

Assignment #2

Problem 2.1: riscv stack frame

a)

1. main: ; function entry label
2. addi sp,sp,-48
 - 48-byte allocation stack frame
 - $sp := sp - 48$
3. sd s0,40(sp)
 - sd = store doubleword (64-bits)
 - save the old frame pointer (s0) into the stack at address $sp + 40$
 - s0 is placed at $sp+40$
4. addi s0,sp,48
 - set up a frame pointer at $s0 := sp + 48$
 - Since sp was decremented earlier by 48, s0 takes its old value
5. mv a5,a0
 - mv is a pseudoinstruction;
 - $mv\ a5,\ a0 = addi\ a5,\ a0,\ 0$
 - copy the incoming a0 (argc) into temporary register a5
6. sd a1,-48(s0)
 - sd = Store Doubleword
 - Store the incoming a1 (argv pointer) at the address $(s0 - 48)$
 - Because $s0 = old_sp$, $s0 - 48 = new_so$, this stores the argv at the bottom of the frame
7. sw a5,-36(s0)
 - sw = store word (32 bits)
 - stores the 32 bit value in a5 (argc) at $s0 - 36$
 - $s0 - 36$ is a slot reserved for local variables / saved arg
8. sw zero,-20(s0)

- store 0 (register) into s0 - 20
 - The local variable 'rc' is initialized to zero in memory
9. lw a5,-20(s0)
- lw = load word
 - load the word from s0 - 20 into the register a5
 - This loads 0 into the variable a5
10. mv a0,a5
- Move the value in a5 into the a0 register
 - a0 is the standard register to hold a functions return value
11. ld s0,40(sp)
- Restore old s0 from the stack at s0 + 40 (64 bits)
12. addi sp,sp,48
- deallocate stack frame $sp := sp + 48$
13. jr ra
- jalr x0, ra, 0
 - jump to the address in ra; return

b)

Higher Addresses			
(above-frame / caller area)			
saved s0 (sd s0,40(sp))	<--- Address = s0 - 8 = sp + 40 (8 bytes)		
(padding / alignment)			
rc (sw zero,-20(s0))	<--- Address = s0 - 20 = sp + 28 (4 bytes)		
argc (sw a5,-36(s0))	<--- Address = s0 - 36 = sp + 12 (4 bytes)		
argv (sd a1,-48(s0))	<--- Address = s0 - 48 = sp + 0 (bottom)		
(pointer, 8 bytes, 64-bit)	<--- Address = sp + 0 (8 bytes)		
Lower addresses			