CSCI 476: Computer Security

Computer Systems Review

Reese Pearsall

Fall 2024

Announcements

TA

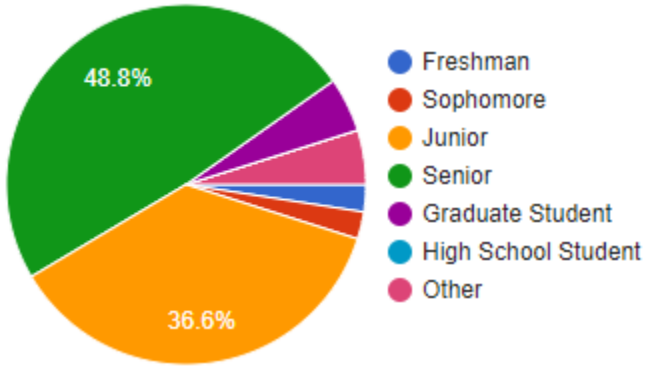
- **Nishu Nath**
- nishu.nath@student.montana.edu
- Office Hours: Tuesdays 3pm – 5pm
- Location: Barnard 259

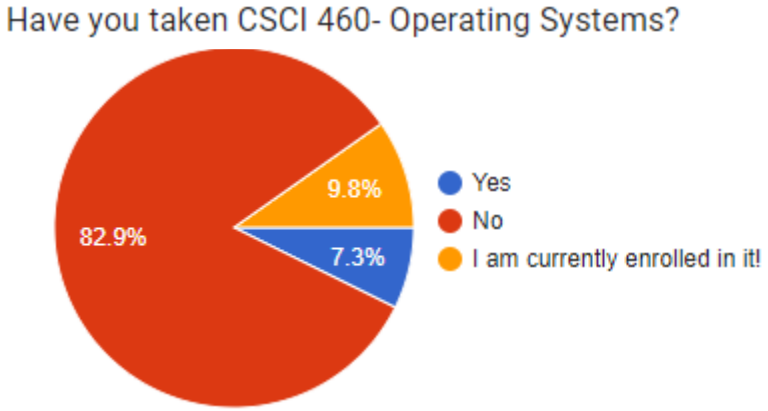
Lab 0 posted→ Due **Sunday** September 8th

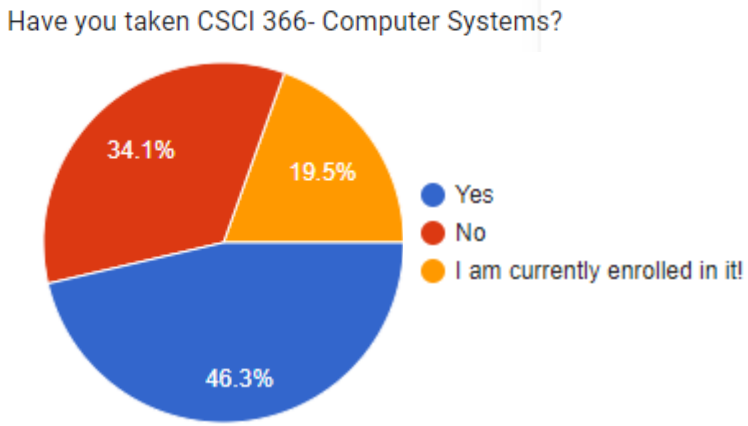
- I want to make sure *everyone* can get something working

Stuff this week is posted in a Brightspace announcement (course website is broken atm)

Course Questionnaire







Most people are comfortable with C and the command line, which is good ☺

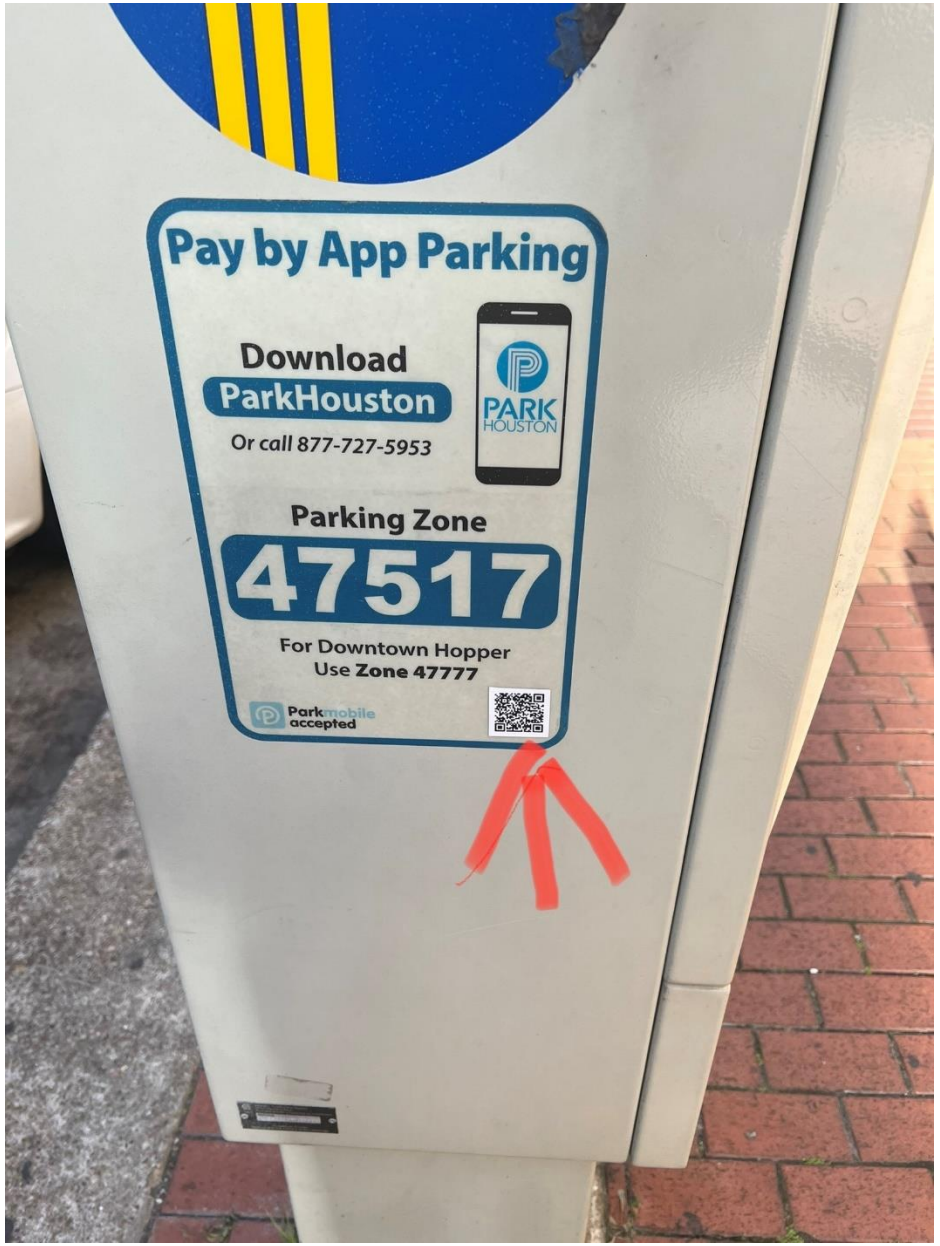
"This class seems relevant and interesting"

"I am Stressed"

"I am working full-time"

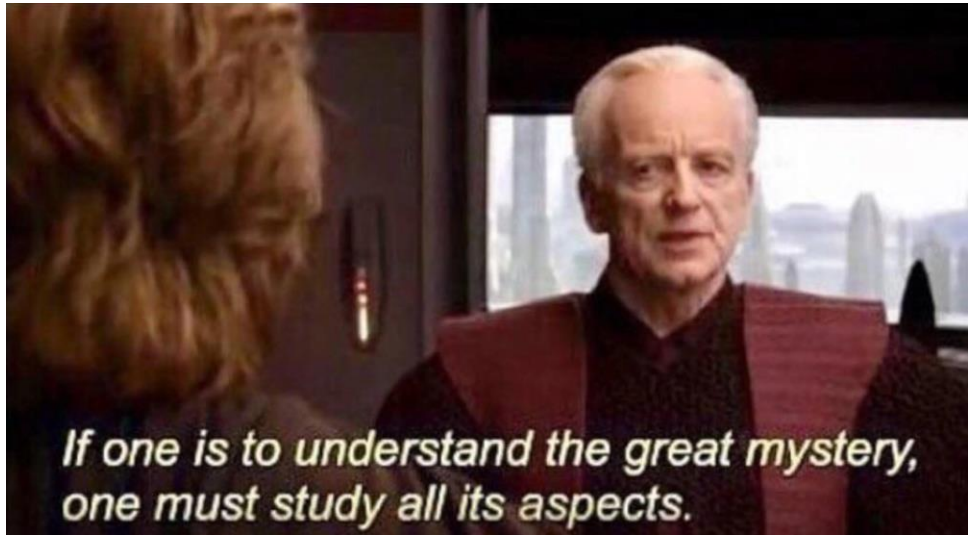
In the news: QR Codes

- Malicious Actors are placing QR code stickers near public parking structures and tricking people to enter credit card information or download malware
    - URLs are embedded in QR codes, and by clicking on a URL, that is enough to initiate a download or send them to a fake website
    - QR codes are used for many things (some good things), but be careful when scanning an unknown QR code.



Computer Systems Review

To understand the technical aspects of security, we must have a good understanding of how computers work



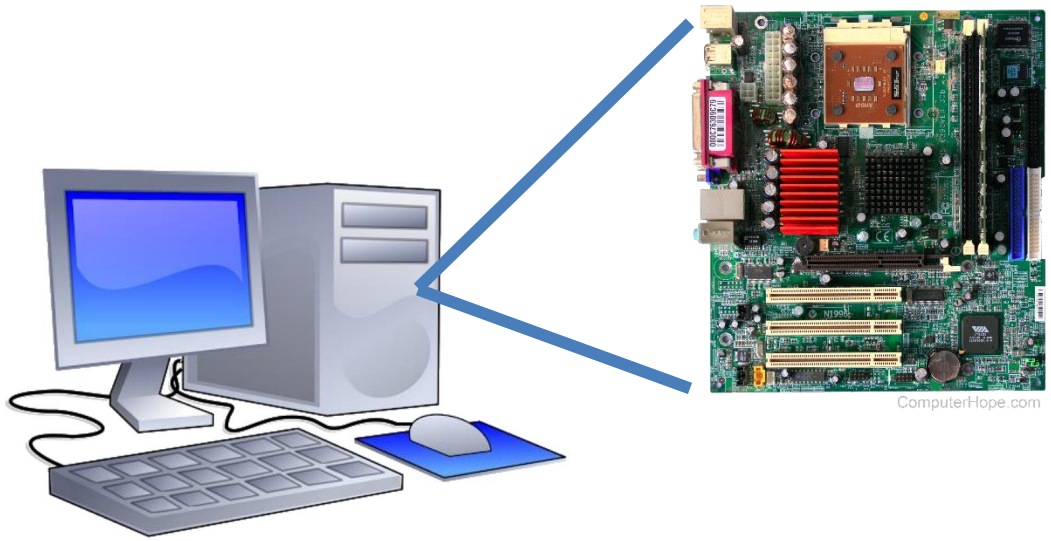
What is a computer?

What is a computer?

A magical box that does stuff

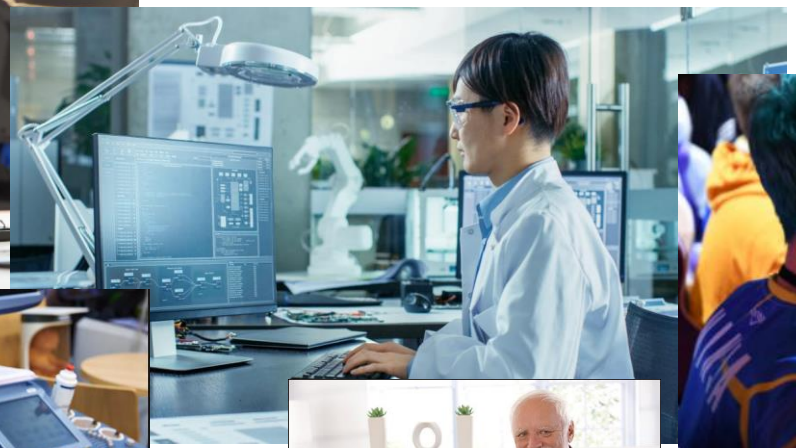
What is a computer?

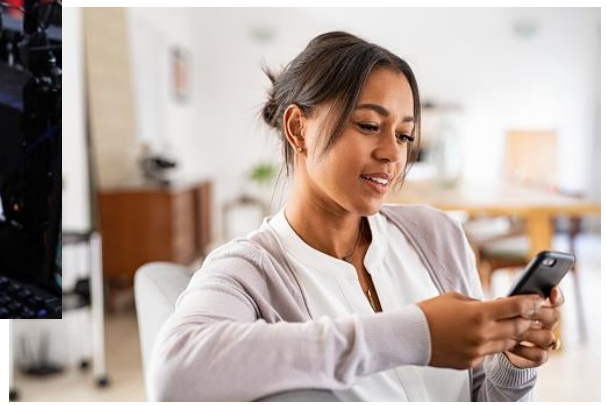
A **semi-**magical box that ~~does stuff~~ **executes instructions**



What is so magical about a computer?

We use computers every day for many different things





What is so magical about a computer?

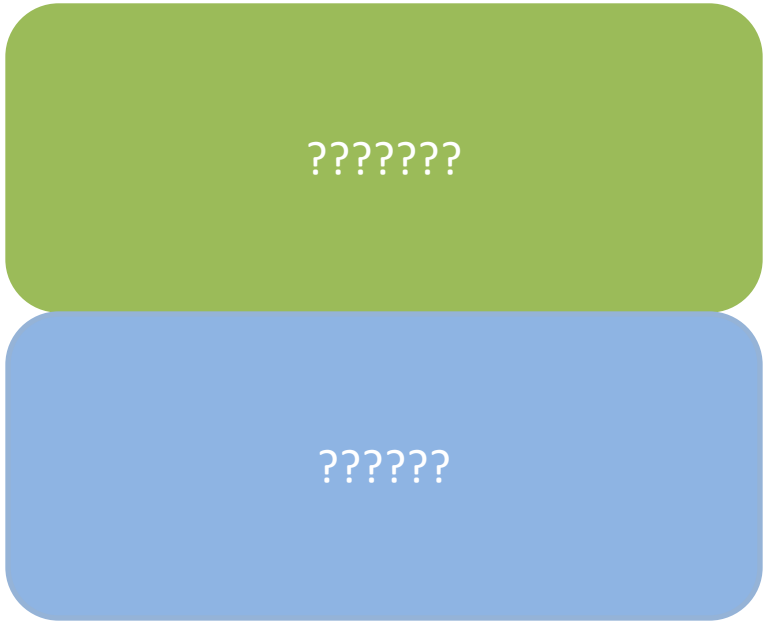
Big Idea

Computers only understand instructions in the form of 0s and 1s (binary)

~~Welcome to CSCI 476~~ 

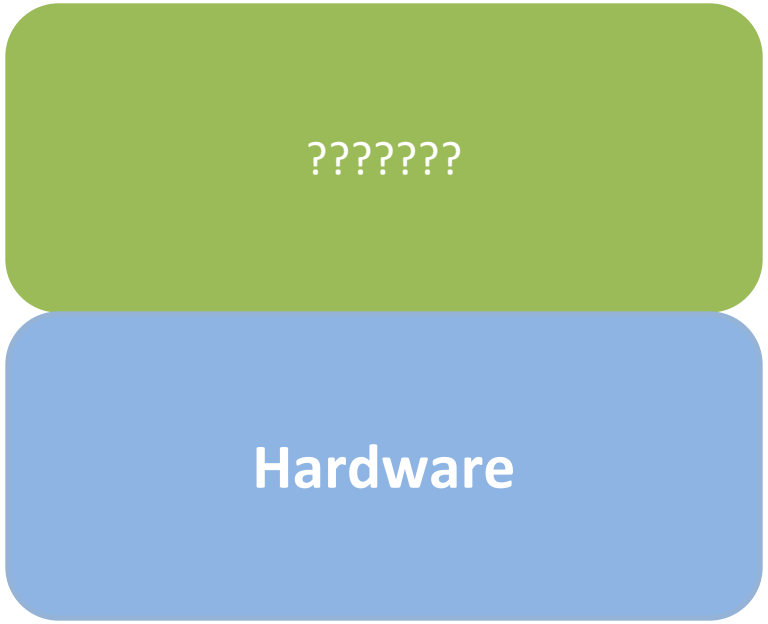
01010111 01100101 01101100 01100011 01101111
01101101 01100101 00100000 01110100 01101111
00100000 01000011 01010011 01000011 01001001
00100000 00110100 00110111 00110110

How does this happen?



From a high level, we will divide a computer system into two parts

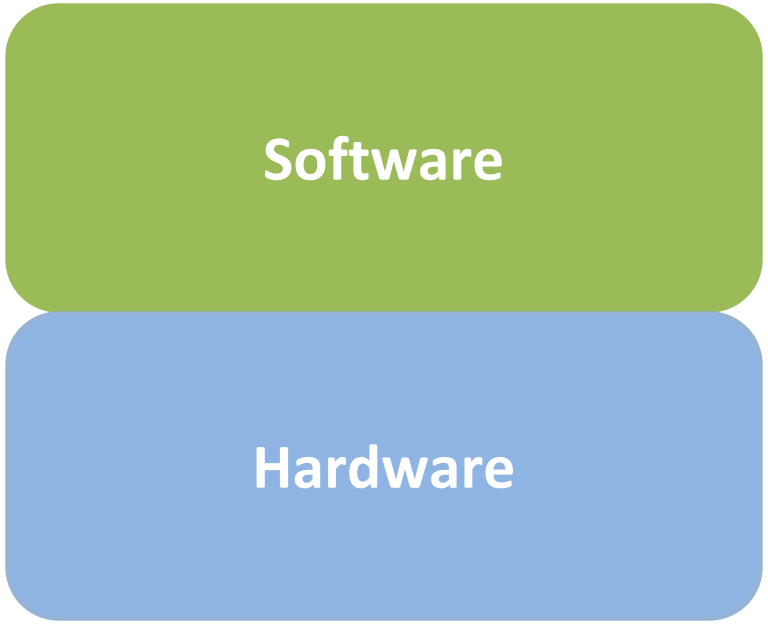
How does this happen?



From a high level, we will divide a computer system into two parts

I. Hardware

How does this happen?



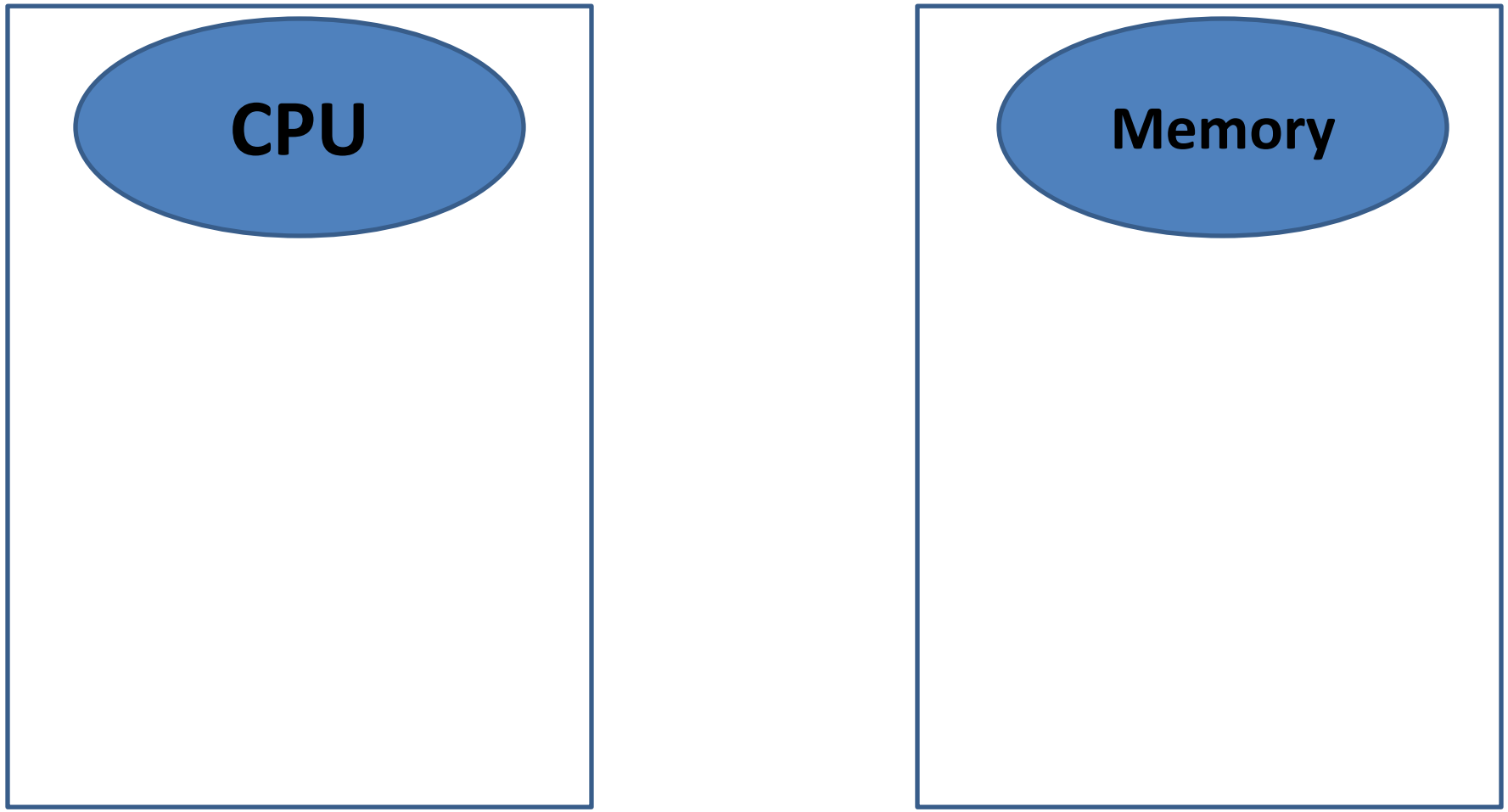
From a high level, we will divide a computer system into two parts

I. **Hardware**
II. **Software**

I. Hardware

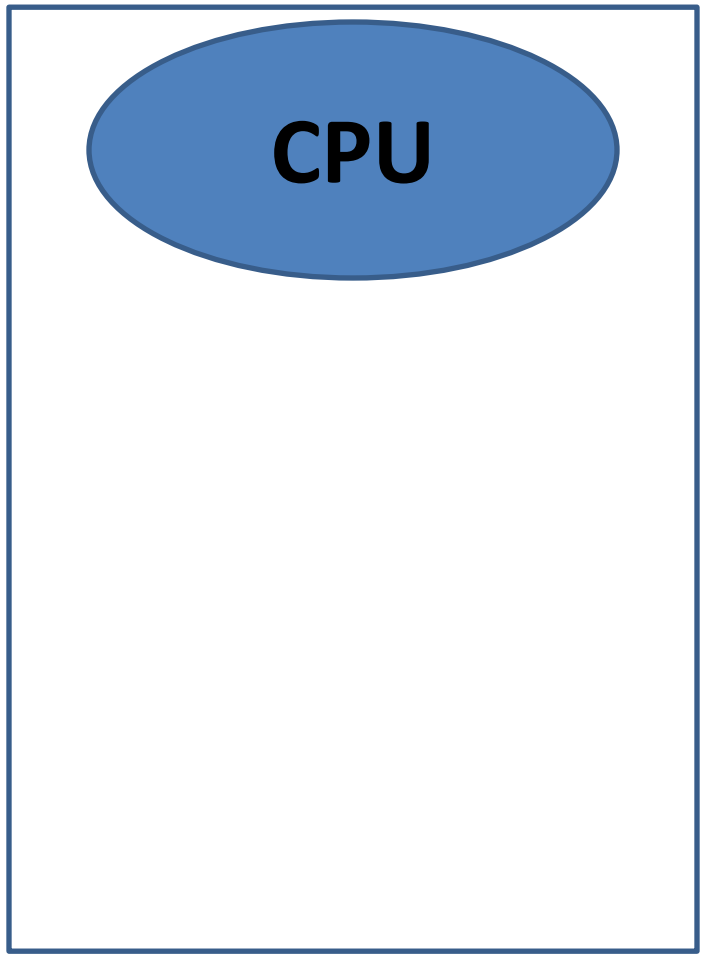
The **physical** parts of a computer

I. Hardware

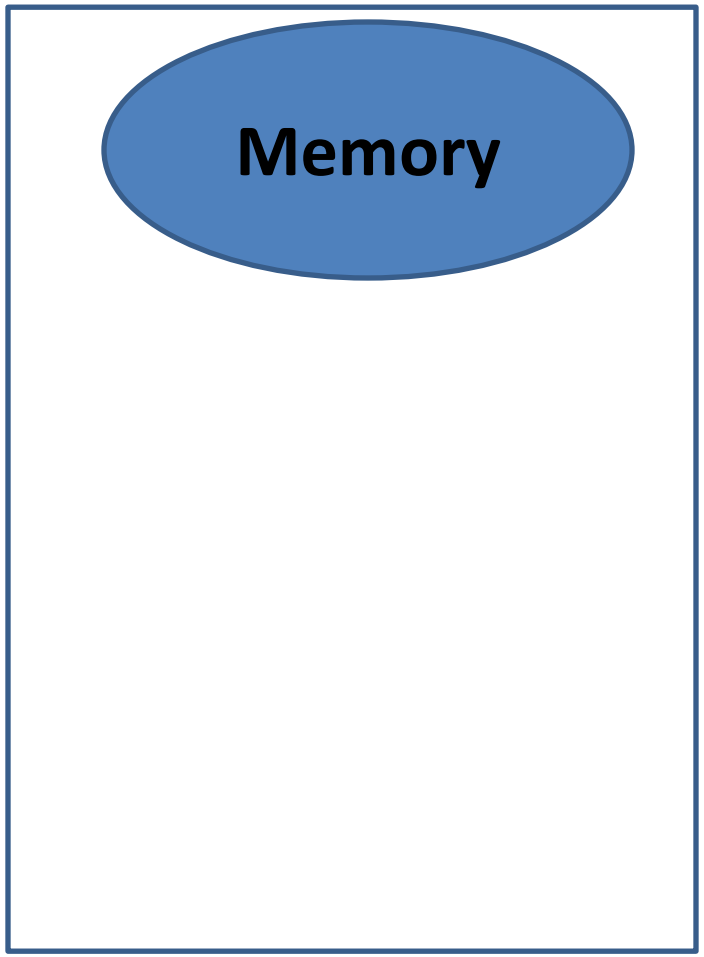


I. Hardware

Brain with no short-term memory



Scratch Pad



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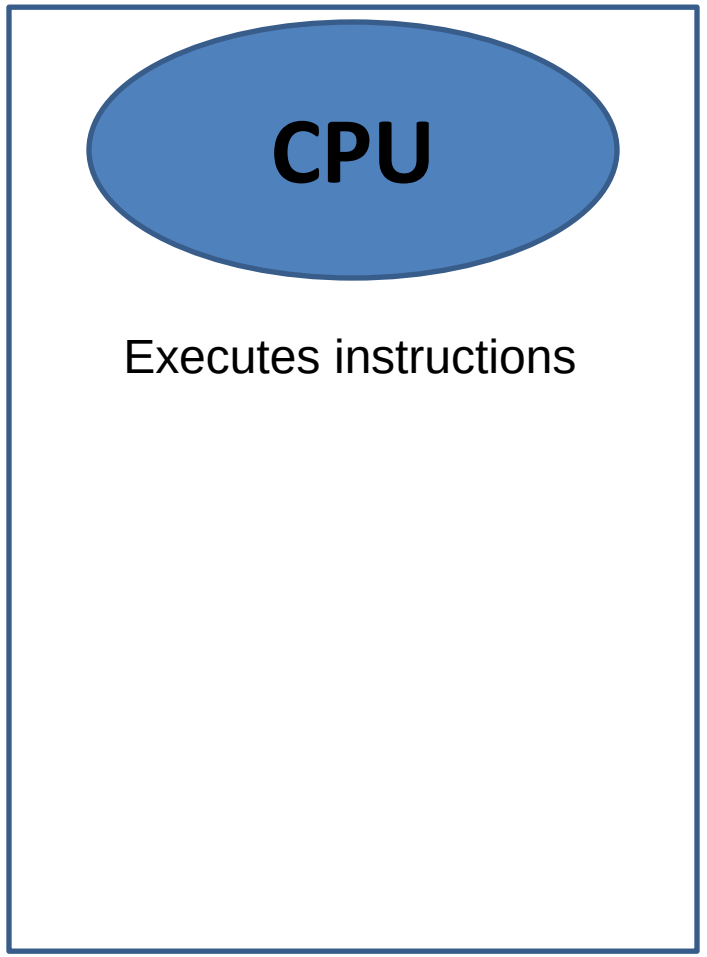

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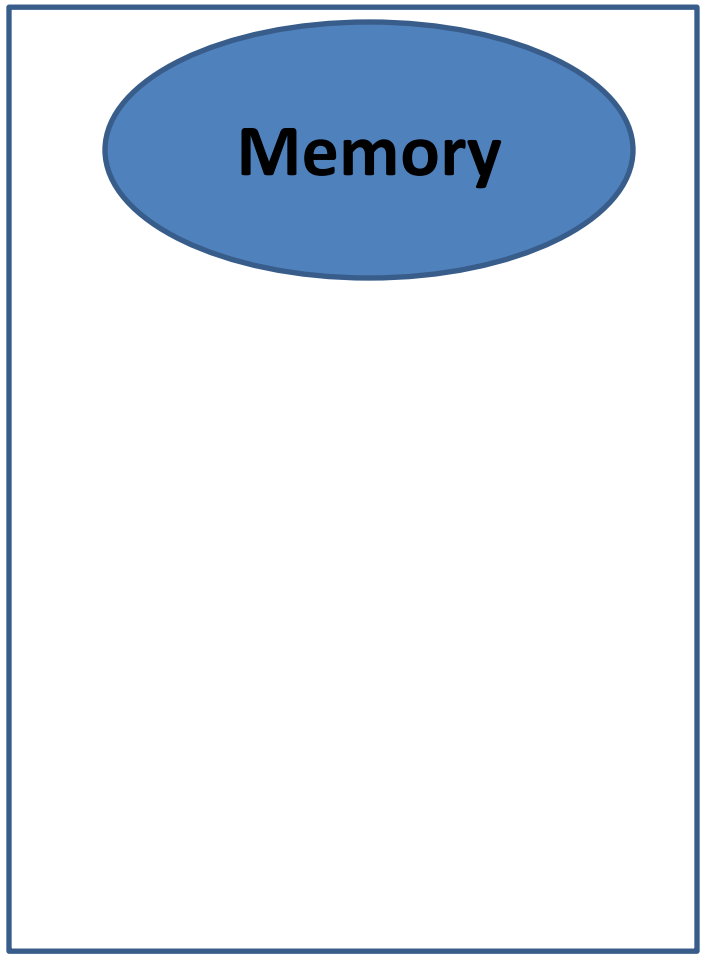

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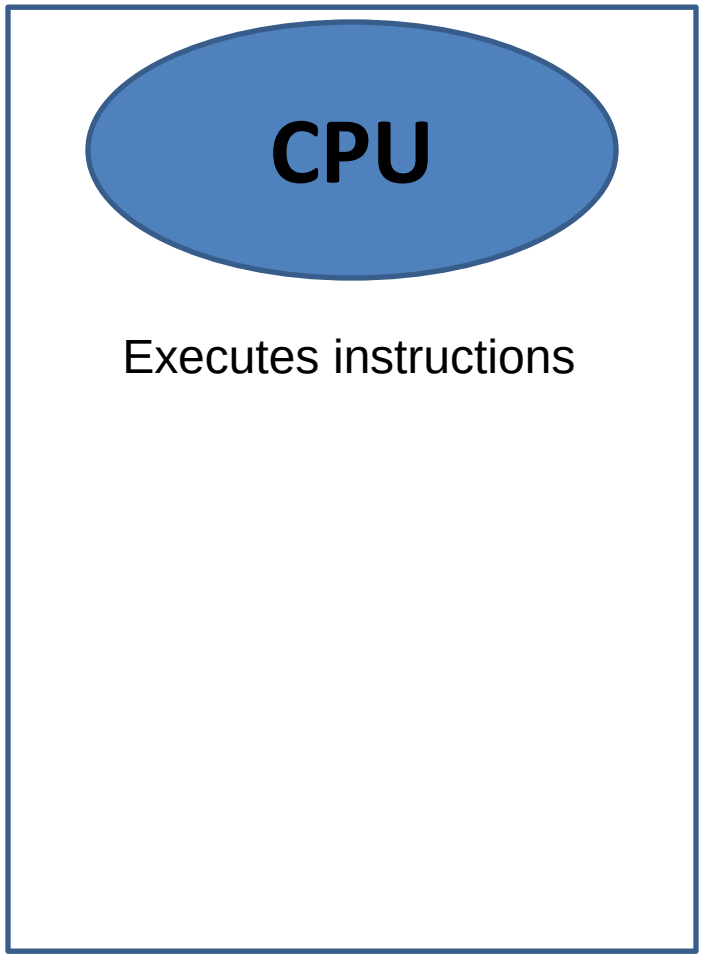
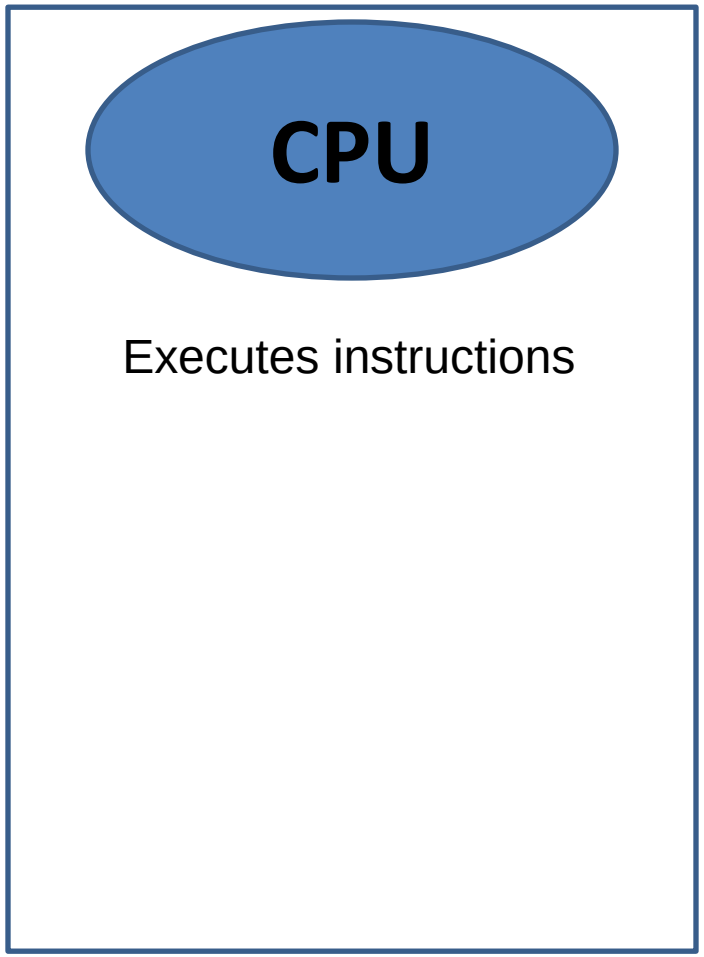


How does it “execute” instructions?

It is sent instructions from another part of the computer

I. Hardware

Brain with no short-term memory



How does it "execute" instructions?

It is sent instructions from another part of the computer

01001100000000110100011100001010

I. Hardware

Brain with no short-term memory

CPU

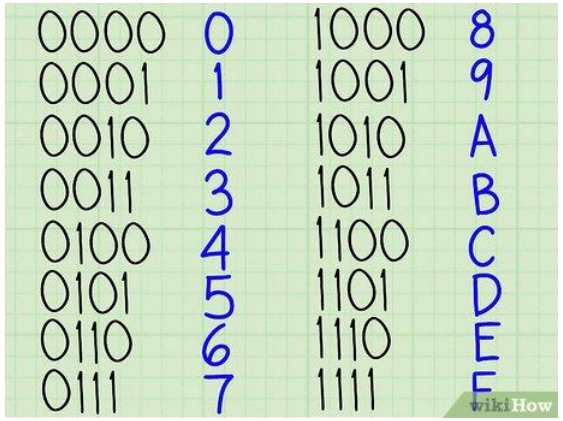
Executes instructions

How does it "execute" instructions?

It is sent instructions from *another part of the computer*

00000000101000010001100000100000 → 00 A1 18 20

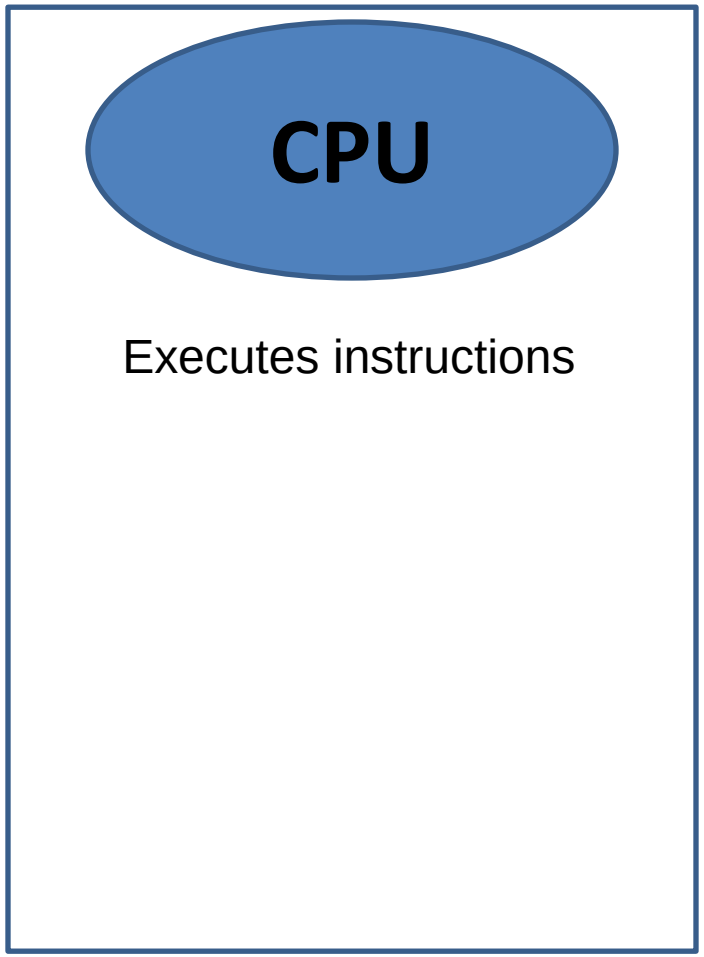
00 A1 18 20



Hex(hexadecimal) is a common representation for binary

I. Hardware

Brain with no short-term memory



How does it "execute" instructions?

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00000000101000100011000001000000

What does this instruction do?????

I. Hardware

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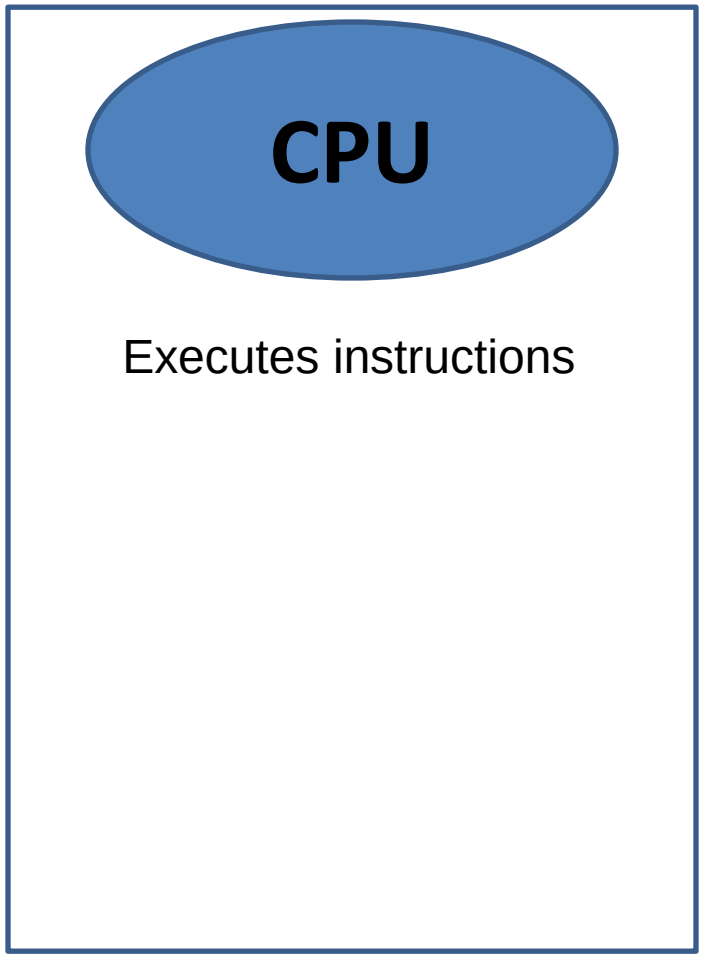

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00000000101000010001100000100000

Opcode

Common MIPS instructions

Notes: *op, funct, rd, rs, rt, imm, address, shamt* refer to fields in the instruction format
PC is assumed to point to the next instruction, **Mem** is the byte addressed main memory

Assembly Instruction	Instr Format	op op/funct	Meaning	Comments
add *$rd, $rs, $rt*	R	0/32	*$rd = $rs + $rt*	Add contents of two registers
sub *$rd, $rs, $rt*	R	0/34	*$rd = $rs - $rt*	Subtract contents of two registers
addi *$rt, $rs, imm*	I	8	*$rt = $rs + imm*	Add signed constant
addu *$rd, $rs, $rt*	R	0/33	*$rd = $rs + $rt*	Unsigned, no overflow
subu *$rd, $rs, $rt*	R	0/35	*$rd = $rs - $rt*	Unsigned, no overflow
addiu *$rt, $rs, imm*	I	9	*$rt = $rs + imm*	Unsigned, no overflow
mfc0 *$rt, $rd*	R	16	*$rt = $rd*	*rd* = coprocessor register (e.g. epc, cause, status)
mult *$rs, $rt*	R	0/24	Hi, Lo = *$rs * $rt*	64 bit signed product in Hi and Lo
multu *$rs, $rt*	R	0/25	Hi, Lo = *$rs * $rt*	64 bit unsigned product in Hi and Lo
div *$rs, $rt*	R	0/26	Lo = *$rs / $rt*, Hi = *$rs* mod *$rt*	
divu *$rs, $rt*	R	0/27	Lo = *$rs / $rt*, Hi = *$rs* mod *$rt* (unsigned)	
mfhi *$rd*	R	0/16	*$rd* = Hi	Get value of Hi
mflo *$rd*	R	0/18	*$rd* = Lo	Get value of Lo
and *$rd, $rs, $rt*	R	0/36	*$rd = $rs & $rt*	Logical AND
or *$rd, $rs, $rt*	R	0/37	*$rd = $rs | $rt*	Logical OR
andi *$rt, $rs, imm*	I	12	*$rt = $rs & imm*	Logical AND, unsigned constant
ori *$rt, $rs, imm*	I	13	*$rt = $rs | imm*	Logical OR, unsigned constant
sll *$rd, $rs, shamt*	R	0/0	*$rd = $rs << shamt*	Shift left logical (shift in zeros)
srl *$rd, $rs, shamt*	R	0/2	*$rd = $rs >> shamt*	Shift right logical (shift in zeros)

Common MIPS instructions

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sub *$rd, $rs, $rt*	R	0/34	*$rd = $rs - $rt*	Subtract contents of two registers
addi *$rt, $rs, imm*	I	8	*$rt = $rs + imm*	Add signed constant

MIPS Instruction formats

Format	Bits 31-26	Bits 25-21	Bits 20-16	Bits 15-11	Bits 10-6	Bits 5-0
R	op	rs	rt	rd	shamt	funct
I	op	rs	rt	imm		
J	op	address				

multu *$rs, $rt*	R	0/25	Hi, Lo = *$rs * $rt*	64 bit unsigned product in Hi and Lo
div *$rs, $rt*	R	0/26	Lo = *$rs / $rt*, Hi = *$rs* mod *$rt*	
divu *$rs, $rt*	R	0/27	Lo = *$rs / $rt*, Hi = *$rs* mod *$rt* (unsigned)	
mfhi *$rd*	R	0/16	*$rd* = Hi	Get value of Hi
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and *$rd, $rs, $rt*	R	0/36	*$rd = $rs & $rt*	Logical AND
or *$rd, $rs, $rt*	R	0/37	*$rd = $rs | $rt*	Logical OR
andi *$rt, $rs, imm*	I	12	*$rt = $rs & imm*	Logical AND, unsigned constant
ori *$rt, $rs, imm*	I	13	*$rt = $rs | imm*	Logical OR, unsigned constant
sll *$rd, $rs, shamt*	R	0/0	*$rd = $rs << shamt*	Shift left logical (shift in zeros)
srl *$rd, $rs, shamt*	R	0/2	*$rd = $rs >> shamt*	Shift right logical (shift in zeros)

I. Hardware

Brain with no short-term memory

CPU

Executes instructions

Must decipher what instruction to execute

How does it “execute” instructions?

It is sent instructions from another part of the computer

00000000101000010001100000100000

Opcode

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Opcode $rs

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Opcode $rs $rt $rd

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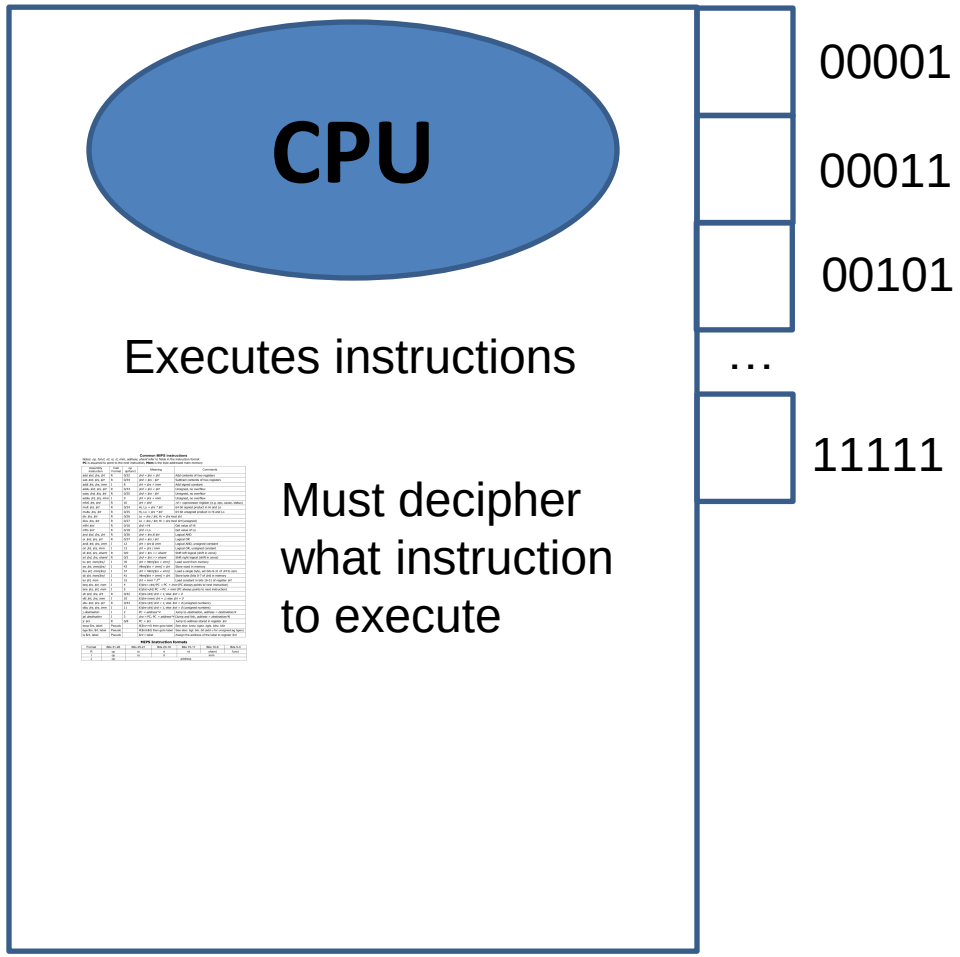

How does it "execute" instructions?

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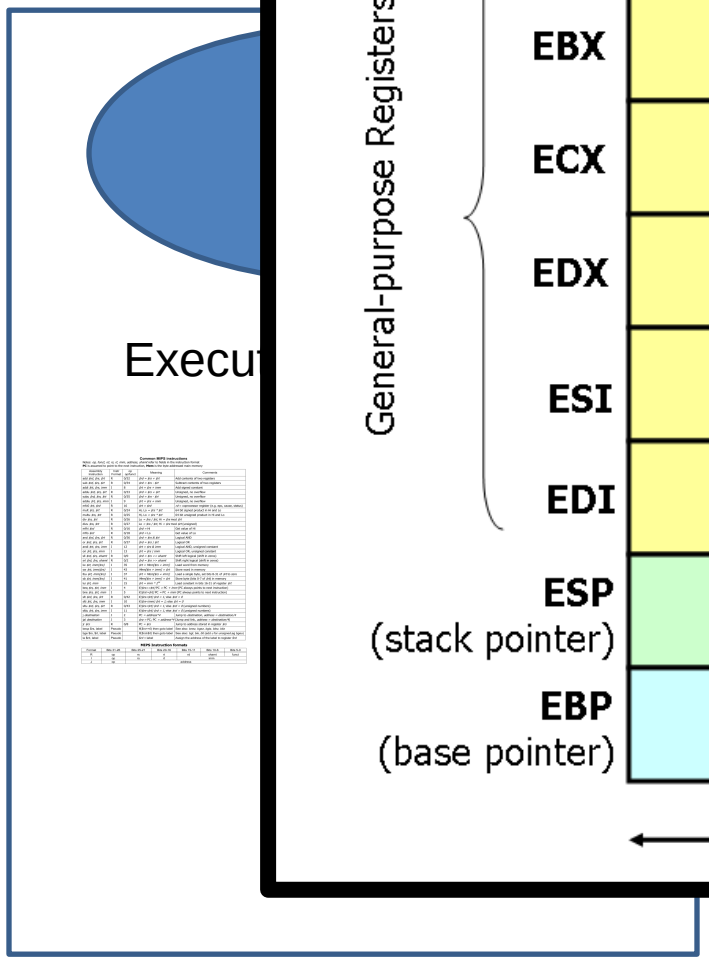
00000000101000010001100000100000

Opcode $rs $rt $rd shamt funct

$ denotes that it is a **register**

I. Hardware

Brain with nc



Execut

ructions?

rt of the computer

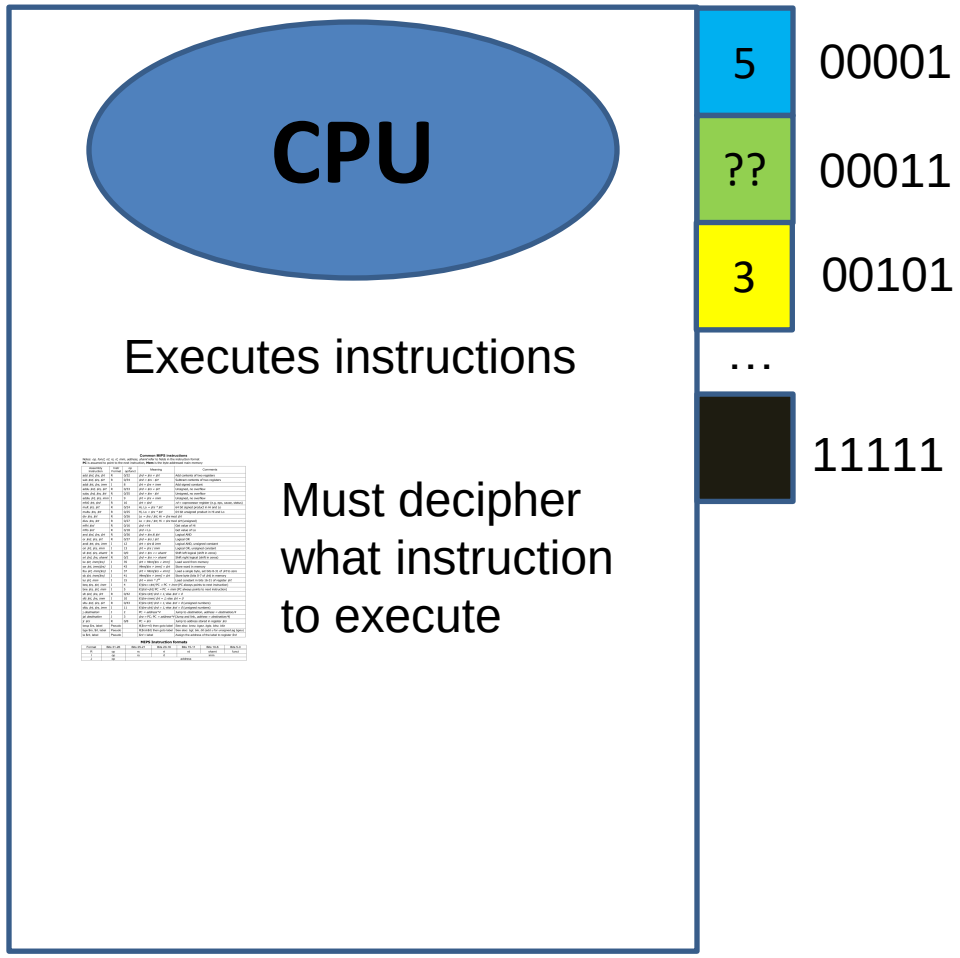
00000100000

shamt funct

er

I. Hardware

Brain with no short-term memory



How does it "execute" instructions?

It is sent instructions from another part of the computer

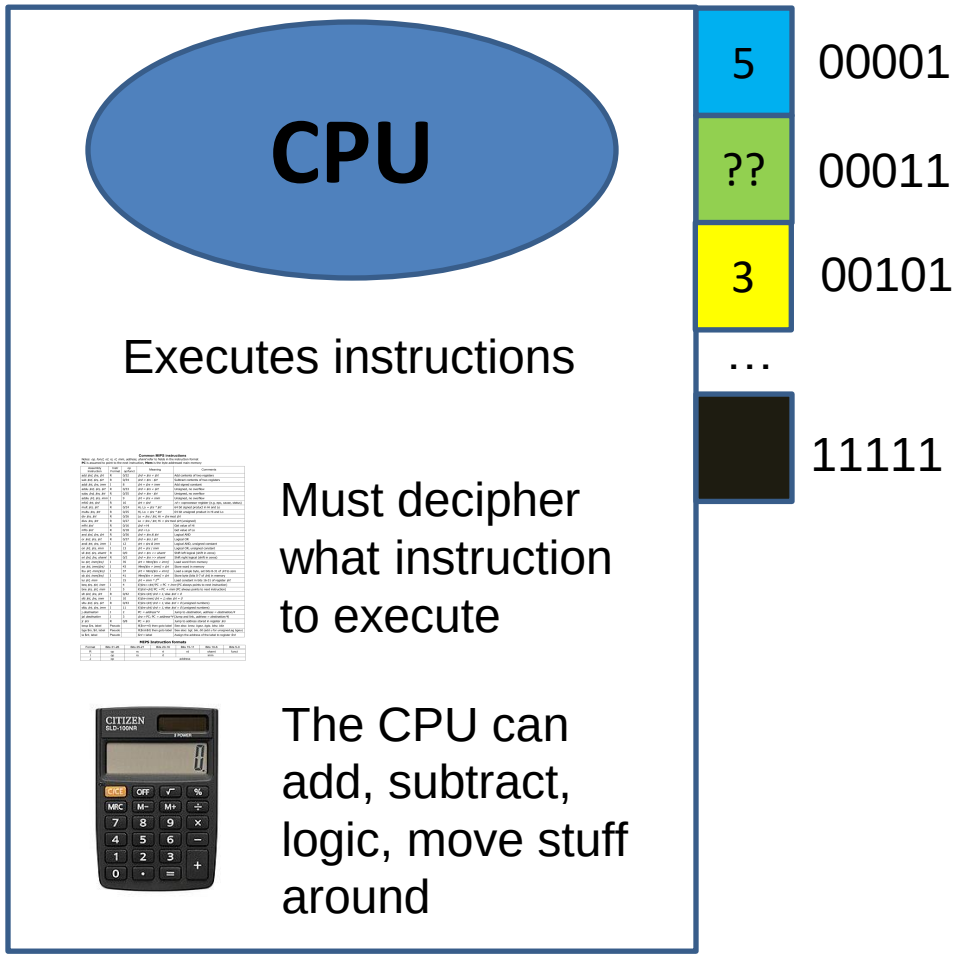
00000000101000010001100000100000

Opcode $rs $rt $rd shamt funct

```
ADD $rs, $rt, $rd
```

I. Hardware

Brain with no short-term memory



Registers

5	00001
??	00011
3	00101
…	
	11111

How does it "execute" instructions?

It is sent instructions from another part of the computer

00000000101000010001100000100000

Opcode $rs $rt $rd shamt funct

```
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```

I. Hardware

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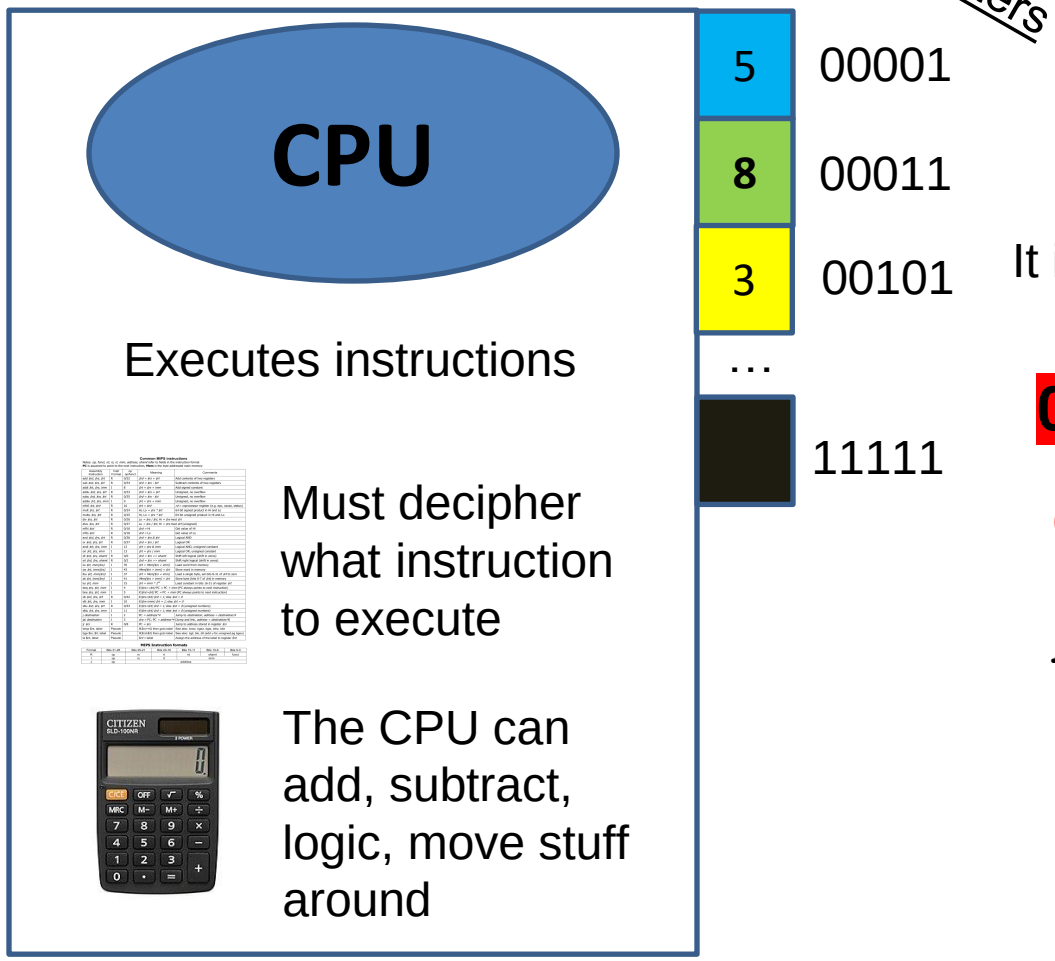

Registers

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8	00011
3	00101
…	
	11111

How does it "execute" instructions?

It is sent instructions from another part of the computer

00000000101000010001100000100000

Opcode $rs $rt $rd shamt funct

```
ADD $rs, $rt, $rd
```

I. Hardware

Brain with no short-term memory

CPU

Executes instructions

Must decipher what instruction to execute

The CPU can add, subtract, logic, move stuff around

Regis

5

8

3

…

0000

00

1

CPU uses

Electricity™

I. Hardware

Brain with no short-term memory

CPU

Executes instructions

Must decipher what instruction to execute

The CPU can add, subtract, logic, move stuff around

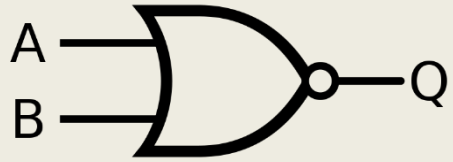
Regis

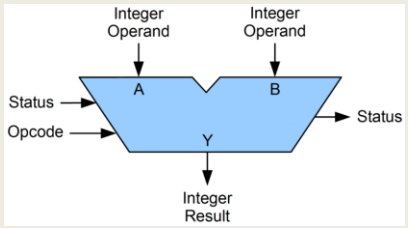
5	0000
8	00
3	
…	
	1

CPU uses

Electricity™

To decipher and execute instructions

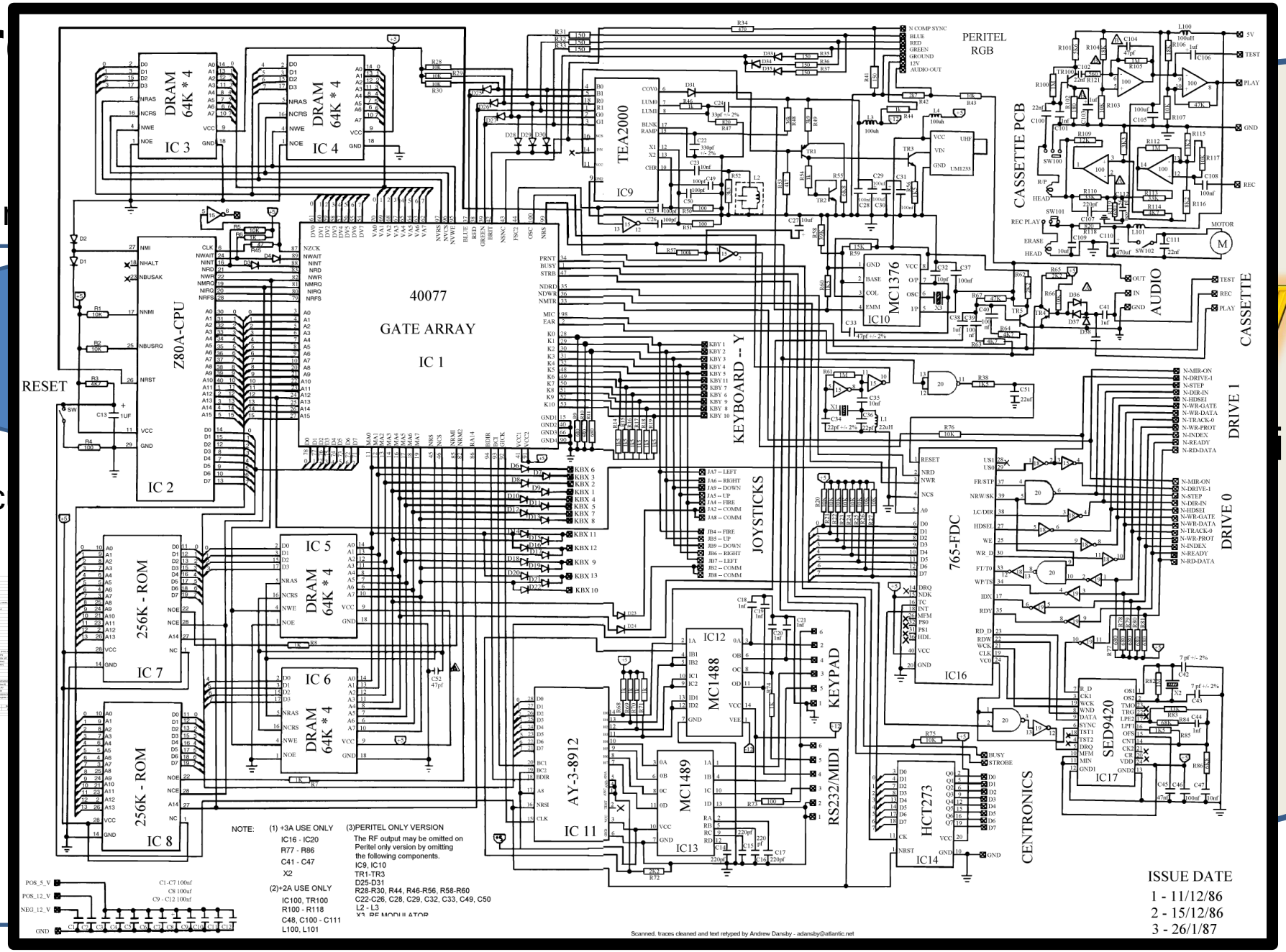




I. Har

Brain with

Exec



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The CPU can add, subtract, logic, move stuff around

Regis

5	0000
8	00
3	
…	
	1

CPU uses

Electricity™

To decipher and execute instructions

*Everything is just **electricity on** or **electricity off***

I. Hardware

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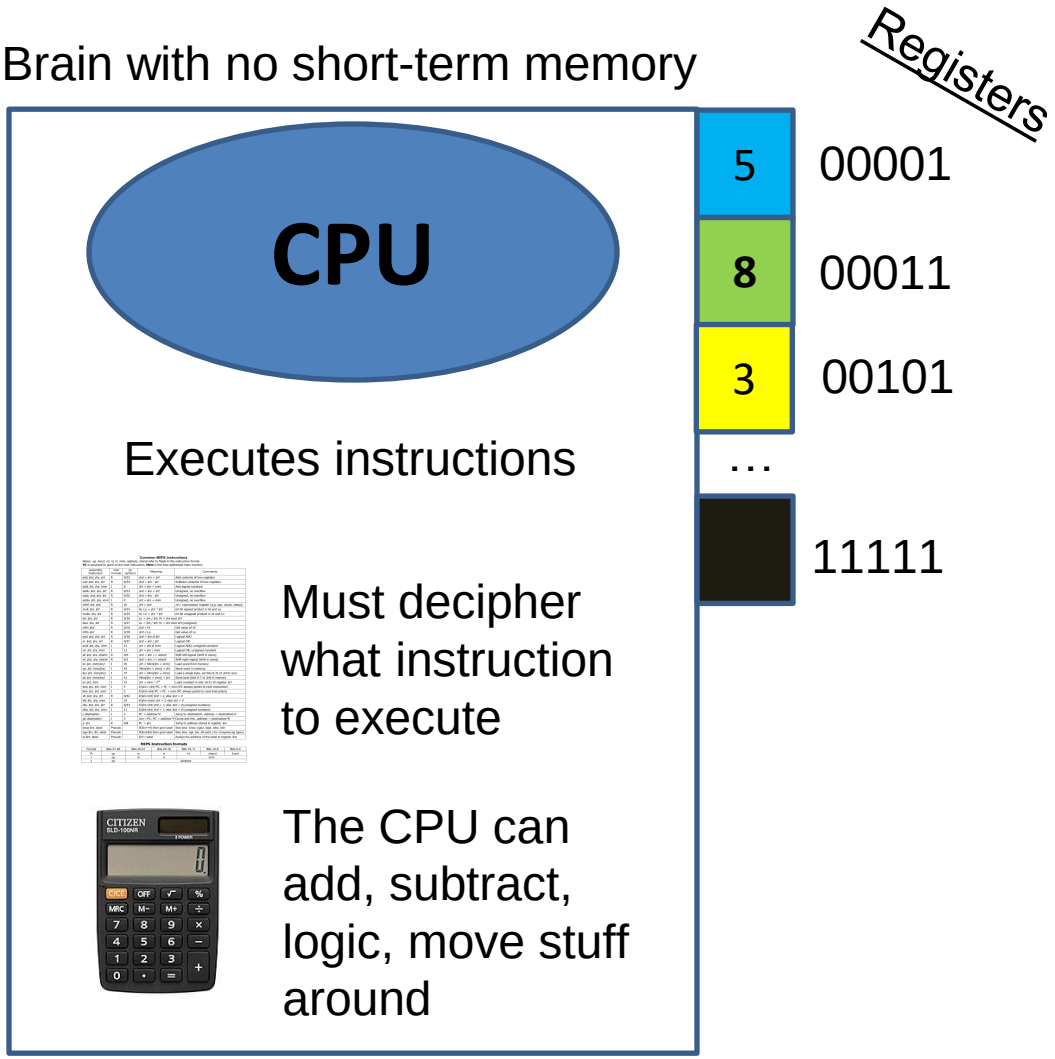
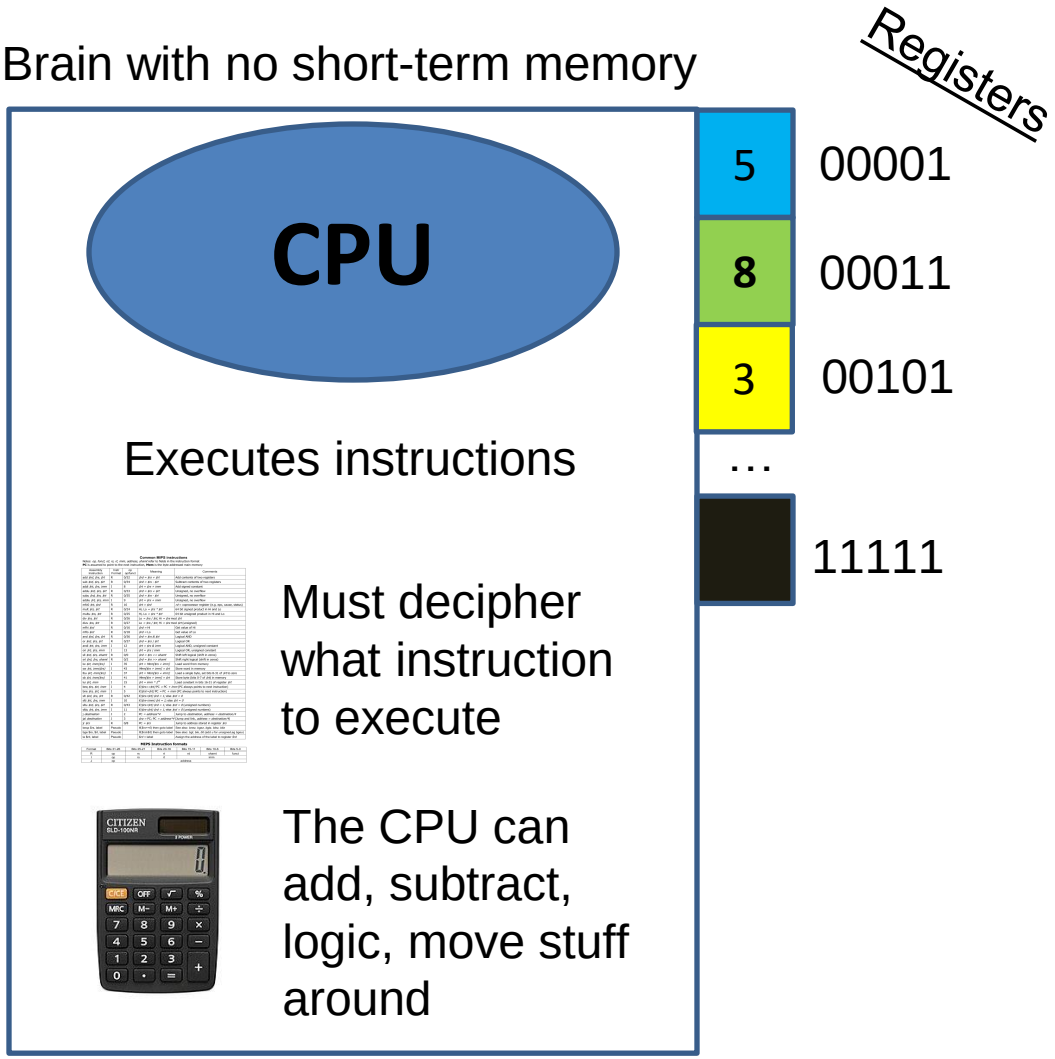


32 bit vs 64 bit?

I. Hardware

Brain with no short-term memory

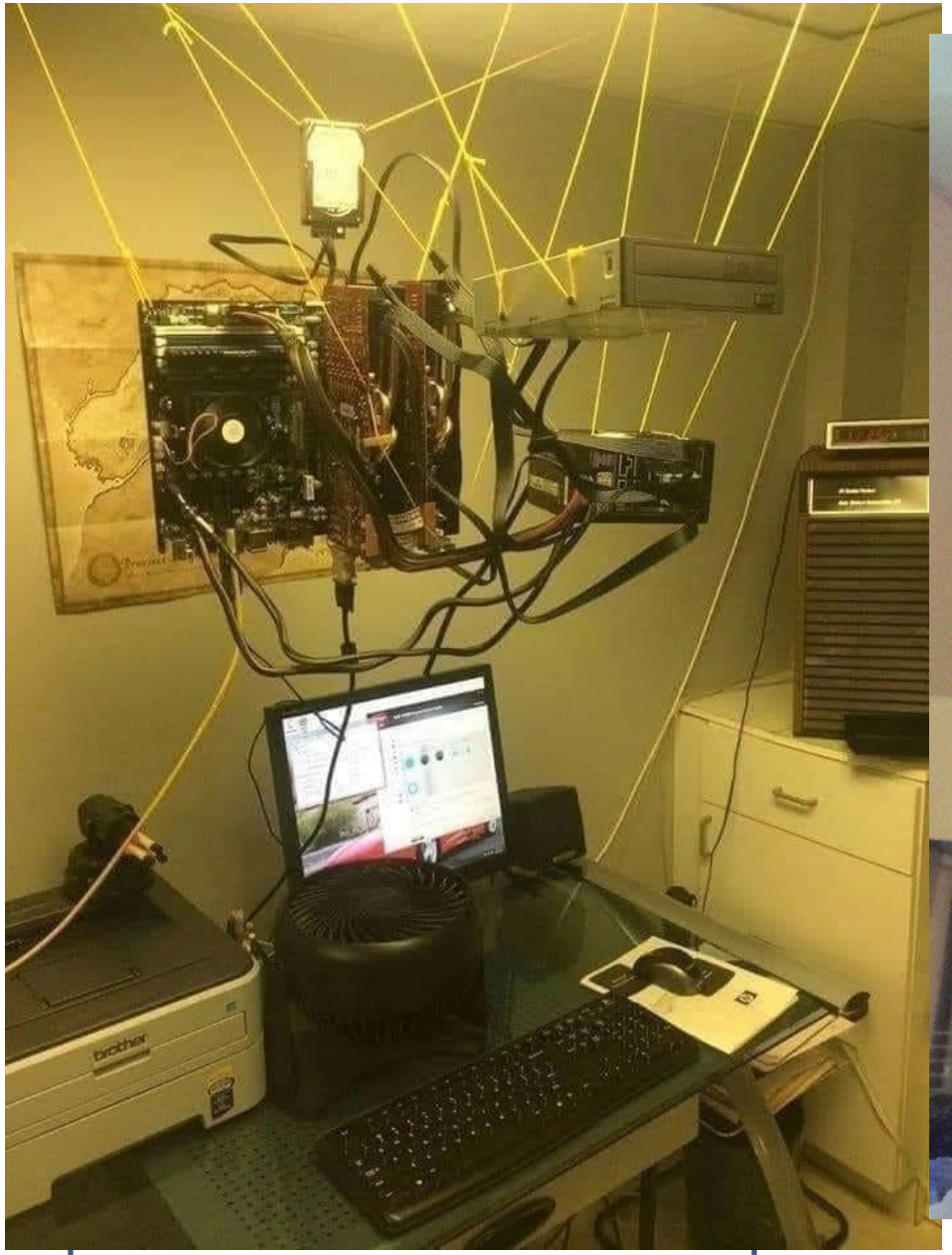


32 bit vs 64 bit?

Parameter	32-bit processors	64-bit processors
Addressable space	It has 4 GB addressable space	64-bit processors have 16 exabytes addressable space
Application support	64-bit applications and programs won't work	32-bit applications and programs will work
OS support	Need a 32-bit operating system.	It can run on 32 and the 64-bit operating system.
Support for multi-tasking	Not an ideal option for stress testing and multi-tasking.	Works best for performing multi-tasking and stress testing.
OS and CPU requirement	32-bit operating systems and applications require 32-bit CPUs	64-bit OS demands 64-bit CPU, and 64-bit applications require 64-bit OS and CPU.
System available	Support Windows 7, 8 Vista, XP, and, Linux.	Windows XP Professional, Windows Vista, Windows 7, Windows 8,Windows 10, Linux, and Mac OS X.

How CPU is made:

https://www.youtube.com/watch?v=vuvckBQ1bME


brother

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Registers

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8	00011
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…	
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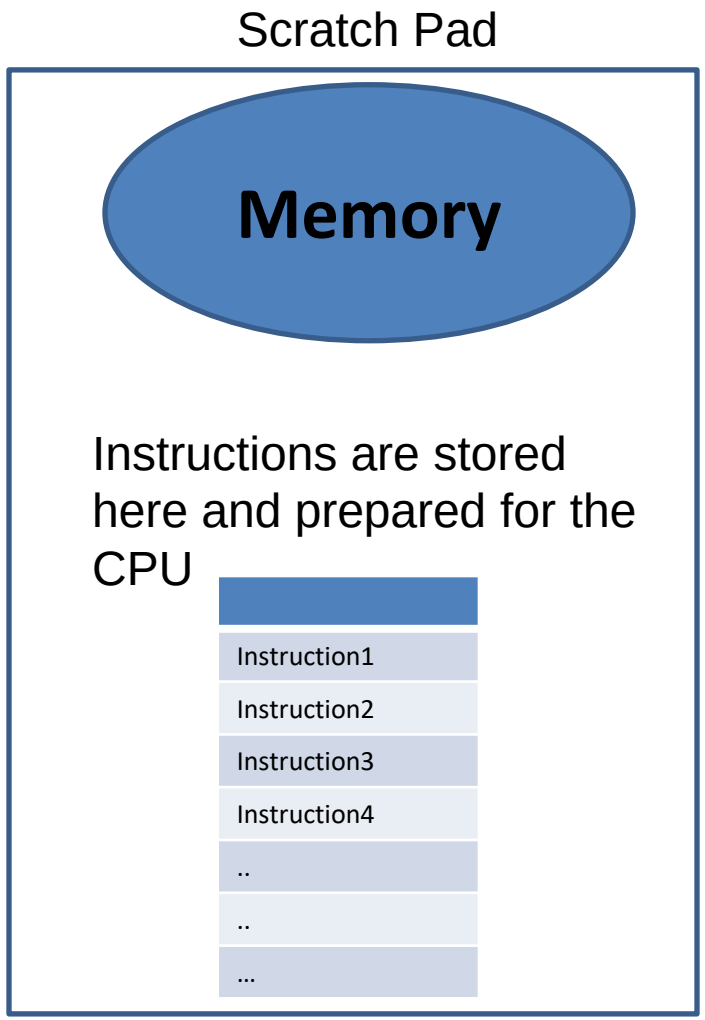
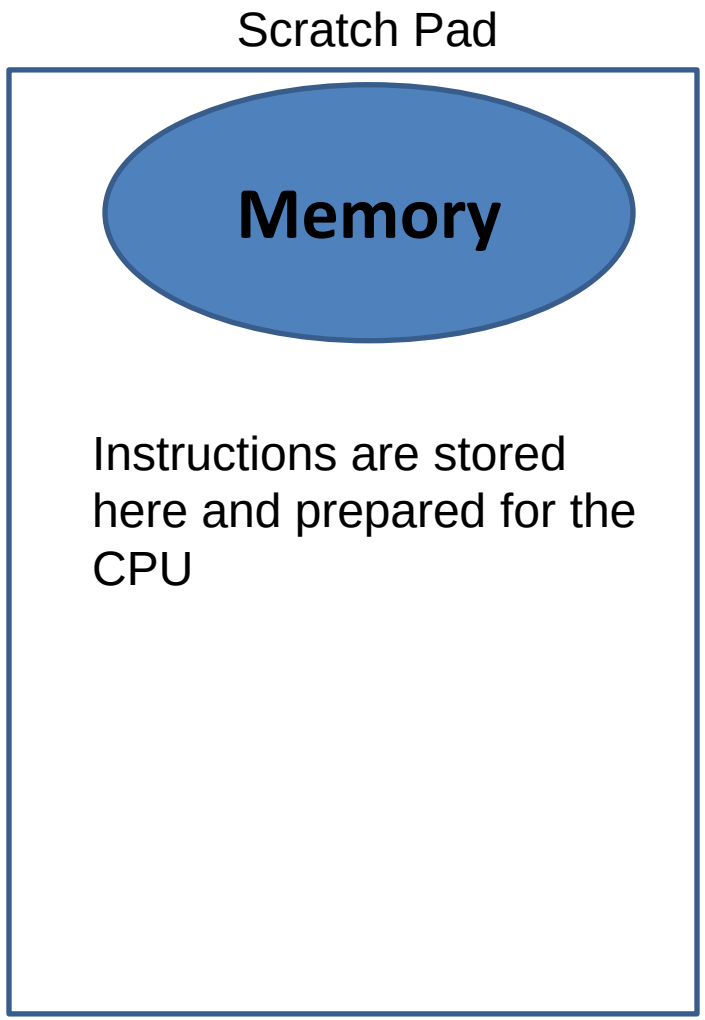
Scratch Pad

Memory

I. Hardware

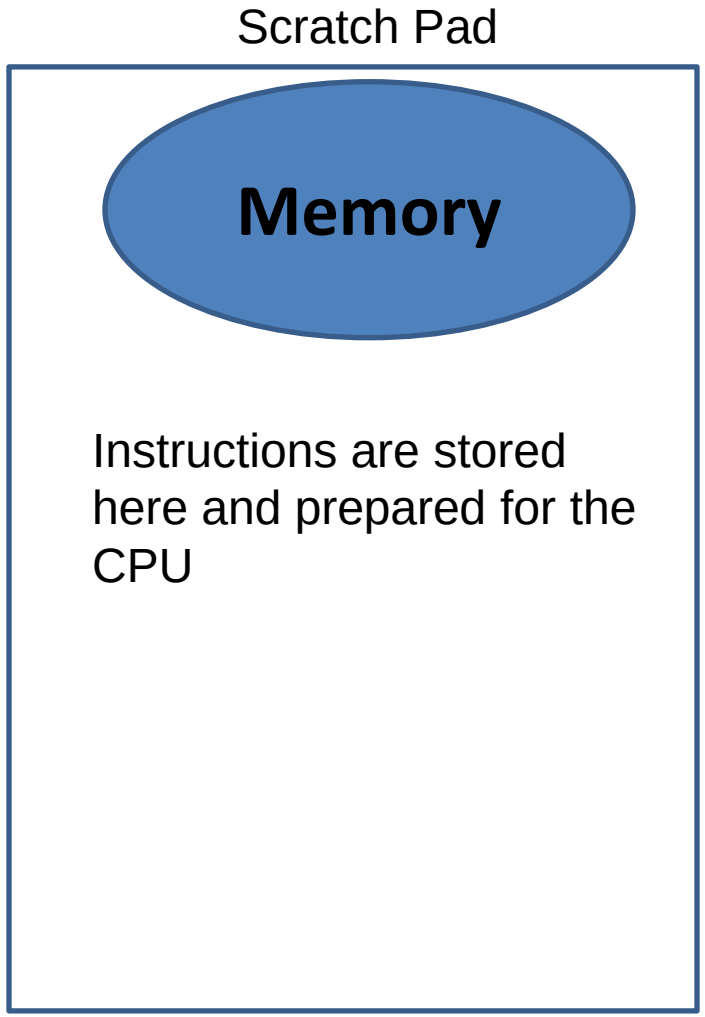

Scratch Pad
Memory
Instructions are stored here and prepared for the CPU
Instruction1
Instruction2
Instruction3
Instruction4
..
..
...

I. Hardware



When computer programs are executed, their instructions will eventually get stored in memory

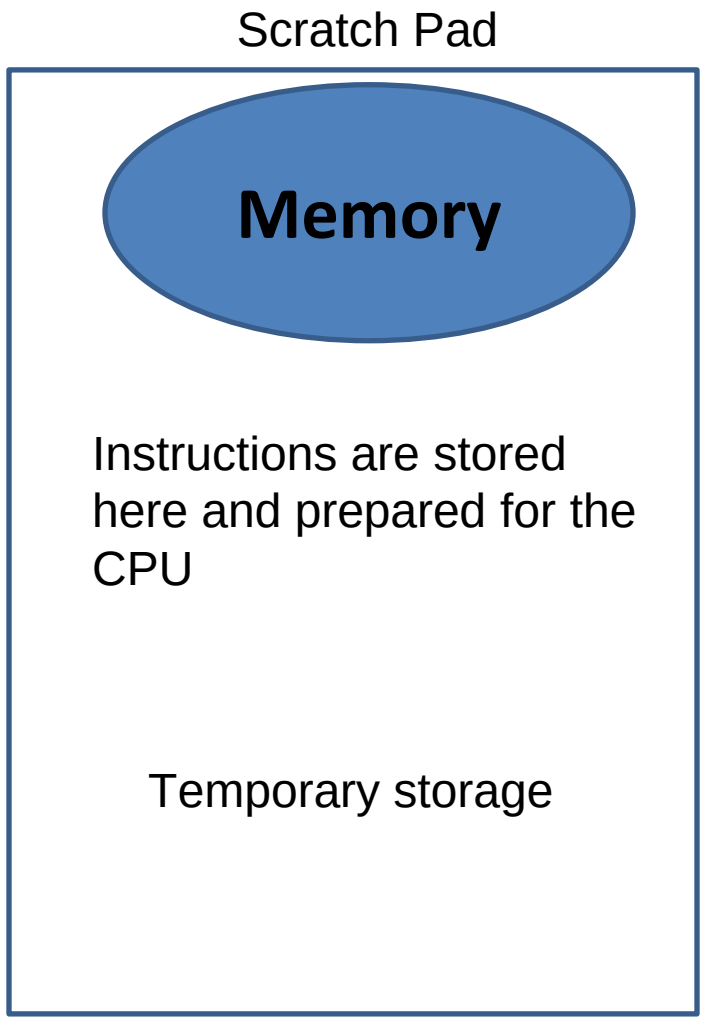
I. Hardware



When computer programs are executed, their instructions will eventually get stored in memory

Permanently?!?!

I. Hardware



When computer programs are executed, their instructions will eventually get stored in memory

Main memory is **volatile**

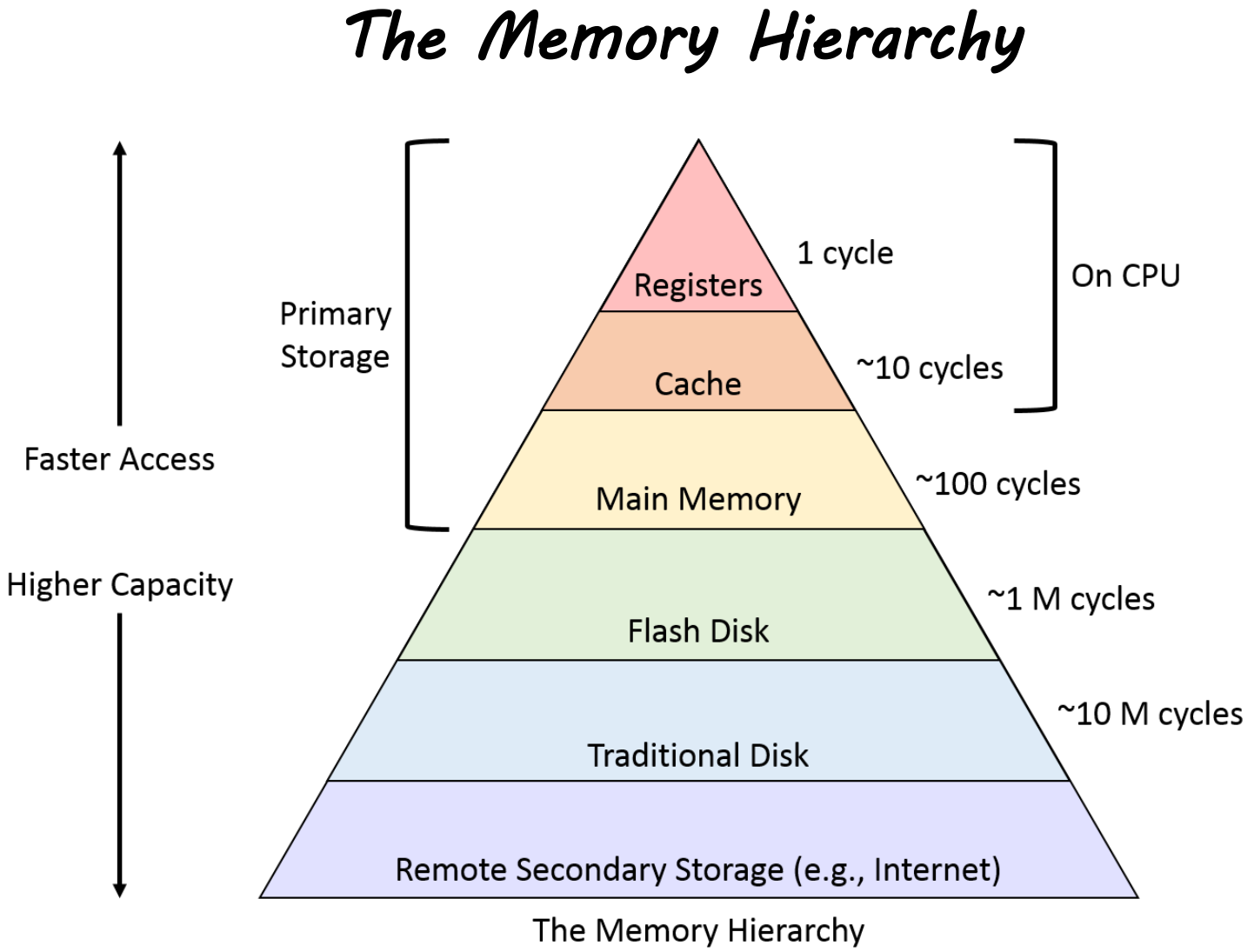
I. Hardware

Scratch Pad

Memory

Instructions are stored here and prepared for the CPU

Temporary storage



The Memory Hierarchy

I. Hardware

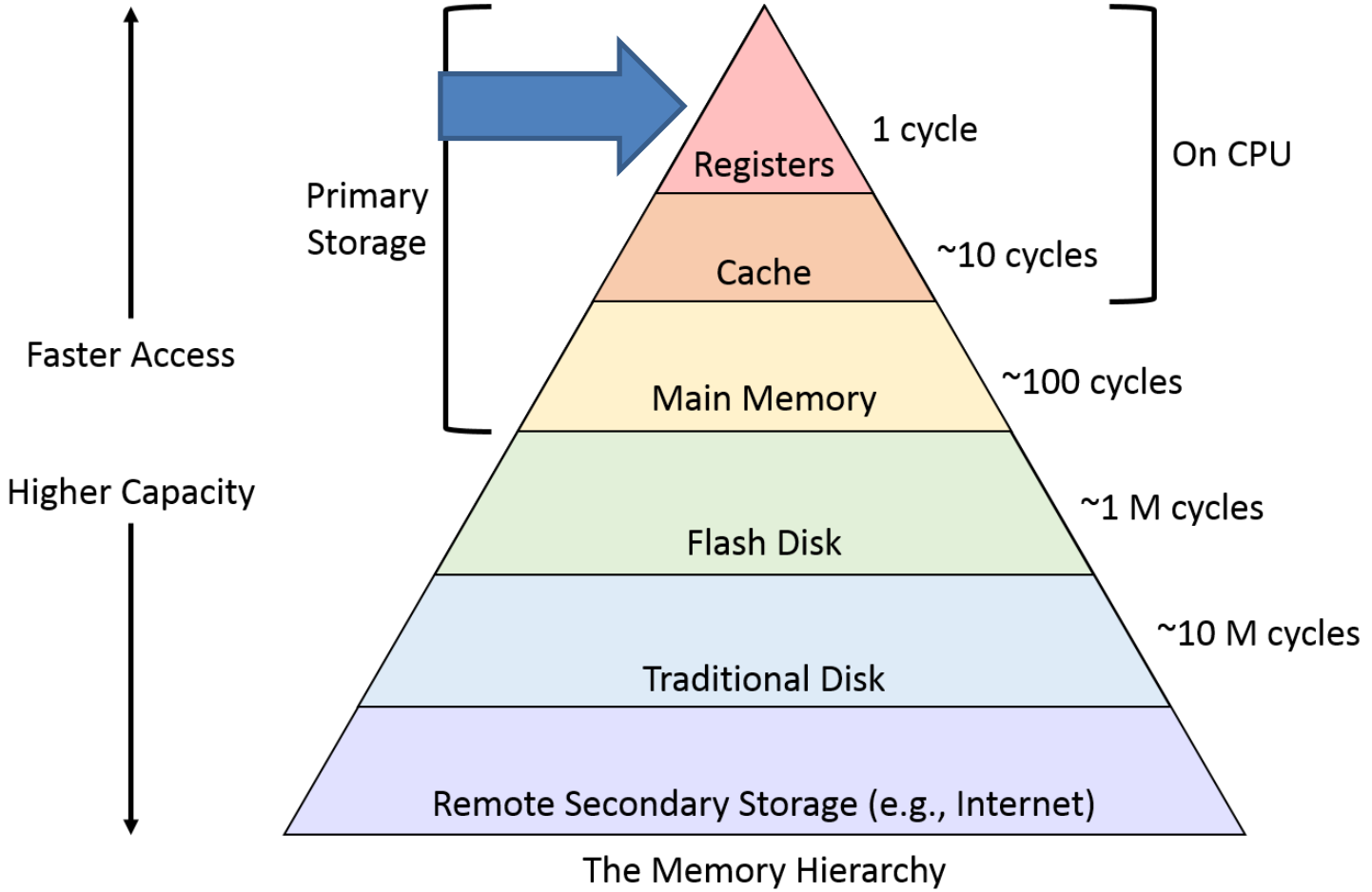
Scratch Pad

Memory

Instructions are stored here and prepared for the CPU

Temporary storage

The Memory Hierarchy

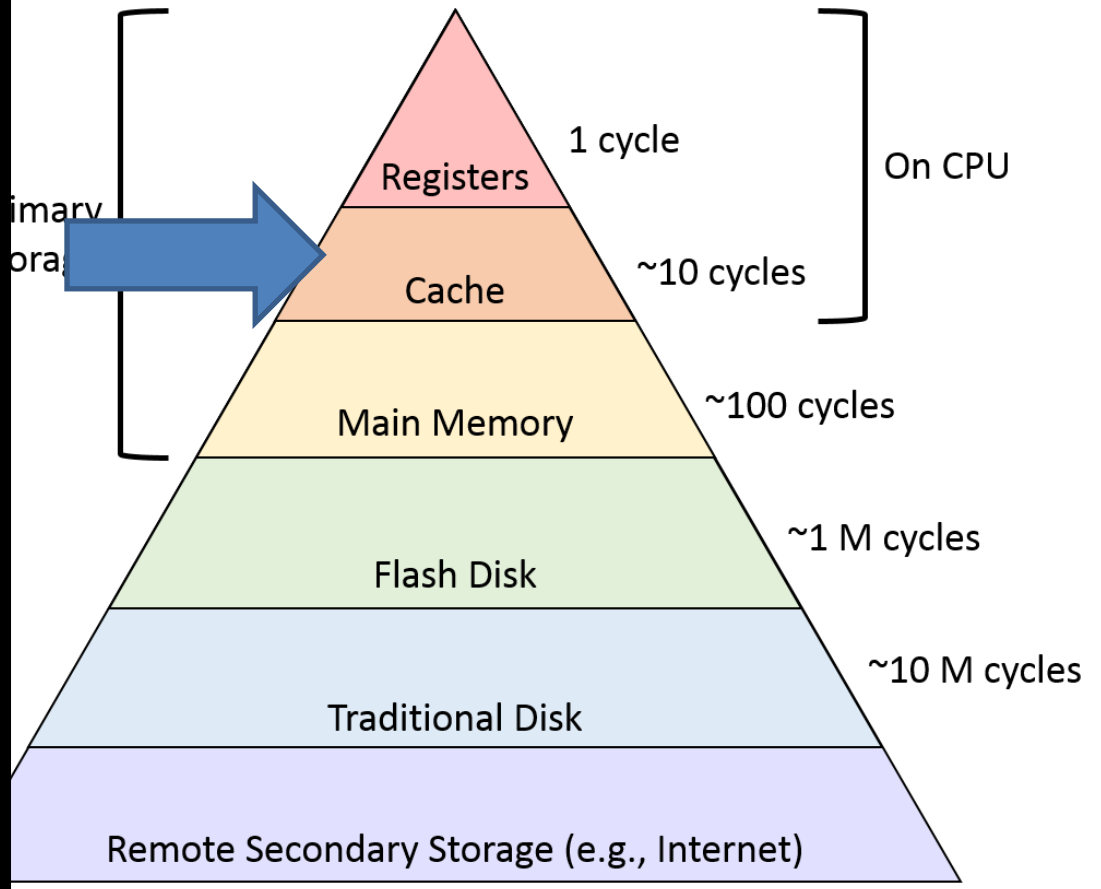


The Memory Hierarchy

I. Hardware

My CPU when the L1 cache misses

The Memory Hierarchy



The Memory Hierarchy

I. Hardware

The Memory Hierarchy

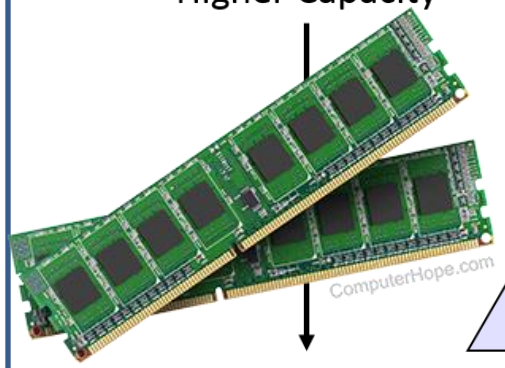
Scratch Pad

Memory

Instructions are stored

Faster Access

Higher Capacity



Primary Storage

Registers	1 cycle	On CPU
Cache	~10 cycles	
RAM Main Memory	~100 cycles	
Flash Disk	~1 M cycles	
Traditional Disk	~10 M cycles	
Remote Secondary Storage (e.g., Internet)		

The Memory Hierarchy

I. Hardware

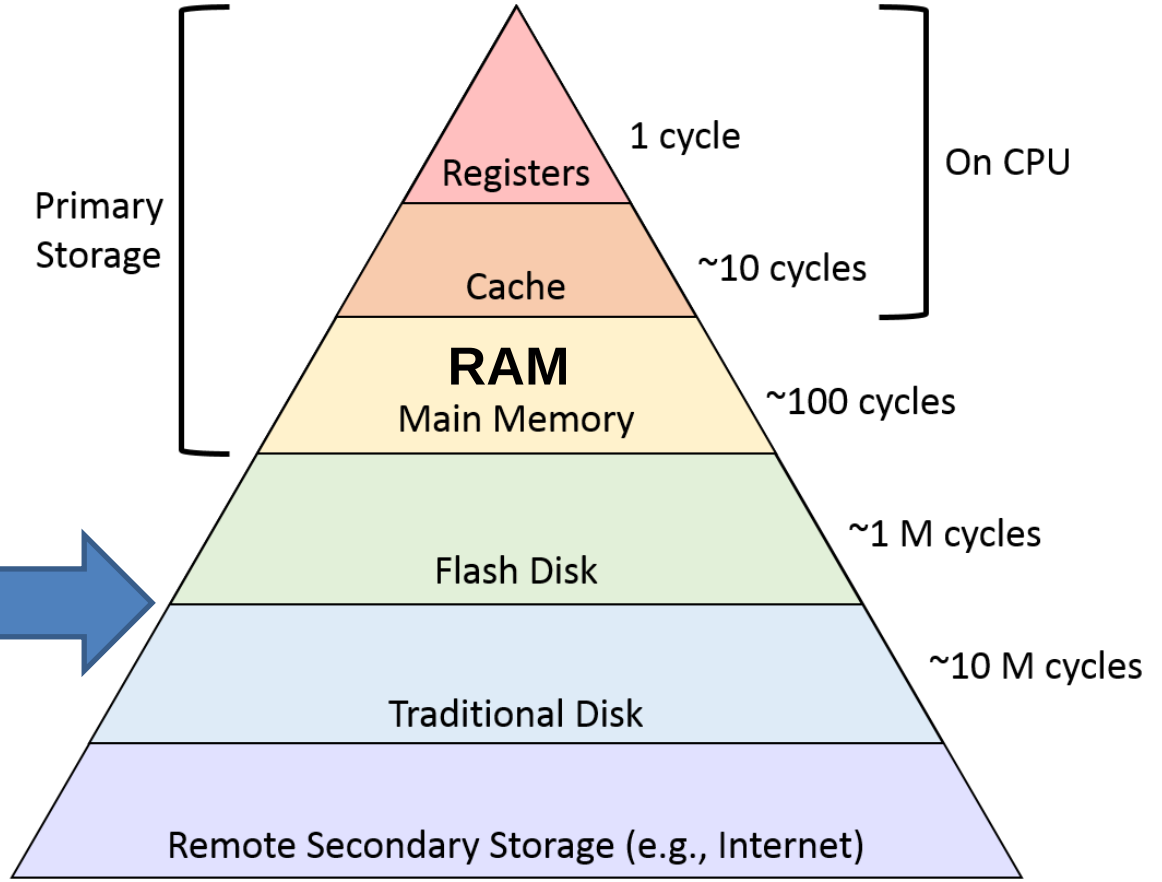
Scratch Pad

Memory

Instructions are stored here and prepared for the CPU

Tempora

The Memory Hierarchy



The Memory Hierarchy

I. Hardware

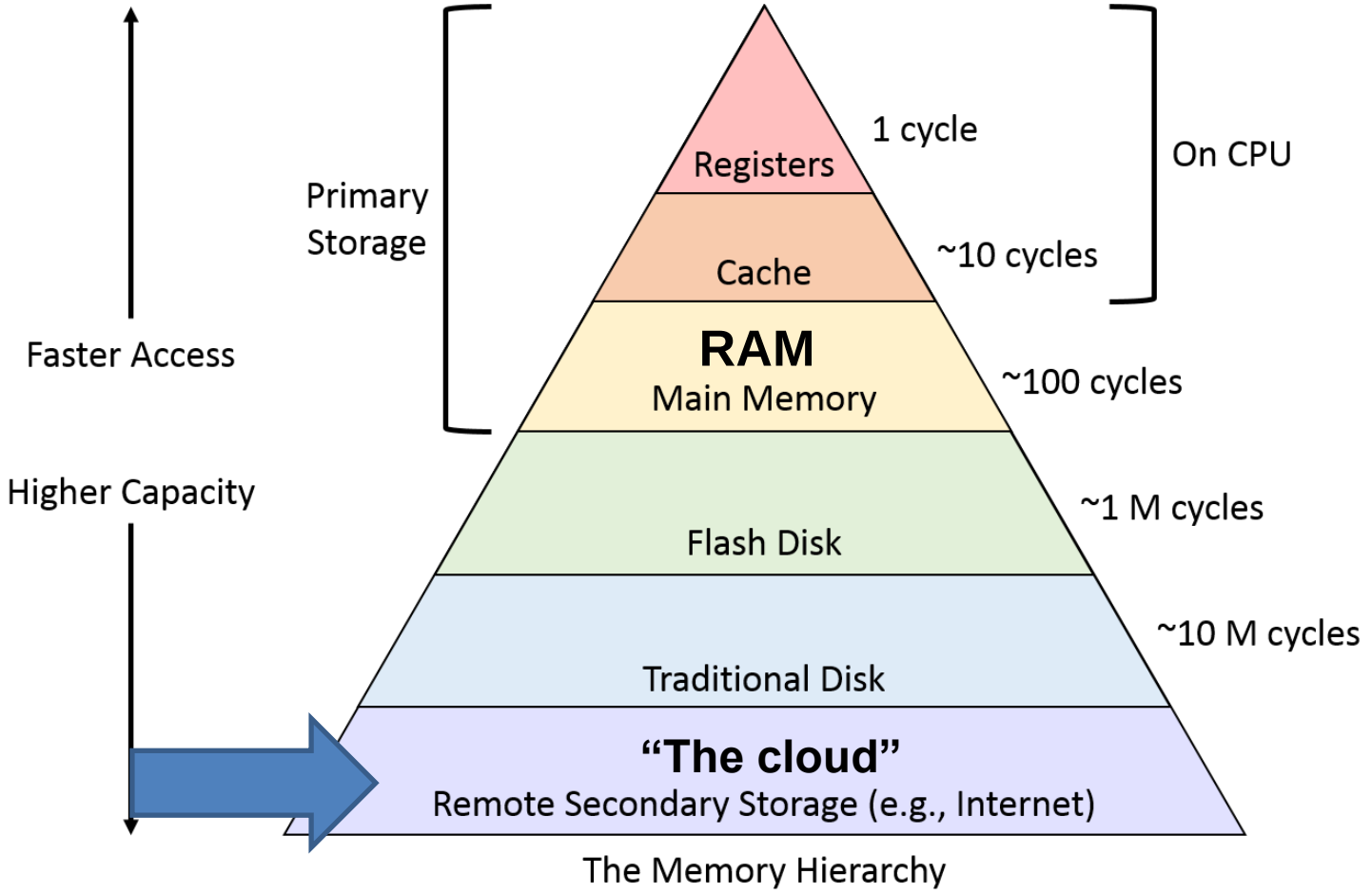
Scratch Pad

Memory

Instructions are stored here and prepared for the CPU

Temporary storage

The Memory Hierarchy



The Memory Hierarchy

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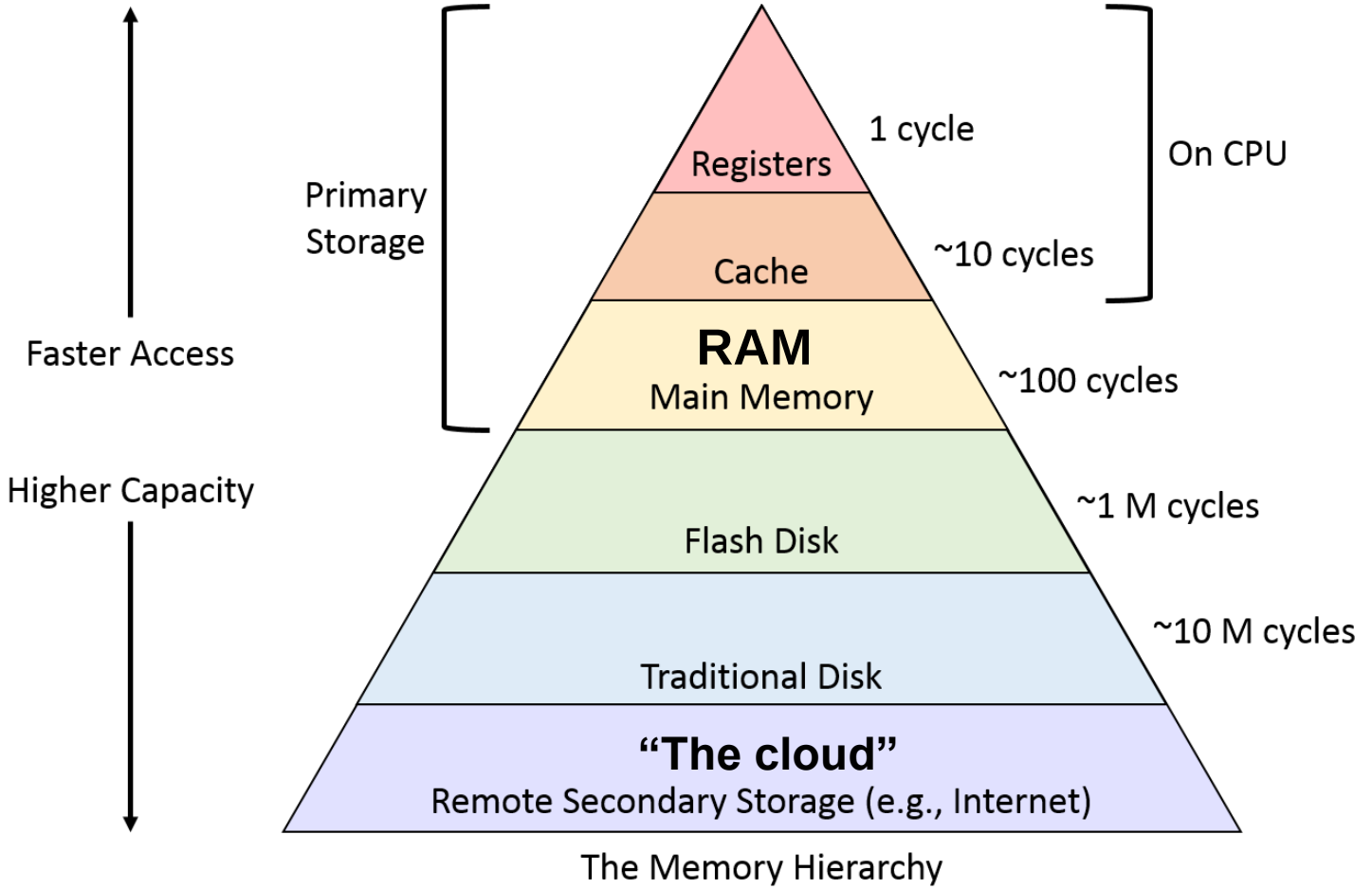
The Memory Hierarchy



The Memory Hierarchy

I. Hardware

Scratch Pad

Memory

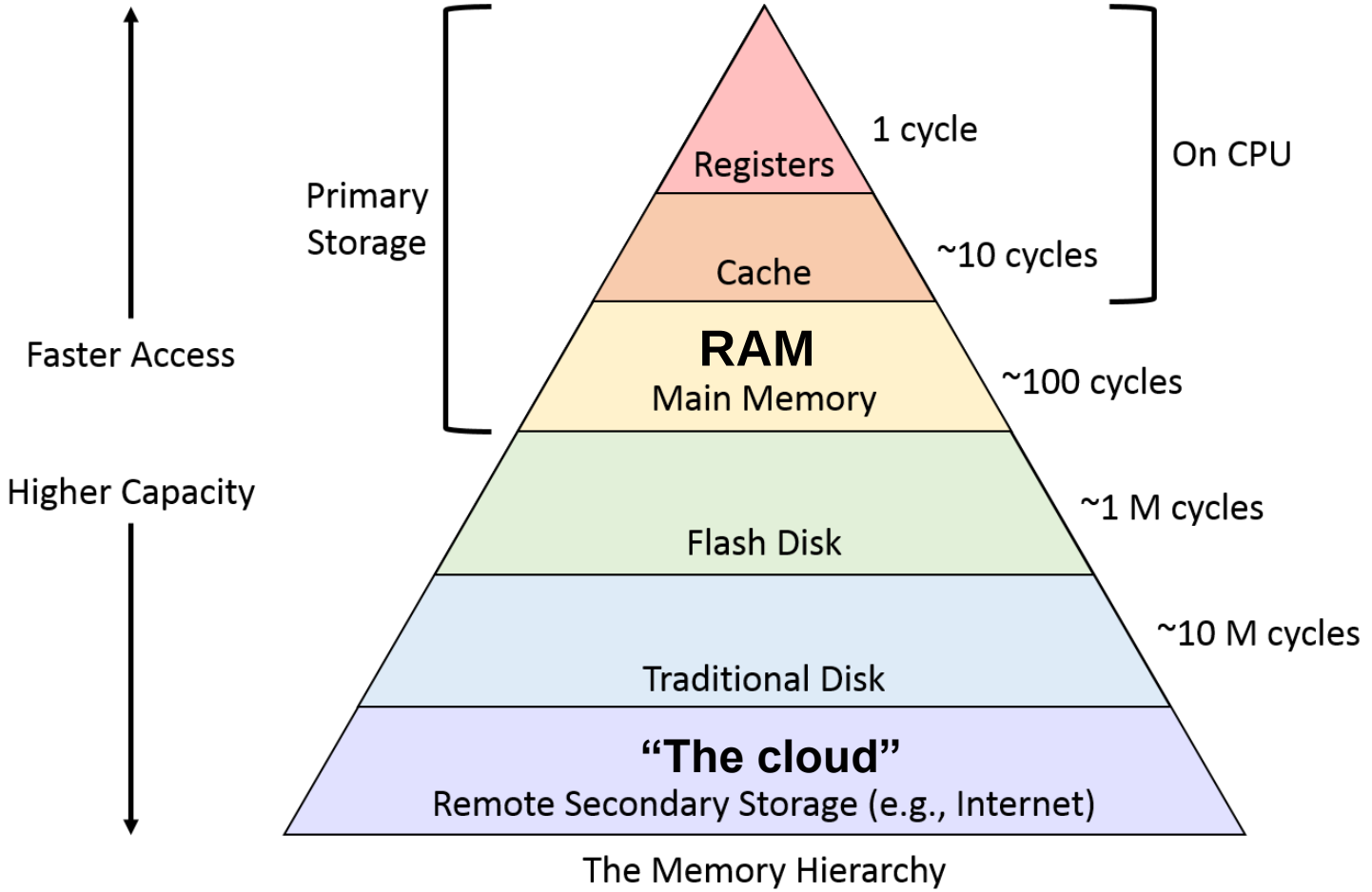
Instructions are stored here and prepared for the CPU

Temporary storage

RAM
Main Memory

Why not use memory/registers for everything??

The Memory Hierarchy



The Memory Hierarchy

I. Hardware

Brain with no short-term memory

CPU

Executes instructions

Must decipher what instruction to execute

The CPU can add, subtract, logic, move stuff around

Registers

5	00001
8	00011
3	00101
…	
	11111

Scratch Pad

Memory

Instructions are stored here and prepared for the CPU

Temporary storage

RAM
Main Memory

I. Hardware

Brain with no short-term memory

CPU

Executes instructions

Must decipher what instruction to execute

The CPU can add, subtract, logic, move stuff around

Registers

5	00001
8	00011
3	00101
…	
	11111

1. Fetch next instruction

Scratch Pad

Memory

Instructions are stored here and prepared for the CPU

Temporary storage

RAM
Main Memory

I. Hardware

Brain with no short-term memory

CPU

Executes instructions

Must decipher what instruction to execute

The CPU can add, subtract, logic, move stuff around

Registers

5	00001
8	00011
3	00101
…	
	11111

1. **Fetch next instruction**
2. **Decode Instruction**

Scratch Pad

Memory

Instructions are stored here and prepared for the CPU

Temporary storage

RAM
Main Memory

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Brain with no short-term memory

CPU

Executes instructions

Must decipher what instruction to execute

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1. **Fetch next instruction**
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3. **Execute Instruction**

Scratch Pad

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Instructions are stored here and prepared for the CPU

Temporary storage

RAM
Main Memory

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Executes instructions

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Registers

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3	00101
…	
	11111

1. **Fetch next instruction**
2. **Decode Instruction**
3. **Execute Instruction**
4. **Store result***

Scratch Pad

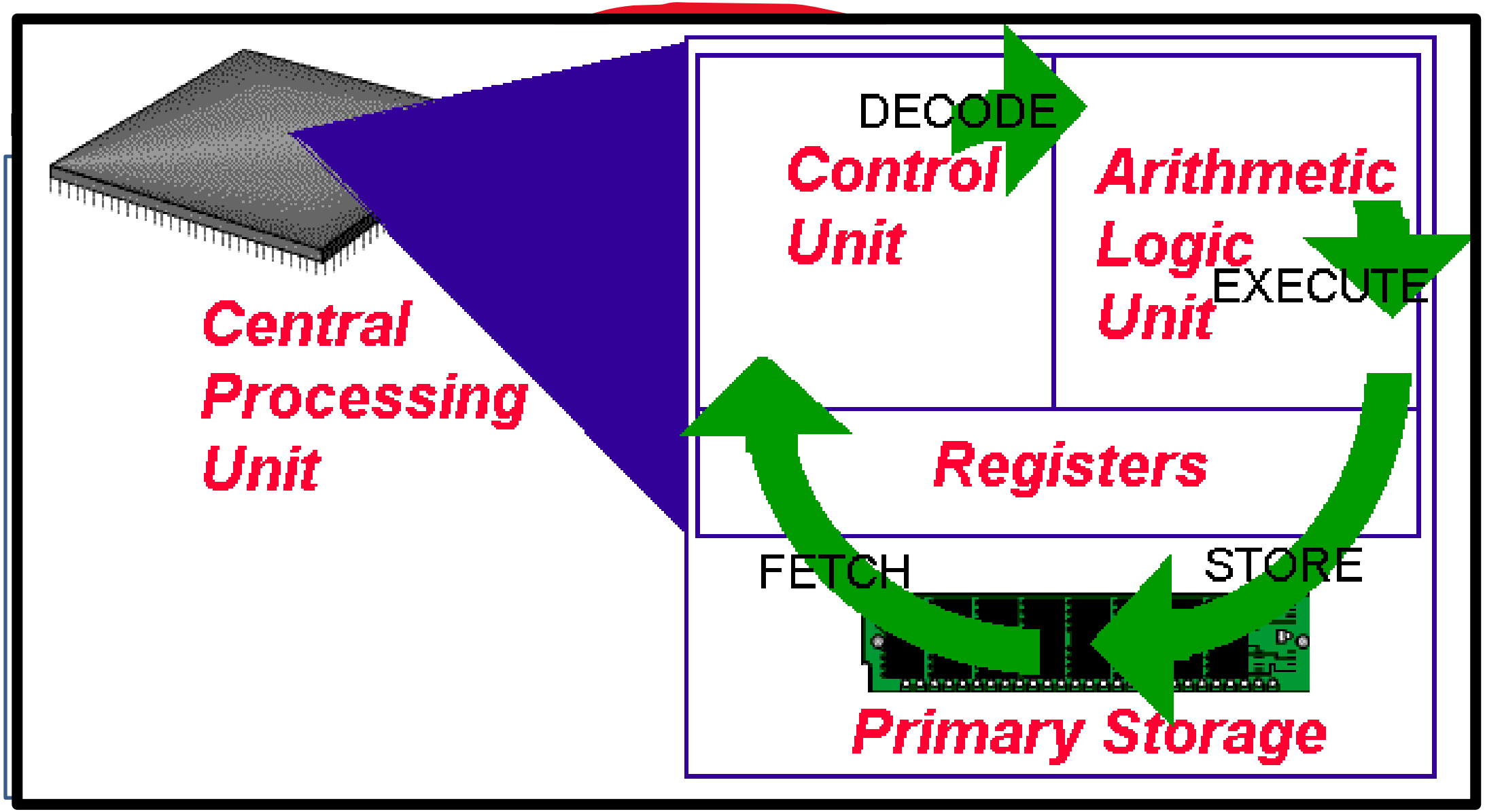
Memory

Instructions are stored here and prepared for the CPU

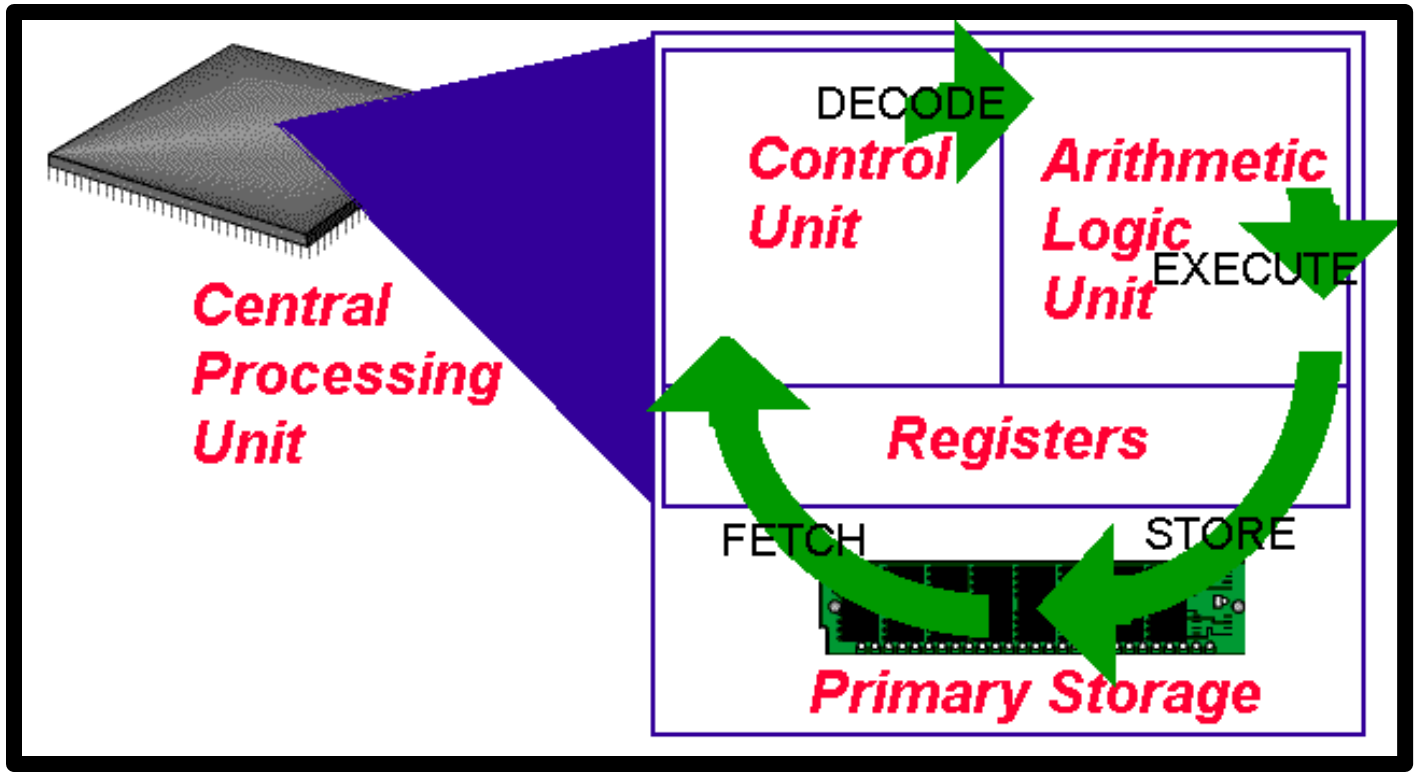
Temporary storage

RAM
Main Memory

I. Hardware

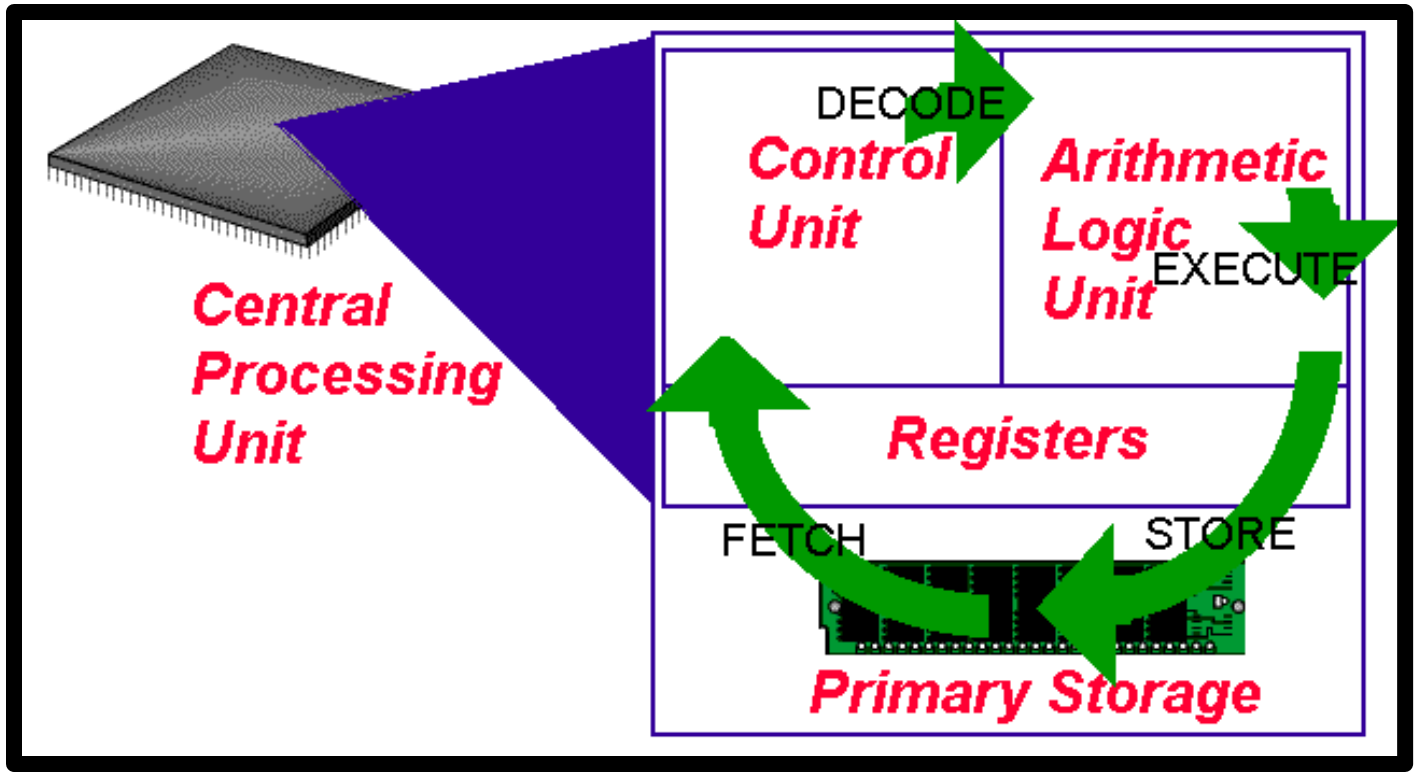

DECODE
Control Unit
Arithmetic Logic Unit
EXECUTE
Central Processing Unit
Registers
FETCH
STORE
Primary Storage

I. Hardware



This process happens really fast

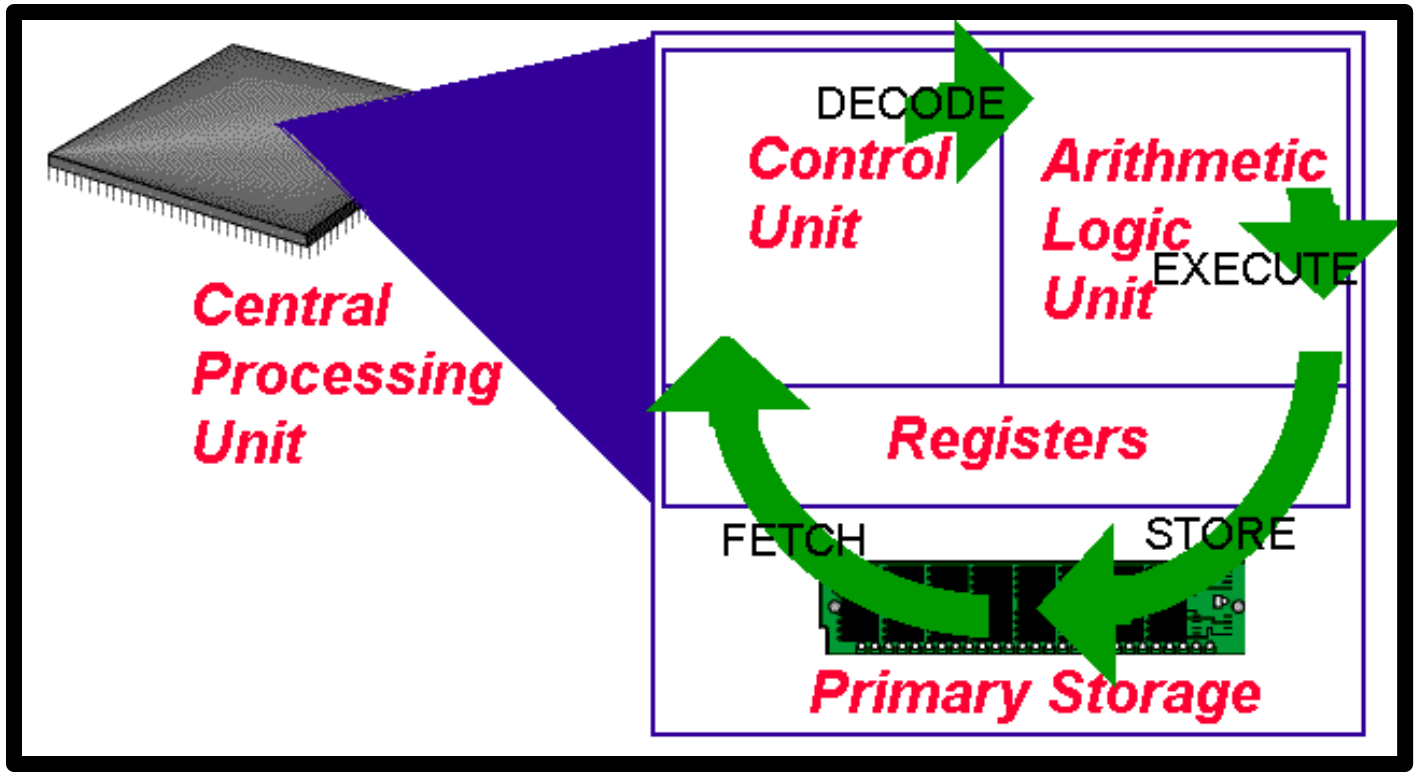
I. Hardware



This process happens really fast

… like ***really*** fast

I. Hardware

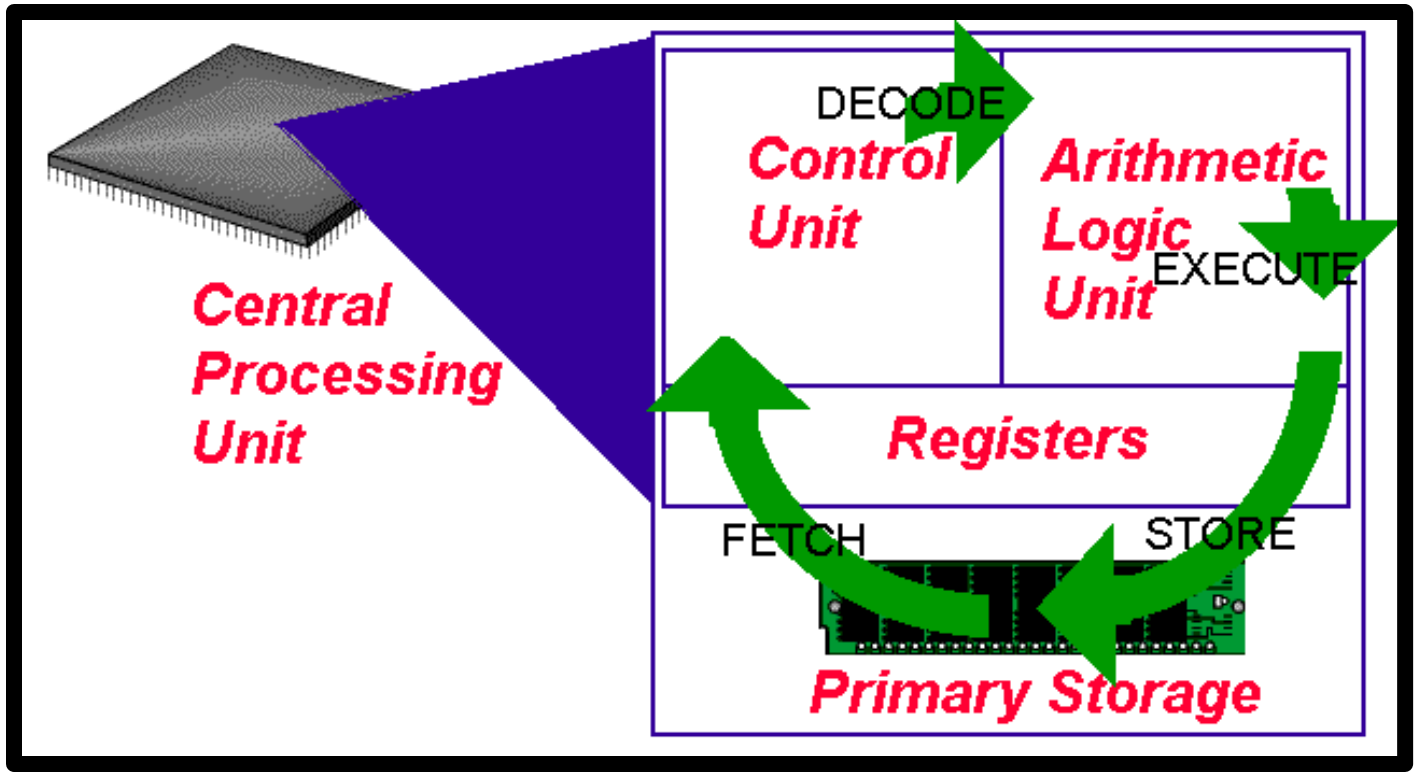


This process happens really fast

… like ***really*** fast

Computers can execute one instructions per clock cycle*

I. Hardware



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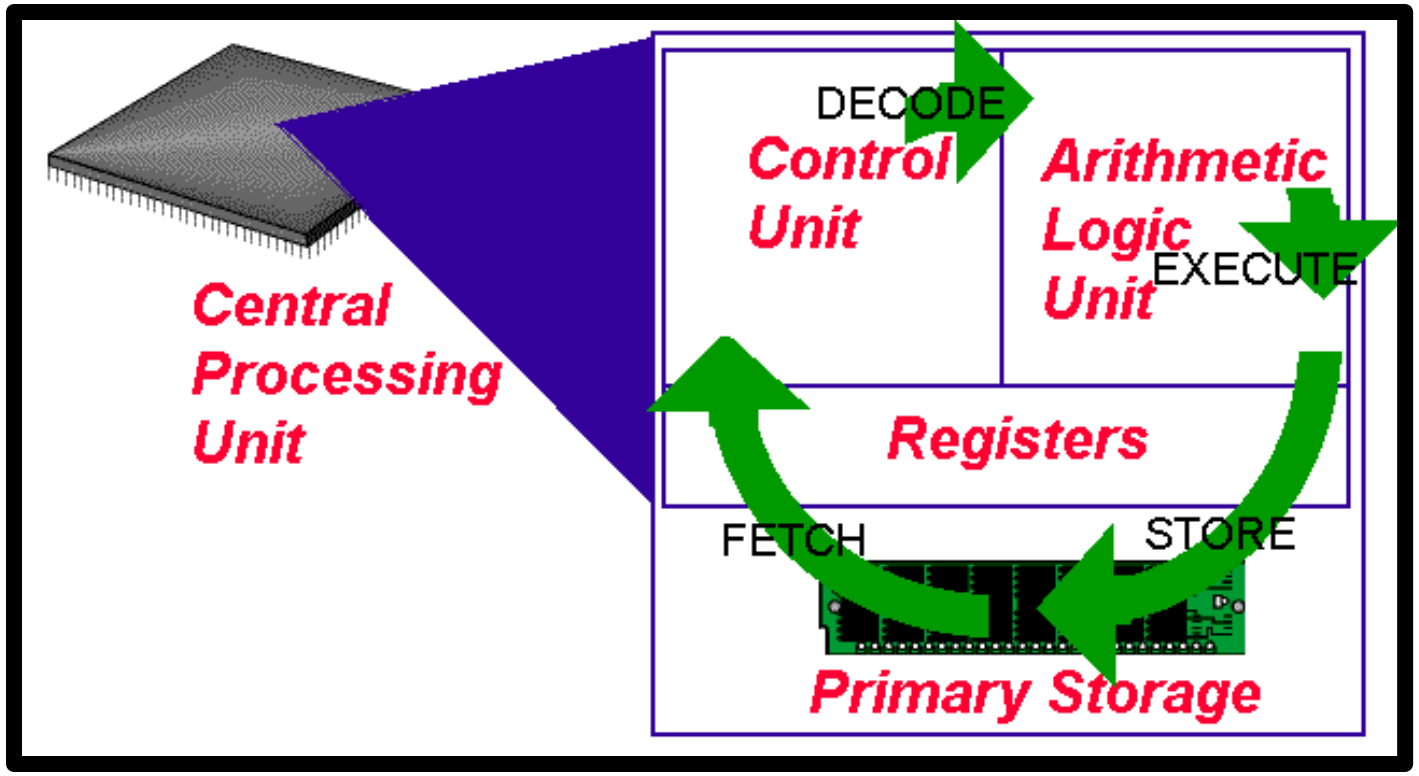
Computers can execute one instructions per clock cycle*

4GHz CPU speed = 4,000,000,000 clock cycles per second

I. Hardware



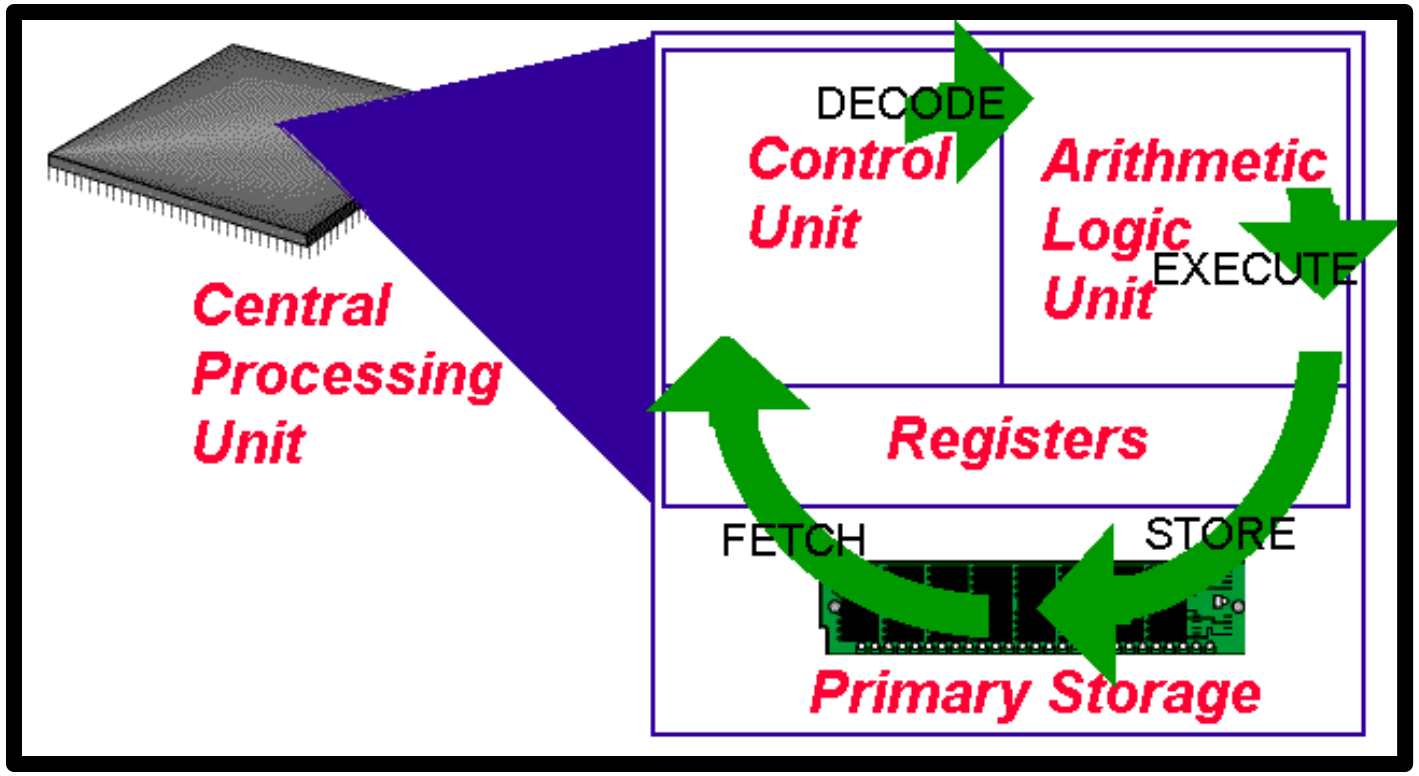
This process happens really fast

… like ***really*** fast

Computers can execute one or more instructions per clock cycle*

4GHz CPU speed = 4,000,000,000 clock cycles per second

I. Hardware



This process happens really fast

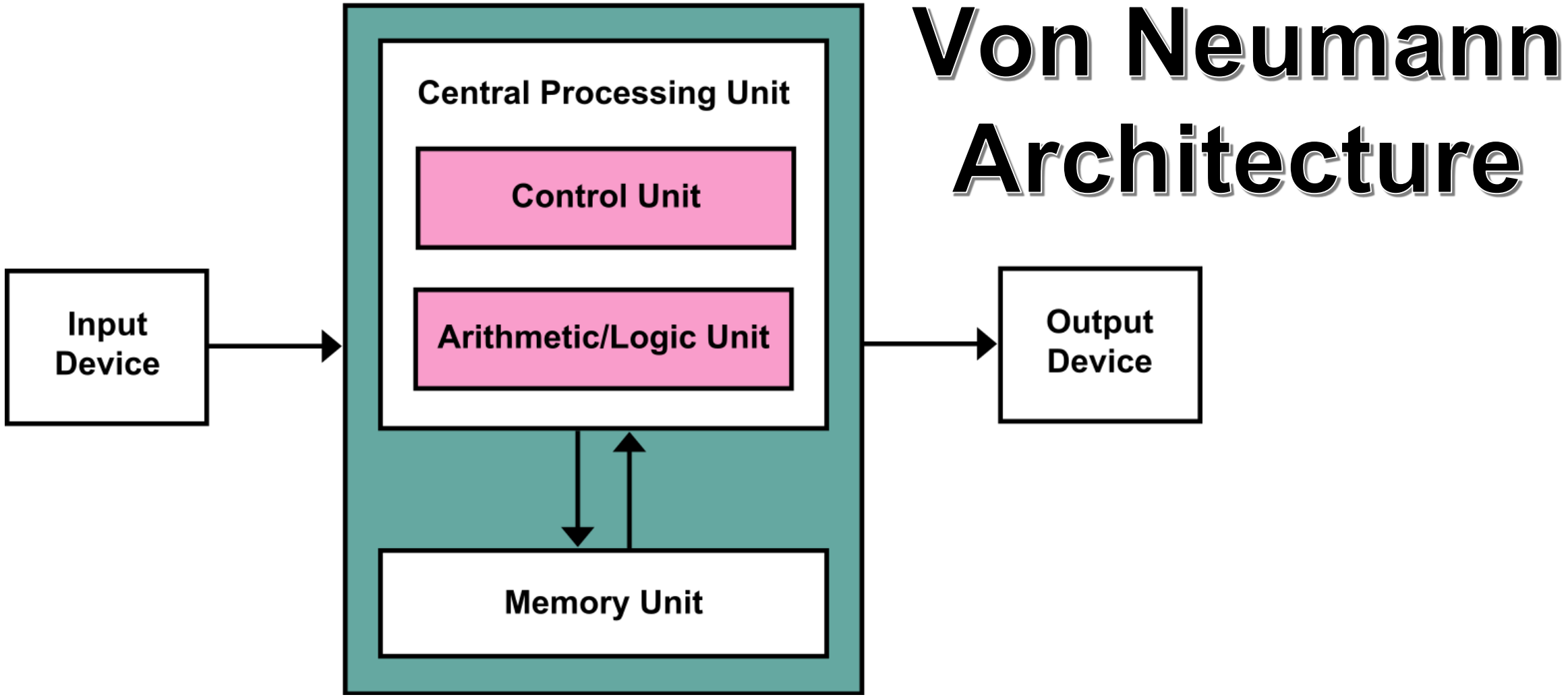
… like ***really*** fast

Computers can execute one instructions per clock cycle*

4GHz CPU speed = 4,000,000,000 clock cycles per second

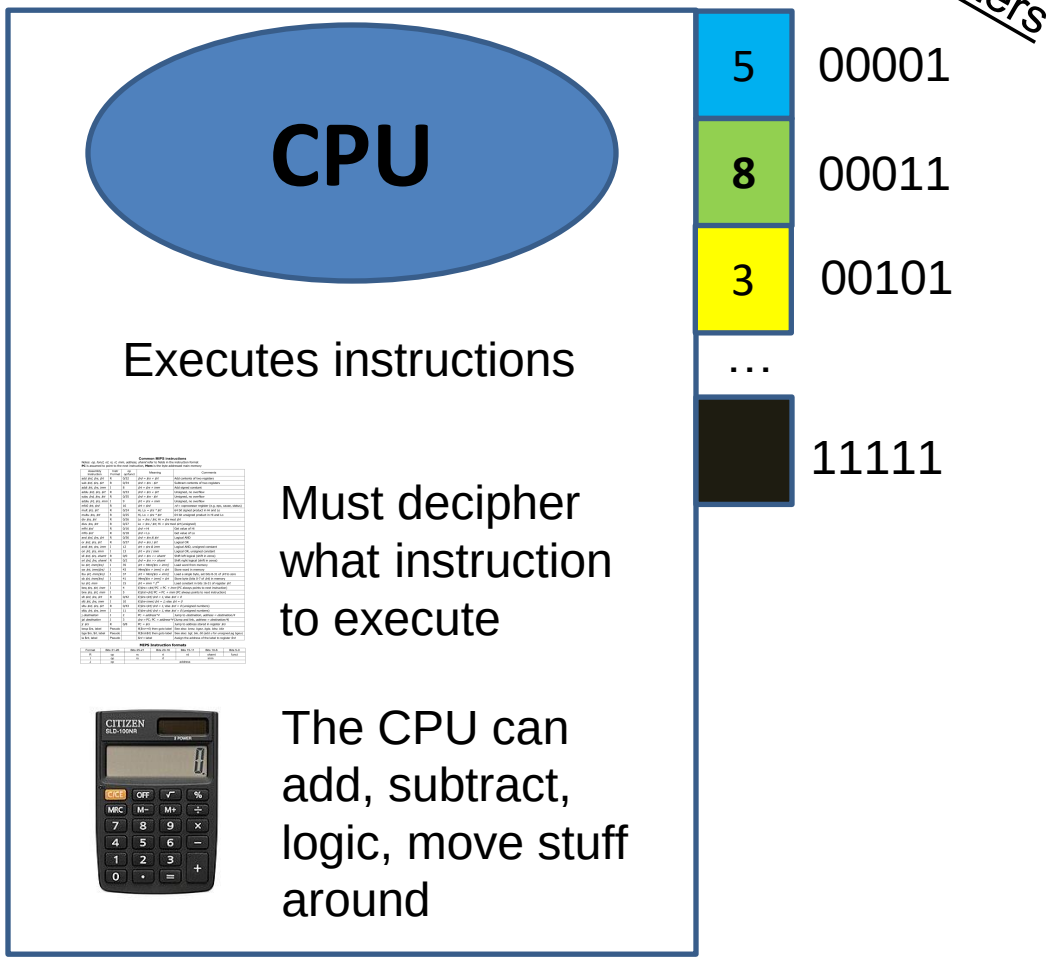
Multi-core systems are even faster

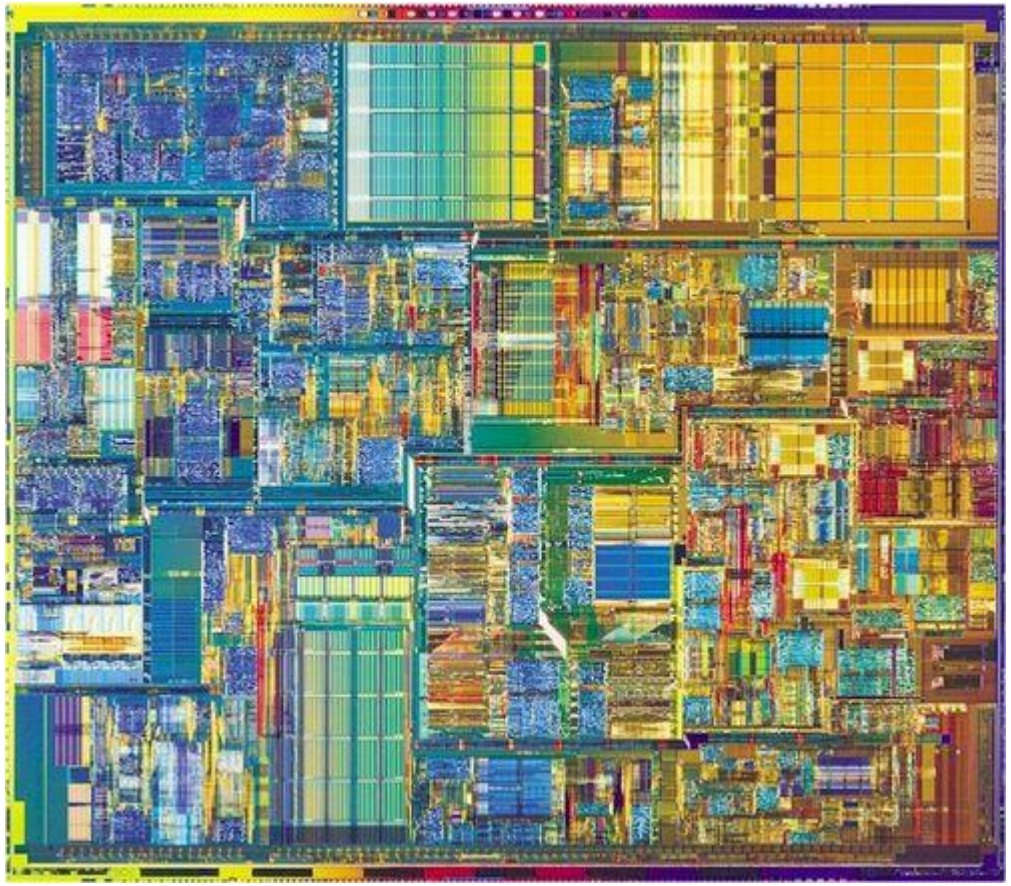
I. Hardware


Von Neumann Architecture
Central Processing Unit
Control Unit
Arithmetic/Logic Unit
Input Device
Output Device
Memory Unit

I. Hardware

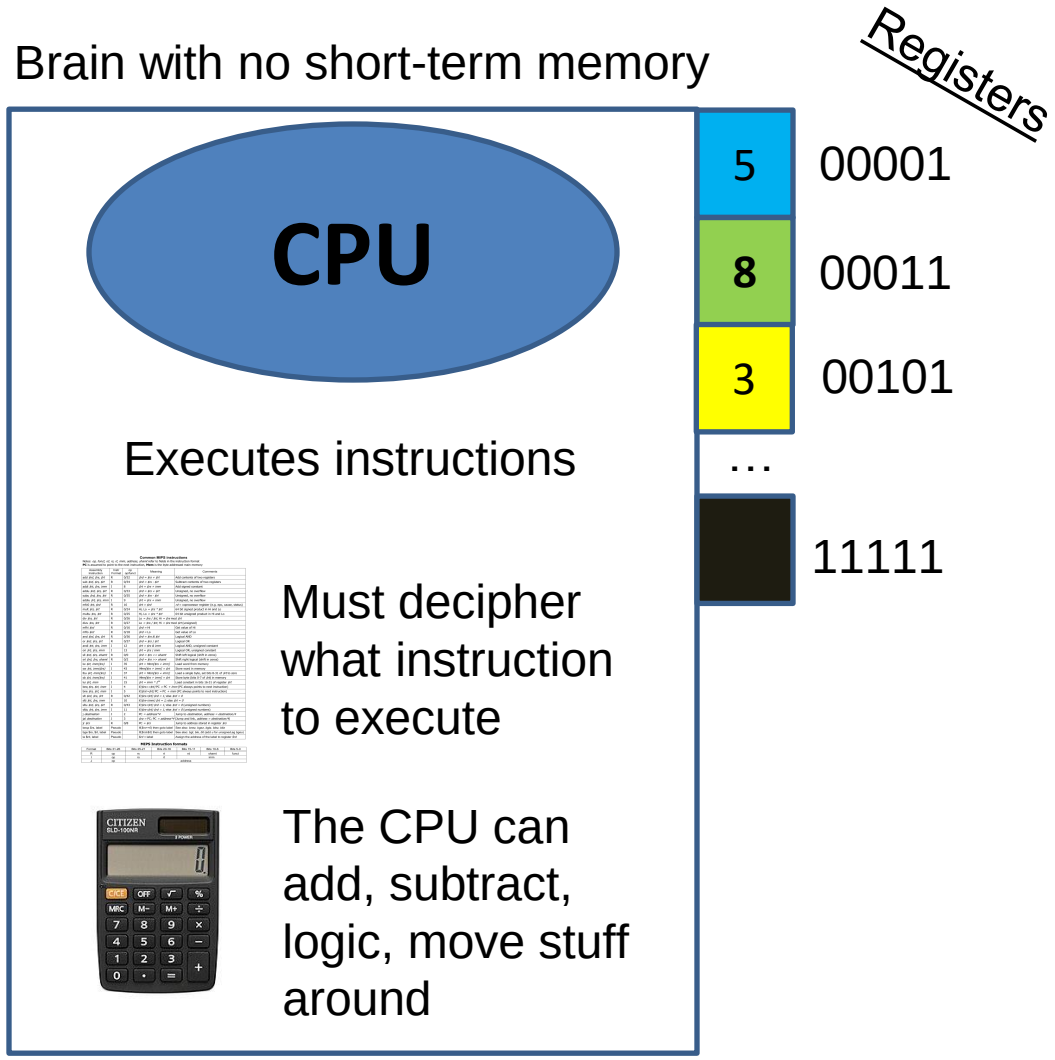
Brain with no short-term memory

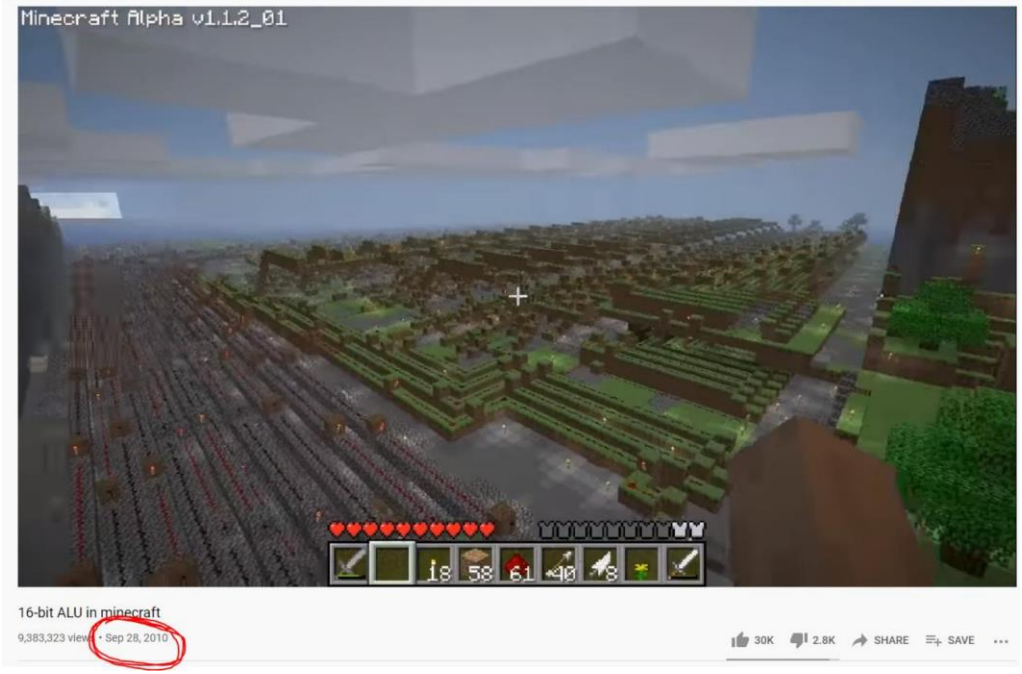




I. Hardware

Brain with no short-term memory





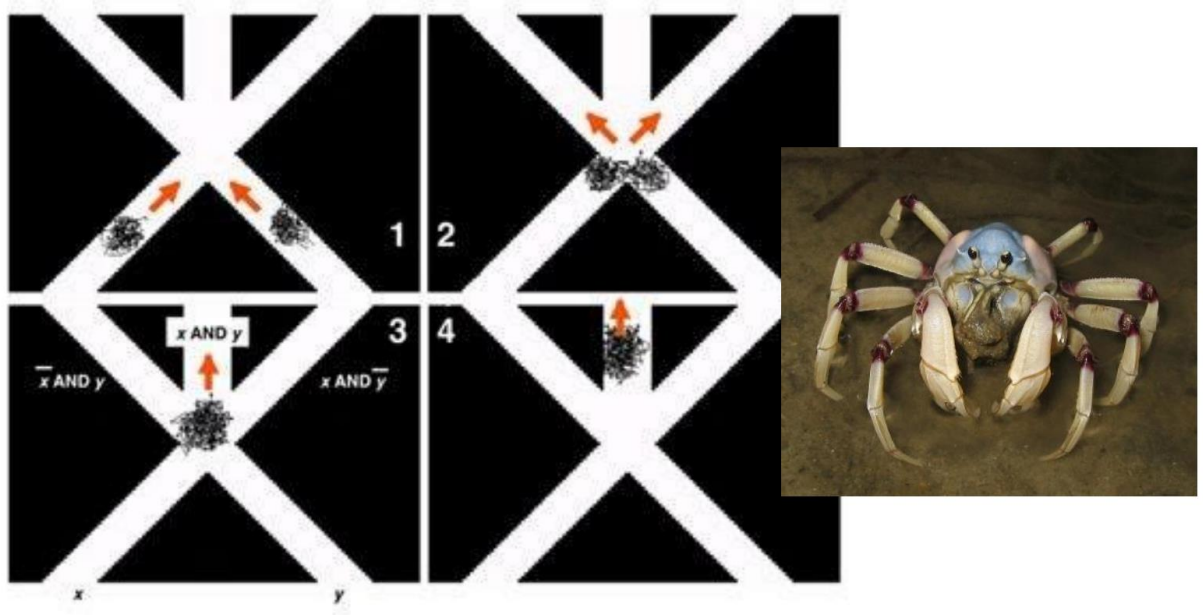
People have been able to create CPU components and fully functional, multi-core computers in games such as Minecraft

I. Hardware

WIRED STAFF BUSINESS 04.14.2012 03:20 PM

Computer Built Using Swarms Of Soldier Crabs

Computer scientists at Kobe University in Japan have built a computer that draws inspiration from the swarming behavior of soldier crabs. The computer is based on theories from the early 1980s that studies how it could be possible to build a computer out of billiard balls. Proposed by Edward Fredkin and Tommaso Toffoli, the mechanical [...]



I. Hardware

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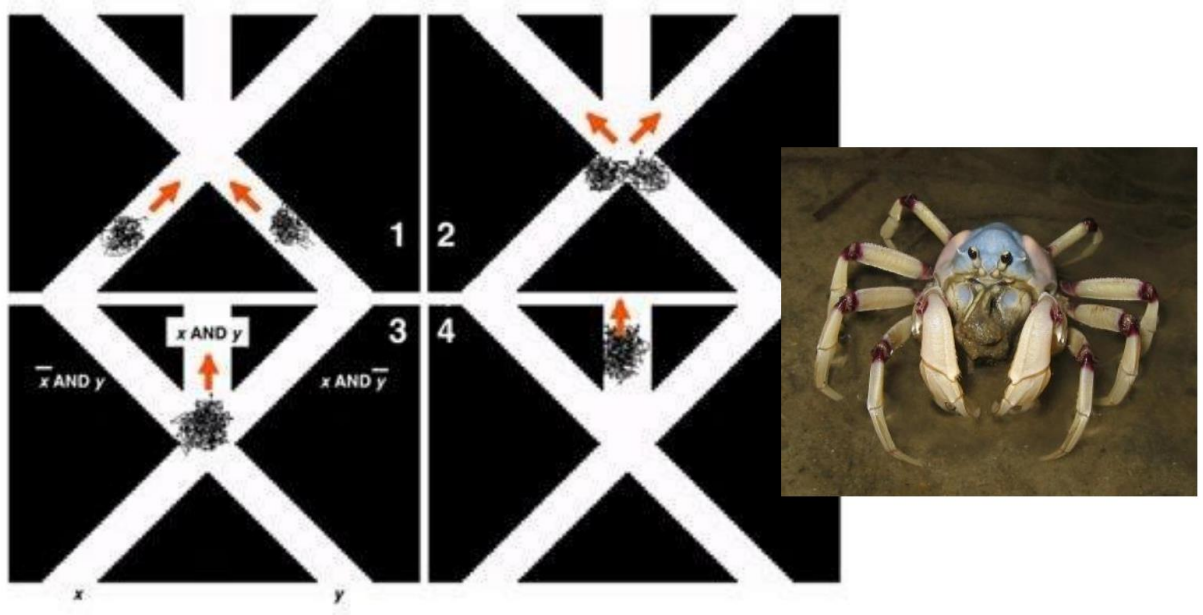

This is very real

Robust Soldier Crab Ball Gate

Yukio-Pegio Gunji
Yuta Nishiyama
Department of Earth and Planetary Sciences
Kobe University
Kobe 657-8501, Japan

Andrew Adamatzky
Unconventional Computing Centre
University of the West of England
Bristol, United Kingdom

Soldier crabs *Mictyris guinotae* exhibit pronounced swarming behavior. Swarms of the crabs are tolerant of perturbations. In computer models and laboratory experiments we demonstrate that swarms of soldier crabs can implement logical gates when placed in a geometrically constrained environment.

1. Introduction

All natural processes can be interpreted in terms of computations. To implement a logical gate in a chemical, physical, or biological spatially extended medium, Boolean variables must be assigned to disturbances, defects, or localizations traveling in the medium. These traveling patterns collide and the outcome of their collisions are converted

https://wpmedia.wolfram.com/uploads/sites/13/2018/02/20-2-2.pdf

I. Hardware

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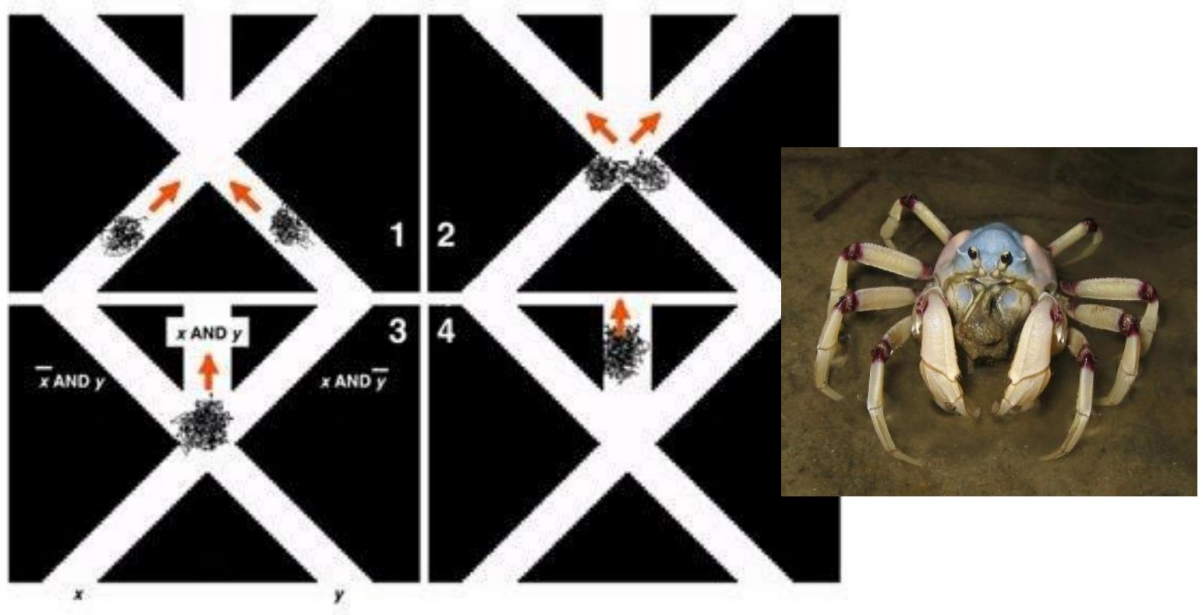
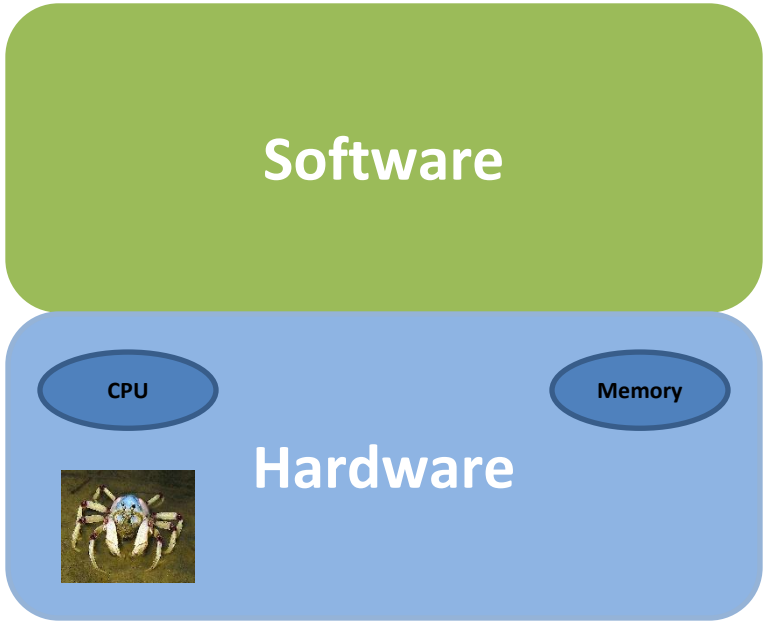


(In theory) If you wanted to play Doom (1993) using a CPU made from soldier crabs, you would need 22 million crabs

(source: twitter)

How does this happen?



From a high level, we will divide a computer system into two parts

I. **Hardware**
II. **Software**

II. Software

A sequence of instructions, or **program**,
that tells the computer how to work






II. Software

A sequence of instructions, or **program**, that tells the computer how to work

Humans write code in binary ?

01110011 00111011 01101111 01110101 01101100 01101110 01100111 01101001
01100001 00111011 01110011 01100111 01101110 00111011 01110011 01101011
01100100 01110110 01101110 01100001 00111011 01100100 01100111 01110100
00111011 01001111 01001100 01001011 01010011 01000100 01000110 00111011
01001100 01010011 01000100 01001000 01000111 01000100 01001100 00111011
01000110 01010011 01001011 01000111 01100110 01100100 01101000 01100100
01100110 01110011 01101000 01110011 01100110 01110111 01100101 01110111
01100101 01110010 01110111 01100101 01110010 01100110 01110110 01100111
01110011 01100100 01100010 01100100 01100110 01101000 01100010 01100100
01100110 01100100 01110011 01100100 01100110 01110011 01100100 01100110
01110011 01100100 01100110 01110011 01100100 01101000 01100100 01100110
01101000 01110011 01100100 01100110 01110011 01100100 01100110 01110011
01100100 01100110 01101000 01110100 01110010 01100100 01100110 01101000
01100100 01100110 01100111 01101000 01100110 01100111 01100111 01110100
01110010 01111001 01101000 01110100 01110010 01101000 01110010 01110011
01100110 01101000 01100111 01100110 01101000 01110011 01110010 01110100
01101000 01110100 01110010 01101000 01110011 01110100 01110010 01100100
01100110 01101000 01100100 01100110 01101000 01100100 01100110 01101000
01100100 01100110 01101000 01100101 01110111 01100110 01110011 01100100
01100110 01110010 01100101 01110111 01110100 01100111 01100101 01100111
01100100 01100110 01110011 01100111 01110011 01100110 0110010001110011
00111011 01101111 01110101 01101100 01101110 01100111 01101001 01100001
00111011 01110011 01100111 01101110 00111011 01110011 01101011 01100100
01110110 01101110 01100001 00111011 01100100 01100111 01110100 00111011
01001111 01001100 01001011 01010011 01000100 01000110 00111011 01001100
01010011 01000100 01001000 01000111 01000100 01001100 00111011 01000110
01010011 01001011 01000111 01100110 01100100 01101000 01100100 01100110
01110011 01101000 01110011 01100110 01110111 01100101 01110111 01100101
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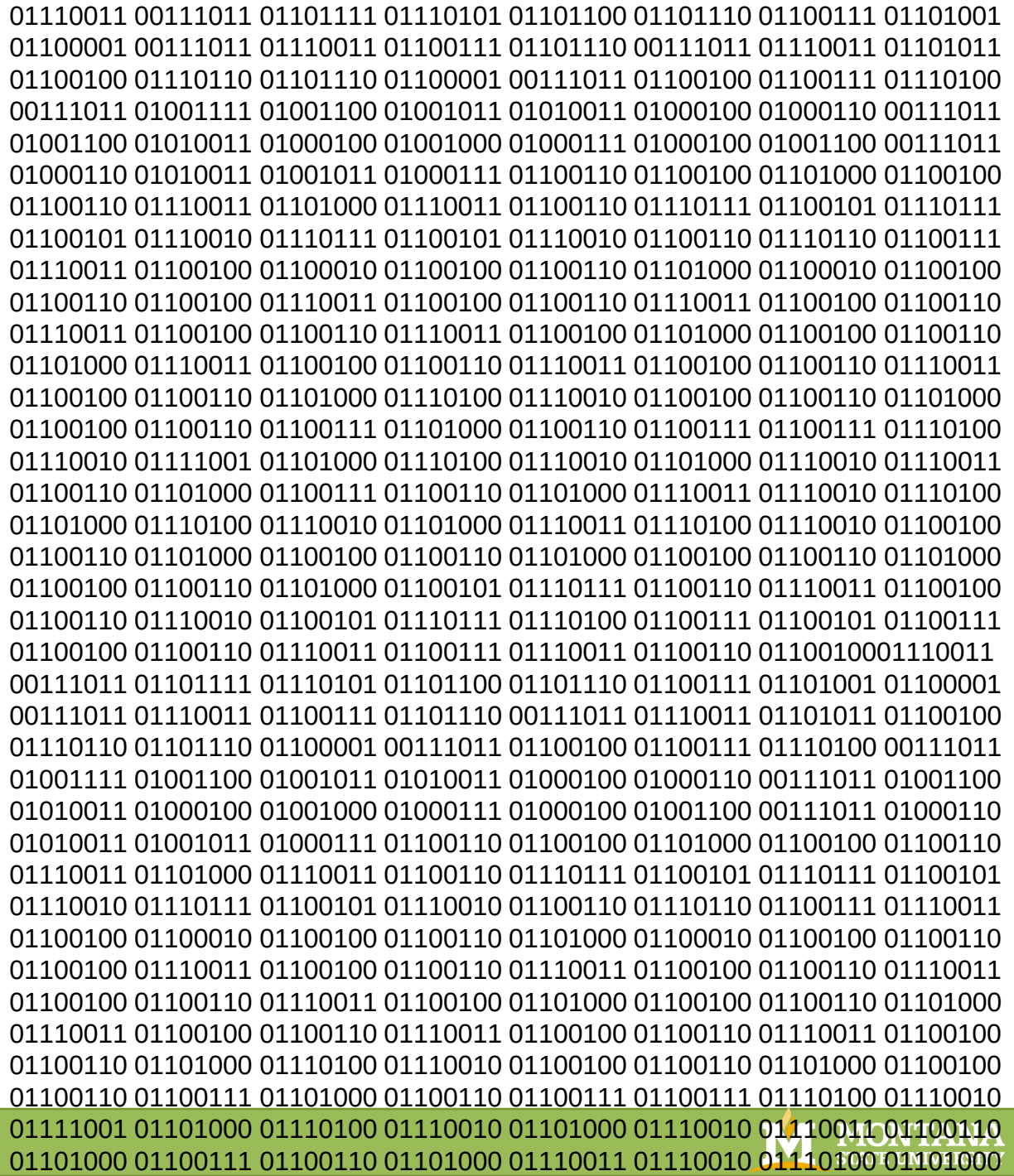
01100100 01110011 01100100 01100110 01110011 01100100 01100110 01110011
01100100 01100110 01110011 01100100 01101000 01100100 01100110 01101000
01110011 01100100 01100110 01110011 01100100 01100110 01110011 01100100
01100110 01101000 01110100 01110010 01100100 01100110 01101000 01100100
01100110 01100111 01101000 01100110 01100111 01100111 01110100 01110010
01111001 01101000 01110100 01110010 01101000 01110010 01110011 01100110
01101000 01100111 01100110 01101000 01110011 01110010 01110100 01101000

II. Software

A sequence of instructions, or **program**, that tells the computer how to work

Humans write code in binary ?

NO



II. Software

We write programs in a readable, higher-level language

```c
#include <stdio.h>

int main() {
        printf("Hello WOrld! \n");


        int x = 0;
        int y = 3;

        int z = x + y;

        printf("%d  %d  %d \n",x,y,z);
        return 0;
}
```

```python
class Person():
#method to initialize name and age attributes.
    def __init__(self,name, age):
        self.name = name
        self.age = age
#method to demonstrate what a person eats
    def eat(self):
        print(self.name.title() + "eats Matooke and rice")
        print("She is"+ str(self.age) + " years old")
    def drink(self):
        print("Drinks water")
#instantiating a class.
my_sister = Person("Haniifa", 30)
#Accessing the class method through the class object.
my_sister.eat()
```

II. Software

We need a way to convert **source** code to **binary**

```c
#include <stdio.h>

int main() {
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}
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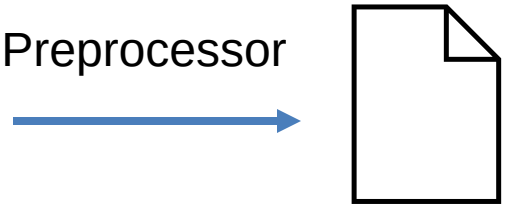
```

Preprocessor



- Removal of comments
- Expand Macros

**THIS PROCESS IS NOT TRUE FOR EVERY LANGUAGE

II. Software

We need a way to convert **source** code to **binary**

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int main() {
        printf("Hello WOrld! \n");


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        int z = x + y;

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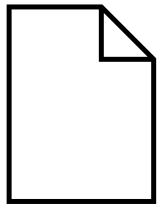
        return 0;

}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
   0:   f3 0f 1e fa             endbr64
   4:   55                      push   %rbp
   5:   48 89 e5                mov    %rsp,%rbp
   8:   48 83 ec 10             sub    $0x10,%rsp
   c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 13 <main+0x13>
  13:   e8 00 00 00 00          callq  18 <main+0x18>
  18:   c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%rbp)
  1f:   c7 45 f8 03 00 00 00    movl   $0x3,-0x8(%rbp)
  26:   8b 55 f4                mov    -0xc(%rbp),%edx
  29:   8b 45 f8                mov    -0x8(%rbp),%eax
  2c:   01 d0                   add    %edx,%eax
  2e:   89 45 fc                mov    %eax,-0x4(%rbp)
  31:   8b 4d fc                mov    -0x4(%rbp),%ecx
  34:   8b 55 f8                mov    -0x8(%rbp),%edx
  37:   8b 45 f4                mov    -0xc(%rbp),%eax
  3a:   89 c6                   mov    %eax,%esi
  3c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 43 <main+0x43>
  43:   b8 00 00 00 00          mov    $0x0,%eax
  48:   e8 00 00 00 00          callq  4d <main+0x4d>
  4d:   b8 00 00 00 00          mov    $0x0,%eax
  52:   c9                      leaveq
  53:   c3                      retq
```

- Converted to assembly code
- .s file

II. Software

We need a way to convert **source** code to **binary**

```c
#include <stdio.h>

int main() {
        printf("Hello WOrld! \n");


        int x = 0;
        int y = 3;

        int z = x + y;

        printf("%d  %d  %d \n",x,y,z);
        return 0;

}
```

Preprocessor

- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
   0:   f3 0f 1e fa             endbr64
   4:   55                      push   %rbp
   5:   48 89 e5                mov    %rsp,%rbp
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   c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 13 <main+0x13>
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  18:   c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%rbp)
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  26:   8b 55 f4                mov    -0xc(%rbp),%edx
  29:   8b 45 f8                mov    -0x8(%rbp),%eax
  2c:   01 d0                   add    %edx,%eax
  2e:   89 45 fc                mov    %eax,-0x4(%rbp)
  31:   8b 4d fc                mov    -0x4(%rbp),%ecx
  34:   8b 55 f8                mov    -0x8(%rbp),%edx
  37:   8b 45 f4                mov    -0xc(%rbp),%eax
  3a:   89 c6                   mov    %eax,%esi
  3c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 43 <main+0x43>
  43:   b8 00 00 00 00          mov    $0x0,%eax
  48:   e8 00 00 00 00          callq  4d <main+0x4d>
  4d:   b8 00 00 00 00          mov    $0x0,%eax
  52:   c9                      leaveq
  53:   c3                      retq
```

- Converted to assembly code
- .s file

Assembler

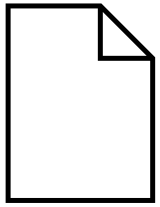
```
1           00000000 00000100 0000000000000000
2  01011110 00001100 11000010 0000000000000010
3           11101111 00010110 0000000000000101
4           11101111 10011110 0000000000001011
5  11111000 10101101 11011111 0000000000010010
6           01100010 11011111 0000000000010101
7  11101111 00000010 11111011 0000000000010111
8  11110100 10101101 11011111 0000000000011110
9  00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15          01010000 11010100 0000000000111011
16                   00000100 0000000000111101
```

II. Software

We need a way to convert **source** code to **binary**

```c
#include <stdio.h>

int main() {
        printf("Hello WOrld! \n");


        int x = 0;
        int y = 3;

        int z = x + y;

        printf("%d  %d  %d \n",x,y,z);
        return 0;

}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
   0:   f3 0f 1e fa             endbr64
   4:   55                      push   %rbp
   5:   48 89 e5                mov    %rsp,%rbp
   8:   48 83 ec 10             sub    $0x10,%rsp
   c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 13 <main+0x13>
  13:   e8 00 00 00 00          callq  18 <main+0x18>
  18:   c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%rbp)
  1f:   c7 45 f8 03 00 00 00    movl   $0x3,-0x8(%rbp)
  26:   8b 55 f4                mov    -0xc(%rbp),%edx
  29:   8b 45 f8                mov    -0x8(%rbp),%eax
  2c:   01 d0                   add    %edx,%eax
  2e:   89 45 fc                mov    %eax,-0x4(%rbp)
  31:   8b 4d fc                mov    -0x4(%rbp),%ecx
  34:   8b 55 f8                mov    -0x8(%rbp),%edx
  37:   8b 45 f4                mov    -0xc(%rbp),%eax
  3a:   89 c6                   mov    %eax,%esi
  3c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 43 <main+0x43>
  43:   b8 00 00 00 00          mov    $0x0,%eax
  48:   e8 00 00 00 00          callq  4d <main+0x4d>
  4d:   b8 00 00 00 00          mov    $0x0,%eax
  52:   c9                      leaveq
  53:   c3                      retq
```

- Converted to assembly code
- .s file

Assembler

```
1           00000000 00000100 0000000000000000
2  01011110 00001100 11000010 0000000000000010
3           11101111 00010110 0000000000000101
4           11101111 10011110 0000000000001011
5  11111000 10101101 11011111 0000000000010010
6           01100010 11011111 0000000000010101
7  11101111 00000010 11111011 0000000000010111
8  11110100 10101101 11011111 0000000000011110
9  00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15          01010000 11010100 0000000000111011
16                   00000100 0000000000111101
```

Linker

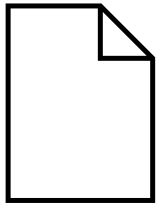
We still need to resolve function calls
i.e. `printf, sleep, sqrt, etc`

II. Software

We need a way to convert **source** code to **binary**

```c
#include <stdio.h>

int main() {
        printf("Hello WOrld! \n");


        int x = 0;
        int y = 3;

        int z = x + y;

        printf("%d  %d  %d \n",x,y,z);
        return 0;

}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
   0:   f3 0f 1e fa             endbr64
   4:   55                      push   %rbp
   5:   48 89 e5                mov    %rsp,%rbp
   8:   48 83 ec 10             sub    $0x10,%rsp
   c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 13 <main+0x13>
  13:   e8 00 00 00 00          callq  18 <main+0x18>
  18:   c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%rbp)
  1f:   c7 45 f8 03 00 00 00    movl   $0x3,-0x8(%rbp)
  26:   8b 55 f4                mov    -0xc(%rbp),%edx
  29:   8b 45 f8                mov    -0x8(%rbp),%eax
  2c:   01 d0                   add    %edx,%eax
  2e:   89 45 fc                mov    %eax,-0x4(%rbp)
  31:   8b 4d fc                mov    -0x4(%rbp),%ecx
  34:   8b 55 f8                mov    -0x8(%rbp),%edx
  37:   8b 45 f4                mov    -0xc(%rbp),%eax
  3a:   89 c6                   mov    %eax,%esi
  3c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 43 <main+0x43>
  43:   b8 00 00 00 00          mov    $0x0,%eax
  48:   e8 00 00 00 00          callq  4d <main+0x4d>
  4d:   b8 00 00 00 00          mov    $0x0,%eax
  52:   c9                      leaveq
  53:   c3                      retq
```

Assembler

- Converted to assembly code
- .s file

Two methods:

Program A

Library 1 string.h

Library 2 stdio.h

Linker

```
1           00000000 00000100 0000000000000000
2  01011110 00001100 11000010 0000000000000010
3           11101111 00010110 0000000000000101
4           11101111 10011110 0000000000001011
5  11111000 10101101 11011111 0000000000010010
6           01100010 11011111 0000000000010101
7  11101111 00000010 11111011 0000000000010111
8  11110100 10101101 11011111 0000000000011110
9  00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
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13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15          01010000 11010100 0000000000111011
16                   00000100 0000000000111101
```
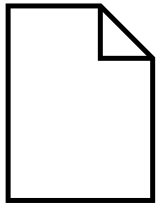

II. Software

We need a way to convert **source** code to **binary**

```c
#include <stdio.h>

int main() {
        printf("Hello WOrld! \n");


        int x = 0;
        int y = 3;

        int z = x + y;

        printf("%d  %d  %d \n",x,y,z);
        return 0;

}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
   0:   f3 0f 1e fa             endbr64
   4:   55                      push   %rbp
   5:   48 89 e5                mov    %rsp,%rbp
   8:   48 83 ec 10             sub    $0x10,%rsp
   c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 13 <main+0x13>
  13:   e8 00 00 00 00          callq  18 <main+0x18>
  18:   c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%rbp)
  1f:   c7 45 f8 03 00 00 00    movl   $0x3,-0x8(%rbp)
  26:   8b 55 f4                mov    -0xc(%rbp),%edx
  29:   8b 45 f8                mov    -0x8(%rbp),%eax
  2c:   01 d0                   add    %edx,%eax
  2e:   89 45 fc                mov    %eax,-0x4(%rbp)
  31:   8b 4d fc                mov    -0x4(%rbp),%ecx
  34:   8b 55 f8                mov    -0x8(%rbp),%edx
  37:   8b 45 f4                mov    -0xc(%rbp),%eax
  3a:   89 c6                   mov    %eax,%esi
  3c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 43 <main+0x43>
  43:   b8 00 00 00 00          mov    $0x0,%eax
  48:   e8 00 00 00 00          callq  4d <main+0x4d>
  4d:   b8 00 00 00 00          mov    $0x0,%eax
  52:   c9                      leaveq
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```

Assembler

- Converted to assembly code
- .s file

```
1           00000000 00000100 0000000000000000
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9  00000011 10100010 11011111 0000000000100001
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11 01111110 11110100 10101101
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13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15          01010000 11010100 0000000000111011
16                   00000100 0000000000111101
```

Linker

.exe

Static Linking- required code and data copied into executable at compile time

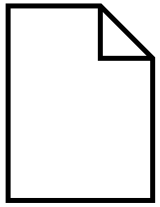
**THIS PROCESS IS NOT TRUE FOR EVERY LANGUAGE

II. Software

We need a way to convert **source** code to **binary**

```c
#include <stdio.h>

int main() {
        printf("Hello WOrld! \n");


        int x = 0;
        int y = 3;

        int z = x + y;

        printf("%d  %d  %d \n",x,y,z);
        return 0;

}
```

Preprocessor



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Compiler

```
0000000000000000 <main>:
   0:   f3 0f 1e fa             endbr64
   4:   55                      push   %rbp
   5:   48 89 e5                mov    %rsp,%rbp
   8:   48 83 ec 10             sub    $0x10,%rsp
   c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 13 <main+0x13>
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  18:   c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%rbp)
  1f:   c7 45 f8 03 00 00 00    movl   $0x3,-0x8(%rbp)
  26:   8b 55 f4                mov    -0xc(%rbp),%edx
  29:   8b 45 f8                mov    -0x8(%rbp),%eax
  2c:   01 d0                   add    %edx,%eax
  2e:   89 45 fc                mov    %eax,-0x4(%rbp)
  31:   8b 4d fc                mov    -0x4(%rbp),%ecx
  34:   8b 55 f8                mov    -0x8(%rbp),%edx
  37:   8b 45 f4                mov    -0xc(%rbp),%eax
  3a:   89 c6                   mov    %eax,%esi
  3c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 43 <main+0x43>
  43:   b8 00 00 00 00          mov    $0x0,%eax
  48:   e8 00 00 00 00          callq  4d <main+0x4d>
  4d:   b8 00 00 00 00          mov    $0x0,%eax
  52:   c9                      leaveq
  53:   c3                      retq
```

Assembler

- Converted to assembly code
- .s file

Two methods:

Program A

Library 1 string.h

Library 2 stdio.h

Linker

```
1           00000000 00000100 0000000000000000
2  01011110 00001100 11000010 0000000000000010
3           11101111 00010110 0000000000000101
4           11101111 10011110 0000000000001011
5  11111000 10101101 11011111 0000000000010010
6           01100010 11011111 0000000000010101
7  11101111 00000010 11111011 0000000000010111
8  11110100 10101101 11011111 0000000000011110
9  00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15          01010000 11010100 0000000000111011
16                   00000100 0000000000111101
```

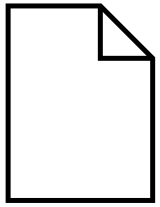
**THIS PROCESS IS NOT TRUE FOR EVERY LANGUAGE

II. Software

We need a way to convert **source** code to **binary**

```c
#include <stdio.h>

int main() {
        printf("Hello WOrld! \n");


        int x = 0;
        int y = 3;

        int z = x + y;

        printf("%d  %d  %d \n",x,y,z);
        return 0;

}
```

Preprocessor



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Compiler

```
0000000000000000 <main>:
   0:   f3 0f 1e fa             endbr64
   4:   55                      push   %rbp
   5:   48 89 e5                mov    %rsp,%rbp
   8:   48 83 ec 10             sub    $0x10,%rsp
   c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 13 <main+0x13>
  13:   e8 00 00 00 00          callq  18 <main+0x18>
  18:   c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%rbp)
  1f:   c7 45 f8 03 00 00 00    movl   $0x3,-0x8(%rbp)
  26:   8b 55 f4                mov    -0xc(%rbp),%edx
  29:   8b 45 f8                mov    -0x8(%rbp),%eax
  2c:   01 d0                   add    %edx,%eax
  2e:   89 45 fc                mov    %eax,-0x4(%rbp)
  31:   8b 4d fc                mov    -0x4(%rbp),%ecx
  34:   8b 55 f8                mov    -0x8(%rbp),%edx
  37:   8b 45 f4                mov    -0xc(%rbp),%eax
  3a:   89 c6                   mov    %eax,%esi
  3c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 43 <main+0x43>
  43:   b8 00 00 00 00          mov    $0x0,%eax
  48:   e8 00 00 00 00          callq  4d <main+0x4d>
  4d:   b8 00 00 00 00          mov    $0x0,%eax
  52:   c9                      leaveq
  53:   c3                      retq
```

Assembler

- Converted to assembly code
- .s file

Two methods:

string.h — Library 1

stdio.h — Library 2

Program A

.exe

Linker

```
1           00000000 00000100 0000000000000000
2  01011110 00001100 11000010 0000000000000010
3           11101111 00010110 0000000000000101
4           11101111 10011110 0000000000001011
5  11111000 10101101 11011111 0000000000010010
6           01100010 11011111 0000000000010101
7  11101111 00000010 11111011 0000000000010111
8  11110100 10101101 11011111 0000000000011110
9  00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15          01010000 11010100 0000000000111011
16                   00000100 0000000000111101
```

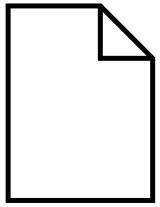
Dynamic Linking - required code and data is linked to executable at runtime

II. Software

We need a way to convert **source** code to **binary**

```c
#include <stdio.h>

int main() {
        printf("Hello WOrld! \n");


        int x = 0;
        int y = 3;

        int z = x + y;

        printf("%d  %d  %d \n",x,y,z);
        return 0;

}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
   0:   f3 0f 1e fa             endbr64
   4:   55                      push   %rbp
   5:   48 89 e5                mov    %rsp,%rbp
   8:   48 83 ec 10             sub    $0x10,%rsp
   c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 13 <main+0x13>
  13:   e8 00 00 00 00          callq  18 <main+0x18>
  18:   c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%rbp)
  1f:   c7 45 f8 03 00 00 00    movl   $0x3,-0x8(%rbp)
  26:   8b 55 f4                mov    -0xc(%rbp),%edx
  29:   8b 45 f8                mov    -0x8(%rbp),%eax
  2c:   01 d0                   add    %edx,%eax
  2e:   89 45 fc                mov    %eax,-0x4(%rbp)
  31:   8b 4d fc                mov    -0x4(%rbp),%ecx
  34:   8b 55 f8                mov    -0x8(%rbp),%edx
  37:   8b 45 f4                mov    -0xc(%rbp),%eax
  3a:   89 c6                   mov    %eax,%esi
  3c:   48 8d 3d 00 00 00 00    lea    0x0(%rip),%rdi        # 43 <main+0x43>
  43:   b8 00 00 00 00          mov    $0x0,%eax
  48:   e8 00 00 00 00          callq  4d <main+0x4d>
  4d:   b8 00 00 00 00          mov    $0x0,%eax
  52:   c9                      leaveq
  53:   c3                      retq
```

- Converted to assembly code
- .s file

Assembler

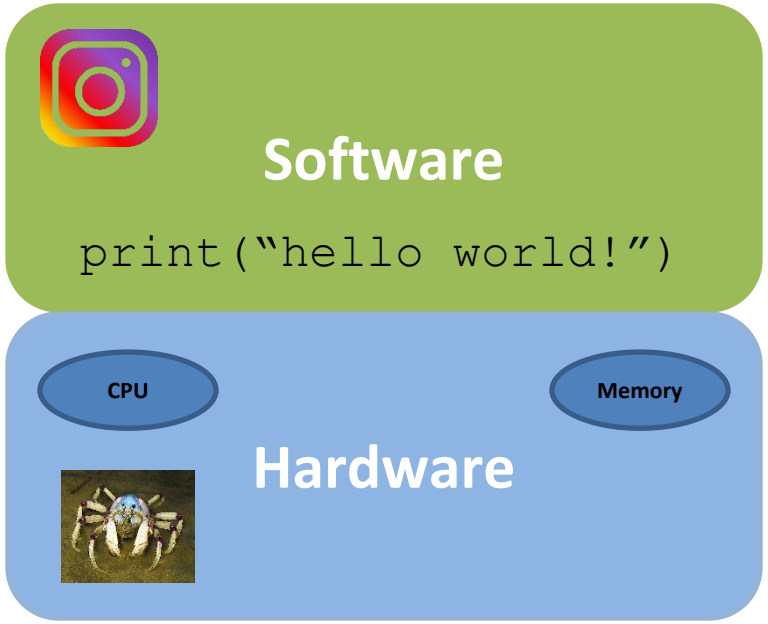
```
1           00000000 00000100 0000000000000000
2  01011110 00001100 11000010 0000000000000010
3           11101111 00010110 0000000000000101
4           11101111 10011110 0000000000001011
5  11111000 10101101 11011111 0000000000010010
6           01100010 11011111 0000000000010101
7  11101111 00000010 11111011 0000000000010111
8  11110100 10101101 11011111 0000000000011110
9  00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15          01010000 11010100 0000000000111011
16                   00000100 0000000000111101
```

Linker

.exe

To be continued

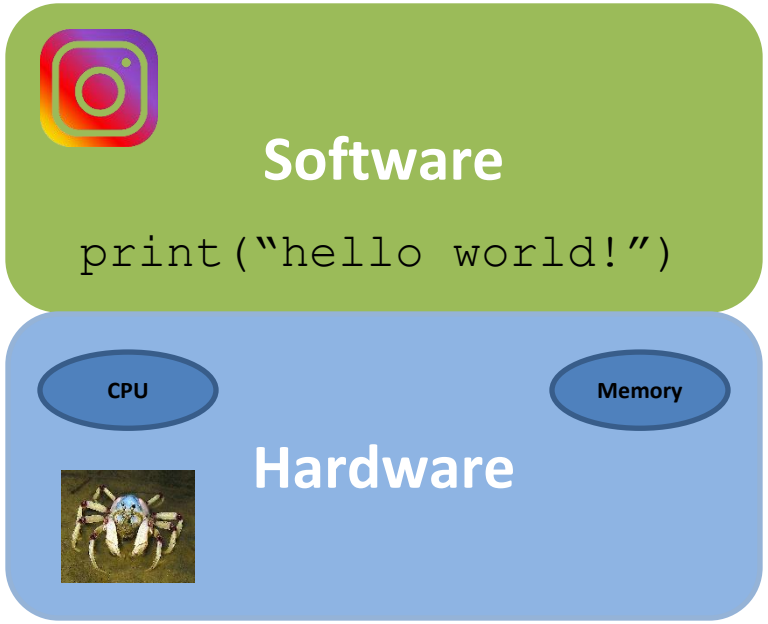
How does this happen?



From a high level, we will divide a computer system into two parts

I. **Hardware**
II. **Software**

How does this happen?

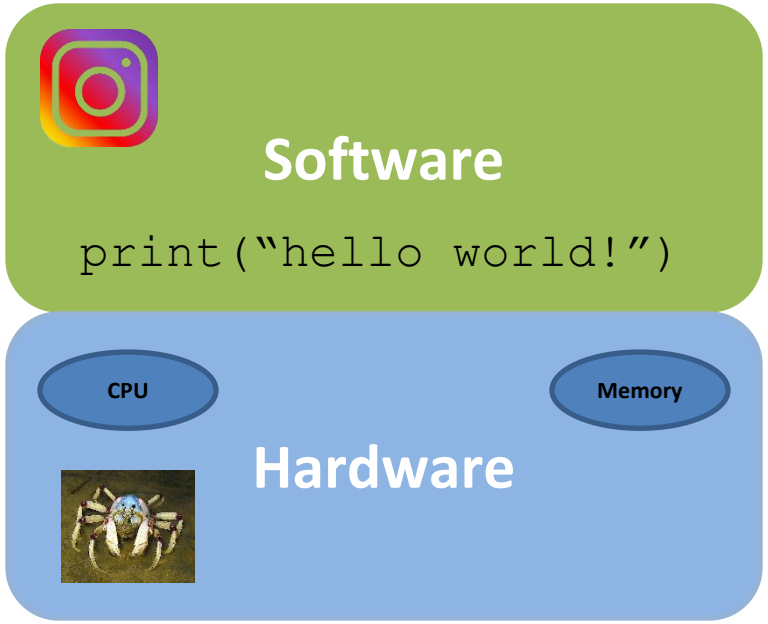


From a high level, we will divide a computer system into two parts

I. **Hardware**
II. **Software**

Software is nothing without hardware

How does this happen?

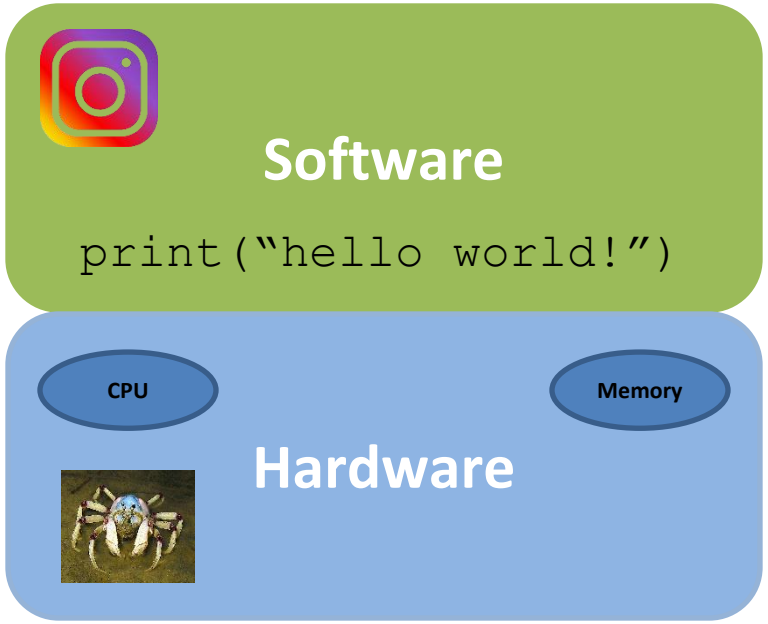


From a high level, we will divide a computer system into two parts

I. **Hardware**
II. **Software**

Hardware is *mostly* nothing without software

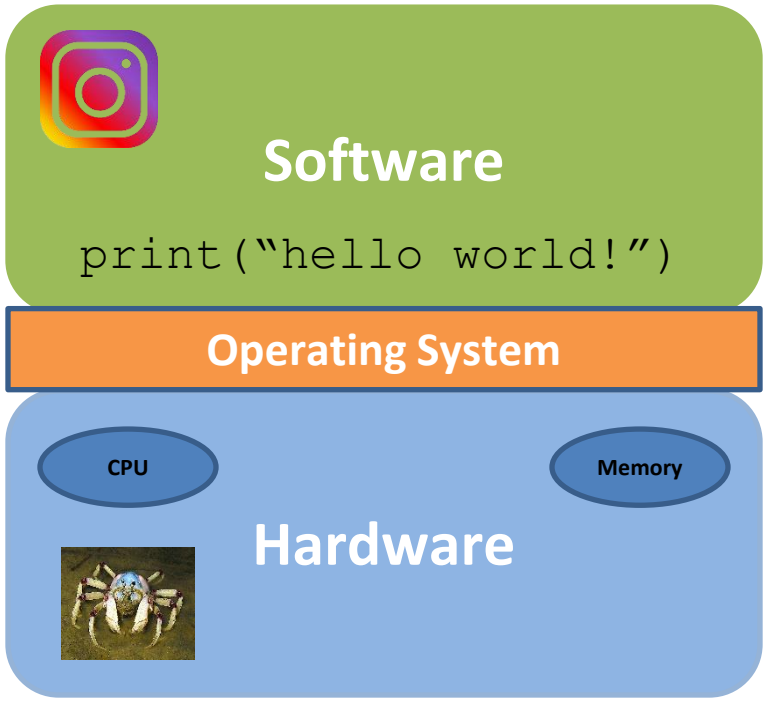
How does this happen?



From a high level, we will divide a computer system into two parts

I. **Hardware**
II. **Software**
III. **???**

How does this happen?

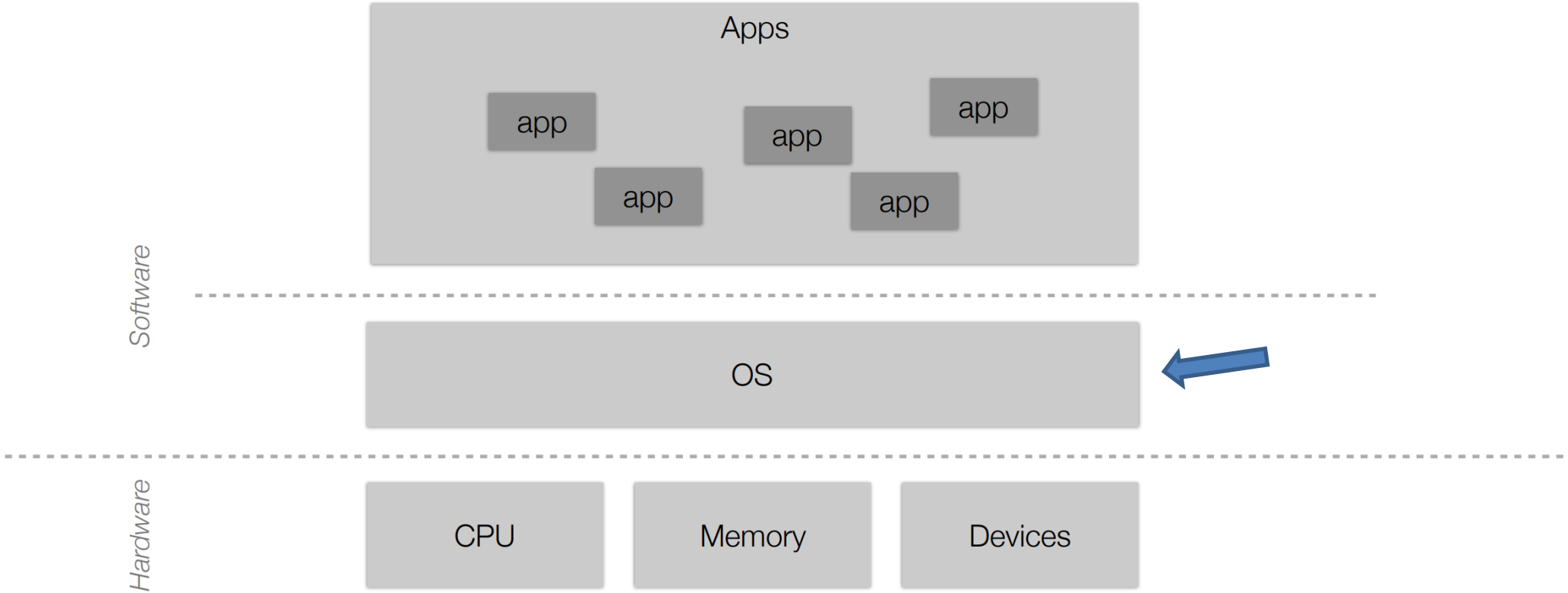


From a high level, we will divide a computer system into two parts

I. **Hardware**
II. **Software**
III. **Operating System**

Typical Layers of a Computer

The **operating system** is a vital component of a computer

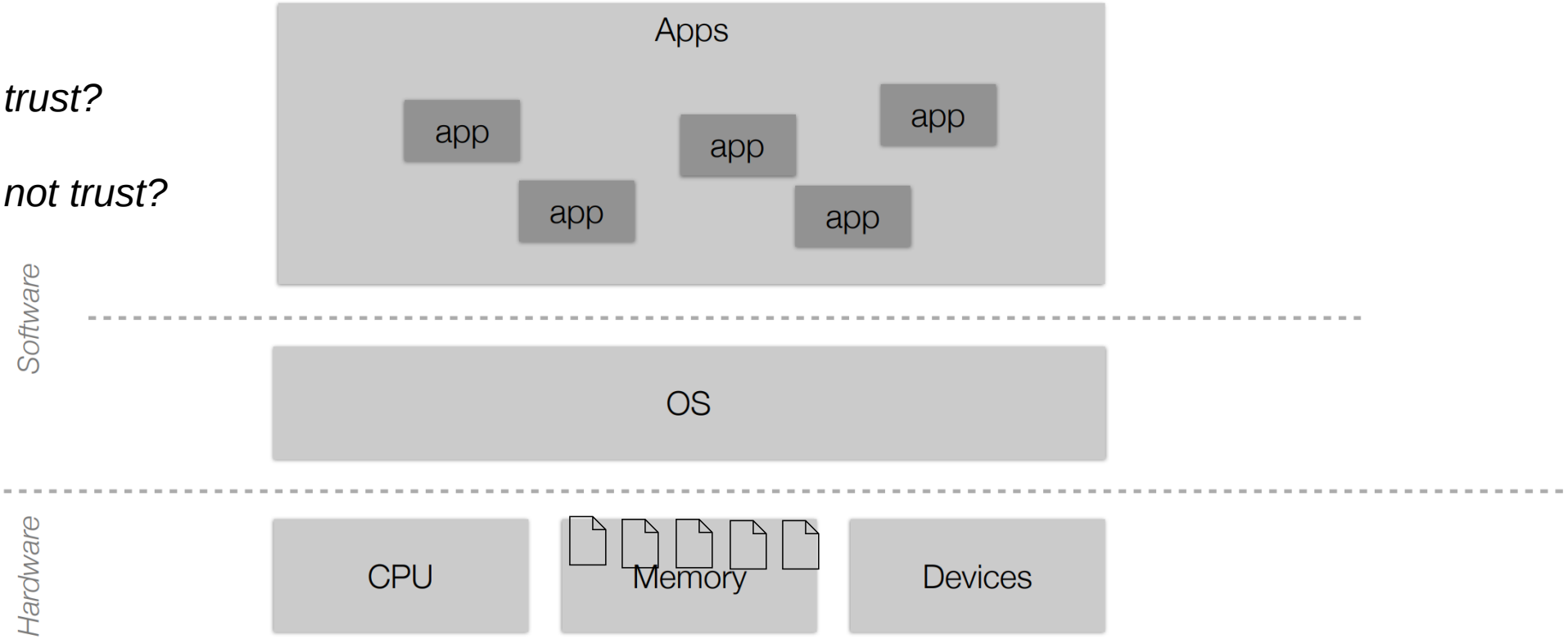


Typical Layers of a Computer

The **operating system** is a vital component of a computer

What do we trust?

What do we not trust?



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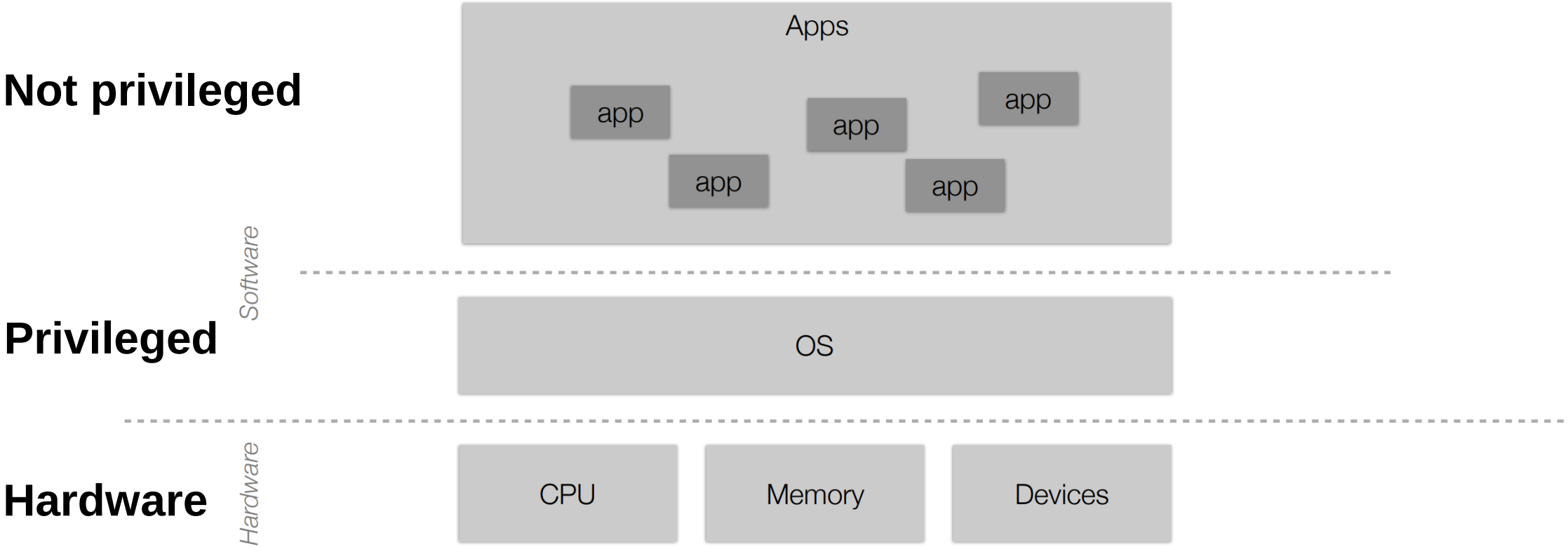

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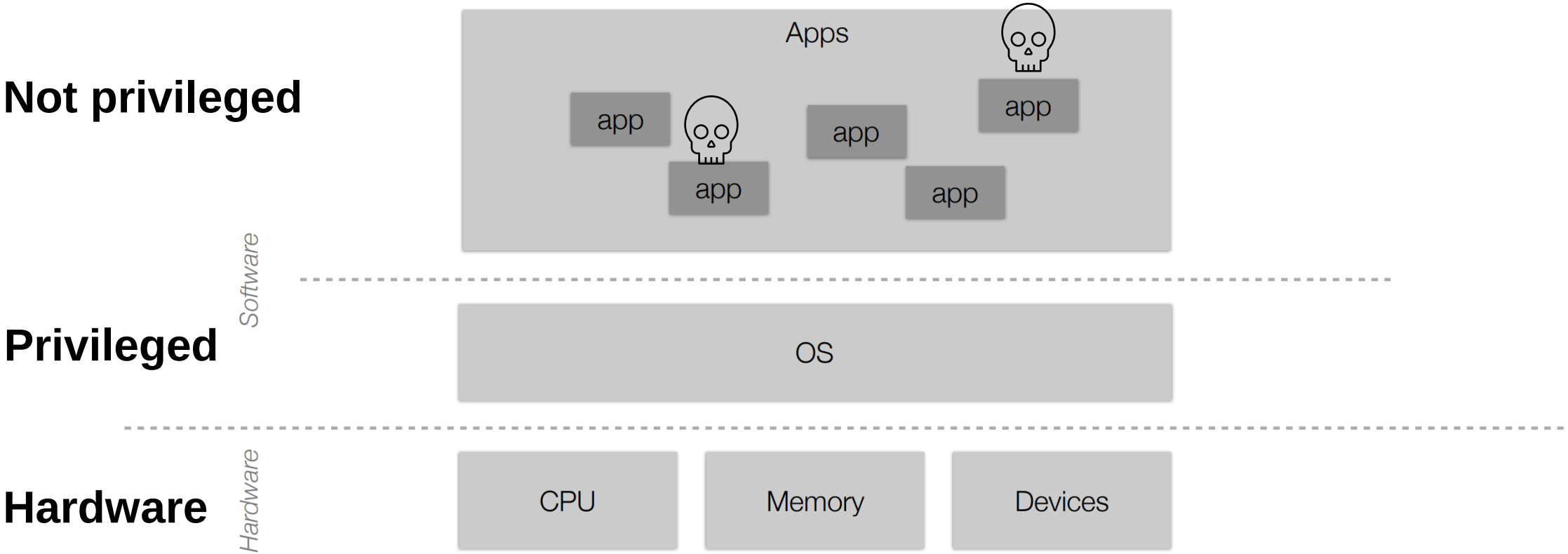

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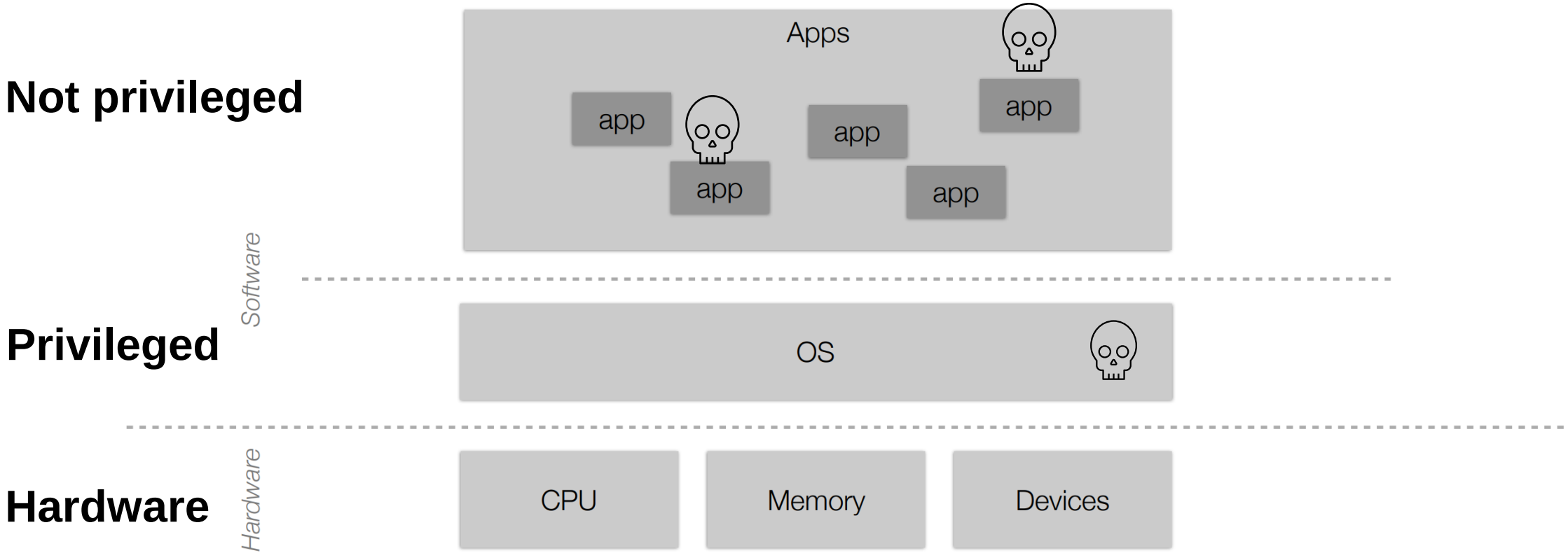
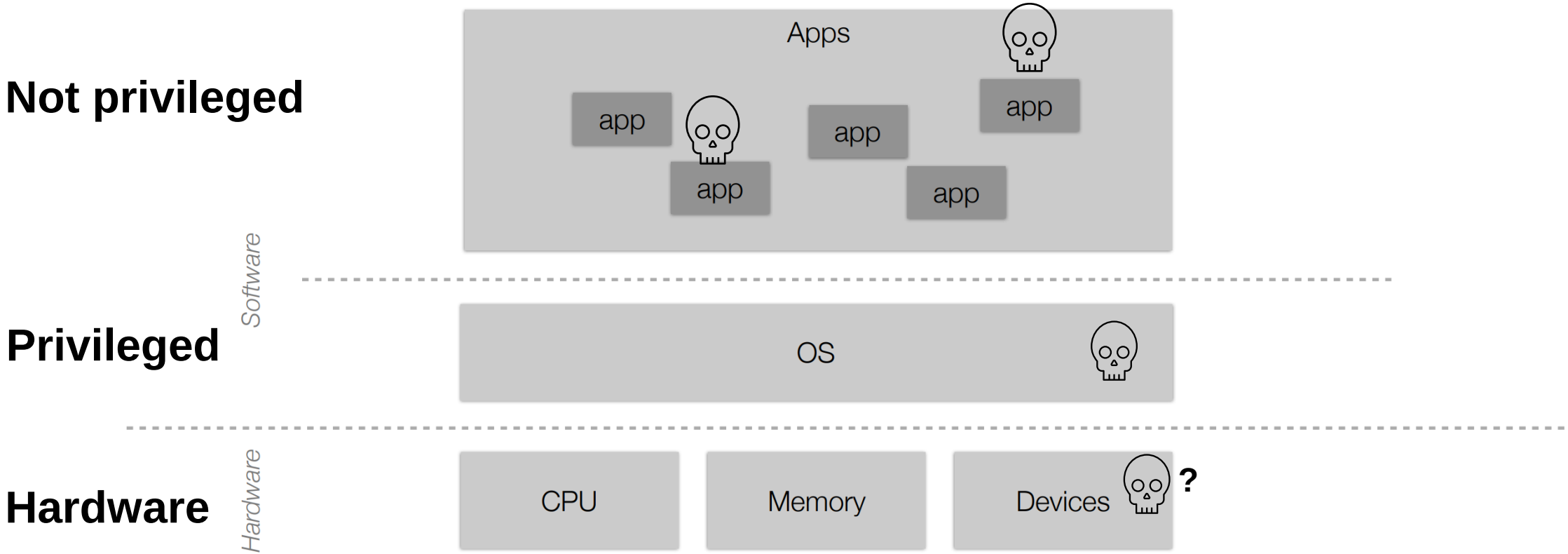
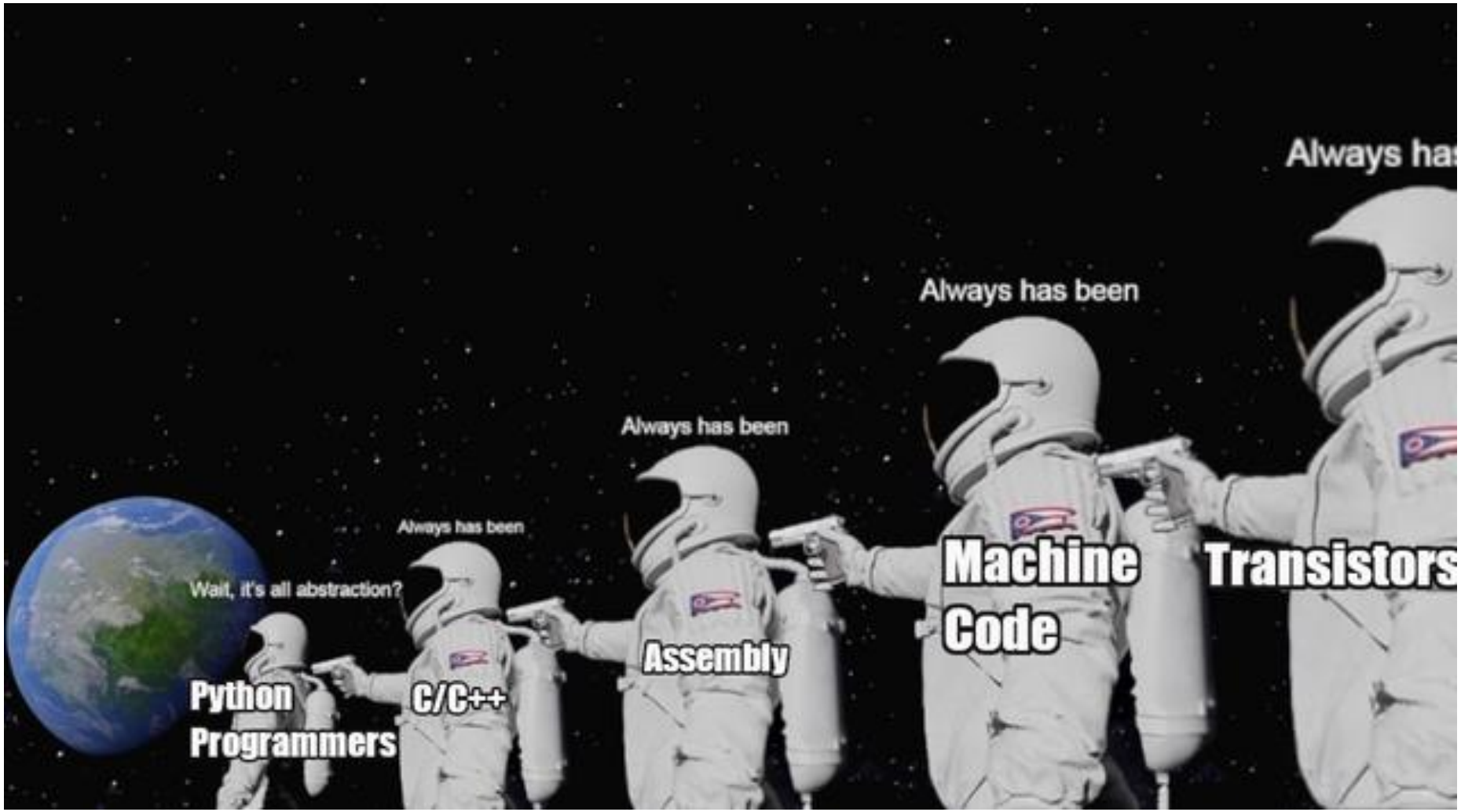


Typical Layers of a Computer

The **operating system** is a vital component of a computer





Meme credit: Carson Gross