

***CSE-312***  
***OPERATING SYSTEM***  
***HW-3 part1***

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# The Second Chance (SC)

.The referenced variable is equal to the second chance bit.

```
int replaceIndex = pt.queue.at(0), pageTableIndex = 0;
for (pageTableIndex = 0; pageTableIndex < pt.tableSize; ++pageTableIndex)
    if (replaceIndex == pt.table[pageTableIndex].page_index)
        break;

PageTableEntry ptReplace = pt.table[pageTableIndex];
virtualArray->pm->valids[ptReplace.pmpage_index]=1;

int validPMIndex = getValidpage_index();
virtualArray->pr.page_indexPM = validPMIndex;

auto it = pt.queue.begin();

if (ptReplace.referenced)
{
    replaceIndex = pt.queue.at(1);
    ptReplace.referenced=0;
    ++it;
}

for (pageTableIndex = 0; pageTableIndex < pt.tableSize; ++pageTableIndex)
    if (replaceIndex == pt.table[pageTableIndex].page_index)
        break;

pt.queue.erase(it);
pt.queue.push_back(page_index);
ptReplace.page_index=page_index;
ptReplace.pmpage_index=validPMIndex;

writeDisk(ptReplace.pmpage_index);

pt.table[pageTableIndex]= ptReplace;
```

.If the second chance bit is ONE, reset its second chance bit (to ZERO) and continue.

. If the second chance bit is ZERO, replace the page in that memory frame.

.The page to be changed has been found, deleted, and a new page has been added.

.New page write disk

# Least-Recently-Used (LRU)

```
int replaceIndex = pt.queue.at(0), pageTableIndex = 0;
for (pageTableIndex = 0; pageTableIndex < pt.tableSize; ++pageTableIndex)
    if (replaceIndex == pt.table[pageTableIndex].page_index)
        break;

PageTableEntry ptReplace = pt.table[pageTableIndex];
virtualArray->pm->valids[ptReplace.pmpage_index]= 1;

int validPMIndex =getValidpage_index();
virtualArray->pr.page_indexPM = validPMIndex;

writeDisk(ptReplace.pmpage_index);

ptReplace.page_index=page_index;
ptReplace.pmpage_index=validPMIndex;

auto it = pt.queue.begin();
pt.queue.erase(it);
pt.queue.push_back(page_index);

pt.table[pageTableIndex]= ptReplace;
```

- Whichever was used last will be replaced.
- The last used page is found and replaced.
- The most recently used page is kept in the queue so it was easy to reach.
- Queues are based on the FIFO principle, So, 0.index This is the least used page.

# Working Set Clock (WSClock)

```
int i = pt.queue.size()-1, pageTableIndex = 0, replaceIndex = 0;
for (; i >= 0; --i)
{
    replaceIndex = pt.queue.at(i);
    for (pageTableIndex = 0; pageTableIndex < pt.tableSize; ++pageTableIndex)
        if (replaceIndex == pt.table[pageTableIndex].page_index)
            break;
    if (pt.table[pageTableIndex].referenced)
    {
        pt.table[pageTableIndex].referenced=0;
    }
    else
    {
        virtualArray->pm->valids[pt.table[pageTableIndex].pmpage_index]= 1;
        int validPMIndex = getValidpage_index();
        virtualArray->pr.page_indexPM = validPMIndex;

        pt.table[pageTableIndex].page_index=page_index;
        pt.table[pageTableIndex].pmpage_index=validPMIndex;
        auto it = pt.queue.begin() + i;
        pt.queue.erase(it);
        pt.queue.push_back(page_index);

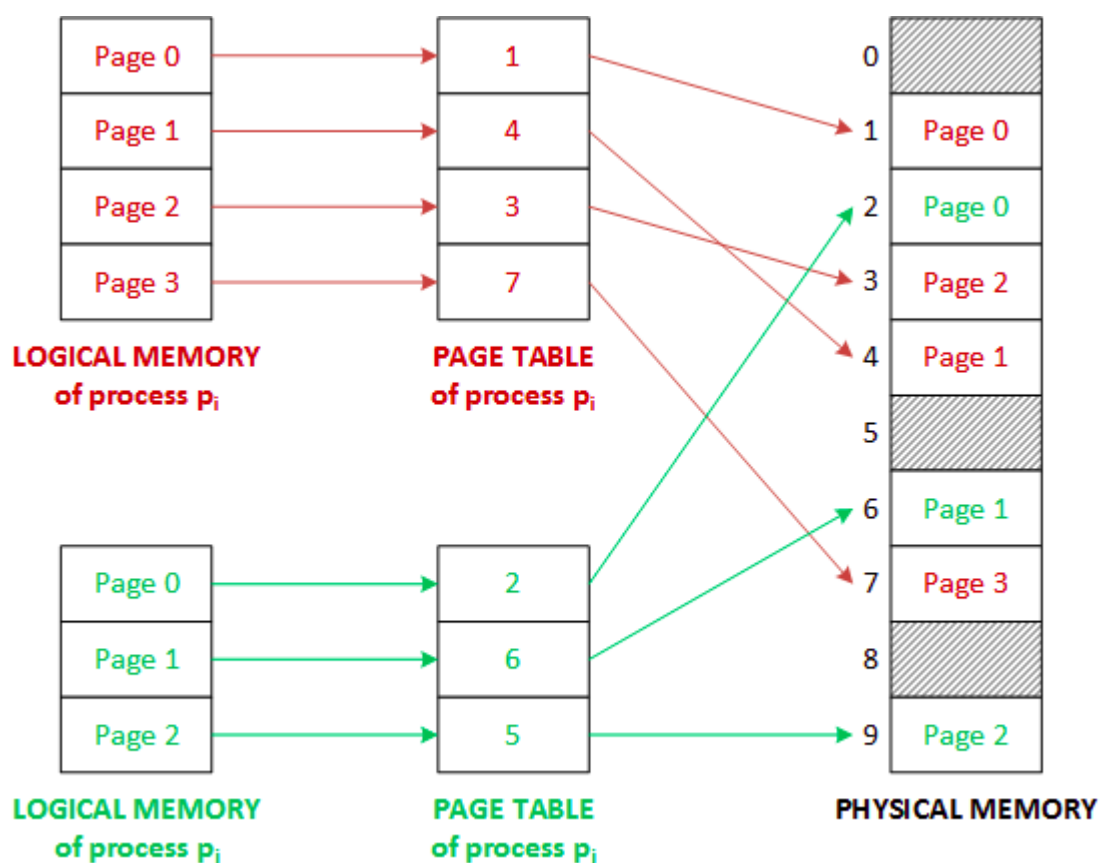
        replaceIndex = pt.queue.at(pt.queue.size()-1);
        for (pageTableIndex = 0; pageTableIndex < pt.tableSize; ++pageTableIndex)
            if (replaceIndex == pt.table[pageTableIndex].page_index)
                break;
        pt.table[pageTableIndex].referenced=1;
        break;
    }
}
```

*.Here the referenced is set to the R bit.*

*. If the R bit is set to 1, the page has been used during the current tick so it is not an ideal candidate to remove. The R bit is then set to 0, the hand advanced to the next page, and the algorithm repeated for that page.*

*.If the page pointed to has  $R = 0$ , it is not in the working set and a valid copy exists on the disk. That page call on disk.*

## Regular Page Table

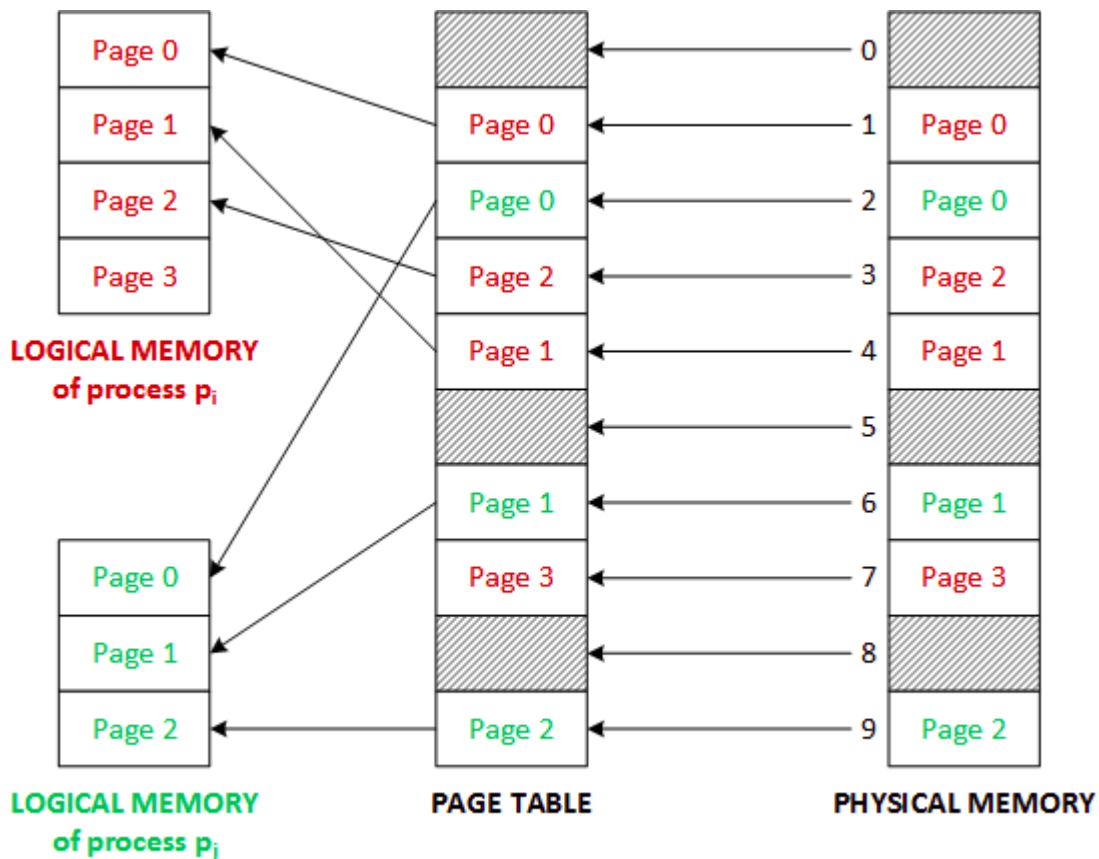


.A search is made within the page and accordingly we find whether it is a miss or a hit.

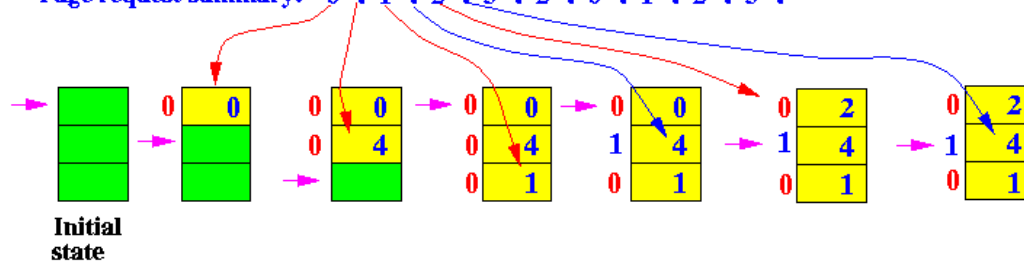
.When hit, referenced and modified is changed according to which page change happened, something is assigned to the queue.

.If miss, the page is retrieved from your physical memory.

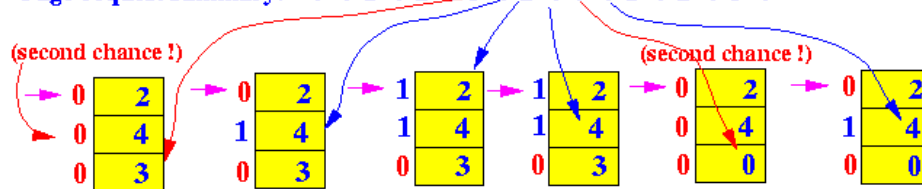
### *Inverted Page Table*



Page request summary: 0 4 1 4 2 4 3 4 2 4 0 4 1 4 2 4 3 4



Page request summary: 0 4 1 4 2 4 3 4 2 4 0 4 1 4 2 4 3 4



Page request summary: 0 4 1 4 2 4 3 4 2 4 0 4 1 4 2 4 3 4

