

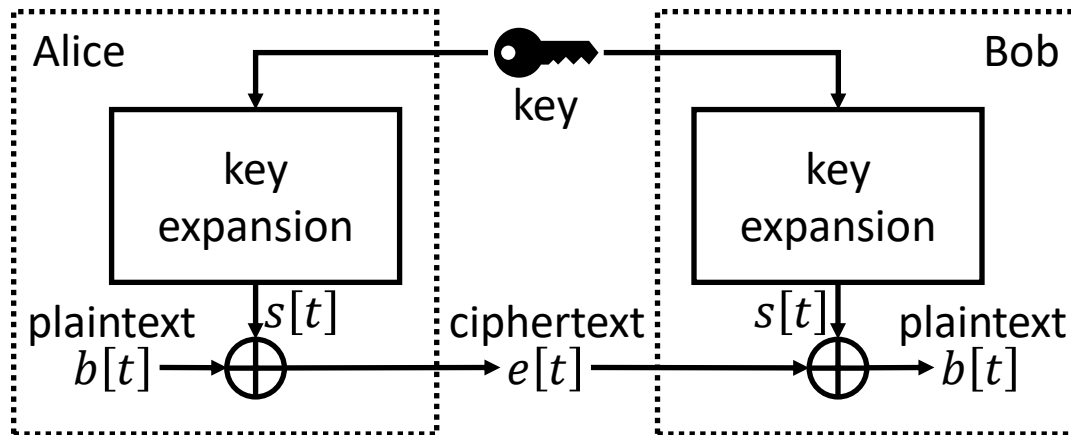
Stream Ciphers

Elements of Applied Data Security

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Stream Cipher

A stream cipher is a **symmetric key** cipher where the plaintext is encrypted (and ciphertext is decrypted) one digit at a time. A digit usually is either a bit or a byte.



Encryption (decryption) is achieved by xoring the plaintext (ciphertext) with a stream of pseudorandom digits obtained as an expansion of the key.

Task

1. Define an object that, starting from a random bit generator, generates random bytes.
2. Define an object implementing a Stream Cipher that, given a Pseudo Random Byte Generator and a key, can encrypt and decrypt a message.
3. Define an iterator that implements the A5/1 architecture and use it as bit generator in a Stream Cipher.
4. Define an iterator that implements the RC4 architecture and use it in the Stream Cipher object to encrypt/decrypt a message

Task 1: Pseudo Random Byte Generator

Pseudo Random Byte Generator

Inputs:

- **seed** (optional, default None): initial state
- **bit_generator** (optional, default None): Iterable that yields a pseudo-random bit starting from an initial seed. If None an LFSR is employed.

Yield:

- **byte**: pseudo-random byte.

Pseudo-Random Byte Generator

Template:

```
def pseudo_random_byte_generator(seed=None, bit_generator=None, **kwargs):  
    ''' generator docstring '''  
    # do stuff  
    while True:  
        # do stuff  
        yield byte
```

```
class PseudoRandomByteGenerator(object):  
    ''' class docstring '''  
  
    def __init__(self, seed=None, bit_generator=None, **kwargs):  
        ''' constructor docstring '''  
        # do stuff  
  
    def __iter__(self):  
        return self  
  
    def __next__(self):  
        # do stuff  
        return byte
```

Stream Cipher Example

- Example:

```
prbg = pseudo_random_byte_generator() # generator
digits = [byte for byte in islice(prbg, 10)]
print(bytes(digits)) # -> b'\x00x\xe7/\xdd7\xe1\\t\xc5'
```

```
prbg = PseudoRandomByteGenerator() # iterator
digits = [byte for byte in islice(prbg, 10)]
print(bytes(digits)) # -> b'\x00x\xe7/\xdd7\xe1\\t\xc5'
```

Python Module

- You may want to use a handy function in several notebooks without copying its definition into each notebook.
- You can put definitions in a `.py` file and use them in a notebook. Such file is a ***module*** and definitions can be *imported* into your notebook or in other modules.
- We can collect all the function and objects employed to implement an LFSR in a python module `lfsr.py` and import them in a new notebook.

LFSR module

Content of the module `lfsr.py`:

```
''' module docstring '''
# import modules
import math
...

# functions definition
def lfsr_generator(poly, state=None):
    ...

def berlekamp_massey(bits):
    ... ..

# classes definition
class LFSR():
    ...

class BerlekampMassey():
    ...
```

LFSR module

Example of usage in a notebook

```
import lfsr          # entire module imported
from lfsr import LFSR # direct import of the class named LFSR

# usage of the class LFSR
mylfsr = LFSR([16, 12, 3, 1, 0], state=0xFA42)
bitstream = mylfsr.cycle()

# usage of the function berlekamp_massey defined in the module
bm_poly = lfsr.berlekamp_massey(bitstream) # -> [16, 12, 3, 1, 0]
```

Task 2: Stream Cipher

Stream Cipher

Inputs:

- **key**: integer representing the shared secret key.
- **PRNG** (optional, default None): Iterator implementing a PRNG that produce a pseudorandom byte stream starting from an initial seed. If None an LFSR is used as PRNG.

Methods:

- **encrypt**: encrypts a plaintext (bytes) and returns the corresponding cyphertext (bytes);
- **decrypt**: decrypts a cypertext (bytes) and returns the corresponding plaintext (bytes);

Stream Cipher

Template:

```
class StreamCipher():
    ''' class docstring '''

    def __init__(self, key, prng=None, **kwargs):
        ''' constructor docstring '''
        # do stuff
        self.prng = ...

    def encrypt(self, plaintext):
        # do stuff
        return ciphertext

    def decrypt(self, ciphertext):
        # do stuff
        return plaintext
```

Stream Cipher Example

```
message = 'hello world!'
key = 0x0123456789ABCDEF

# create a StreamCipher instance for Alice and Bob
# no PRNG is specified, then LFSR is employed
alice = StreamCipher(key)
bob    = StreamCipher(key)

plaintextA = message.encode('utf-8')    # string to bytes
ciphertext = alice.encrypt(plaintextA)  # encryption by Alice
plaintextB = bob.decrypt(ciphertext)    # decryption by Bob

print(plaintextA) # -> b'hello world!'
print(ciphertext) # -> b'\x9f\x03\xbcf\xfa\xdb`\xf6\x17\xce7\x88'
print(plaintextB) # -> b'hello world!'
```

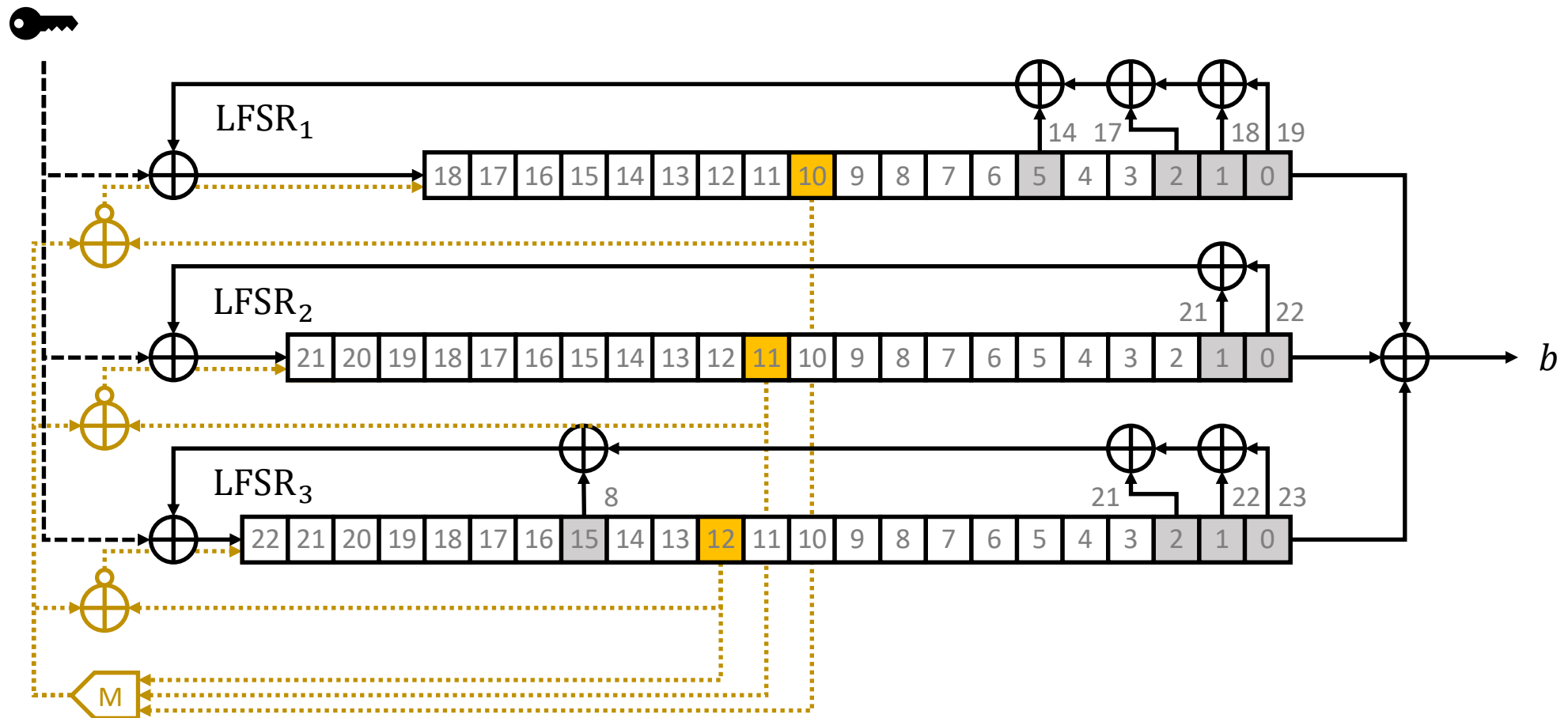
Task 3: A5/1

A5/1

- The **stream cipher** used to provide privacy in the **GSM** cellular telephone standard.
- Based on a combination of **three LFSRs with irregular clocking**.
At each cycle, the clocking bit of all three registers is examined and the majority bit is determined. A register is clocked if the clocking bit agrees with the majority bit.

LFSR	Length	Feedback Polynomial	Clocking bit
1	19	$x^{19} + x^{18} + x^{17} + x^{14} + 1$	10
2	22	$x^{22} + x^{21} + 1$	11
3	23	$x^{23} + x^{22} + x^{21} + x^8 + 1$	12

A5/1



A5/1

- A5/1 is initialised using a 64-bit private **key** together with a publicly known 22-bit **frame** number.
 - Initially, the registers are set to zero.
 - The 64-bit secret key is mixed: in cycle i (with $0 \leq i < 64$), the i -th key bit is added to the most significant bit of each register using XOR.
 - Similarly, the 22-bits frame number is added in 22 cycles.
- Then, the cipher is clocked using the normal majority clocking mechanism for 100 cycles, with the output discarded.
- Then, the cipher is ready to produce two 114 bit sequences of output keystream, first 114 for downlink, last 114 for uplink.

A5/1

- Define a Iterator that implements the A5/1 stream cipher.
- Given a key and a frame number, encrypt and decrypt a message.

```
message = 'hello world!'
key, frame = 0x0123456789ABCDEF, 0x2F695A

# create a StreamCipher instance for both Alice and Bob
alice = StreamCipher(key, PseudoRandomByteGenerator, bit_generator=A5_1, frame=frame)
bob    = StreamCipher(key, PseudoRandomByteGenerator, bit_generator=A5_1, frame=frame)

plaintextA = message.encode('utf-8')    # -> b'Hello world!'
ciphertext = alice.encrypt(plaintextA)  # -> b'\xec\xe3zM;@\x08\xf3r\xa0\x14\x1c'
plaintextB = bob.decrypt(ciphertext)    # -> b'Hello world!'
```

A5/1 Key insertion

<u>iter</u>	<u>key</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>	<u>iter</u>	<u>key</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>	<u>iter</u>	<u>key</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>
1	1	40000	200000	400000	0	0	22	1	141bd	0bcdef	5245de	0	0	43	1	5f949	376a7c	5ab312	1	1
2	1	60000	300000	600000	0	0	23	0	4a0de	05e6f7	2922ef	0	0	44	0	6fca4	1bb53e	2d5989	1	1
3	1	70000	380000	700000	0	0	24	1	6506f	22f37b	149177	0	1	45	0	37e52	2dda9f	56acc4	1	1
4	1	78000	3c0000	780000	0	0	25	1	72837	3179bd	4a48bb	0	1	46	0	5bf29	16ed4f	2b5662	1	0
5	0	3c000	1e0000	3c0000	0	0	26	0	3941b	38bcde	25245d	1	0	47	1	6df94	2b76a7	15ab31	0	0
6	1	5e000	2f0000	5e0000	0	0	27	0	1ca0d	3c5e6f	12922e	1	0	48	0	76fca	15bb53	0ad598	1	1
7	1	6f000	378000	6f0000	0	0	28	1	4e506	3e2f37	094917	1	0	49	1	3b7e5	2adda9	056acc	1	0
8	1	77800	3bc000	778000	0	0	29	0	27283	1f179b	44a48b	0	1	50	1	1dbf2	156ed4	02b566	1	0
9	1	7bc00	3de000	3bc000	0	0	30	0	13941	0f8bcd	625245	1	1	51	0	0edf9	0ab76a	415ab3	1	0
10	0	3de00	1ef000	5de000	0	0	31	0	49ca0	27c5e6	312922	0	0	52	0	076fc	255bb5	20ad59	1	0
11	1	5ef00	2f7800	2ef000	1	0	32	1	24e50	13e2f3	189491	1	0	53	0	03b7e	32adda	1056ac	1	0
12	1	6f780	37bc00	177800	1	0	33	1	52728	29f179	4c4a48	0	1	54	1	01dbf	1956ed	082b56	0	0
13	0	37bc0	1bde00	0bbc00	1	0	34	1	29394	14f8bc	662524	0	0	55	0	00edf	2cab76	0415ab	1	0
14	0	1bde0	0def00	45de00	1	0	35	1	149ca	2a7c5e	331292	1	0	56	0	4076f	3655bb	020ad5	0	1
15	1	0def0	26f780	22ef00	0	0	36	0	4a4e5	353e2f	598949	1	1	57	1	603b7	3b2add	41056a	0	0
16	1	06f78	337bc0	117780	1	0	37	0	65272	1a9f17	2cc4a4	0	1	58	0	301db	3d956e	6082b5	0	0
17	1	037bc	39bde0	48bbc0	1	0	38	1	72939	2d4f8b	566252	0	0	59	0	180ed	3ecab7	70415a	0	0
18	1	41bde	3cdef0	245de0	1	0	39	1	7949c	36a7c5	2b3129	1	0	60	0	4c076	1f655b	7820ad	0	0
19	0	20def	1e6f78	122ef0	1	1	40	0	7ca4e	3b53e2	559894	0	0	61	0	6603b	0fb2ad	3c1056	0	0
20	1	506f7	2f37bc	491778	1	1	41	1	7e527	1da9f1	6acc4a	1	0	62	0	7301d	27d956	1e082b	0	0
21	0	2837b	179bde	248bbc	0	1	42	0	3f293	2ed4f8	356625	0	0	63	0	3980e	33ecab	0f0415	0	0
														64	0	1cc07	19f655	07820a	0	0

A5/1 Frame number insertion

<u>iter</u>	<u>frame</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>
1	0	4e603	2cfb2a	03c105	1	0
2	1	67301	167d95	01e082	0	0
3	0	73980	2b3eca	00f041	1	1
4	1	79cc0	159f65	407820	1	1
5	1	7ce60	0acfb2	603c10	1	0
6	0	7e730	2567d9	301e08	1	1
7	1	3f398	12b3ec	580f04	0	0
8	0	1f9cc	0959f6	6c0782	0	0
9	1	0fce6	04acfb	3603c1	1	0
10	0	47e73	02567d	5b01e0	0	0
11	0	63f39	212b3e	2d80f0	1	1
12	1	71f9c	10959f	16c078	0	1
13	0	78fce	084acf	4b603c	1	1
14	1	7c7e7	242567	25b01e	1	0
15	1	7e3f3	3212b3	12d80f	0	1
16	0	7f1f9	190959	096c07	0	1
17	1	7f8fc	0c84ac	04b603	0	1
18	1	7fc7e	264256	025b01	1	1
19	1	3fe3f	13212b	012d80	0	0
20	1	5ff1f	299095	4096c0	1	0
21	0	6ff8f	34c84a	604b60	1	1
22	1	37fc7	1a6425	7025b0	0	0

A5/1 Key mixing (1)

<u>iter</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>	<u>iter</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>	<u>iter</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>
1	37fc7	2d3212	3812d8	1	1	21	10adb	356376	35c9c0	0	1	41	2ab10	34b375	190eda	0	1
2	5bfe3	2d3212	1c096c	0	1	22	0856d	3ab1bb	5ae4e0	0	0	42	15588	3a59ba	4c876d	1	1
3	5bfe3	369909	4e04b6	1	0	23	0856d	1d58dd	6d7270	1	0	43	0aac4	3d2cdd	4c876d	0	0
4	6dff1	3b4c84	4e04b6	1	1	24	442b6	2eac6e	36b938	1	0	44	45562	3d2cdd	6643b6	1	1
5	36ff8	1da642	4e04b6	0	0	25	442b6	375637	5b5c9c	0	1	45	22ab1	3e966e	6643b6	0	1
6	36ff8	2ed321	27025b	0	0	26	6215b	1bab1b	5b5c9c	1	0	46	11558	3f4b37	3321db	1	0
7	36ff8	376990	13812d	1	1	27	6215b	0dd58d	6dae4e	0	0	47	08aac	1fa59b	3321db	0	0
8	5b7fc	1bb4c8	13812d	0	1	28	310ad	26eac6	76d727	1	0	48	04556	0fd2cd	1990ed	1	0
9	5b7fc	0dda64	49c096	1	0	29	310ad	337563	3b6b93	0	1	49	022ab	0fd2cd	4cc876	0	0
10	2dbfe	06ed32	49c096	0	0	30	58856	19bab1	1db5c9	1	0	50	41155	27e966	66643b	0	0
11	56dff	06ed32	64e04b	1	0	31	58856	2cdd58	0edae4	1	0	51	208aa	27e966	33321d	1	1
12	2b6ff	237699	64e04b	0	1	32	58856	166eac	076d72	0	0	52	208aa	33f4b3	19990e	0	1
13	2b6ff	31bb4c	727025	1	0	33	2c42b	166eac	43b6b9	1	0	53	10455	19fa59	19990e	1	0
14	15b7f	18dda6	393812	1	1	34	56215	0b3756	21db5c	0	1	54	0822a	2cfd2c	4ccc87	0	1
15	15b7f	2c6ed3	5c9c09	1	1	35	2b10a	259bab	21db5c	1	1	55	04115	2cfd2c	266643	0	0
16	15b7f	163769	2e4e04	0	0	36	2b10a	12cdd5	10edae	0	1	56	0208a	2cfd2c	133321	1	1
17	0adbf	2b1bb4	572702	1	1	37	55885	12cdd5	4876d7	1	1	57	0208a	167e96	499990	1	0
18	056df	158dda	572702	1	1	38	55885	2966ea	643b6b	0	0	58	0208a	2b3f4b	64ccc8	0	1
19	42b6f	2ac6ed	572702	0	0	39	2ac42	34b375	643b6b	1	0	59	41045	2b3f4b	726664	0	0
20	215b7	356376	6b9381	1	0	40	55621	34b375	321db5	1	1	60	20822	2b3f4b	793332	1	1

A5/1 Key mixing (2)

<u>iter</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>	<u>iter</u>	<u>LFSR1</u>	<u>LFSR2</u>	<u>LFSR3</u>	<u>maj</u>	<u>out</u>
61	20822	159fa5	7c9999	1	0	81	3cf48	141dd5	39caf9	1	0
62	20822	2acfd2	3e4ccc	0	0	82	1e7a4	2a0eea	39caf9	1	1
63	10411	2acfd2	5f2666	1	1	83	0f3d2	350775	39caf9	0	0
64	48208	3567e9	5f2666	0	1	84	479e9	3a83ba	1ce57c	0	1
65	24104	3ab3f4	2f9333	0	1	85	23cf4	3d41dd	0e72be	1	1
66	52082	1d59fa	2f9333	1	1	86	11e7a	3d41dd	07395f	1	0
67	52082	2eacfd	57c999	0	0	87	08f3d	3d41dd	439caf	1	1
68	69041	2eacfd	2be4cc	0	0	88	4479e	3d41dd	21ce57	0	0
69	74820	2eacfd	15f266	1	1	89	4479e	3ea0ee	10e72b	0	1
70	74820	37567e	4af933	0	1	90	4479e	3f5077	487395	1	0
71	7a410	3bab3f	4af933	1	0	91	223cf	3f5077	2439ca	0	0
72	3d208	1dd59f	657c99	0	0	92	511e7	1fa83b	2439ca	1	0
73	1e904	0eeacf	657c99	1	0	93	511e7	0fd41d	521ce5	0	1
74	1e904	077567	72be4c	0	1	94	288f3	27ea0e	521ce5	1	0
75	4f482	03bab3	72be4c	1	1	95	288f3	33f507	290e72	0	0
76	67a41	01dd59	395f26	1	0	96	54479	19fa83	548739	1	1
77	67a41	20eeac	1caf93	0	0	97	2a23c	0cfd41	548739	0	0
78	73d20	20eeac	4e57c9	1	1	98	1511e	0cfd41	2a439c	0	1
79	79e90	107756	672be4	0	0	99	0a88f	0cfd41	5521ce	0	0
80	79e90	283bab	7395f2	1	1	100	45447	0cfd41	2a90e7	1	1

Task 4: RC4

Rivest Cipher 4 (RC4)

RC4 is a stream cipher that generates the keystream from a secret internal state which consists of two parts:

- A permutation P of all 256 possible bytes.
- Two 8-bit index-pointers (denoted i and j).

P is initialized with a variable length key by means of the key-scheduling algorithm (KSA).

Then, the keystream is generated using the pseudo-random generation algorithm (PRGA) that updates the indexes i and j , modifies the permutation P and generates a random byte.

Key Scheduling Algorithm (KSA)

The KSA is used to initialize the permutation P starting from a key composed by L bytes. Typical values for L range from 40 to 256.

```
Input key =  $[k_0, k_1, \dots, k_{L-1}]$ ,  
         with  $k_i \in \{0, 1, \dots, 255\}$   
 $j \leftarrow 0$   
for  $i = 0, 1, \dots, 255$   
     $P[i] \leftarrow i$   
endfor  
for  $i = 0, 1, \dots, 255$   
     $j \leftarrow (j + P[i] + \text{key}[i \bmod L]) \bmod 256$   
     $P[i], P[j] \leftarrow P[j], P[i]$   
endfor  
Output  $P$ 
```

P is initialized with an identity permutation ($P[i] = i$).

Then, bytes of P are mixed iteratively in a way that depends on the key.

Pseudo-random generation Algorithm (PRGA)

For each iteration, PRGA modifies the state (represented by the permutation P and the pair of indexes i, j) and outputs a byte.

State P, i, j

$i \leftarrow (i + 1) \bmod 256$

$j \leftarrow (j + P[i]) \bmod 256$

$P[i], P[j] \leftarrow P[j], P[i]$

$K \leftarrow P[(P[i] + P[j]) \bmod 256]$

Output K

In each iteration,

- i is incremented,
- j is updated by adding the value $P[i]$,
- $P[i]$ and $P[j]$ are swapped.
- The output byte is element of P at the location $P[i] + P[j] \pmod{256}$

RC4-drop[n]

RC4 has many known vulnerabilities mainly related to the correlation between the key and the first bytes of the permutation P .

Most of them can be avoided by discarding the first n bytes of the output stream, from where it becomes RC4-drop[n].

Typical values for n are:

- $n = 768$
- $n = 3072$ (more conservative value)

RC4

- Define a Iterator that implements the RC4 stream cipher.
- Given a key and a drop number, encrypt and decrypt a message.

```
message = 'hello world!'
key = b'0123456789ABCDEF'

# create a StreamCipher instance for both Alice and Bob
alice = StreamCipher(key, prng=RC4, drop=10)
bob    = StreamCipher(key, prng=RC4, drop=10)

plaintextA = message.encode('utf-8')    # -> b'Hello world!'
ciphertext = alice.encrypt(plaintextA)  # -> b' /\x9e\xfb\x83@\x81}\xa9\xd0\xd4\xd5\xfb'
plaintextB = bob.decrypt(ciphertext)    # -> b'Hello world!'
```