

Theory

Task 4.

The given **for loop** (*) we have converted to the following **while loop** (**), which is considered to be the same. The reason for the conversion is that we already have type-checking rules for the **while loop**:

$$for (int\ i = e_0; e_1; s_1)\ s_2\ (*)$$

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int i; i = e0; while(e1) { s2; s1; }
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$$\frac{\Gamma' = \Gamma, (\ell, i: \text{int}) \quad \frac{\frac{\frac{\Gamma' \vdash_e i: \text{int}}{\Gamma' \vdash_s i = e_0;}}{\Gamma' \vdash_e e_0: \text{int}} \text{ (int literal)} \quad \frac{\frac{\frac{\Gamma' \vdash_e e_1: \text{bool}}{\Gamma' \vdash_s \text{while}(e_1)\{s2; s1;\}} \text{ (bool)} \quad \frac{\frac{\frac{\Gamma' \vdash_{sl} s2; s1;}{\Gamma' \vdash_s \{s2; s1;\}} \text{ (seq)} \quad \frac{\frac{\frac{\Gamma' \vdash_s s2; s1;}{\Gamma' \vdash_{sl} s2; s1;}}{\Gamma' \vdash_s \text{while}(e_1)\{s2; s1;\}} \text{ (block)}}{\Gamma' \vdash_{sl} i = e_0; \text{while}(e_1)\{s2; s1;\}} \text{ (while)}}{\Gamma' \vdash_{sl} \text{int } i; i = e_0; \text{while}(e_1)\{s2; s1;\}} \text{ (seq)} \quad \forall \text{int}:(\ell, i: \text{int}) \notin \Gamma \text{ (var decl)}}{\Gamma' \vdash_{sl} \text{int } i; i = e_0; \text{while}(e_1)\{s2; s1;\}} \text{ (seq)}$$

Following the above derivation, we can state that the formal rule for the expression (*) will look like this:

for $\Gamma' = \Gamma, (\ell, i: \text{int})$

$$\frac{\Gamma' \vdash_e i: \text{int} \quad \Gamma' \vdash_e e_0: \text{int} \quad \Gamma' \vdash_e e_1: \text{bool} \quad \Gamma' \vdash_s s_1 \quad \Gamma' \vdash_s s_2 \quad \forall_{\text{int}}: (\ell, i: \text{int}) \notin \Gamma}{\Gamma \vdash_s \text{for} (\text{int } i = e_0; e_1; s_1) s_2}$$

Task 5.

$$\frac{\frac{\Gamma' = \Gamma, (\ell, y, \text{int}[]) \quad \frac{\frac{\frac{\Gamma' \vdash_e y : \text{int}[]}{} (int\ array) \quad \frac{\Gamma' \vdash_e 2 : \text{int}}{\Gamma' \vdash_e \text{new int}[2] : \text{int}[]} (new\ array) \text{int}[] \prec \text{int}[]}{\Gamma' \vdash_s y = \text{new int}[2];} \quad \frac{\frac{\frac{\Gamma' \vdash_e y : \text{int}[]}{} (int\ array) \quad \frac{\Gamma' \vdash_e 1 : \text{int}}{\Gamma' \vdash_e y[1] : \text{int}} (array\ lookup) \quad \frac{\Gamma' \vdash_e 1 : \text{int}}{} (int\ literal)}{\Gamma' \vdash_s y[1] = 1;} \quad \frac{\frac{\Gamma' \vdash_e y : \text{int}[]}{} (int\ array) \text{int}[] \prec \text{int}[] \quad (\text{return} : \text{int}[]) \epsilon}{\Gamma' \vdash_s \text{return } y;}}{(\text{assign})} \quad \frac{\Gamma' \vdash_{sl} \text{int}[] \text{ } y; y = \text{new int}[2]; y[1] = 1; \text{return } y; \quad \Gamma = (\text{return} : \text{int}[]) \vdash_{sl} \text{int}[] \text{ } y; y = \text{new int}[2]; y[1] = 1; \text{return } y;}{(\text{return})} \quad \frac{\Gamma' \vdash_{sl} \text{int}[] \text{ } y; y = \text{new int}[2]; y[1] = 1; \text{return } y; \quad \Gamma = (\text{return} : \text{int}[]) \vdash_{sl} \text{int}[] \text{ } y; y = \text{new int}[2]; y[1] = 1; \text{return } y;}{(\text{seq } \forall_{\text{int}[]} (\ell, y : \text{int}[]) \notin \Gamma} \quad (\text{var decl})}$$