
Congestion Control

CS 168 – Fall 2024 – Discussion 9

Agenda

- Logistics
- Congestion Control
- Worksheet

Logistics

- Homework 3 was due Monday, October 28th
- Project 3a due Friday, November 1st
- Mid-Semester Form for extra credit!
 - Will try to reduce answering time for questions during discussions

TCP Circa 1986

- What happens when router buffers fill up?
 - Packets get dropped
- Then what happens at the **sender**?
 - Increased RTT, timeouts, retransmits
 - → **More packets in the network!!**
 - ... **So more retransmits!**
- Eventually, useful throughput approaches **zero**
- This is the **congestion collapse** of 1986!

Congestion Control

Goal of Congestion Control

- Limit the # of packets in flight
 - Utilize our fair share of bandwidth...
 - But don't overload the network
- Adapt to the right bandwidth
- Be fair
 - Links are shared among many hosts

A naïve solution

- **Overwhelmed? Tell the sender to slow down!**
 - Both routers & receiver
- “Source Quench”
 - signals the originating device to reduce the amount of traffic it is sending
- Problems?
 - If the link is already overwhelmed, these extra messages may be dropped too! (or add more traffic)

A host-based solution sketch

1. Pick an initial rate R
2. Try sending at rate R for some period of time
 - a. If congestion: reduce R
 - b. Else: increase R
3. Repeat step 2

A host-based solution sketch

1. Pick an initial rate $R \rightarrow$ How to pick initial rate
2. Try sending at rate R for some period of time \rightarrow How to detect congestion
 - a. If congestion: reduce $R \rightarrow$ How much to increase/decrease
 - b. Else: increase R
3. Repeat step 2

TCP: loss-based feedback

Idea: drop implies congestion

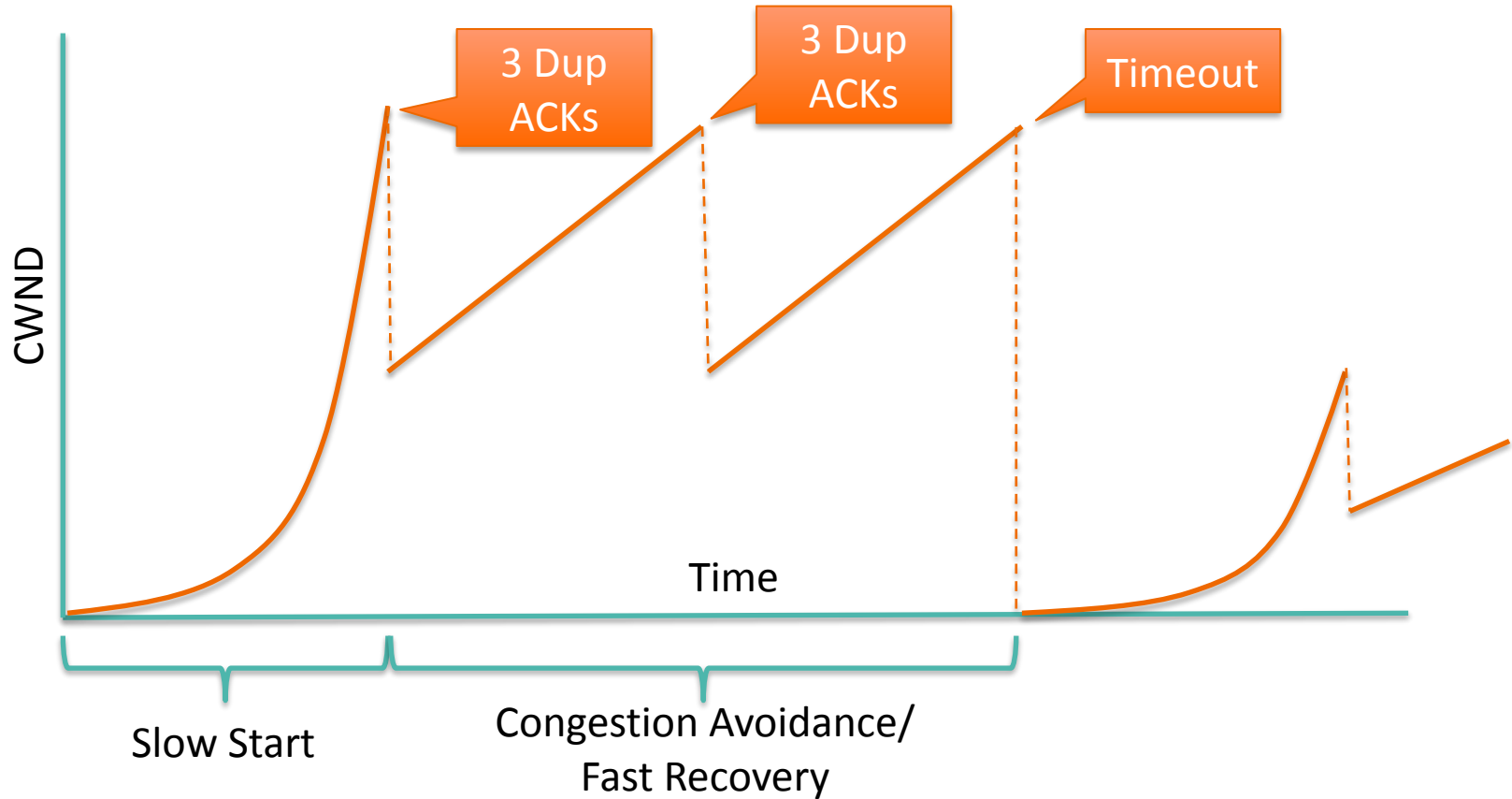
- 3 dupACK: minor congestion (ACKs get through)
- Timeout: major congestion (*nothing* gets through)

TCP's response depends on the *kind* of loss.

Congestion Control: Windows

- Receive Window (RWND)
 - What rate at which the **receiver** can process packets
- Congestion Window (CWND)
 - What rate at which the **network** can process packets
- Sending rate
 - Smaller of the two
- In this class, we assume $CWND \ll RWND$
 - **Network** will determine our sending rate

TCP Sawtooth

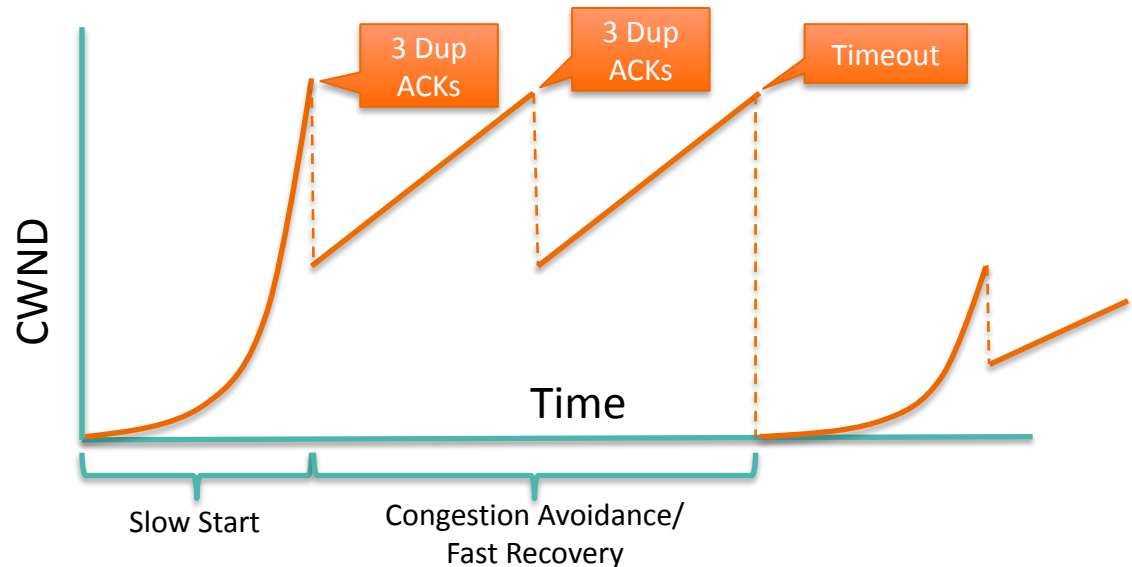


Utilization

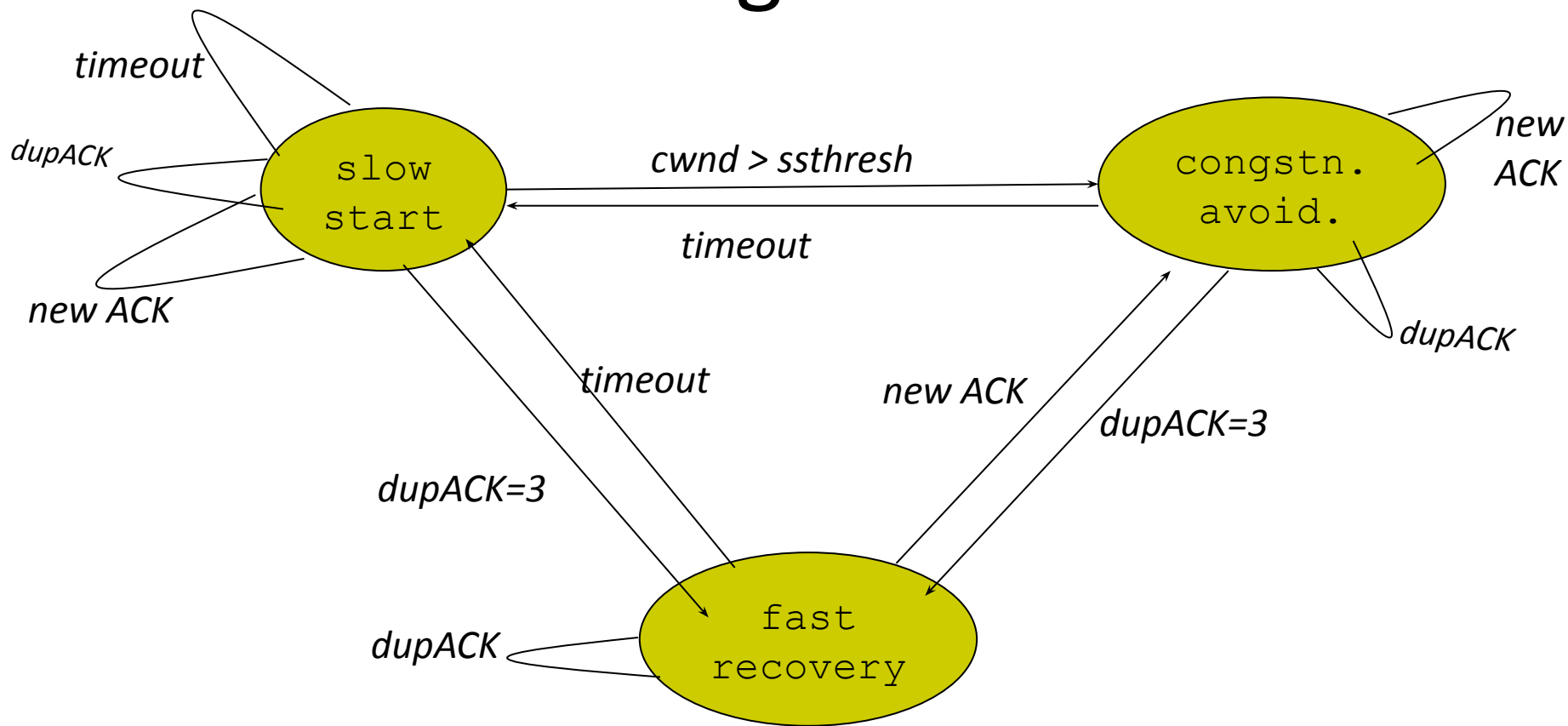
- The TCP sawtooth alternates between:
 - *Over-utilizing* bandwidth (causing drops)
 - *Under-utilizing* bandwidth
- Smart choices around buffering can result in higher utilization by absorbing the increase in window size

Three States

1. Slow Start
2. Congestion Avoidance
3. Fast Recovery



The Big Picture



Implementation

- State at sender
 - CWND
 - Max sending rate without congesting network (assuming $CWND \ll RWND$)
 - Ssthresh (slow-start threshold)
 - Threshold CWND for exiting slow start
 - dupACKcount
 - Count of contiguous duplicate ACKs received
 - timer

Congestion Control Mechanics

1. Slow Start
 - *Rapidly* increase our initial sending rate until we hit bottleneck
2. Congestion Avoidance
 - Adapting our sending rate to current network conditions
 - AIMD (**A**dditive **I**ncrease, **M**ultiplicative **D**ecrease)
3. Fast Recovery
 - Optimizing recovery from isolated loss
 - Detected through Duplicate ACKs

Implementation

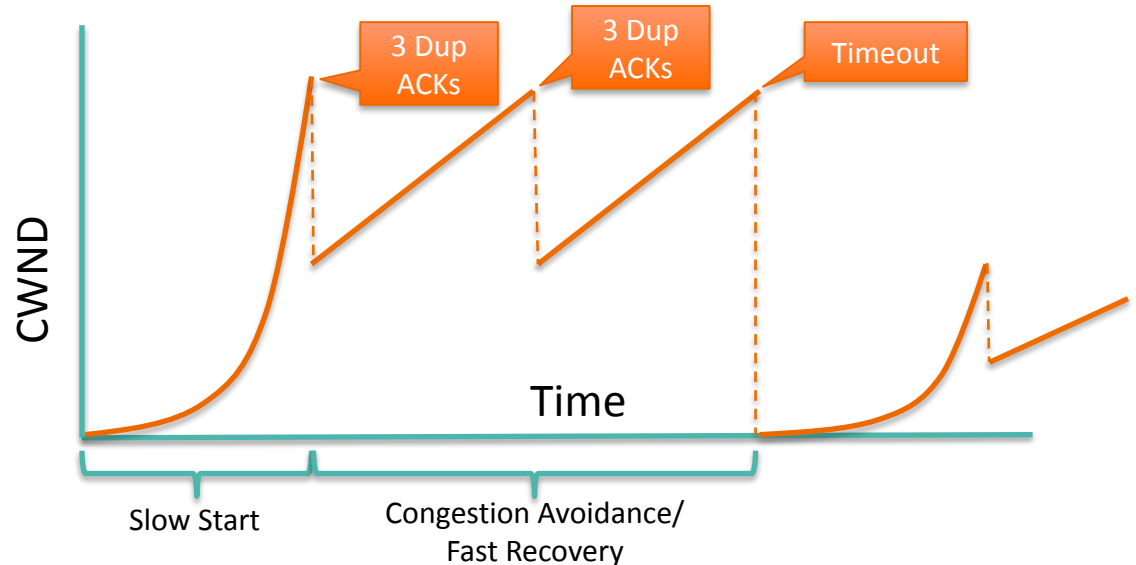
- Events at sender
 - ACK (new data)
 - dupACK (duplicate ACK for old data)
 - Timeout
- ... receiver just receives packets and sends ACKs

Now the Details

- Thanks Alex Triana, our amazing F'15 TA!

Three States

1. Slow Start
2. Congestion Avoidance
3. Fast Recovery



Slow Start

- Value of CWND starts at (small constant) * MSS
- For each packet that is acknowledged, increase the CWND by 1
 - Effectively **doubles** CWND every **RTT**!
- Window goes from $1 \rightarrow 2 \rightarrow 4 \rightarrow \dots$

Slow Start -- Intuition

- Instead of blasting packets based on the receive window
- Build up initial transmission rate *slowly*
- Back off when we've exceeded the capacity
- It's exponential growth, but still a slow start

Slow Start – When Does It End?

- 2 Ways
 - 1) If $CWND > ssthresh$
 - Enter congestion avoidance
 - 2) If we get 3 duplicate ACKs
 - Enter Fast Recovery
- If timeout:
 - Restart slow start, $ssthresh = cwnd/2$, $CWND = 1$

DupACK

- DupAckCount++

Timeout

- ssthresh = cwnd / 2
- cwnd = 1
- dupAckCount = 0

New ACK

- cwnd = cwnd + 1
- dupAckCount = 0

Slow
Start



dupAckCount == 3

Move to Fast Recovery

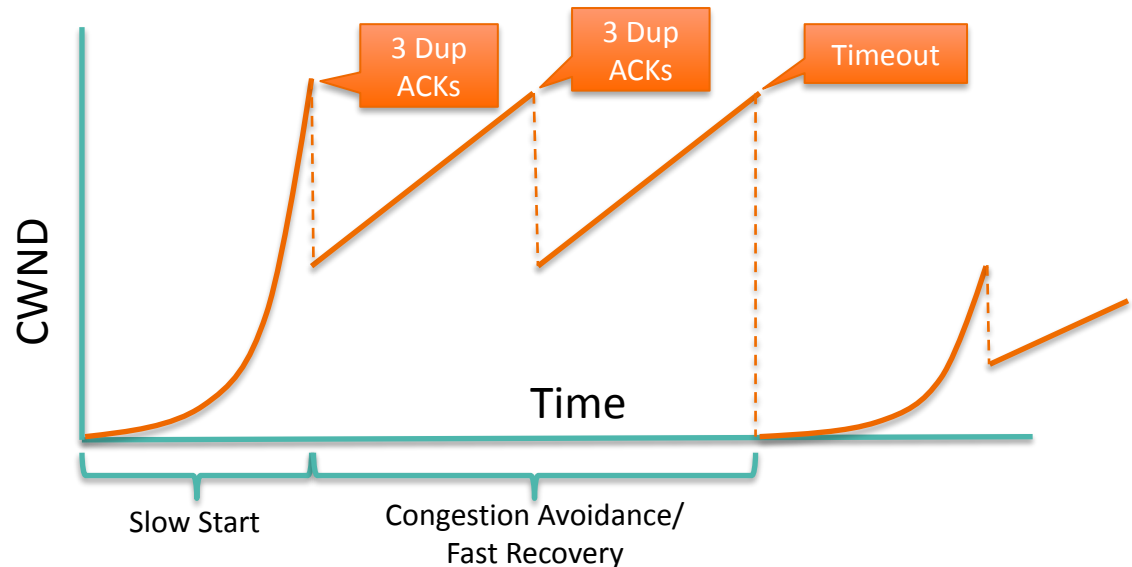
- ssthresh = cwnd / 2
- cwnd = ssthresh + 3
- Retransmit missing packet

cwnd > ssthresh

Move to Congestion Avoidance

Three States

1. Slow Start
2. Congestion Avoidance
3. Fast Recovery



Congestion Avoidance -- Intuition

- In the steady state
- Constantly probe for more bandwidth
- When we've exceeded – back off aggressively

Congestion Avoidance

- Growth is more conservative than slow start
- After each **new** ACK, increase CWND by $1 / \text{CWND}$
 - After one **RTT**, CWND will have **increased by ~1**
- When does it stop?
 - 1) Timeout → back to slow start
 - 2) 3 duplicate ACKS → Fast recovery

DupACK

-DupAckCount++

New ACK

- $cwnd = cwnd + 1/cwnd$
- $dupAckCount = 0$

Congestion Avoidance

Timeout

Move to Slow Start

- $ssthresh = cwnd / 2$
- $cwnd = 1$
- $dupAckCount = 0$

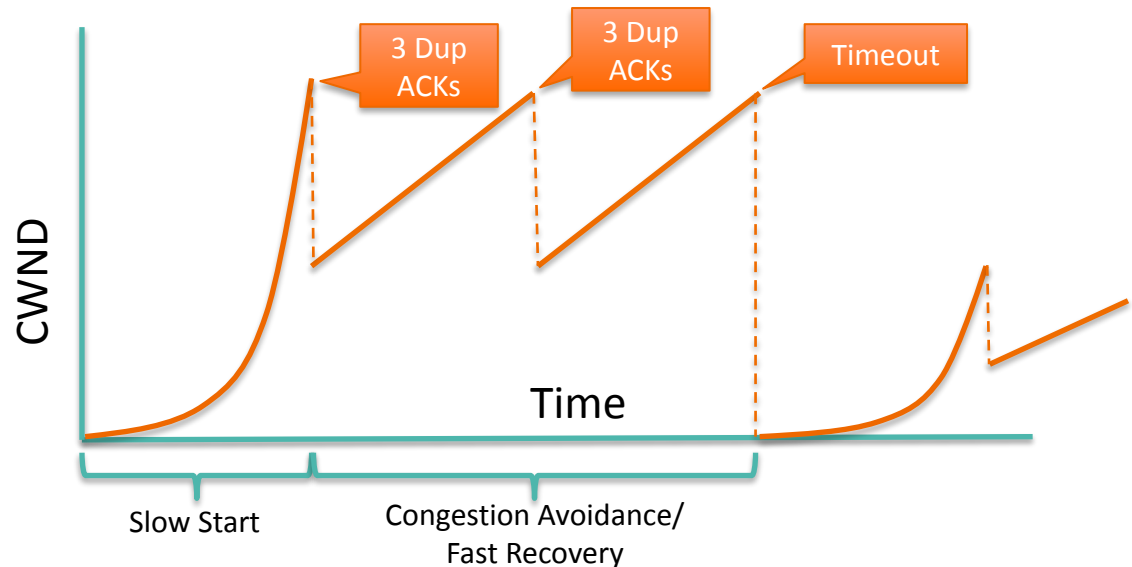
$dupACKCount == 3$

Move to Fast Recovery

- $ssthresh = cwnd / 2$
- $cwnd = ssthresh + 3$
- Retransmit missing packet

Three States

1. Slow Start
2. Congestion Avoidance
3. Fast Recovery



Fast Recovery – Intuition

- A single lost packet
 - May just be a fluke
- Resetting CWND may be too aggressive
- Instead just retransmit that single packet
 - And continue as if nothing happened

Fast Recovery

- Every **duplicate** ACK increases the window by 1
- When does it stop?
 - 1) Timeout → back to Slow Start
 - 2) New ACK → back to Congestion Avoidance

Timeout

Move to Slow Start

- $\text{ssthresh} = \text{cwnd} / 2$
- $\text{cwnd} = 1$
- $\text{dupAckCount} = 0$

DupACK

- $\text{cwnd} = \text{cwnd} + 1$

Fast Recovery

A central light blue oval labeled "Fast Recovery" has three orange arrows pointing away from it: one to the top-left, one to the top-right, and one that loops back to the bottom. A blue arrow points into the oval from the bottom.

New ACK

Move to Congestion Avoidance

- $\text{cwnd} = \text{ssthresh}$
- $\text{dupAckCount} = 0$

Big Ideas

- Fundamental concepts:
 - Slow Start
 - AIMD
- Hack
 - Fast Recovery
- Lesson
 - Sometimes, BAND-AIDs scale remarkably well!

Feedback Form:
<https://tinyurl.com/cs168-disc-fa24>



Worksheet

Question 1

1. UDP uses congestion control.
2. Flow control slows down the sender when the network is congested.
3. For TCP timer implementations, every time the sender receives an ACK for a previously unACKed packet, it will recalculate ETO.
4. CWND (congestion window) is usually smaller than RWND (receiver window).
5. AIMD is the only "fair" option among MIMD, AIAD, MIAD, and AIMD.

1) Without Fast Recovery

- On new ACK, $CWND = CWND + 1/Floor(CWND)$
- On triple duplicate ACKs, $SSTHRESH = Floor(CWND/2)$, then $CWND = SSTHRESH$

CWND=102-111

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
1.0	102 (101)	$10 + 1/Floor(10) = 10.1$	111 (No)
1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/

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CWND=102-106

Ssthresh=5

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
1.0	102 (101)	$10 + 1/\text{Floor}(10) = 10.1$	111 (No)
1.2	102 (103)	10.1	/
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1.4	102 (105)	$\text{Floor}(10.1/2) = 5$	102 (Yes)

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1.5	102 (106)	5	/

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1.5	102 (106)	5	/
1.6	102 (107)	5	/

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1.7	102 (108)	5	/

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1.6	102 (107)	5	/
1.7	102 (108)	5	/
1.8	102 (109)	5	/

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1.7	102 (108)	5	/
1.8	102 (109)	5	/
1.9	102 (110)	5	/

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1.6	102 (107)	5	/
1.7	102 (108)	5	/
1.8	102 (109)	5	/
1.9	102 (110)	5	/
2.0	102 (111)	5	/

1) Without Fast Recovery

- On new ACK, **CWND = CWND + 1/Floor(CWND)**
- On triple duplicate ACKs, **SSTHRESH = Floor(CWND/2)**, then **CWND = SSTHRESH**

CWND=112-116

Ssthresh=5

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
1.0	102 (101)	$10 + 1/\text{Floor}(10) = 10.1$	111 (No)
1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2) = 5$	102 (Yes)
1.5	102 (106)	5	/
1.6	102 (107)	5	/
1.7	102 (108)	5	/
1.8	102 (109)	5	/
1.9	102 (110)	5	/
2.0	102 (111)	5	/
2.4	112 (102)	$5 + 1/\text{Floor}(5) = 5.2$	112 – 116 (No)

2) With Fast Recovery

- On triple duplicate ACK, $SSTHRESH = CWND/2$, then $CWND = SSTHRESH + 3$, and enter fast recovery
- In fast recovery, $CWND += 1$ on every duplicate ACK
- Exit fast recovery on new ACK, setting $CWND = SSTHRESH$ CWND=102-111

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
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$CWND=102-109$

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1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2)+3=8$	102 (Yes)

Ssthresh=5

2) With Fast Recovery

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- Exit fast recovery on new ACK, setting $CWND = SSTHRESH$ CWND=102-110

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1.0	102 (101)	$10+1/\text{Floor}(10)=10.1$	111 (No)
1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2)+3=8$	102 (Yes)
1.5	102 (106)	9	/

Ssthresh=5

2) With Fast Recovery

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1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2) + 3 = 8$	102 (Yes)
1.5	102 (106)	9	/
1.6	102 (107)	10	/

Ssthresh=5

2) With Fast Recovery

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- In fast recovery, $CWND += 1$ on every duplicate ACK
- Exit fast recovery on new ACK, setting $CWND = SSTHRESH$ CWND=102-112

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
1.0	102 (101)	$10+1/\text{Floor}(10)=10.1$	111 (No)
1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2)+3=8$	102 (Yes)
1.5	102 (106)	9	/
1.6	102 (107)	10	/
1.7	102 (108)	11	112 (No)

Ssthresh=5

2) With Fast Recovery

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- In fast recovery, $CWND += 1$ on every duplicate ACK
- Exit fast recovery on new ACK, setting $CWND = SSTHRESH$ CWND=102-113

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
1.0	102 (101)	$10+1/\text{Floor}(10)=10.1$	111 (No)
1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2)+3=8$	102 (Yes)
1.5	102 (106)	9	/
1.6	102 (107)	10	/
1.7	102 (108)	11	112 (No)
1.8	102 (109)	12	113 (No)

Ssthresh=5

2) With Fast Recovery

- On triple duplicate ACK, $SSTHRESH = CWND/2$, then $CWND = SSTHRESH + 3$, and enter fast recovery
- In fast recovery, $CWND += 1$ on every duplicate ACK
- Exit fast recovery on new ACK, setting $CWND = SSTHRESH$ CWND=102-114

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
1.0	102 (101)	$10+1/\text{Floor}(10)=10.1$	111 (No)
1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2)+3=8$	102 (Yes)
1.5	102 (106)	9	/
1.6	102 (107)	10	/
1.7	102 (108)	11	112 (No)
1.8	102 (109)	12	113 (No)
1.9	102 (110)	13	114 (No)

Ssthresh=5

2) With Fast Recovery

- On triple duplicate ACK, $SSTHRESH = CWND/2$, then $CWND = SSTHRESH + 3$, and enter fast recovery
- In fast recovery, $CWND += 1$ on every duplicate ACK
- Exit fast recovery on new ACK, setting $CWND = SSTHRESH$ CWND=102-115

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
1.0	102 (101)	$10+1/\text{Floor}(10)=10.1$	111 (No)
1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2)+3=8$	102 (Yes)
1.5	102 (106)	9	/
1.6	102 (107)	10	/
1.7	102 (108)	11	112 (No)
1.8	102 (109)	12	113 (No)
1.9	102 (110)	13	114 (No)
2.0	102 (111)	14	115 (No)

Ssthresh=5

2) With Fast Recovery

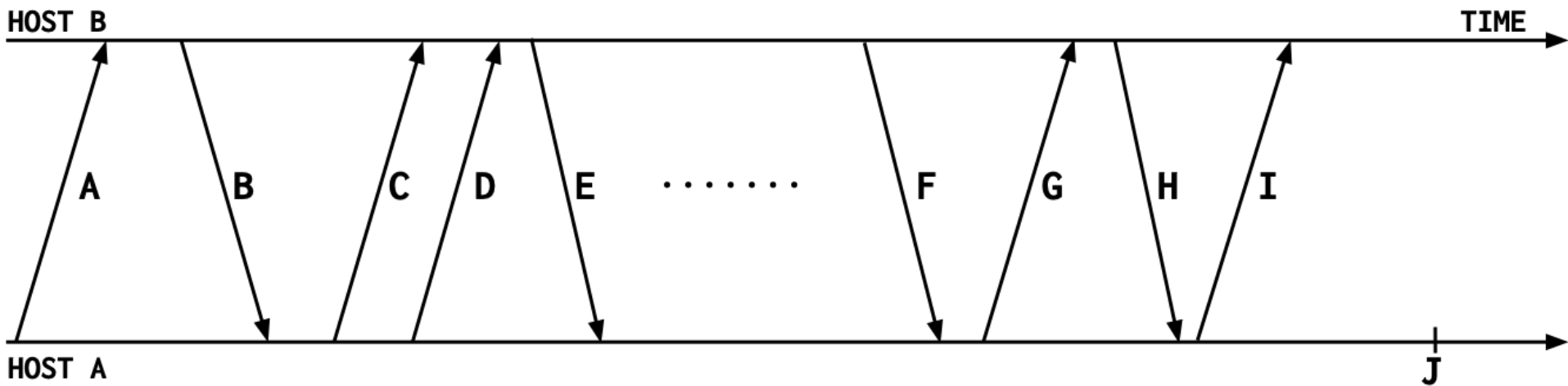
- On triple duplicate ACK, $SSTHRESH = CWND/2$, then $CWND = SSTHRESH + 3$, and enter fast recovery
- In fast recovery, $CWND += 1$ on every duplicate ACK
- Exit fast recovery on new ACK, setting $CWND = SSTHRESH$

CWND=112-116

Time (sec)	Receive ACK (due to)	CWND	Transmit Seq # (yes for retransmit)
1.0	102 (101)	$10+1/\text{Floor}(10)=10.1$	111 (No)
1.2	102 (103)	10.1	/
1.3	102 (104)	10.1	/
1.4	102 (105)	$\text{Floor}(10.1/2)+3=8$	102 (Yes)
1.5	102 (106)	9	/
1.6	102 (107)	10	/
1.7	102 (108)	11	112 (No)
1.8	102 (109)	12	113 (No)
1.9	102 (110)	13	114 (No)
2.0	102 (111)	14	115 (No)
2.4	112 (102)	$SSTHRESH = 5$	116 (No)

Ssthresh=5

Question 3



A:

D:

G:

B:

E:

H:

C:

F:

I:

J

Feedback Form:
<https://tinyurl.com/cs168-disc-fa24>

