

Numerical Analysis Lecture Notes

Rebekah Dix

October 18, 2018

Contents

1	Results from Real Analysis	3
2	Solution of equations by iteration	3
2.1	Simple Iteration	3
2.2	Newton's Method	5
2.3	Secant Method	7
3	Solution of systems of linear equations	9
3.1	LU Decomposition	9
3.2	Least Squares	9
3.3	Gram-Schmidt Orthogonalization	10
3.4	QR Factorization	10
3.4.1	Application to Least Squares	11
3.5	Norms and Condition Numbers	11
3.5.1	Conditioning	15
4	Special Matrices	17
4.1	Symmetric Positive Definite Matrices	17
4.2	Cholesky Factorization	19
4.3	Banded Matrices and Differential Equations	19
5	Simultaneous nonlinear equations	21
5.1	Analysis Preliminaries	21
5.2	Simultaneous iteration	21
6	Eigenvalues of Eigenvectors of a symmetric matrix	24
6.1	Why we use iteration to calculate eigenvalues/eigenvectors	24
6.2	Power Iteration	25
6.3	Inverse Iteration	26
6.4	Simultaneous Iteration	26

6.5	Shifted Power Iteration	27
6.6	QR Algorithm	27
6.7	Simultaneous Iteration equivalent to QR Algorithm	27

1 Results from Real Analysis

Theorem 1. (The Mean Value Theorem) Suppose f is a real-valued function, defined and continuous on the closed interval $[a, b] \in \mathbb{R}$ and f differentiable on the open interval (a, b) . Then there exists a number $\xi \in (a, b)$ such that

$$f(b) - f(a) = f'(\xi)(b - a) \quad (1)$$

Theorem 2. (Taylor's Theorem) Suppose that n is a nonnegative integer, and f is a real-valued function, defined and continuous on the closed interval $[a, b]$ of \mathbb{R} , such that the derivatives of f of order up to and including n are defined and continuous on the closed interval $[a, b]$. Suppose further that $f^{(n)}$ is differentiable on the open interval (a, b) . Then, for each value of $x \in [a, b]$, there exists a number $\xi = \xi(x)$ in the open interval (a, b) such that

$$f(x) = f(a) + (x - a)f'(a) + \cdots + \frac{(x - a)^n}{n!}f^{(n)}(a) + \frac{(x - a)^{n+1}}{(n + 1)!}f^{(n+1)}(\xi) \quad (2)$$

2 Solution of equations by iteration

2.1 Simple Iteration

Theorem 3. (Existence of Root) Let f be a real-valued function, defined and continuous on a bounded closed interval $[a, b]$ of the real line. Assume further, that $f(a)f(b) \leq 0$; then, there exists ξ in $[a, b]$ such that $f(\xi) = 0$.

Proof. The condition $f(a)f(b) \leq 0$ implies that $f(a)$ and $f(b)$ have opposite signs, or one of them is 0. If either $f(a)$ or $f(b)$ is 0, then we've found a root. Suppose that both endpoints are non-zero (in which case they have opposite signs). In this case, 0 must belong to the open interval whose endpoints are $f(a)$ and $f(b)$. The intermediate value theorem gives the existence of a root in the open interval (a, b) . Thus, in both cases, a zero is guaranteed. \square

- The converse of Theorem 3 is clearly false.

Theorem 4. (Brouwer's Fixed Point Theorem) Suppose that g is a real-valued function, defined and continuous on a bounded closed interval $[a, b]$ of the real line, and let $g(x) \in [a, b]$ for all $x \in [a, b]$. Then, there exists $\xi \in [a, b]$ such that $\xi = g(\xi)$. ξ is called a fixed point of the function g .

Proof. Define a function $f(x) = x - g(x)$. If we find a root ξ of f , then ξ is a fixed point of g . Then,

$$f(a)f(b) = (a - g(a))(b - g(b)) \leq 0 \quad (3)$$

By assumption, $a \leq g(a), g(b) \leq b$. Therefore, the first term is negative and the second term is positive. Therefore, $f(a)f(b) \leq 0$. By Theorem 3, there exists a $\xi \in [a, b]$ such that $f(\xi) = 0$. Then, for this ξ , $g(\xi) = \xi$. \square

Definition 1. (Simple Iteration) Suppose that g is a real-valued function, defined and continuous on a bounded closed interval $[a, b]$ of the real line, and let $g(x) \in [a, b]$ for all $x \in [a, b]$. Given that $x_0 \in [a, b]$, the recursion defined by

$$x_{k+1} = g(x_k) \quad (4)$$

is called simple iteration; the numbers $x_k, k \geq 0$, are referred to as iterates.

- If this sequence converges, the limit must be a fixed of g , since g is continuous on a closed interval. Note that

$$\xi = \lim_{k \rightarrow \infty} x_{k+1} = \lim_{k \rightarrow \infty} g(x_k) = g\left(\lim_{k \rightarrow \infty} x_k\right) = g(\xi) \quad (5)$$

Definition 2. (Contraction) Let g be a real-valued function, defined and continuous on a bounded closed interval $[a, b]$ of the real line. Then, g is said to be a contraction on $[a, b]$ if there exists a constant L such that $0 < L < 1$ and

$$|g(x) - g(y)| \leq L|x - y| \quad \forall x, y \in [a, b] \quad (6)$$

Theorem 5. (Contraction Mapping Theorem) Suppose that g is a real-valued function, defined and continuous on a bounded closed interval $[a, b]$ of the real line, and let $g(x) \in [a, b]$ for all $x \in [a, b]$. Suppose g is a contraction on $[a, b]$. Then, g has a unique fixed point ξ in the interval $[a, b]$. Moreover, the sequence (x_k) defined by simple iteration converges to ξ as $k \rightarrow \infty$ for any starting value x_0 in $[a, b]$.

Let $\epsilon > 0$ be a certain tolerance, and let $k_0(\epsilon)$ denote the smallest positive integer such that x_k is no more than ϵ away from the fixed point ξ (i.e. $|x_k - \xi| \leq \epsilon$) for all $k \geq k_0(\epsilon)$. Then,

$$k_0(\epsilon) \leq \left\lfloor \frac{\ln|x_1 - x_0| - \ln(\epsilon(1 - L))}{\ln(1/L)} \right\rfloor + 1 \quad (7)$$

Proof. Let $E_k = |x_k - \xi|$ be the error at k . Then

$$\begin{aligned} |x_{k+1} - \xi| &= |g(x_k) - g(\xi)| && \text{(definition of } g \text{ and } \xi \text{ a fixed point)} \\ &< L|x_k - \xi| && \text{(} g \text{ a contraction)} \end{aligned}$$

Therefore by induction

$$E_k \leq L^k E_0 \quad (8)$$

Since $L < 1$, $L^k \rightarrow 0$ as $k \rightarrow \infty$, so that $\lim_{k \rightarrow \infty} |x_k - \xi| = 0$. \square

Theorem 6. (Contraction Mapping Theorem when Differentiable) Suppose that g is a real-valued function, defined and continuous on a bounded closed interval $[a, b]$ of the real line, and let $g(x) \in [a, b]$ for all $x \in [a, b]$. Let $\xi = g(\xi) \in [a, b]$ be a fixed point of g (the existence of this point is guaranteed by Brouwer's fixed point theorem). Assume g has a continuous derivative in some neighborhood of ξ with $|g'(\xi)| < 1$. Then the sequence

(x_k) defined by simple iteration $x_{k+1} = g(x_k)$, $k \geq 0$, converges to ζ as $k \rightarrow \infty$, provided that x_0 is close to ζ .

Definition 3. (Stable, Unstable Fixed Point) Suppose that g is a real-valued function, defined and continuous on a bounded closed interval $[a, b]$ of the real line, and let $g(x) \in [a, b]$ for all $x \in [a, b]$, and let ζ denote a fixed point of g . ζ is a stable fixed point of g if the sequence (x_k) defined by the iteration $x_{k+1} = g(x_k)$, $k \geq 0$, converges to ζ whenever the starting value x_0 is sufficiently close to ζ . Conversely, if no sequence (x_k) defined by this iteration converges to ζ for any starting value x_0 close to ζ , except for $x_0 = \zeta$, then we say that ζ is an unstable fixed point of g .

- With this definition, a fixed point may be neither stable nor unstable.
- If $|g'(\zeta)| < 1$, then ζ is a stable fixed point (provided g is continuous, differentiable etc.)

Theorem 7. (Unstable Fixed Points) Suppose that $\zeta = g(\zeta)$, where the function g has a continuous derivative in some neighborhood of ζ , and let $|g'(\zeta)| > 1$ (thus ζ is an unstable fixed point). Then the sequence (x_k) defined by simple iteration $x_{k+1} = g(x_k)$, $k \geq 0$, does not converge to ζ from any starting value x_0 , $x_0 \neq \zeta$.

Definition 4. (Rate of Convergence) Suppose $\zeta = \lim_{k \rightarrow \infty} x_k$. Define $E_k = |x_k - \zeta|$.

- The sequence (x_k) converges to ζ linearly if there exists a number $\mu \in (0, 1)$ such that

$$\lim_{k \rightarrow \infty} \frac{E_{k+1}}{E_k} = \mu \quad (9)$$

- The sequence (x_k) converges to ζ superlinearly if $\mu = 0$. That is, the sequence of μ_k generated at each step $\rightarrow 0$ as $k \rightarrow \infty$.
- The sequence (x_k) converges to ζ with order q if there exists a $\mu > 0$ such that

$$\lim_{k \rightarrow \infty} \frac{E_{k+1}}{E_k^q} = \mu \quad (10)$$

In particular, if $q = 2$, then the sequence converges quadratically.

2.2 Newton's Method

Definition 5. (Newton's Method) Newton's method for the solution of $f(x) = 0$ is defined by

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)} \quad (11)$$

Geometrically, $(x_{n+1}, 0)$ is the intersection of the x -axis and the tangent of the graph of f at $(x_n, f(x_n))$.

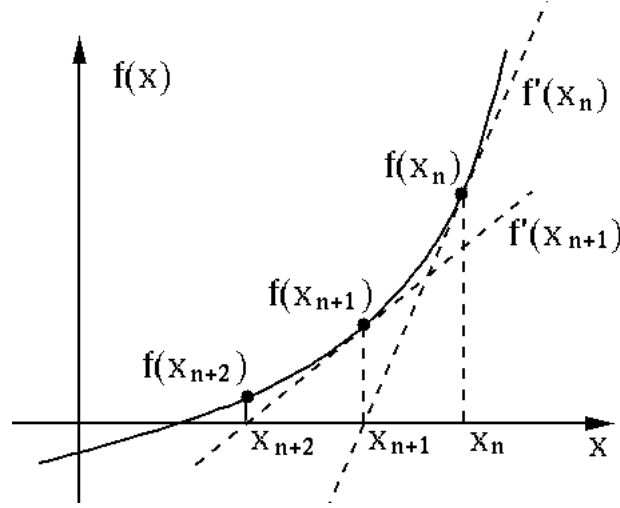


Figure 1: Geometric Interpretation of Newton's Method in \mathbb{R}

Intuitively, the fixed points of this iteration g will be stable.

We can show that $|g'(\xi)| < 1$.

$$\begin{aligned} g'(x) &= 1 - \frac{f' \cdot f' - f \cdot f''}{(f')^2} \\ &= 1 - \left(1 - \frac{f(x) \cdot f''(x)}{(f'(x))^2} \right) \\ &= \frac{f(x) \cdot f''(x)}{(f'(x))^2} \end{aligned}$$

Therefore

$$|g'(\xi)| = \left| \frac{f(\xi) \cdot f''(\xi)}{(f'(\xi))^2} \right| = 0 < 1 \quad (12)$$

Theorem 8. (Convergence of Newton's Method) Suppose that f is a continuous real-valued function with continuous second derivative f'' defined on the closed interval $I_\delta = [\xi - \delta, \xi + \delta]$, $\delta > 0$, such that $f(\xi) = 0$ and $f''(\xi) \neq 0$. Additionally suppose that there exists a positive constant A such that

$$\frac{|f''(x)|}{|f'(y)|} \leq A \quad \forall x, y \in I_\delta \quad (13)$$

If initially

$$|\xi - x_0| \leq h = \min\left(\delta, \frac{1}{A}\right) \quad (14)$$

then the sequence (x_k) defined by Newton's method converges quadratically to ξ .

Proof. We first compute the Taylor expansion of $f(\xi)$, expanding about the point $x_k \in I_\delta$, where $|\xi - x_k| \leq h = \min(\delta, \frac{1}{A})$. Thus

$$f(\xi) = f(x_k) + (\xi - x_k)f'(x_k) + \frac{(\xi - x_k)^2}{2}f''(\eta_k) \quad (15)$$

where η_k is between ξ and x_k . Recall that $f(\xi) = 0$. We can use this fact and the definition of Newton's iteration to rearrange the above expansion as

$$\xi - x_{k+1} = -\frac{(\xi - x_k)^2 f''(\eta_k)}{2f'(x_k)} \quad (16)$$

A small modification to this equation allows us to derive a relationship between adjacent errors

$$E_{k+1} = \frac{f''(\eta_k)}{2f'(x_k)} E_k^2 \quad (17)$$

Recall by assumption we have that $|\xi - x_k| \leq h = \min(\delta, \frac{1}{A})$ and $\frac{|f''(x)|}{|f'(y)|} \leq A \quad \forall x, y \in I_\delta$. Therefore,

$$|E_{k+1}| = \frac{1}{2} \left| \frac{f''(\eta_k)}{f'(x_k)} \right| |E_k|^2 \leq \frac{1}{2} |E_k| \quad (18)$$

We are given that $|\xi - x_0| \leq h = \min(\delta, \frac{1}{A})$, so that induction gives that

$$|E_k| = |\xi - x_k| \leq \frac{1}{2^k} h \quad (19)$$

Therefore (x_k) converges to ξ as $k \rightarrow \infty$.

To show convergence is quadratic, notice that

$$\begin{aligned} \lim_{k \rightarrow \infty} \frac{|E_{k+1}|}{|E_k|} &= \lim_{k \rightarrow \infty} \frac{1}{2} \frac{|f''(\eta_k)|}{|f'(x_k)|} \\ &= \frac{1}{2} \frac{|f''(\xi)|}{|f'(\xi)|} = \mu \leq \frac{A}{2}. \end{aligned}$$

This shows that convergence is quadratic. □

2.3 Secant Method

Observe that Newton's method requires us to know the first derivative f' of f . In applications, we might not know f' or it could be expensive to calculate. This motivates approximating the $f'(x_k)$ in Newton's method with

$$f'(x_k) \approx \frac{f(x_k) - f(x_{k-1})}{x_k - x_{k-1}} \quad (20)$$

Definition 6. (Secant Method) The secant method is defined by

$$x_{k+1} = x_k - f(x_k) \frac{x_k - x_{k-1}}{f(x_k) - f(x_{k-1})} \quad (21)$$

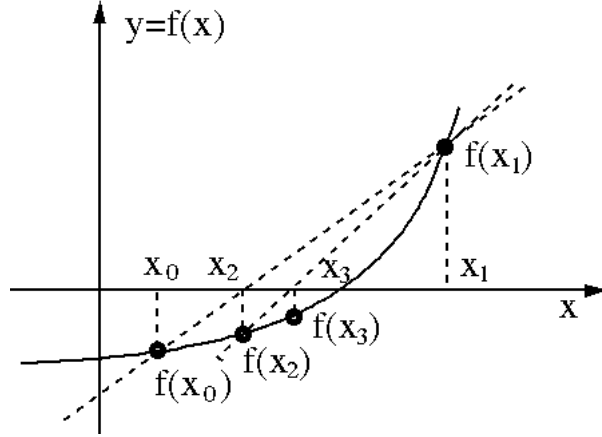


Figure 2: Geometric Interpretation of Secant Method in \mathbb{R}

Theorem 9. (Convergence of Secant Method) Suppose that f is a real-valued function, defined and continuously differentiable on an interval $I = [\zeta - h, \zeta + h]$, $h > 0$, with center point ζ . Suppose further that $f(\zeta) = 0$, $f'(\zeta) \neq 0$. Then, the sequence (x_k) defined by the secant method converges at least linearly to ζ provided that x_0 and x_1 are sufficiently close to ζ .

Proof. Without loss of generality, assume that $\alpha = f'(\zeta) > 0$ in a small neighborhood of ζ . We'll choose this neighborhood such that

$$0 < \frac{3}{4}\alpha < f'(x) < \frac{5}{4}\alpha \quad (22)$$

for all x in the interval.

Recall that the secant method is defined by

$$x_{k+1} = x_k - f(x_k) \frac{x_k - x_{k-1}}{f(x_k) - f(x_{k-1})} \quad (23)$$

We can repeatedly use the mean value theorem to approximate each of these terms. First observe that

$$\frac{f(x_k) - f(\zeta)}{x_k - \zeta} = f'(\eta_k) \quad (24)$$

for some η_k between x_k and ζ . Since $f(\zeta) = 0$, this equation implies that

$$f'(x_k) = f'(\eta_k)(x_k - \zeta) \quad (25)$$

Next observe that

$$\frac{f(x_k) - f(x_{k-1})}{x_k - x_{k-1}} = f'(\theta_k) \quad (26)$$

for some θ_k between x_k and x_{k-1} . Therefore, we can put these pieces together to observe that,

$$x_{k+1} = x_k - \frac{f'(\eta_k)(x_k - \xi)}{f'(\theta_k)} \quad (27)$$

To show convergence, we can compare successive error terms.

$$\begin{aligned} E_{k+1} &= x_{k+1} - \xi \\ &= E_k - \frac{f'(\eta_k)}{f'(\theta_k)} E_k \\ &= \left(1 - \frac{f'(\eta_k)}{f'(\theta_k)}\right) E_k \end{aligned}$$

Therefore

$$\begin{aligned} \frac{E_{k+1}}{E_k} &= 1 - \frac{f'(\eta_k)}{f'(\theta_k)} \\ &< 1 - \frac{5\alpha/4}{3\alpha/4} \\ &= \frac{2}{3} \\ &< 1 \end{aligned}$$

Therefore the secant method converges at least linearly. □

3 Solution of systems of linear equations

3.1 LU Decomposition

3.2 Least Squares

Given a system of equations $Ax = b$, the least squares problem is

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_2^2 \quad (28)$$

We can expand the objective function out as

$$\begin{aligned} \|Ax - b\|_2^2 &= (Ax - b)^T (Ax - b) \\ &= x^T A^T A x - 2b^T A x + b^T b \end{aligned}$$

To find the x that minimizes this expression we find the x that satisfies $\nabla_x F = 0$. That is

$$\nabla_x F = 0 = 2A^T Ax - 2A^T b \quad (29)$$

Therefore the minimizer is $x = (A^T A)^{-1} A^T b$. $(A^T A)^{-1} A^T$ is called the pseudo-inverse of A . If A is square and invertible, then the pseudo-inverse equals A^{-1} .

3.3 Gram-Schmidt Orthogonalization

Algorithm: Denote the columns of A by a_i .

1. $q_1 = a_1$. Then normalized by $q_1 = \frac{q_1}{\|q_1\|}$.
2. $q_2 = a_2 - \langle q_1, a_2 \rangle q_1$. Then normalize by $q_2 = \frac{q_2}{\|q_2\|}$. It's simple to verify that $q_2 \perp q_1$.
3. For an arbitrary k , $q_k = a_k - \langle a_k, q_1 \rangle q_1 - \langle a_k, q_2 \rangle q_2 - \dots - \langle a_k, q_{k-1} \rangle q_{k-1}$. Then normalize by $q_k = \frac{q_k}{\|q_k\|}$.

We can observe the following properties:

1. $\|q_i\| = 1$ (this follows directly)
2. $q_i \perp q_j$ for all $i \neq j$
3. $q_k \in \text{span}(a_1, \dots, a_k)$ and $a_k \in \text{span}(q_1, \dots, q_k)$ so that $\text{span}(a_1, \dots, a_k) = \text{span}(q_1, \dots, q_k)$.

[[Write proof for 2]].

3.4 QR Factorization

Definition 7. (Unitary Matrix) A matrix $Q = [q_1 \dots q_n] \in \mathbb{R}^{m \times n}$ is unitary if and only if $\langle q_i, q_j \rangle = \delta_{ij}$.

Observations about this definition:

1. $Q^T Q = I$
2. If Q is square, then $Q^T = Q^{-1}$.

To calculate the QR decomposition, we can find Q by using the Gram-Schmidt process. Then R can be found as

$$R = \begin{pmatrix} \langle e_1, a_1 \rangle & \langle e_1, a_2 \rangle & \langle e_1, a_3 \rangle & \dots \\ 0 & \langle e_2, a_2 \rangle & \langle e_2, a_3 \rangle & \dots \\ 0 & 0 & \langle e_3, a_3 \rangle & \dots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix} \quad (30)$$

3.4.1 Application to Least Squares

Suppose that we can write $A = QR$, where $A \in \mathbb{R}^{m \times n}$, $Q \in \mathbb{R}^{m \times n}$ and unitary, and $R \in \mathbb{R}^{n \times n}$ and upper triangular. Then the least squares solution to $Ax = b$ is given by

$$\begin{aligned}
 x &= (A^T A)^{-1} A^T b \\
 &= (R^T Q^T Q R)^{-1} R^T Q^T b \\
 &= (R^T R)^{-1} R^T Q^T b \\
 \implies (R^T R)x &= R^T Q^T b \\
 Rx &= Q^T b \quad (\text{assume } R \text{ is invertible (i.e. no zeros on the diagonal)})
 \end{aligned}$$

We can then solve for x using back substitution, which is $\mathcal{O}(n^2)$.

3.5 Norms and Condition Numbers

Definition 8. (Norm) Suppose that \mathcal{V} is a linear space over the field \mathbb{R} . The *nonnegative* real-valued function $\|\cdot\|$ is a norm on \mathcal{V} if the following axioms are satisfied: Fix $v \in \mathcal{V}$

1. Positivity: $\|v\| = 0$ if and only if $v = 0$
2. Scale Preservation: $\|\alpha v\| = |\alpha| \|v\|$ for all $\alpha \in \mathbb{R}$
3. Triangle Inequality: $\|v + w\| \leq \|v\| + \|w\|$.

Example 1. (Examples of Norms)

1. 1-norm:

$$\|v\|_1 = \sum_{i=1}^n |v_i| = |v_1| + \cdots + |v_n| \quad (31)$$

2. 2-norm:

$$\|v\|_2 = \left(\sum_{i=1}^n v_i^2 \right)^{\frac{1}{2}} = \sqrt{v_1^2 + \cdots + v_n^2} = \sqrt{v^T v} \quad (32)$$

3. ∞ -norm

$$\|x\|_\infty = \max_{i=1, \dots, n} |v_i| \quad (33)$$

4. p -norm

$$\|v\|_p = \left(\sum_{i=1}^n |v_i|^p \right)^{\frac{1}{p}} \quad (34)$$

For the p -norm, proving the triangle inequality follows from the Minkowski's inequality.

Definition 9. (Operator Norm) Let A be an $m \times n$ matrix. That is, A is a linear transformation from \mathbb{R}^n to \mathbb{R}^m . Then the operator norm (or subordinate matrix norm) of A is

$$\|A\|_{p,q} = \sup_{x \in \mathbb{R}^n, x \neq 0} \frac{\|Ax\|_q}{\|x\|_p}. \quad (35)$$

Observations about this definition:

1. It's easy to check that this definition of the operator norm satisfies the properties of a norm given in Definition 8. For the triangle inequality, observe that

$$\begin{aligned} \|(A+B)x\|_p &\leq \|Ax\|_p + \|Bx\|_p && \text{(from Minkowski's inequality)} \\ \implies \frac{\|(A+B)x\|_p}{\|x\|_p} &\leq \frac{\|Ax\|_p}{\|x\|_p} + \frac{\|Bx\|_p}{\|x\|_p} \end{aligned}$$

Taking the supremum of both sides over x shows that $\|A+B\|_p \leq \|A\|_p + \|B\|_p$.

2. The definition immediately implies that for an arbitrary $x \in \mathbb{R}^n, x \neq 0$,

$$\|Ax\|_q \leq \|A\|_{p,q} \|x\|_p \quad (36)$$

We can generalize this inequality to claim that

$$\|AB\| \leq \|A\| \|B\| \quad (37)$$

for conformable matrices A, B . Indeed, fix $0 \neq x \in \mathbb{R}^n$. Then

$$\|ABx\| \leq \|A\| \|Bx\| \leq \|A\| \|B\| \|x\| \quad (38)$$

We can divide all inequalities by $\|x\|$ to see that for all $x \neq 0$,

$$\frac{\|ABx\|}{\|x\|} \leq \|A\| \|B\| \quad (39)$$

Taking the supremum over x on the left hand side shows that $\|AB\| \leq \|A\| \|B\|$.

Theorem 10. (The 1-norm of a matrix is the largest absolute-value column sum) Let $A \in \mathbb{R}^{m \times n}$ and denote the columns of A by $a_j, j = 1, \dots, n$. Then $\|A\|_1 = \max_{j=1, \dots, n} \sum_{i=1}^m |a_{ij}| = \max_{j=1, \dots, n} \|a_j\|$.

Proof. Fix $x \in \mathbb{R}^n$. Let $C = \max_{j=1, \dots, n} \sum_{i=1}^m |a_{ij}|$. First consider the product $A \cdot x$. The i th

element is $\sum_{j=1}^n a_{ij}x_j$. Then

$$\begin{aligned}
\|Ax\|_1 &= \sum_{i=1}^m |(Ax)_i| = \sum_{i=1}^m \left| \sum_{j=1}^n a_{ij}x_j \right| \\
&\leq \sum_{i=1}^m \sum_{j=1}^n |a_{ij}| |x_j| && \text{(triangle inequality)} \\
&= \sum_{j=1}^n |x_j| \left(\sum_{i=1}^m |a_{ij}| \right) && \text{(interchange order of summation, assumed finite)} \\
&\leq C \|x\|_1
\end{aligned}$$

Therefore $\frac{\|Ax\|_1}{\|x\|_1} \leq C$ for all x . Next, we find an x such we achieve equality with C . Call index J the index such that $\|a_J\|_1 = C = \max_{j=1,\dots,n} \sum_{i=1}^m |a_{ij}|$. Then let e_J be the n -vector of zeros with a 1 in the J th entry. Clearly $\|e_J\|_1 = 1$. But then

$$\|Ae_J\|_1 = \|a_J\|_1 = C \quad (40)$$

In sum, we first showed that for all $x \in \mathbb{R}^n$

$$\frac{\|Ax\|_1}{\|x\|_1} \leq C \quad (41)$$

We then found an $x \in \mathbb{R}^n$ such that $\frac{\|Ax\|_1}{\|x\|_1} = C$. Therefore

$$\|A\|_1 = \sup_{x \in \mathbb{R}^n, x \neq 0} \frac{\|Ax\|_1}{\|x\|_1} = C = \max_{j=1,\dots,n} \sum_{i=1}^m |a_{ij}| = \max_{j=1,\dots,n} \|a_j\|_1 \quad (42)$$

□

Theorem 11. (The ∞ -norm of a matrix is the largest absolute-value row sum) Let $A \in \mathbb{R}^{m \times n}$ and denote the rows of A by b_i , $i = 1, \dots, m$. Then $\|A\|_\infty = \max_{i=1,\dots,m} \sum_{j=1}^n |a_{ij}| = \max_{i=1,\dots,m} \|b_i\|_1$.

Proof. Fix $x \in \mathbb{R}^n$. Let $C = \max_{i=1,\dots,m} \sum_{j=1}^n |a_{ij}|$.

$$\begin{aligned}
\|Ax\|_\infty &= \max_{i=1,\dots,m} \left| \sum_{j=1}^n a_{ij}x_j \right| \\
&\leq \max_{i=1,\dots,m} \sum_{j=1}^n |a_{ij}| |x_j| && \text{(by the triangle inequality)} \\
&\leq \max_{i=1,\dots,m} \sum_{j=1}^n |a_{ij}| \|x\|_\infty && \text{(since } |x_j| \leq \|x\|_\infty \text{ for all } j) \\
&= C \|x\|_\infty
\end{aligned}$$

Next, we find an x such we achieve equality with C . Call I the index for which $\|b_I\|_\infty = C$. Define

$$x_j = \begin{cases} 1 & a_{Ij} > 0 \\ -1 & a_{Ij} < 0 \end{cases} \quad (43)$$

Observe that $\|x\|_\infty = 1$. Then

$$\begin{aligned}
|A \cdot x|_I &= |b_I^T \cdot x| \\
&= \left| \sum_{j=1}^m a_{Ij}x_j \right| \\
&= \left| \sum_{j=1}^m |a_{Ij}| \right| \\
&= C
\end{aligned}$$

We then found an $x \in \mathbb{R}^n$ such that $\frac{\|Ax\|_\infty}{\|x\|_\infty} = C$. Therefore

$$\|A\|_\infty = \sup_{x \in \mathbb{R}^n, x \neq 0} \frac{\|Ax\|_\infty}{\|x\|_\infty} = C = \max_{i=1,\dots,m} \sum_{j=1}^n |a_{ij}| = \max_{i=1,\dots,m} \|b_i\| \quad (44)$$

□

Theorem 12. (The 2-norm of a symmetric positive definite matrix is the maximum absolute value of its eigenvalues) Let A be a positive definite $n \times n$ matrix. Then

$$\|A\|_2 = \max_{i=1,\dots,n} |\lambda_i| \quad (45)$$

Proof. Since A is positive definite, A has n distinct eigenvalues, which implies that it has n linearly independent eigenvectors. Therefore, for an arbitrary $x \in \mathbb{R}^n$, we can write x

as a linearly combination of the eigenvectors x_1, \dots, x_n . Then

$$\begin{aligned} x &= c_1 x_1 + \dots + c_n x_n \\ Ax &= c_1 A x_1 + \dots + c_n A x_n \\ &= c_1 \lambda_1 x_1 + \dots + c_n \lambda_n x_n \end{aligned}$$

We can normalize the eigenvectors of A so that $x_i^T x_i = 1$. Then $\|Ax\|_2 = \sqrt{\sum_{i=1}^n c_i^2 \lambda_i^2}$ and $\|x\|_2 = \sqrt{\sum_{i=1}^n c_i^2}$. Therefore

$$\frac{\|Ax\|_2}{\|x\|_2} = \sqrt{\frac{\sum_{i=1}^n c_i^2 \lambda_i^2}{\sum_{i=1}^n c_i^2}} \leq \max_i |\lambda_i| = |\lambda_I| \quad (46)$$

Now we'll find an x such that we actually achieve equality. Call I the index of the maximum absolute value of an eigenvalue. Then, consider the eigenvector associated with this eigenvalue, called x_I . Then

$$\frac{\|Ax_I\|_2}{\|x_I\|_2} = \frac{|\lambda_I| \|x_I\|}{\|x_I\|} = |\lambda_I| \quad (47)$$

This shows that $\|A\|_2 = \max_i |\lambda_i|$. □

Theorem 13. (The 2-norm of a matrix $A_{m \times n}$ equals its largest singular value) Let A be an $m \times n$ matrix and denote the eigenvalues of the matrix $B = A^T A$ by $\lambda_i, i = 1, \dots, n$. Then

$$\|A\|_2 = \max_i \sqrt{\lambda_i} \quad (48)$$

The square roots of the (nonnegative) eigenvalues of $A^T A$ are referred to as the singular values of A .

3.5.1 Conditioning

Conditioning helps us quantify the sensitivity of the output to perturbations of the input. In what follows, let f be a mapping from a subset D of a normed linear space \mathcal{V} to another normed linear space \mathcal{W} .

Definition 10. (Absolute Condition Number)

$$\text{Cond}(f) = \sup_{x, y \in D, x \neq y} \frac{\|f(x) - f(y)\|}{\|x - y\|} \quad (49)$$

Definition 11. (Absolute Local Condition Number)

$$\text{Cond}_x(f) = \sup_{x + \delta x \in D, \delta x \neq 0} \frac{\|f(x + \delta x) - f(x)\|}{\|\delta x\|} \quad (50)$$

The previous two definitions depend on the magnitudes of $f(x)$ and x . In applications, it's often better to rescale as follows

Definition 12. (Relative Local Condition Number)

$$\text{cond}_x(f) = \sup_{x+\delta x \in D, \delta x \neq 0} \frac{\|f(x+\delta x) - f(x)\| / \|f(x)\|}{\|\delta x\| / \|x\|} \quad (51)$$

In these definitions, if f is differentiable then we can replace the differences with the appropriate derivatives.

Example 2. (Example of condition numbers) Let D be a subinterval of $[0, \infty)$ and $f(x) = \sqrt{x}$. Then $f'(x) = \frac{1}{2\sqrt{x}}$.

1. If $D = [1, 2]$, then $\text{Cond}(f) = \frac{1}{2}$.
2. If $D = [0, 1]$, then $\text{Cond}(f) = \infty$.
3. If $D = (0, \infty)$, then the absolute local condition number of f at $x \in D$ is

$$\text{Cond}_x(f) = \frac{1}{2\sqrt{x}} \quad (52)$$

Thus as $x \rightarrow 0$, $\text{Cond}_x(f) \rightarrow \infty$, and as $x \rightarrow \infty$, $\text{Cond}_x(f) \rightarrow 0$.

4. If $D = (0, \infty)$, then the relative local condition number of f is $\text{cond}_x(f) = 1/2$ for all $x \in D$.

Definition 13. (Condition Number of a Nonsingular Matrix) The condition number of a nonsingular matrix A is defined by

$$\kappa(A) = \|A\| \|A^{-1}\| \quad (53)$$

If $\kappa(A) \gg 1$, the matrix is said to be ill-conditioned.

Observations about this definition:

1. $\kappa(A) = \kappa(A^{-1})$
2. For all A , $\kappa(A) \geq 1$. This follows because

$$1 = \|I\| = \|AA^{-1}\| \leq \|A\| \|A^{-1}\| \quad (54)$$

3. The condition number of a matrix is unaffected by scaling all its elements by multiplying by a nonzero constant.
4. There is a condition number for each norm, and the size of the condition number is strongly dependent on the choice of norm.

4 Special Matrices

4.1 Symmetric Positive Definite Matrices

Definition 14. (Symmetric, Positive Definite, spd) The real matrix A is said to be symmetric if $A = A^T$. A square $n \times n$ matrix is called positive definite if

$$\mathbf{x}^T A \mathbf{x} > 0 \quad (55)$$

for all $\mathbf{x} \in \mathbb{R}^n$, $\mathbf{x} \neq 0$.

Theorem 14. (Properties of spd matrices) Let A be an $n \times n$ real, spd matrix. Then

1. $a_{ii} > 0$ for all $i = 1, \dots, n$ (the diagonal elements of A are positive).
2. $A\mathbf{x}_i = \lambda_i \mathbf{x}_i \implies \lambda_i \in \mathbb{R}_{>0}, \mathbf{x}_i \in \mathbb{R}^n \setminus \{0\}$ (the eigenvalues of A are real and positive, and the eigenvectors of A belong to $\mathbb{R}^n \setminus \{0\}$).
3. $\mathbf{x}_i \perp \mathbf{x}_j$ if $\lambda_i \neq \lambda_j$ (the eigenvectors of distinct eigenvalues of A are orthogonal)
4. $\det(A) > 0$ (the determinant of A is positive)
5. Every submatrix B of A obtained by deleting any set of rows and the corresponding set of columns from A is symmetric and positive definite (in particular, every principal submatrix is positive definite).

Proof. We prove each claim in the theorem as follows

1. Let \mathbf{e}_i be the i th canonical basis vector in \mathbb{R}^n . Then

$$a_{ii} = \mathbf{e}_i^T A \mathbf{e}_i > 0 \quad (56)$$

since A is pd. A few observations: this only relies on A being pd. $\mathbf{e}_i^T A$ picks out the i th row of A . $A \mathbf{e}_i$ picks out the i th column of A .

2. We'll first show that the eigenvalues of A are real. Suppose λ, \mathbf{x} are an eigenvalue/vector pair of A . Thus $A\mathbf{x} = \lambda\mathbf{x}$. We can conjugate this equation to find that $\bar{A}\bar{\mathbf{x}} = A\bar{\mathbf{x}} = \bar{\lambda}\bar{\mathbf{x}}$ (thus complex eigenvalues of real valued matrices come in conjugate pairs). Then

$$\begin{aligned} \mathbf{x}^T A \bar{\mathbf{x}} &= \bar{\lambda} \mathbf{x}^T \bar{\mathbf{x}} \\ \mathbf{x}^T A^T \bar{\mathbf{x}} &= (A\mathbf{x})^T \bar{\mathbf{x}} = \lambda \mathbf{x}^T \bar{\mathbf{x}} \end{aligned}$$

Since $A = A^T$, we know that $\lambda \mathbf{x}^T \bar{\mathbf{x}} = \bar{\lambda} \mathbf{x}^T \bar{\mathbf{x}}$. As long as $\mathbf{x} \neq 0$, then $\mathbf{x}^T \bar{\mathbf{x}} \neq 0$. Therefore $\bar{\lambda} = \lambda$, which shows $\lambda \in \mathbb{R}$.

The fact that the eigenvector associated with λ has real elements follows by noting that all elements of the singular matrix $A - \lambda I$ are real numbers. Therefore, the

columns of $A - \lambda I$ are linearly dependent in \mathbb{R}^n . Hence there exists an $x \in \mathbb{R}^n$ such that $(A - \lambda I)x = 0$.

This proof only requires that A is symmetric – therefore any real, symmetric matrix has real eigenvalues/vectors.

Next we'll show the eigenvalues of A are positive. Suppose λ, x are an eigenvalue/vector pair of A . Then

$$0 < x^T A x = \lambda x^T x \quad (57)$$

Since $x \neq 0$ and $x^T x$ is positive (it's actually the squared 2-norm of x), then $\lambda > 0$. Note that this part of the proof requires A be pd.

3. Let λ_i, λ_j be distinct eigenvalues of A , and x_i, x_j the corresponding eigenvectors. Then

$$\begin{aligned} x_i^T A x_j &= \lambda_j x_i^T x_j \\ x_i^T A^T x_j &= (A x_i)^T x_j = \lambda_i x_i^T x_j \end{aligned}$$

Since A is symmetric, these two string of equalities must be equal. We can subtract them to find that

$$(\lambda_i - \lambda_j) x_i^T x_j = 0 \quad (58)$$

Since we assumed $\lambda_i \neq \lambda_j$, then it must be that $x_i^T x_j = 0$. Therefore $x_i \perp x_j$. This proof again only relies on the symmetry of A , which is the product of the diagonal elements of the matrix (the eigenvalues).

4. This follows from the fact that the determinant of A is equal to the product of its eigenvalues. Or, we can write A in terms of its eigenvalue decomposition. Thus

$$A = X \Lambda X^{-1} \quad (59)$$

Therefore

$$\det(A) = \det(X) \det(\Lambda) \det(X)^{-1} = \det(\Lambda) \quad (60)$$

Or, we can write A in terms of its eigenvalue decomposition. Thus

$$A = X \Lambda X^{-1} \quad (61)$$

Therefore

$$\det(A) = \det(X) \det(\Lambda) \det(X)^{-1} = \det(\Lambda) \quad (62)$$

5. Let $I \subset \{1, 2, \dots, n\}$ be a subset of indices and let $B = A_{II}$. A is symmetric, so that $A_{II} = A_{II}^T$. Therefore B is symmetric. Let $x \in \mathbb{R}^n$ and define a vector y that is 0 for the indices not included in I and follows the value of x for the indices included in I . Therefore, $x^T B x = y^T A y > 0$ since A is pd.

□

4.2 Cholesky Factorization

Theorem 15. If A is spd, then there exists a lower diagonal matrix L such that $A = LL^T$. This is called the Cholesky decomposition.

Algorithm 1 Cholesky Factorization

Require: $A \in \mathbb{R}^{n \times n}$, SPD

```

 $L_1 \leftarrow \sqrt{a_{11}}$ 
for  $k \leftarrow 2, 3, \dots, n$  do
    Solve  $L_{k-1}l_k = a_k$  for  $l_k$ 
     $l_{kk} \leftarrow \sqrt{a_{kk} - l_k^T l_k}$ 
     $L_k \leftarrow \begin{pmatrix} L_{k-1} & 0 \\ l_k^T & l_{kk} \end{pmatrix}$ 
end for

```

Notation:

- L_{k-1} : the first $k-1 \times k-1$ upper left corner of L
- a_k : the first $k-1$ entries in column k of A
- l_k : the first $k-1$ entries in column k of L^T [[?]]
- a_{kk}, l_{kk} : the kk entries of A and L , respectively

4.3 Banded Matrices and Differential Equations

Consider the two-point boundary value problem

$$u'' + 2u' = -1, \quad u(x=0) = 0, u(x=1) = 0 \quad (63)$$

where $x \in [0, 1]$.

Define a sequence of grid points $\{x_i\}_{i=0}^{N+1}$. We can approximate the derivative of u at each point on the grid as follows

$$\begin{aligned}
 u'(x_j) &= \lim_{\delta \rightarrow 0} \frac{u(x_j + \delta) - u(x_j - \delta)}{2\delta} \\
 &\approx \frac{u(x_{j-1}) - u(x_{j+1})}{2\Delta x}
 \end{aligned}$$

where we use the centered difference quotient of order 2. Similarly, we can approximate

the second derivative of u at each point in the domain as

$$\begin{aligned}
u''(x_j) &= \lim_{\delta \rightarrow 0} \frac{u'(x_j + \delta) - u'(x_j - \delta)}{2\delta} \\
&\approx \frac{u'(x_{j+1}) - u'(x_{j-1}))}{2\Delta x} \\
&= \frac{\frac{u(x_{j+2}) - u(x_j)}{2\Delta x} - \frac{u(x_j) - u(x_{j-2}))}{2\Delta x}}{2\Delta x} \\
&= \frac{u(x_{j+2}) - 2u(x_j) + u(x_{j-2}))}{4\Delta x}
\end{aligned}$$

Let's instead use the grid points adjacent to x_j :

$$u''(x_j) \approx \frac{u(x_{j+1}) - 2u(x_j) + u(x_{j-1}))}{\Delta x} \quad (64)$$

Then, going back to the initial differential equation, for x_j , we have

$$\frac{u(x_{j+1}) - 2u(x_j) + u(x_{j-1}))}{\Delta x} + 2 \frac{u(x_{j-1}) - u(x_{j+1}))}{2\Delta x} = -1 \quad (65)$$

Let

$$U = \begin{bmatrix} u(x_1) \\ u(x_2) \\ \vdots \\ u(x_N) \end{bmatrix} \quad (66)$$

and let $u_i = u(x_i)$. In this notation, the differential equation at x_j can be written as

$$\frac{u_{j+1} - 2u_j + u_{j-1}}{\Delta x^2} + \frac{u_{j-1} - u_{j+1}}{\Delta x} = -1 \quad (67)$$

We can put these equations together into a matrix. Each row will only have 3 non-zero entries at $j - 1$, j , and $j + 1$. Thus the j th row is

$$\left(0 \quad 0 \quad \dots \quad \frac{1}{\Delta x^2} + \frac{1}{\Delta x} \quad \frac{-2}{\Delta x^2} \quad \frac{1}{\Delta x^2} - \frac{1}{\Delta x} \quad 0 \quad 0 \quad \dots \right) \quad (68)$$

Thus stacking these rows together will give a tridiagonal matrix. Call this matrix A . Then we have that

$$AU = -1 \quad (69)$$

5 Simultaneous nonlinear equations

5.1 Analysis Preliminaries

Definition 15. (Cauchy Sequence) A sequence $(x^{(k)}) \subset \mathbb{R}^n$ is called a Cauchy sequence in \mathbb{R}^n if for any $\epsilon > 0$ there exists a positive integer $k_0 = k_0(\epsilon)$ such that

$$\|x^{(k)} - x^{(m)}\|_\infty < \epsilon \quad \forall k, m \geq k_0(\epsilon) \quad (70)$$

Remark 1. \mathbb{R}^n is **complete** in the sense that every Cauchy sequence $(x^{(k)})$ converges to some $\xi \in \mathbb{R}^n$.

Definition 16. (Continuous function) Let $D \subset \mathbb{R}^n$ be nonempty and $f : D \rightarrow \mathbb{R}^n$. Given $\xi \in D$, f is continuous at ξ if for every $\epsilon > 0$, there exists a $\delta = \delta(\epsilon) > 0$ such that for every $x \in B(\xi; \delta) \cap D$

$$\|f(x) - f(\xi)\|_\infty < \epsilon \quad (71)$$

Lemma 1. Let $D \subset \mathbb{R}^n$ be nonempty and $f : D \rightarrow \mathbb{R}^n$ be defined and continuous on D . If $(x^{(k)}) \subset D$ converges in \mathbb{R}^n to $\xi \in D$, then $f(x^{(k)})$ also converges to $f(\xi)$.

We want to find a vector $x \in \mathbb{R}^n$ such that $f(x) = 0$.

Example 3. Consider the linear system

$$Ax = b \quad (72)$$

Then $A : \mathbb{R}^n \rightarrow \mathbb{R}^m$. Let $f(x) = Ax - b$.

Example 4. Let

$$f = \begin{bmatrix} x_1^2 + x_2^2 - 1 \\ 5x_1^2 + 21x_2^2 - 9 \end{bmatrix} \quad (73)$$

Note that $x_1^2 + x_2^2 = 1$ is the 0 level set of f , and is the unit circle. $5x_1^2 + 21x_2^2 = 9$ is the 0 level set of f and is an ellipse.

This function has four zeros

$$\begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} \pm\sqrt{3}/2 \\ \pm\frac{1}{2} \end{bmatrix} \quad (74)$$

5.2 Simultaneous iteration

Example 5.

Definition 17. (Lipschitz condition, constant, and contraction) Let D be a closed subset of \mathbb{R}^n and $g : D \rightarrow D$. If there exists a positive constant L such that

$$\|g(x) - g(y)\|_\infty \leq L\|x - y\|_\infty \quad (75)$$

for all $x, y \in D$, then g satisfies the Lipschitz condition on D in the ∞ -norm. L is called the Lipschitz constant. If $L \in (0, 1)$, then g is called a contraction on D in the ∞ -norm.

Observations about this definition:

- Any function g that satisfies the Lipschitz condition on D is continuous on D (to see this, set $\delta = \frac{\epsilon}{L}$).
- If g satisfies the Lipschitz condition on D in the ∞ -norm, then it also does in the p -norm for $p \in [1, \infty)$ and vice-versa. However the size of L depends on the choice of norm.

Theorem 16. (Contraction Mapping Theorem in \mathbb{R}^n) Suppose D is a closed subset of \mathbb{R}^n and $g : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is defined on D , and $g(D) \subset D$. Suppose further that g is a contraction on D in the ∞ -norm. Then,

1. g has a unique fixed point $\xi \in D$
2. The sequence $(x^{(k)})$ defined by $x^{(k+1)} = g(x^{(k)})$ converges to ξ for any starting value $x^{(0)} \in D$.

Proof. The proof has three parts:

1. First prove uniqueness, assuming existence of a fixed point.
2. Prove the iteration generates a Cauchy sequence (then convergence to some ξ follows from the completeness of the space).
3. Show ξ is indeed the fixed point.

Uniqueness: Suppose ξ, η are both fixed points of g in D . Then,

$$\begin{aligned} \|\xi - \eta\|_\infty &= \|g(\xi) - g(\eta)\| && (\xi, \eta \text{ are fixed points}) \\ &\leq L\|\xi - \eta\|_\infty && (g \text{ is a contraction on } D) \end{aligned}$$

We can rearrange this to see that $(1 - L)\|\xi - \eta\|_\infty \leq 0$. By assumption, $L \in (0, 1)$, and the norm of a quantity is always weakly positive. Therefore, $\|\xi - \eta\|_\infty = 0$ which implies $\xi = \eta$.

Convergence: Assuming g has a fixed point $\xi \in D$, the sequence $x^{(k+1)} = g(x^{(k)})$ will converge to ξ for any $x^{(0)} \in D$. This follows because

$$\|x^{(k)} - \xi\|_\infty \leq L^k \frac{1}{1 - L} \|x^{(1)} - x^{(0)}\|_\infty \quad (76)$$

Since $L \in (0, 1)$, $\lim_{k \rightarrow \infty} L^k = 0$, and therefore

$$\lim_{k \rightarrow \infty} \|x^{(k)} - \xi\|_\infty = 0 \quad (77)$$

Existence: First observe that if $\mathbf{x}^{(0)}$ belongs to D , then $\mathbf{x}^{(k+1)} = g(\mathbf{x}^{(k)}) \in D$ for all $k \geq 1$ since $g(D) \subset D$ (this is important since the proof relies on g being a contraction on D). Next, consider the distance between two adjacent terms on the sequence $\mathbf{x}^{(k+1)} = g(\mathbf{x}^{(k)})$

$$\begin{aligned}\|\mathbf{x}^{(k)} - \mathbf{x}^{(k-1)}\|_\infty &= \|g(\mathbf{x}^{(k-1)}) - g(\mathbf{x}^{(k-2)})\|_\infty && \text{(definition of } g) \\ &\leq L\|\mathbf{x}^{(k-1)} - \mathbf{x}^{(k-2)}\|_\infty && (g \text{ is a contraction on } D) \\ &\leq L^{k-1}\|\mathbf{x}^{(1)} - \mathbf{x}^{(0)}\|_\infty && \text{(induction)}\end{aligned}$$

Now, fix positive integers m, k such that $m > k$. Then

$$\begin{aligned}\|\mathbf{x}^{(m)} - \mathbf{x}^{(k)}\|_\infty &= \|\mathbf{x}^{(m)} - \mathbf{x}^{(m-1)} + \mathbf{x}^{(m-1)} - \mathbf{x}^{(m-2)} + \dots + \mathbf{x}^{(k+1)} - \mathbf{x}^{(k)}\|_\infty \\ &\leq \|\mathbf{x}^{(m)} - \mathbf{x}^{(m-1)}\|_\infty + \dots + \|\mathbf{x}^{(k+1)} - \mathbf{x}^{(k)}\|_\infty && \text{(triangle inequality)} \\ &\leq (L^{m-1} + \dots + L^k)\|\mathbf{x}^{(1)} - \mathbf{x}^{(0)}\|_\infty && (g \text{ a contraction}) \\ &= L^k(L^{m-k-1} + \dots + 1)\|\mathbf{x}^{(1)} - \mathbf{x}^{(0)}\|_\infty \\ &\leq L^k \frac{1}{1-L} \|\mathbf{x}^{(1)} - \mathbf{x}^{(0)}\|_\infty && \text{(geometric series)}\end{aligned}$$

Since $L \in (0, 1)$, $\lim_{k \rightarrow \infty} L^k = 0$. Therefore, $\mathbf{x}^{(k)}$ is a Cauchy sequence in \mathbb{R}^n , that is for all $\epsilon > 0$, there exists a k_0 such that

$$\|\mathbf{x}^{(m)} - \mathbf{x}^{(k)}\|_\infty < \epsilon \quad \forall m, k \geq k_0 \quad (78)$$

Any Cauchy sequence in \mathbb{R}^n is convergent in \mathbb{R}^n . Thus, there exists some $\boldsymbol{\xi} \in \mathbb{R}^n$ such that $\boldsymbol{\xi} = \lim_{k \rightarrow \infty} \mathbf{x}^{(k)}$.

$\boldsymbol{\xi}$ is indeed the fixed point: Since g satisfies the Lipschitz condition on D , g is continuous on D . Therefore,

$$\boldsymbol{\xi} = \lim_{k \rightarrow \infty} \mathbf{x}^{(k+1)} = \lim_{k \rightarrow \infty} g(\mathbf{x}^{(k)}) = g\left(\lim_{k \rightarrow \infty} \mathbf{x}^{(k)}\right) = g(\boldsymbol{\xi}) \quad (79)$$

therefore $\boldsymbol{\xi}$ is a fixed point of g , and observe that $\boldsymbol{\xi} \in D$ since D is closed. \square

Definition 18. (Jacobian) Let $g = (g_1, \dots, g_n)^T : \mathbb{R}^n \rightarrow \mathbb{R}^n$ be a function defined and continuous in an (open) neighborhood of $\boldsymbol{\xi} \in \mathbb{R}^n$. Suppose the first partial derivatives of each g_i exist at $\boldsymbol{\xi}$. The Jacobian matrix $J_g(\boldsymbol{\xi})$ of g at $\boldsymbol{\xi}$ is the $n \times n$ matrix with elements

$$J_g(\boldsymbol{\xi})_{ij} = \frac{\partial g_i}{\partial x_j}(\boldsymbol{\xi}) \quad (80)$$

Theorem 17. Let $g = (g_1, \dots, g_n)^T : \mathbb{R}^n \rightarrow \mathbb{R}^n$ be a function defined and continuous on a closed set $D \subset \mathbb{R}^n$. Let $\boldsymbol{\xi} \in D$ be a fixed point of g . Suppose the first partial derivatives of

each g_i are defined and continuous in some (open) neighborhood $N(\xi) \in D$ of ξ , with

$$\|J_g(\xi)\|_\infty < 1 \quad (81)$$

Then there exists $\epsilon > 0$ such that $g(\bar{B}_\epsilon(\xi)) \subset \bar{B}_\epsilon(\xi)$, and the sequence $x^{(k+1)} = g(x^k)$ converges to ξ for all $x^{(0)} \in \bar{B}_\epsilon(\xi)$ (in other words, the sequence converges to ξ as long as $x^{(0)}$ is close enough to ξ).

Example 6.

Newton's Method

Definition 19. (Newton's Method) The sequence defined by

$$x^{(k+1)} = x^{(k)} - [J_f(x^{(k)})]^{-1}f(x^{(k)}) \quad (82)$$

where $x^{(0)} \in \mathbb{R}^n$, is called Newton's method.

Theorem 18. Suppose $f(\xi) = 0$, that in some (open) neighborhood $N(\xi)$ of ξ , where f is defined and continuous, all the second-order partial derivatives of f are defined and continuous, and that the Jacobian matrix $J_f(x^{(k)})$ of f at the point ξ is nonsingular. Then the sequence defined by Newton's method converges to ξ provided that $x^{(0)}$ is sufficiently close to ξ .

6 Eigenvalues of Eigenvectors of a symmetric matrix

Another matrix decomposition is

$$A = X\Lambda X^{-1} \quad (83)$$

where X is a matrix of the eigenvectors and Λ is a diagonal matrix with the eigenvalues.

6.1 Why we use iteration to calculate eigenvalues/eigenvectors

We call λ an eigenvalue and $x \neq 0$ an eigenvector of A if $Ax = \lambda x$. Thus, $(Ax - \lambda I) = 0$. Therefore, $x \in \text{Null}(A - \lambda I)$. Since $x \neq 0$, $A - \lambda I$ has a non-trivial nullspace, so we must have $\det(A - \lambda I) = 0$. This suggests a way to transform an eigenvalue finding problem to a root finding problem. Define

$$\rho(\lambda) = \det(A - \lambda I) \quad (84)$$

Recall that the determinant of a matrix is the product of its eigenvalues. If A is a $n \times n$ real, symmetric matrix, then $\rho(\lambda)$ is an n -th order polynomial in λ , whose roots are the eigenvalues of A .

Theorem 19. (Abel(-Ruffini) Theorem, or “No-go Theorem”) There is no algebraic solution (that is, a solution expressed in terms of radicals) to general polynomial equations of degree five or higher with arbitrary coefficients.

Therefore, there is no finite-number operation procedure that provides an eigenvalue decomposition.

6.2 Power Iteration

Find the biggest eigenvalue/vector.

Algorithm 2 Power Iteration

Require: $v^{(0)}$ = some vector with $\|v^{(0)}\| = 1$

- 1: **for** $k \leftarrow 1, 2, \dots$ **do**
 - 2: $w \leftarrow Av^{(k-1)}$ ▷ Apply A
 - 3: $v^{(k)} \leftarrow w / \|w\|$ ▷ Normalize
 - 4: $\lambda^{(k)} \leftarrow (v^{(k)})^T Av^{(k)} = \langle v^{(k)}, Av^{(k)} \rangle$ ▷ Rayleigh Quotient
 - 5: **end for**
-

Theorem 20. (Convergence of Power Iteration) Suppose $|\lambda_1| > |\lambda_2| \geq \dots \geq |\lambda_n|$ and $q_1^T v^{(0)} \neq 0$. Then the iterates of power iteration satisfy

$$\|v^{(k)} - (\pm q_1)\| = \mathcal{O} \left(\left| \frac{\lambda_2}{\lambda_1} \right|^k \right) \quad (\text{error of eigenvector})$$

$$|\lambda^{(k)} - \lambda_1| = \mathcal{O} \left(\left| \frac{\lambda_2}{\lambda_1} \right|^{2k} \right) \quad (\text{error of eigenvalue})$$

Proof. Convergence of eigenvector: Write $v^{(0)} = v$ as a linear combination of the orthonormal eigenvectors q_i :

$$v = c_1 q_1 + \dots + c_n q_n \quad (85)$$

$v^{(k)}$ is a scalar multiple of $A^k v^{(0)}$. Therefore

$$\begin{aligned} v^{(k)} &= \alpha_k A^k v^{(0)} && (\alpha_k \text{ a normalization constant}) \\ &= \alpha_k \left(\sum_{i=1}^n \lambda_i^k a_i q_i \right) \\ &= \alpha_k \lambda_1^k \left(c_1 q_1 + c_2 \left(\frac{\lambda_2}{\lambda_1} \right)^k q_2 + \dots + c_n \left(\frac{\lambda_n}{\lambda_1} \right)^k q_n \right) \end{aligned}$$

We can choose α_k such that $\alpha_k \lambda_1^k$ is 1. Therefore, $c_1 q_1$ is dominating (as long as $c_1 \neq 0$). The other terms are of order $\mathcal{O}\left(\left|\frac{\lambda_2}{\lambda_1}\right|^k\right)$.

Convergence of eigenvalue: see proposition below. □

Theorem 21. (Error of Rayleigh Quotient) Let x_1 be the eigenvector that corresponds to the largest (in absolute value) eigenvalue. If $\|x - x_1\| = \mathcal{O}(\epsilon)$, then

$$\left| \frac{\langle x, Ax \rangle}{\langle x, x \rangle} - \lambda_1 \right| = \mathcal{O}(\epsilon^2) \quad (86)$$

Proof. **TODO.** □

6.3 Inverse Iteration

Find the smallest eigenvalue/vector.

6.4 Simultaneous Iteration

Obtain the full set of eigenvalues/vectors simultaneously.

Algorithm 3 Simultaneous Iteration

Require: $Q^{(0)} = V = I$, a list of vectors V , which we choose to be the identity

- 1: **for** $k \leftarrow 1, 2, \dots$ **do**
 - 2: $Z \leftarrow A Q^{(k-1)}$ ▷ Apply A
 - 3: $Z \leftarrow \underline{Q}^{(k)} R^{(k)}$ ▷ QR factorization of Z
 - 4: $A^{(k)} \leftarrow (\underline{Q}^{(k)})^T A Q^{(k)}$ ▷ $A_{ii}^{(k)} = \langle q_i^{(k)}, A q_i^{(k)} \rangle$
 - 5: **end for**
-

Intuitively,

$$A^K \cdot V = [\sum_i \lambda_i^k c_{1i} \tilde{q}_i \mid \sum_i \lambda_i^k c_{2i} \tilde{q}_i \mid \dots] \quad (87)$$

The first column vector will converge to \tilde{q}_1 . The second vector will converge to $\tilde{q}_1 + \mathcal{O}\left(\left|\frac{\lambda_2}{\lambda_1}\right|^k\right) \tilde{q}_2$.

$\underline{Q}^{(k)}$ will converge to the matrix of eigenvectors:

$$X = [x_1 \mid x_2 \mid \dots \mid x_n] \quad (88)$$

$\underline{A}^{(k)}$ will converge to a diagonal matrix containing the eigenvalues.

6.5 Shifted Power Iteration

Find the eigenvalue close to a specific number.

6.6 QR Algorithm

The QR can be viewed as a stable procedure for computing QR factorizations of the matrix powers A, A^2, A^3, \dots

Algorithm 4 QR Algorithm (without shifts)

Require: $A^{(0)} = A$

1: **for** $k \leftarrow 1, 2, \dots$ **do**

2: $Q^{(k)} R^{(k)} \leftarrow A^{(k-1)}$

▷ QR factorization of $A^{(k-1)}$

3: $A^{(k)} \leftarrow R^{(k)} Q^{(k)}$

▷ Recombine factors in reverse order

4: **end for**

6.7 Simultaneous Iteration equivalent to QR Algorithm

The QR algorithm is equivalent to simultaneous iteration applied to a full set of initial vectors, namely, $\hat{Q}^{(0)} = I$. Summary of each algorithm:

Simultaneous Iteration

$$\underline{Q}^{(0)} = I \quad \text{(initial condition)}$$

$$Z = A \underline{Q}^{(k-1)} \quad \text{(apply } A)$$

$$Z = \underline{Q}^{(k)} R^{(k)} \quad \text{(resemblance of normalization, QR factorization of } Z)$$

$$A^{(k)} = (\underline{Q}^{(k)})^T A \underline{Q}^{(k)} \quad \text{(resemblance of Rayleigh quotient)}$$

QR Algorithm

$$A^{(0)} = A \quad \text{(initial condition)}$$

$$A^{(k-1)} = Q^{(k)} R^{(k)} \quad \text{(compute QR factorization)}$$

$$A^{(k)} = R^{(k)} Q^{(k)} \quad \text{(reverse order of factors)}$$

$$\underline{Q}^{(k)} = Q^{(1)} Q^{(2)} \dots Q^{(k)} \quad \text{(definition of } \underline{Q}^{(k)})$$

and

$$\underline{R}^{(k)} = R^{(k)} R^{(k-1)} \dots R^{(1)} \quad \text{(definition of } \underline{R}^{(k)})$$

Theorem 22. (Equivalence of Simultaneous Iteration and the QR Algorithm) Simultaneous Iteration and the QR Algorithm generate identical sequences of matrices $\underline{R}^{(k)}, \underline{Q}^{(k)}, A^{(k)}$. Both give

$$\begin{aligned} (a) : A^{(k)} &= \underline{Q}^{(k)} \underline{R}^{(k)} && \text{(QR factorization of the } k\text{th power of } A) \\ (b) : A^{(k)} &= (\underline{Q}^{(k)})^T A \underline{Q}^{(k)} && \text{(projection)} \end{aligned}$$

Proof. By induction on k (number of iterations). The base case $k = 0$ is trivial.

1. QR gives (a): Assume $A^{(k-1)} = \underline{Q}^{(k-1)} \underline{R}^{(k-1)}$. The inductive hypothesis for (b) gives that $A^{(k-1)} = (\underline{Q}^{(k-1)})^T A \underline{Q}^{(k-1)}$ or that $\underline{Q}^{(k-1)} A^{(k-1)} = A \underline{Q}^{(k-1)}$. Then

$$\begin{aligned} A^{(k)} &= A A^{(k-1)} && \text{(decompose to use inductive hypothesis)} \\ &= A \underline{Q}^{(k-1)} \underline{R}^{(k-1)} && \text{(inductive hypothesis)} \\ &= \underline{Q}^{(k-1)} A^{(k-1)} \underline{R}^{(k-1)} && \text{(inductive hypothesis from (b))} \\ &= \underline{Q}^{(k-1)} R^{(k-1)} \underline{Q}^{(k-1)} \underline{R}^{(k-1)} && \text{(from algorithm)} \\ &= \underline{Q}^{(k)} \underline{R}^{(k)} && \text{(from definitions of } \underline{Q}^{(k)}, \underline{R}^{(k)}) \end{aligned}$$

2. QR gives (b): Assume $A^{(k-1)} = (\underline{Q}^{(k-1)})^T A \underline{Q}^{(k-1)}$. From the relationship $A^{(k-1)} = \underline{Q}^{(k-1)} \underline{R}^{(k-1)}$ and the fact that $\underline{Q}^{(k)}$ is orthogonal, we can apply $(\underline{Q}^{(k)})^T$ to both sides (on the left) to get that $(\underline{Q}^{(k)})^T A^{(k-1)} = \underline{R}^{(k)}$. Then

$$\begin{aligned} A^{(k)} &= R^{(k)} \underline{Q}^{(k)} \\ &= (\underline{Q}^{(k)})^T A^{(k-1)} \underline{Q}^{(k)} \\ &= (\underline{Q}^{(k)})^T (\underline{Q}^{(k-1)})^T A \underline{Q}^{(k-1)} \underline{Q}^{(k)} && \text{(inductive hypothesis)} \\ &= (\underline{Q}^{(k)})^T A \underline{Q}^{(k)} && \text{(definition of } \underline{Q}^{(k)}) \end{aligned}$$

□