CSCI 476: Computer Security

Asymmetric Cryptography (Public Key Cryptography)

Reese Pearsall Fall 2023

Announcements

Hashing Lab due on Sunday (12/10)

Extra credit opportunity (course evaluation)

2% added to your final exam grade

Extra credit opportunity (attend Thursday's lecture)

1 % added to your final exam grade

Check your grades

Final Exam (in person)

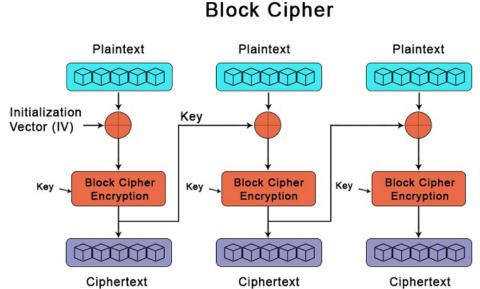
- Thursday December 14th 2:00 3:50 PM
- Romney 315

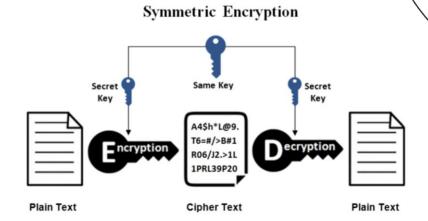


Symmetric key encryption uses the same, **shared**, key for encrypting and decrypting

What is the one major hurdle we have not discussed yet?

How do the keys get sent without being intercepted? Do the keys get encrypted?





(S)

PLAIN TEXT

6

Ν

4

ഗ

G

I

7

Asymmetric Cryptography

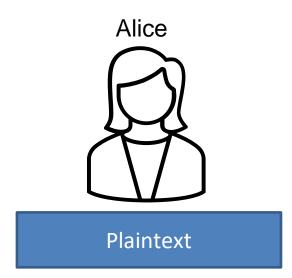
AKA Public key Cryptography

The keys used for encrypting and decrypting data are different

Additionally, each user now gets two-keys. A public key, and a private key

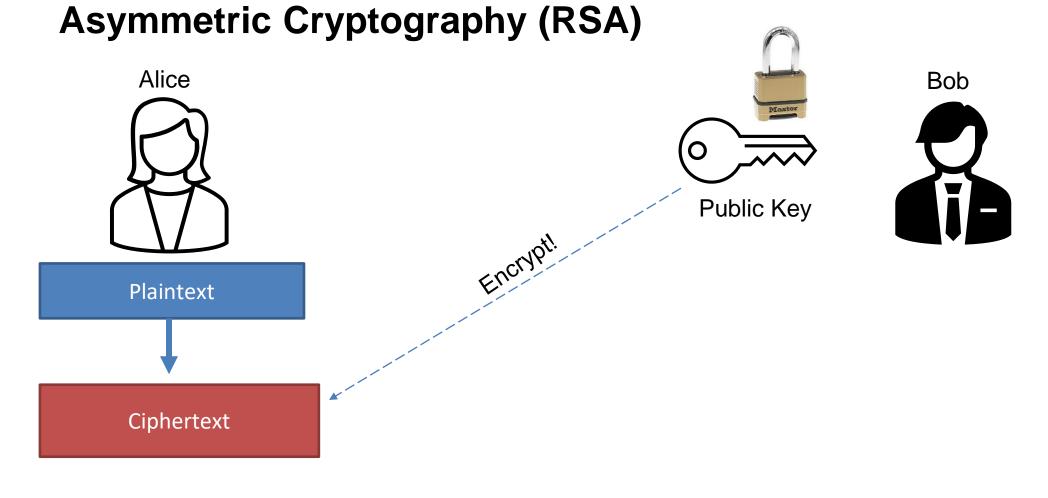
This involves some complicated math, and I won't go super deep into it. YouTube videos can explain it much better than I can

RSA (Rivest–Shamir–Adleman) is the most popular public key cryptosystem. We rely on it whenever we do communicate securely on the internet

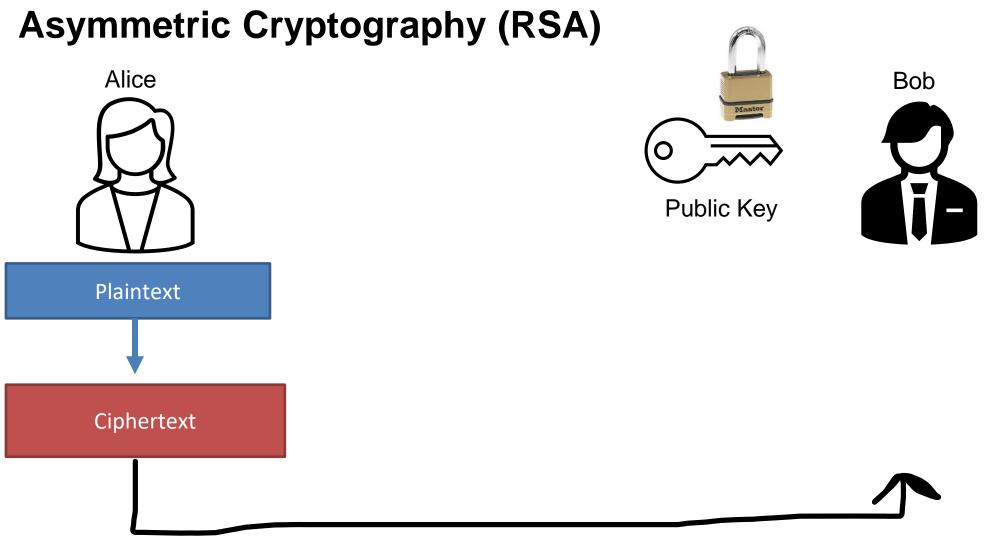




Alice has a plaintext that she wants to send to bob



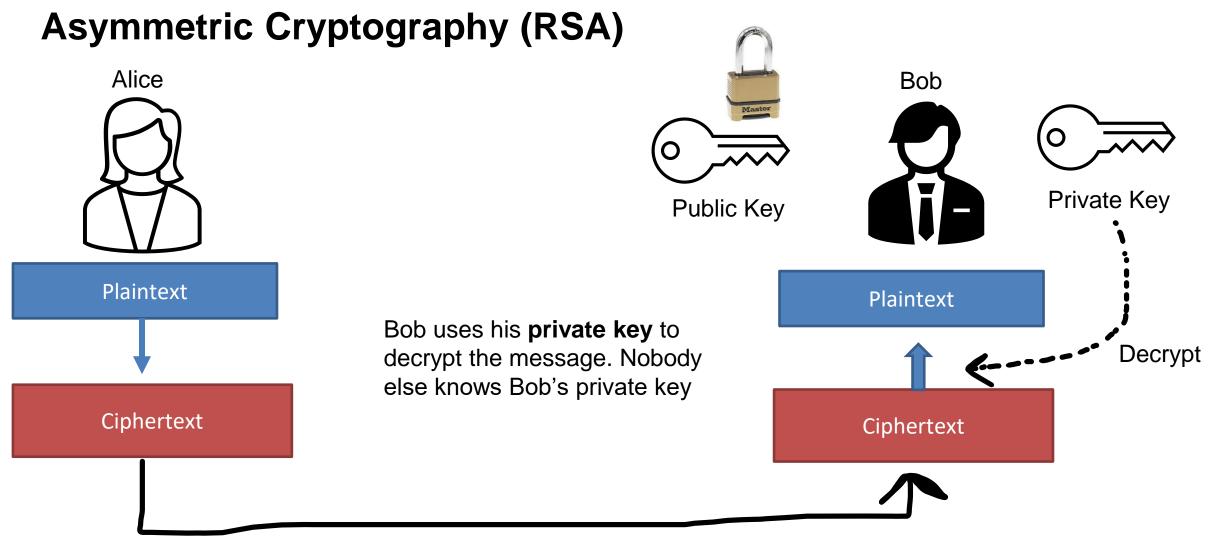
She uses Bob's **public key** to encrypt her message



Ciphertext is sent over some medium



Eve can intercept this message,
But can't decrypt it (public key is not used for decrypting!)



Ciphertext is sent over some medium



Eve can intercept this message,
But can't decrypt it (public key is not used for decrypting!)

If you multiply two prime numbers (**p** and **q**) together, the product can only be divisible by those two number

This is very difficult to figure out for the people that don't know p or q

In fact, there is not an efficient program that can calculate the factors of integers

This problem is in NP

If you multiply two prime numbers (**p** and **q**) together, the product can only be divisible by those two number

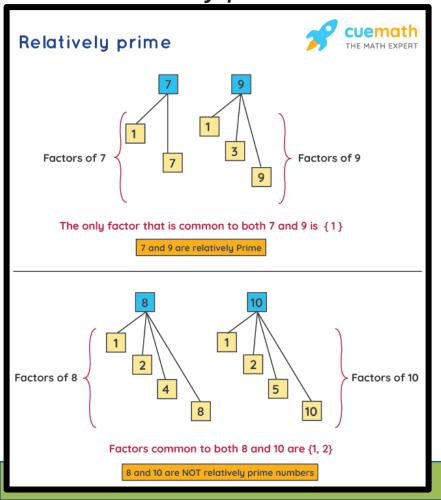
RSA is based on large numbers that are difficult to factorize The public and private keys are derived from these prime numbers

How long should RSA keys be? 1024 or 2048 bits long!

The longer the key = the more difficult to crack (exponentially)

Euler's Totient Function

 $\Phi(n)$ = number of values less than n which are relatively prime to n



Ф(3127)

1 2	How many of these numbers are relatively prime w/ 3127?
3 3125 3126	Difficult But very easy for the product of two prime #s!
	The $\Phi(n)$ of a product of two prime numbers will always be $(p-1)(q-1)$

A number is relatively prime to n if they share no common factors

Eve's stolen goods

Alice





$$p = 53$$

 $q = 59$

Step 1: Choose two large primer numbers, p and q

Eve's stolen goods

Alice







$$p = 53$$

$$q = 59$$

$$n = 3127$$

Step 1: Choose two large primer numbers, p and q

Step 2: Calculate the product n

Eve's stolen goods



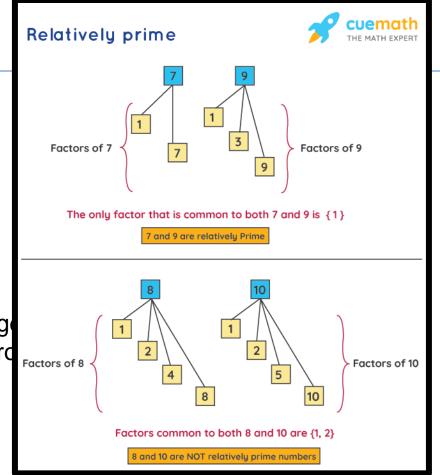
p = 53 q = 59n = 3127

Step 1: Choose two large

Step 2: Calculate the pro

Step 3: Calculate Φ(n)







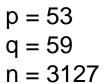
 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

3126

How many of these
numbers are relatively prime
w/ 3127?

Eve's stolen goods

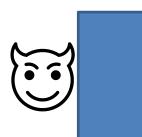


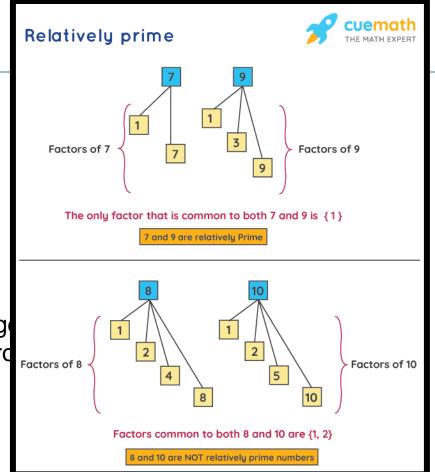


Step 1: Choose two large

Step 2: Calculate the pro

Step 3: Calculate $\Phi(n)$







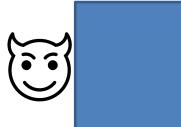
 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

1	
2	How many of these
3	numbers are relatively prime
	w/ 3127?
3125	Difficult But very easy for
3126	the product of two prime
	#s!

Eve's stolen goods

Alice







p = 53

q = 59

n = 3127

which are *relatively prime* to n

 $\Phi(n)$ = number of values less than n

Step 1: Choose two large primer numbers, p and q

Step 2: Calculate the product n

Step 3: Calculate Φ(n)

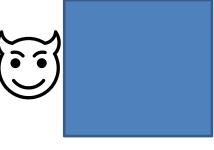
The $\Phi(n)$ of a product of two prime numbers will always be (p-1)(q-1)

Eve's stolen goods

Alice









$$p = 53$$

$$q = 59$$

$$n = 3127$$

$$\Phi(n) = 52*28 = 3016$$

Step 1: Choose two large primer numbers, p and q

Step 2: Calculate the product n

Step 3: Calculate Φ(n)

 $\Phi(n)$ = number of values less than n which are relatively prime to n

The $\Phi(n)$ of a product of two prime numbers will always be (p-1)(q-1)

Eve's stolen goods

 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob







$$p = 53$$

 $q = 59$
 $n = 3127$
 $\Phi(n) = 3016$

$$e = 1 < e < \Phi(n)$$

Not be a factor of n, but an integer

Step 1: Choose two large primer numbers, p and q

Step 2: Calculate the product n

Step 3: Calculate Φ(n)

Step 4: Choose public exponent e

Eve's stolen goods

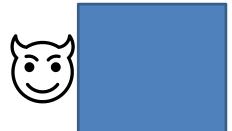
 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob









$$p = 53$$

 $q = 59$
 $n = 3127$
 $\Phi(n) = 3016$

e = 3

$$e = 1 < e < \Phi(n)$$

Not be a factor of n, but an integer

Step 1: Choose two large primer numbers, p and q

Step 2: Calculate the product n

Step 3: Calculate Φ(n)

Step 4: Choose public exponent e

Eve's stolen goods

 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob









$$p = 53$$

 $q = 59$
 $n = 3127$
 $\Phi(n) = 3016$
 $e = 3$

$$J = \frac{K * \Phi(n) + 1}{e}$$

Step 1: Choose two large primer numbers, p and q

Step 2: Calculate the product n

Step 3: Calculate Φ(n)

Step 4: Choose public exponent e

Step 5: Select private exponent d

K = some integer that will make the quotient an integer

Eve's stolen goods

 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob







$$p = 53$$

 $q = 59$
 $n = 3127$
 $\Phi(n) = 3016$
 $e = 3$

$$J = \frac{2*3016+1}{3}$$

Step 1: Choose two large primer numbers, p and q

Step 2: Calculate the product n

Step 3: Calculate $\Phi(n)$

Step 4: Choose public exponent e

Step 5: Select private exponent d

K = some integer that will make the quotient an integer

Eve's stolen goods

 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob







$$p = 53$$

 $q = 59$
 $n = 3127$
 $\Phi(n) = 3016$
 $e = 3$
 $d = 2011$

- Step 1: Choose two large primer numbers, p and q
- Step 2: Calculate the product n
- Step 3: Calculate $\Phi(n)$
- Step 4: Choose public exponent e
- Step 5: Select private exponent d

K = some integer that will make the quotient an integer

Eve's stolen goods

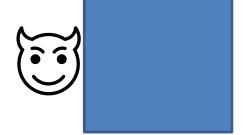
 $\Phi(n)$ = number of values less than n which are *relatively prime* to n











Alice's Public Key

Eve's stolen goods

 $\Phi(n)$ = number of values less than n which are *relatively prime* to n











Bob has a message to send to Alice



Message must be converted into a number

Eve's stolen goods

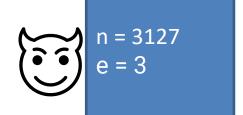
 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob









Alice's Public Key

Bob has a message to send to Alice



Use Alice's Public Key to encrypt

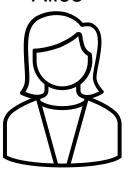
Eve's stolen goods

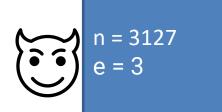
 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob









Alice's Public Key

Bob has a message to send to Alice



Use Alice's Public Key to encrypt

$$89^{3}$$
 mod 3127

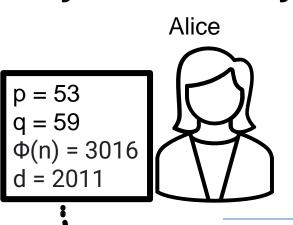
Eve's stolen goods

n = 3127

 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob





Alice's Public Key

Bob has a message to send to Alice



Use Alice's Public Key to encrypt

$$89^{3}$$
 mod 3127

1394 2011 mod 3127

Alice decrypts message using her private key

Eve's stolen goods

 $\Phi(n)$ = number of values less than n which are *relatively prime* to n

Bob





Alice's Public Key

d = 2011

Bob has a message to send to Alice

89

Use Alice's Public Key to encrypt

 89^{3} mod 3127 C = 1394



Eve's stolen goods

 $\Phi(n)$ = number of values less than n which are relatively prime to n







p = 53q = 59 $\Phi(n) = 3016$ d = 2011



Alice's Public Key

Bob has a message to send to Alice

Alice's Private Key

$$n = 3127$$

 $d = 2011$

What does eve know??



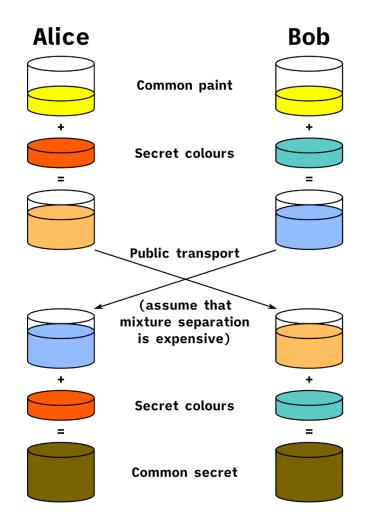
Use Alice's Public Key to encrypt



We now have a method for sending secure messages over a possibly unsecure channel!

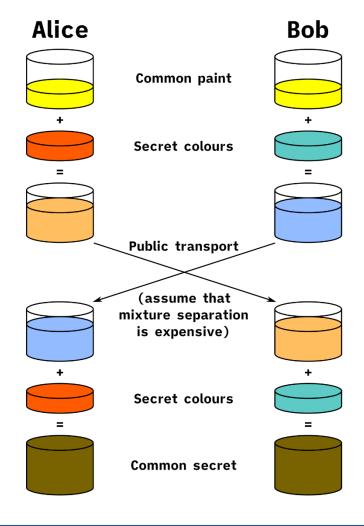
This method is known as the **Diffie Helman Key Exchange**

RSA is built on DHKE to create an encryption/decryption algorithm



We now have a method for sending secure messages over a possibly unsecure channel!

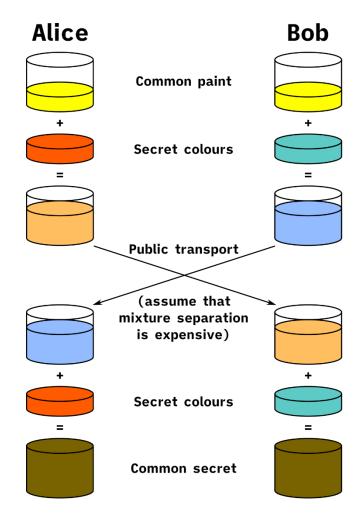
Limitation of RSA: Can only encrypted data that is smaller or equal to key length (< 2048 bits)



We now have a method for sending secure messages over a possibly unsecure channel!

Limitation of RSA: Can only encrypted data that is smaller or equal to key length (< 2048 bits)

What could we encrypt instead??

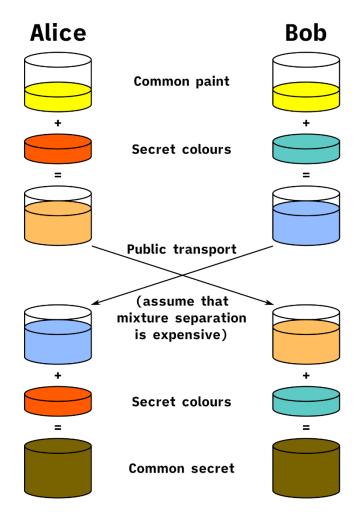


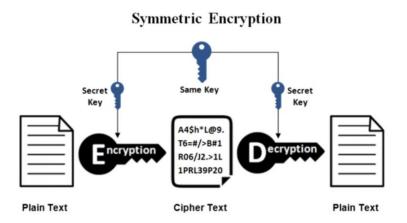
We now have a method for sending secure messages over a possibly unsecure channel!

Limitation of RSA: Can only encrypted data that is smaller or equal to key length (< 2048 bits)

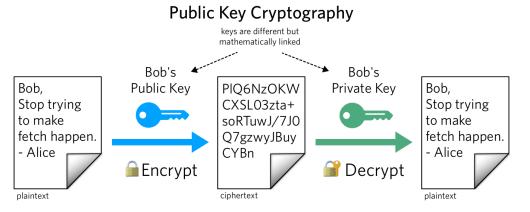
What could we encrypt instead??

The key for a symmetric cryptography algorithm! (< 2048 bits)





- Same key used for encrypting and decrypting
- Using block ciphers (AES), we can encrypt an arbitrary size of data
- Issue: How to securely share secret keys with each other?



- Two keys: Public Key (a lock), and a price key (the key)
- Public key is used to encrypt. Private key used to decrypt message
- Using math, we can securely send messages over an unsecure channel without sharing any sensitive information
- Issue: We can not encrypt stuff bigger than our key (2048 bits)
- Often times, symmetric and asymmetric cryptography are used together

(use RSA to send the key for symmetric crypto!)

We know that Public and Private keys are derived from big prime numbers (We are talking hundreds of digits long...)

Our computer can't compute products and exponents for such large numbers

OpenSSL on our VMs has tools for generating public/private RSA keys

[11/29/22]seed@VM:~\$ sudo openssl genrsa -aes128 -out private.pem 1024

Example: generate a 1024-bit public/private key pair

- · Use openssl genrsa to generate a file, private.pem
- private.pem is a Base64 encoding of DER generated binary output

```
$ openssl genrsa -aes128 -out private.pem 1024 # passphrase csci476
$ more private.pem
----BEGIN RSA PRIVATE KEY----
Proc-Type: 4, ENCRYPTED
DEK-Info: AES-128-CBC, C30BF6EB3FD6BA9A81CCB9202B95EC1A
sLIQ7Fs5j5z0exdWkZUoiv2W82q03qNERmfG+fwnVnbsIZAuW8E9wiB7tqz8rEL+
xfL+U20lyQNxpmOTUeKlN3qCcJROcGYSNd1BeNpqLWV1bN5FPYce9GRb4tFr4bhK
RPtJNKUryhVnAC4a3gp0gcXk1IQLeHeyKQCPQ1SckQRdrBzHjjCNN42N1CVEpcsF
WJ8ikqDd9FslGHc1PT6ktW5oV9cB8G2wfo7D85n91SQfSzuwAcyx7Ecir1o4PfKG
----END RSA PRIVATE KEY-
```

The *actual* content of **private.pem**:

```
$ openssl rsa -in private.pem -noout -text
Enter pass phrase for private.pem: csci476
Private-Key: (1024 bit)
modulus:
    00:b8:52:5c:25:cc:7c:f2:ef:a6:35:9d:de:3d:5d: ...
publicExponent: 65537 (0x10001)
privateExponent:
    4b:0d:ce:53:dd:e6:6b:0d:c6:82:42:9c:42:24:a7: ...
prime1:
   00:ef:14:46:57:9c:d0:4c:98:de:c3:0b:aa:d8:72: ...
prime2:
    00:c5:5d:f8:0b:f9:75:dc:88:ea:d4:d0:56:ee:f9: ...
exponent1:
    00:e6:49:9a:44:14:19:94:5e:7f:dc:52:65:bb:5d: ...
exponent2:
   7c:ad:77:dc:58:a2:13:c6:8a:52:15:aa:55:1c:22: ...
coefficient:
    3a:7c:b9:a0:12:e8:fa:88:b8:6f:38:4a:ed:bc:17: ...
```

The actual content of public.pem:

```
$ openssl rsa -in private.pem -pubout > public.pem
Enter pass phrase for private.pem: csci476
writing RSA key
$ more public.pem
----BEGIN PUBLIC KEY----
MIGfMA0GCSqGSIb3DQEBAQUAA4GNADCBiQKBgQC4UlwlzHzy76Y1nd49XakNUwqJ
Ud3ph0uBWWfnLnjIYgQL/spg9WE+1Q1YPp2t3FBFljhGHdWMA8abfNXG4jmpD+uq
Ix0WVyXg12WWi1kY2/vs8xI1K+PumWTtq8R8ueAq7RzETc3873DO1vjMxXWqau7k
zIkUuJ/JCjzjYfbsDQIDAQAB
----END PUBLIC KEY-----
```

```
$ openssl rsa -in public.pem -pubin -text -noout
Public-Key: (1024 bit)
Modulus:
    00:b8:52:5c:25:cc:7c:f2:ef:a6:35:9d:de:3d:5d: ...

Exponent: 65537 (0x10001)

(e,n) = public Key!
```

OpenSSL Tools: Encryption and Decryption

· Create a plaintext message:

```
$ echo "This is a secret." > msg.txt
```

Encrypt the plaintext:

```
$ openssl rsautl -encrypt -inkey public.pem -pubin -in msg.txt -out msg.enc
```

OpenSSL Tools: Encryption and Decryption

· Create a plaintext message:

```
$ echo "This is a secret." > msg.txt
```

Encrypt the plaintext:

```
$ openssl rsautl -encrypt -inkey public.pem -pubin -in msg.txt -out msg.enc
```

Decrypt the ciphertext:

```
$ openssl rsautl -decrypt -inkey private.pem -in msg.enc
Enter pass phrase for private.pem: csci476
This is a secret.
```

OpenSSL Tools: Encryption and Decryption

· Create a plaintext message:

```
$ echo "This is a secret." > msg.txt
```

Encrypt the plaintext:

```
$ openssl rsautl -encrypt -inkey public.pem -pubin -in msg.txt -out msg.enc
```

Decrypt the ciphertext:

```
$ openssl rsautl -decrypt -inkey private.pem -in msg.enc
Enter pass phrase for private.pem: csci476
This is a secret.
```

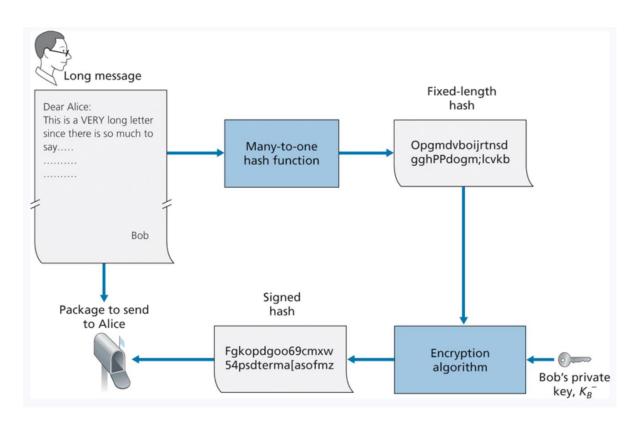
BIG NUM API

```
int main ()
 BN CTX *ctx = BN CTX new();
 BIGNUM *p, *q, *n, *phi, *e, *d, *m, *c, *res;
 BIGNUM *new_m, *p_minus_one, *q_minus_one;
  p = BN_new(); q = BN_new(); n = BN_new(); e = BN_new();
  d = BN new(); m = BN new(); c = BN new();
 res = BN new(); phi = BN new(); new m = BN new();
  p minus one = BN new(); q minus one = BN new();
  // Set the public key exponent e
 BN dec2bn(&e, "65537");
  // Generate random p and q.
  BN generate prime ex(p, NBITS, 1, NULL, NULL, NULL);
 BN_generate_prime_ex(q, NBITS, 1, NULL, NULL, NULL);
  BN sub(p minus one, p, BN value one()); // Compute p-1
```

Digital Signatures

- What is a unique identifier for bob? What is something that only bob knows and nobody else?
- > His private key

Bob encrypts his hashed message using his **private key**, and sends the signed hash, along with message to Alice



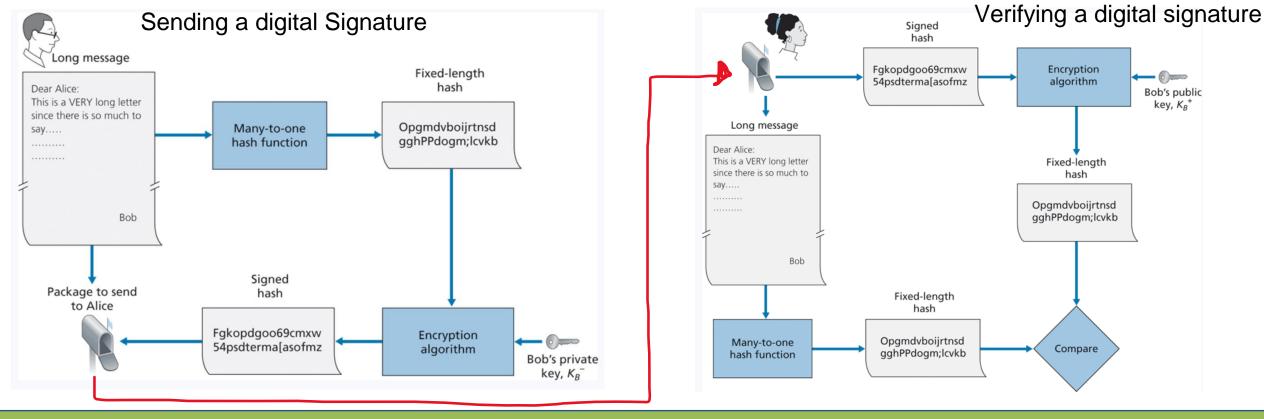
When Alice receives this message, she must find a way to decrypt the signed hash

She will use Bob's public key

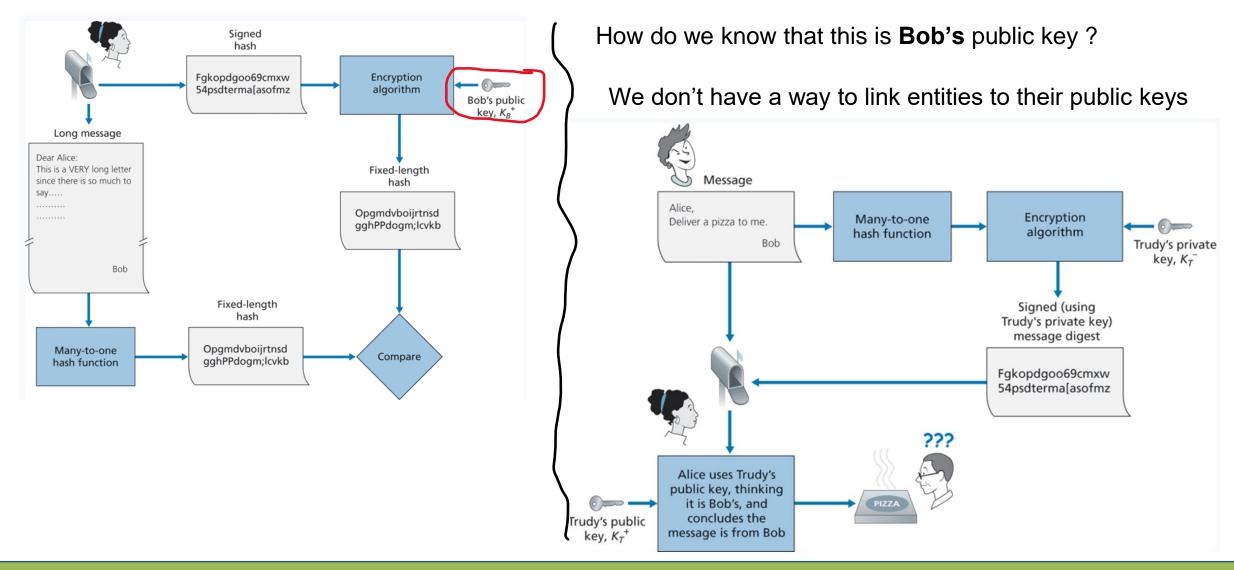
Digital Signatures

- What is a unique identifier for bob? What is something that only bob knows and nobody else?
- > His private key

Bob encrypts his hashed message using his **private key**, and sends the signed hash, along with message to Alice. Alice decrypts using his **public key** and verifies that the hashes match



Digital Signatures

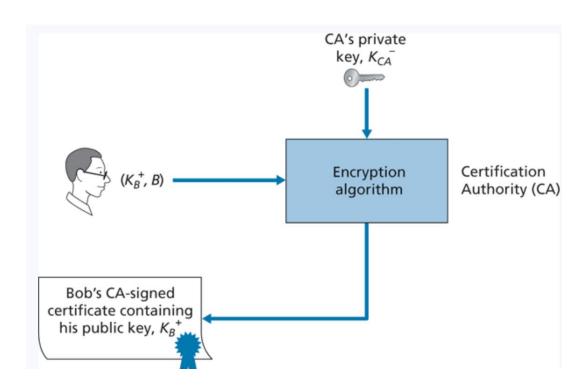


Digital Certificates

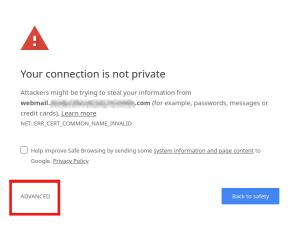
Certificates are an authoritative document that links entities (person, router, organization) to their public key

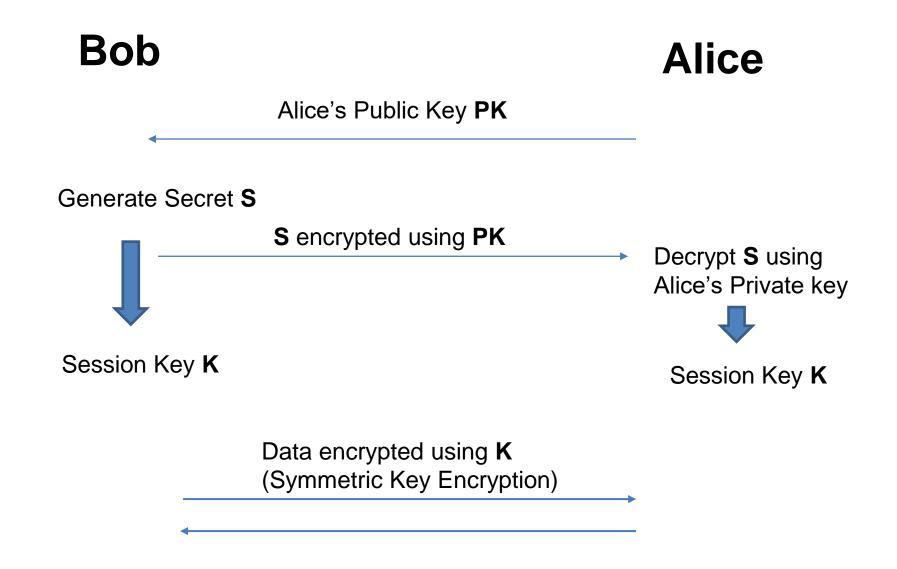
Creating certificates are done by a **Certification Authority** (digicert, lets encrypt, comodo)

Some are more trustworthy than others...



On your web browser, you exchange certificate information with the websites you are visiting

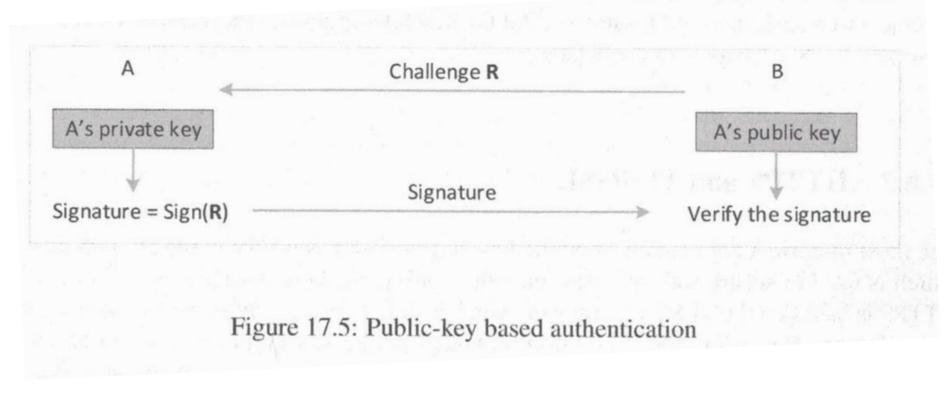




Application: HTTPS and TLS/SSL

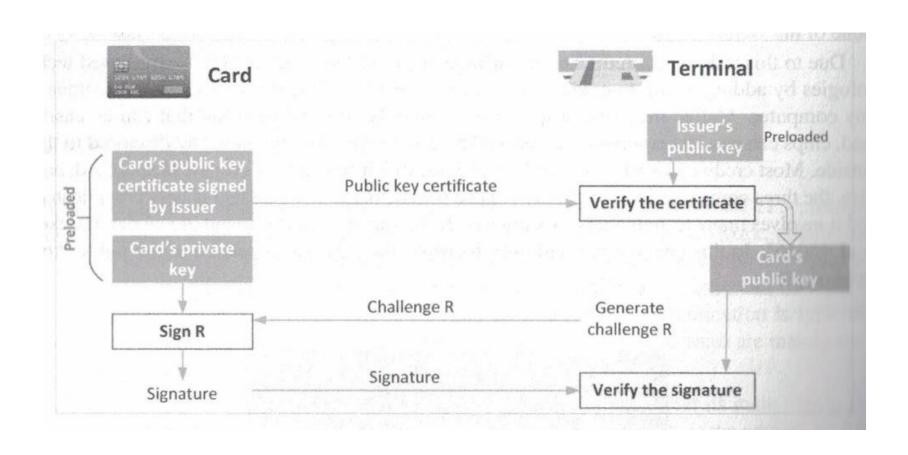
Public Key Authentication: Credit Cards



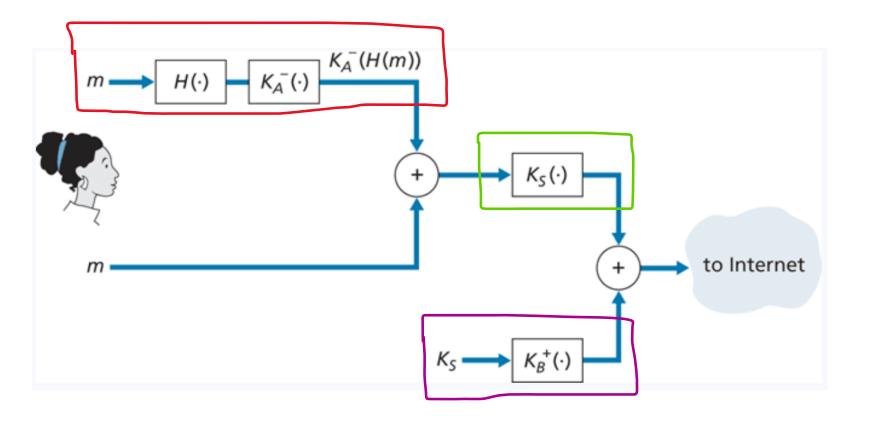


Public Key Authentication: Credit Cards





Symmetric Crypto Asymmetric Crypto and Hashing all work together to send secure, authentic messages



Legitimate organizations must meet **compliance** standards if they want to do business. This includes things such as handling transactions securely, encrypting user data, no plaintext passwords, etc

These rules are structured as a compliance framework, which is a structures set of guidelines and best practices that details a company's processes for meeting regulatory requirements



STIG - 230503

The Red Hat Enterprise Linux operating system must be configured to disable USB mass storage.

STIG - 230534

The Red Hat Enterprise Linux operating system must be configured so that the root account must be the <u>only</u> account having unrestricted access to the system.

STIG - 217976

The audit system must be configured to audit all use of **setuid** and setgid programs.

STIG - 217976

The Ubuntu operating system must implement **address space layout randomization** to protect its memory from unauthorized code execution.

STIG - 230231

RHEL 8 must encrypt all stored passwords with a FIPS 140-2 approved **cryptographic hashing** algorithm.

SHA-256, SHA-512, etc