

- recap
- impl of multi-threaded code
 - executor service
 - Callable & Future
 - Synchronisation problem

⇒ Synchronization

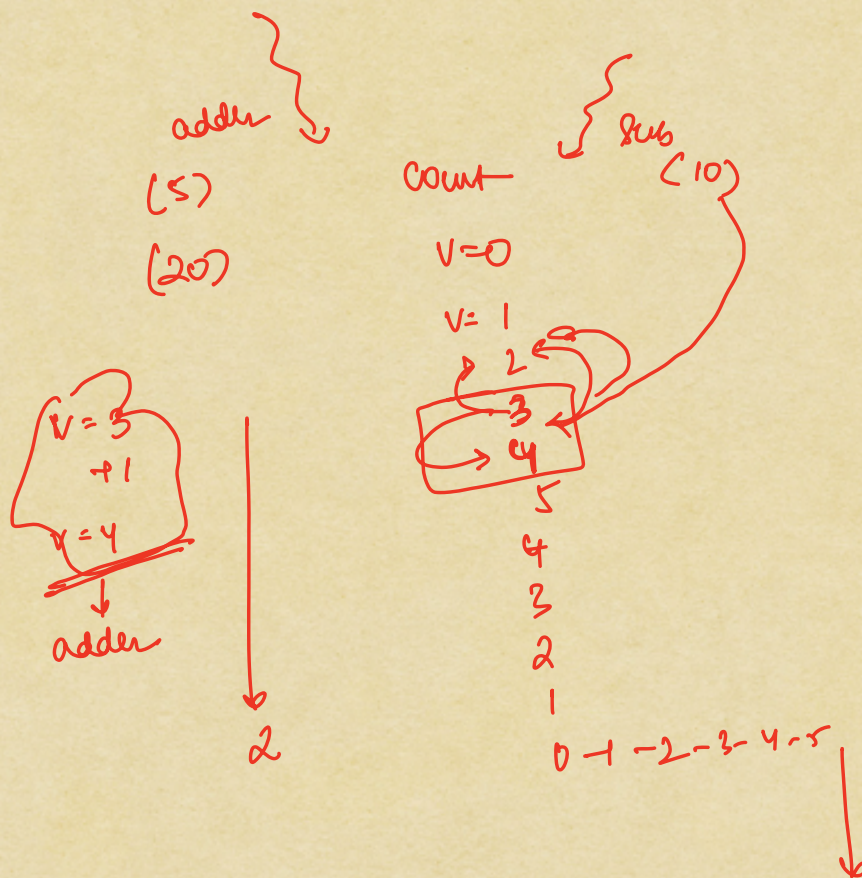
- what
- why
- How we can solve
- mutex
- synchronized keyword
- Semaphore

```
public static void main(String[] args) throws ExecutionException {  
    /*...*/  
    Count c = new Count(value: 0);  
    Adder adder = new Adder(c);  
    Subtractor subtractor = new Subtractor(c);  
    Thread adderThread = new Thread(adder);  
    Thread subtractorThread = new Thread(subtractor);  
    adderThread.start();  
    subtractorThread.start();  
    adderThread.join();  
    subtractorThread.join();  
    System.out.println(c.getValue());  
}
```

main
waits

← makes your current thread wait until the
thread at which join is called
is completed & killed

↘ prints the value inside c. (count)



⇒ Properties of a good solⁿ of synchronisation prob.

Critical section:- part of the code that updates the shared data

```

public void run() {
    → sout("Hello")
    for(i → 1 to 1000) {
        → sout("incr count")
        → c.set(c.get + 1);
        → sout("done")
    }
    → sout("bye")
}

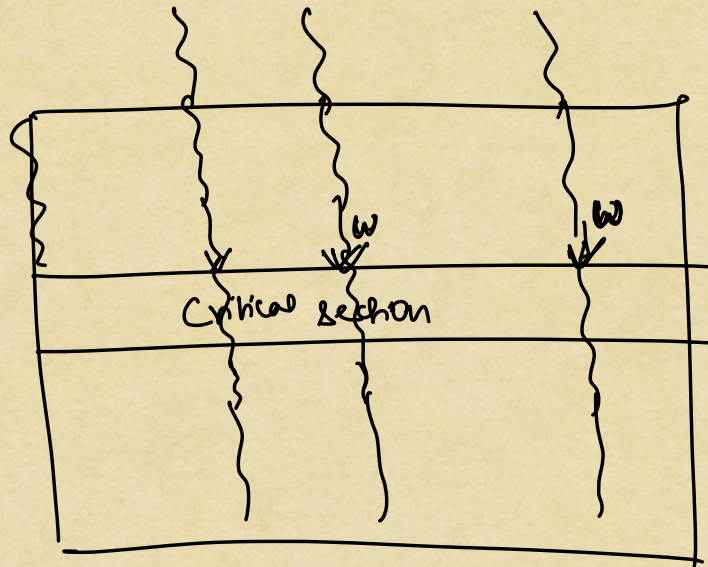
```

critical section

Count
↓
[set value()]

1) Mutually exclusive:- Only 1 thread can execute the critical section at a time.

2) Progress :-



Overall system should keep progressing only critical section area threads should wait.

3) Bounded Waiting:-

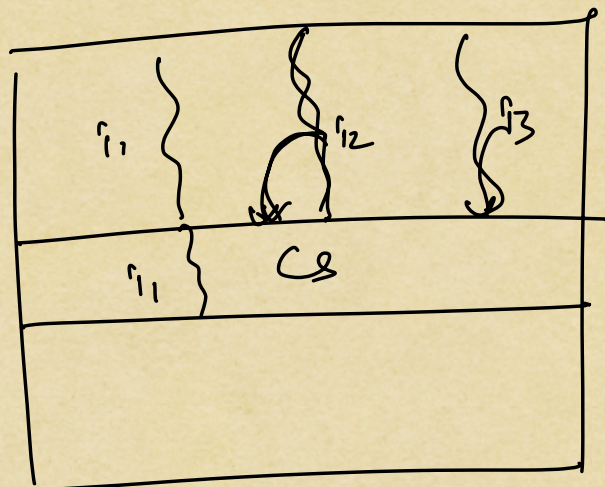
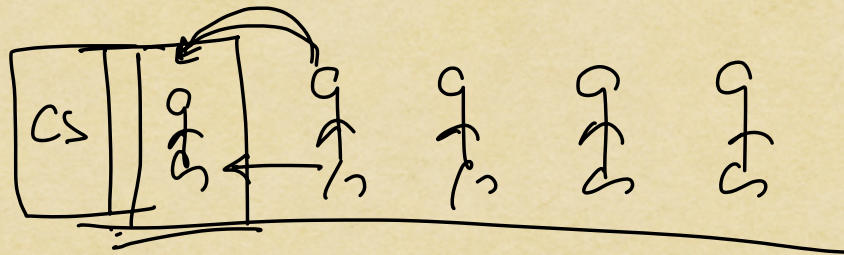
→ 1 thread enters CS and 2 threads are waiting

→ for some reason, thread in CS gets stuck

→ so, we shouldn't have unlimited waiting there

should be bounded waiting

4) NO BUSY WAITING



When a thread is waiting to enter the CS, it should not continuously keep checking if CS is available to enter.

↓
ideal soln is that the thread ^{or something else} that enters the CS should notify other waiting threads.

i) MUTEX ii) Synchronised keyword

⇒ Mutex

⇒ Mutual Exclusion lock | Mutex lock

method() {

 cout("Hello")

1) lock.lock()

2) $x \leftarrow \text{read}(\text{obj})$

3) $x \leftarrow x + 1$

4) $\text{save}(\text{obj}) \leftarrow x$

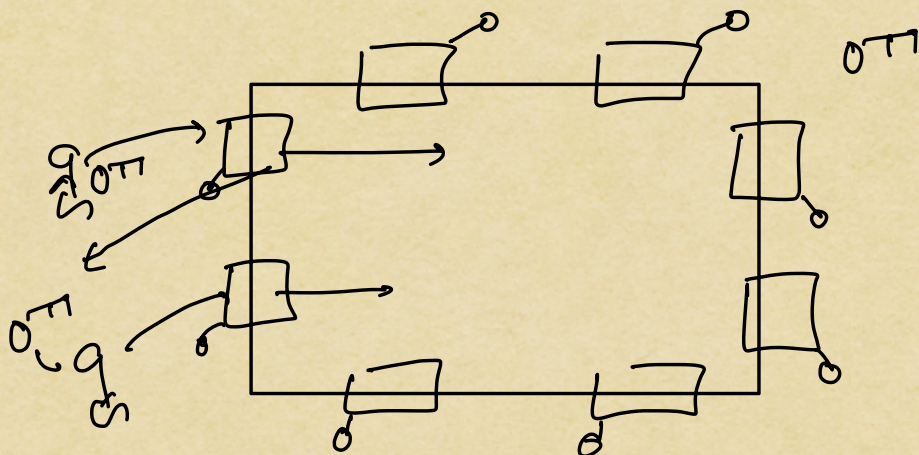
 lock.unlock()

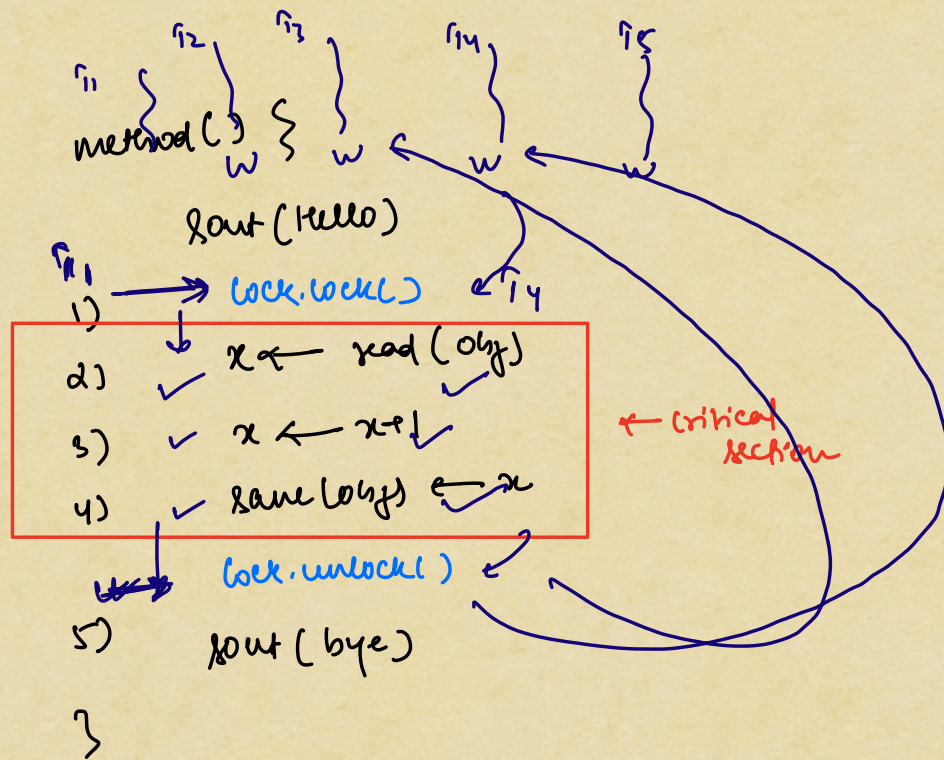
5) cout("bye")

}

← critical section

⇒ multiple locks but single key





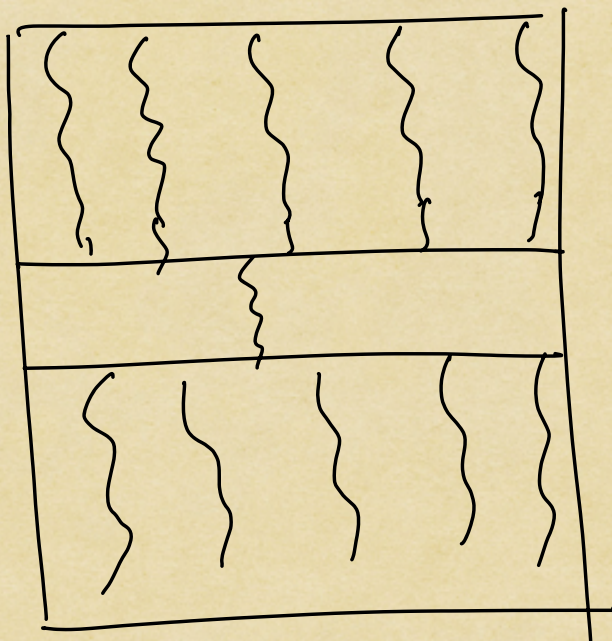
- Mutual exclusion ✓
- progress ✓
- bounded waiting → default wait time ✓
- ✓ → no busy waiting ⇒ unlock() notifies the threads and allows any one of them to enter.

①

Nonces

CS

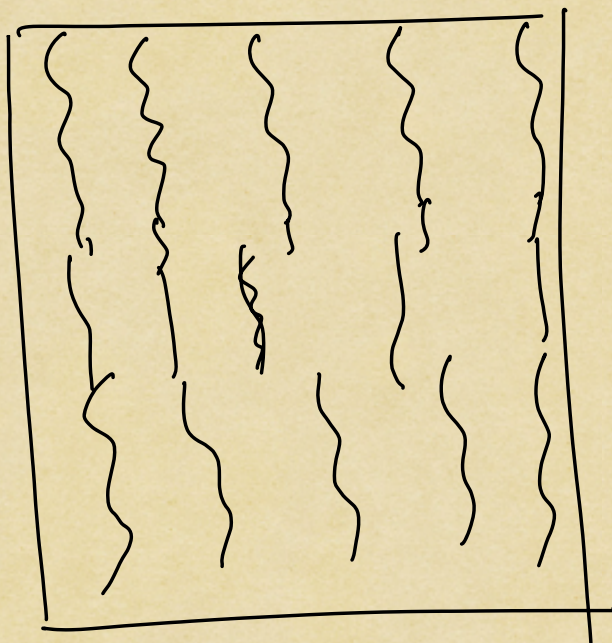
Nonces



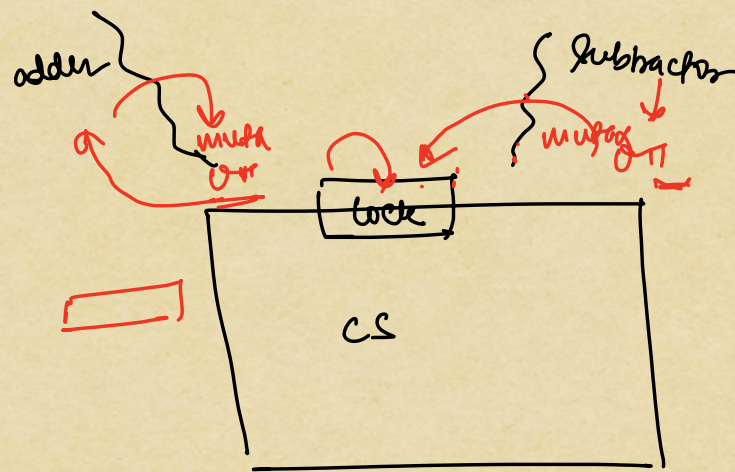
②

Nonces

Nonces



* Solving synchronisation issues reduce performance.



⇒ Synchronised

↙ ↘
 Synchronised Synchronised
 method block-

↓ ↓ ↓ ↓ ↓ ↓ CS

```

public synchronised void doSomething() {
    _____
    _____
    _____
    _____
    _____
    _____
    _____
}
  
```

A red box highlights the body of the `doSomething()` method, and a red arrow labeled 'CS' points to it, indicating that the entire method body is executed within a critical section.

public void doSomething() {

~~CS~~

synchronized(this) {

CS

CS

}