CSCI 466: Networks

Lecture 5: Peer 2 Peer Networks (P2P), Content Distribution Networks (CDN)

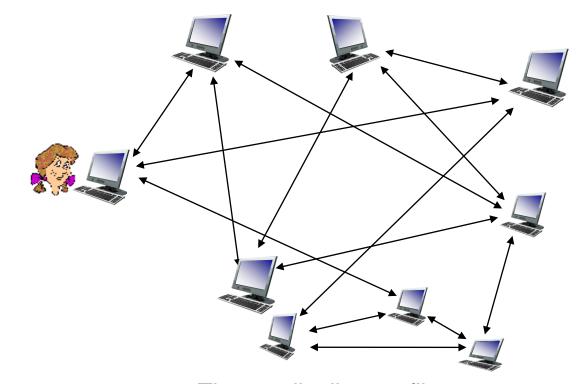
Reese Pearsall Fall 2022

Announcements

there are no announcements

P2P Networks

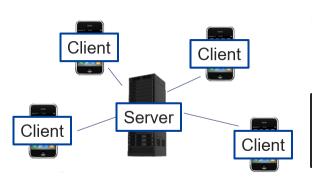
- No always-on server
- Arbitrary end systems directly communicate
- Peers are intermittently connected and change IP addresses



Time to distribute a file of size F to N clients

$$D_{P2P} = \max \left\{ \frac{F}{u_s}, \frac{F}{d_{\min}}, \frac{NF}{u_s + \sum_{i=1}^{N} u_i} \right\}$$

Client-Server Architecture



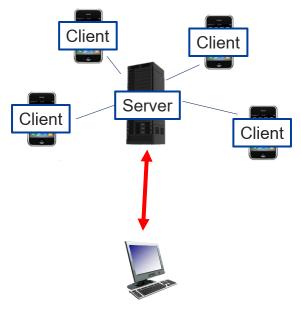
- Server can upload data at rate u_s
- Clients download data at rates d₁, d₂, ..., d_N

$$D_{CS} = \max\left\{\frac{NF}{u_s}, \frac{F}{d_{\min}}\right\}$$

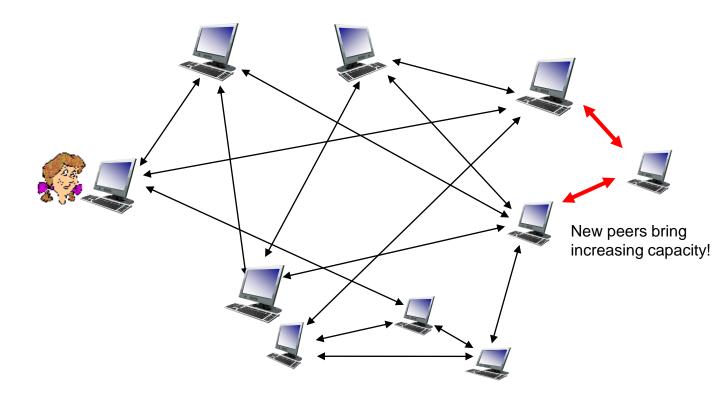
Time to distribute a clients

P2P Networks

Client-Server Architecture



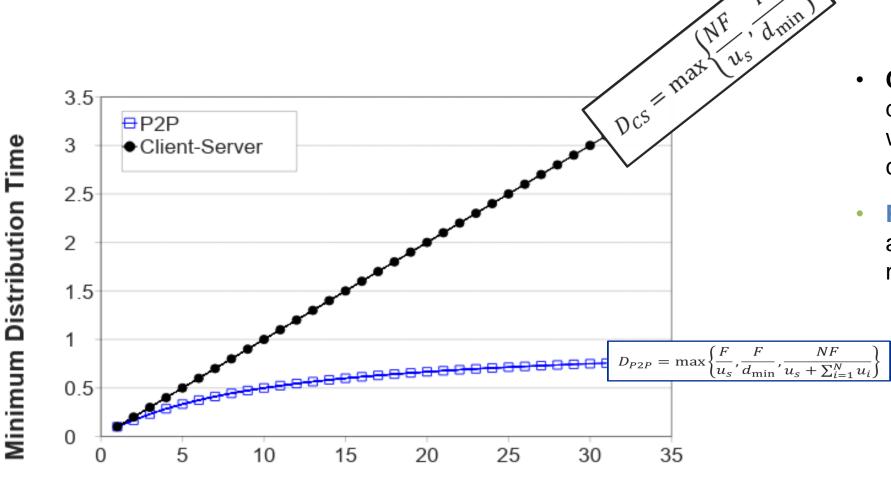
Existing clients have to share resources with new users



New peers are both a client and a server. There will not be a negative impact on current peers

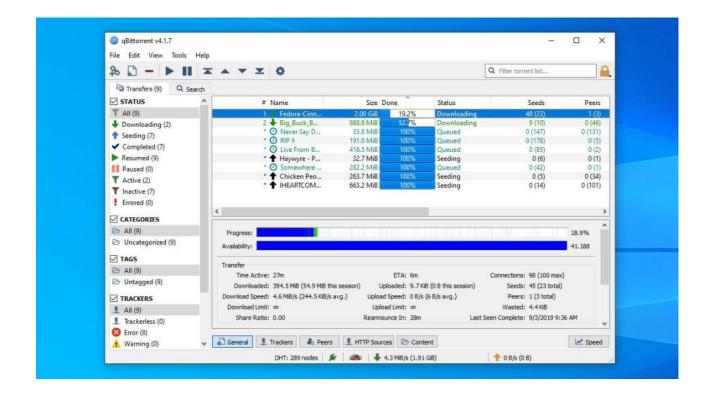
P2P architectures are self-scaling





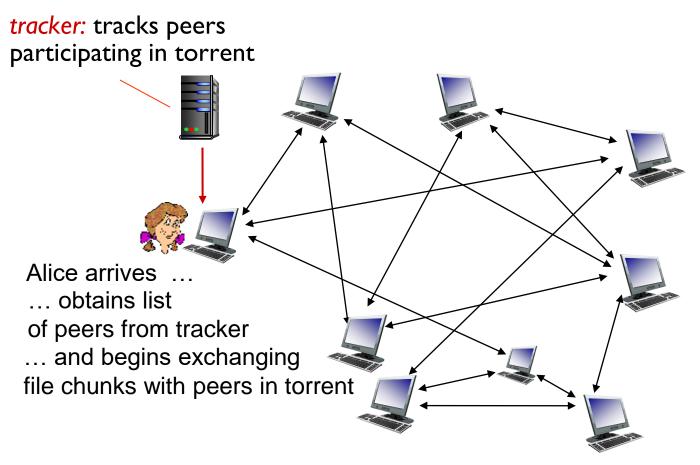
Ν

- Client server
 distribution time grows
 with the number of
 clients
- P2P distribution time approaches 1 hour as number of clients grows



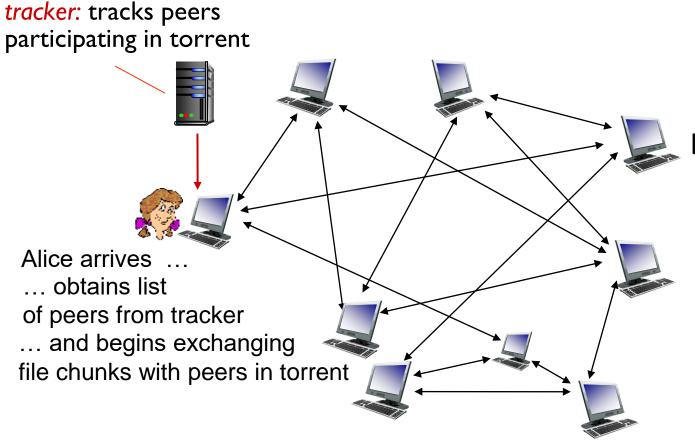
Service for sharing files over the internet in a decentralized fashion

- Files are divided into chunks
- Peers in torrent send/receive file chunks



torrent: group of peers exchanging chunks of a file

- Files are divided into chunks
- Peers in torrent send/receive file chunks

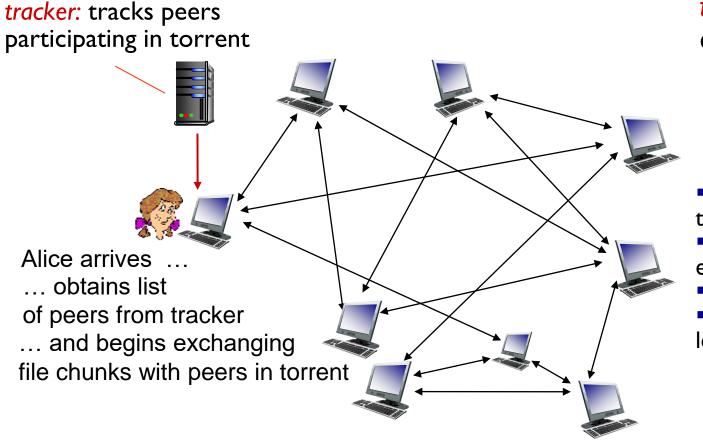


torrent: group of peers exchanging chunks of a file

peer joining torrent:

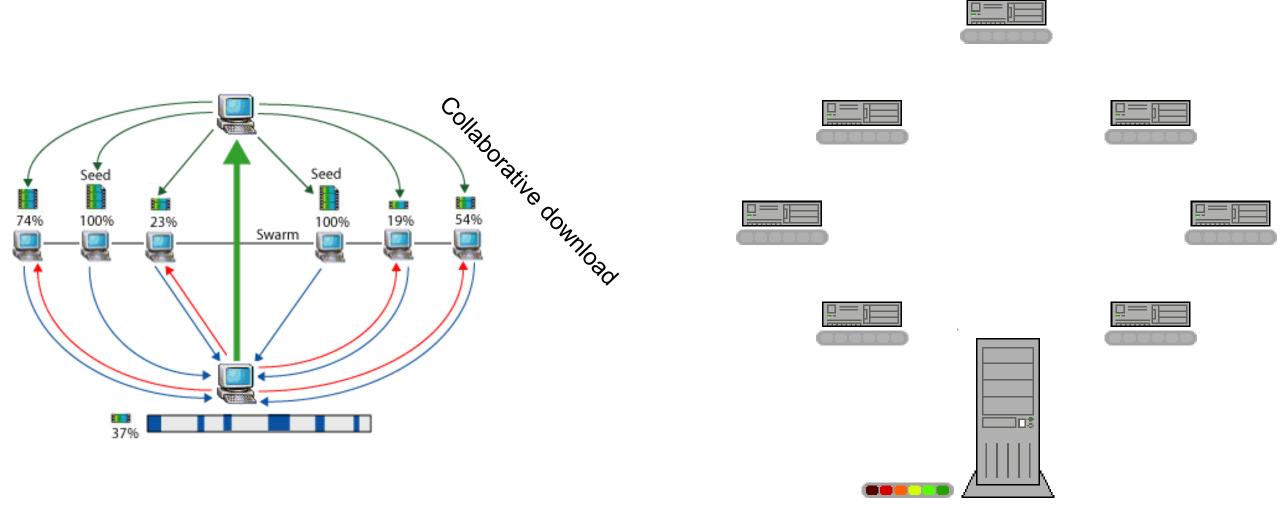
- has no chunks, but will accumulate them over time from other peers
- registers with tracker to get list of peers, connects to subset of peers ("neighbors")

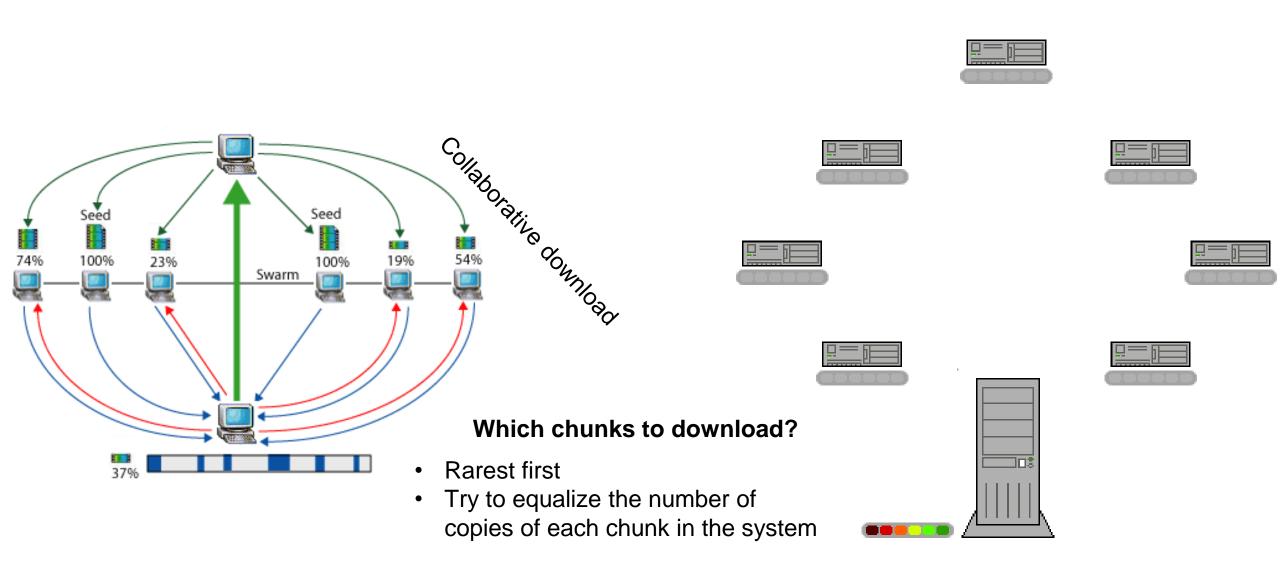
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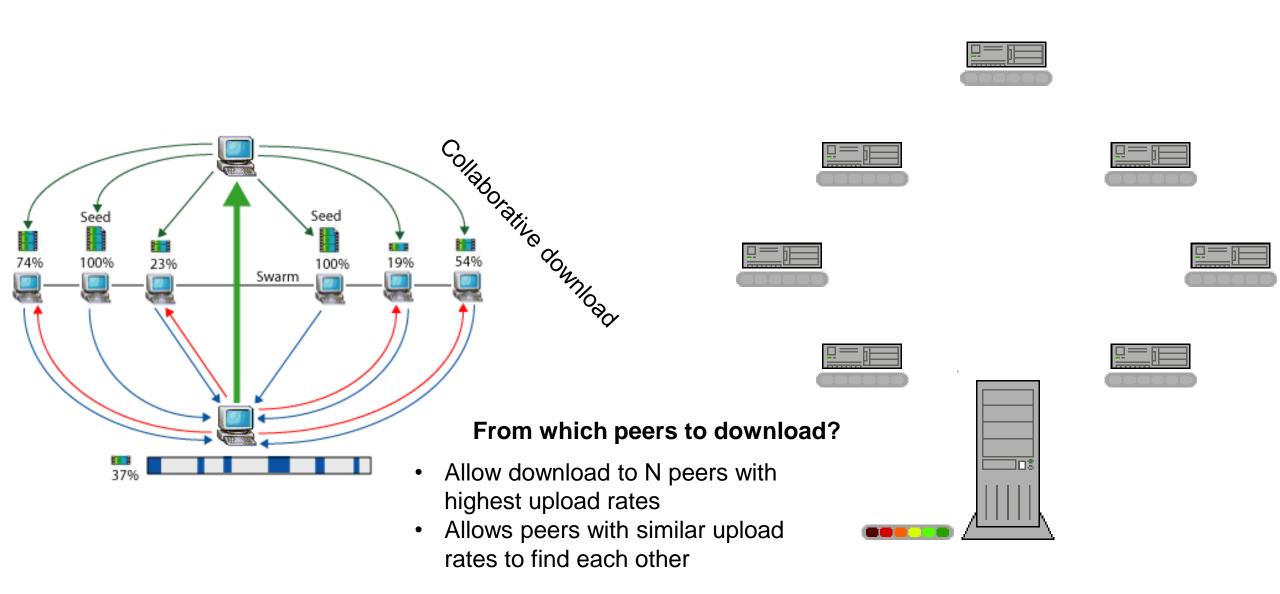


torrent: group of peers exchanging chunks of a file

- while downloading, peer uploads chunks to other peers
- •peer may change peers with whom it exchanges chunks
- **churn**: peers may come and go
- •once peer has entire file, it may (selfishly) leave or (altruistically) remain in torrent

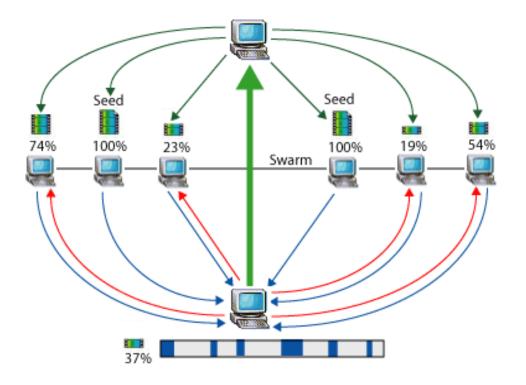


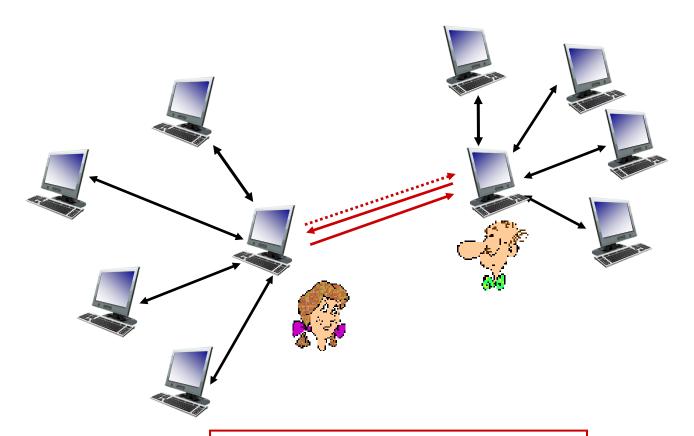




requesting chunks:

- at any given time, different peers have different subsets of file chunks
- periodically, Alice asks each peer for list of chunks that they have
- Alice requests missing chunks from peers, rarest first





higher upload rate: find better trading partners, get file faster!

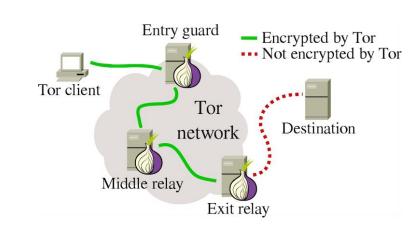
sending chunks: tit-for-tat

- Alice sends chunks to those four peers currently sending her chunks at highest rate
 - other peers are *choked* by Alice (do not receive chunks from her)
 - re-evaluate top 4 every 10 secs
- every 30 secs: randomly select another peer, starts sending chunks
 - "optimistically unchoke" this peer
 - newly chosen peer may join top 4

BitTorrent is referred to as a unstructured P2P

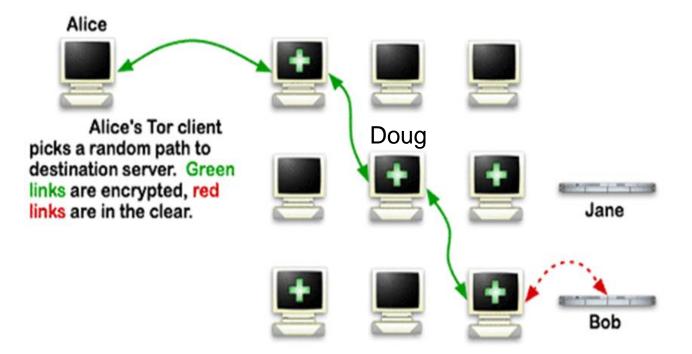
P2P file distribution mechanisms

To help with anonymity and privacy, some P2P may implement an onion router/TOR router



Functionality

- Sender obtains a set of router keys
- Each router only knows next hop
- Intermediate routers cannot read message



Structured P2P Example

Content is placed into the P2P network based on a Globally Unique Identifier (GUID)

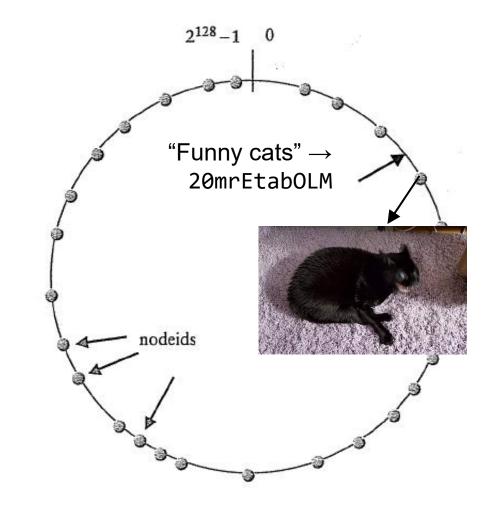
• GUID = H(key) where H is some hash function

$$P(H(a) = H(b))$$
 is very low

- Distributed <key, value> search
- > Fully Distributed and self-organzing

Distributed Hash Table (DHT)

- Join: join ring with GUID from public key
- Publish: put(GUID, value)
- Search/Fetch: value = get(GUID)



Each node maintains a link to the 4 **closest** neighbors and to nodes **half** way across the ring, **quarter** way across the ring, **eight** across the way... O(logN)

Structured P2P Example

	Unstructured P2P	Structured P2P
Advantages	Self-organizingNaturally resilient to node failure	 Guaranteed to locate objects (if exist) Time and complexity bounds Low message overhead
Disadvantages	 Cannot offer guarantees on locating objects Prone to excessive messaging that limits scalability 	 Need to maintain complex overlays High control traffic overhead in dynamic environments

- video traffic: major consumer of Internet bandwidth
 - Netflix, YouTube: 37%, 16% of downstream residential ISP traffic
 - ~1B YouTube users, ~75M Netflix users
- challenge: scale how to reach ~1B users?
 - single mega-video server won't work (why?)
- challenge: heterogeneity
 - different users have different capabilities (e.g., wired versus mobile; bandwidth rich versus bandwidth poor)
- solution: distributed, application-level infrastructure

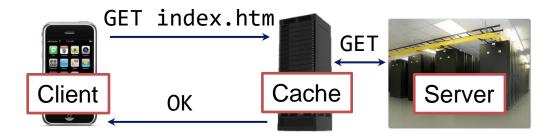






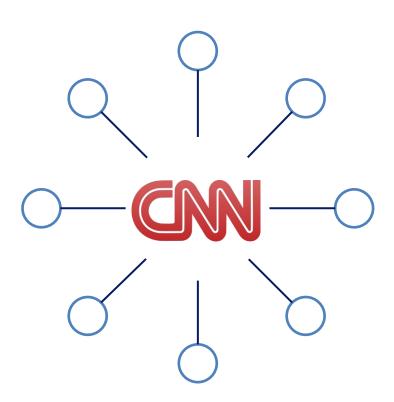


- Caching
 - Save previously delivered data
 - Subsequent requests served from cache on the browser, or in the access network
- Applications
 - Reduce response time for client request
 - Reduce ISP traffic costs
- Content distribution networks
 - Distributed caches
 - Web objects addressed to CDN server
 - CDN server fetches from content provider on first access



- Conditional GET
 - Cache: specify date of cached copy in HTTP request
 - If-modified-since: <date>
 - Server: response contains no object if cached copy is up-to-date: HTTP/1.0 304 Not Modified

- Challenge:
 - How to stream content to millions of users?
- Option 1:
 - Single mega-datacenter
 - Pros: Simple
 - Cons:
 - Single point of failure
 - Point of network congestion
 - Long path to distant clients
 - Multiple copies of video sent over outgoing link



- Option 2:
 - Store/serve multiple copies of videos at multiple geographically distributed sites

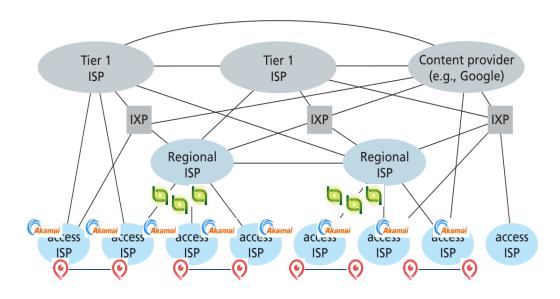
A **content delivery network (CDN)** refers to a geographically distributed group of cache servers which work together to provide fast delivery of Internet content. (not a web host)



34 DNS lookups

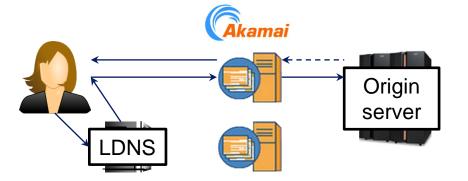
204 HTTP requests

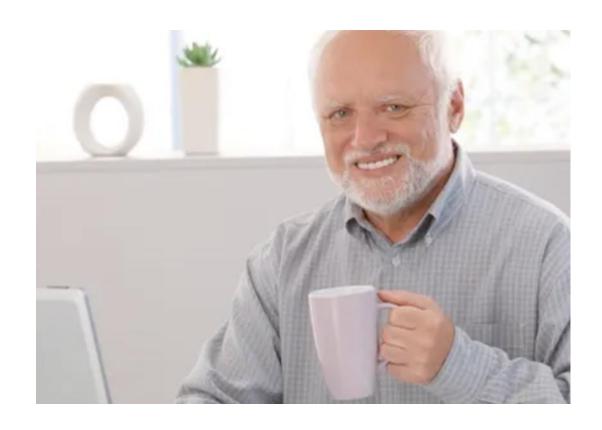
520 KB of data downloaded





56% of domains resolve to a CDN

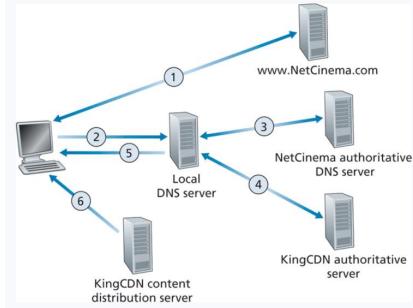




https://**Cdn**.discordapp.com/attachments/101763035 8621143110/1017849417216294972/unknown.png

- Challenge: how does CDN DNS select "good" CDN node to stream to client
 - Pick CDN node geographically closest to client's local DNS

 Pick CDN node with shortest delay (or min # hops) to client (CDN nodes periodically ping access ISPs, reporting results to CDN DNS)



In HTTP steaming, video is stored at an HTTP server and retrieved with a GET request

Frames are sent to a client buffered and played back after a certain threshold

All clients receiving the same encoding of the video despite widely different bandwidth

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All clients receiving the same encoding of the video despite widely different bandwidth

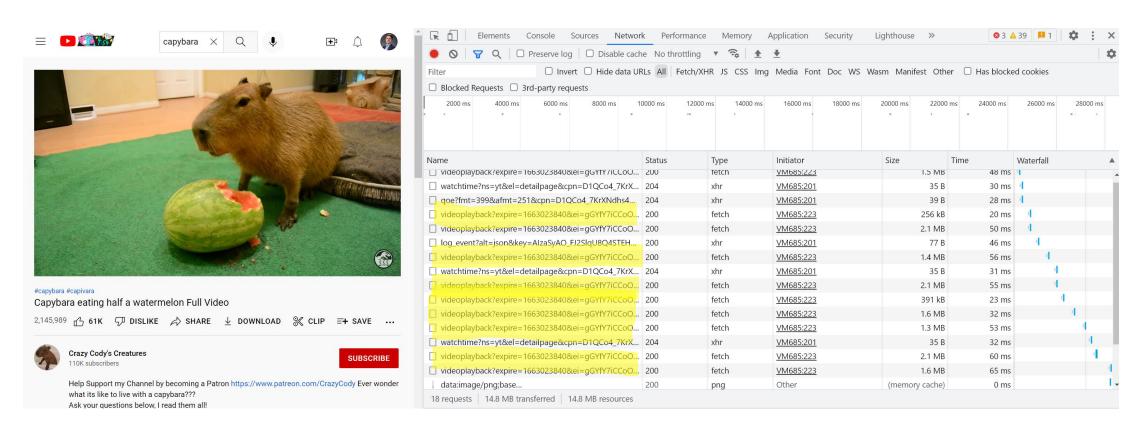
Dynamic Adaptive Streaming over HTTP (DASH)- video is encoded into several different versions

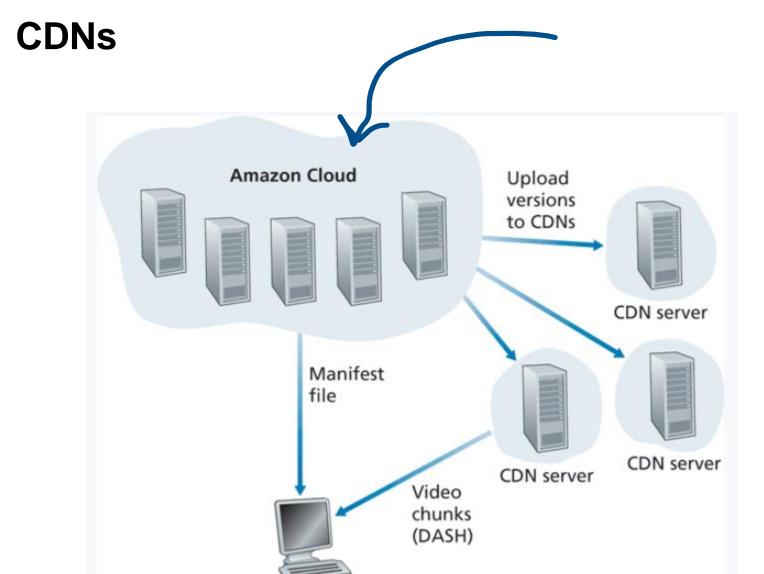
Different encoding rate, different quality, etc

Clients dynamically request chunks of video segments every few seconds via GET requests

- Is the current bandwidth good? → Retrieve good quality
- Is the current bandwidth bad? → Retrieve ok quality

When streaming, you are consistently issuing GET requests

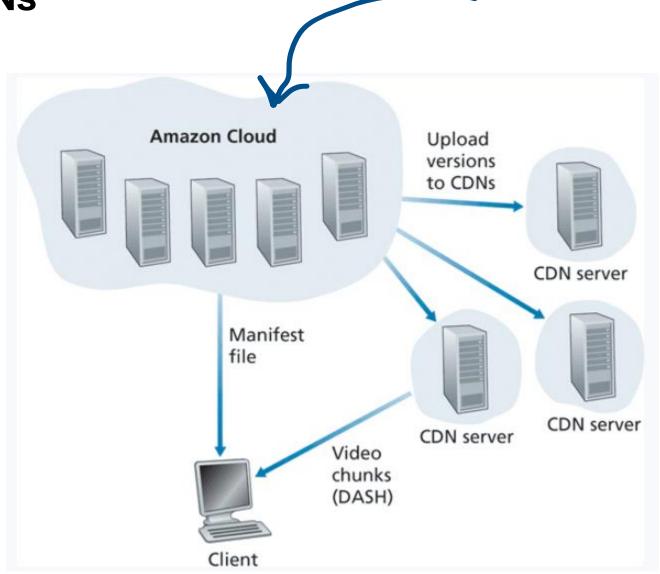




Client

Studio maser versions are uploaded to a private Amazon cloud

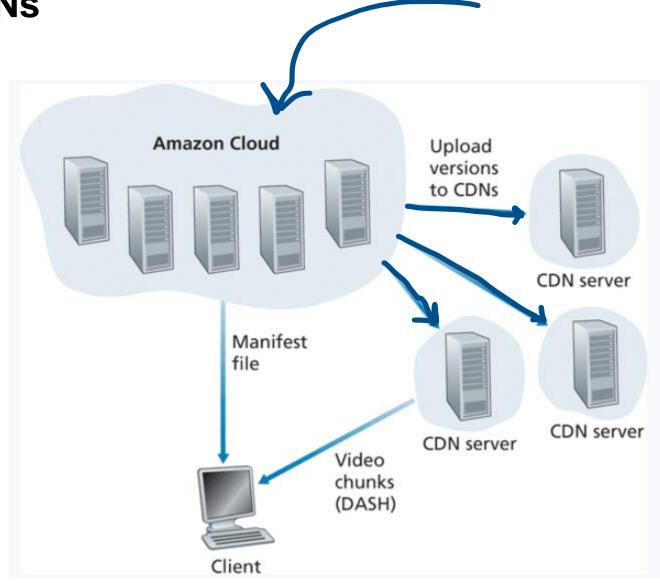




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Videos are processed into many different formats, allowing for DASH





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Videos are processed into many different formats, allowing for DASH

Versions are uploaded to Netflix's CDNs