
An introduction to recursion and induction

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Predefined lists: [*False*, *True*]

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`"..."` normally not shown on slides

Structural induction on lists

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- and for arbitrary x and xs , $P\ xs$ implies $P\ (Cons\ x\ xs)$

A recursive function: append

Definition by *primitive recursion*:

primrec $app :: 'a\ list \Rightarrow 'a\ list \Rightarrow 'a\ list$ **where**
 $app\ Nil\ ys = ?$ |
 $app\ (Cons\ x\ xs)\ ys = ??$

A recursive function: *append*

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1 rule per constructor

Recursive calls must drop the constructor \implies Termination

Demo: append and reverse

Proofs

General schema:

```
lemma name : " . . . "  
apply ( . . . )  
apply ( . . . )  
:  
done
```

If the lemma is suitable as a simplification rule:

```
lemma name[simp] : " . . . "
```

Proof methods

- Structural induction
 - Format: *(induct x)*
x must be a free variable in the first subgoal.
The type of x must be a datatype.
 - Effect: generates 1 new subgoal per constructor
- Simplification and a bit of logic
 - Format: *auto*
 - Effect: tries to solve as many subgoals as possible using simplification and basic logical reasoning.

Top down proofs

sorry

“completes” any proof.

Suitable for top down developments:

Assume lemmas first, prove them later.

Disproving

quickcheck

tries to find counterexample by random testing