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CS350

HW 3

1.

```
int min, max;
array a;
check if array only have 1 data
if length==1{
    max = a[0];
    min = a[0];}
Initialize min and max
if a[0] > a[1] {
    max = a[0];
    min = a[1];}
else {
    max = a[1];
    min = a[0];}
start loop starting from 2
for (int i = 2; i <= a.length - 2;) {
    if (a[i] > a[i + 1]) {
        min = Math.min(min, a[i + 1]);
        max = Math.max(max, a[i]);
    } else {
        min = Math.min(min, a[i]);
        max = Math.max(max, a[i + 1]);
    }
    i = i + 2;}
Continuation for odd number length
if a.length % 2 == 1 {
    min = Math.min(min, a[a.length - 1]);
    max = Math.max(max, a[a.length - 1]);}
```
2. Let $i = 5$ and $n = 15$. S takes $O(i * n)$ time, which means $(5 * 15)$ in this case. While T takes $O(n \log n)$ time, which mean $(16 * \log 16) = (16 * 4)$ in this case. Hence, the algorithm S performs better than the algorithm T when the value of i is less than the value of $\log n$. On the other hand, Let $i = 12$ and $n = 16$. S takes $O(i * n)$ time to compute the result, which is equal to $(12 * 16)$ in this case. T takes $O(n \log n)$ time, which is equal to $(16 * \log 16) = (16 * 4)$ in this case. Hence, the algorithm T performs better than the algorithm S when the value of i is greater than the value of $\log n$.
3. When $k=3$, number of median & groups is $n/3$. Worst case scenario would be $n - n/3 = 2n/3$
 $Tw(n) \leq \theta(n) + Tw(n/3) + Tw(2n/3)$
When $k=7$, number of median & groups is $n/7$. Worst case scenario would be $n - 2n/8 = 5n/7$
 $Tw(n) \leq \theta(n) + Tw(n/7) + Tw(5n/7)$

4. Since we knew quickselect's worst case is $O(n^2)$, then the worst case of ilselect is $O(n^2)$.

While the average case, $T_{avg}(n)$ can be formulized as:

$$T_{avg}(n) = \theta(n) + \frac{1}{n} \sum (O(1) * \frac{1}{n} + O(r-1) * \frac{(r-1)}{n} + O(n-r) * \frac{(n-r)}{n})$$

Which then simplified

$$T_{avg}(n) = \theta(n) + \frac{2}{n} \sum_{1 \leq r \leq n} O(r-1) * \frac{(r-1)}{n}$$

Then we can simplify and get $T_{avg}(n) = \theta n$

5. Since the operation compares number, so the comparison of N will be
 $(N-1) + (N-2) + \dots + 1 = N(N-1)/2$