LECTURE 9

Decentralised Coordination

Contents

- Working Together
- Decentralised Coordination
- Distributed Constraint Optimisation Problems
 - Max-sum algorithm

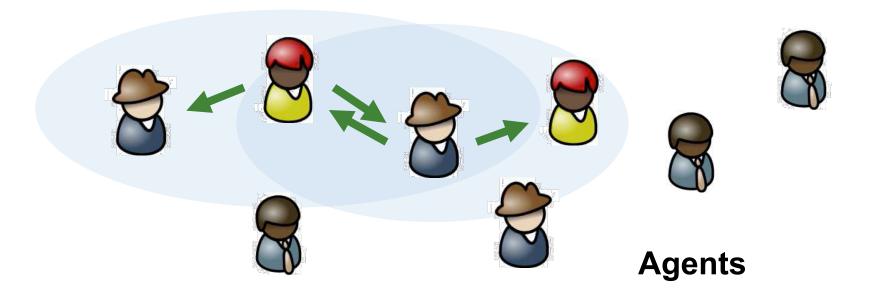
Working Together

- In many settings agents must work together to achieve their goals
 - When filtering options, the choices of other agents must be considered
 - In Robocup rescue the time to rescue a civilian depends on how many agents are working together
 - Depends on number of agents to be considered
 - Many agents model collectively as the environment
 - Few agents interesting difficult problem
 - No agents same case as before

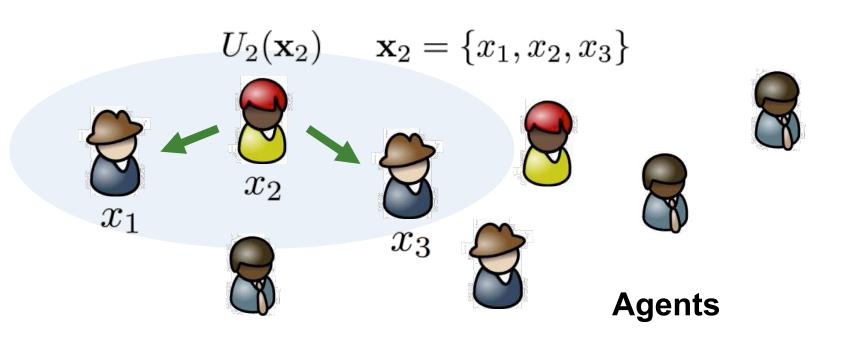
Discrete set of possible actions

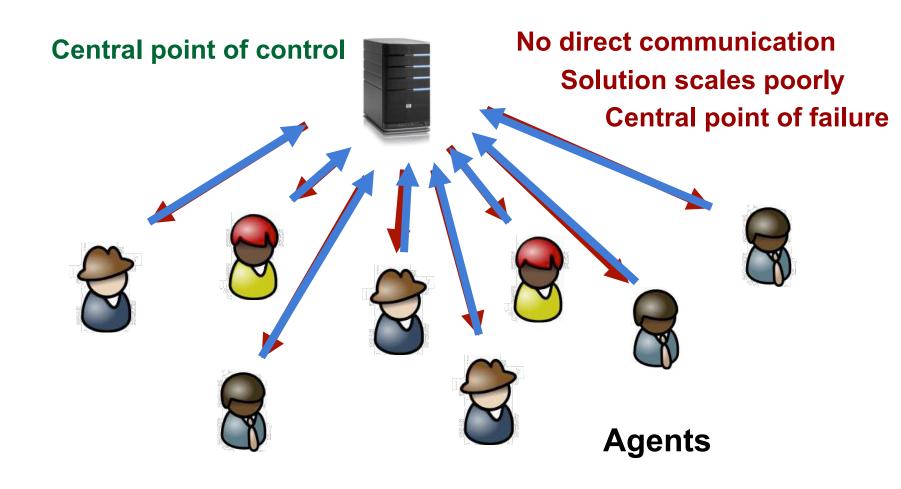


Discrete set of possible actions Some locality of interaction



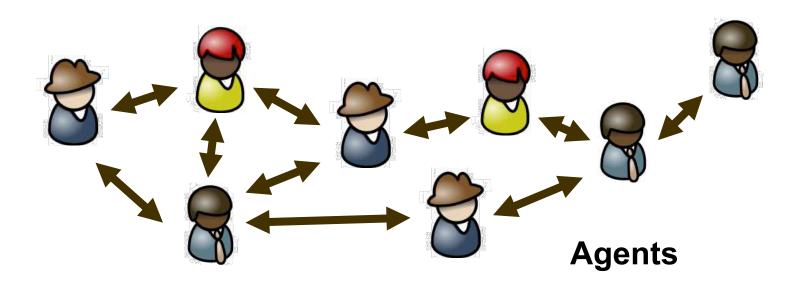
Discrete in the ration $\max_{\mathbf{x}} \sum_{m=1}^{\infty} U_m(\mathbf{x}_m)$ Some locality of interaction





Decentralised control and coordination through local computation and message passing.

 Speed of convergence, guarantees of optimality, communication overhead, computability

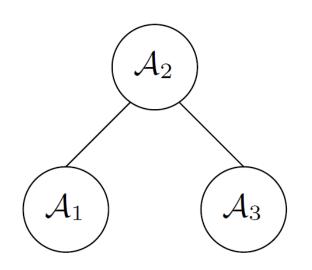


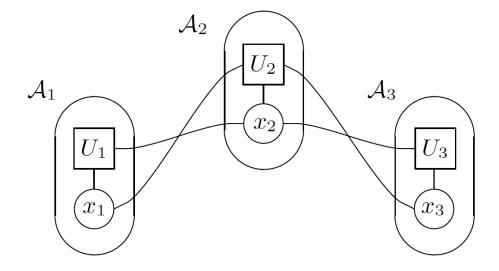
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Factor Graphs

 Can represent utility relationships as bipartite factor graphs



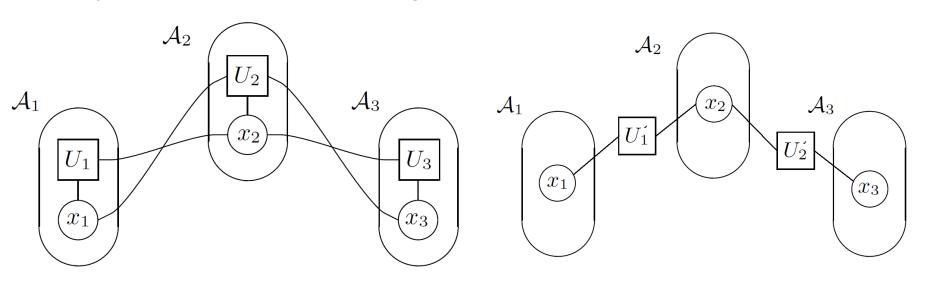


Agent graph

Factor graph

Constraint Graphs

 Often we can rearrange the factor graph to yield a constraint graph



Factor graph

Constraint graph

$$U_1(x_1, x_2) + U_2(x_1, x_2, x_3) + U_3(x_2, x_3) = U'_1(x_1, x_2) + U'_2(x_2, x_3)$$
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Max-sum Algorithm

- Applies to tree-based constraint graphs
- Algorithm proceeds as follows:
 - Leaf nodes send messages
 - If a node has k neighbours, it waits until it has received all k-1 messages, before sending message to each particular neighbour
 - Algorithm ends when each node has received k messages

Max-Sum

 Messages flow between function and variable nodes of the factor graph

From variable to factor:

$$Q_{n\to m}(x_n) = \sum_{m'\in M(n)\backslash m} R_{m'\to n}(x_n)$$

From factor to variable:

$$R_{m\to n}(x_n) = \max_{\mathbf{x}_m \setminus n} \left(U_m(\mathbf{x}_m) + \sum_{n' \in N(m) \setminus n} Q_{n'\to m}(x_{n'}) \right)$$

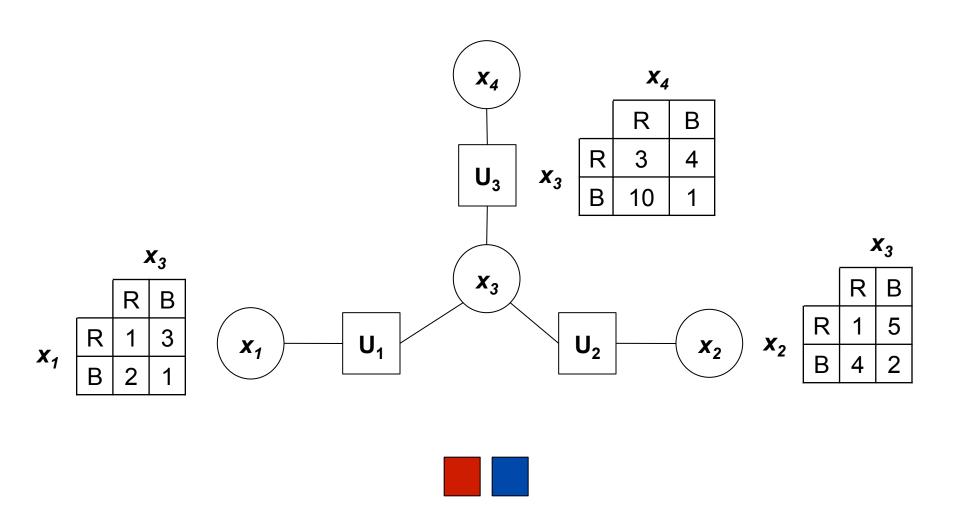
Max-Sum

At completion

$$Z_n(x_n) = \sum_{m \in M(n)} R_{r \to n}(x_n)$$

Solution given by

$$x_n^* = \arg\max_{x_n} Z_n(x_n)$$



Step 1

$$Q_{1\to 1}(x_1) = \{0,0\}$$

$$Q_{2\to 2}(x_2) = \{0,0\}$$

$$Q_{4\to3}(x_4) = \{0,0\}$$

Step 2

$$R_{1\to3}(x_3)=\{2,3\}$$

$$R_{2\to3}(x_3)=\{4,5\}$$

$$R_{3\rightarrow 3}(x_4) = \{4, 10\}$$

Step 3

$$Q_{3\rightarrow 1}(x_3)=\{8,15\}$$

$$Q_{3\rightarrow 2}(x_3)=\{6,13\}$$

$$Q_{3\rightarrow 3}(x_3) = \{6, 8\}$$

Step 4

$$R_{1\to 1}(x_1) = \{18, 16\}$$

$$R_{2\to 2}(x_2) = \{18, 15\}$$

$$R_{3\to4}(x_4)=\{18,10\}$$

Step 5

$$Z_1(x_1) = \{18, 16\}$$

$$Z_2(x_2) = \{18, 15\}$$

$$Z_3(x_3) = \{10, 18\}$$

$$Z_1(x_1) = \{18, 16\}$$
 $Z_2(x_2) = \{18, 15\}$ $Z_3(x_3) = \{10, 18\}$ $Z_4(x_4) = \{18, 10\}$

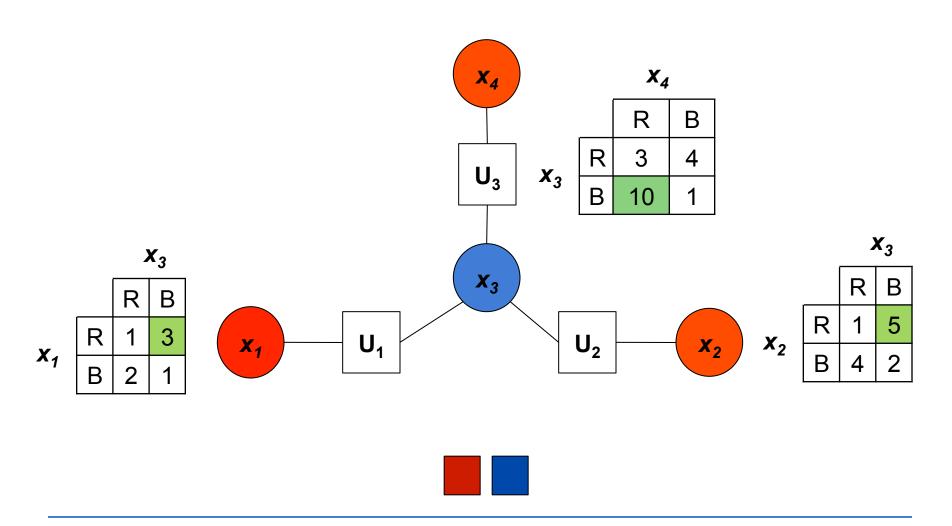
Step 7

$$X_1 = R$$

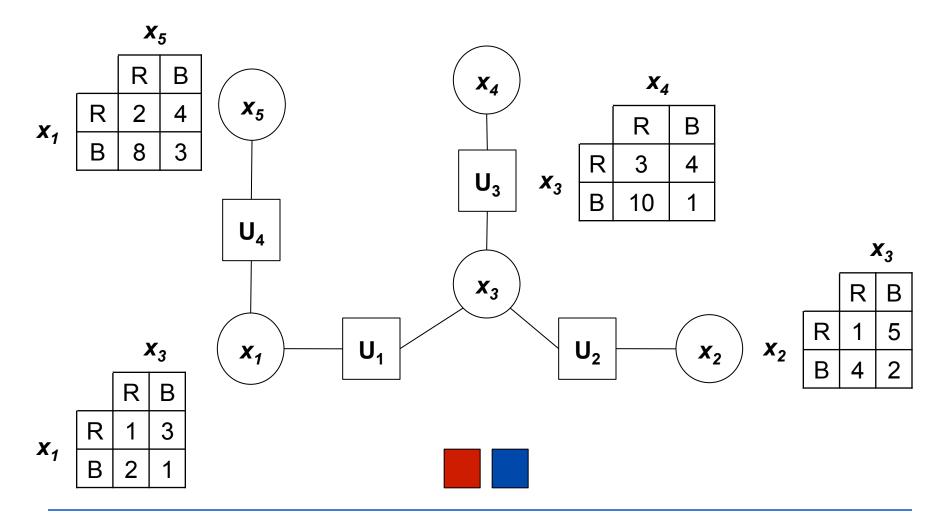
$$x_2 = R$$

$$x_3 = B$$

$$x_1 = R$$
 $x_2 = R$ $x_3 = B$ $x_4 = R$



Exercise



Max-Sum Properties

- Optimal algorithm
 - When constraint graph is a tree
- Solution converges after 2 x (n-1) messages
- Can propagate messages iteratively and asynchronously
 - Will still converge to the optimal solution

DPOP Algorithm

- Extends max-sum to constraint graphs with cycles or loops
 - Create pseudo-tree
 - Extend variables to combinations of variables
 - This is like a junction tree in graphical models
 - Optimal algorithm
 - Exponentially sized messages!

Other Approaches

- Bounded max-sum
 - Prune cyclic graph back to a tree
 - Optimally solve the new problem
 - Calculate maximum bound from solution of the original problem.
- Loopy max-sum
 - Use max-sum as is even though the graph has cycles
 - Normalise messages to stop them exploding
 - Works surprisingly well

Reading

- Weiss: Multiagent Systems
 - Chapter 12