



K. J. Somaiya College of Engineering
(A Constituent College of Somaiya Vidyavihar University)

Batch: A1 Roll No.: 16010120015

Experiment No. ____9____

Grade: AA / AB / BB / BC / CC / CD / DD

Signature of the Staff In-charge with date

Title: Implementation of Graph Colouring Backtracking Algorithm

Objective: To learn the Backtracking strategy of problem solving for Graph Colouring problem

CO to be achieved:

CO 2 Analyze and solve problems for divide and conquer strategy, greedy method, dynamic programming approach and backtracking and branch & bound policies.

Books/ Journals/ Websites referred:

1. Ellis horowitz, Sarataj Sahni, S.Rajsekaran,” Fundamentals of computer algorithm”, University Press
2. T.H.Cormen ,C.E.Leiserson,R.L.Rivest and C.Stein,” Introduction to algortihmts”,2nd Edition ,MIT press/McGraw Hill,2001
3. <http://www.math.utah.edu/~alfeld/queens/queens.html>
4. <http://www-isl.ece.arizona.edu/ece175/assignments275/assignment4a/Solving%208%20queen%20problem.pdf>
5. http://www.slideshare.net/Tech_MX/8-queens-problem-using-back-tracking
6. <http://www.mathcs.emory.edu/~cheung/Courses/170.2010/Syllabus/Backtracking/8queens.html>
7. <http://www.geeksforgeeks.org/backtracking-set-3-n-queen-problem/>
8. <http://www.hbmeyer.de/backtrack/achtdamen/eight.htm>

Pre Lab/ Prior Concepts:

Data structures, Concepts of algorithm analysis

Historical Profile:



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Given an undirected graph and a number m , determine if the graph can be colored with at most m colors such that no two adjacent vertices of the graph are colored with same color. Here coloring of a graph means assignment of colors to all vertices.

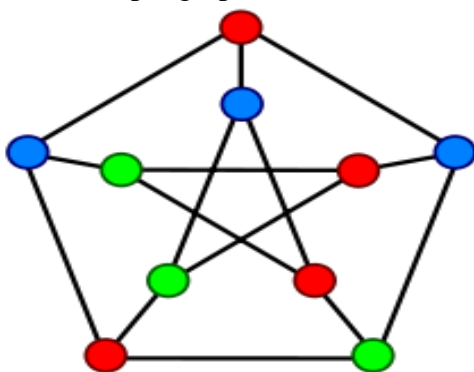
Input:

1) A 2D array $graph[V][V]$ where V is the number of vertices in graph and $graph[V][V]$ is adjacency matrix representation of the graph.

Output:

An array $color[V]$ that should have numbers from 1 to m . $color[i]$ should represent the color assigned to the i th vertex. The code should also return false if the graph cannot be colored with m colors.

Following is an example graph can be colored with 3 colors.



New Concepts to be learned:

Application of algorithmic design strategy to any problem, Backtracking method of problem solving Vs other methods of problem solving problem graph colouring and its applications.

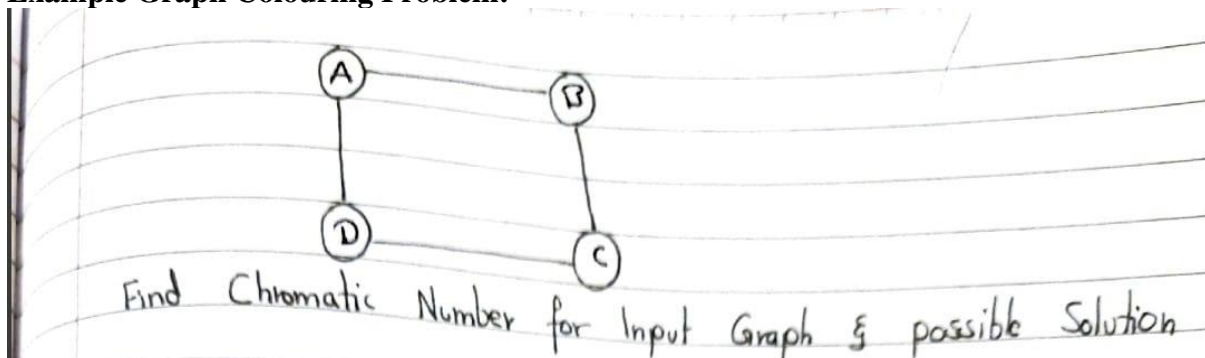
Algorithm Graph colouring Problem:-



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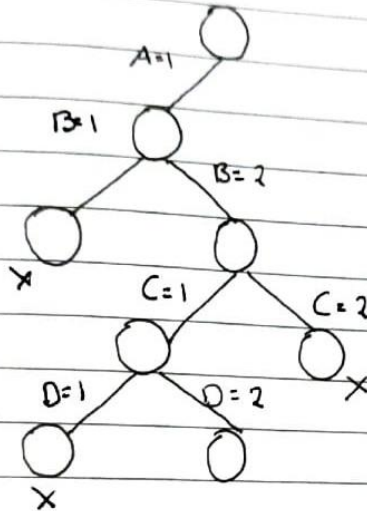
```
1  Algorithm mColoring( $k$ )
2  // This algorithm was formed using the recursive backtracking
3  // schema. The graph is represented by its boolean adjacency
4  // matrix  $G[1 : n, 1 : n]$ . All assignments of  $1, 2, \dots, m$  to the
5  // vertices of the graph such that adjacent vertices are
6  // assigned distinct integers are printed.  $k$  is the index
7  // of the next vertex to color.
8  {
9      repeat
10     { // Generate all legal assignments for  $x[k]$ .
11         NextValue( $k$ ); // Assign to  $x[k]$  a legal color.
12         if ( $x[k] = 0$ ) then return; // No new color possible
13         if ( $k = n$ ) then // At most  $m$  colors have been
14             // used to color the  $n$  vertices.
15             write ( $x[1 : n]$ );
16             else mColoring( $k + 1$ );
17     } until (false);
18 }
```

Example Graph Colouring Problem:



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Solution :-



While Solving the above tree, first Solution comes which says us the minimum no of Solutions Required

\therefore Chromatic Number for Input Graph = 2.

CODE:

```
#include<stdio.h>
int G[50][50],x[50];
void next_color(int k){
    int i,j;
    x[k]=1;
    for(i=0;i<k;i++){
        if(G[i][k]!=0 && x[k]==x[i])
            x[k]=x[i]+1;
    }
}
int main(){
    int n,e,i,j,k,l;
    printf("Enter no. of vertices : ");
    scanf("%d",&n);
    printf("Enter no. of edges : ");
```



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```
scanf("%d",&e);

for(i=0;i<n;i++)
for(j=0;j<n;j++)
G[i][j]=0;

printf("Enter indexes where value is 1-->\n");
for(i=0;i<e;i++){
scanf("%d %d",&k,&l);
G[k][l]=1;
G[l][k]=1;
}

for(i=0;i<n;i++)
next_color(i);
printf("Colors of vertices -->\n");
for(i=0;i<n;i++)
printf("Vertex[%d] : %d\n",i+1,x[i]);
return 0;
}
```

OUTPUT:

 C:\Users\yashg\OneDrive\Desktop\AOA\Expt2\main\bin\Debug\main.exe

```
Enter no. of vertices : 5
Enter no. of edges : 5
Enter indexes where value is 1-->
0 1
1 2
1 3
2 3
3 0
Colors of vertices -->
Vertex[1] : 1
Vertex[2] : 2
Vertex[3] : 1
Vertex[4] : 3
Vertex[5] : 1

Process returned 0 (0x0)   execution time : 36.892 s
Press any key to continue.
_
```



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Analysis of Backtracking solution for Graph Colouring Problem:

If all colors are assigned,
print vertex assigned colors
Else a. Trying all possible colors, assign a color to the vertex
b. If color assignment is possible, recursively assign colors to next vertices
c. If color assignment is not possible, de-assign color, return False

Complexity Analysis:

Time Complexity: $O(m^V)$.

There are total $O(m^V)$ combination of colors.
So time complexity is $O(m^V)$. The upperbound time complexity remains the same but the average time taken will be less.

Space Complexity: $O(V)$.

Conclusion:

- Backtracking provides the hope to solve some problem instances of nontrivial sizes by pruning non-promising branches of the state-space tree.
- The success of backtracking varies from problem to problem and from instance to instance.
- Backtracking possibly generates all possible candidates in an exponentially growing state-space tree.

Hence , We were able to implement C program to implement the Backtracking strategy of problem solving for Graph Colouring problem.