



## Problem Set #2

Quiz, 5 questions

5/5 points (100%)

Congratulations! You passed!

Next Item

1 / 1  
point

1.

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence  $T(n) = 7 * T(n/3) + n^2$ . What's the overall asymptotic running time (i.e., the value of  $T(n)$ )?

☐  $\theta(n^{2.81})$ ☒  $\theta(n^2)$ **Correct**

$a=7$ ,  $b=3$ ,  $d=2$ . Since  $b^d > a$ , this is case 2 of the Master Method.

☐  $\theta(n^2 \log n)$ ☐  $\theta(n \log n)$ 1 / 1  
point

2.

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence  $T(n) = 9 * T(n/3) + n^2$ . What's the overall asymptotic running time (i.e., the value of  $T(n)$ )?

☐  $\theta(n^{3.17})$ ☐  $\theta(n^2)$ ☒  $\theta(n^2 \log n)$ **Correct**

$a = b^d = 9$ , so this is case 1 of the Master Method.

☐  $\theta(n \log n)$

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3.

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence  $T(n) = 5 * T(n/3) + 4n$ . What's the overall asymptotic running time (i.e., the value of  $T(n)$ )?

☐  $\theta(n^2)$

☐  $\theta(n \log(n))$

☐  $\theta(n^{2.59})$

☐  $\theta(n^{\frac{\log 3}{\log 5}})$

☒  $\theta(n^{\log_3(5)})$

**Correct**

$a = 5, b = 3, d = 1$ . Since  $a > b^d$ , this is case 3 of the Master Method.

☐  $\theta(n^{5/3})$

1 / 1  
point

4.

Consider the following pseudocode for calculating  $a^b$  (where  $a$  and  $b$  are positive integers)

```

1 FastPower(a,b) :
2   if b = 1
3     return a
4   else
5     c := a*a
6     ans := FastPower(c, [b/2])
7   if b is odd
8     return a*ans
9   else return ans
10 end

```

Here  $[x]$  denotes the floor function, that is, the largest integer less than or equal to  $x$ .

Now assuming that you use a calculator that supports multiplication and division (i.e., you can do multiplications and divisions in constant time), what would be the overall asymptotic running time of the above algorithm (as a function of  $b$ )?

☒  $\Theta(\log(b))$

**Correct**

Constant work per digit in the binary expansion of  $b$ .

☐  $\Theta(b)$



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5.

Choose the smallest correct upper bound on the solution to the following recurrence:  $T(1) = 1$  and  $T(n) \leq T(\lfloor \sqrt{n} \rfloor) + 1$  for  $n > 1$ . Here  $\lfloor x \rfloor$  denotes the "floor" function, which rounds down to the nearest integer. (Note that the Master Method does not apply.)

 $O(\sqrt{n})$  $O(\log n)$  $O(1)$  $O(\log \log n)$ **Correct**

Bingo! This answer may be easiest to see by writing  $n$  as  $2^{\log n}$  and then noting that every square-root operation cuts the exponent in half.

