

# CSCI 466: Networks

Reliable Data Transfer (RDT)

Reese Pearsall  
Fall 2024

# PA1 Due **Sunday**

**NO CLASS ON  
FRIDAY**

### Quiz 2 on Friday

- DNS
- SMTP
- FTP
- P2P
- CDNs/Caches
- Transport Layer
- Reliable Data Transfer

*Opens at 8AM, closes at 11:59 PM*

# Announcements

Tuesday, September 26, 2023 at 3:00:02 PM

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**Reese Pearsall**

unified\reese.pearsall | reesepearsall@montana.edu

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- Quiz 3 (
- DNS
  - SMTF
  - DTP
  - P2P
  - Trans
  - Reliak

# OSI Model

Application Layer

Presentation Layer \*

Session Layer \*

Transport Layer

Network Layer

Data Link Layer

Physical Layer

Application Layer

Messages from Network Applications



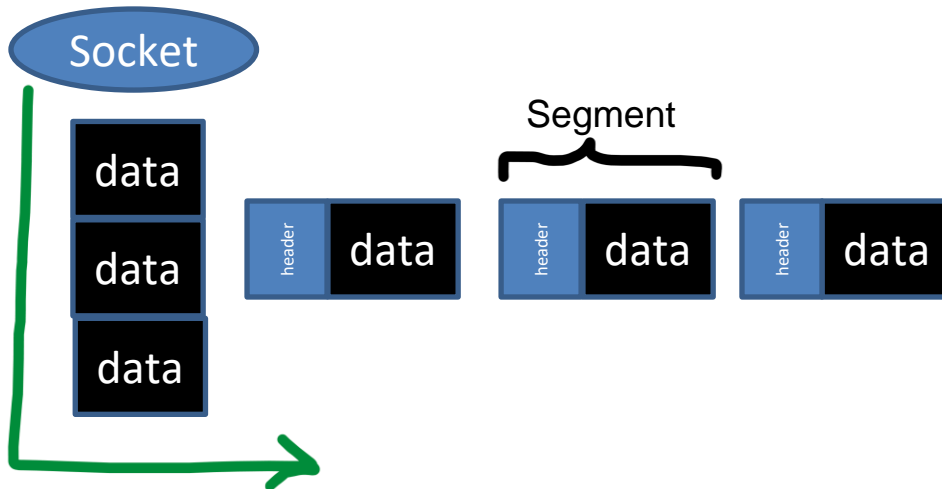
Physical Layer

Bits being transmitted over some medium

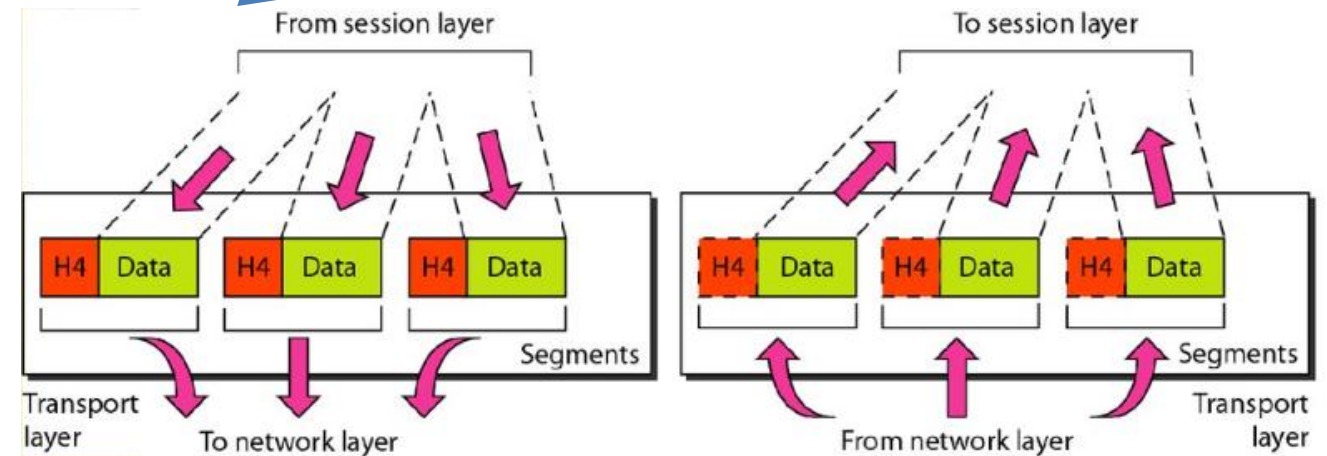
*\*In the textbook, they condense it to a 5-layer model, but 7 layers is what is most used*

# Transport Layer

**Multiplexing** is the process of gathering chunks from sockets, encapsulating chunks with header information, and passing the segment into the network layer

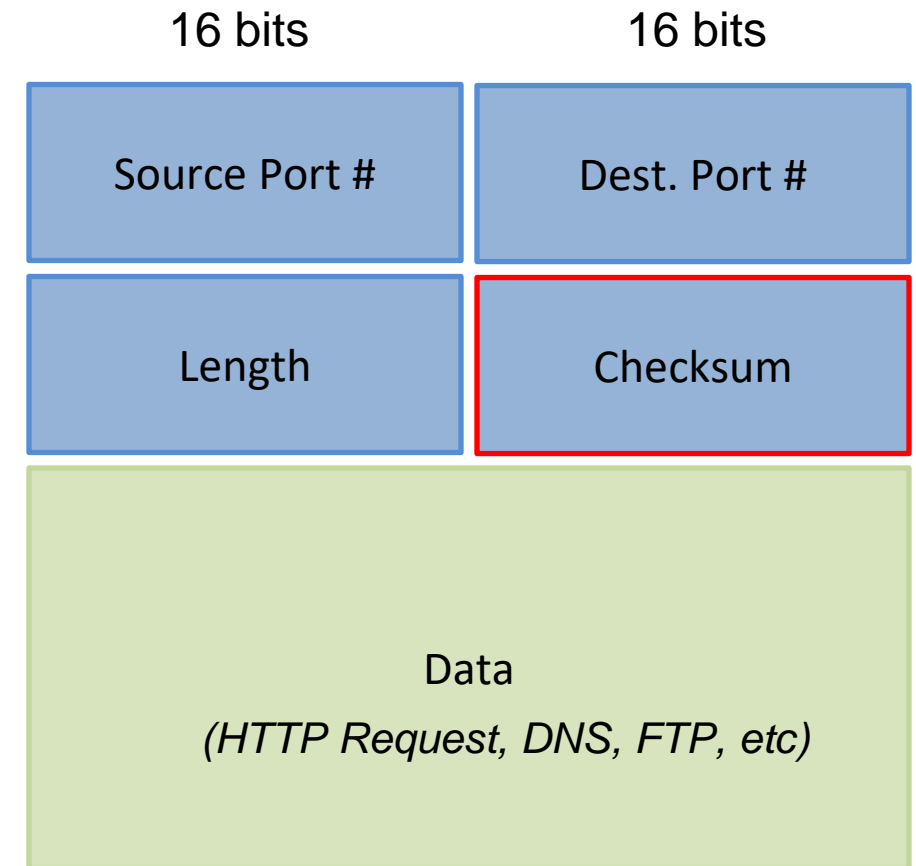


**Demultiplexing** is the receiving segments from the transport layer and delivering the segment to the correct socket.

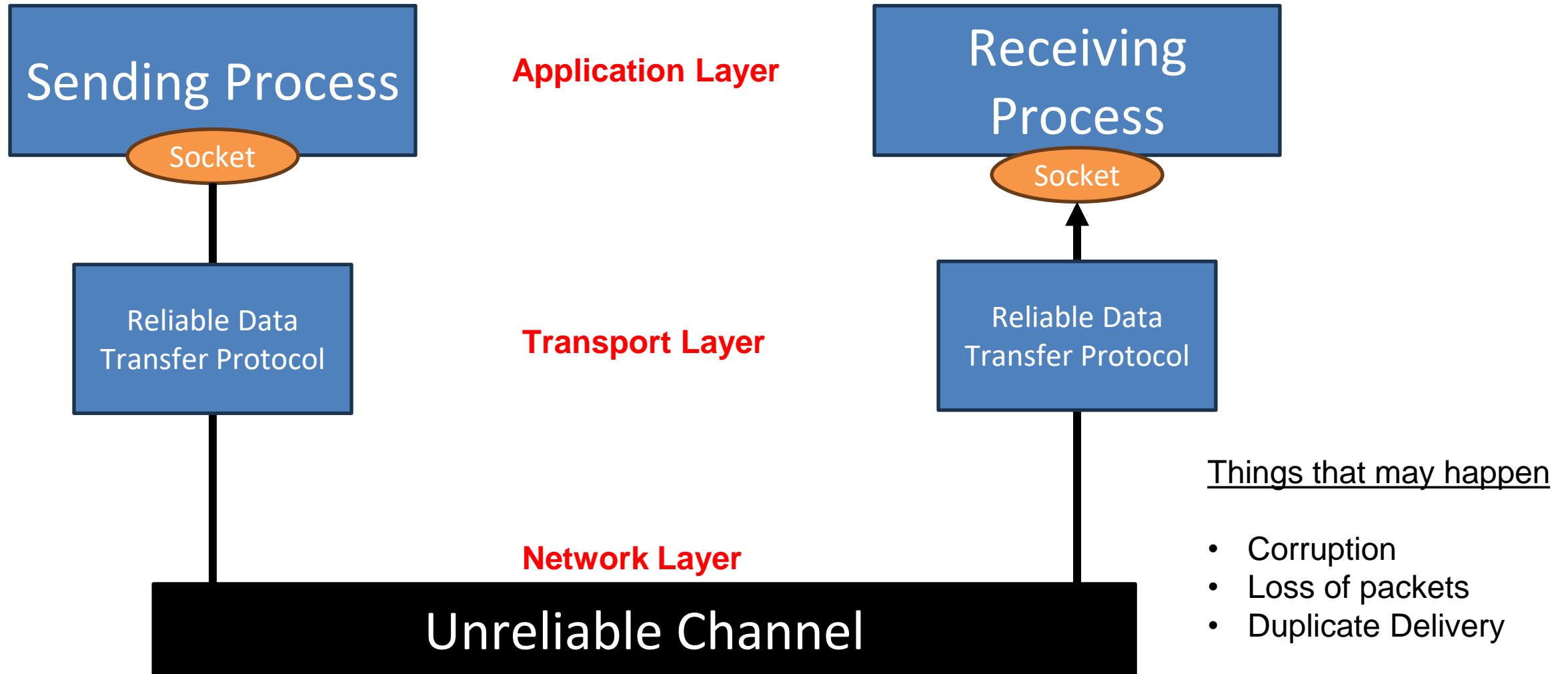


# Transport Layer

Transport Layer provides a **checksum** that is used to determine whether bits within a segment have been altered



# Transport Layer



# Transport Layer

Sending Process

Socket

Reliable Data  
Transfer Protocol

Application Layer

Transport Layer

Network Layer

Unreliable Channel

Receiving Process

Socket

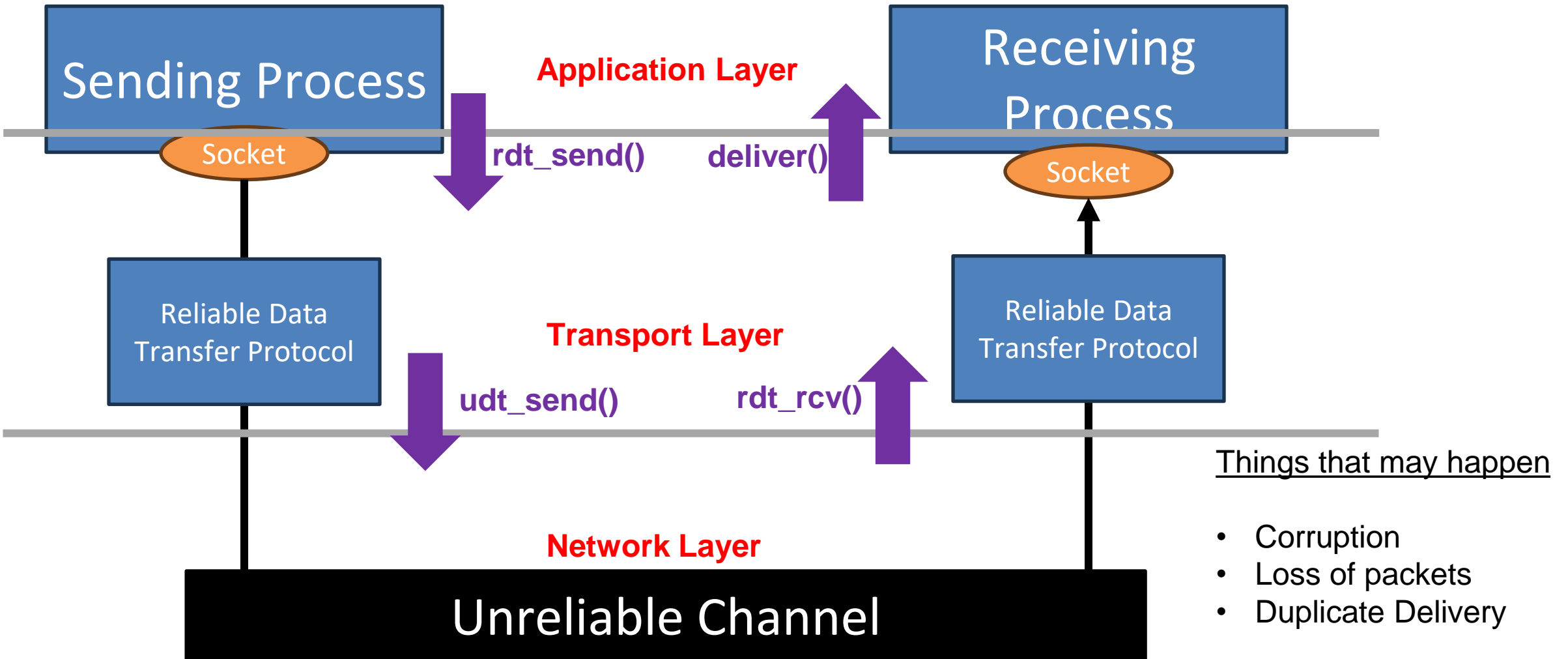
Reliable Data  
Transfer Protocol

Things that may happen

- Corruption
- Loss of packets
- Duplicate Delivery

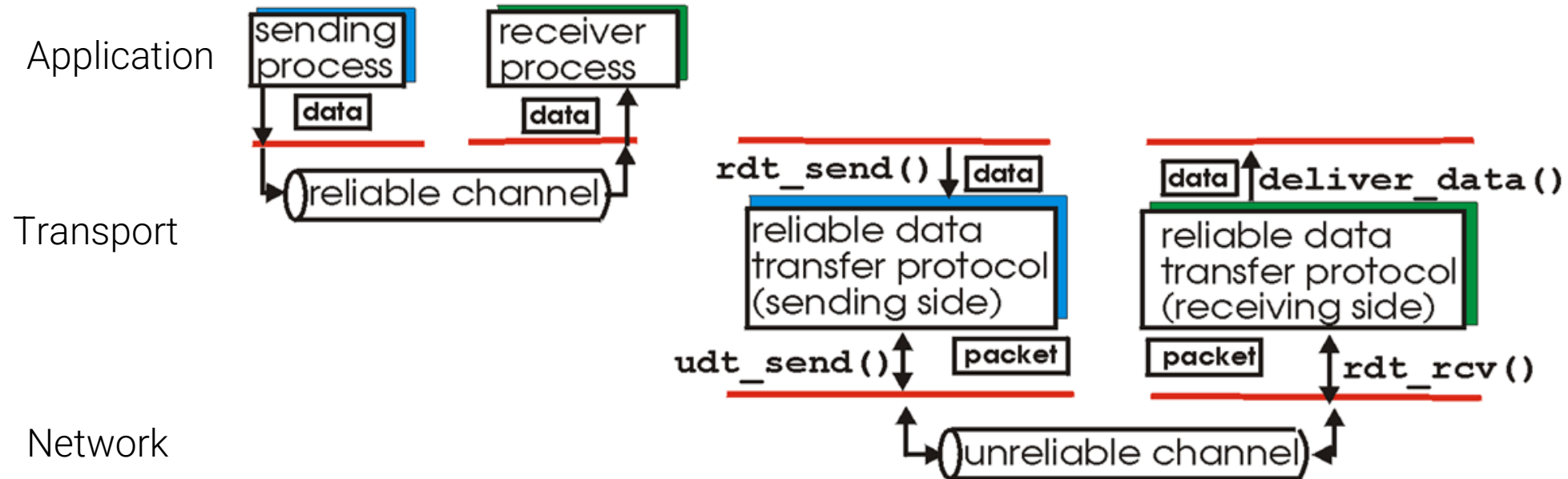


# Transport Layer



# Transport Layer

## Reliable Data Transfer



Characteristics of unreliable channel will determine complexity of reliable data transfer protocol (rdt)

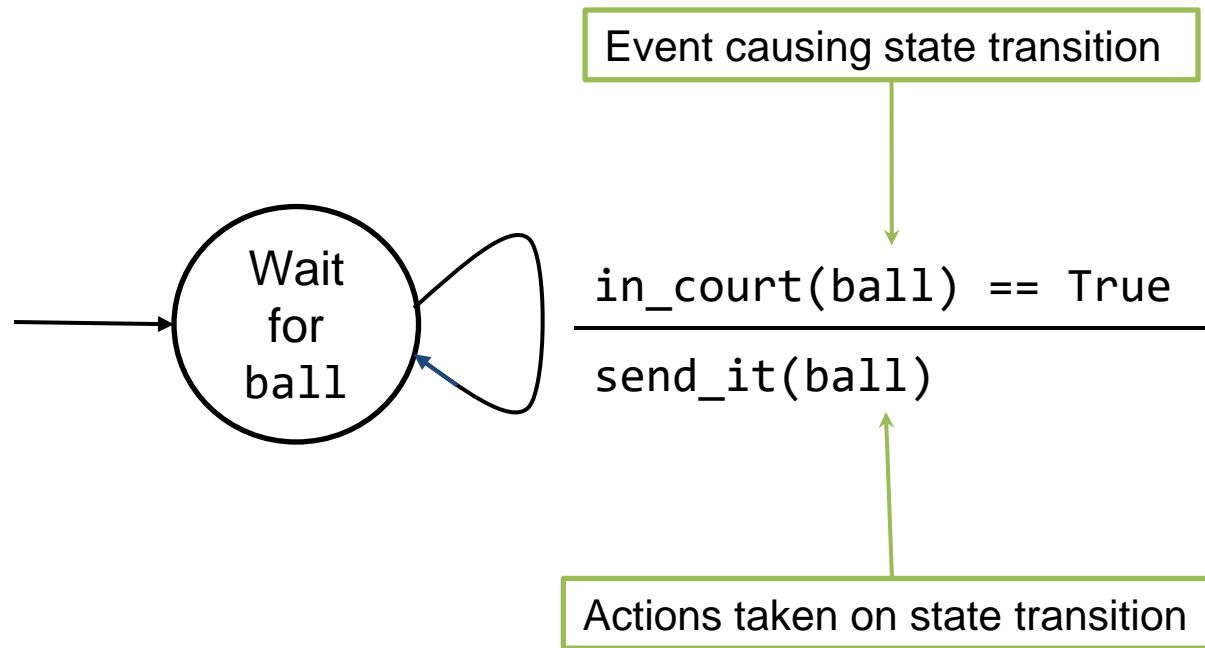
What are some ways in which the network channel can be unreliable?

### Things to consider:

- Corruption
- Loss of packets
- Duplicate delivery



Bruce Lee FSM



# Transport Layer

## Reliable Data Transfer 1.0

### RDT 1.0

#### Assumptions:

- Unidirectional long data flows
- Perfectly reliable channel
- No bit errors
- No packet loss
- No packet reordering



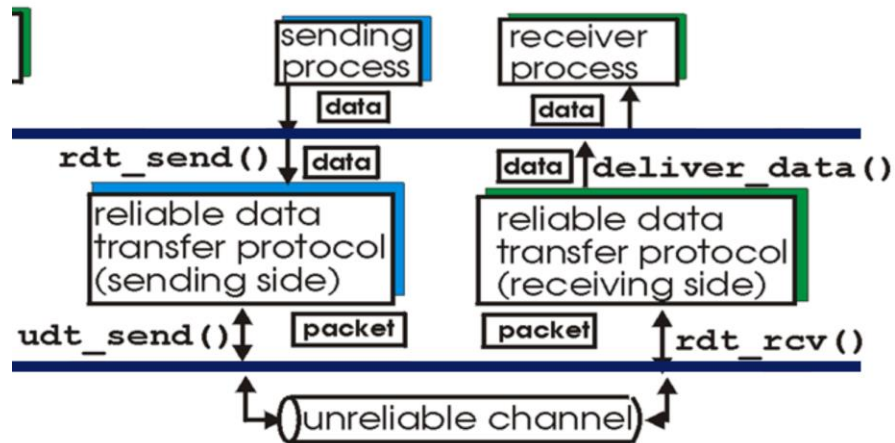
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## Reliable Data Transfer 1.0

### RDT 1.0

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# Transport Layer

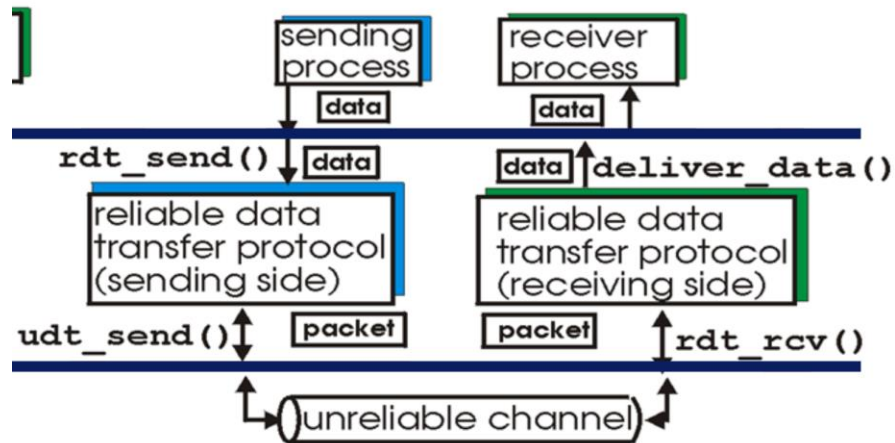
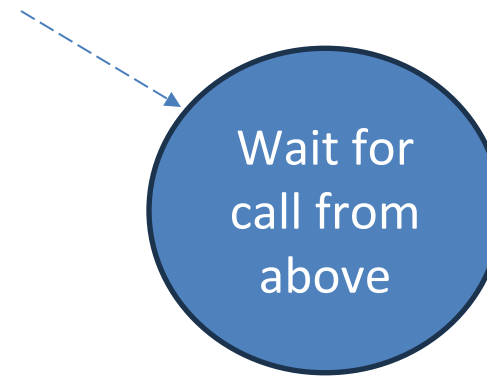
## Reliable Data Transfer 1.0

RDT 1.0

Assumptions:

- Unidirectional long data flows
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Sender



# Transport Layer

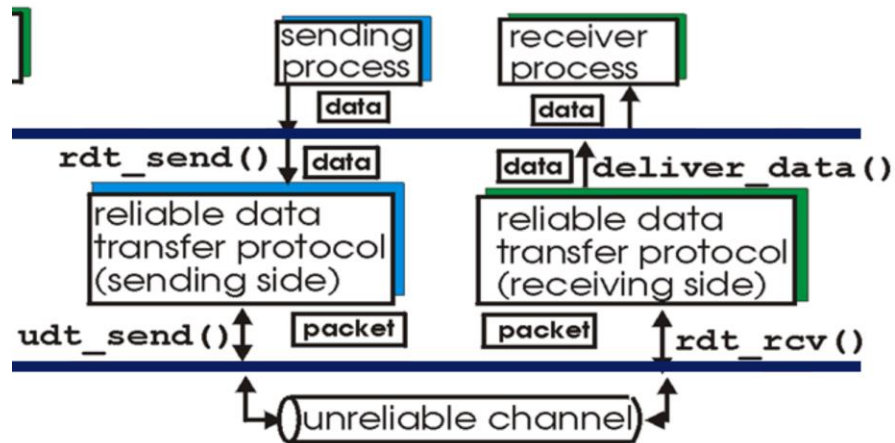
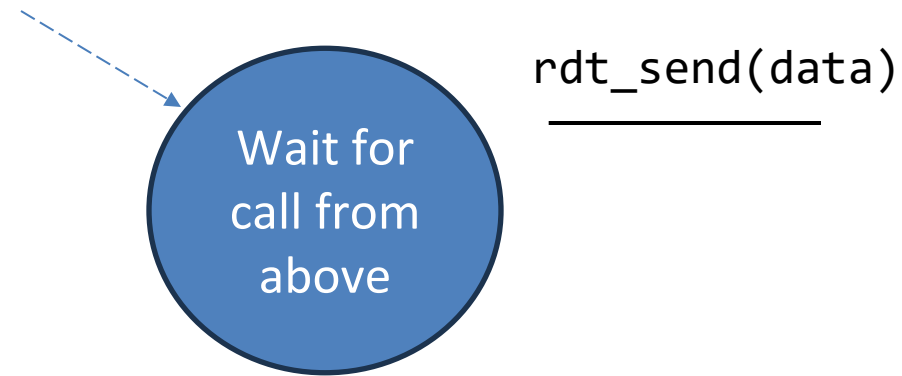
## Reliable Data Transfer 1.0

RDT 1.0

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### Sender





# Transport Layer

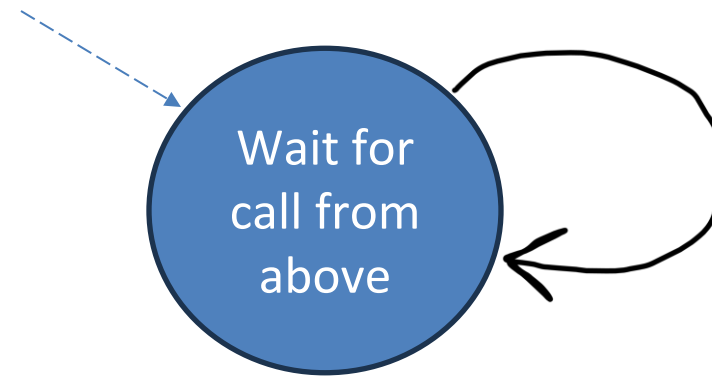
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RDT 1.0

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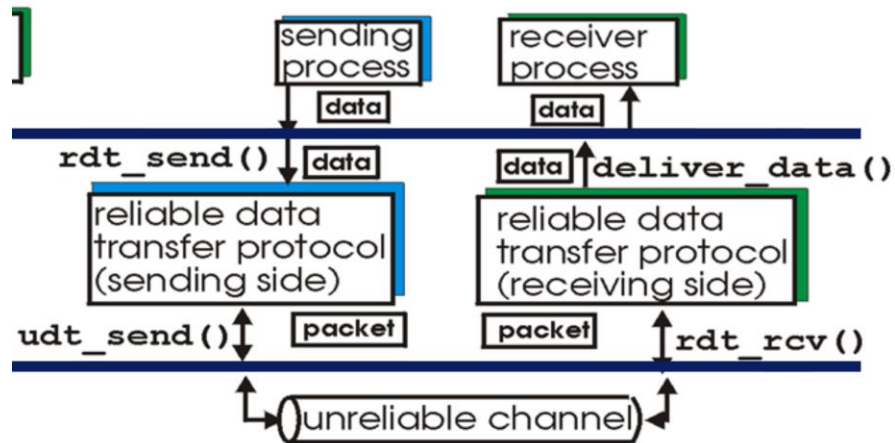
### Sender



rdt\_send(data)

packet = make\_pkt(data)

udt\_send(packet)

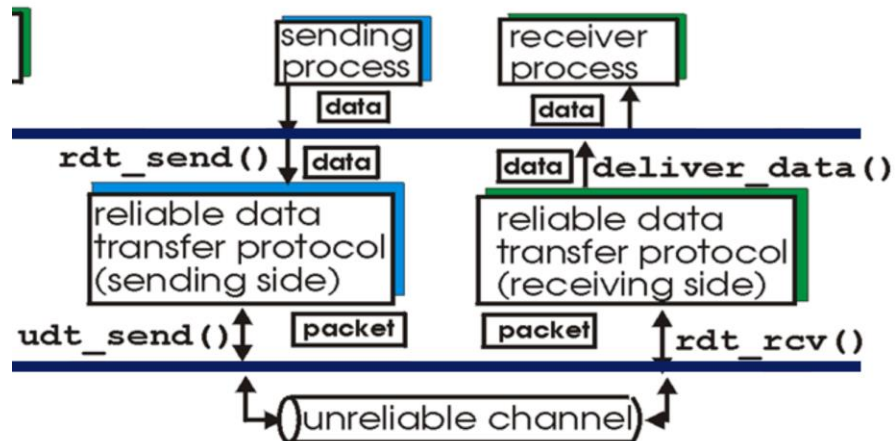


# Transport Layer

## RDT 1.0

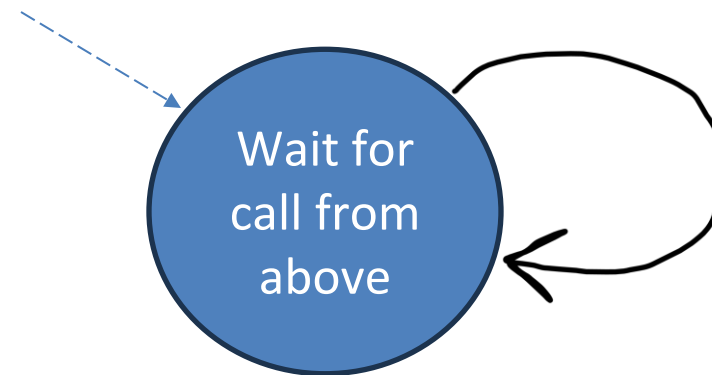
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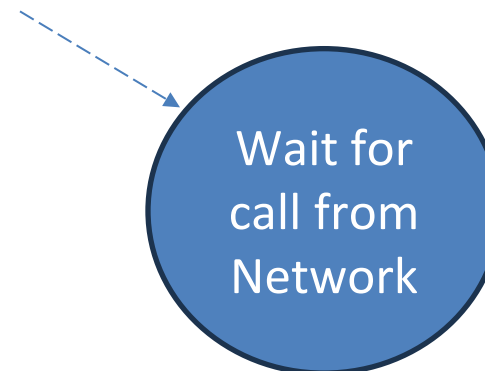
## Reliable Data Transfer 1.0

### Sender



```
rdt_send(data)
    packet = make_pkt(data)
    udt_send(packet)
```

### Receiver



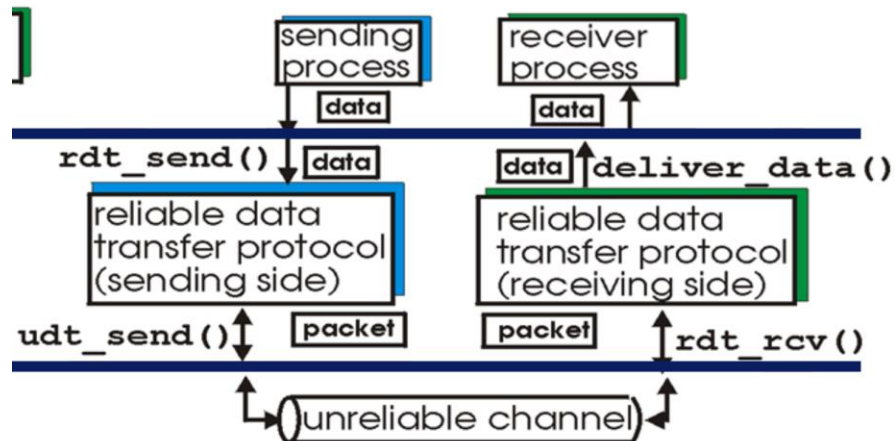
# Transport Layer

## Reliable Data Transfer 1.0

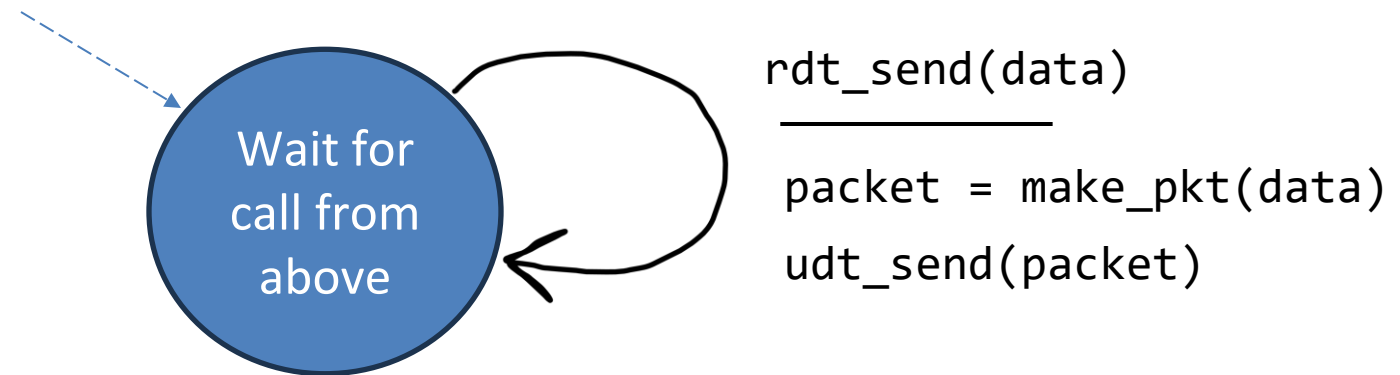
RDT 1.0

Assumptions:

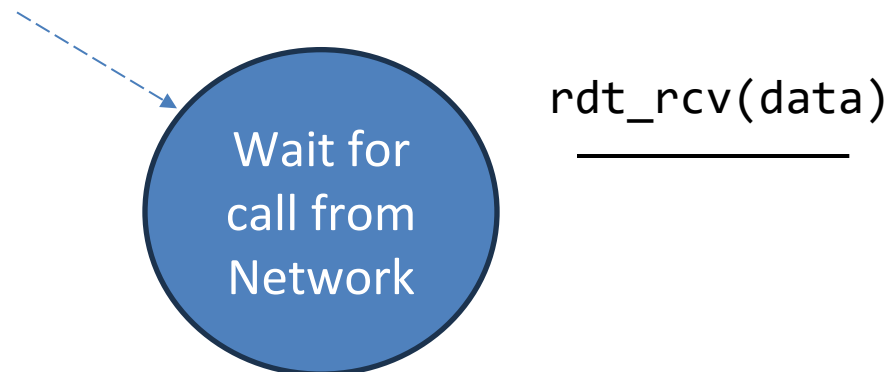
- Unidirectional long data flows
- Perfectly reliable channel
- No bit errors
- No packet loss
- No packet reordering



### Sender



### Receiver



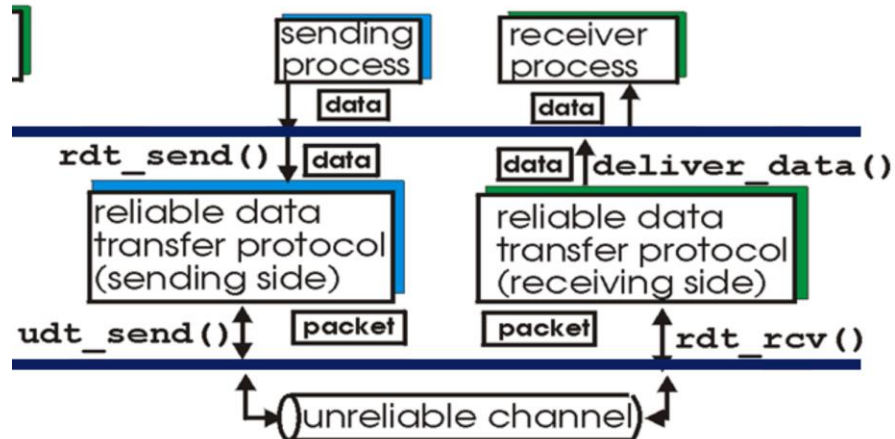
# Transport Layer

## Reliable Data Transfer 1.0

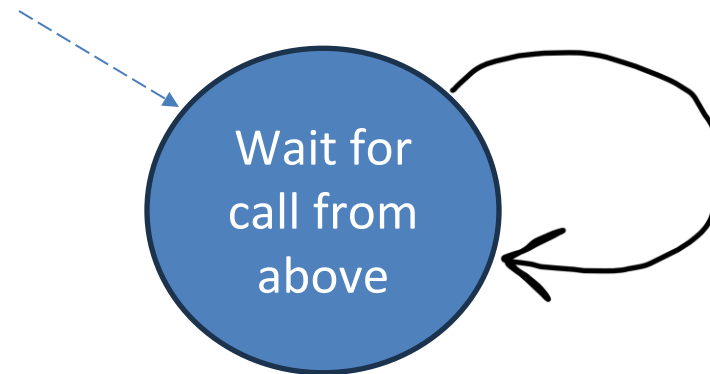
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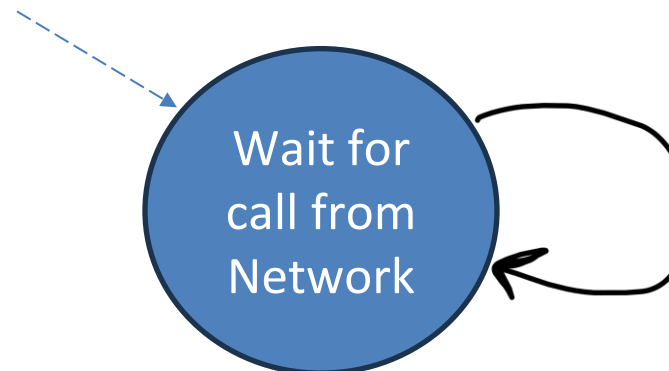


### Sender



```
rdt_send(data)
packet = make_pkt(data)
udt_send(packet)
```

### Receiver



```
rdt_rcv(data)
data = extract(packet)
deliver_data(data)
```

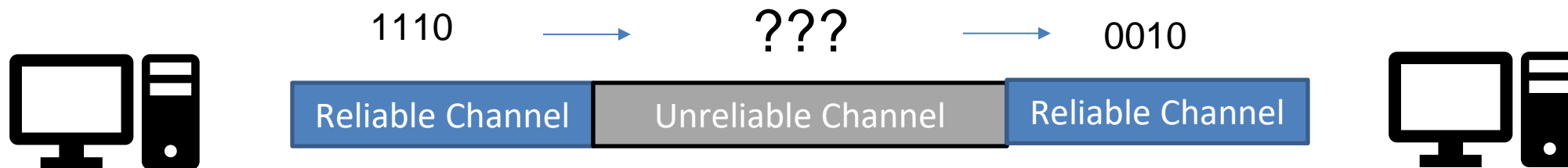
# Transport Layer

## Reliable Data Transfer 2.0

RDT 2.0

Potential for bit errors

How can we detect errors?



# Transport Layer

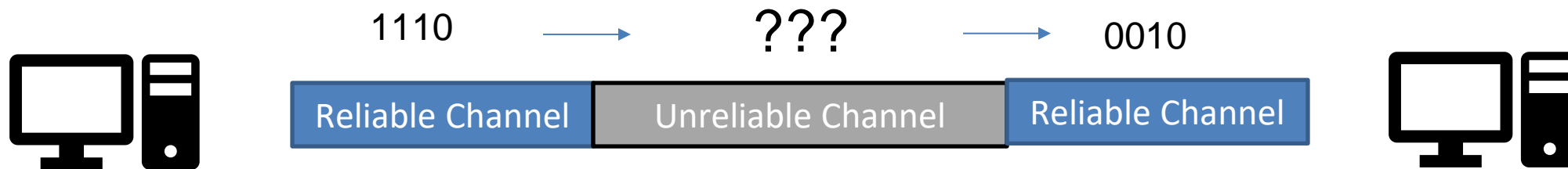
## Reliable Data Transfer 2.0

RDT 2.0

Potential for bit errors

How can we detect errors?

-Checksum



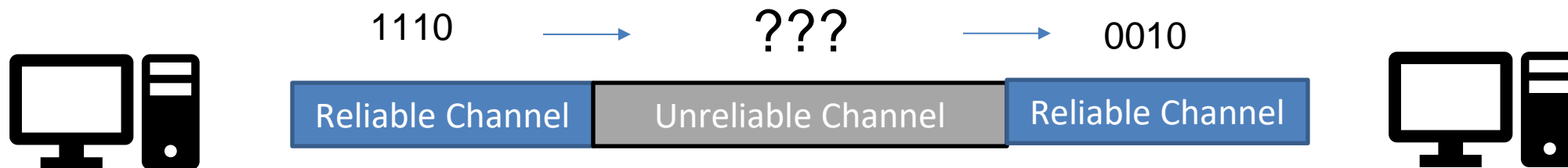
RDT 2.0

Potential for bit errors

How can we detect errors?

-Checksum

What is a good way to handle and prevent errors?



RDT 2.0

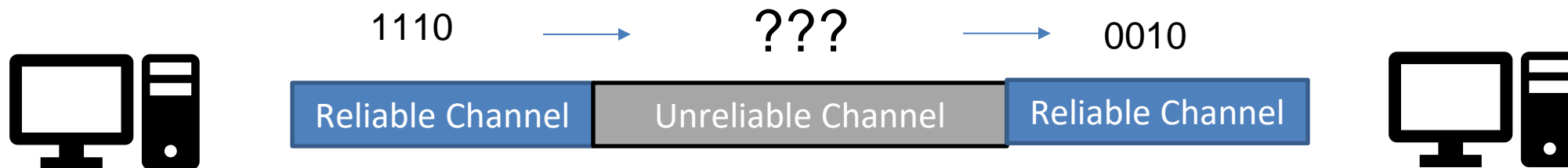
Potential for bit errors

How can we detect errors?

- Checksum

What is a good way to handle and prevent errors?

- Acknowledged packet, Ask for retransmit if needed





# Transport Layer

## Reliable Data Transfer 2.0

RDT 2.0

Potential for bit errors

How can we detect errors?

-Checksum

What is a good way to handle and prevent errors?  
- Acknowledged packet, Ask for retransmit if needed

Stop-and-wait: sender sends one packet, then waits for receiver response

Sender

`rdt_send(data)`

`sndpkt = make_pkt(data, checksum)`

`udt_send(sndpkt)`

Wait for call  
from appl

# Transport Layer

## Reliable Data Transfer 2.0

RDT 2.0

Potential for bit errors

How can we detect errors?

-Checksum

What is a good way to handle and prevent errors?  
- Acknowledged packet, Ask for retransmit if needed

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`rdt_send(data)`

`sndpkt = make_pkt(data, checksum)`

`udt_send(sndpkt)`

Wait for call  
from appl



Wait for ACK  
or NAK

## Transport Layer

# Reliable Data Transfer 2.0

Sender

rdt\_send(data)

sndpkt = make\_pkt(data, checksum)

udt\_send(sndpkt)



Receiver

Sender

rdt\_send(data)

sndpkt = make\_pkt(data, checksum)

udt\_send(sndpkt)

Wait for call  
from appl

Wait for ACK  
or NAK

Receiver

Wait for call  
from below

Sender

rdt\_send(data)

sndpkt = make\_pkt(data, checksum)

udt\_send(sndpkt)

Wait for call  
from appl

Wait for ACK  
or NAK

Receiver

Wait for call  
from below

rdt\_rcv(rcvpkt) &&  
corrupt(rcvpkt)

udt\_send(NAK)

Sender

rdt\_send(data)

sndpkt = make\_pkt(data, checksum)

udt\_send(sndpkt)

Wait for call  
from appl

Wait for ACK  
or NAK

Receiver

Wait for call  
from below

rdt\_rcv(rcvpkt) &&  
corrupt(rcvpkt)

udt\_send(NAK)

rdt\_rcv(rcvpkt) &&  
!corrupt(rcvpkt)  
data = extract(packet)  
deliver\_data(data)  
udt\_send(ACK)

### Sender

rdt\_send(data)

sndpkt = make\_pkt(data, checksum)

udt\_send(sndpkt)

Wait for call  
from appl

Wait for ACK  
or NAK

### Receiver

rdt\_rcv(rcvpkt) &&  
corrupt(rcvpkt)

udt\_send(NAK)

rdt\_rcv(rcvpkt) &&  
!corrupt(rcvpkt)

data = extract(packet)  
deliver\_data(data)  
udt\_send(ACK)

Wait for call  
from below

# Transport Layer

## Reliable Data Transfer 2.0

Sender

rdt\_send(data)

sndpkt = make\_pkt(data, checksum)

udt\_send(sndpkt)

rdt\_rcv(rcvpkt) &&

isNAK(rcvpkt)

udt\_send(sndpkt)

Wait for call  
from appl

Wait for ACK  
or NAK

Receiver

rdt\_rcv(rcvpkt) &&

corrupt(rcvpkt)

udt\_send(NAK)

rdt\_rcv(rcvpkt) &&

!corrupt(rcvpkt)

data = extract(packet)

deliver\_data(data)

udt\_send(ACK)

Wait for call  
from below



### Sender

rdt\_send(data)

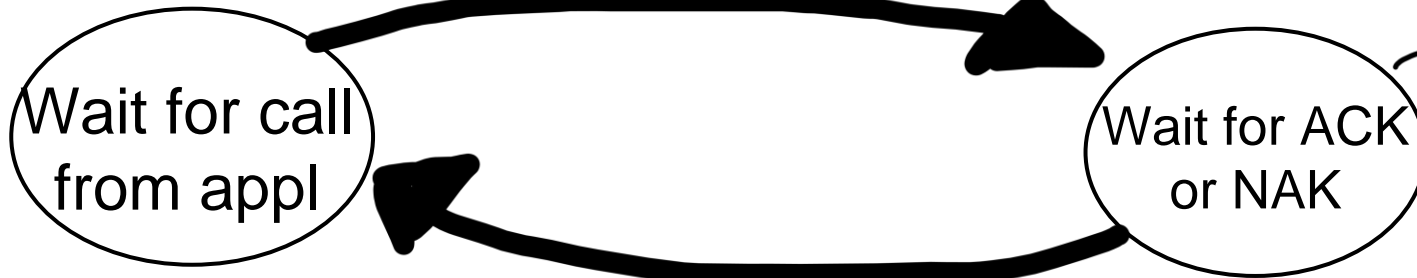
sndpkt = make\_pkt(data, checksum)

udt\_send(sndpkt)

rdt\_rcv(rcvpkt) &&

isNAK(rcvpkt)

udt\_send(sndpkt)



rdt\_rcv(rcvpkt) && isACK(rcvpkt)

$\Lambda$

### Receiver

rdt\_rcv(rcvpkt) &&

corrupt(rcvpkt)

udt\_send(NAK)

rdt\_rcv(rcvpkt) &&

!corrupt(rcvpkt)

data = extract(packet)

deliver\_data(data)

udt\_send(ACK)

$\Lambda$  = do no action, follow arrow to next state

Sender

What happens if ACK/NAK  
Corrupted?

→ Duplicate delivery, or no  
retransmission

Solution?

→ Retransmit if CORRUPT packet  
received

How to deal we duplicate Packets?

→ ???

Wait for  
from ap

Λ = do no a

Receiver

```
rdt_rcv(rcvpkt) &&  
corrupt(rcvpkt)
```

```
udt_send(NAK)
```

```
rdt_rcv(rcvpkt) &&  
!corrupt(rcvpkt)
```

```
data = extract(packet)  
deliver_data(data)  
udt_send(ACK)
```

call  
flow

Sender

What happens if ACK/NAK  
Corrupted?

→ Duplicate delivery, or no  
retransmission

Solution?

→ Retransmit if CORRUPT packet  
received

How to deal we duplicate Packets?

→ **Sequence Numbers**

Wait for  
from ap

Λ = do no a

Receiver

```
rdt_rcv(rcvpkt) &&  
corrupt(rcvpkt)
```

```
udt_send(NAK)
```

```
rdt_rcv(rcvpkt) &&  
!corrupt(rcvpkt)
```

```
data = extract(packet)  
deliver_data(data)  
udt_send(ACK)
```

call  
flow

## Transport Layer

# Reliable Data Transfer 2.1

Potential for bit errors and garbled ACKs

(We only need to use 0 or 1 for the sequence number)

Sender

Wait for  
call 0 from  
above

## Transport Layer

# Reliable Data Transfer 2.1

Potential for bit errors and garbled ACKs

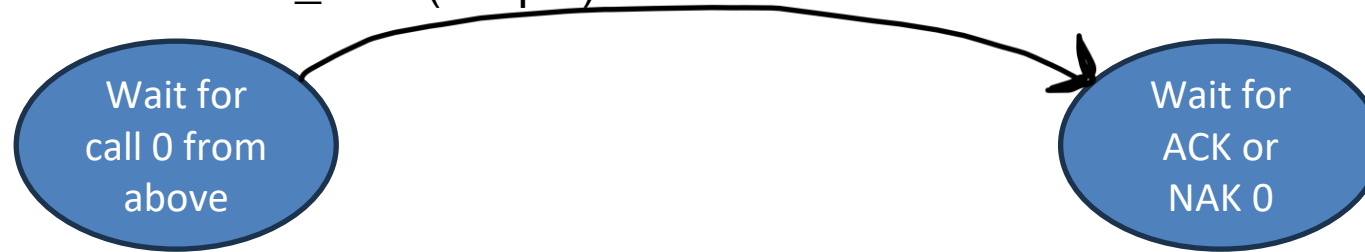
(We only need to use 0 or 1 for the sequence number)

Sender

`rdt_send(data)`

`sndpkt = make_pkt(0, data, checksum)`

`udt_send(sndpkt)`



# Transport Layer

## Reliable Data Transfer 2.1

Potential for bit errors and garbled ACKs

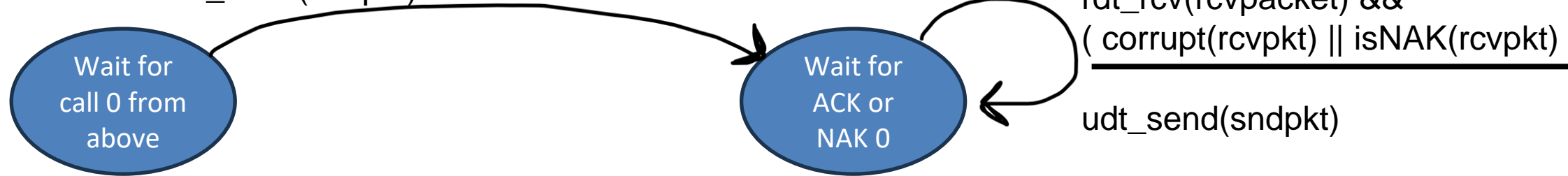
(We only need to use 0 or 1 for the sequence number)

Sender

rdt\_send(data)

sndpkt = make\_pkt(0, data, checksum)

udt\_send(sndpkt)



# Transport Layer

## Reliable Data Transfer 2.1

Potential for bit errors and garbled ACKs

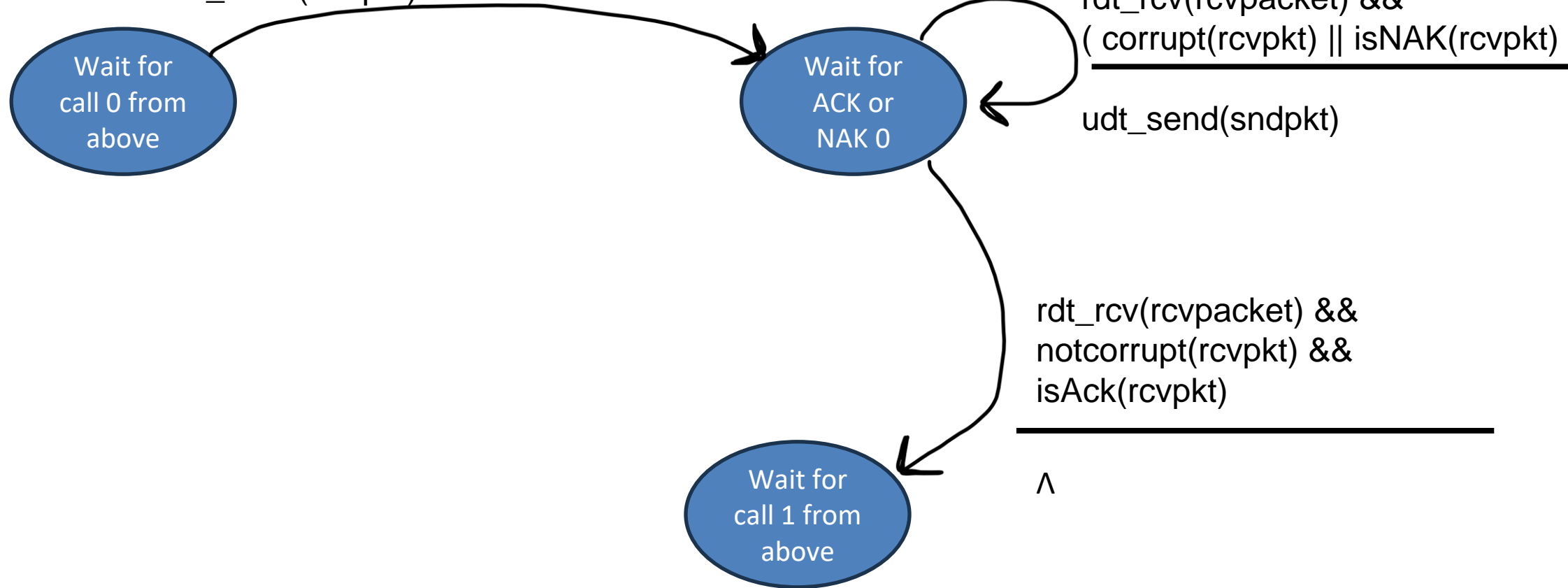
(We only need to use 0 or 1 for the sequence number)

Sender

rdt\_send(data)

sndpkt = make\_pkt(0, data, checksum)

udt\_send(sndpkt)



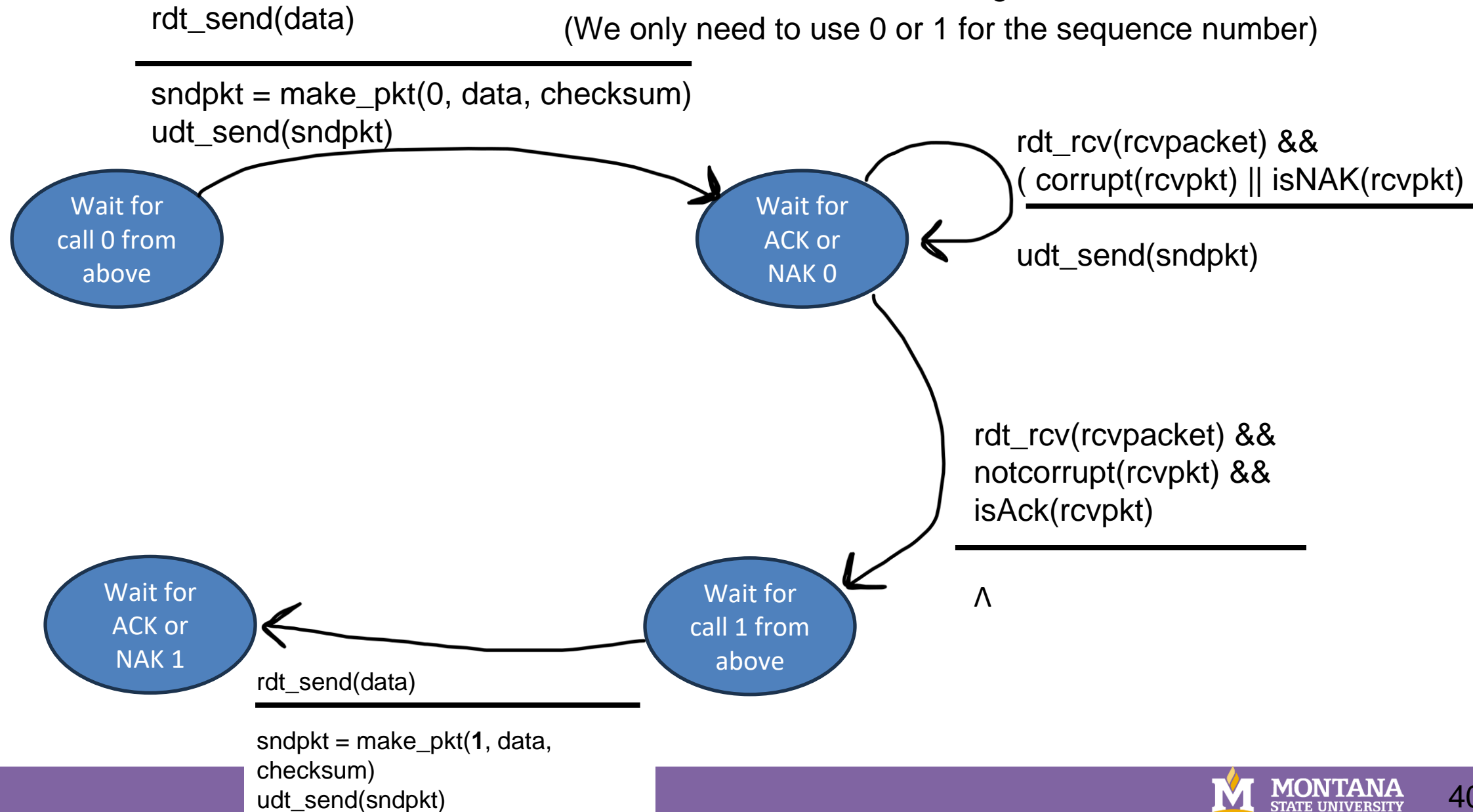
# Transport Layer

## Reliable Data Transfer 2.1

Potential for bit errors and garbled ACKs

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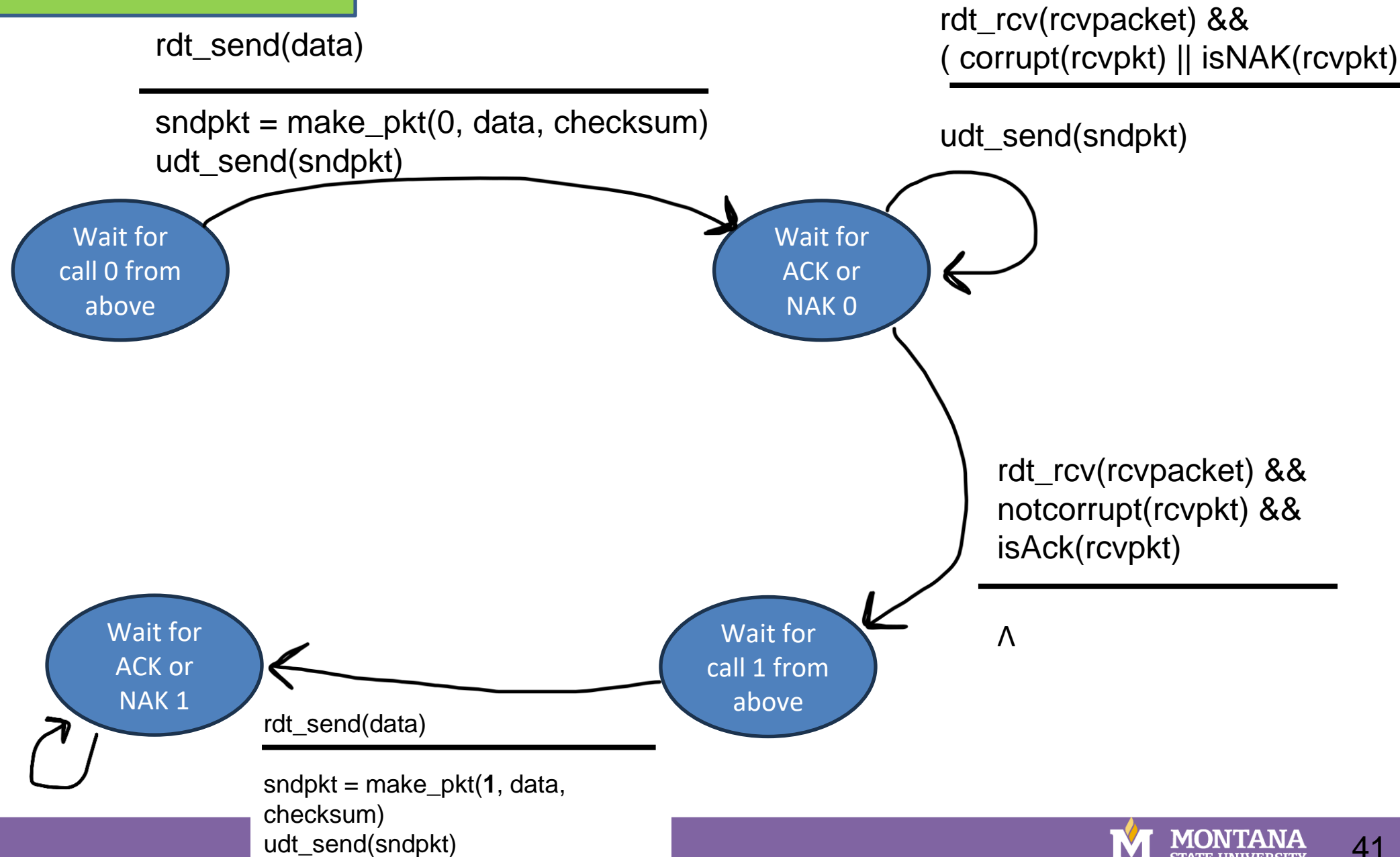
Sender





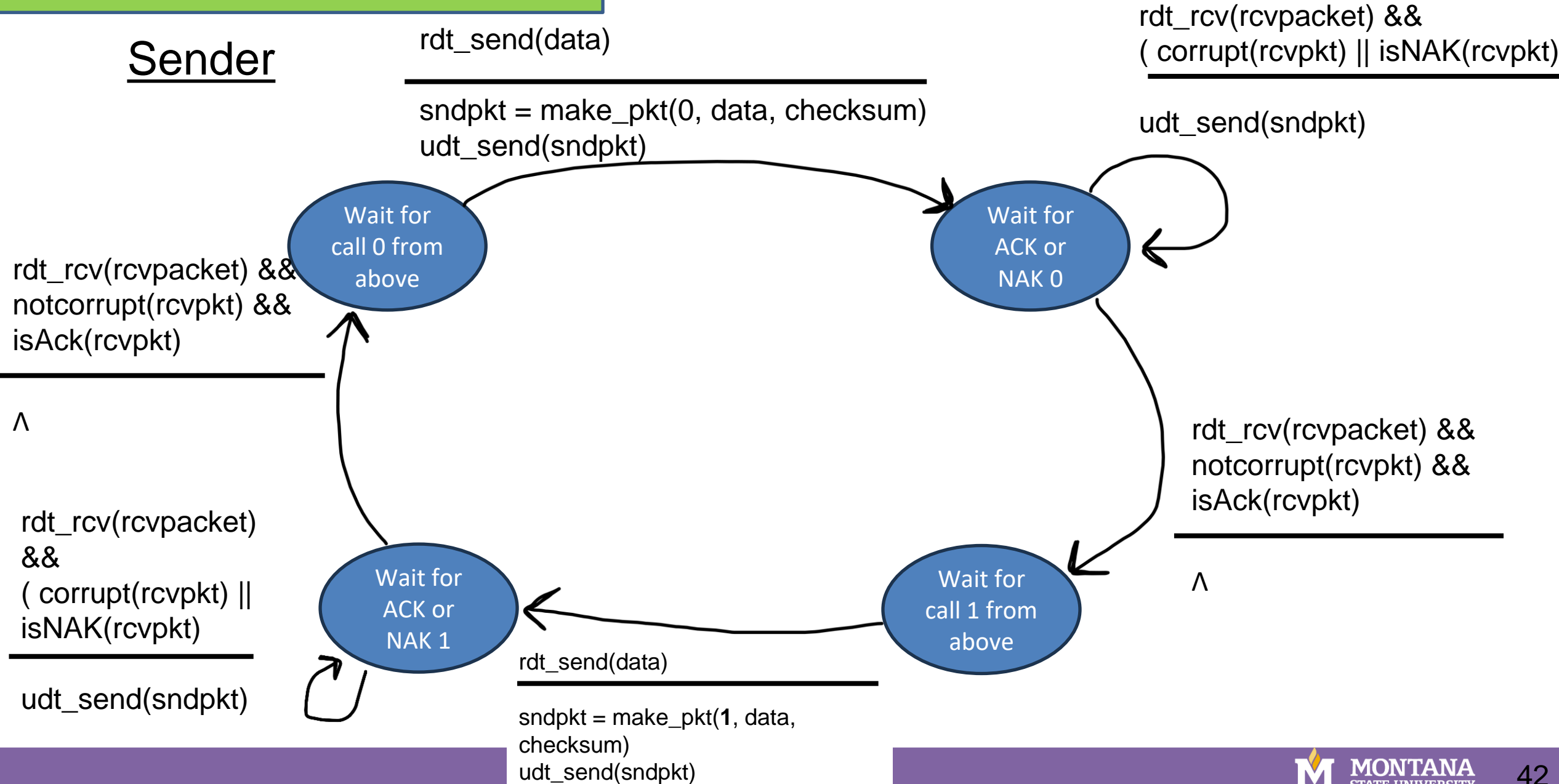
# Transport Layer

## Sender



# Transport Layer

## Sender



# Transport Layer

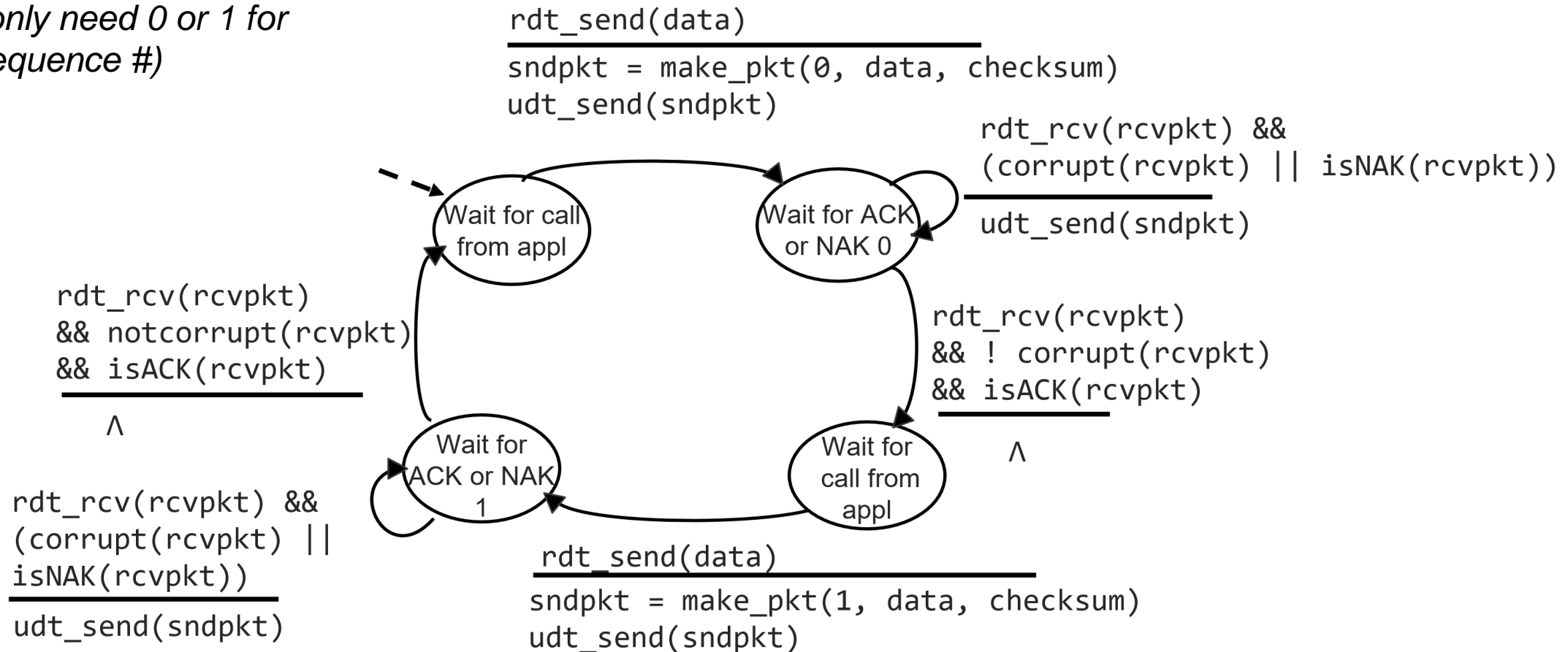
## Reliable Data Transfer 2.1

### RDT 2.1

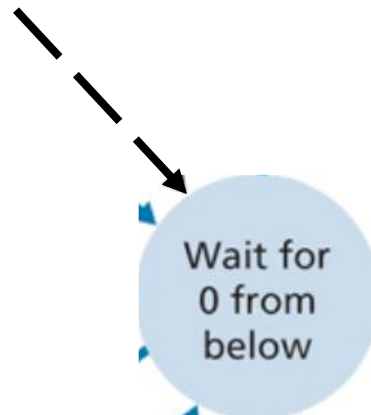
Potential for bit errors and garbled ACKs

Stop-and-wait: sender sends one packet, then waits for receiver response

*(We only need 0 or 1 for the sequence #)*

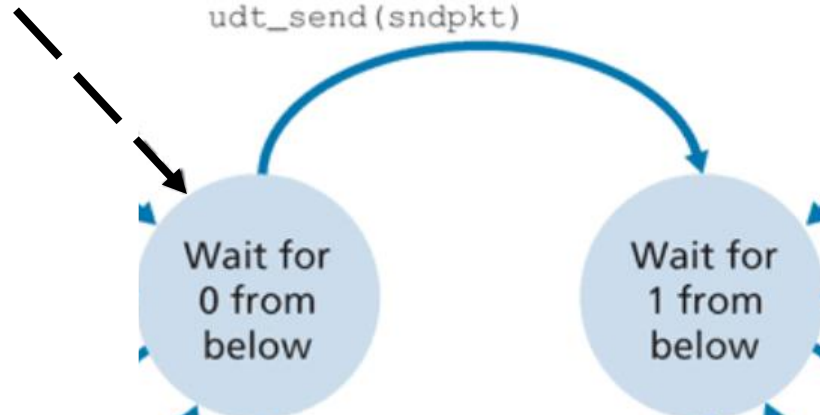


Receiver

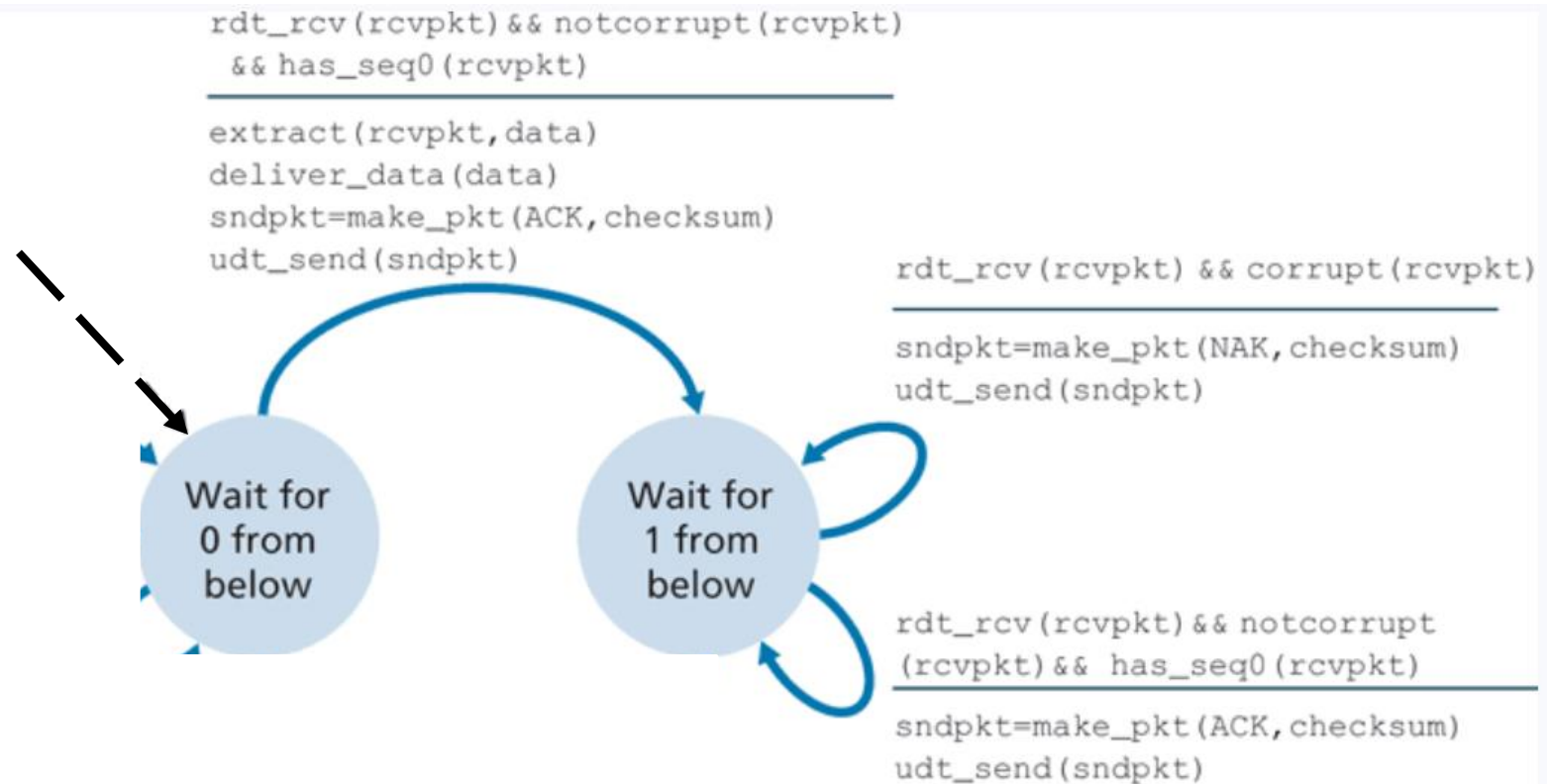


### Receiver

```
rdt_rcv(rcvpkt) && notcorrupt(rcvpkt)
&& has_seq0(rcvpkt)
-----
extract(rcvpkt, data)
deliver_data(data)
sndpkt=make_pkt(ACK, checksum)
udt_send(sndpkt)
```



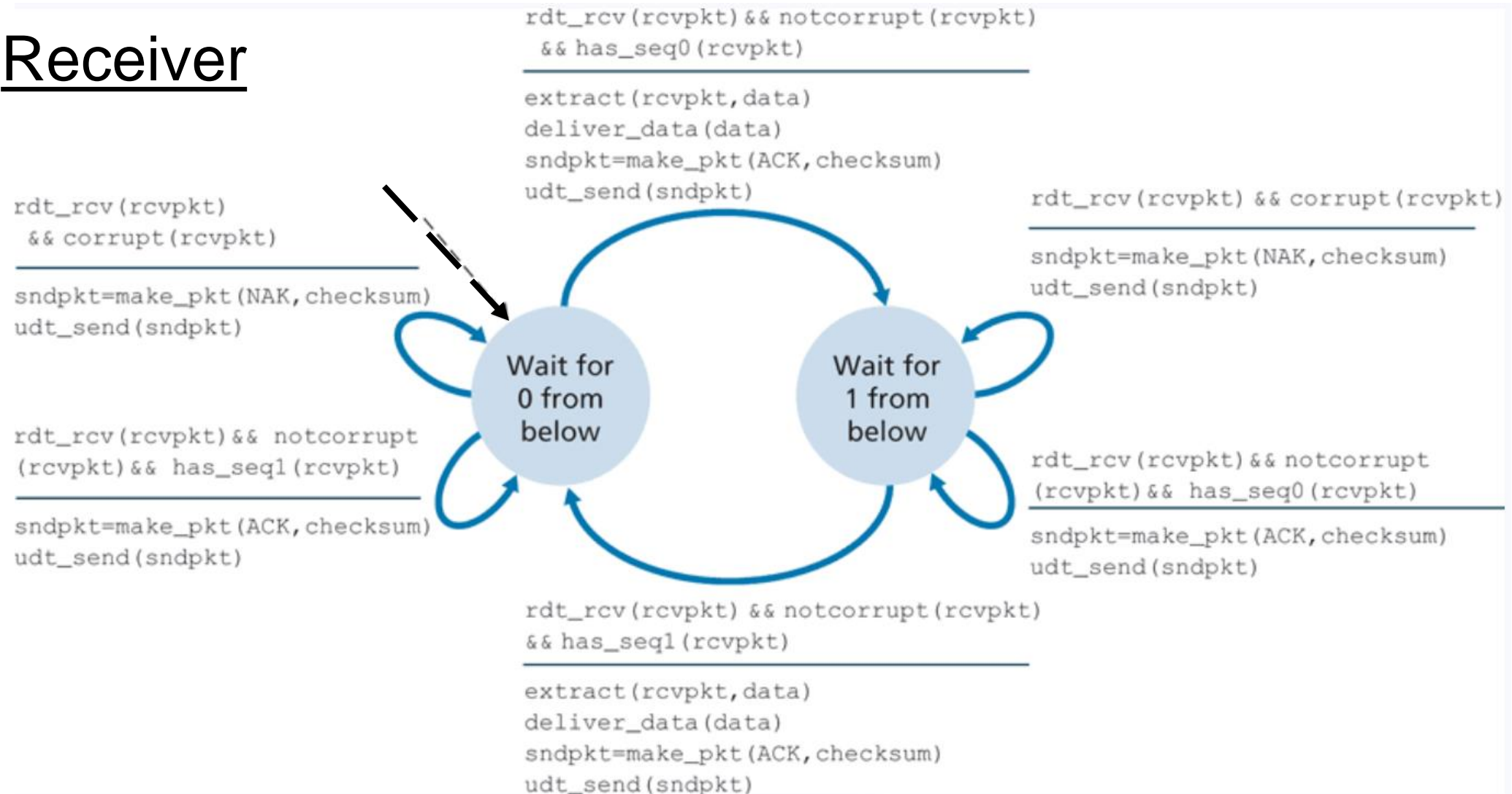
### Receiver



### Receiver



### Receiver

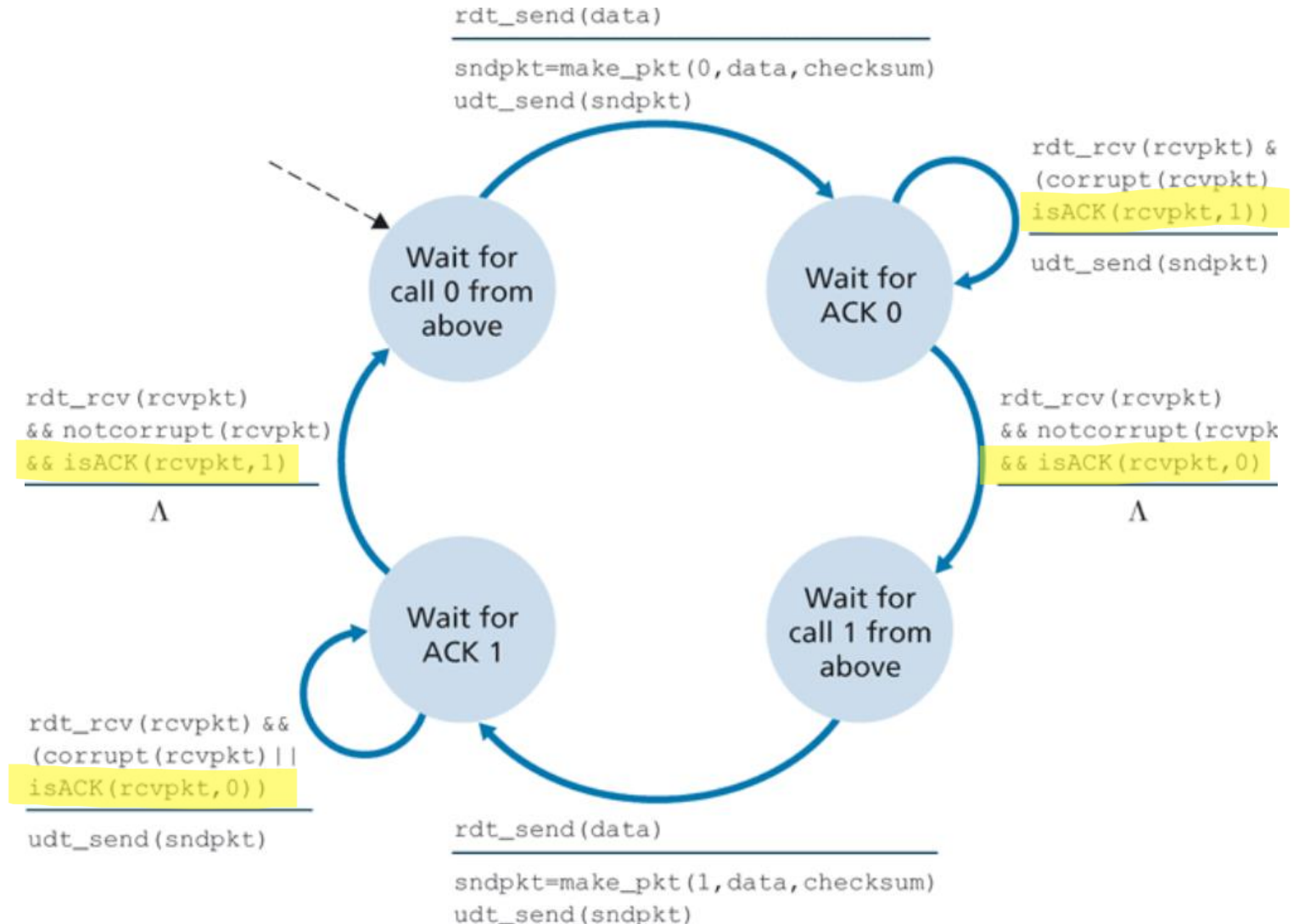




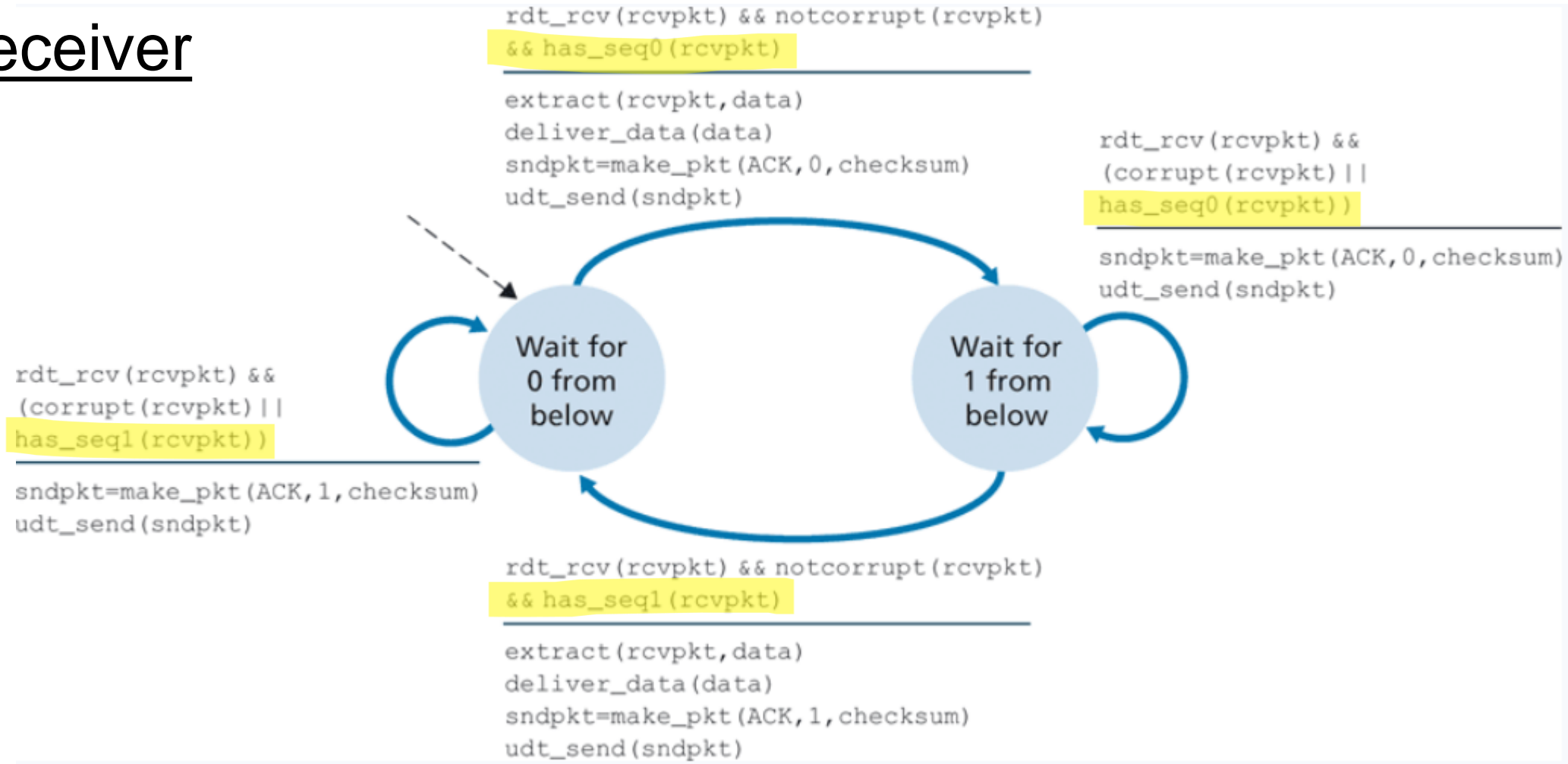
### Sender

Instead of NAKs, we provide the sequence number of the packet for our ACK (ACK1, ACK0)

Duplicate ACK at sender results in same action as NAK: retransmit current pkt



### Receiver



# Transport Layer



A quick meatball break...

We will still need  
checksums, Seq #'s,  
ACKS, but we need more

Potential for:

- bit errors
- garbled ACKs
- **Lost/dropped packets**

What if an ACK gets dropped? Sender is stuck waiting forever

# Reliable Data Transfer 3.0

We will still need  
checksums, Seq #'s,  
ACKS, but we need more

Potential for:

- bit errors
- garbled ACKs
- **Lost/dropped packets**

What if an ACK gets dropped? Sender is stuck waiting forever

**Solution?**

We will still need  
checksums, Seq #'s,  
ACKS, but we need more

Potential for:

- bit errors
- garbled ACKs
- **Lost/dropped packets**

What if an ACK gets dropped? Sender is stuck waiting forever

**A timer**

Retransmit if ACK received in X amount of time

**What if the ACK is just taking a really long time to arrive?**

We will still need  
checksums, Seq #'s,  
ACKS, but we need more

Potential for:

- bit errors
- garbled ACKs
- **Lost/dropped packets**

What if an ACK gets dropped? Sender is stuck waiting forever

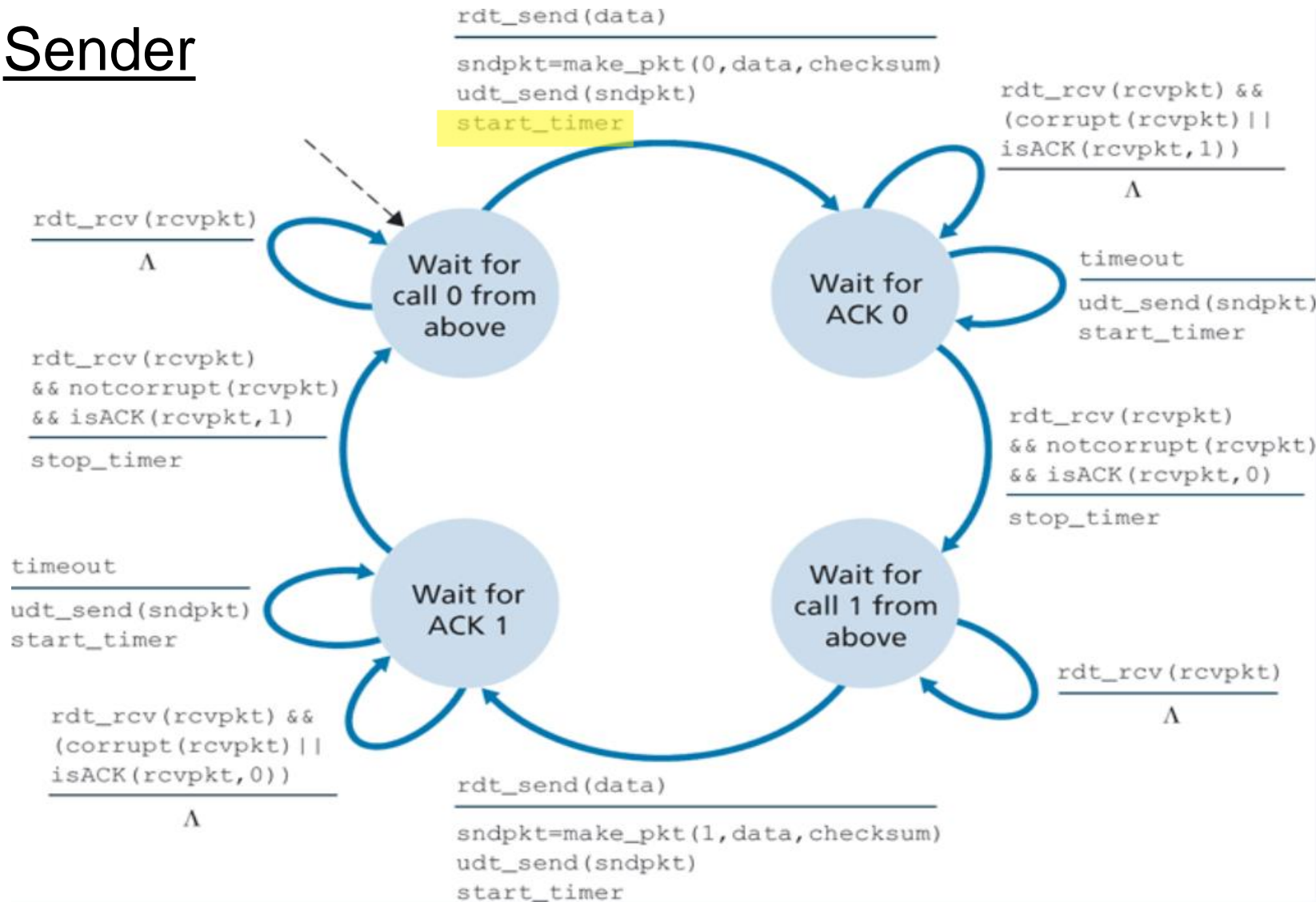
**A timer**

Retransmit if ACK received in X amount of time

What if the ACK is just taking a really long time to arrive? This will be duplicate data, but we have a Seq # to handle that

(Packets can get dropped or lost)

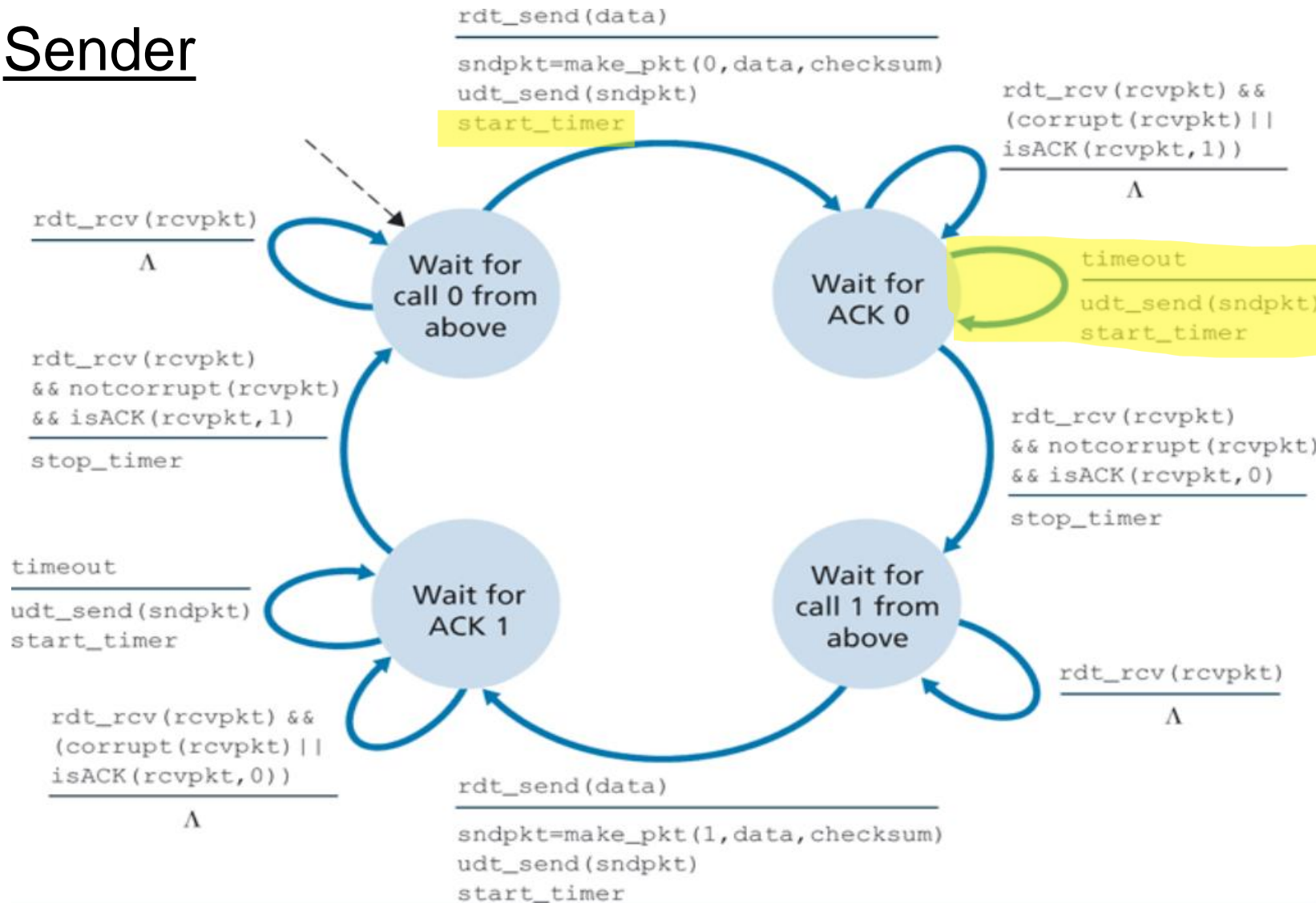
### Sender





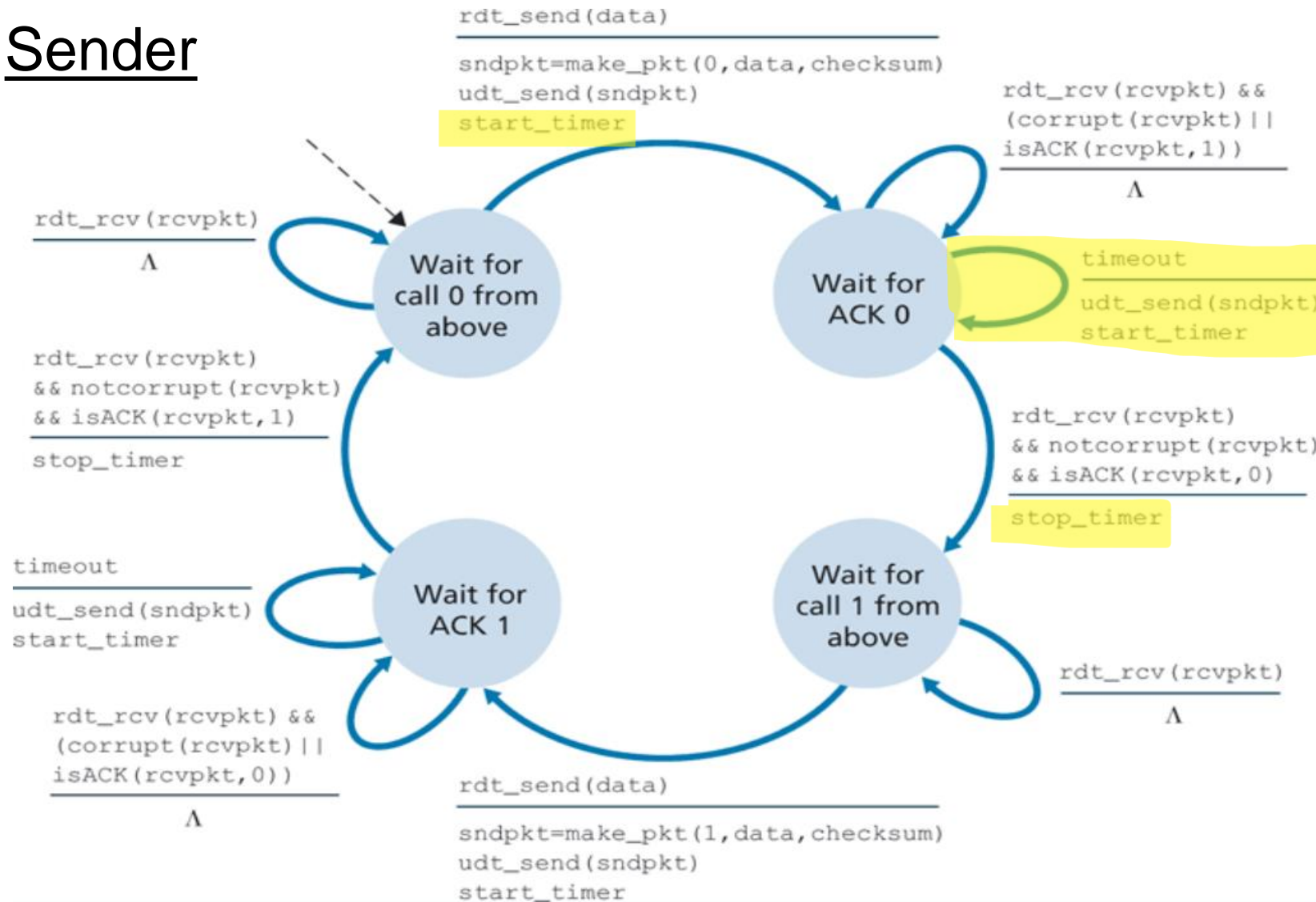
(Packets can get dropped or lost)

### Sender



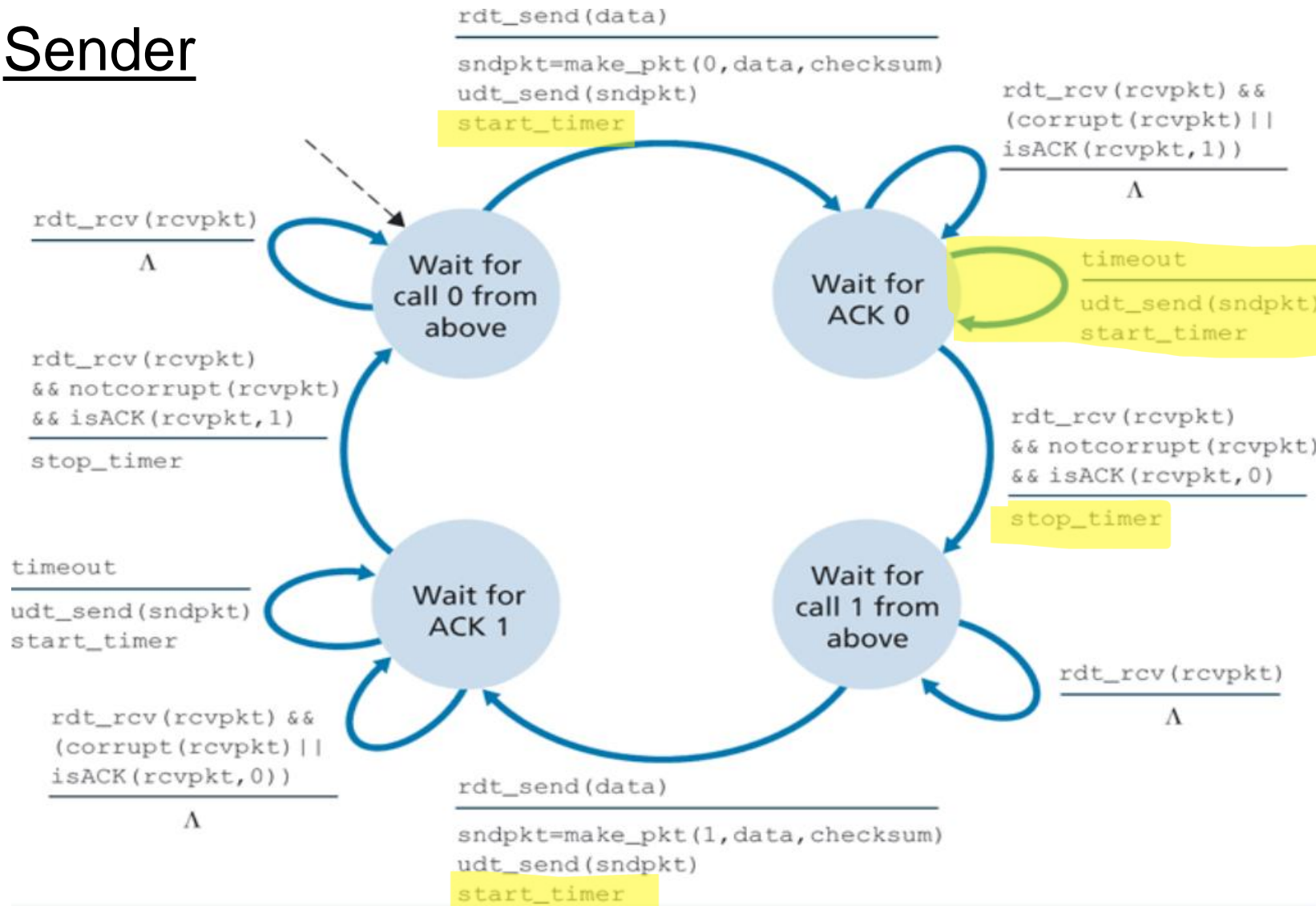
(Packets can get dropped or lost)

### Sender



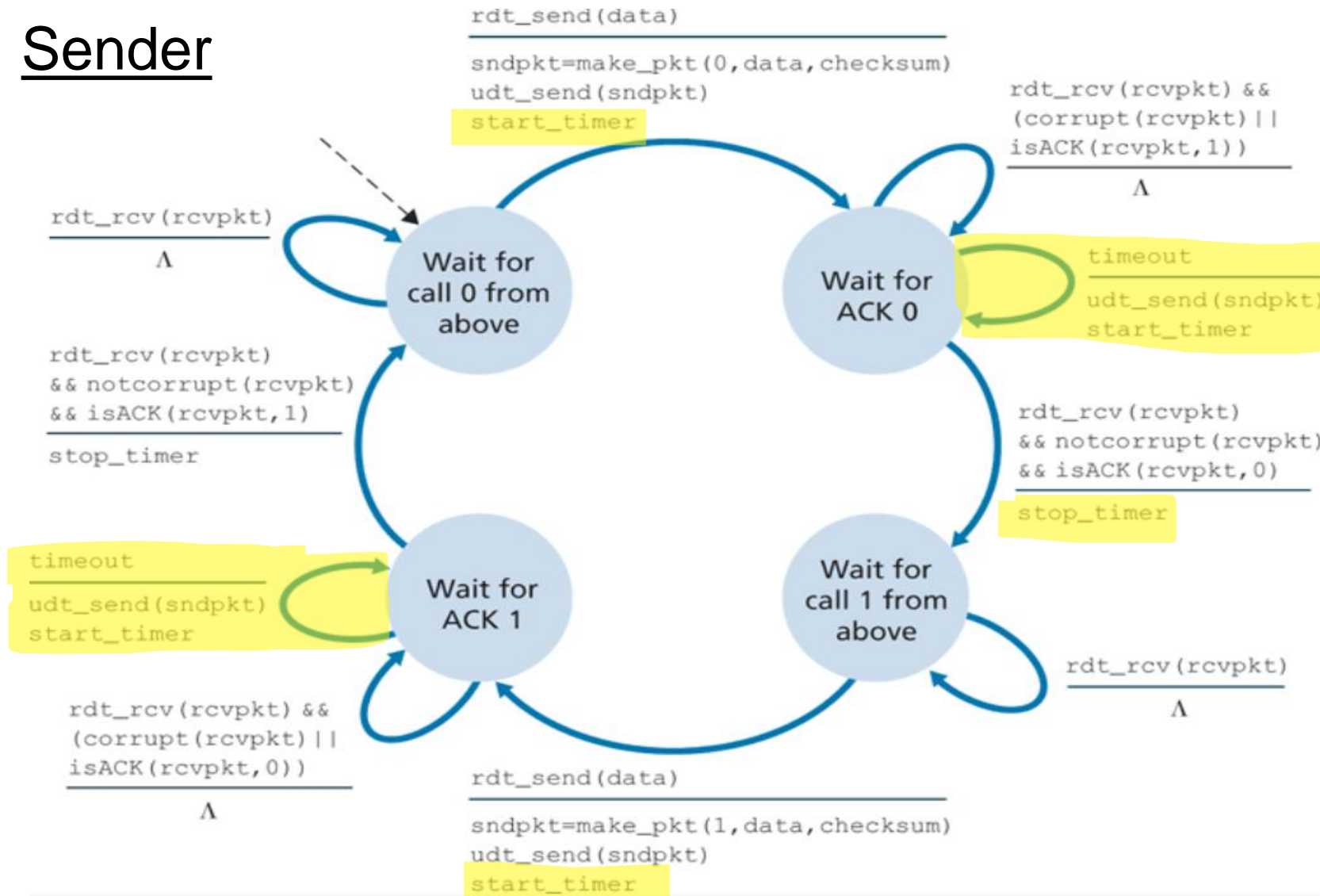
(Packets can get dropped or lost)

### Sender



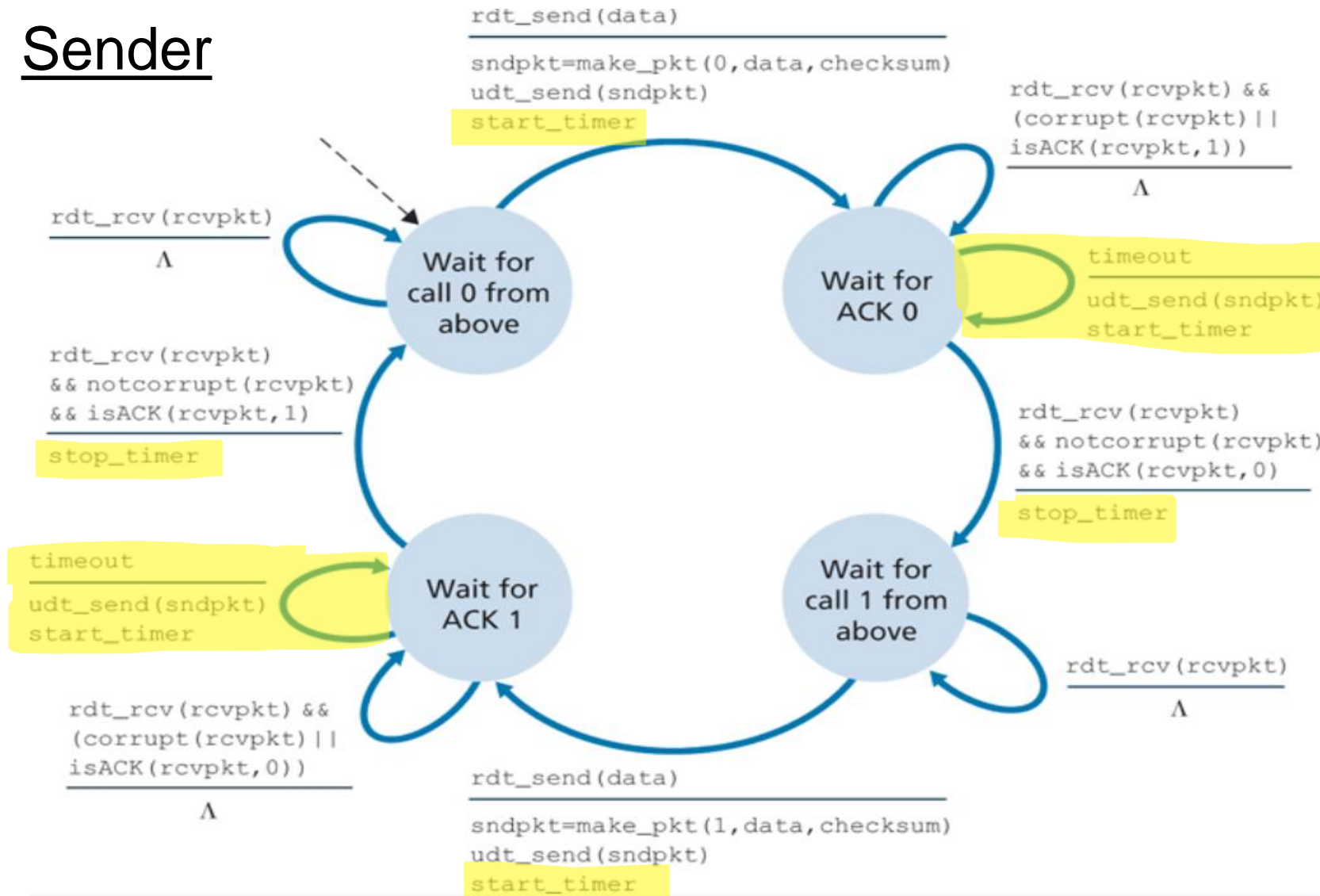
(Packets can get dropped or lost)

### Sender



(Packets can get dropped or lost)

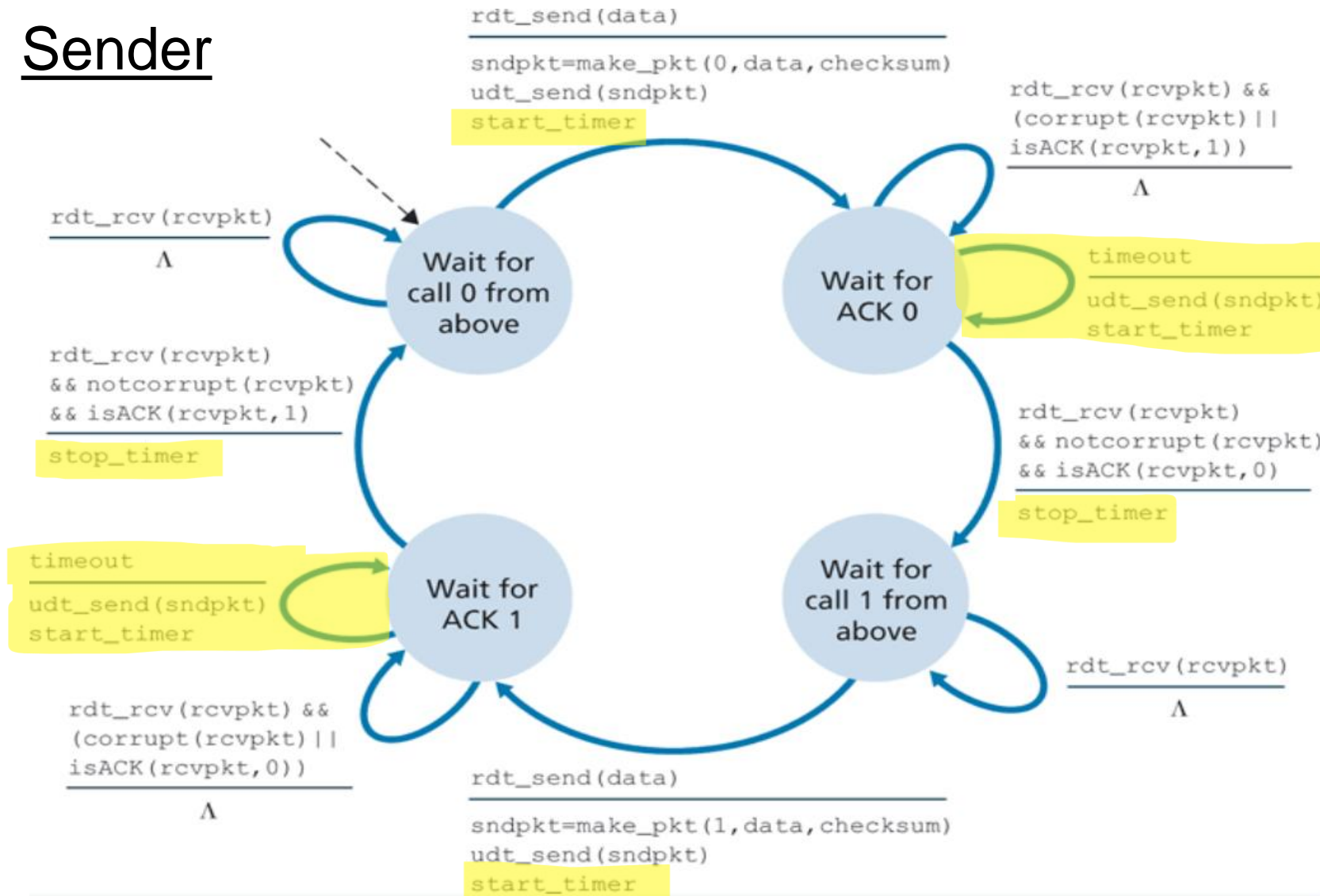
### Sender





(Packets can get dropped or lost)

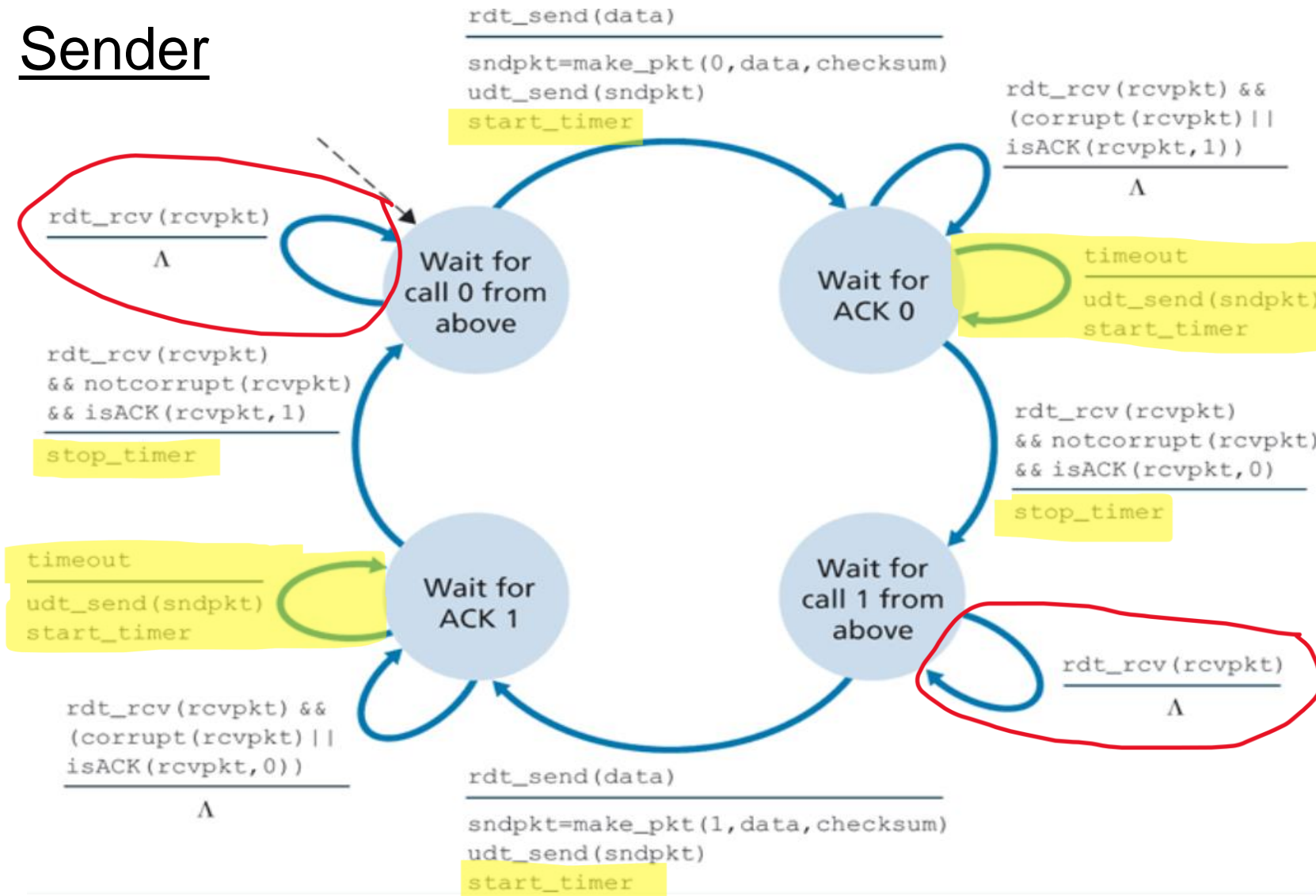
### Sender



*Receiver will be same as 2.2*

(Packets can get dropped or lost)

### Sender

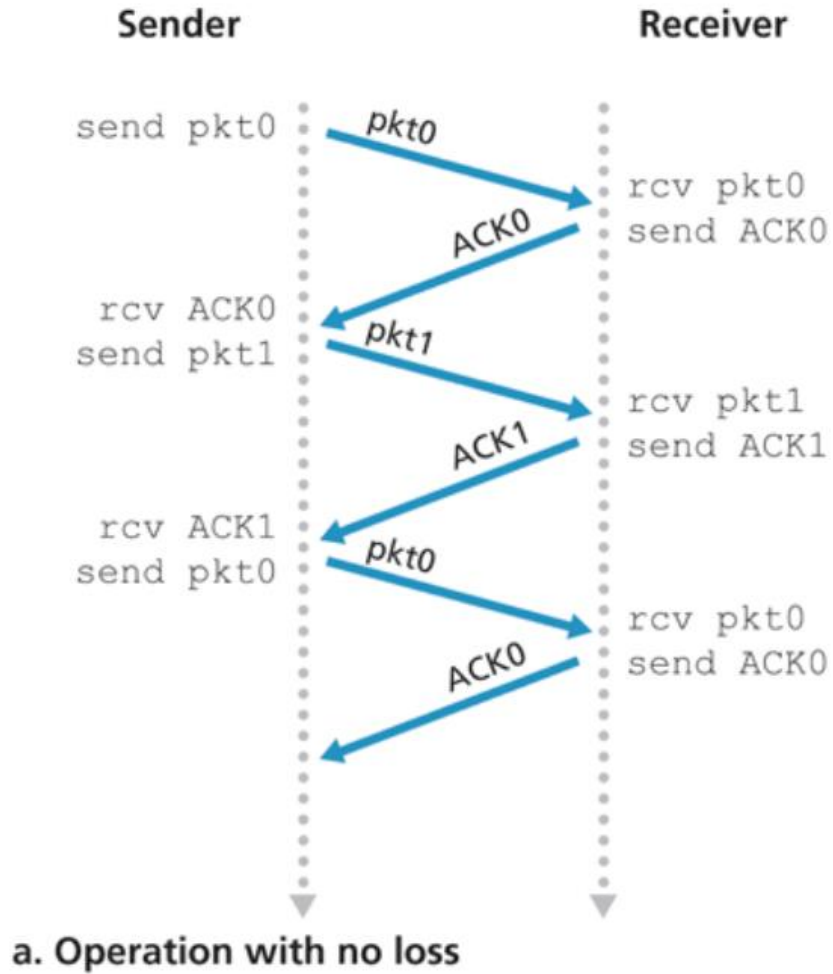


If we are waiting to send packet, but we receive a packet for some reason, this packet must be a duplicate packet, so do nothing

# Transport Layer

## Reliable Data Transfer 3.0

(Packets can get dropped or lost)

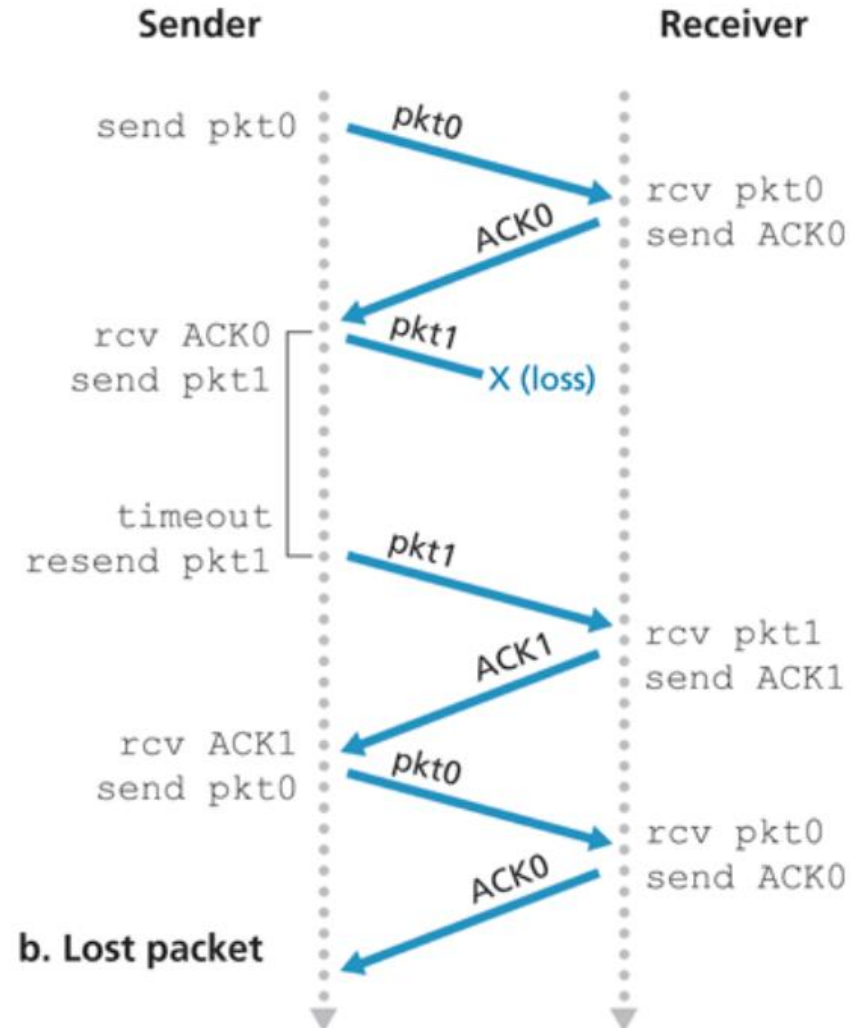
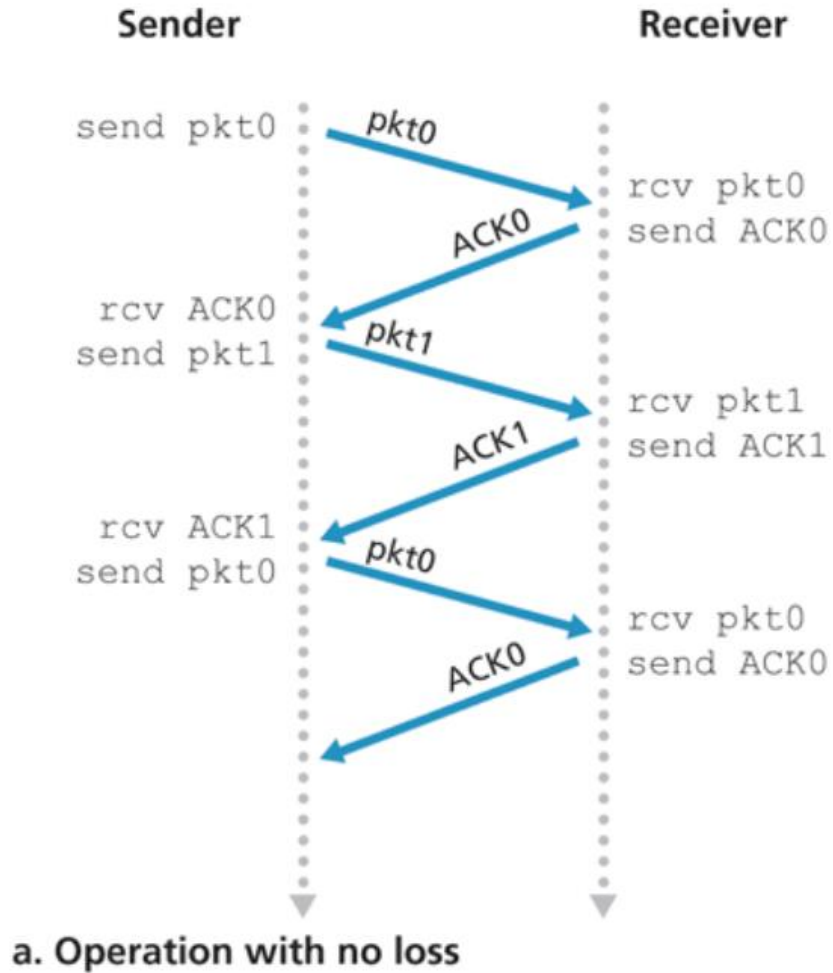




# Transport Layer

## Reliable Data Transfer 3.0

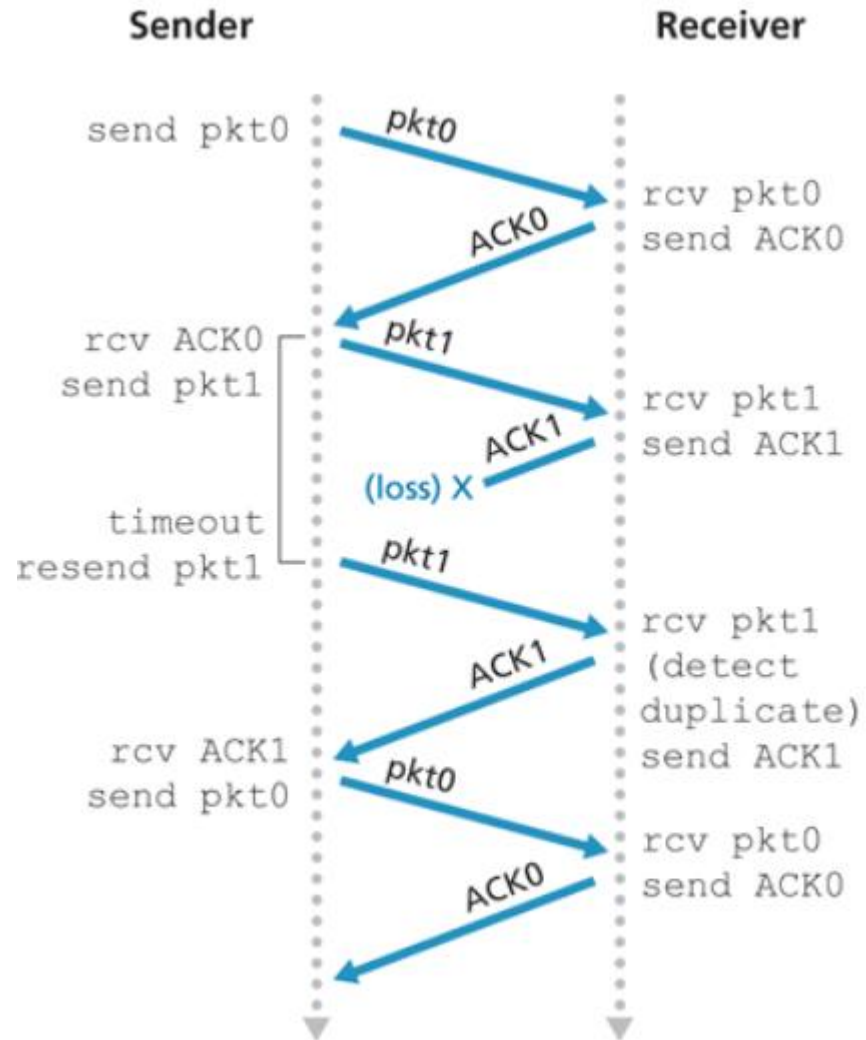
(Packets can get dropped or lost)



# Transport Layer

## Reliable Data Transfer 3.0

(Packets can get dropped or lost)

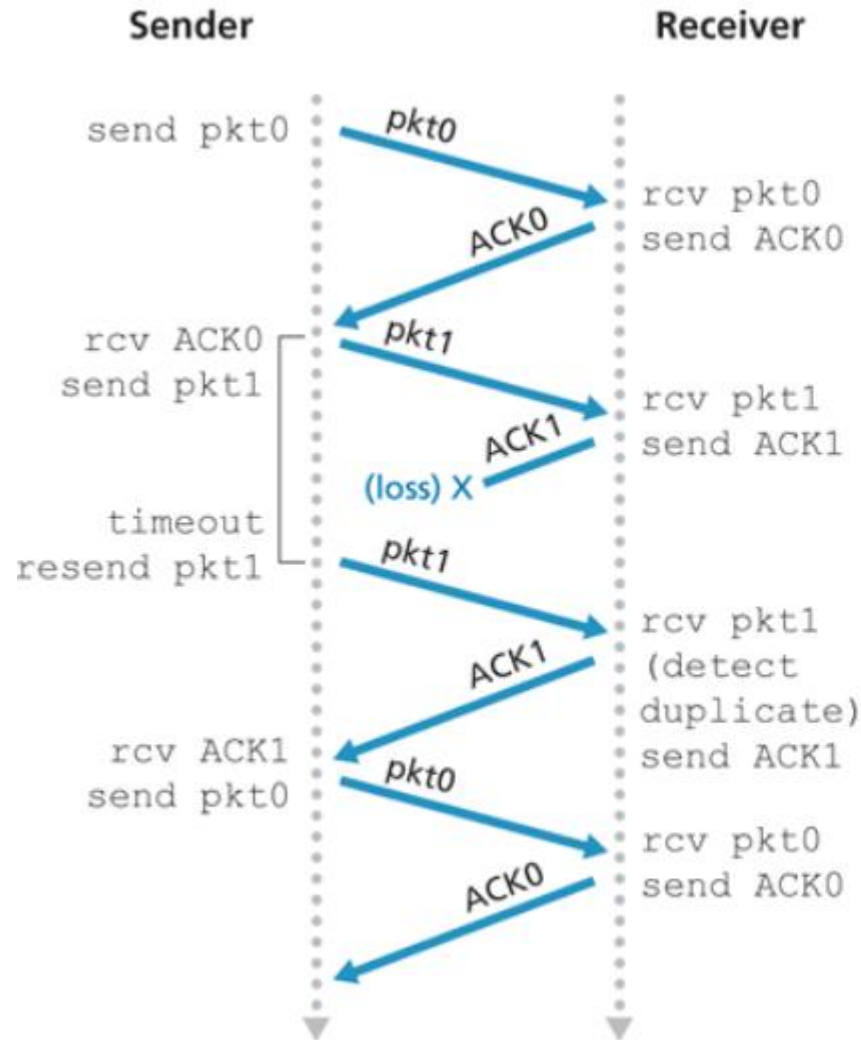


c. Lost ACK

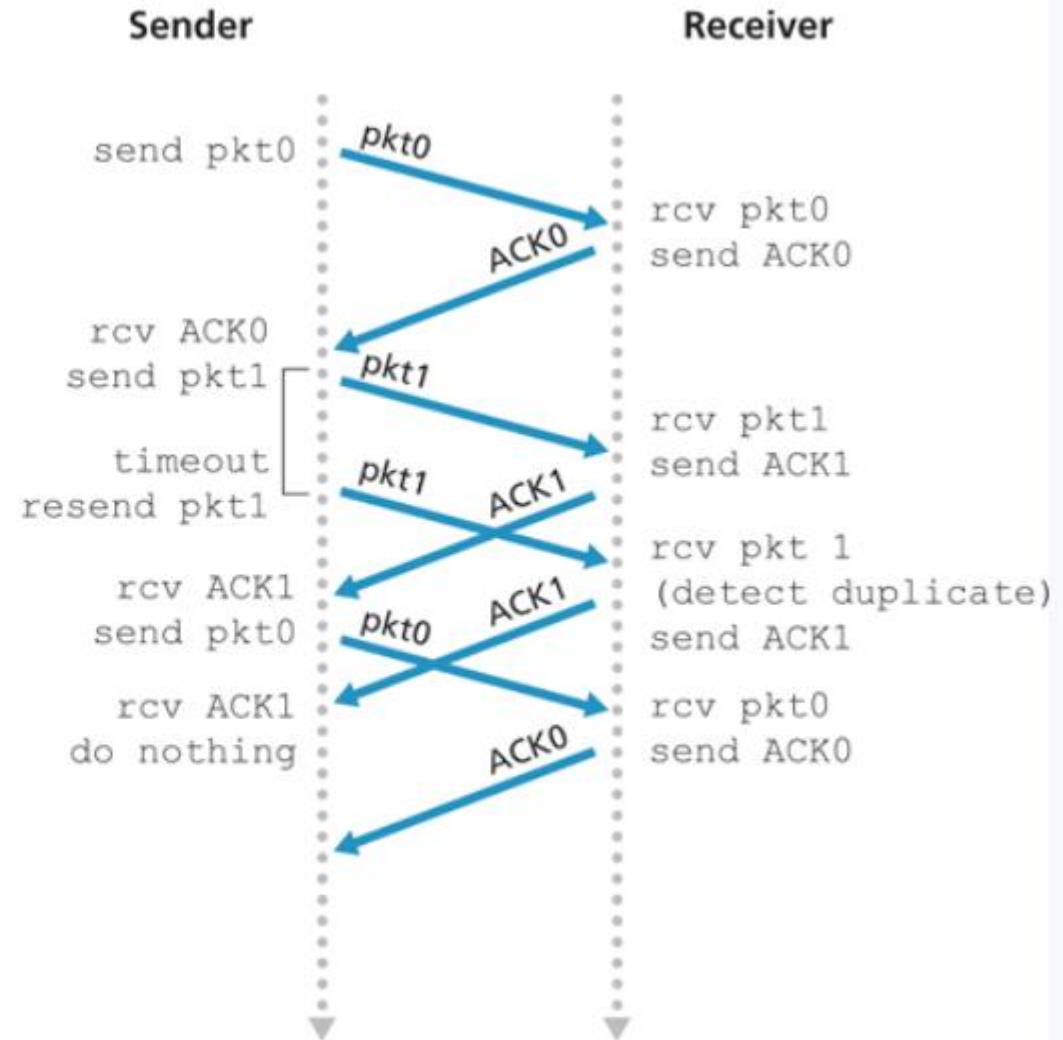
# Transport Layer

## Reliable Data Transfer 3.0

(Packets can get dropped or lost)



c. Lost ACK



d. Premature timeout