CSCI 466: Networks

Transport Layer (Multiplexing, Error Checking)

Reese Pearsall Fall 2023

Announcements

PA1 Due Wednesday @ 11:59 PM

- Files must be pushed to a PA1 folder on your GitHub Repo
- Video demo is required
- Submit your repo link to D2L when finished
- Adding time.sleep() after your .send calls can resolve timing issues

Quiz on Friday

Application Layer

Presentation Layer

Session Layer

Transport Layer

Network Layer

Data Link Layer

Physical Layer



Application Layer

Messages from Network Applications



Physical Layer

Bits being transmitted over a copper wire

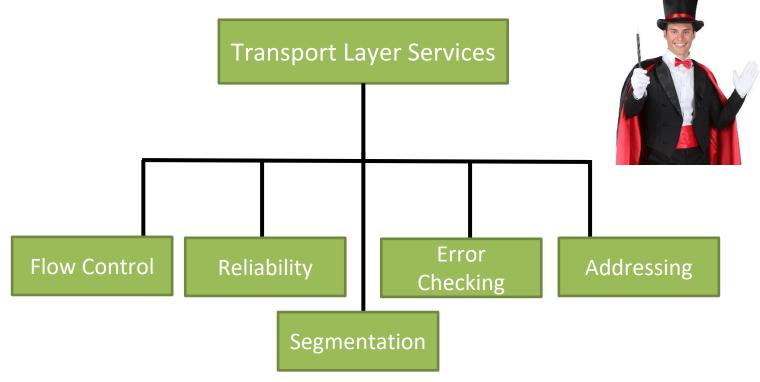
*In the textbook, they condense it to a 5-layer model, but 7 layers is what is most used

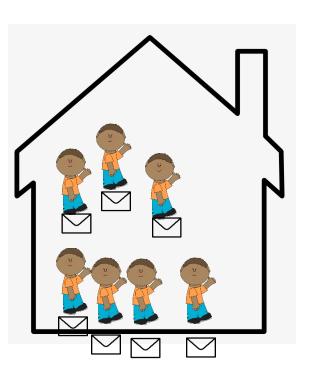
Provides logical end-to-end communication between application processes

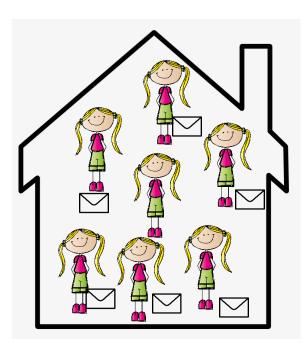
(Network layer handles delivery between hosts)

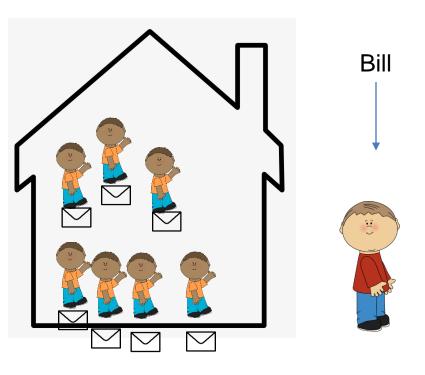
Important Services

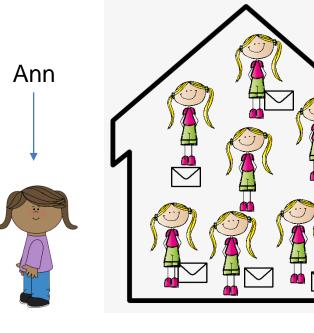
- Flow Control
- Reliability
- Segmentation
- Error Checking
- Addressing

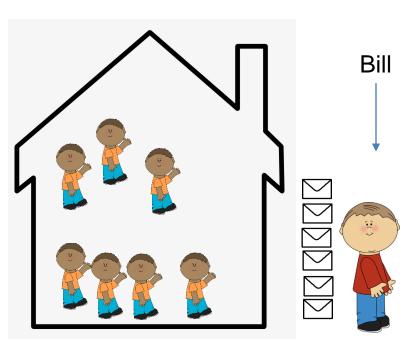


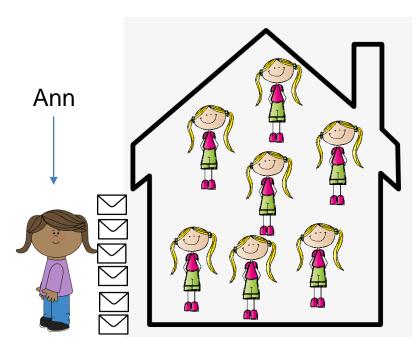




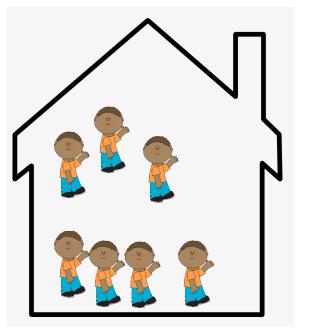


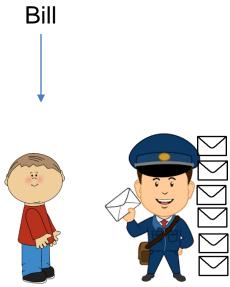


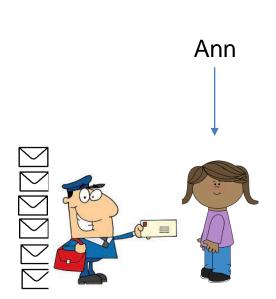


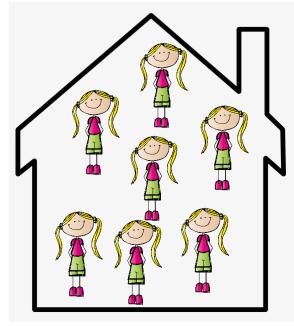


Bill and Ann are responsible for collecting their siblings mail...



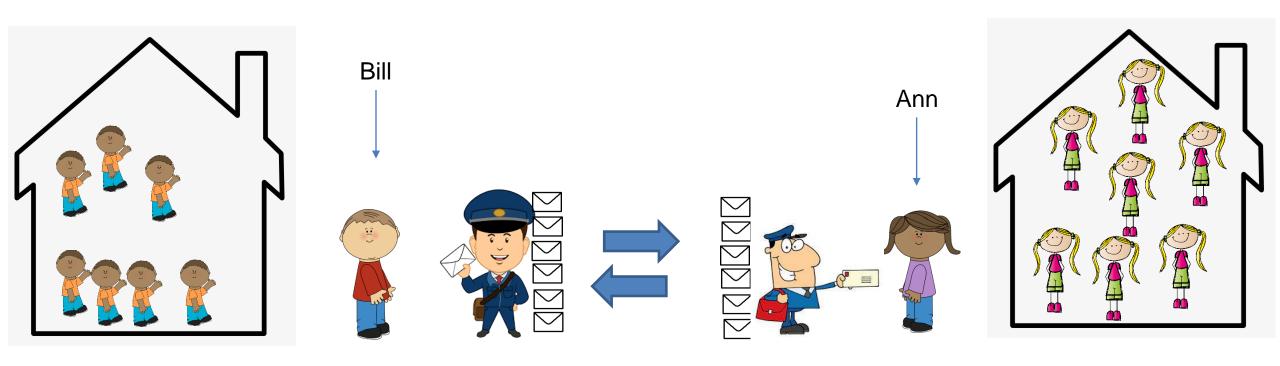






Bill and Ann are responsible for collecting their siblings mail...

And delivering it to the postal service worker



Bill and Ann are responsible for collecting their siblings mail...

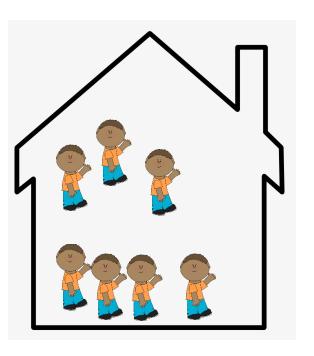
And delivering it to the postal service worker



Bill and Ann are responsible for collecting their siblings mail...

And delivering it to the postal service worker

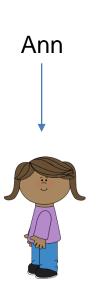
Letters in envelopes = Application messages

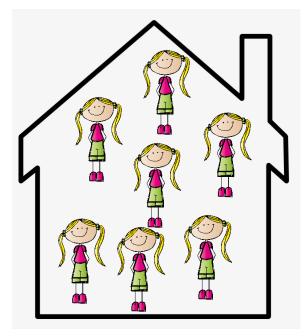




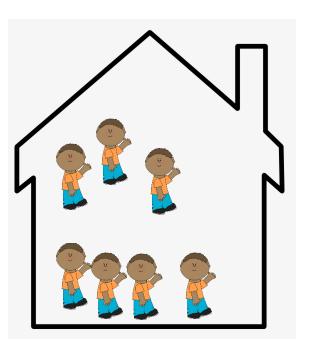








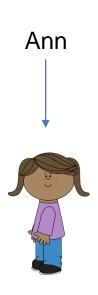
Letters in envelopes = Application messages Cousins = Processes

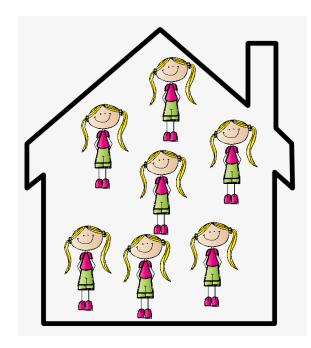




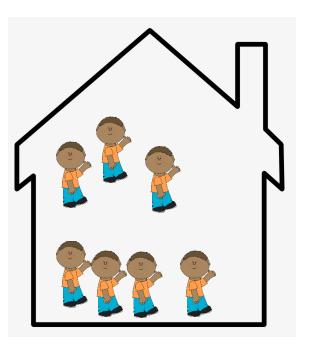








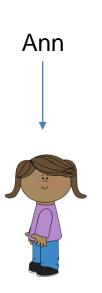
Bill

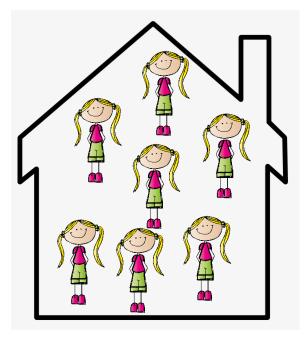


Letters in envelopes = Application messages Cousins = Processes Houses = Hosts/End systems

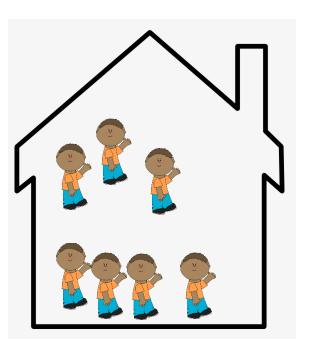






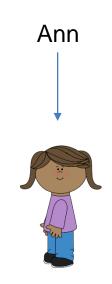


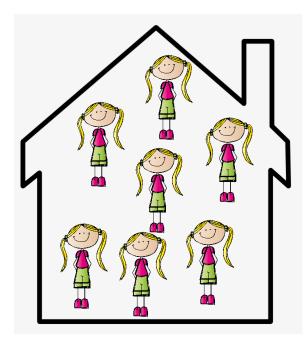
Bill

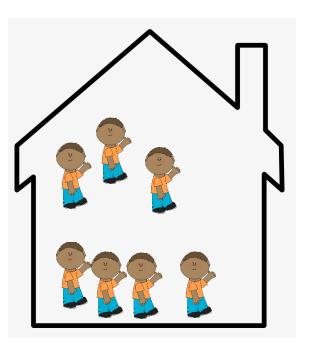


Letters in envelopes = Application messages
Cousins = Processes
Houses = Hosts/End systems
Ann and Bill = Transport Layer Protocol







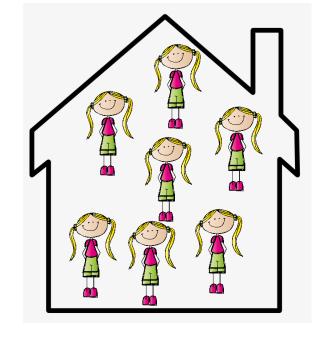


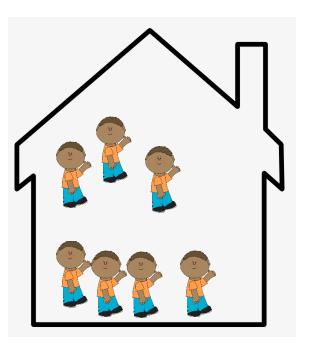
Letters in envelopes = Application messages
Cousins = Processes
Houses = Hosts/End systems
Ann and Bill = Transport Layer Protocol
Postal Service = Network Layer/Network core
Ann



Bill





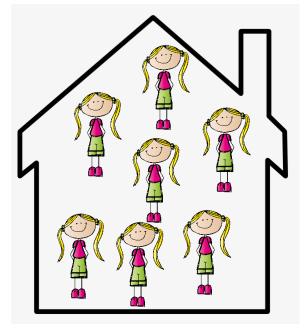


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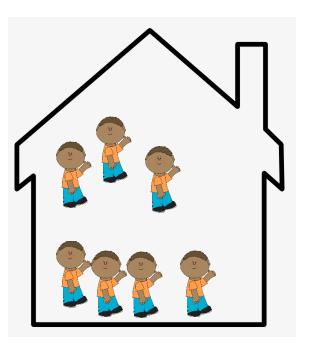


Bill





From the perspective of the cousins, Bill and Ann are the postal service

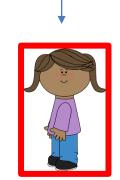


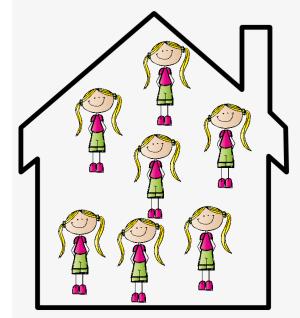
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Bill







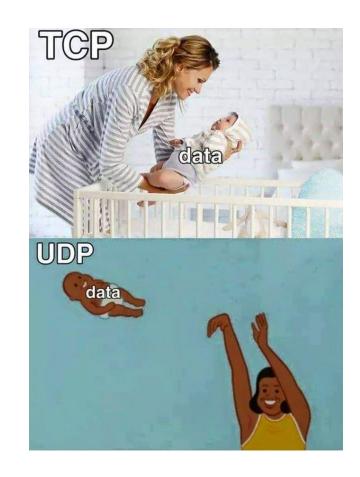
From the perspective of the cousins, Bill and Ann are the postal service

What if Bill and Ann are sick?

TCP – Reliable, Connection-oriented transport layer protocol

UDP – Unreliable, connectionless transport layer protocol

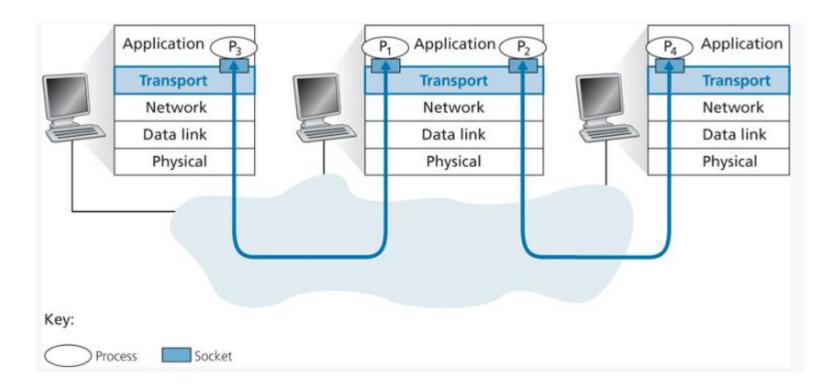
Application	Application-Layer Protocol	Underlying Transport Protocol
Electronic mail	SMTP	TCP
Remote terminal access	Telnet	TCP
Secure remote terminal access	SSH	TCP
Web	HTTP, HTTP/3	TCP (for HTTP), UDP (for HTTP/3)
File transfer	FTP	TCP
Remote file server	NFS	Typically UDP
Streaming multimedia	DASH	TCP
Internet telephony	typically proprietary	UDP or TCP
Network management	SNMP	Typically UDP
Name translation	DNS	Typically UDP



TCP – Reliable, Connection-oriented transport layer protocol

UDP – Unreliable, connectionless transport layer protocol

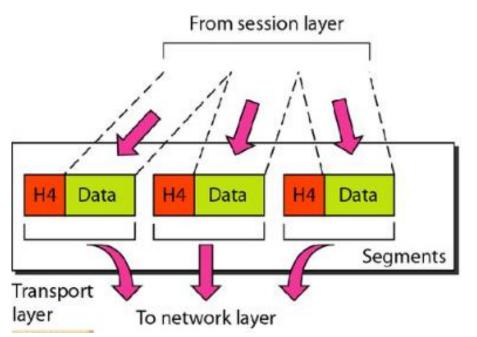
Transport layer delivers to **sockets**

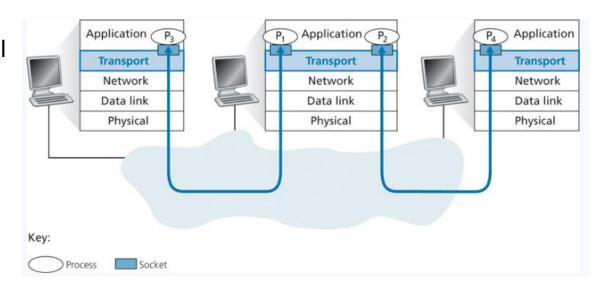


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Transport layer delivers to **sockets**



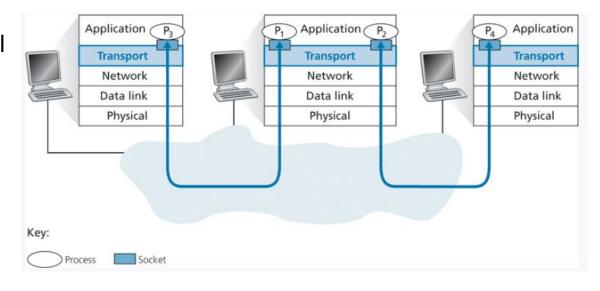


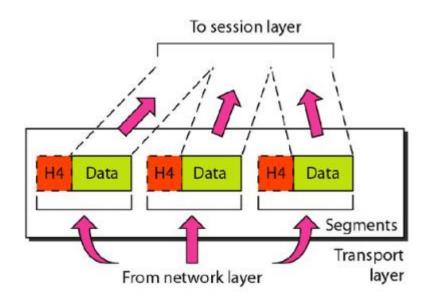
Messages from the Application Layer/Session Layer are split into smaller chunks called **segments**, and passed into the network layer

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Transport layer delivers to **sockets**



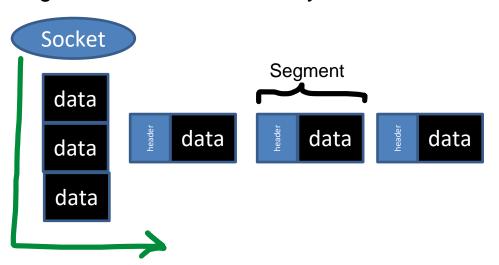


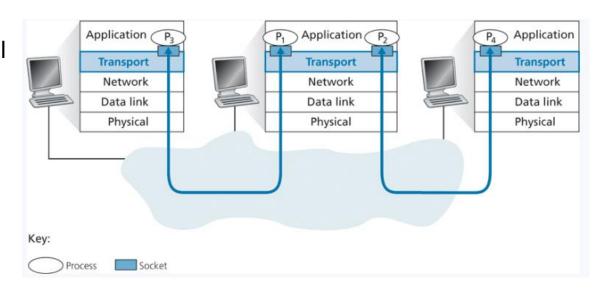
Messages from the network layer arrive as segments. Transport layer must reassemble to send it to the correct process

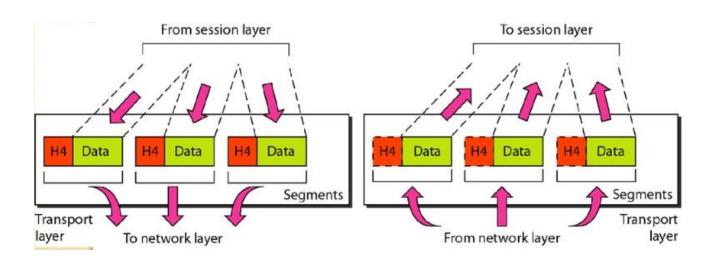
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Multiplexing is the process of gathering chunks from sockets, encapsulating chunks with header information, and passing the segment into the network layer



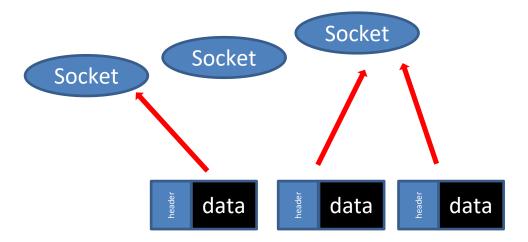


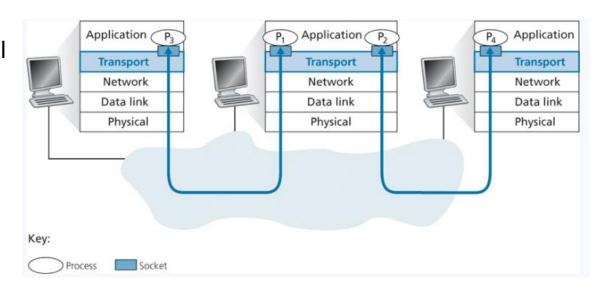


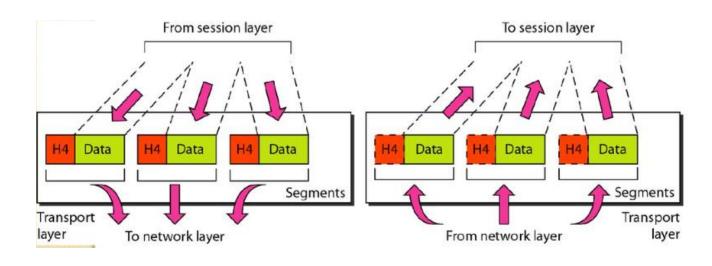
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Demultiplexing is the receiving segments from the transport layer and delivering the segment to the correct socket.



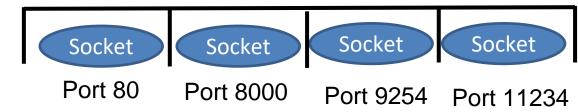


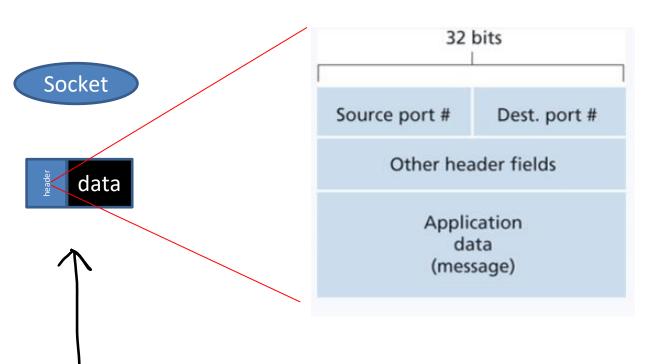




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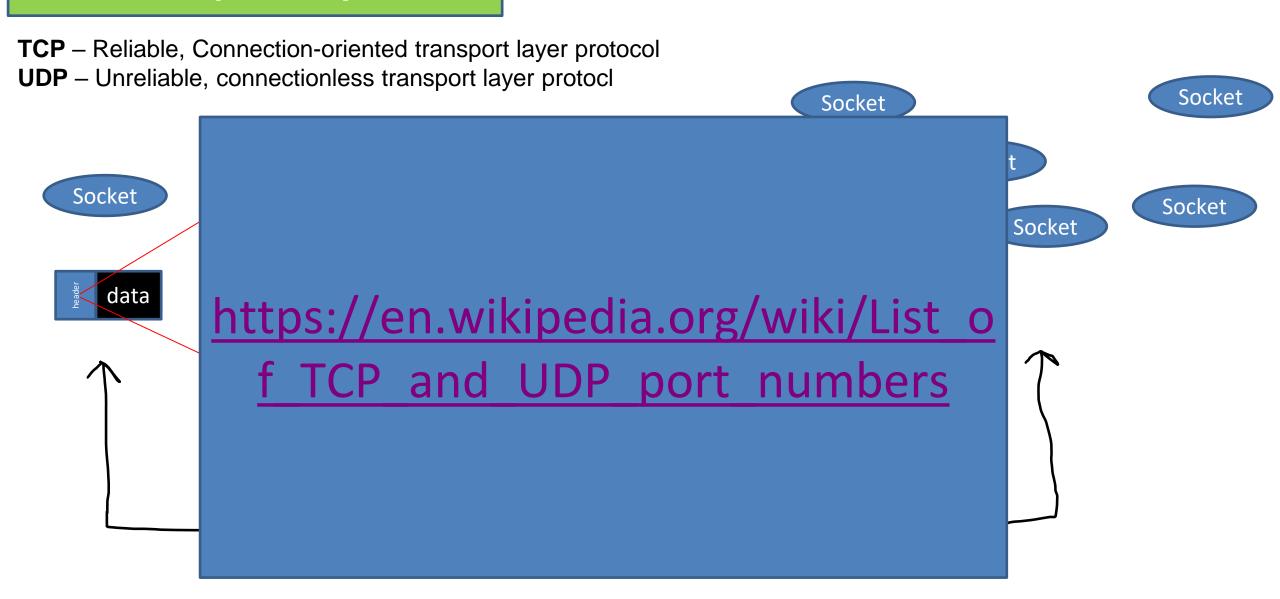




A **port** is a 16-bit number that is an entrance/exit point of a machine

An active port is associated with a specific process or service

Network Layer

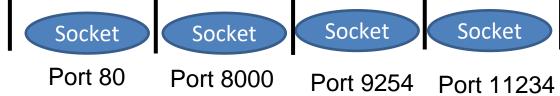


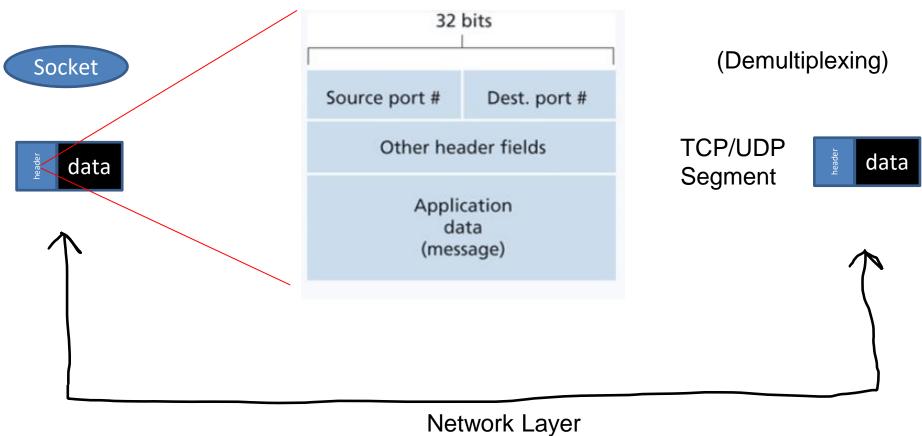
Port	Request type
7	ЕСНО
20	FTP Data
21	FTP Control
22	SSH Remote Login Protocol
23	Telnet
25	Simple Mail Transfer Protocol (SMTP)
37	Time
53	Domain Name System (DNS)
69	Trivial File Transfer Protocol (TFTP)
79	Finger
80	HTTP
110	POP3
115	Simple File Transfer Protocol (SFTP)
137	NetBIOS Name Service
139	NetBIOS Datagram Service
143	Interim Mail Access Protocol (IMAP)
156	SQL Server
161	SNMP
194	Internet Relay Chat (IRC)
389	Lightweight Directory Access Protocol (LDAP)
443	HTTPS
445	Microsoft-DS
458	Apple QuickTime
546	DHCP Client
547	DHCP Server



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UDP – Unreliable, connectionless transport layer protocl



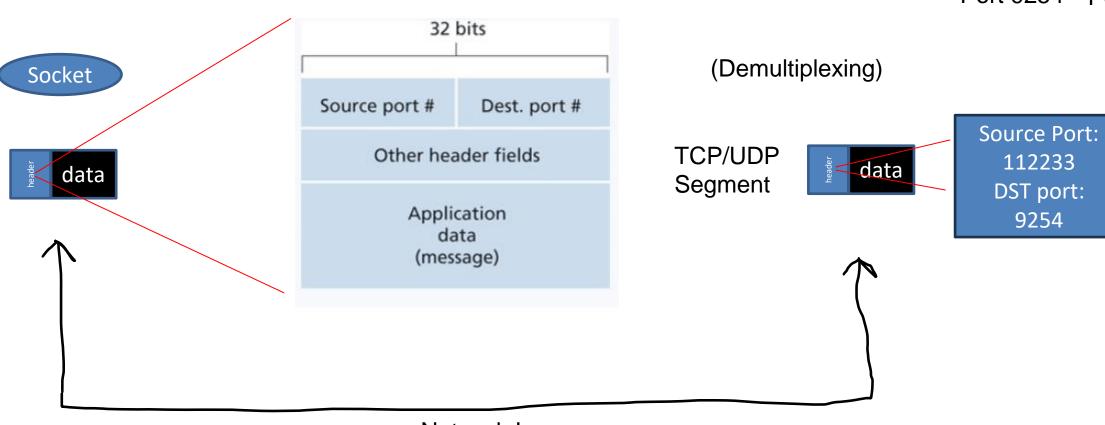




TCP – Reliable, Connection-oriented transport layer protocol **UDP** – Unreliable, connectionless transport layer protocl

Socket Socket Socket Socket

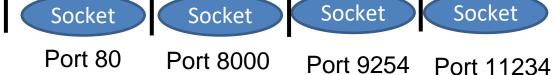
Port 80 Port 8000 Port 9254 Port 11234

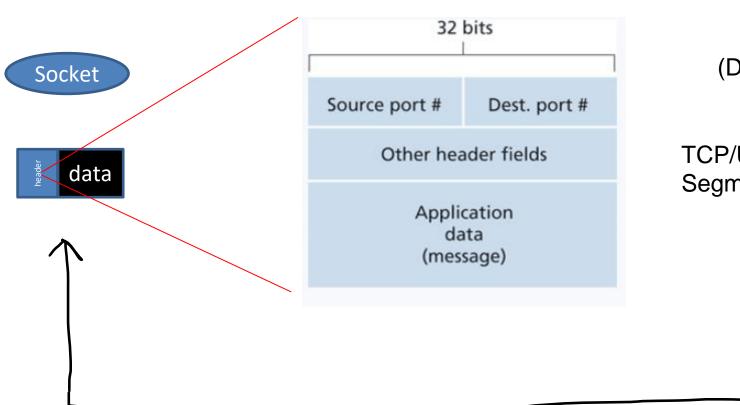


Network Layer

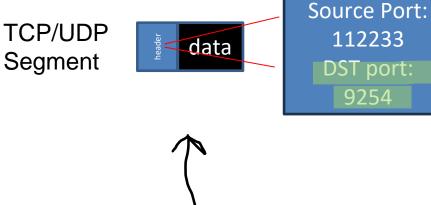


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(Demultiplexing)



Network Layer



Socket

Socket

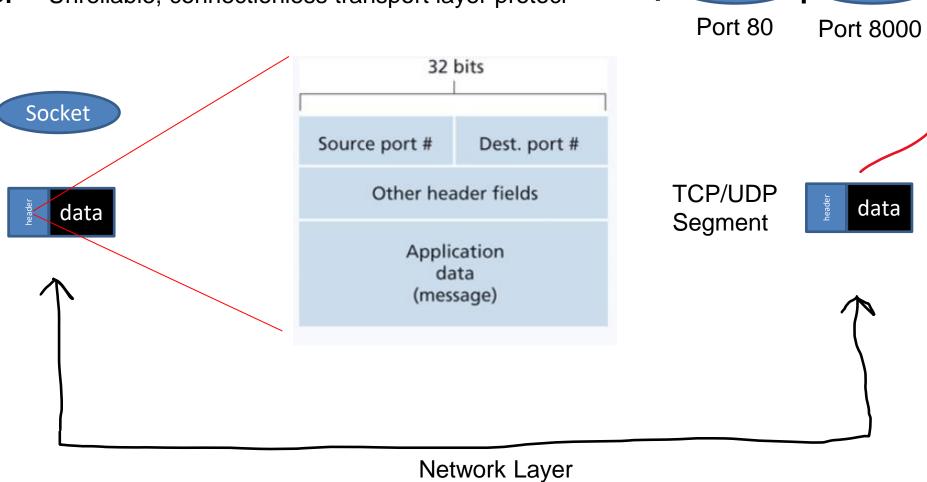
Socket

Port 9254

Socket

Port 11234

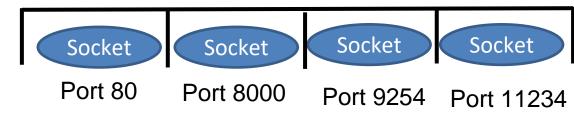
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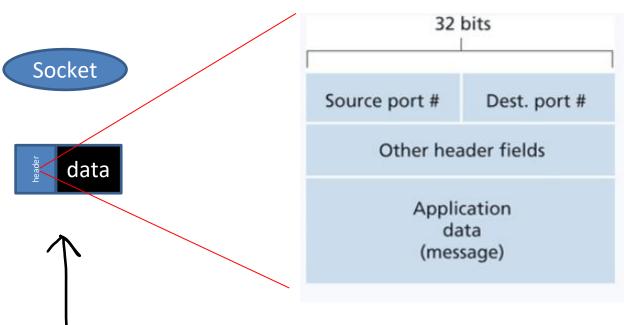




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An active port is associated with a specific process or service

When developing a new application, we need to assigned the application a port number (greater than 1024)

Network Layer

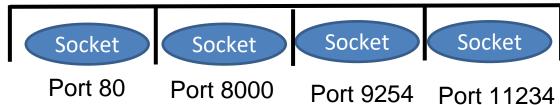
clientSocket.bind(('', 19157))

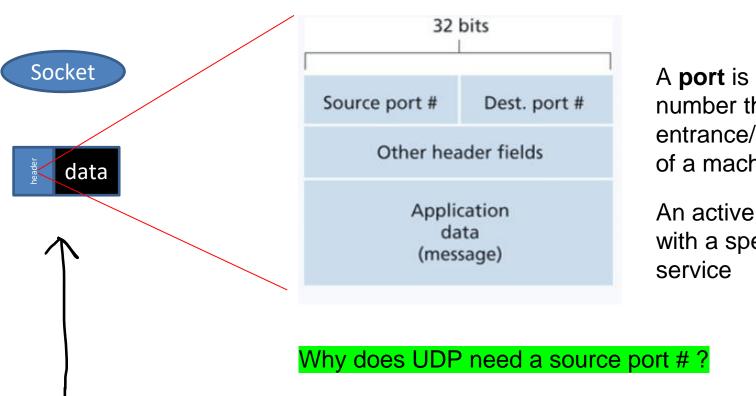
clientSocket = socket(AF_INET, SOCK_DGRAM)



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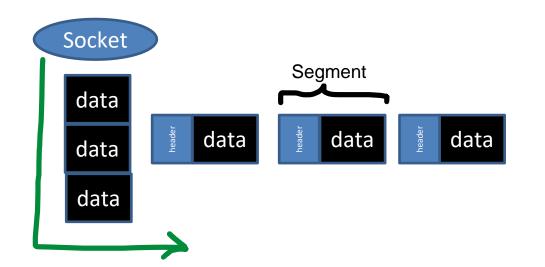
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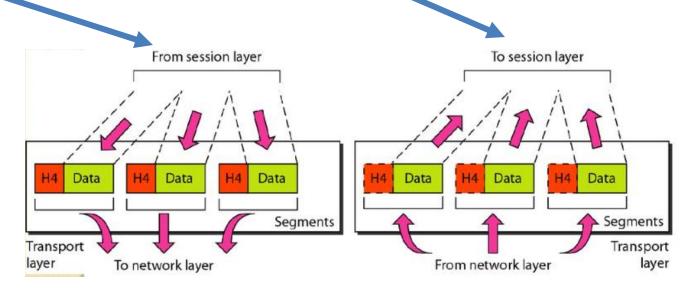
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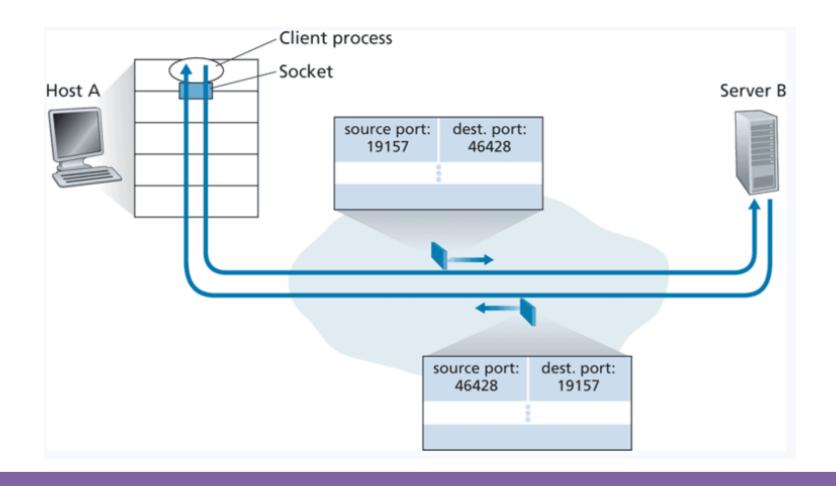
Network Layer

Multiplexing is the process of gathering chunks from sockets, encapsulating chunks with header information, and passing the segment into the network layer

Demultiplexing is the receiving segments from the transport layer and delivering the segment to the correct socket.

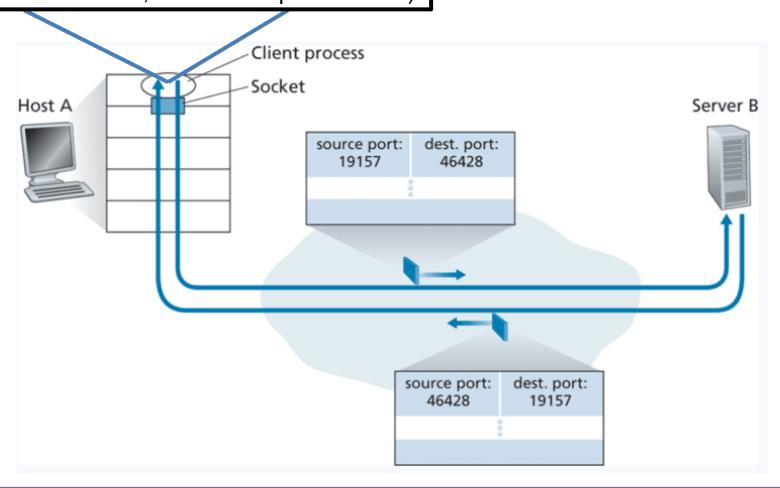


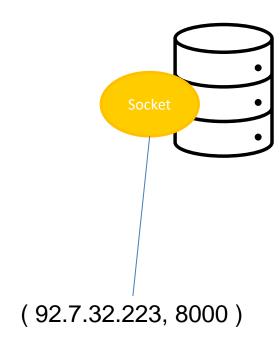


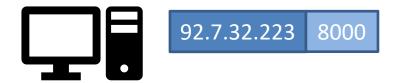


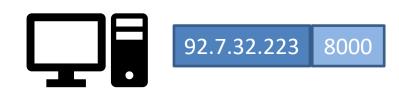
UDP sockets are identified by a two-tuple

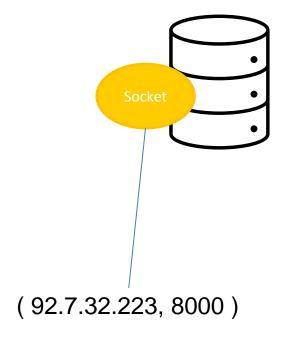
(destination IP address, destination port number)

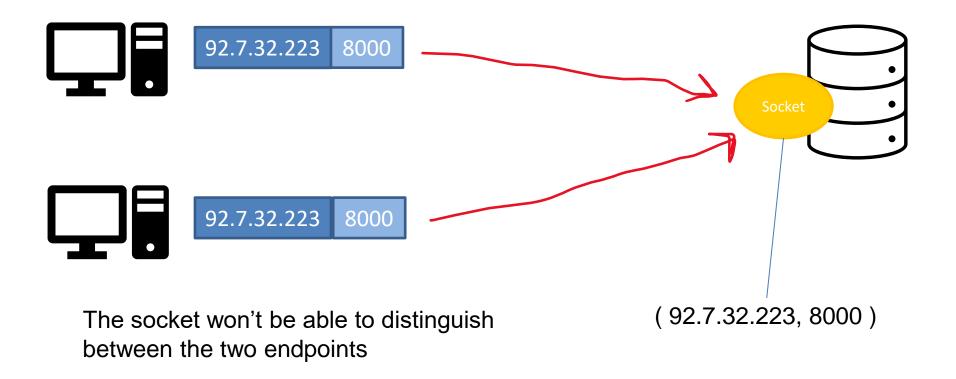










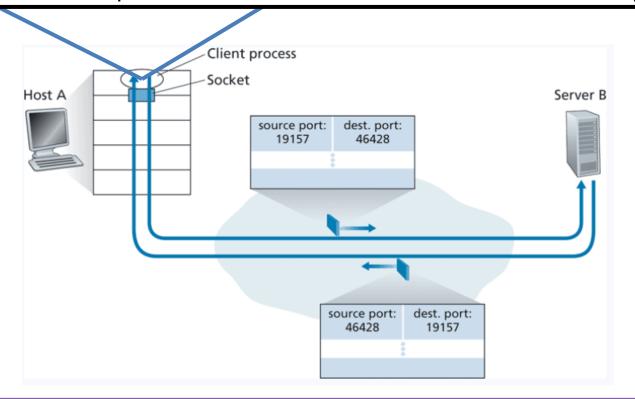


The two segments will be directed to the same process

TCP sockets are identified by a four-tuple

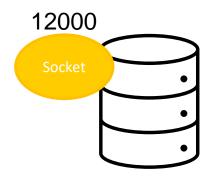
A host will demultiplex segments using all of these values

(source IP address, source port number, destination IP address, destination port number)

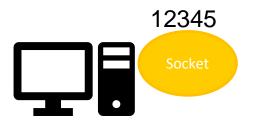


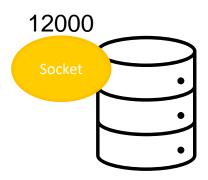
TCP servers will have a "welcoming socket" before creating the processes' socket (12000)





The TCP server application has a "welcoming socket," that waits for connection-establishment requests from TCP clients on port number 12000.

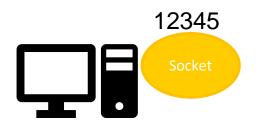


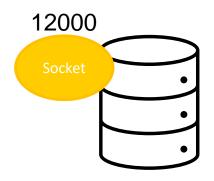


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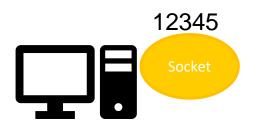
The TCP client creates a socket and sends a connection establishment request segment with the lines:

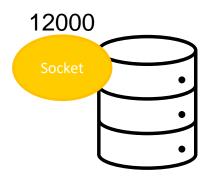
```
clientSocket = socket(AF_INET, SOCK_STREAM)
clientSocket.connect((serverName,12000))
```





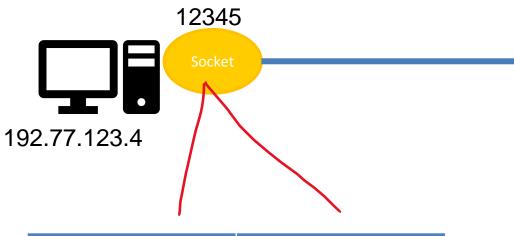
A connection-establishment request is nothing more than a TCP segment with destination port number 12000 and a special connection-establishment bit set in the TCP header. The segment also includes a source port number that was chosen by the client.



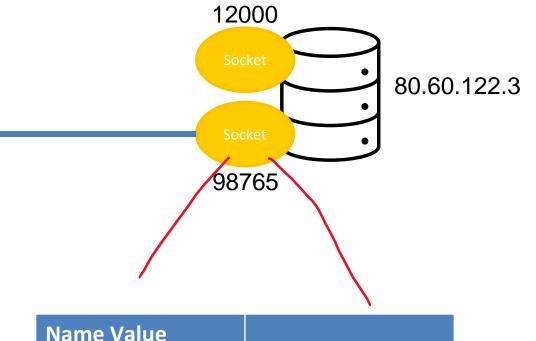


When the host operating system of the computer running the server process receives the incoming connection-request segment with destination port 12000, it locates the server process that is waiting to accept a connection on port number 12000. The server process then creates a new socket:

connectionSocket, addr = serverSocket.accept()

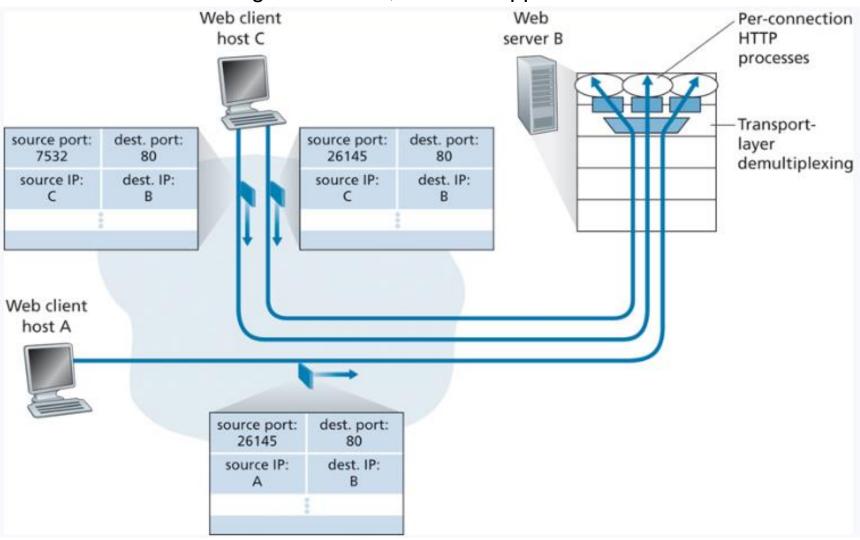


Name	Value
Src IP	192.77.123.4
Src Port	12345
Dst IP	80.60.122.3
Dst Port	98765

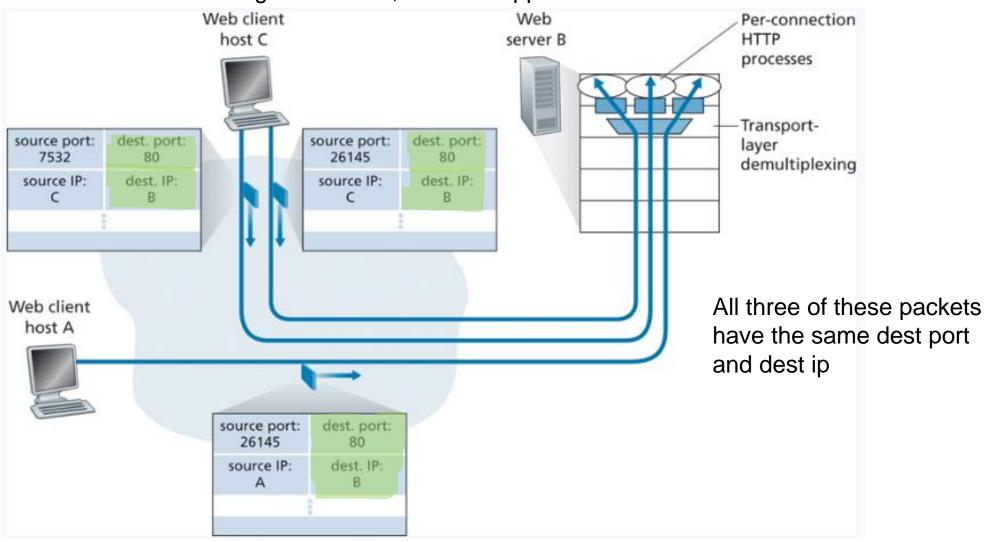


Name Value	
Src IP	80.60.122.3
Src Port	98765
Dst IP	192.77.123.4
Dst Port	12345

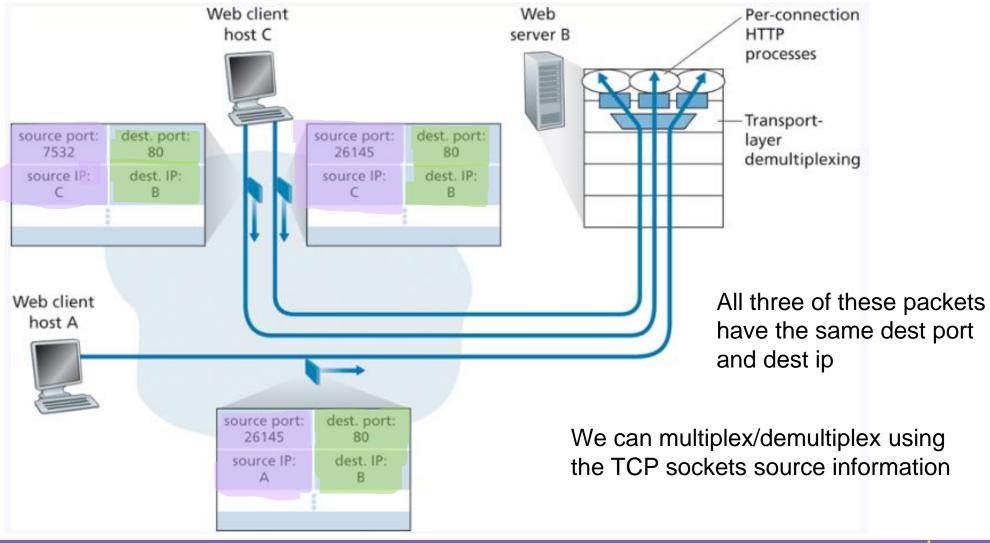




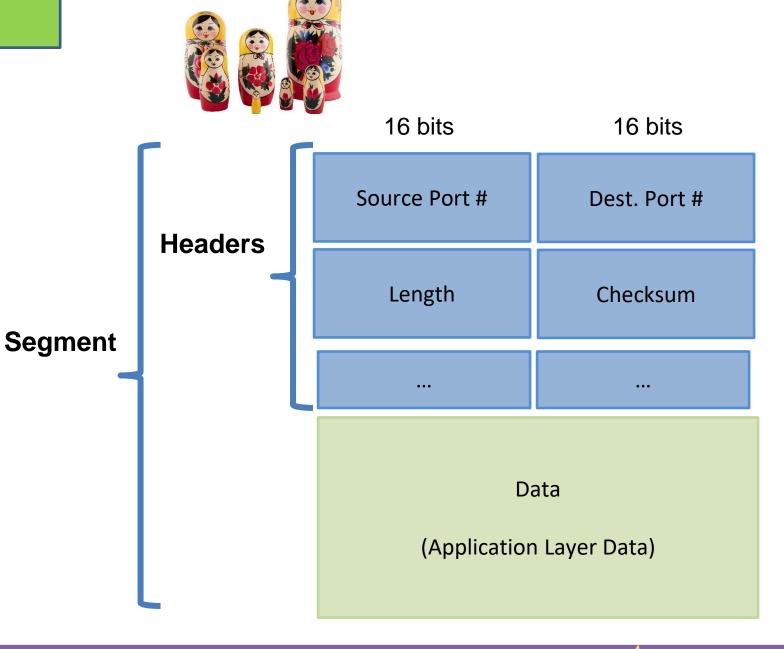
Now that we are using four values, we can support simultaneous connections



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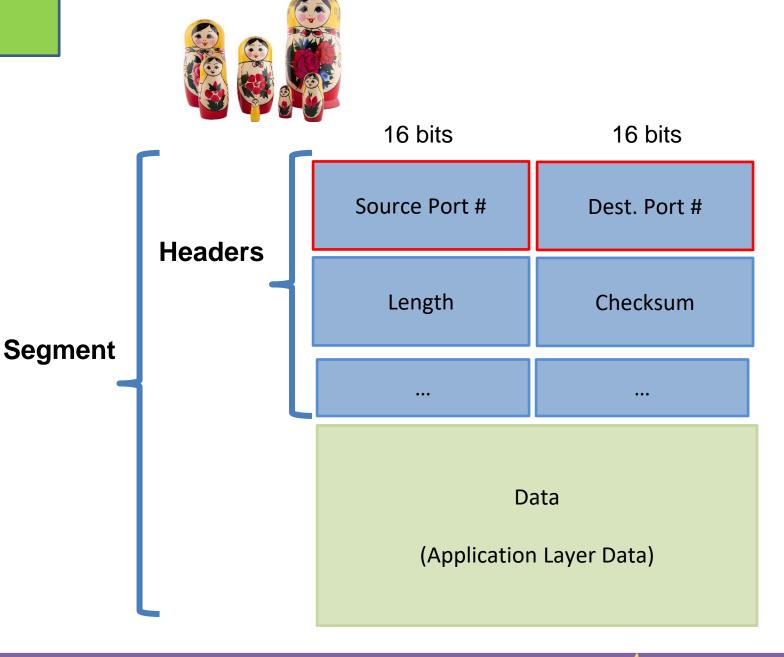


Our application layer message (HTTP request, DNS Query, FTP data) gets split into chunks, and each chunk is encapsulated in a transport layer header



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Source port and Dest Port are attached to our packet

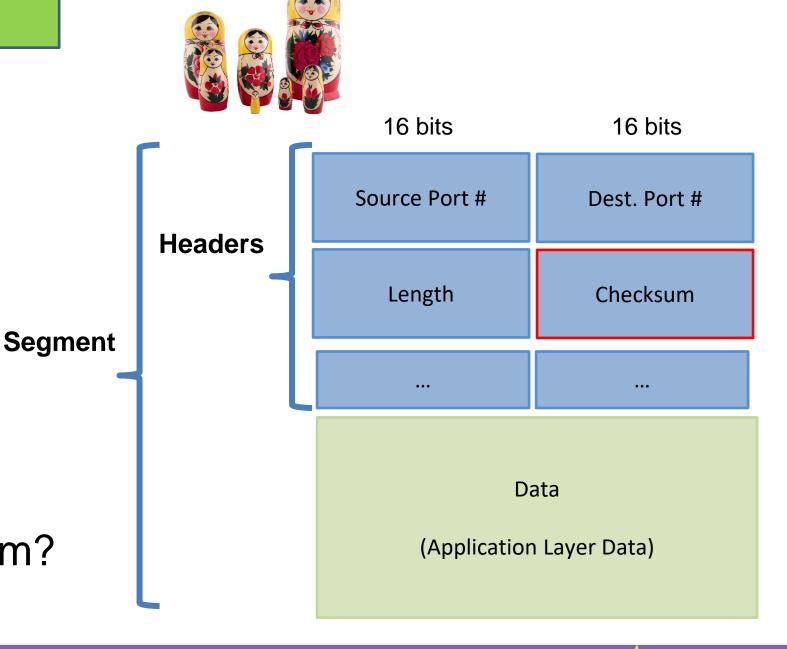


Our application layer message (HTTP request, DNS Query, FTP data) gets split into chunks, and each chunk is encapsulated in a transport layer header

attached to our packet

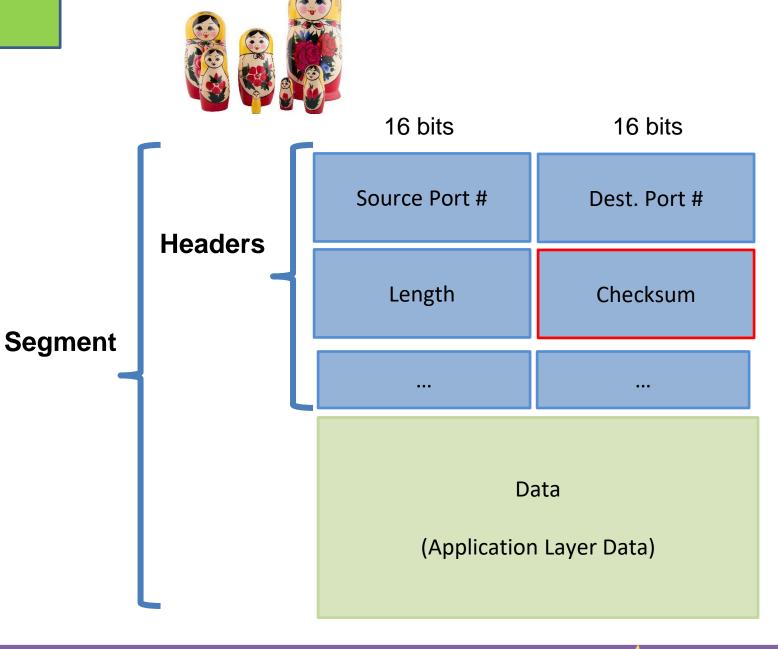
Source port and Dest Port are

What is the checksum?

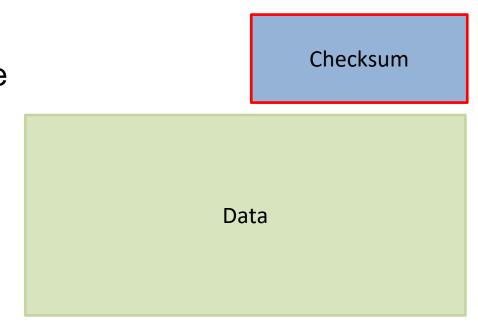


Important Services

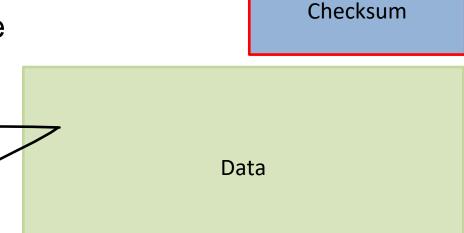
- Flow Control
- Reliability
- Segmentation
- Error Checking
- Addressing



Transport layer provides a **checksum** that is used to determine whether bits within the segment have been altered/corrupted



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O110011001100000

O101010101010101

1000111100001100

Split up data into 16-bit "words"

Checksum

 Split up data into 16-bit "words" Sum up all words

Checksum

0110011001100000 0101010101010101 1000111100001100

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0110011001100000 0101010101010101 1000111100001100

Split up data into 16-bit "words" Sum up all words Compute the ones compliment

 Checksum

0110011001100000 0101010101010101 31000111100001100

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Checksum

0110011001100000 0101010101010101 1000111100001100

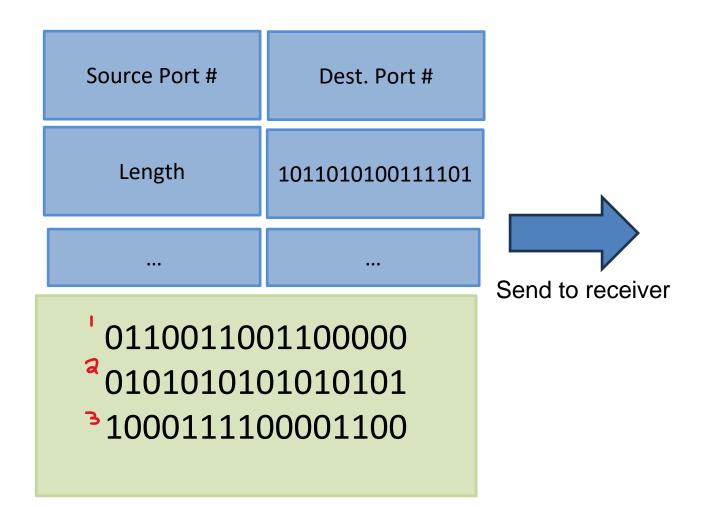
This is our checksum value!

Split up data into 16-bit "words" Sum up all words Compute the ones compliment

1011010100111101

0110011001100000 0101010101010101 1000111100001100

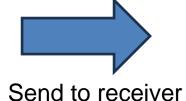
This is our checksum value!



Source Port # Dest. Port #

Length 1011010100111101

...



0110011001100000

²01010101010101

31000111100001100

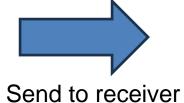
When the receiver gets the segment, it will compute the sum of all four 16-bit words (including the checksum)

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Source Port # Dest. Port #

Length 10110100111101

... ...



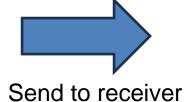
0110011001100000 0101010101010101 1000111100001100

When the receiver gets the segment, it will compute the sum of all four 16-bit words (including the checksum)

Source Port # Dest. Port #

Length 101101010111101

... ...



Source Port # Dest. Port #

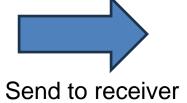
Length

1011010100111101

•••

•••

0110011001100000 0101010101010101 1000111100001100

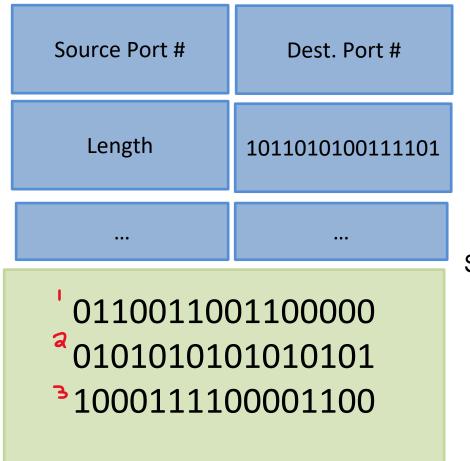


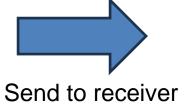


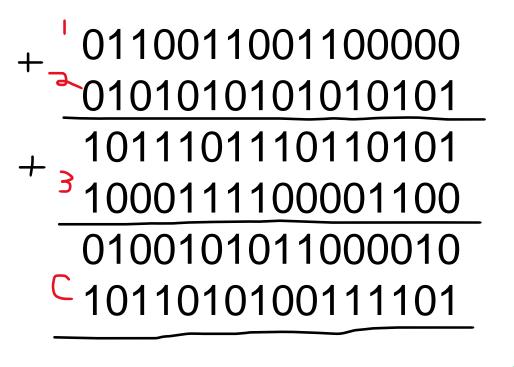
When the receiver gets the segment, it will compute the sum of all four 16-bit words (including the checksum)

+ 01100110010000 01010101010101 + 10111011101101 1000111100001100 0100101011000010 - 1011010100111101

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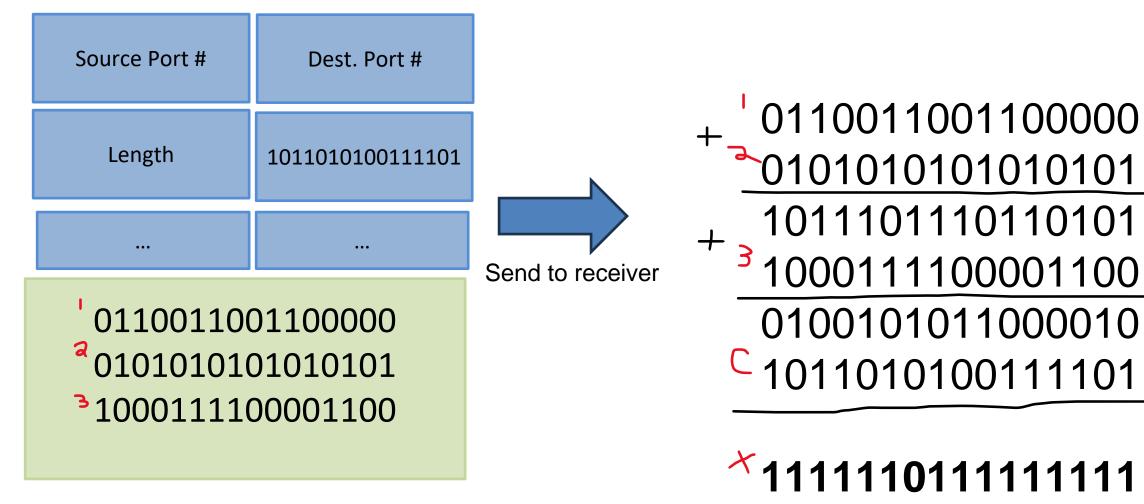




1111111111111111

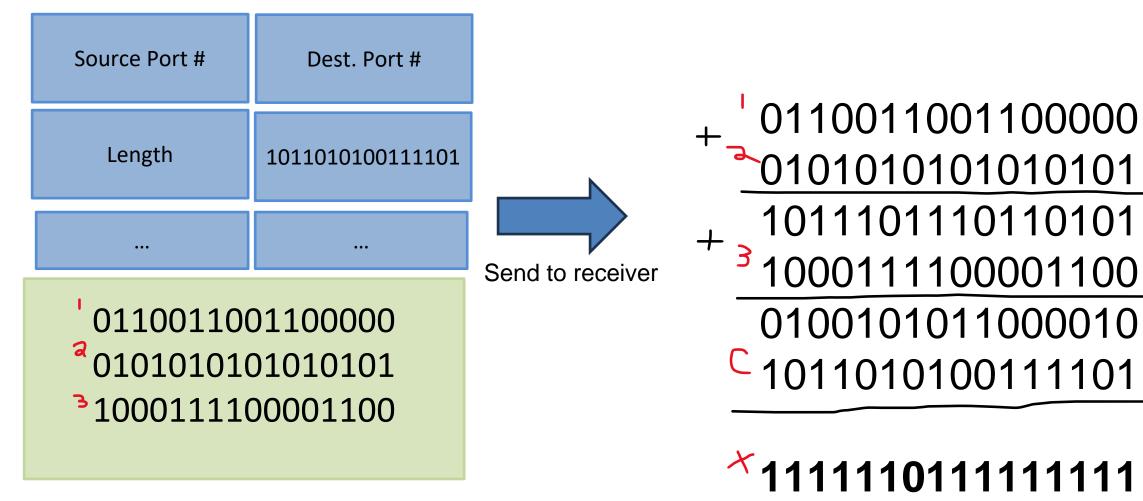
If the result of this operation is **not** all 1s, then data **must** have been corrupted in the segment

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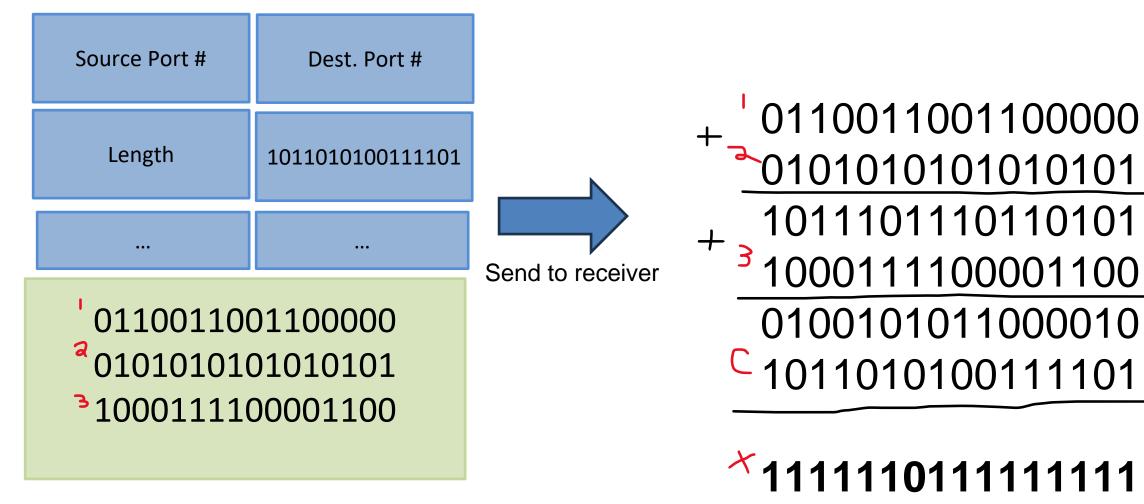
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The receiver will do that for each packet, if the checksum is invalid, it may ask for a retransmission

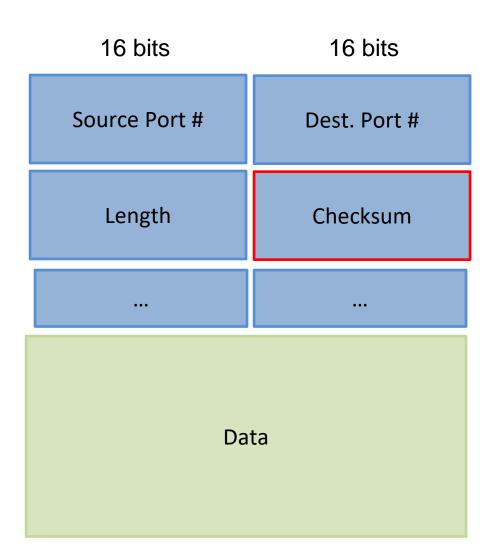
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Why do error checking here?

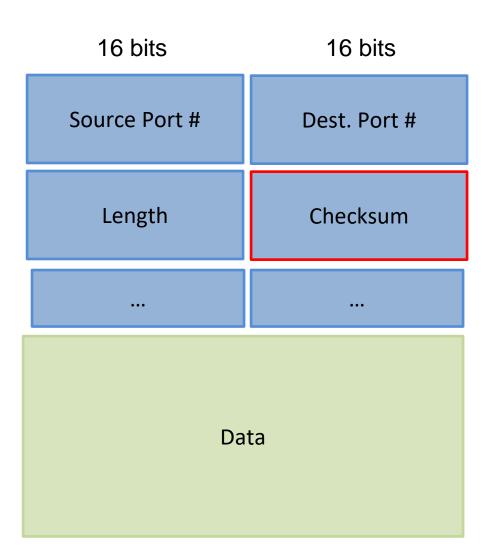
Protocols at other layers (link layer) can also do error checking



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However, there is no guarantee that our packet of information will use these protocols



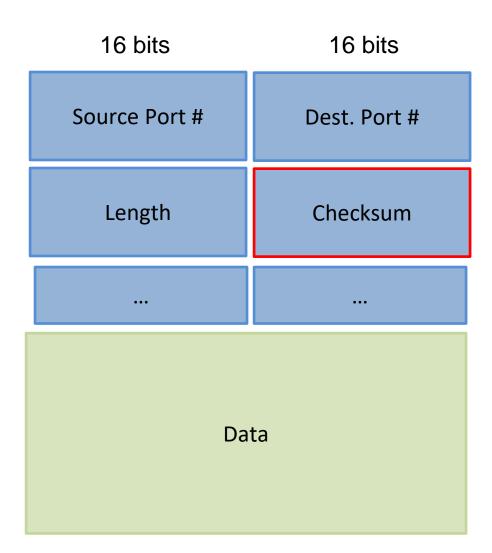
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Protocols at other layers (link layer) can also do error checking

However, there is no guarantee that our packet of information will use these protocols

End-to-end principle states that since certain functionality such as error detection, must be implemented on an end-end bases

Functionality places at the lower levels may be redundant or of little value when compared to the cost of providing them at a higher level



Important Services

- Flow Control
- Reliability
- Segmentation ✓
- Error Checking
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