Combinatorics + Lucas + Chinese Remainder Theorem

Combinatorics & Counting Template

```
const int N = 1e6 + 9, oo = 1e17, MOD = 1e9+7;
// f[i] = i!
// inv[i] = modular inverse of i
// finv[i] = modular inverse of i!
int f[N], inv[N], finv[N];
// Precompute factorials, inverse of numbers, and inverse factorials
void prec() {
   f[0] = 1;
   for (int i = 1; i < N; i++)
        f[i] = 1LL * i * f[i - 1] % MOD;
    inv[1] = 1;
   for (int i = 2; i < N; i++) {
        // Fermat's trick to compute modular inverse of i
        inv[i] = (-(1LL * MOD / i) * inv[MOD % i]) % MOD;
        inv[i] = (inv[i] + MOD) % MOD; // make sure it's positive
    }
   finv[0] = 1;
   for (int i = 1; i < N; i++)
        finv[i] = 1LL * inv[i] * finv[i - 1] % MOD;
}
// nCr = number of ways to choose r items from n (order doesn't matter)
int ncr(int n, int r) {
   if (n < r || n < 0 || r < 0) return 0;
    return 1LL * f[n] * finv[n - r] % MOD * finv[r] % MOD;
}
// nPr = number of ways to arrange r items from n (order matters)
int npr(int n, int r) {
   if (n < r || n < 0 || r < 0)
        return 0;
    return 1LL * f[n] * finv[n - r] % MOD;
}
// Fast exponentiation: computes base^pow % MOD in O(log pow)
int fastPower(int base, int pow){
   base %= MOD;
    int res = 1;
    while (pow > 0) {
        if (pow & 1)
            res = res * base % MOD;
        base = base * base % MOD;
        pow >>= 1;
}
```

Extra Notes

```
1. When to use nPr (permutations):

If order matters

Keywords: "arrange", "sequence", "line up", "in order"

2. When to use nCr (combinations):

If order doesn't matter

Keywords: "choose", "select", "group of"
```

3. Stars and Bars:
Used when distributing items into groups
Example: distribute 10 candies among 3 kids
Formula: Number of integer solutions for: x1 + x2 + ... + xk = n
→ Answer = C(n + k - 1, k - 1)
4. Pigeonhole Principle:
If you have more items than containers,
→ at least one container has more than one item.
Example: 13 people in 12 months → 2 people share a birthday month.

Lucas's Theorem

 mod کود Lucas's Theorem (لما P - (P \leq N لما

```
const int MOD = 1e6 + 3; // اولي وصغير P لازم يكون // (1e6 + 3)
int fact[MOD];
// Precompute factorials mod P
void init_factorials() {
   fact[0] = 1;
   for (int i = 1; i < MOD; i++)
        fact[i] = 1LL * fact[i - 1] * i % MOD;
}
// Fast power mod P
int power(int base, int exp) {
    int res = 1;
    while (exp > 0) {
        if (exp & 1) res = 1LL * res * base % MOD;
        base = 1LL * base * base % MOD;
        exp >>= 1;
   }
    return res;
// Modular inverse using Fermat's Little Theorem
int mod_inverse(int a) {
    return power(a, MOD - 2);
}
// C(n, k) % P when n, k < P
int nCr_mod_p_small(int n, int k) {
    if (k > n) return 0;
    return 1LL * fact[n] * mod_inverse(fact[k]) % MOD * mod_inverse(fact[n - k]) % MOD;
// Lucas's Theorem: C(n, k) % P for any n, k
int lucas(int n, int k) {
   if (k == 0) return 1;
    return 1LL * lucas(n / MOD, k / MOD) * nCr_mod_p_small(n % MOD, k % MOD) % MOD;
```

Notes

```
• lucas(n, k) بترجع (n, k) % P و n مهما كان حجم k
```

- Prime يكونLight Prime
- بتقسم n و k لـ أرقام صغيرة أقل من P
- بتشتغل کویس حتی لما ۱e۱8 م لو P صغیر

```
typedef long long 11;
11 mulmod(l1 a, l1 b, l1 mod) {
    return a * b % mod;
}
ll power(ll a, ll b, ll mod) {
    11 \text{ res} = 1;
    a %= mod;
    while (b > 0) {
       if (b & 1) res = res * a % mod;
       a = a * a % mod;
        b >>= 1;
    }
    return res;
}
11 modInverse(ll a, ll mod) {
    return power(a, mod - 2, mod); // mod must be prime
}
// compute nCr % p for small n < p</pre>
11 C_mod_p(ll n, ll r, ll p, vector<ll>& fact) {
    if (r > n) return 0;
    return fact[n] * modInverse(fact[r], p) % p * modInverse(fact[n - r], p) % p;
}
// Lucas Theorem for a single prime p
11 Lucas(ll n, ll r, ll p) {
    vector<ll> fact(p, 1);
    for (11 i = 1; i < p; ++i)
        fact[i] = fact[i - 1] * i % p;
    11 result = 1;
    while (n \mid \mid r) {
        11 \text{ ni} = n \% p, ri = r \% p;
        result = result * C_mod_p(ni, ri, p, fact) % p;
        n /= p, r /= p;
    }
    return result;
}
// Chinese Remainder Theorem: combine values mod p1, p2, ...
11 CRT(const vector<11>& rem, const vector<11>& mod) {
    11 \text{ prod} = 1, \text{ res} = 0;
    int k = rem.size();
    for (int i = 0; i < k; ++i)
        prod *= mod[i];
    for (int i = 0; i < k; ++i) {
        11 pp = prod / mod[i];
        res = (res + rem[i] * modInverse(pp, mod[i]) % prod * pp % prod) % prod;
    return res;
// Main function: Extended Lucas for mod M = p1 * p2 * ...
11 extendedLucas(11 n, 11 r, const vector<11>& primes) {
    vector<ll> rem;
    for (ll p : primes)
        rem.push_back(Lucas(n, r, p)); // result mod each pi
    return CRT(rem, primes);
                                         // combine using CRT
}
```

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