

CSCI 132:

Basic Data Structures and Algorithms

Recursion (Part 1)

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Spring 2023

Announcements

Program 4 Due April 19th

INTERVIEWER: SORT THE ARRAY

OTHER PEOPLE

```
● ● ●  
arr = [5, 2, 8, 7, 1]  
temp = 0  
  
for i in range(0, len(arr)):  
    for j in range(i+1, len(arr)):  
        if(arr[i] > arr[j]):  
            temp = arr[i]  
            arr[i] = arr[j]  
            arr[j] = temp  
  
print("Array sorted in ascending order: ")  
  
for i in range(0, len(arr)):  
    print(arr[i], end=" ")
```



ME

```
● ● ●  
arr = [5, 2, 8, 7, 1]  
arr.sort()  
print(arr)
```



Recursion is a problem-solving technique that involves a method calling itself to solve some smaller problem

```
static int factorial(int n)
{
    if (n == 0)
        return 1;

    return n * factorial(n - 1);
}
```

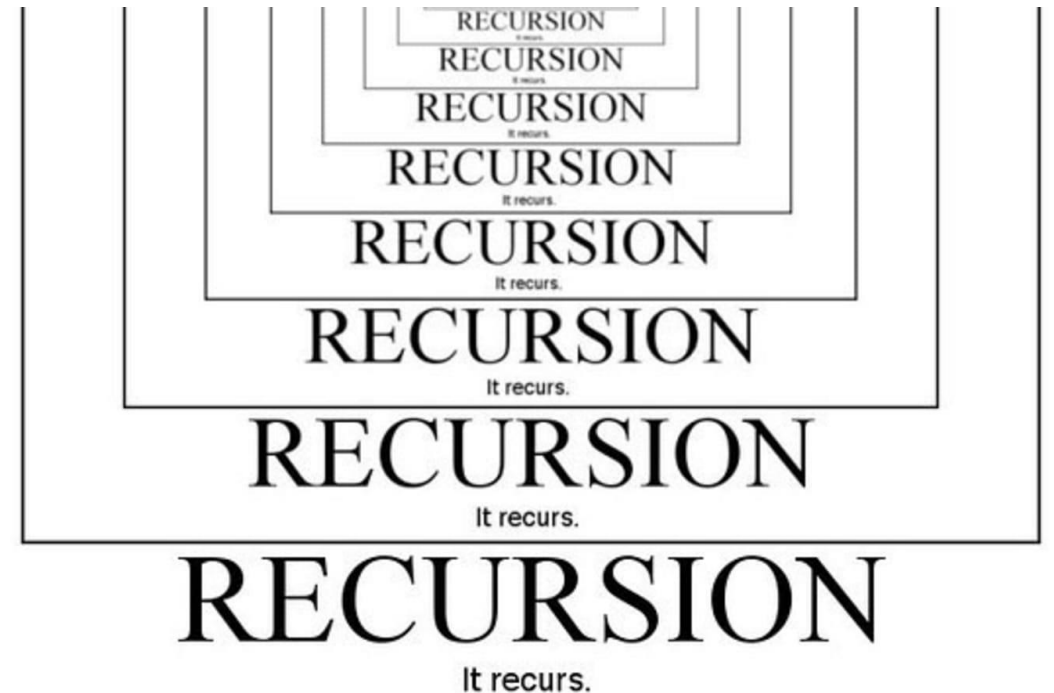
TOP DEFINITION

recursion

See recursion.

by [Anonymous](#) December 05, 2002

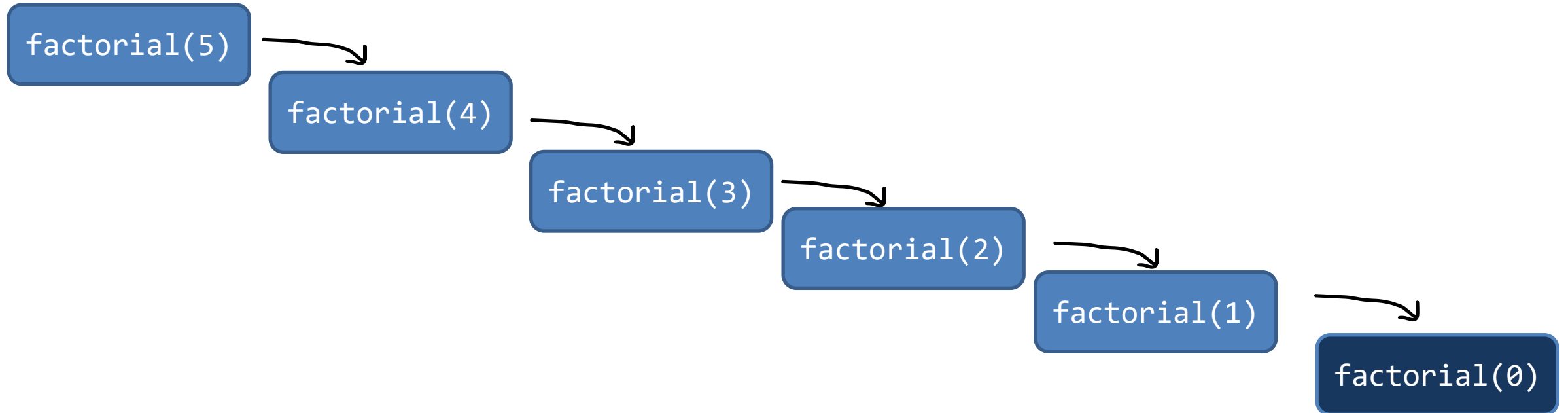
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```
static int factorial(int n)
{
    if (n == 0)
        return 1;

    return n * factorial(n - 1);
}
```

We can solve the factorial for n by solving smaller problems (factorial of $n-1$) !



```
static int factorial(int n)
```

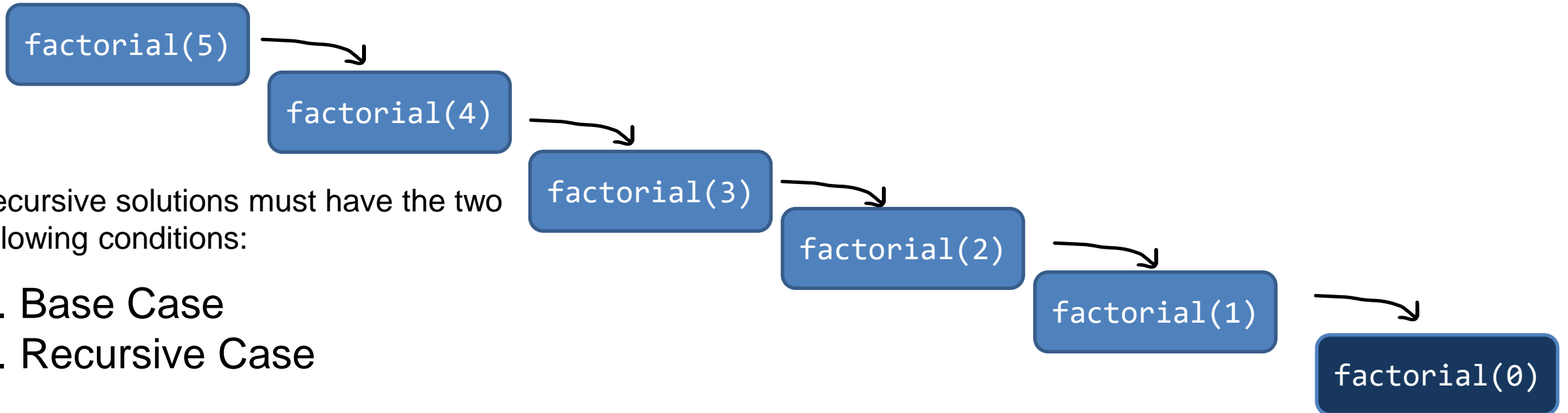
```
{
```

```
    if (n == 0)           (base case)  
        return 1;
```

```
    return n * factorial(n - 1); (recursive case)
```

```
}
```

We can solve the factorial for n by solving smaller problems (factorial of $n-1$) !



Recursive solutions must have the two following conditions:

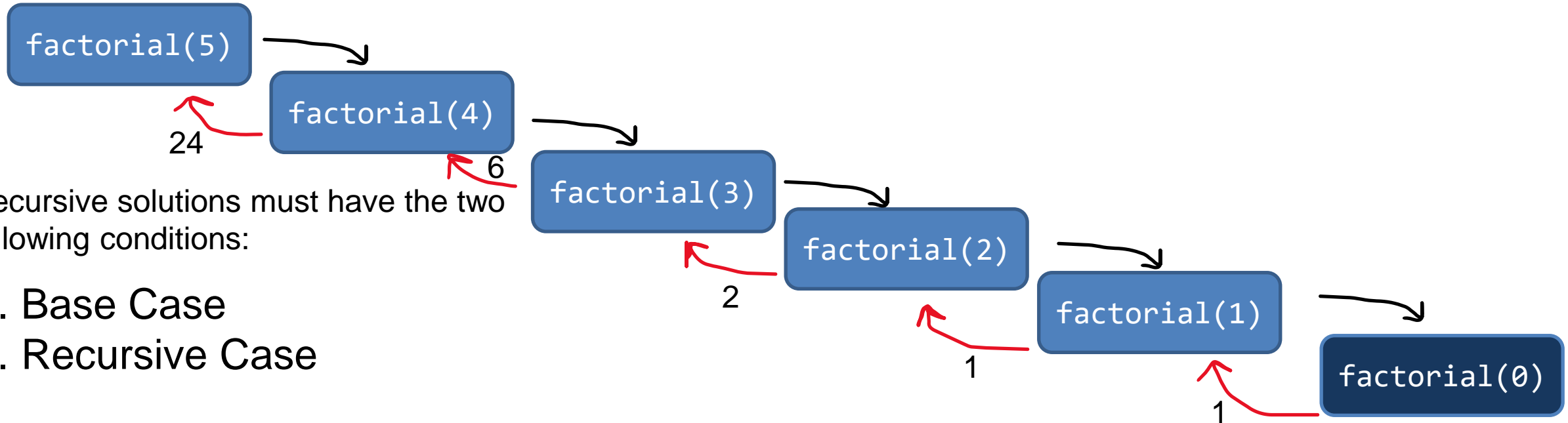
1. Base Case
2. Recursive Case

```
static int factorial(int n)
{
    if (n == 0)           (base case)
        return 1;
```

```
    return n * factorial(n - 1); (recursive case)
}
```

We can solve the factorial for n by solving smaller problems (factorial of $n-1$) !

120



Recursive solutions must have the two following conditions:

1. Base Case
2. Recursive Case

The Fibonacci sequence is a sequence in which each number is the sum of the two preceding ones

So, the N^{th} digit of the Fibonacci Sequence = $f(N-1) + f(N-2)$

The Fibonacci Sequence

1,1,2,3,5,8,13,21,34,55,89,144,233,377...

$$1+1=2$$

$$1+2=3$$

$$2+3=5$$

$$3+5=8$$

$$5+8=13$$

$$8+13=21$$

$$13+21=34$$

$$21+34=55$$

$$34+55=89$$

$$55+89=144$$

$$89+144=233$$

$$144+233=377$$

Because the solution to some problem can be expressed in terms of some smaller problem(s), recursion may be a good fit here

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$$55+89=144$$

$$89+144=233$$

$$144+233=377$$

Base Case?

Recursive Case?

Calculate

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So, the N^{th} digit of the Fibonacci Sequence = $f(N-1) + f(N-2)$

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Base Case?

If finding the 1st or 2nd digit, return 1

Recursive Case?

Calculate the previous two digits, $f(n-1)$, $f(n-2)$

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Base Case?

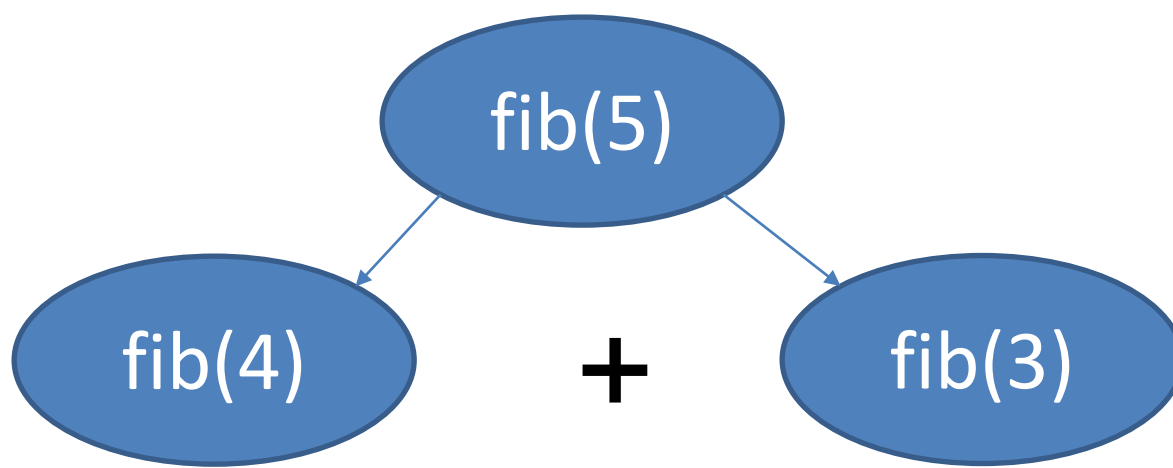
If finding the 1st or 2nd
digit, return 1

Recursive Case?

Calculate the previous
two digits, $f(n-1)$, $f(n-2)$

fib(5)

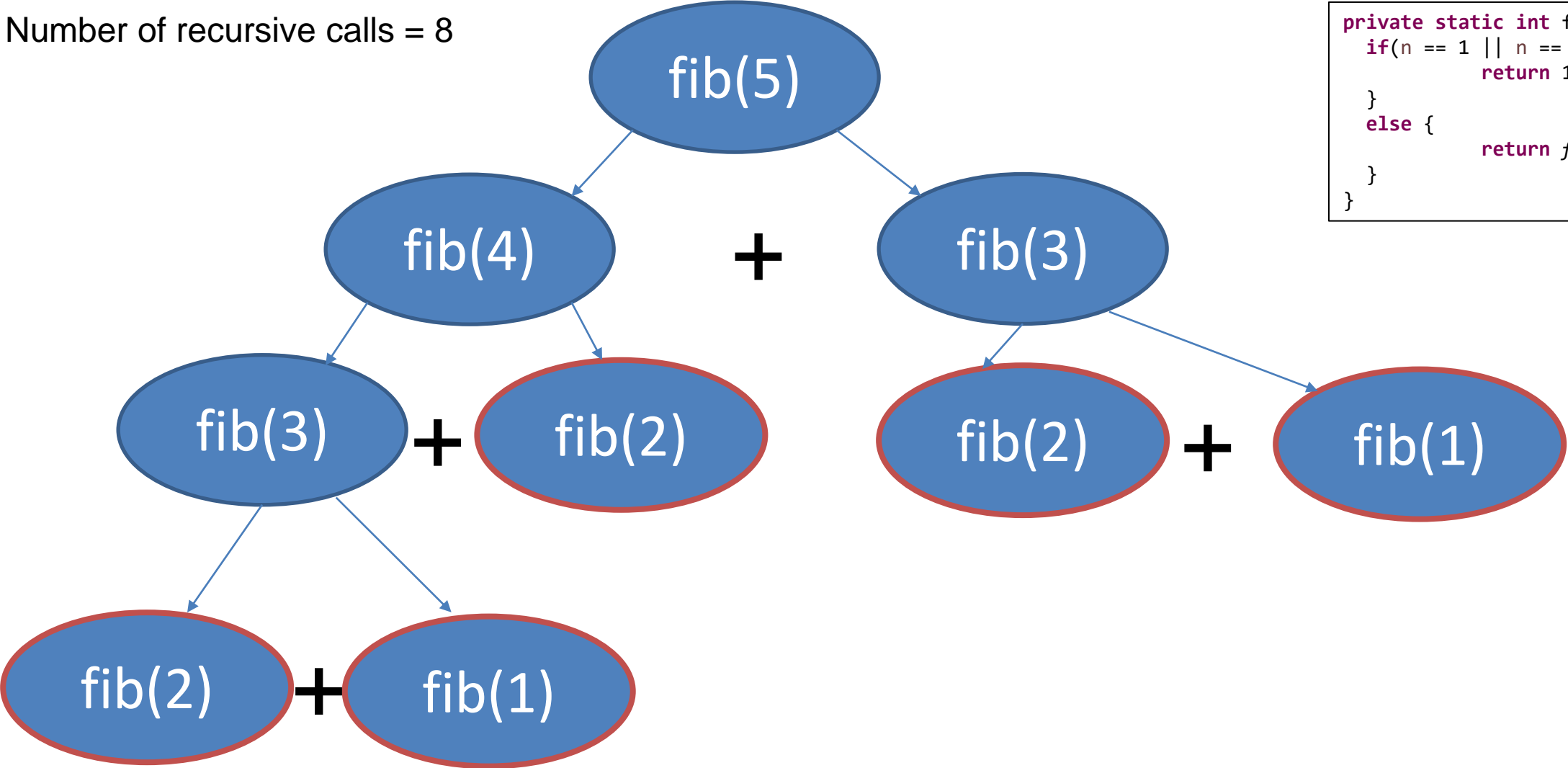
```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```



```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

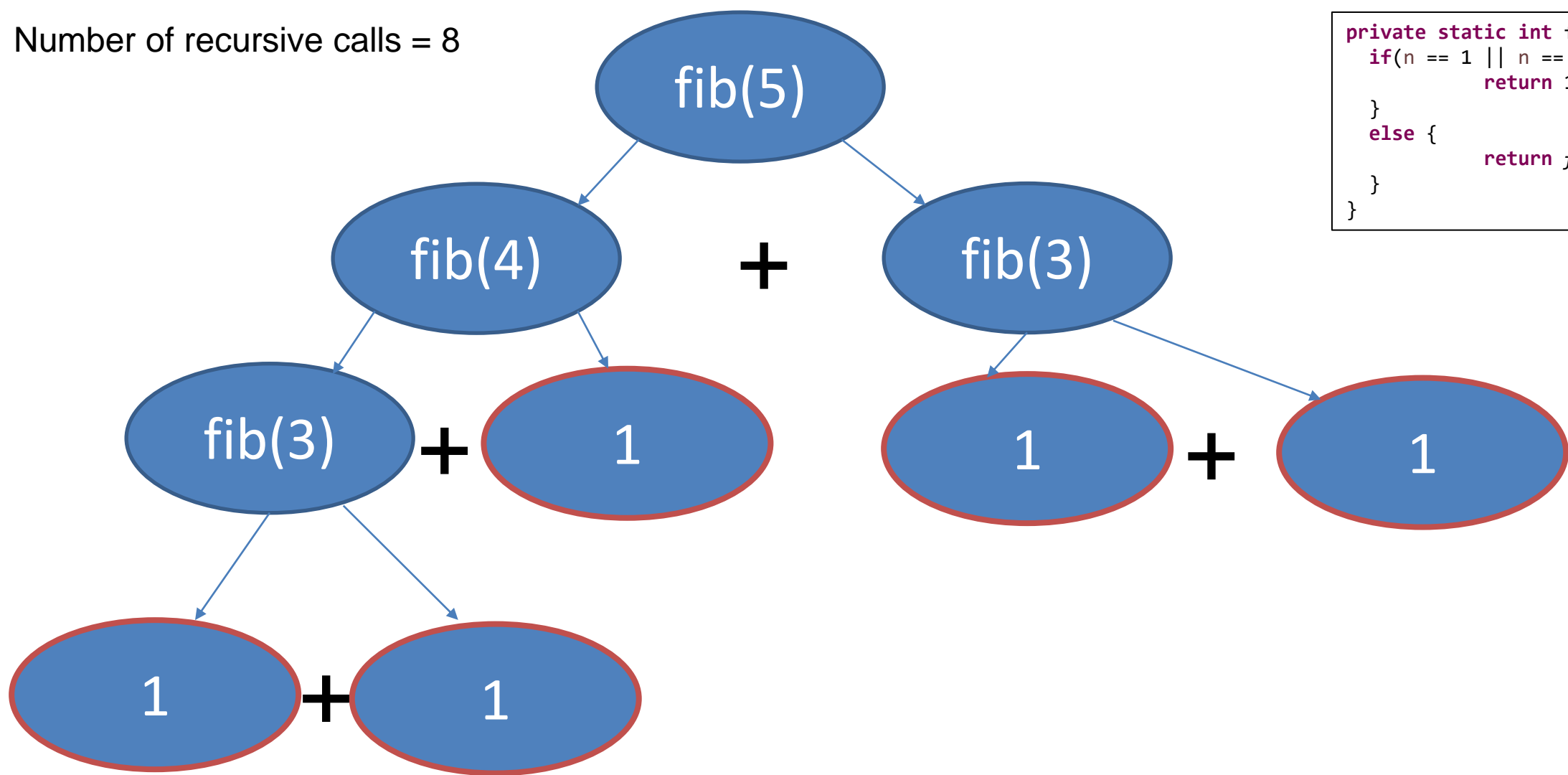
Number of recursive calls = 8

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

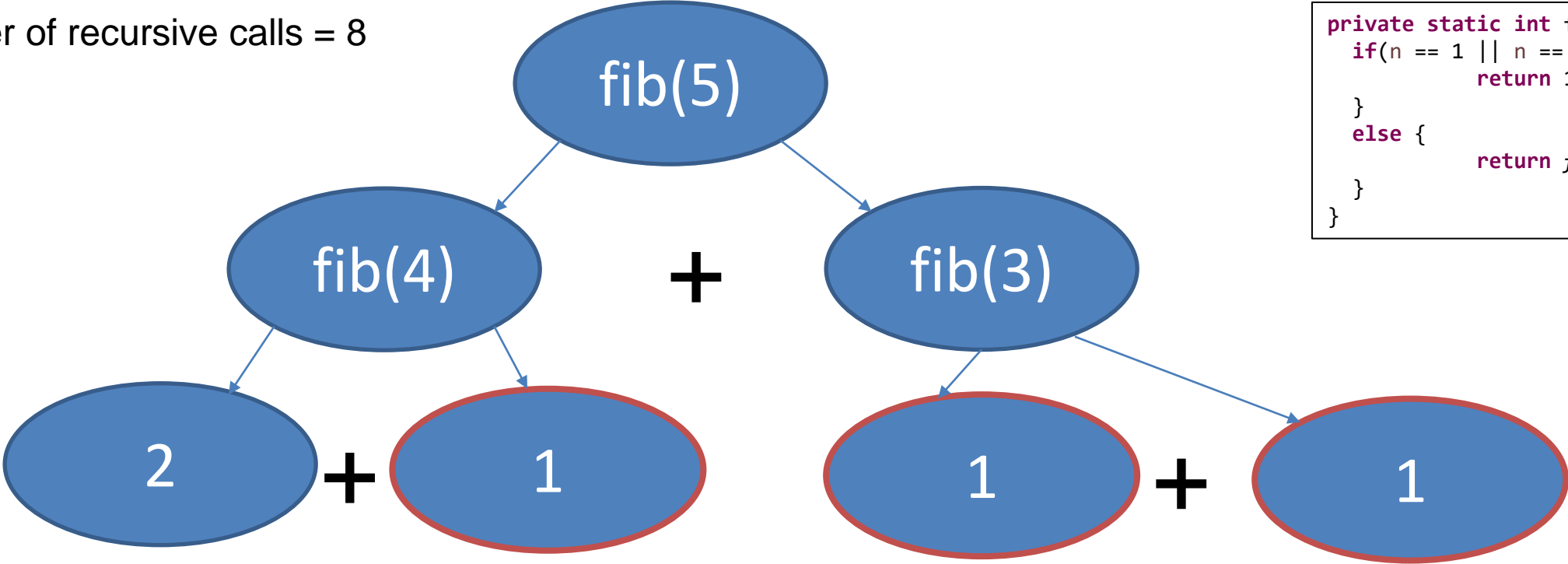


Number of recursive calls = 8

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

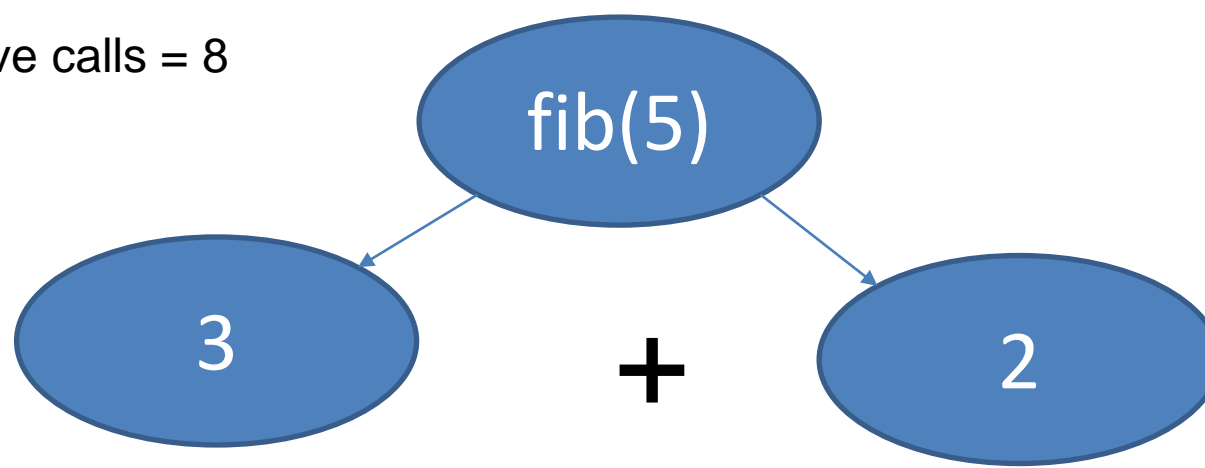


Number of recursive calls = 8



```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

Number of recursive calls = 8



```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```


Number of recursive calls = 8

5

Final answer!

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

Running Time?

```
private static int fib(int n) {  
    if(n == 1 || n == 2) { O(1)  
        return 1; O(1)  
    }  
    else {  
        O(1)          O(1)  
        return fib(n-1) + fib(n-2);  
    }  
}
```

Running Time?

O(1) ?

```
private static int fib(int n) {  
    if(n == 1 || n == 2) { O(1)  
        return 1; O(1)  
    }  
    else {  
        O(1)          O(1)  
        return fib(n-1) + fib(n-2);  
    }  
}
```

Running Time?

~~$\Theta(1)$~~ ?

No!

When we are analyzing recursive algorithms, we have to calculate running time slightly different

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

Generally speaking, we can compute the running time of a recursive algorithm by using the following formula:

Running time = # of recursive calls made * amount of work done in each call

```
private static int fib(int n) {  
    if(n == 1 || n == 2) { O(1)  
        return 1; O(1)  
    }  
    else {  
        O(1) O(1)  
        return fib(n-1) + fib(n-2);  
    }  
}
```

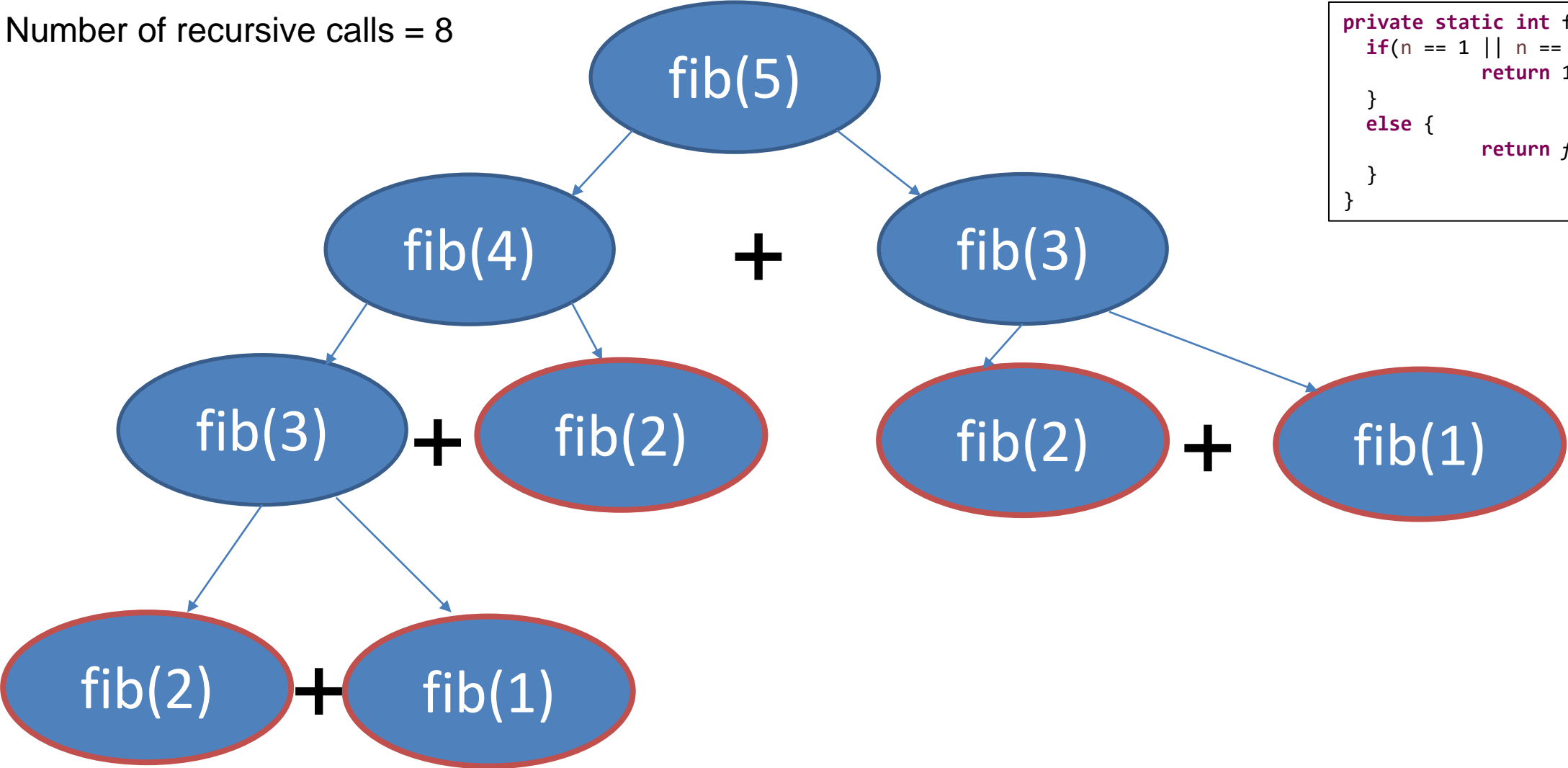
Generally speaking, we can compute the running time of a recursive algorithm by using the following formula:

Running time = # of recursive calls made * **amount of work done in each call**

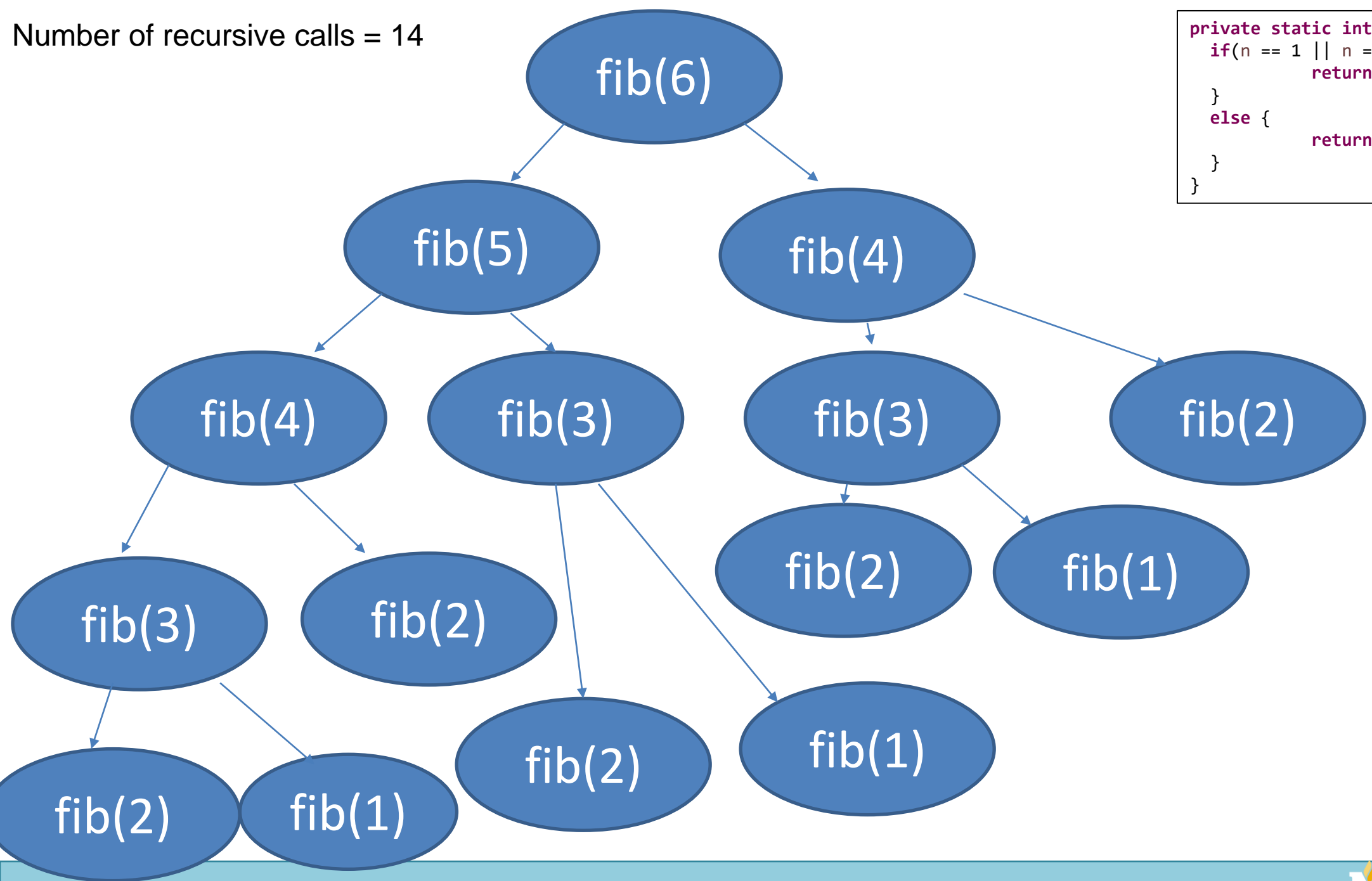
Running time = ??? * **O(1)**

Number of recursive calls = 8

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```



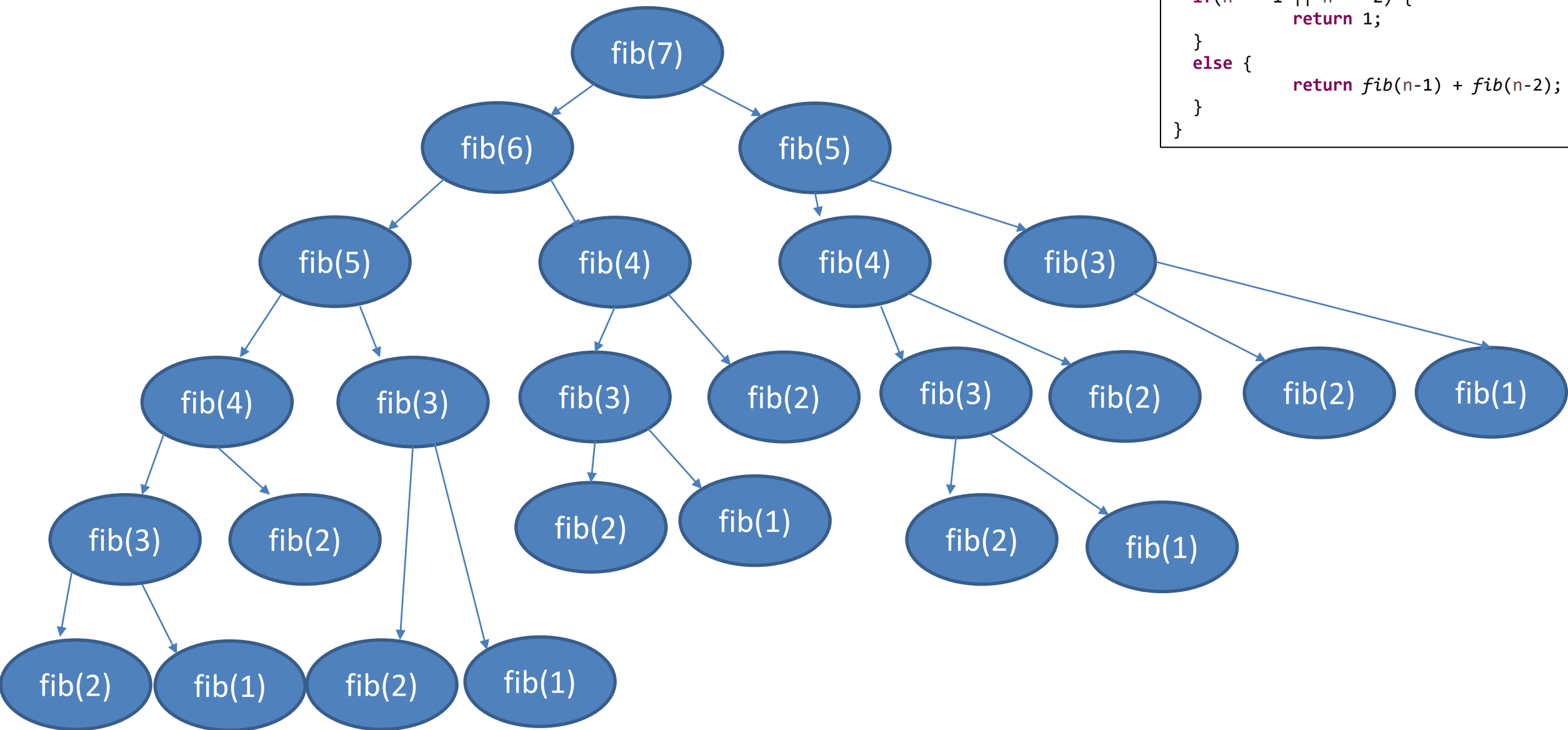
Number of recursive calls = 14



```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

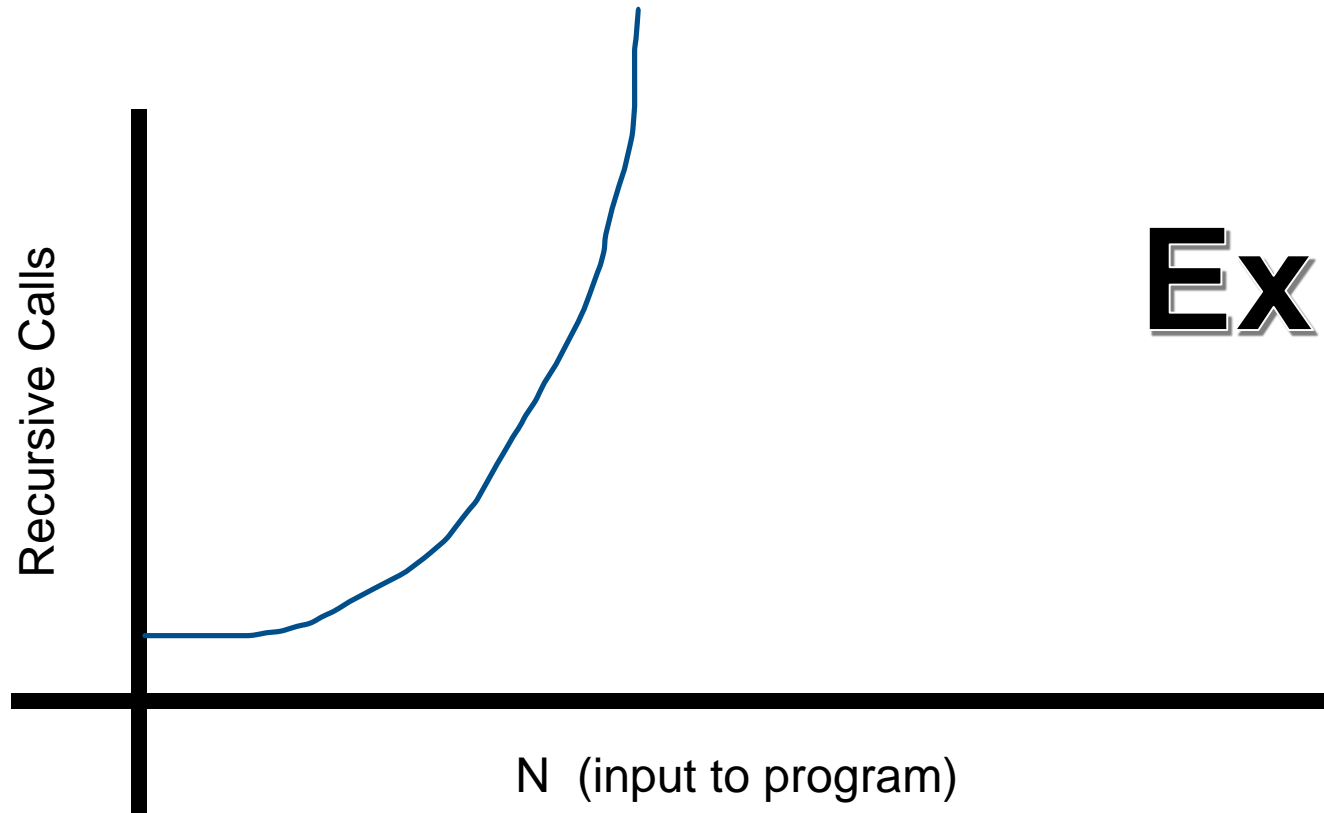

Number of recursive calls = 24

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```



If we were to plot the number of recursive calls made as n increases, it would look something like this:

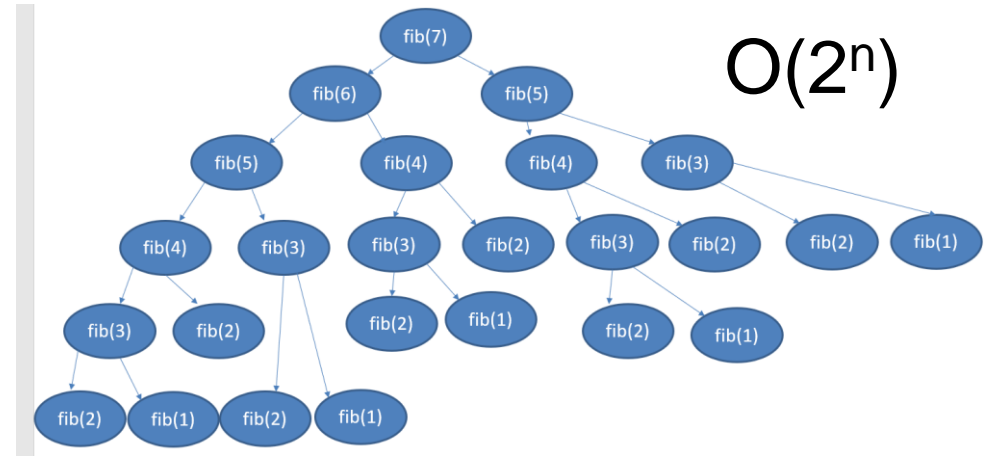
```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```



Exponential

Aka. $O(2^n)$

```
private static int fib(int n) {
    if(n == 1 || n == 2) { O(1)
        return 1; O(1)
    }
    else {
        return fib(n-1) + fib(n-2); O(1)
    }
}
```



Generally speaking, we can compute the running time of a recursive algorithm by using the following formula:

Running time = # of recursive calls made * amount of work done in each call

Running time = **$O(2^n)$ * $O(1)$**

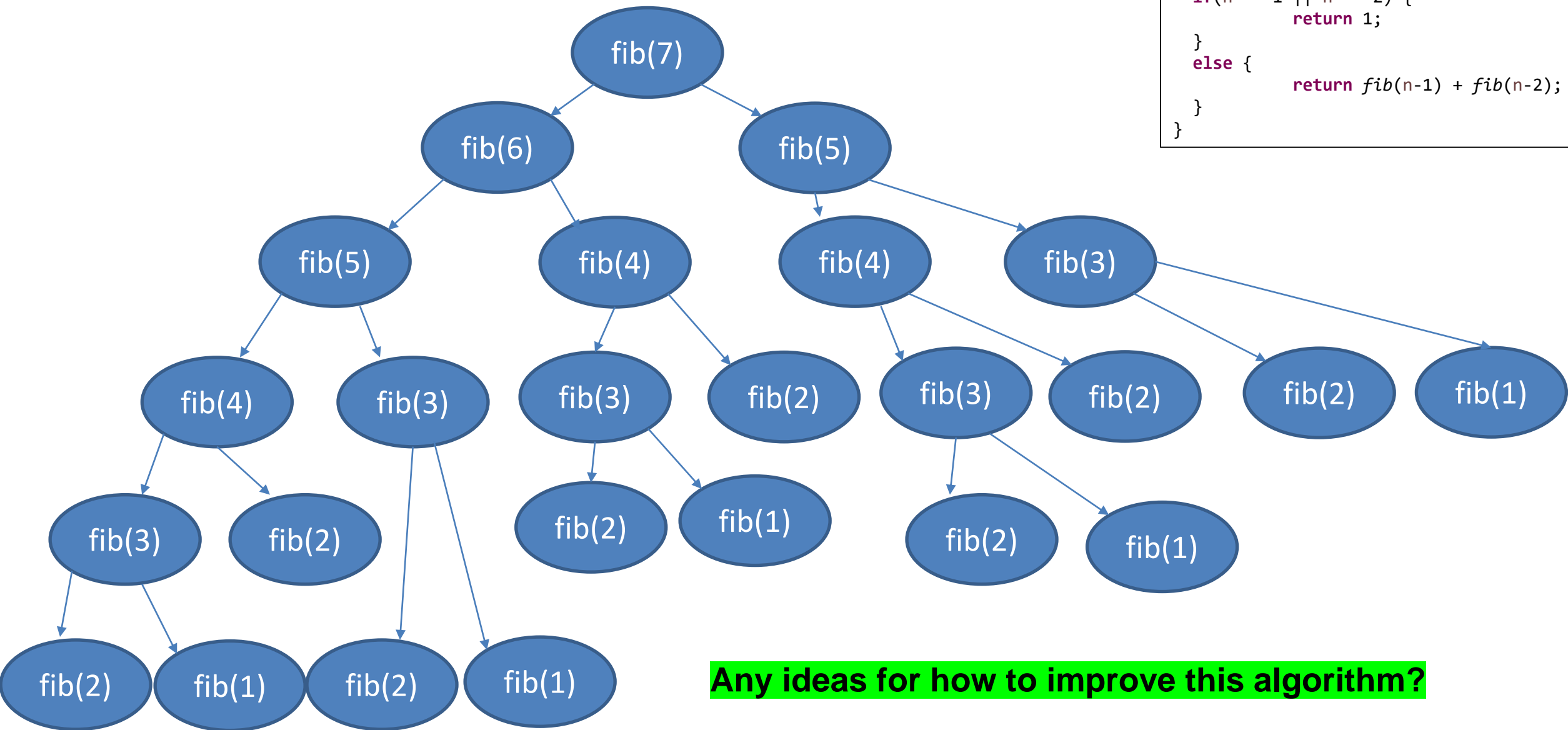
Total running time = **$O(2^n)$**

n = requested Fibonacci digit

$O(2^n)$ is very bad...

Number of recursive calls = 24

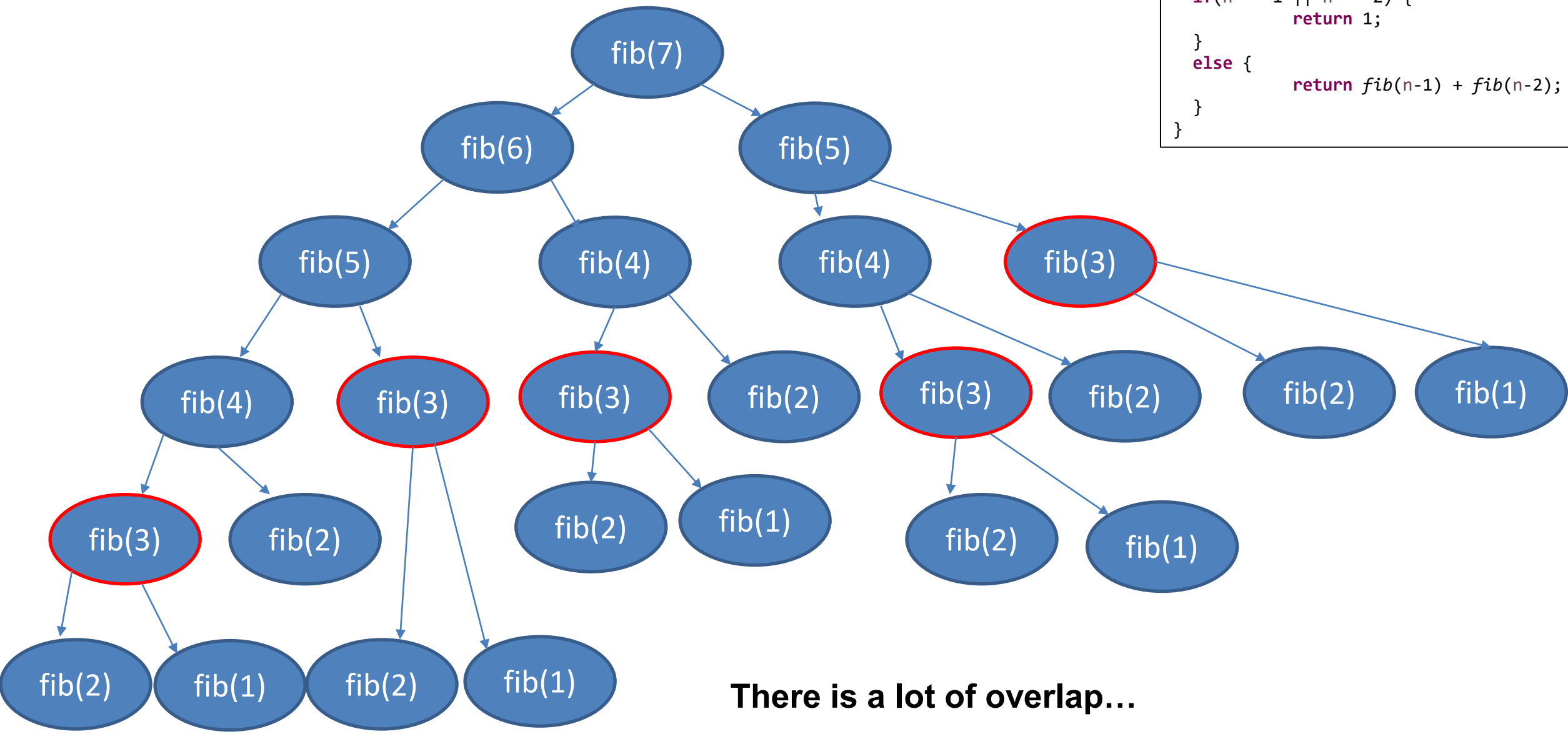
```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```



Any ideas for how to improve this algorithm?

Number of recursive calls = 24

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```

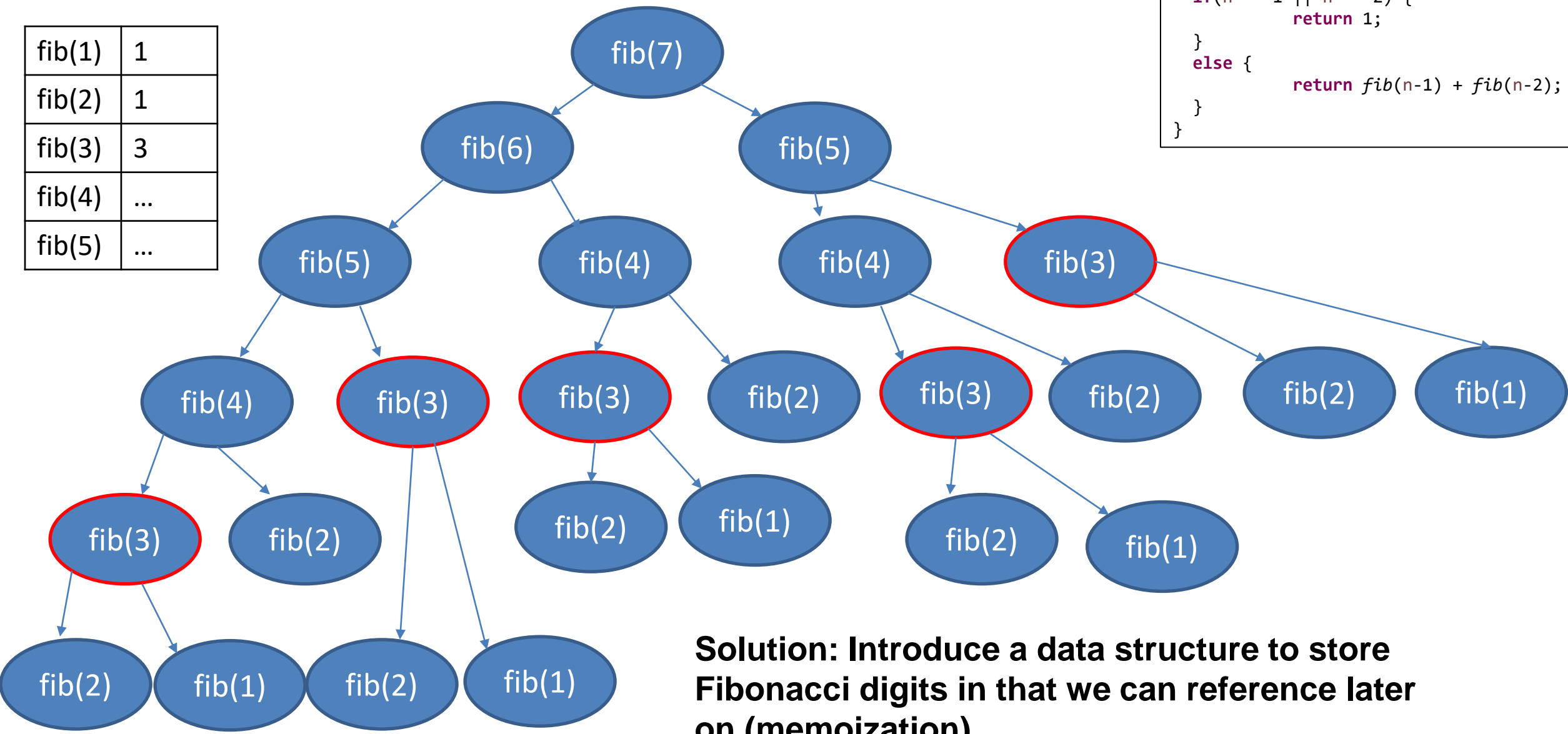


There is a lot of overlap...

Number of recursive calls = 24

fib(1)	1
fib(2)	1
fib(3)	3
fib(4)	...
fib(5)	...

```
private static int fib(int n) {  
    if(n == 1 || n == 2) {  
        return 1;  
    }  
    else {  
        return fib(n-1) + fib(n-2);  
    }  
}
```



Solution: Introduce a data structure to store Fibonacci digits in that we can reference later on (memoization)
(These lookups happen in constant time!)

Limitations of recursion?

Bubble Sorting Recursively

2	5	1	3	4
---	---	---	---	---

2	1	3	4	5
---	---	---	---	---

1	2	3	4	5
---	---	---	---	---

1	2	3	4	5
---	---	---	---	---

1	2	3	4	5
---	---	---	---	---

Bubble sort can be solved by solving smaller instances of bubble sort!

Base Case:

If the unsorted portion of the array is of size 1, return current array

Recursive Case:

Do one iteration of bubble sort, recursively call method and pass smaller array


```

static void bubbleSort(int arr[], int n)
{
    // Base case
    if (n == 1)
        return;

    int count = 0;
    for (int i=0; i<n-1; i++)
        if (arr[i] > arr[i+1])
        {
            // swap
            int temp = arr[i];
            arr[i] = arr[i+1];
            arr[i+1] = temp;
            count = count+1;
        }

    if (count == 0)
        return;

    bubbleSort(arr, n-1);
}

```

Recursive case