

# PROGETTO RETI

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#2e3



$$N = \{A, B, C\}$$

$$F = \{1, 2\}; f_1 = f_2 = 2$$

$$C(\emptyset) = 0$$

$$C(\{A\}) = 4$$

$$C(\{B\}) = 3$$

$$C(\{C\}) = 3$$

$$C(\{A, B\}) = 6$$

$$C(\{B, C\}) = 4$$

$$C(\{A, C\}) = 7$$

$$C(N) = 8$$

$$f_3 = 3$$

$$F = \{1, 2, 3\}$$

$$\alpha_A = 4 \quad \alpha_A = 4$$

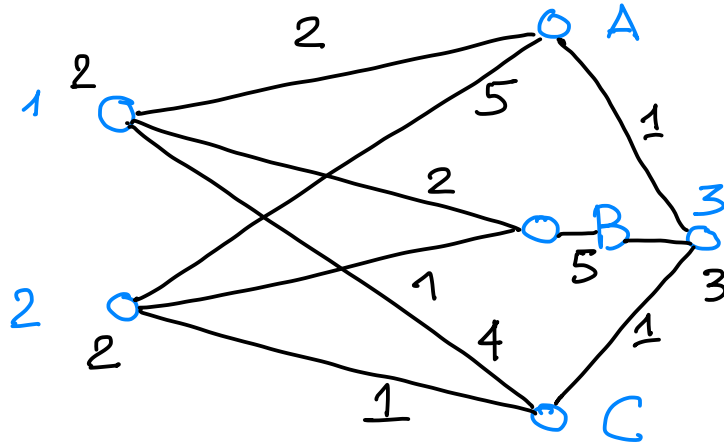
$$\alpha_B = 1 \quad \alpha_B = 2$$

$$\alpha_C = 3 \quad \alpha_C = 2$$

VECTOR COST  
SHARING

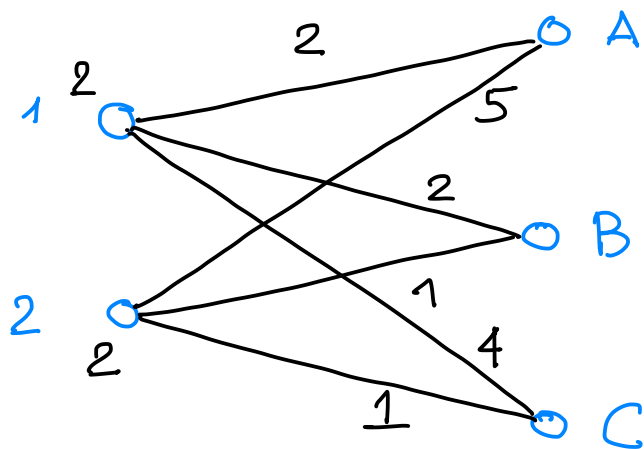
$$3.5 \quad 2.5 \quad 1.5$$

$$\alpha_A + \alpha_B + \alpha_C \leq \frac{15}{2}$$



...  $\left\{ \begin{array}{ll} \alpha_A \leq 4 & 0 \\ \alpha_B \leq 3 & 0 \\ \alpha_C \leq 3 & 0 \\ \alpha_A + \alpha_B \leq 6 & 1/2 \\ \alpha_B + \alpha_C \leq 4 & 1/2 \\ \alpha_A + \alpha_C \leq 5 & 1/2 \\ \alpha_A + \alpha_B + \alpha_C = 8 & \end{array} \right.$

NUCLEO



$$Ax \leq b$$

$$A^T y \geq c$$

$$z^* = \max 4x_1 + x_2 + 5x_3 + 3x_4 = cx$$

$$y_1 \quad 0 \quad 0 \quad x_1 - x_2 - x_3 + 3x_4 \leq 1$$

$$y_2 \quad 1 \quad 5/3 \quad 5x_1 + x_2 + 3x_3 + 8x_4 \leq 55$$

$$y_3 \quad 1 \quad 0 \quad -x_1 + 2x_2 + 3x_3 - 5x_4 \leq 3$$

$$x_1, x_2, x_3, x_4 \geq 0$$

$$\max cx$$

$$Ax \leq b$$

$$\min y_1 + 55y_2 + 3y_3$$

$$\{ \dots A^T y \geq c$$

$$y_1, y_2, y_3 \geq 0$$

$$\frac{25}{3}x_1 + \frac{5}{3}x_2 + 5x_3 + \frac{40}{3}x_4 \leq \frac{275}{3} \rightarrow z^* \leq \frac{275}{3}$$

$$4x_1 + 3x_2 + 6x_3 + 3x_4 \leq 58 \rightarrow z^* \leq 58$$

$$\begin{cases} 1. y_1 + 5y_2 - y_3 \geq 4 \\ -y_1 + y_2 + 2y_3 \geq 1 \\ -y_1 + 3y_2 + 3y_3 \geq 5 \\ 3y_1 + 8y_2 - 5y_3 \geq 3 \end{cases} \rightarrow z^* \leq y_1 + 55y_2 + 3y_3$$

$$y_1, y_2, y_3 \geq 0$$

$\alpha \in \mathbb{R}_+^n$  è nel nucleo  $\gamma$ -~~approssimato~~ ( $\gamma$ -nucleo)  
di un gioco cooperativo  $(N, v)$  se :

$$\sum_{j \in S} \alpha_j \leq c(S) \quad \forall S \subseteq N$$

$$\sum_{j \in N} \alpha_j \geq \gamma \cdot c(N)$$

$$\alpha = (3.5, 2.5, 1.5) \in \frac{15}{16} \text{ nucleo}$$

$$\text{PER DEF } \gamma^* \leq 1$$

TROVARE  $\gamma^* = \max \gamma : \gamma \text{ NUCLEO} \neq \emptyset$

- come individuare un vettore  $x$  che ci permette di recuperare la funzione  $f^*$  di  $C(N)$

$$\begin{array}{ccc} // & // & \equiv \\ & & x \in \gamma^* \text{-NUCLEO} \end{array}$$

TEORIA  
DEI  
GIOCHI

$$Q(N) = \min \sum_{i \in F} f_i \cdot y_i + \sum_{j \in N} \sum_{i \in F} d_{ij} x_{ij}$$

$$\max \sum_{j \in N} x_j \quad \text{QUANTO PAGA } j$$

$$x_j \quad \sum_{i \in F} x_{ij} \geq 1 \quad \forall j \in N$$

$$y_i \quad \sum_{j \in N} B_{ij} \leq f_i \quad i \in F$$

$$B_{ij} \quad y_i - x_{ij} \geq 0 \quad \forall j \in N, i \in F$$

$$x_{ij} \quad x_j \leq B_{ij} + d_{ij} \quad i \in F, j \in N$$

~~$$x_i, x_{ij} \in \{0, 1\}$$~~

CON QUANTO  
CONTRIBUISCE  
J ALL'APERTURA  
DELLA FACILITY i

$$y_i \geq 0$$

$$x_{ij} \geq 0$$

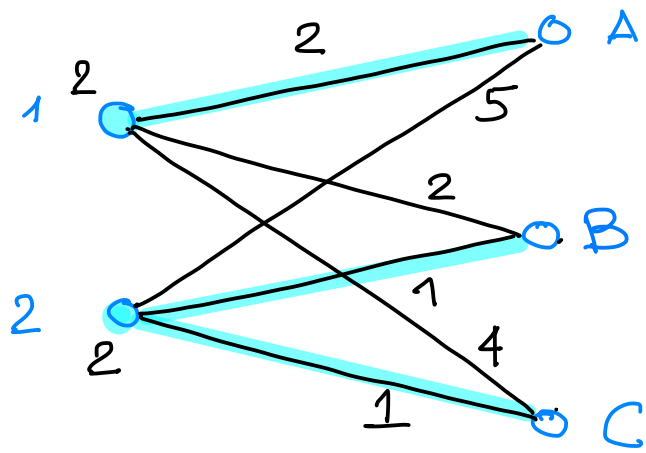
$$x_j \geq 0 \quad j \in N$$

$$B_{ij} \geq 0 \quad j \in N, i \in F$$

$z_{RL}^*$  : valore soluzione rilassamento lineare

$$w_D^* = z_{RL}^* \leq z_I^*$$

$z_I^*$  : valore " " problema intero =  $C(X)$



$$z_I^* = 8$$

$$z_{PL}^* = 8; y_1 = y_2 = 1 \quad \text{else } 0$$

$$x_{1A} = x_{2B} = x_{2C} = 1$$

$$w_D^* = 0; \quad \alpha_A = 4; \alpha_B = 2; \alpha_C = 2$$

$$\beta_{1A} = 2; \beta_{2B} = 1; \beta_{2C} = 1$$

$$\min \sum_{i \in I} f_i \cdot y_i + \sum_{j \in N} \sum_{i \in I} d_{ij} x_{ij}$$

$$\sum_{i \in I} x_{ij} \geq 1 \quad \forall j \in N$$

$$y_i - x_{ij} \geq 0 \quad \forall j \in N, i \in I$$

$$y_i, x_{ij} \geq 0$$

$$\max \sum_{j \in N} \alpha_j$$

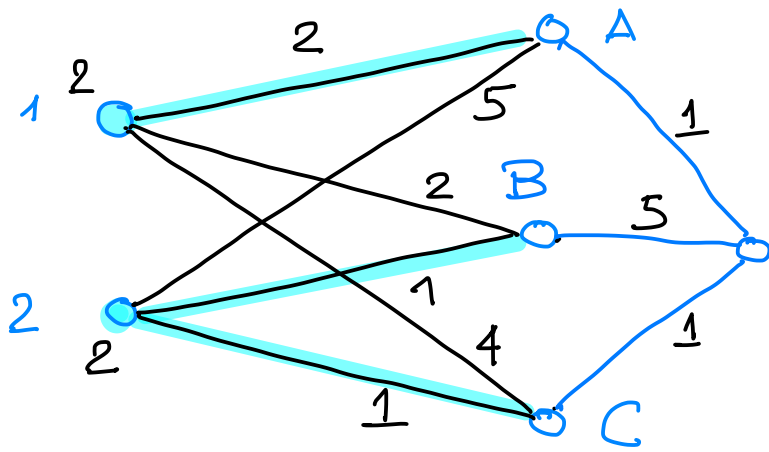
$$\sum_{j \in N} \beta_{ij} \leq f_i \quad i \in I$$

$$\alpha_j \leq \beta_{ij} + d_{ij} \quad i \in I, j \in N$$

$$\alpha_j, \beta_{ij} \geq 0$$







$$\min \sum_{i \in F} f_i \cdot y_i + \sum_{j \in N} \sum_{i \in F} d_{ij} x_{ij}$$

$$\sum_{i \in F} x_{ij} \geq 1 \quad \forall j \in N$$

$$y_i - x_{ij} \geq 0 \quad \forall j \in N, i \in F$$

$$y_i, x_{ij} \geq 0$$

$$z_I^* = 8;$$

$$z_{PL}^* = \frac{15}{2} \quad y_1 = y_2 = y_3 = \frac{1}{2}$$

$$x_{C2} = x_{C3} = \frac{1}{2}$$

$$x_{B1} = x_{B2} = \frac{1}{2}$$

$$x_{A1} = x_{A3} = \frac{1}{2}$$

$$w_D^* = \frac{15}{2}$$

→ DA TROVARE

$$\max \sum_{j \in N} K_j$$

$$\sum_{j \in N} B_{ij} \leq f_i \quad i \in F$$

$$K_j \leq B_{ij} + d_{ij} \quad i \in F, j \in N$$

$$K_j, B_{ij} \geq 0$$