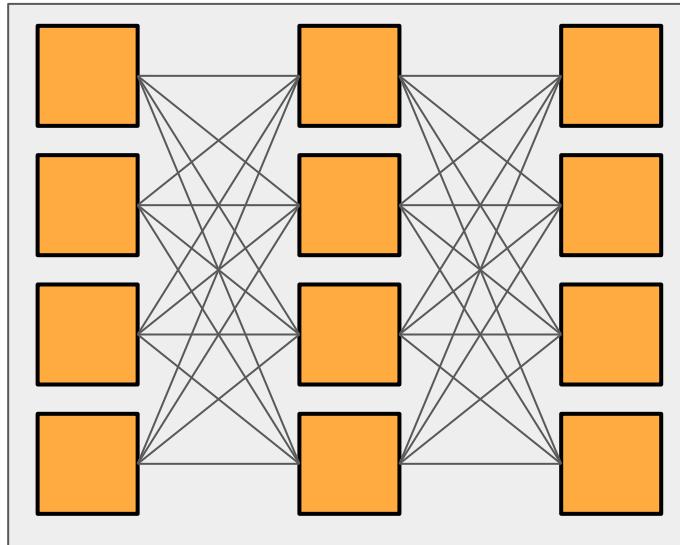
Still not scalable – or very expensive

# A Tree



Problem: low bisection bandwidth → **congestion**

# Clos Networks



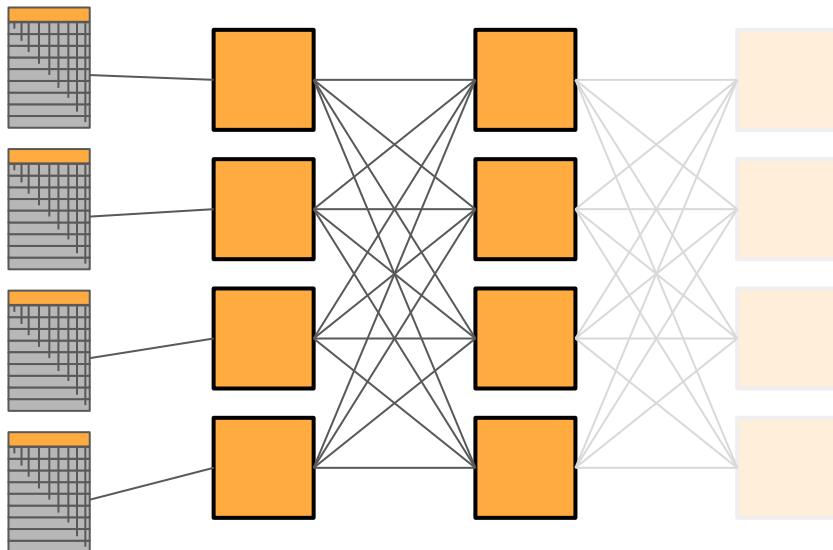
- All switches have same # of ports.
- # of ports per switch is low.
- All link speeds are the same.
- Highly multi-path.

Using small (commodity, cheap!) elements to build large capacity-rich networks.

# Clos Networks

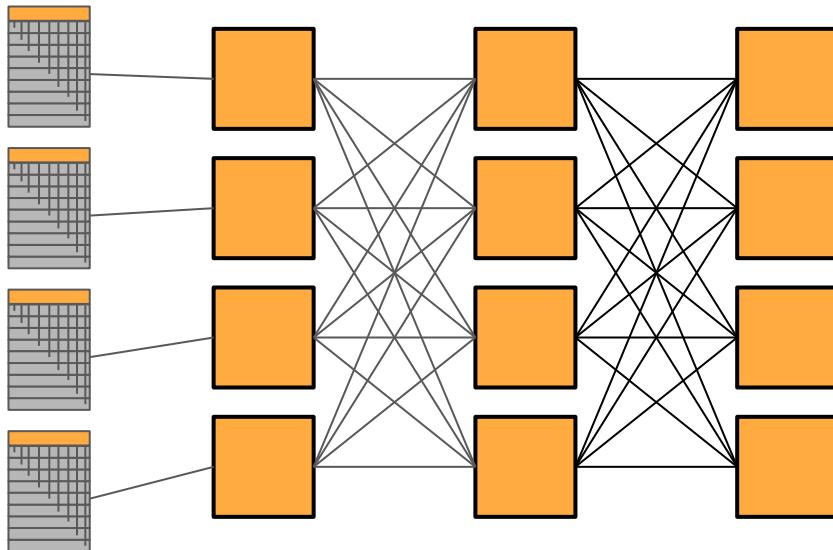
- Not a new idea!
- Formalised by Charles Clos in 1952.
- Networks can be scaled by adding *stages*.

# Clos Networks



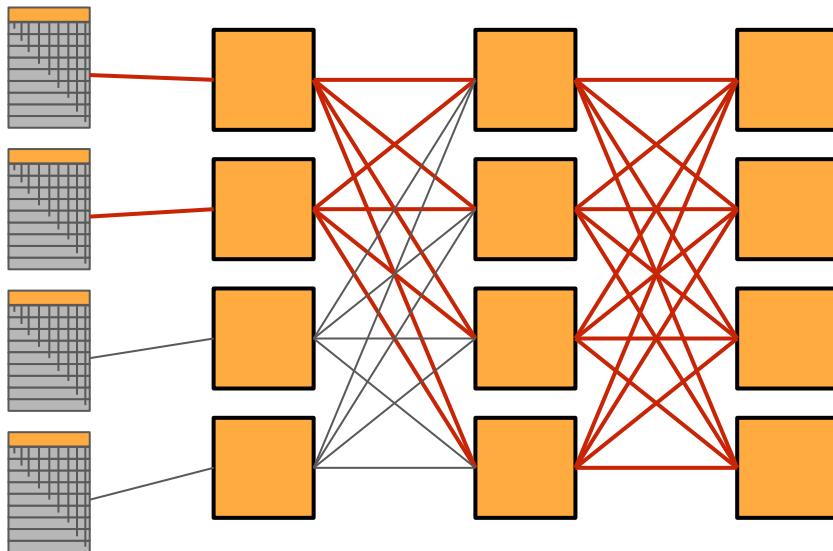
- DC networks tend to be *folded Clos*.
- Input and output switches are the same.
  - Network links are bidirectional

# Clos Networks



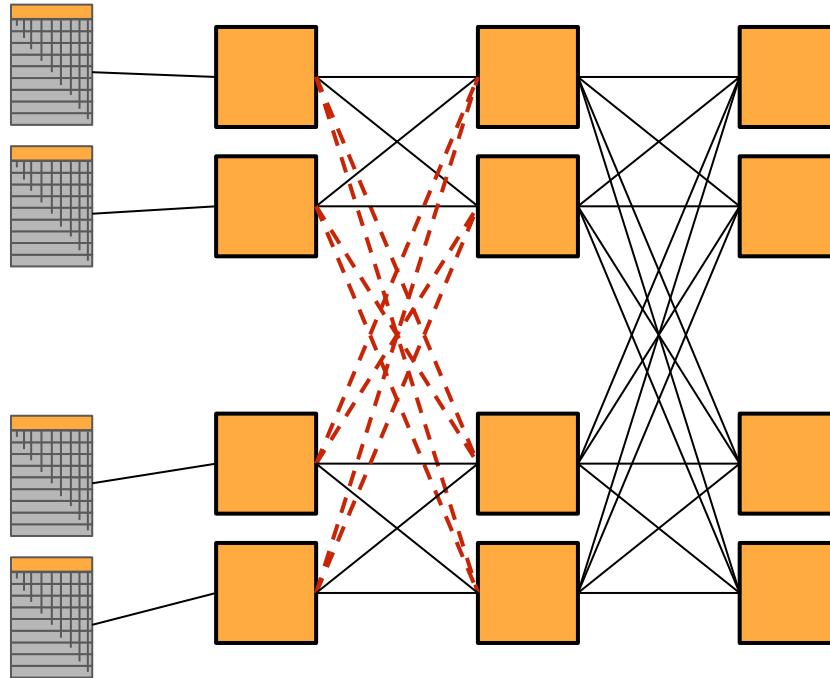
- DC networks tend to be *multi-stage*.
- Allows scaling beyond the radix of the commodity switch platforms being used.

# Clos Networks

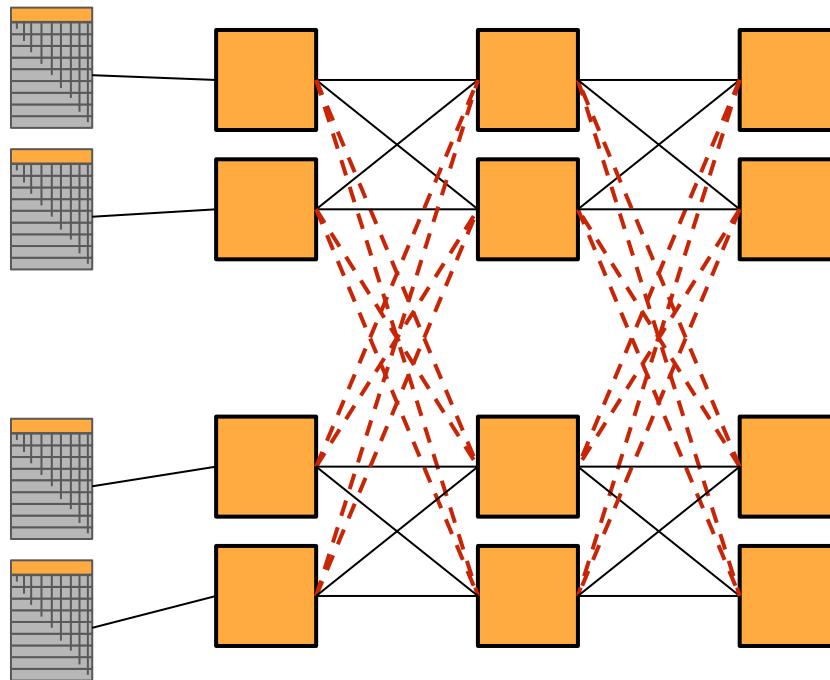


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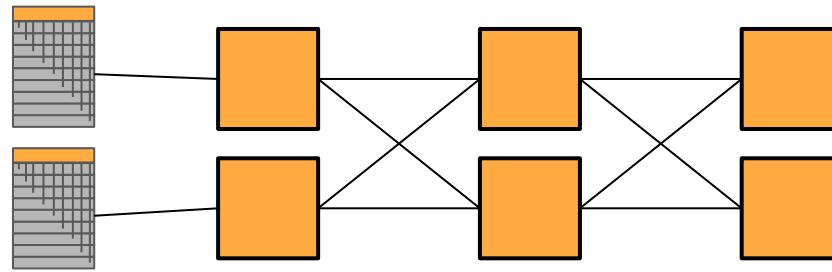
# Clos Networks - Bisection Bandwidth



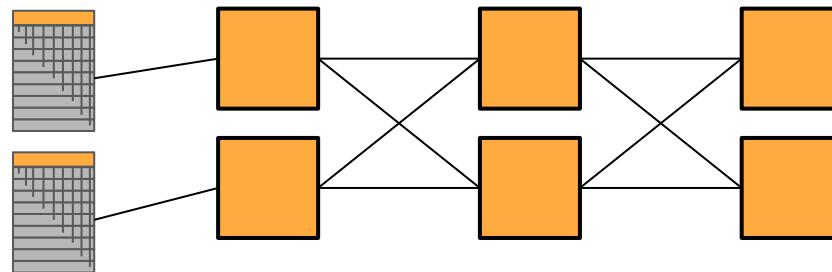
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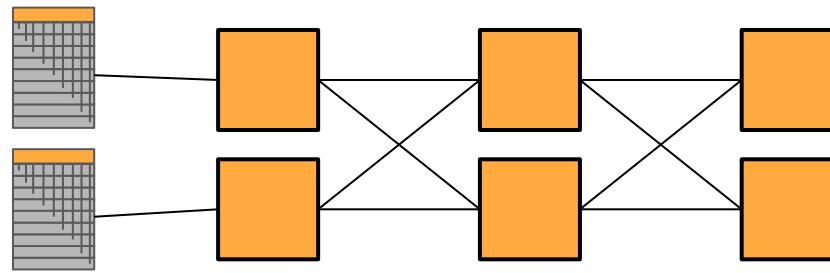
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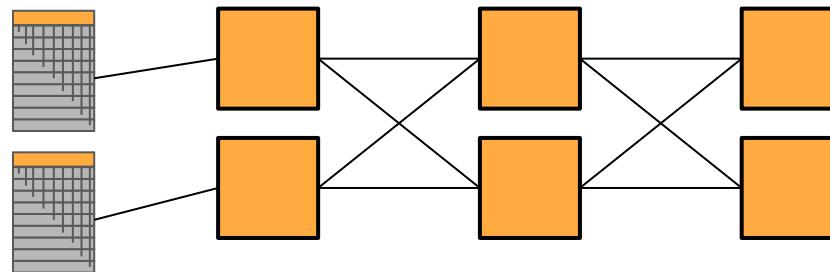
$16 \times 100\text{G}$  links failed to partition =  $1600\text{Gbps}$  bisection bandwidth



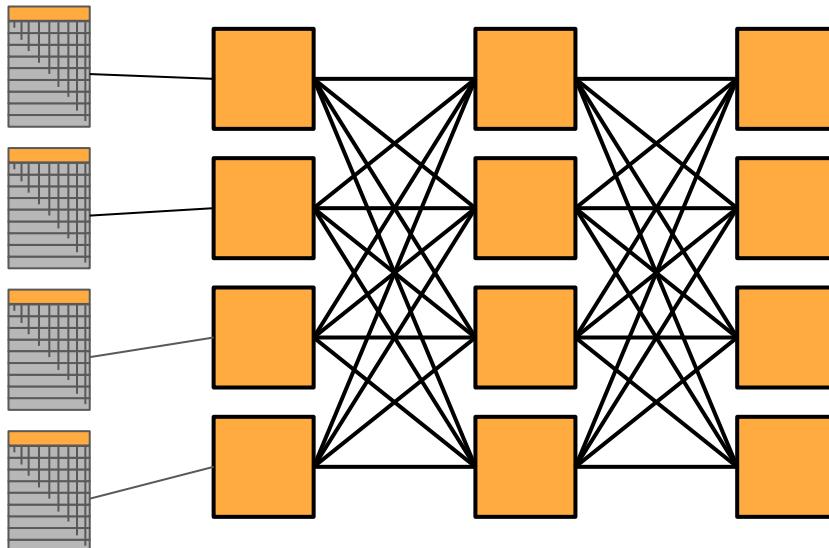
# Clos Networks - Bisection Bandwidth



$$\text{Full bisection bandwidth} = (4 * 80) / 2 * 100\text{G} = \mathbf{1600\text{G}}$$

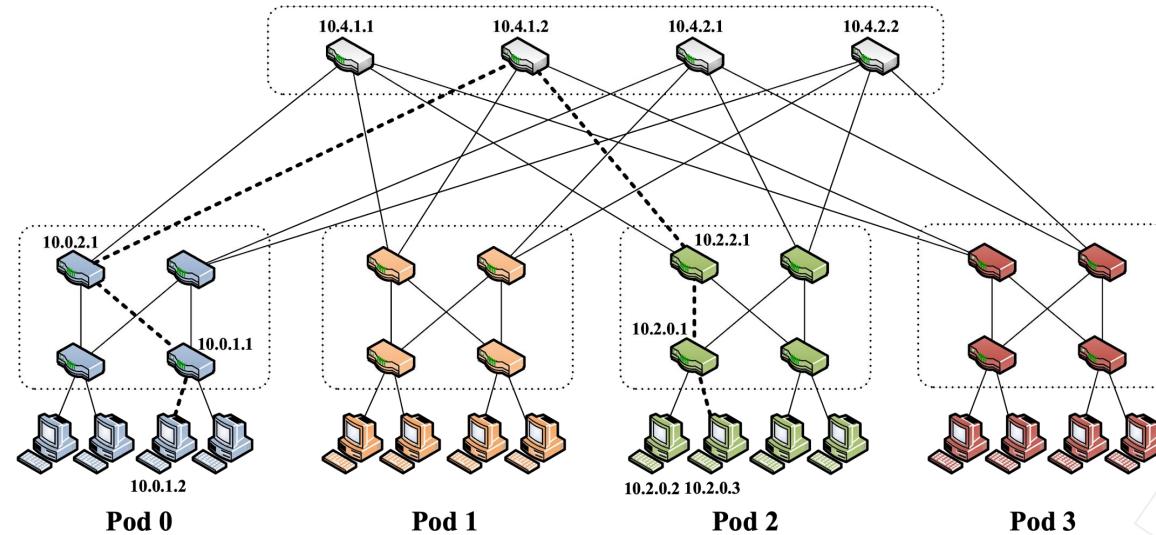


# Mixing Link Speeds



- Need not have all the links be exactly the same capacity.
- Server uplinks/access links can be lower bandwidth than switch to switch links.
- Easy to accomplish where switch chips allow “breaking out” of individual ports.
- e.g., 200G server uplink, 400G switch-to-switch

# Evolution of Clos Networks for DC



A Scalable, Commodity Data Center Network Architecture

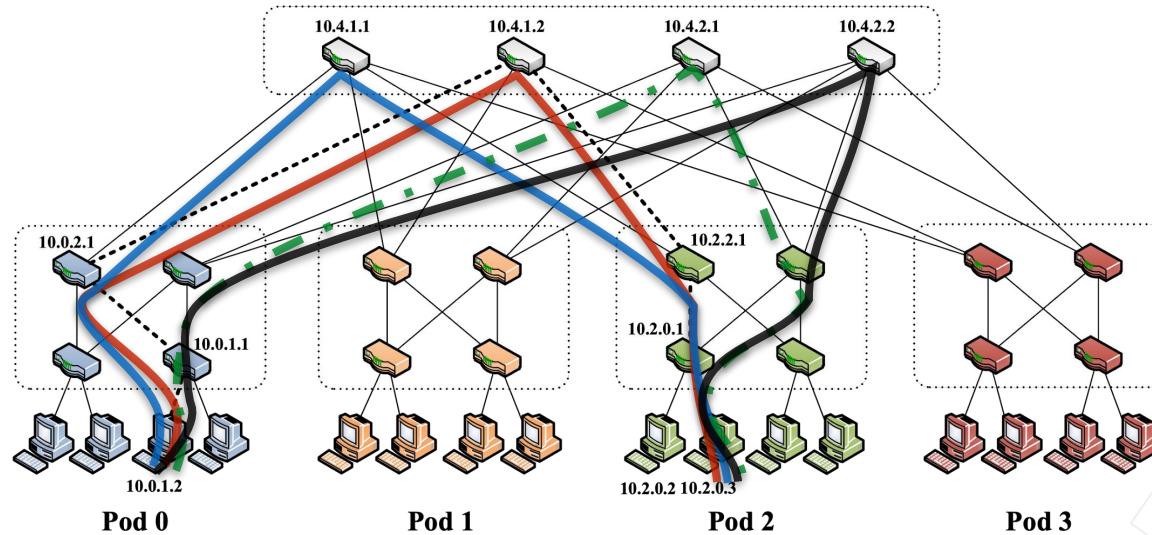
Mohammad Al-Fares

Alexander Loukissas

Amin Vahdat

ACM SIGCOMM 2008

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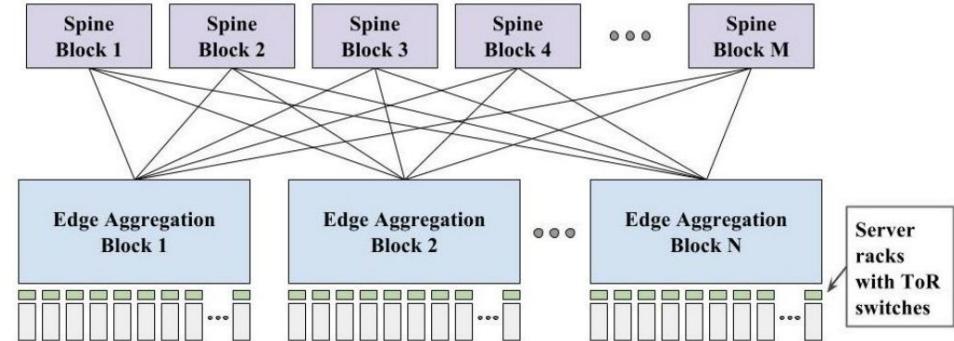
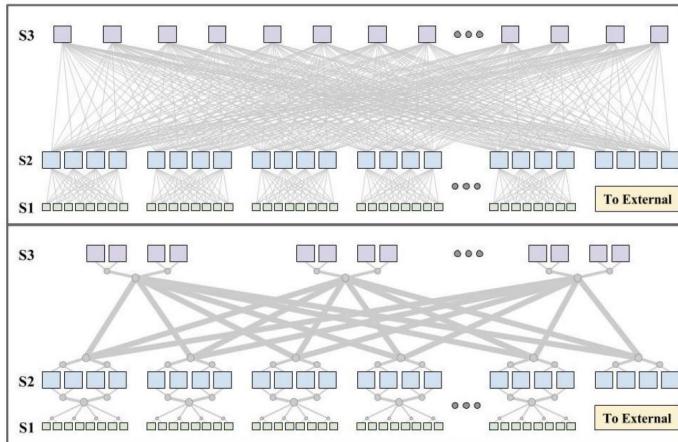
Mohammad Al-Fares

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ACM SIGCOMM 2008

# Evolution of Clos Networks for DC



Jupiter Rising: A **Decade** of Clos Topologies and  
Centralized Control in Google's Datacenter Network

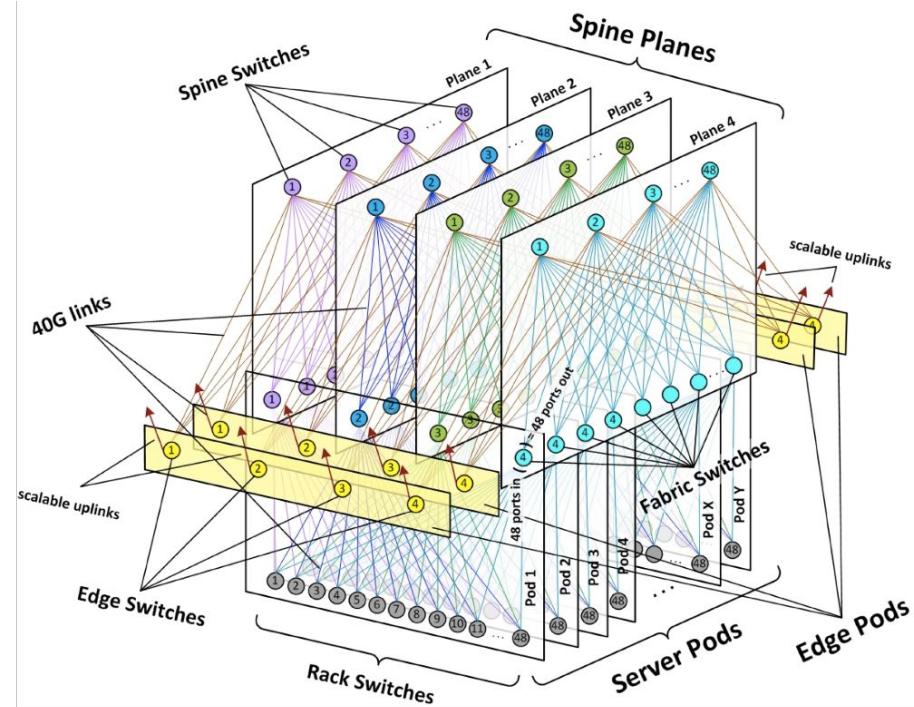
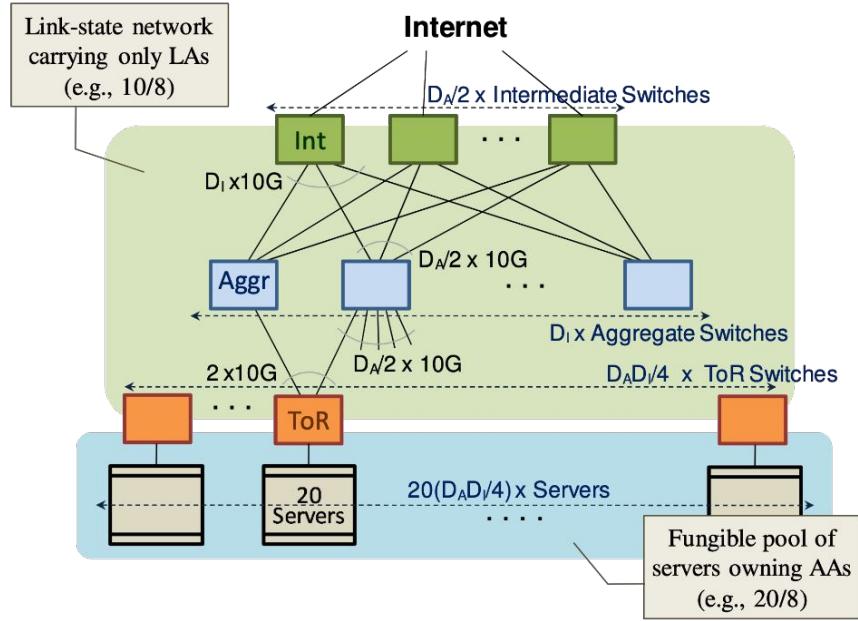
Arjun Singh, Joon Ong, Amit Agarwal, Glen Anderson, Ashby Armistead, Roy Bannon,  
Seb Boving, Gaurav Desai, Bob Felderman, Paulie Germano, Anand Kanagal, Jeff Provost,  
Jason Simmons, Eiichi Tanda, Jim Wanderer, Urs Hözle, Stephen Stuart, and Amin Vahdat

Google, Inc.

jupiter-sigcomm@google.com

ACM SIGCOMM 2015

# Design Variants are Common

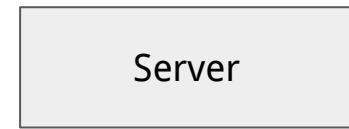


VL2 @ Microsoft, ACM SIGCOMM'09  
Greenburg, Hamilton, Jain, Kandula, Kim, Lahiri, Maltz, Patel, Sengupta

"Introducing data center fabric, the next-generation Facebook data center network", Alexey Andreyev, 2015

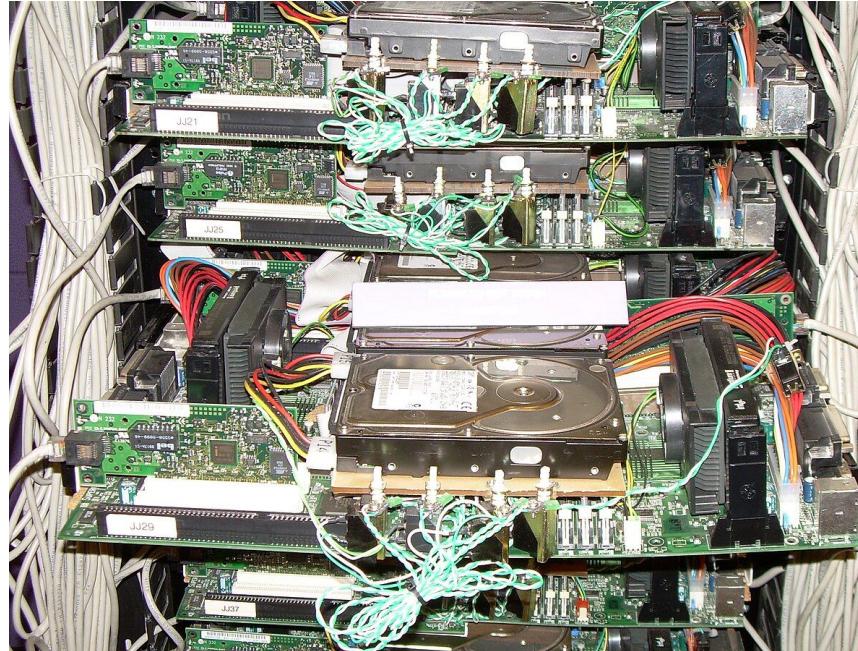
Questions?

# Components of a Datacenter Network



End hosts are servers within the datacenter.

# Components of a Datacenter Network

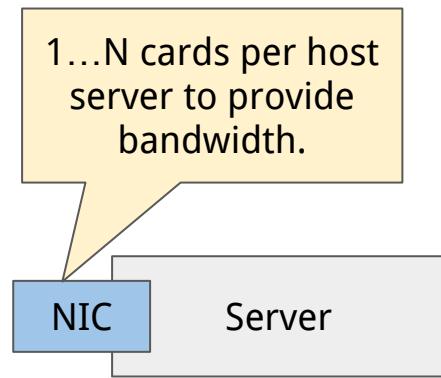
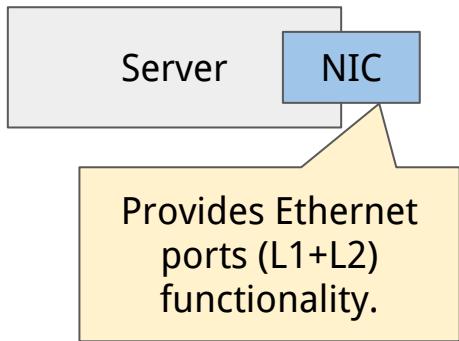


Server

Server

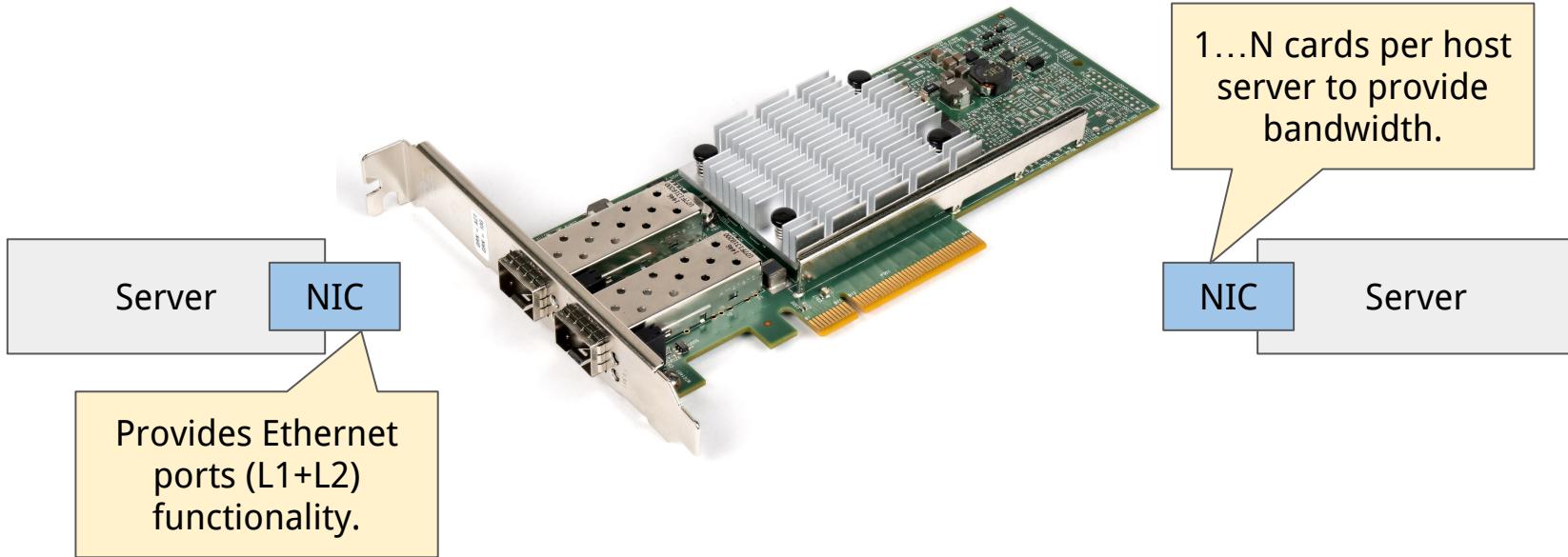
Historical Google servers within a rack.

# Components of a Datacenter Network



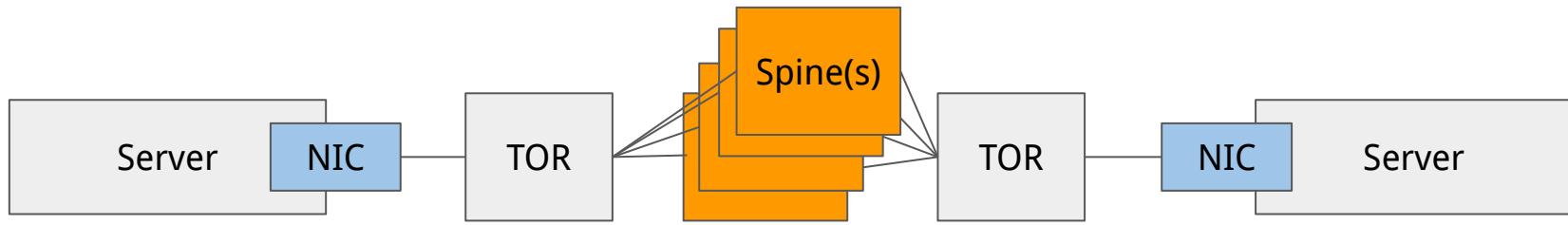
NIC = Network Interface Card.

# Components of a Datacenter Network



NIC = Network Interface Card.

# Components of a Datacenter Network



- Packets flow across multiple layers of switches according to the topology selected.
- All components (including servers, typically) are in the control of one organisation.

# Thinking about *optimisation*

- We identified ways that the network topology can be designed to support the applications that are deployed within a datacenter.
- What other components of the end-to-end system could we think about optimising?
- Next time – can we think about optimising routing.
  - e.g., how to choose paths to carry traffic according to performance requirements?
- Other questions:
  - How should *hosts* be designed to support datacenter applications?
  - How might congestion control be adapted to take advantage of datacenter characteristics?

# Summary

- Datacenters are single organisation, multi-application environments.
- A key criteria is high any-to-any bandwidth.
  - We characterise this as bisection bandwidth.
- The topology of the datacenter must be designed to both be scalable, and cost efficient.

# Next Time

- What else is different in datacenters?
  - Particularly, how does routing work in these topologies?
- How do we address the multi-tenant nature of a DC?