

CO 250: Introduction to Optimization

Module 1: Formulations (LP Models)

Constrained Optimization

- In this course, we consider optimization problems of the following form:

$$\min\{f(x) : g_i(x) \leq b_i, (1 \leq i \leq m), x \in \mathbb{R}^n\},$$

where

- $n, m \in \mathbb{N}$,
- $b_1, \dots, b_m \in \mathbb{R}$, and
- f, g_1, \dots, g_m are functions with from \mathbb{R}^n to \mathbb{R} .
- Problems like the above are **very hard** to solve in general
 \implies we focus on special cases.
- **This class:** All functions are **affine**.

Modeling: Linear Programs

Affine Functions

Definition

A function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is **affine** if $f(x) = a^T x + \beta$ for $a \in \mathbb{R}^n$, $\beta \in \mathbb{R}$. It is **linear** if, in addition, $\beta = 0$.

Example

- (i) $f(x) = 2x_1 + 3x_2 - x_3 + 7$ (**affine**, but **not linear**)
- (ii) $f(x) = -3x_1 + 5x_3$ (**linear**)
- (iii) $f(x) = 5x - 3 \cos(x) + \sqrt{x}$ (**not affine** and **not linear**)

Definition

The optimization problem

$$\min \{f(x) : g_i(x) \leq b_i, \forall 1 \leq i \leq m, x \in \mathbb{R}^n\} \quad (\text{P})$$

is called a **linear program** if f is **affine** and g_1, \dots, g_m is **finite** number of **linear** functions.

Comments:

- Instead of **set notation**, we often write LPs more verbosely.
- Often give **non-negativity** constraints separately
- May use **max** instead of **min**
- Sometimes replace **subject to** by **s.t.**

$$\begin{array}{ll} \text{max} & -2x_1 + 7x_3 \\ \text{subject to} & x_1 + 7x_2 \leq 3 \\ & 3x_2 + 4x_3 \leq 2 \\ & x_1, x_3 \geq 0 \end{array}$$

- We often write $x \geq 0$ as a short form for **all variables are non-negative**.

$$\begin{array}{ll}
 \min & -x_1 - 2x_2 - x_3 \\
 \text{s.t.} & 2x_1 + x_3 \geq 3 \\
 & x_1 + 2x_2 = 2 \\
 & x \geq 0
 \end{array}$$

- The second mathematical program is **not an LP**.

There are three reasons:

- Dividing by variables is not allowed.
- Cannot have **strict** inequalities.
- Must have **finite** number of constraints.

$$\begin{array}{ll}
 \max & -1/x_1 - x_3 \\
 \text{subject to} & 2x_1 + x_3 < 3 \\
 & x_1 + \alpha x_2 = 2 \quad \forall \alpha \in \mathbb{R}
 \end{array}$$

Production Revisited

$$\begin{array}{ll} \max & 300x_1 + 260x_2 + 220x_3 + 180x_4 - 8y_s - 6y_u \\ & 11x_1 + 7x_2 + 6x_3 + 5x_4 \leq 700 \\ & 4x_1 + 6x_2 + 5x_3 + 4x_4 \leq 500 \\ & 8x_1 + 5x_2 + 5x_3 + 6x_4 \leq y_s \\ \text{s.t.} & 7x_1 + 8x_2 + 7x_3 + 4x_4 \leq y_u \\ & y_s \leq 600 \\ & y_u \leq 650 \\ & x_1, x_2, x_3, x_4, y_u, y_s \geq 0. \end{array}$$

The mathematical program for **WaterTech** example from last class is in fact an LP!

Multiperiod Models

A main feature of the WaterTech **production** model:
Decisions about production levels have to be made **once and for all**.

In practice, we often have to make a **series** of decisions that influence each other.

One such example is **multiperiod models**:

- Time is split into **periods**.
- We have to make a **decision in each period**.
- All decisions influence the final outcome.

KW Oil

KW Oil is local **supplier of heating oil**.

It needs to decide on **how much oil to purchase** in order to **satisfy demand** of its customers.

Years of experience give the following **demand forecast for the next 4 months**:

| Month | 1 | 2 | 3 | 4 |
|-------------------|------|------|------|------|
| Demand (ℓ) | 5000 | 8000 | 9000 | 6000 |

The projected **price of oil fluctuates** from month to month:

| Month | 1 | 2 | 3 | 4 |
|---------------------|------|------|------|------|
| Price ($\$/\ell$) | 0.75 | 0.72 | 0.92 | 0.90 |

Question: **When** should we purchase **how much** oil?

Question: When should we purchase how much oil when the goal is to minimize overall total cost?

Additional Complication: The company has a storage tank that

- has a capacity of 4000 litres of oil, and
- currently (beginning of month 1) contains 2000 litres of oil.

Assumption: Oil is delivered at beginning of the month, and consumption occurs in the middle of the month.

KW Oil Model – Variables

| Month | 1 | 2 | 3 | 4 |
|-------------------|------|------|------|------|
| Demand (ℓ) | 5000 | 8000 | 9000 | 6000 |

| Month | 1 | 2 | 3 | 4 |
|---------------------|------|------|------|------|
| Price ($\$/\ell$) | 0.75 | 0.72 | 0.92 | 0.90 |

- (i) Need to decide **how many litres of oil to purchase in each month i** .
→ variable p_i for $i \in [4]$
- (ii) How much **oil is stored in the tank** at beginning of month i ?
→ variable t_i for $i \in [4]$

Objective Function

Minimize **cost of oil** procurement.

$$\min \quad 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$$

Constraints: When do

$$t_1, \dots, t_4, p_1, \dots, p_4$$

correspond to a **feasible purchasing scheme?**

Variables:

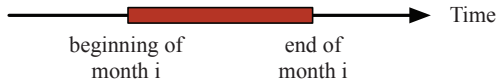
p_i : oil purchase in month i

t_i : tank level in month i

Constraints

Assumptions:

- (i) Oil is purchased at beginning of month.
- (ii) Oil is consumed afterwards.



Variables:

p_i : oil purchased in month i

t_i : tank level in month i

We need: $p_i + t_i \geq \{\text{demand in month } i\}$

[Balance Equation] $p_i + t_i = \{\text{demand in month } i\} + t_{i+1}$

Constraints

[Balance Equation] $p_i + t_i = \{\text{demand in month } i\} + t_{i+1}$

Tank content in month 1: 2000 litres

- Month 1:

$$p_1 + 2000 = 5000 + t_2$$

- Month 2:

$$p_2 + t_2 = 8000 + t_3$$

- Month 3:

$$p_3 + t_3 = 9000 + t_4$$

- Month 4:

$$p_4 + t_4 \geq 6000$$

| Month | 1 | 2 | 3 | 4 |
|-------------------|------|------|------|------|
| Demand (ℓ) | 5000 | 8000 | 9000 | 6000 |

$$\min \quad 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$$

subject to

$$p_1 + t_2 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

$$p_3 + t_3 = 9000 + t_4$$

$$p_4 + t_4 \geq 6000$$

$$t_1 = 2000$$

$$t_i \leq 4000 \quad (i = 2, 3, 4)$$

$$t_i, p_i \geq 0 \quad (i = 1, 2, 3, 4)$$

Solution: $p = (3000, 12000, 5000, 6000)^T$, and
 $t = (2000, 0, 4000, 0)^T$

KW Oil: Add-Ons

Can easily capture **additional features**. E.g.,...

- **Storage comes at a cost:** storage cost is \$.15 per litre/month.

Add $\sum_{i=1}^4 0.15t_i$ to objective.

- Minimize the **maximum # of litres of oil purchased** over all months.
 - (i) We will need a new **variable M** for maximum # of litres purchased.
 - (ii) We will have to add constraints

$$\min \quad 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$$

subject to

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

$$p_3 + t_3 = 9000 + t_4$$

$$p_4 + t_4 \geq 6000$$

$$t_1 = 2000$$

$$t_i \leq 4000 \quad (i = 2, 3, 4)$$

$$t_i, p_i \geq 0 \quad (i = 1, 2, 3, 4)$$

KW Oil: Add-Ons

(i) Add **variable M** for **maximum # of litres purchased** over all months.

(ii) Add constraints

$$p_i \leq M$$

for all $i \in [4]$.

(iii) **Done?** No! We need to replace the objective function with

$$\min \quad M$$

$$\min \quad 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$$

subject to

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

$$p_3 + t_3 = 9000 + t_4$$

$$p_4 + t_4 \geq 6000$$

$$t_1 = 2000$$

$$t_i \leq 4000 \quad (i = 2, 3, 4)$$

$$t_i, p_i \geq 0 \quad (i = 1, 2, 3, 4)$$

Goal: Minimize **maximum # of litres of oil purchased** over all months.

Minimizing the Maximum Purchase: LP

$$\min \quad M$$

s.t.

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

$$p_3 + t_3 = 9000 + t_4$$

$$p_4 + t_4 \geq 6000$$

$$t_1 = 2000$$

$$t_i \leq 4000 \quad (i = 2, 3, 4)$$

$$p_i \leq M \quad (i = 1, 2, 3, 4)$$

$$t_i, p_i \geq 0 \quad (i = 1, 2, 3, 4)$$

KW Oil: Correctness

- Why is this a **correct** model?

- Suppose that

$M, p_1, \dots, p_4, t_1, \dots, t_4$
is an **optimal**
solution to the LP.

- Clearly:

$$M \geq \max_i p_i$$

- Since M, p, t is
optimal we must
have $M = \max_i p_i$.

Why?

$$\min \quad M$$

s.t.

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

$$p_3 + t_3 = 9000 + t_4$$

$$p_4 + t_4 \geq 6000$$

$$t_1 = 2000$$

$$t_i \leq 4000 \quad (i = 2, 3, 4)$$

$$p_i \leq M \quad (i = 1, 2, 3, 4)$$

$$t_i, p_i \geq 0 \quad (i = 1, 2, 3, 4)$$

Otherwise, we could decrease M by a little bit, without violating feasibility. This would contradict optimality because we would get a new feasible solution that has a smaller objective function.