

## **HW1** (due March 21, 2021 Tuesday before recitation)

1- Sketch a composed TM  $M$  that performs the computation :

$$(s, \diamond \# \omega) \xrightarrow{*}_M (h, \diamond \# \omega^c)$$

where  $\omega^c$  is the compressed version of  $\omega$  where all the characters '\$' within  $\omega$  are removed.

2- Problems from the main text book :

4.1.4, 4.1.7, 4.1.12, 4.2.2, 4.3.1, 4.3.2 (a);(b), 4.3.6

An optional HW question

**Question** – A Turing machine  $M_n$  operates on an abstract  **$n$ -dimensional quadrant** given by  $X = N \times \dots \times N = N^n$  where  $N = \{0, 1, 2, \dots\}$  denotes the set of natural numbers. The head can move in positive and negative directions along each of the dimensions; hence it has precisely  **$2n$**  moves at each slot position. Each slot is an  **$n$ -dimensional cube** that carries data represented by a symbol in its alphabet set  $\Sigma$ . A standard TM  $M$  to simulate  $M_n$  must assign a unique slot on its tape corresponding to each  $n$ -dimensional cubic slot of  $M_n$ .

This is accomplished by defining functions  $f_k: N^k \rightarrow N$  defined inductively via functions  $f_1, f_2, \dots, f_n$ , each with the domain  $N^k$  for  $k=1, \dots, n$  as follows :

$$f_1(i_1) = i_1 \text{ for all } i_1 \in N$$

$$f_k(i_1, i_2, \dots, i_k) := (f_{k-1}(i_1, i_2, \dots, i_{k-1}) + i_k)(f_{k-1}(i_1, i_2, \dots, i_{k-1}) + i_k + 1)/2 + i_k, \quad k=2, 3, \dots, n \quad (1)$$

Hence each  $f_k(i_1, i_2, \dots, i_k)$  can be computed sequentially using the inductive recipe above.

For example  $f_2(i_1, i_2) = (i_1 + i_2)(i_1 + i_2 + 1)/2 + i_2$  etc.

### **Theorem**

The function  $f_n: N^n \rightarrow N$  described above by (1) is a **bijection**.

i.e. :

(i)  $f_n(i_1, i_2, \dots, i_n) = f_n(j_1, j_2, \dots, j_n) \Rightarrow i_p = j_p$  for  $p=1, \dots, n$  (hence it is an **injection** (1-to-1))

(ii) for every  $u \in N$  there exist  $(i_1, \dots, i_n) \in N^n$  such that  $f_n(i_1, i_2, \dots, i_n) = u$  (hence it is a **surjection** (onto) )

Using this theorem describe how  $M$  can simulate  $M_n$ . In particular describe how  $M$  moves its head when  $M_n$  moves its head from slot  $(i_1, \dots, i_p, \dots, i_n)$  to  $(i_1, \dots, i_p \pm 1, \dots, i_n)$ .

(To get a feeling for the problem first solve it for  $n=2$ , namely simulating a 2D-TM ! )

**Meanwhile try also to prove the theorem above using induction on  $n$**