



# How we measure and optimize for RAIL in V8's GC

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# Contents





# Who We Are - V8 GC Team Munich



Hannes  
Payer



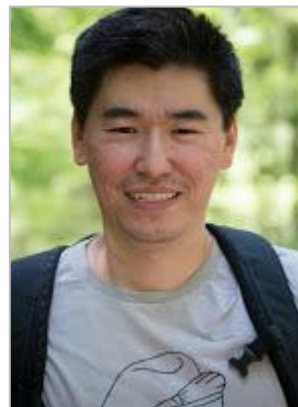
Jochen  
Eisinger



Marcel  
Hlopko



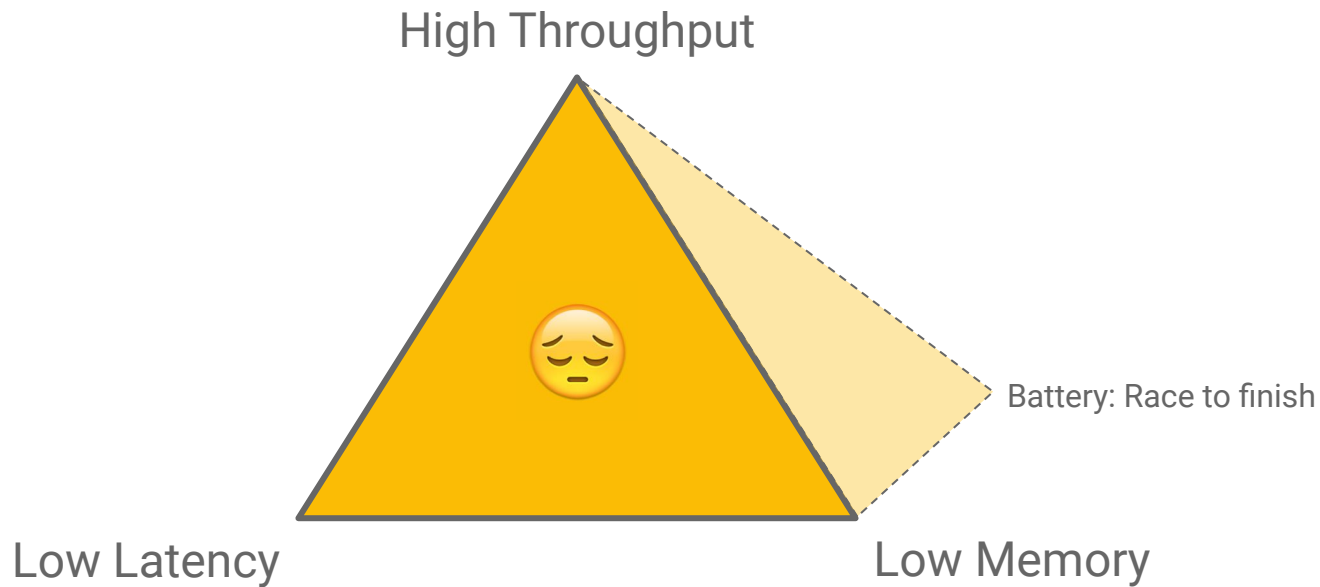
Michael  
Lippautz



Ulan  
Degenbaev

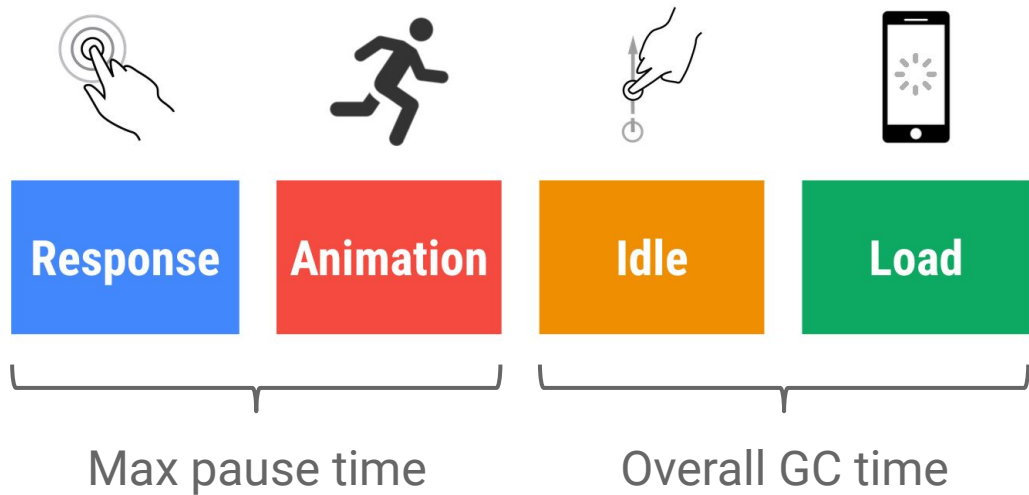
# V8 GC and RAIL

# The Impossible Garbage Collection Triad



# RAIL and GC Metrics

- Ideally: No GC... ever
- In practice: Prioritize

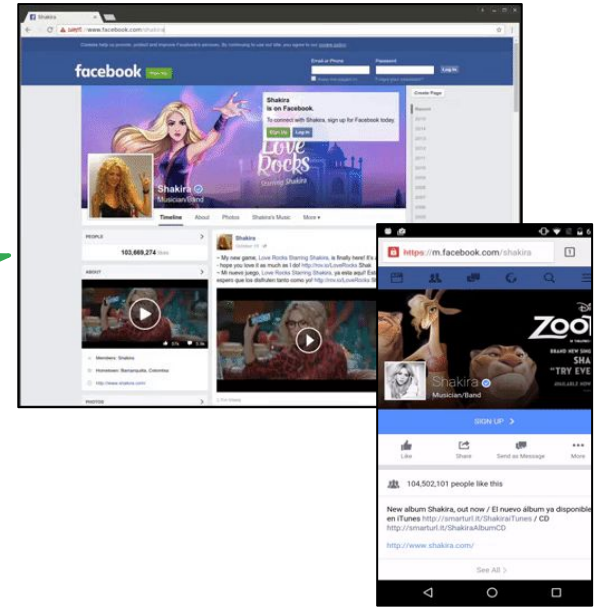


Fine print

*Without consuming too  
much memory 😊*

# Real World Benchmarking

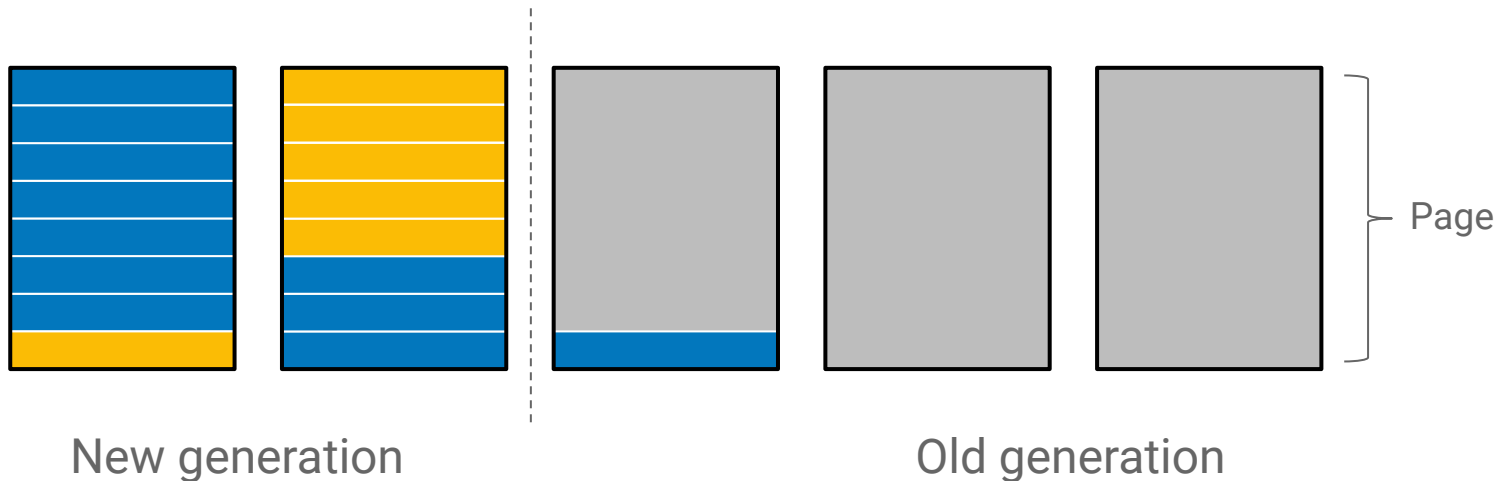
- Catapult (Telemetry)
- Record/replay for real-world websites
- Mobile and desktop



# V8's Generational Garbage Collector

Generational hypothesis: “Most objects die shortly after their allocation”

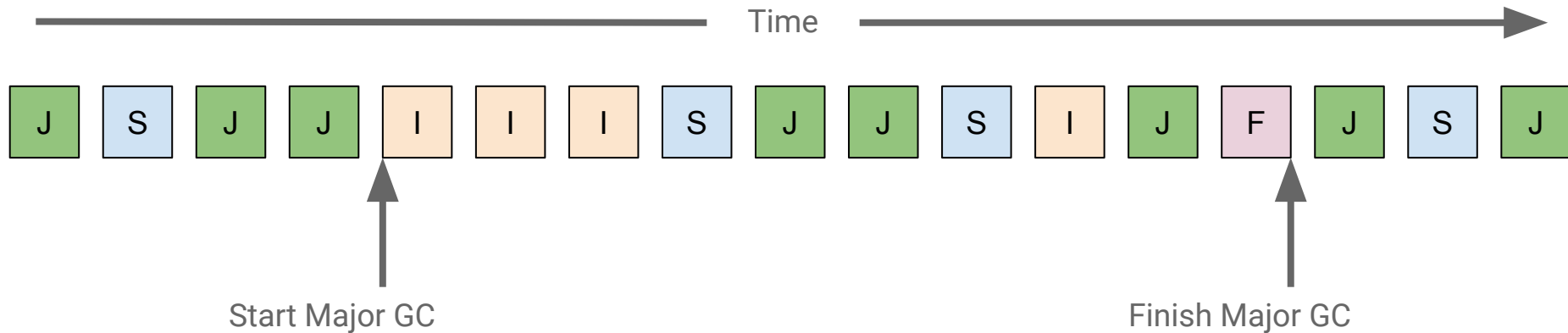
- New generation: Semi-space **Scavenger** (Cheney)
- Old generation: **Mark-Compact** (and Sweep)





# GC Events

- J** JavaScript code
- S** Minor GC: **Scavenger** (~0-10 ms)
- I** Major GC: Incremental Marking (~0.01-*CONFIGURABLE* ms)
- F** Major GC: Final **Mark-Compact** Collection (~4-20 ms)



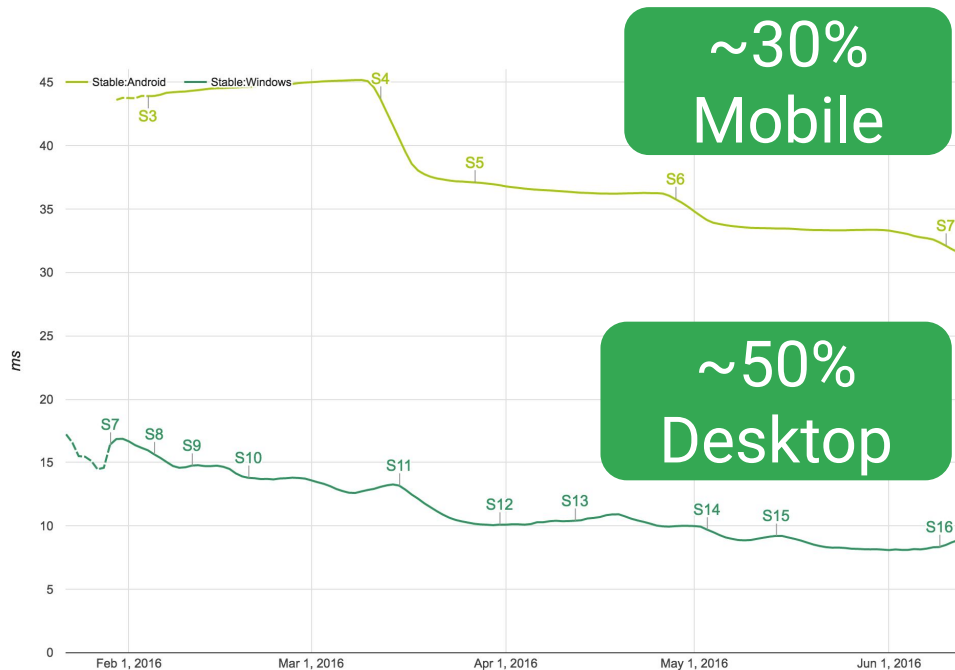
# V8 GC Update

Focusing on **A**

# Orinoco

Mostly parallel and concurrent GC  
without strict generational boundaries  
to  
*reduce jank* and memory consumption  
while providing high throughput

... landing incrementally



UMA: Major GC final pause **F** (50 %-ile)

# Evacuation

- *Copy* objects within semi space of new generation
- *Move* objects from new to old generation
- *Copy* objects within the old generation
- Re-write remembered set pointers

**Writing memory is expensive**

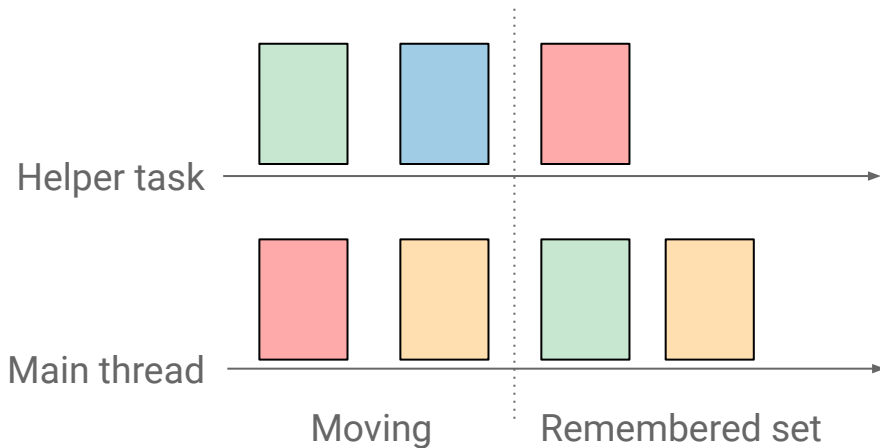
*Orinoco*

**Design GC to utilize available resources**

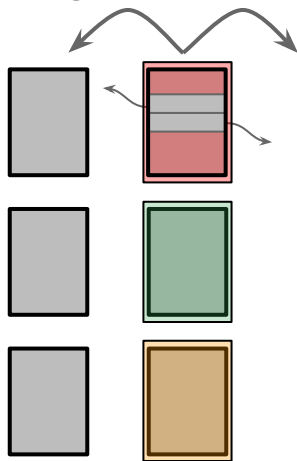
# Parallel Evacuation during Major GC

- Lock-step

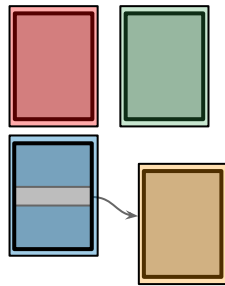
- Parallelize moving of memory based on pages
- Parallelize processing remembered set pointers



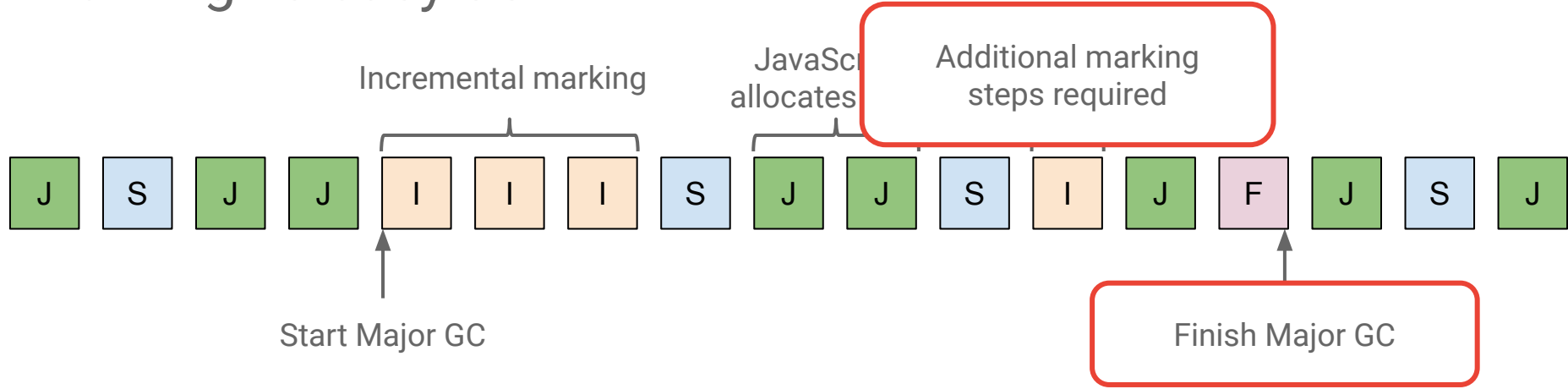
New generation



Old generation



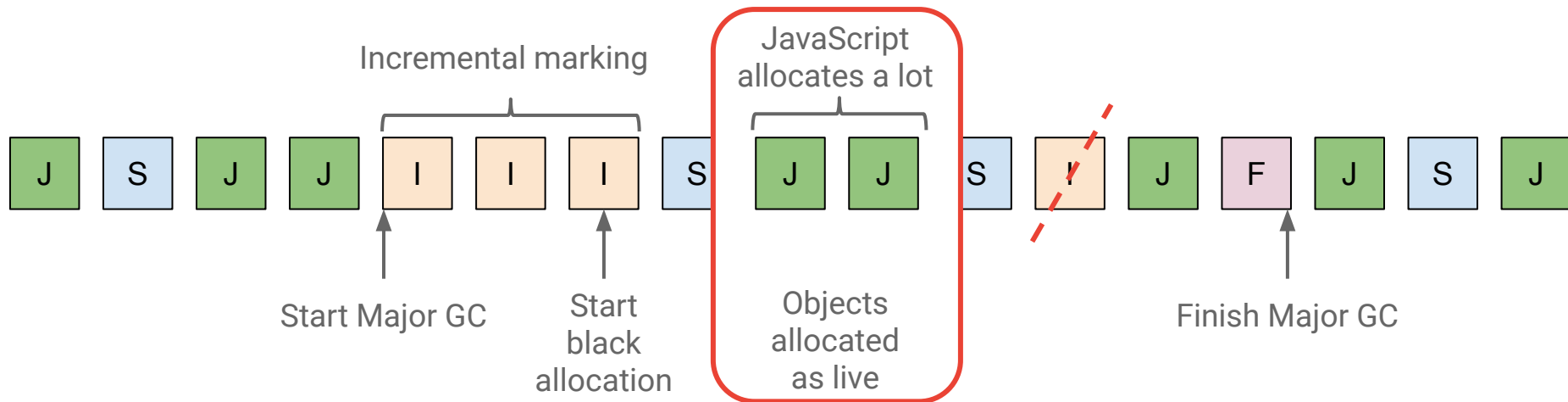
# Marking vs busy JS



Live heap gets bigger  
**More marking steps, longer finalization pause**



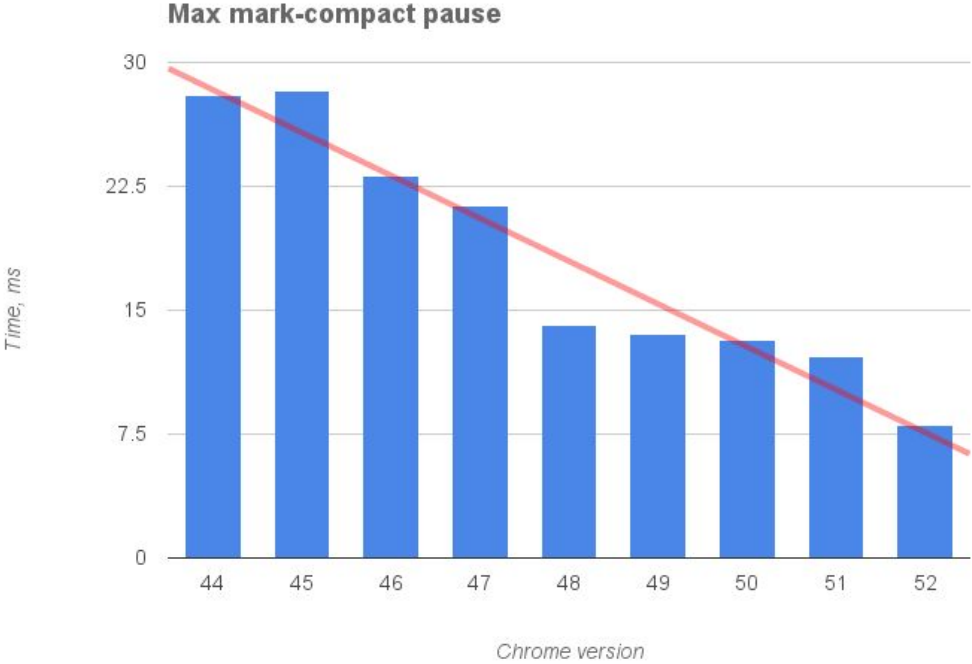
# Black Allocation



Assumption: Objects allocated during marking will survive the following GC

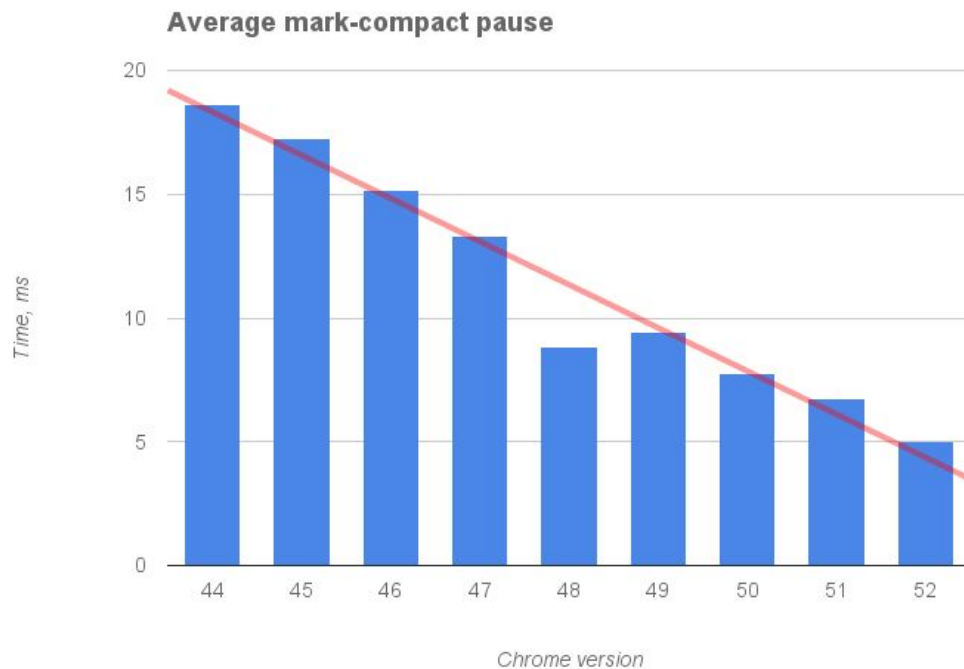
Objects are allocated as live (already marked)

# Some Results: Facebook during A





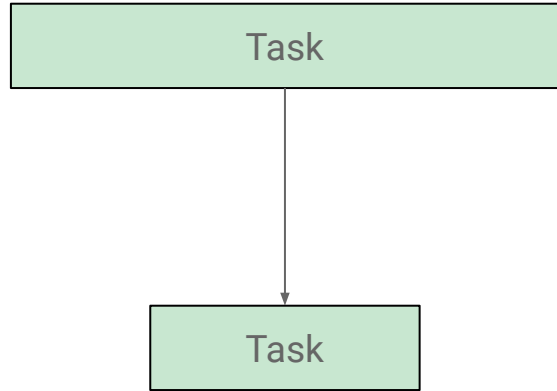
# Some Results: Facebook during A



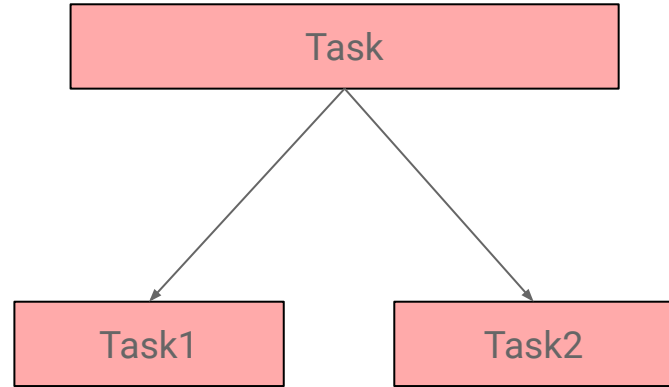
# Reflection and Future

# Two ways to reduce latency

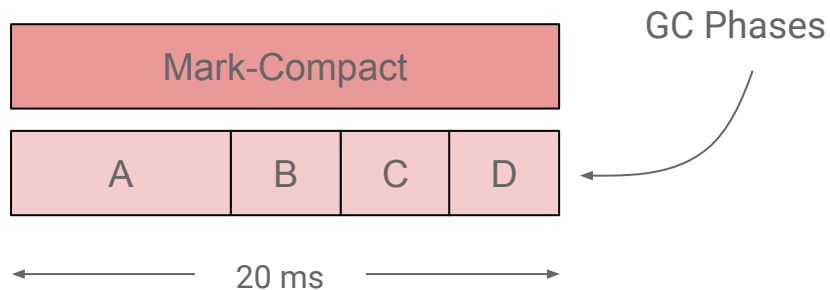
Optimize



Split up



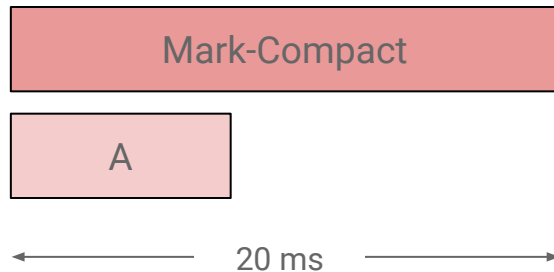
# Making GC more incremental



# Making GC more incremental

Before:

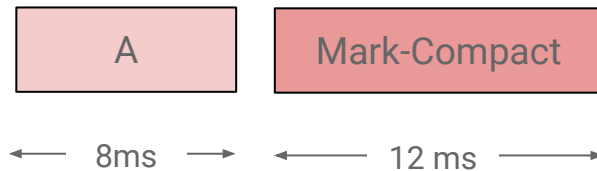
max-pause = 20ms



---

After:

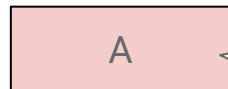
max-pause = 12ms



# Making GC more incremental

Before:

max-pause = 20ms

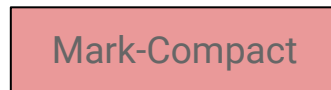
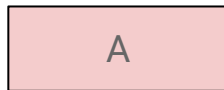


← 20 ms

**40% latency  
improvement?**

After:

max-pause = 12ms



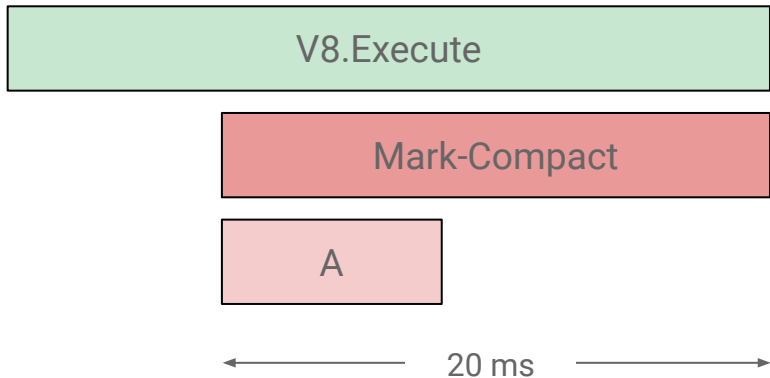
← 8ms →

← 12 ms →

# Making GC more incremental

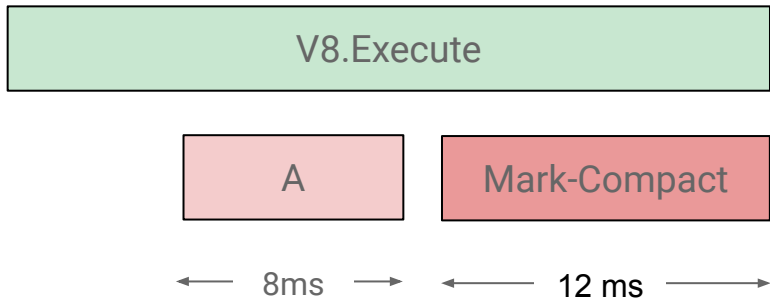
Before:

max-pause = 20ms



After:

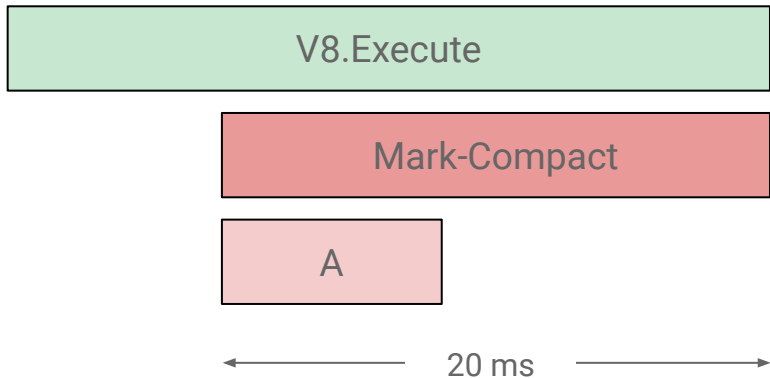
max-pause = 12ms



# Making GC more incremental

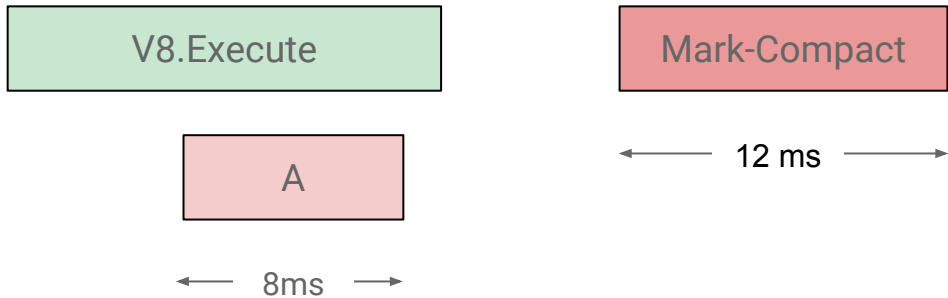
Before:

max-pause = 20ms



After:

max-pause = 12ms





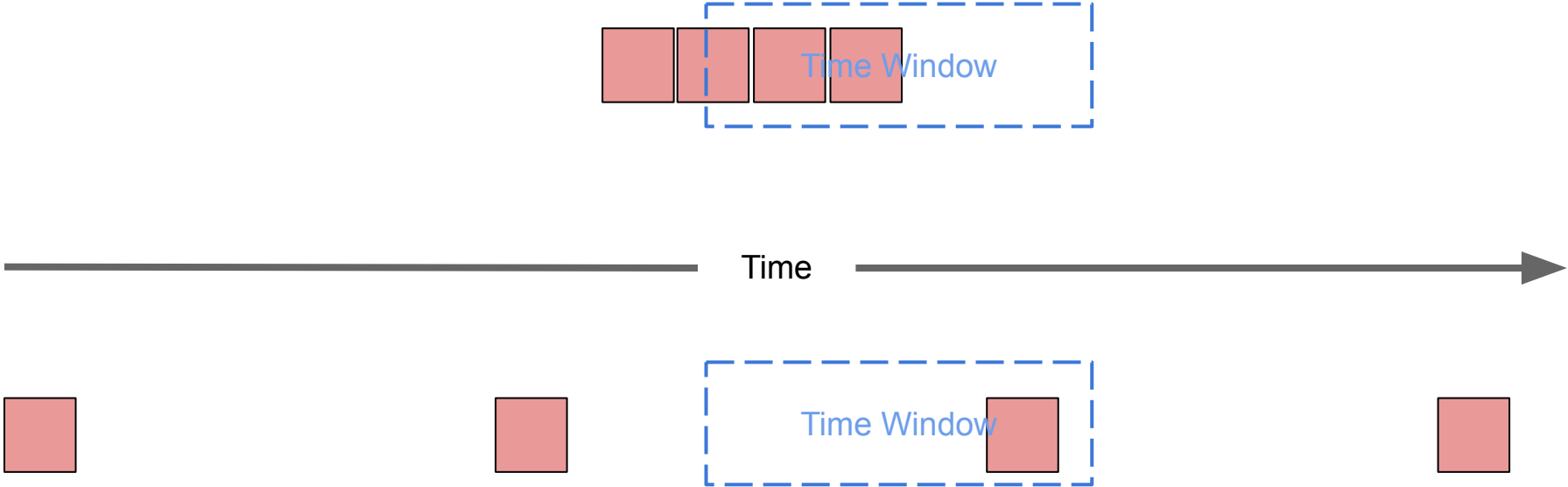
# User facing metrics from PWM and system health

- **Response:** response latency
  - minimize GC work in critical time window
- **Animation:** animation latency
  - minimize GC work in critical time window
- **Idle:** responsiveness risk a.k.a responsiveness hazard
  - minimize GC work in tasks longer than 50ms
- **Load:** time until load finishes
  - minimize GC work during page load

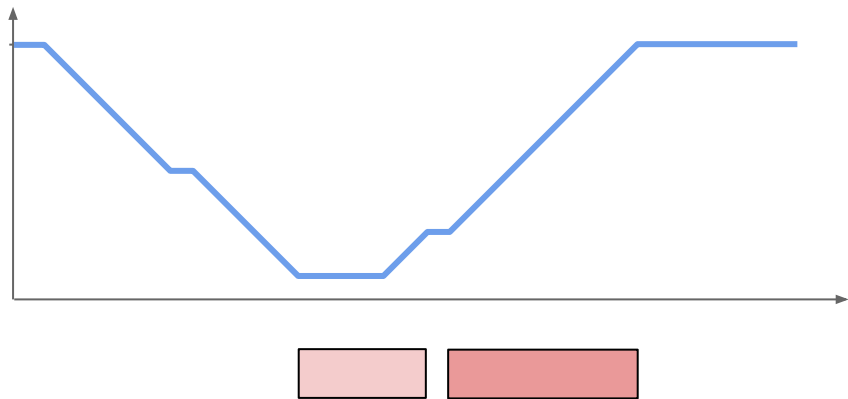
Overfitting: delay GC in RAL, post many 49ms tasks in I

Simpler model: consider only RA, critical time window position is arbitrary

# Uniform incremental steps



# Mutator Utilization



- Introduced by Cheng and Blleloch in 2001 for real-time GC.
- Called ***mutator utilization*** - application *mutates* reachability graph from GC POV
- Can be generalized to other tasks in Chrome

# Utilization function

Free time in window

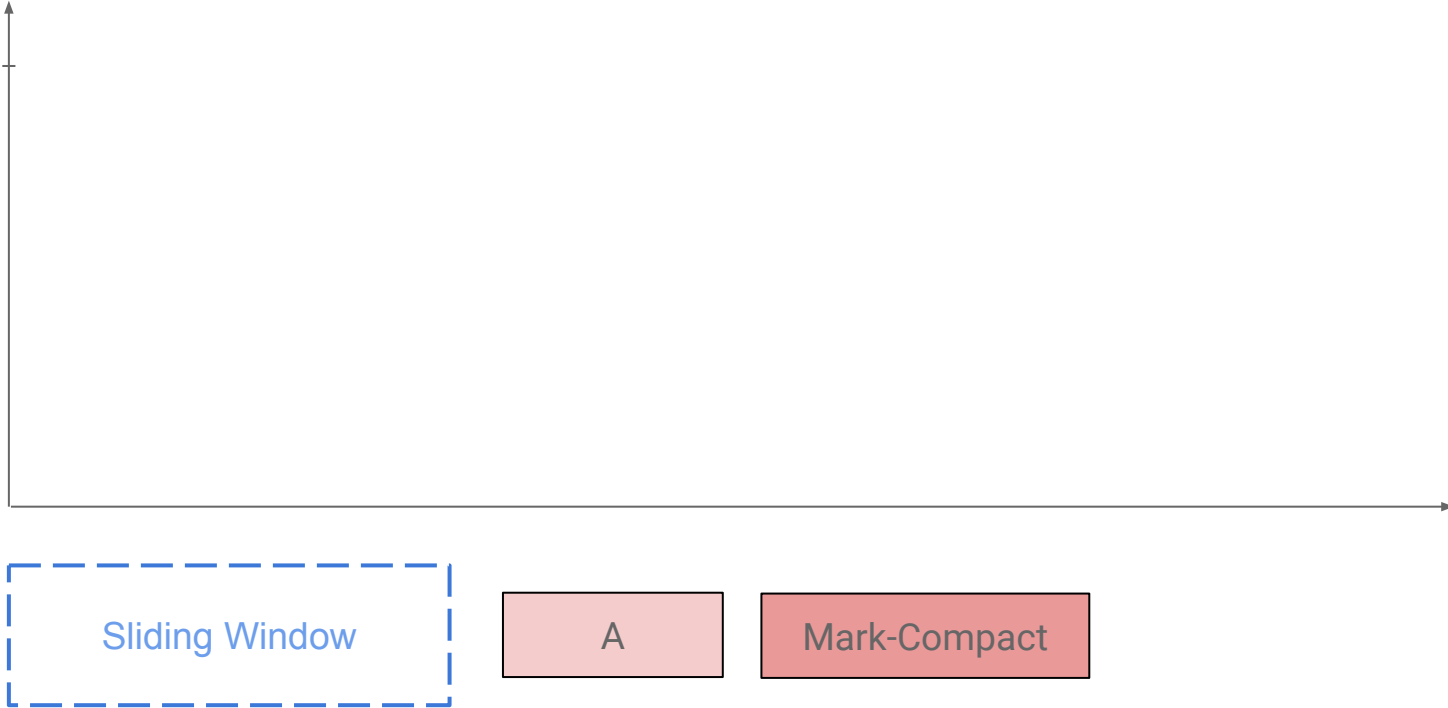
100%

Time

Sliding Window

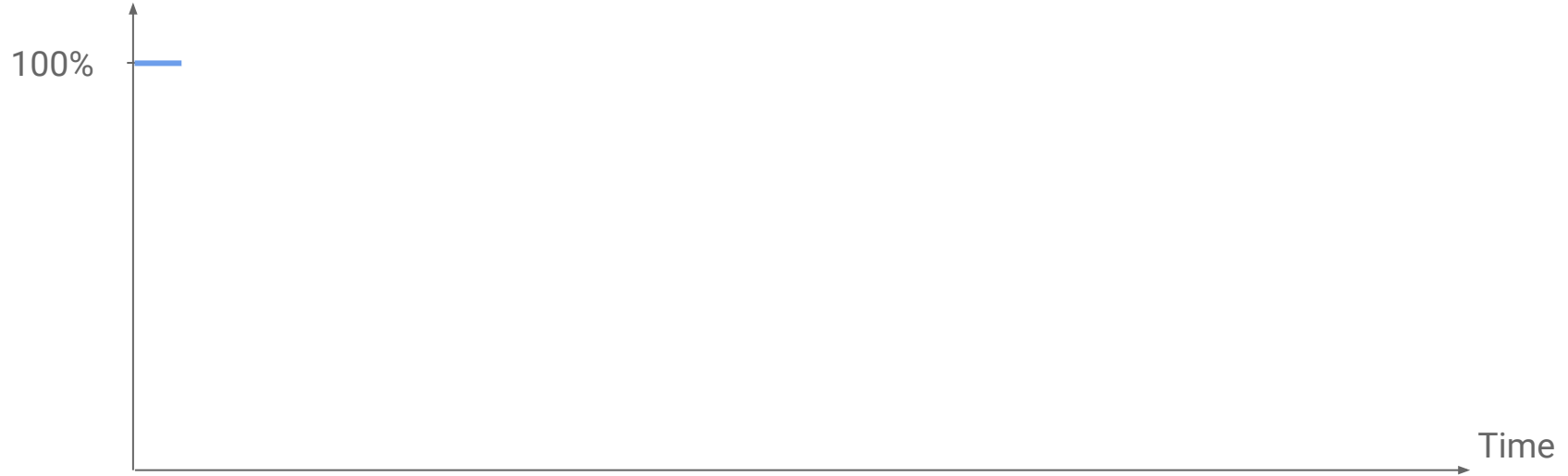
A

Mark-Compact



# Utilization function

Free time in window



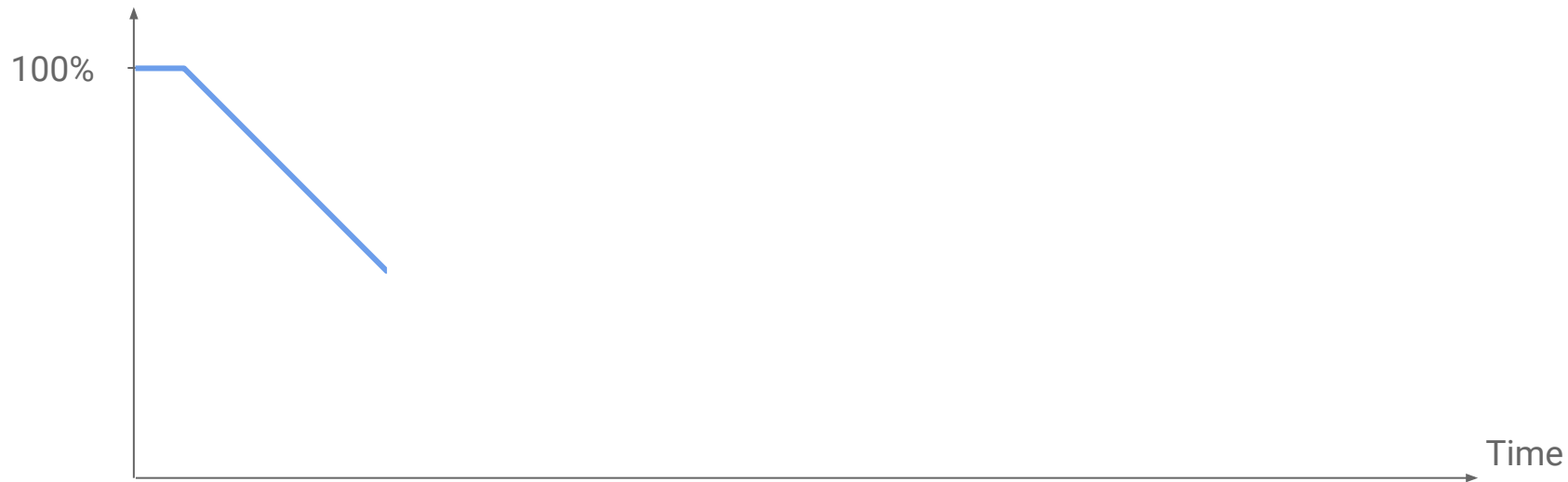
Sliding Window

A

Mark-Compact

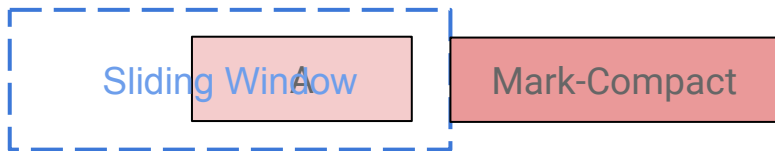
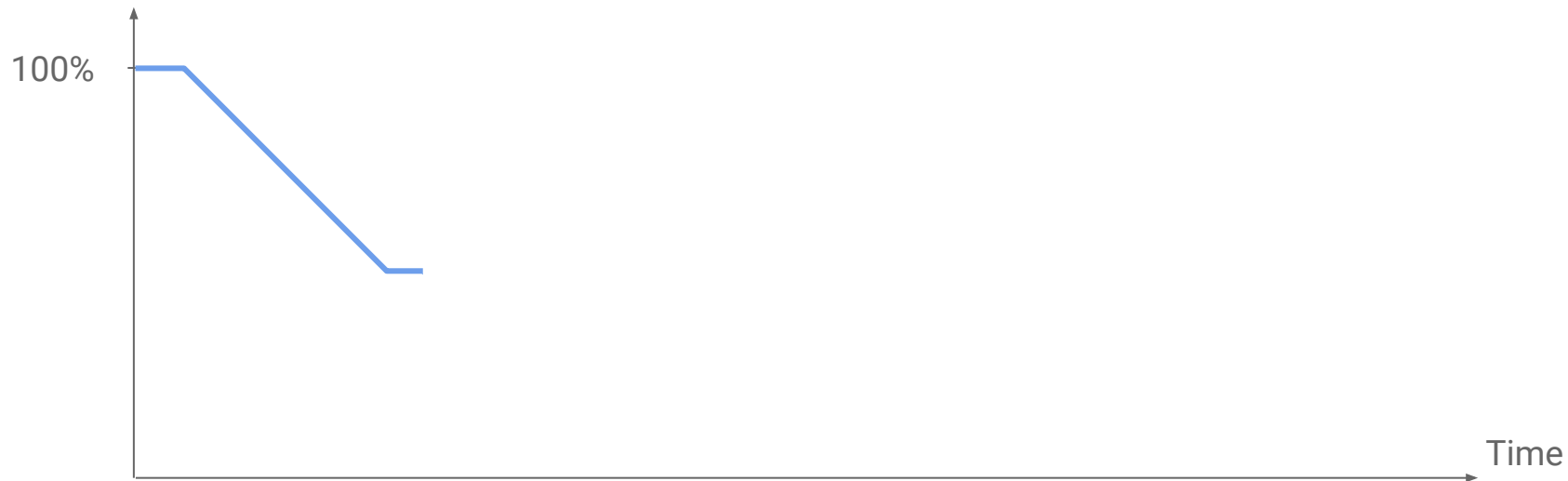
# Utilization function

Free time in window



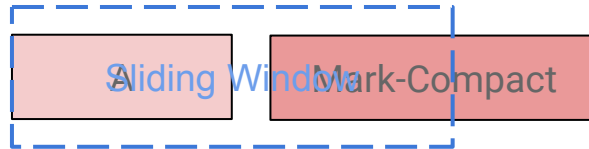
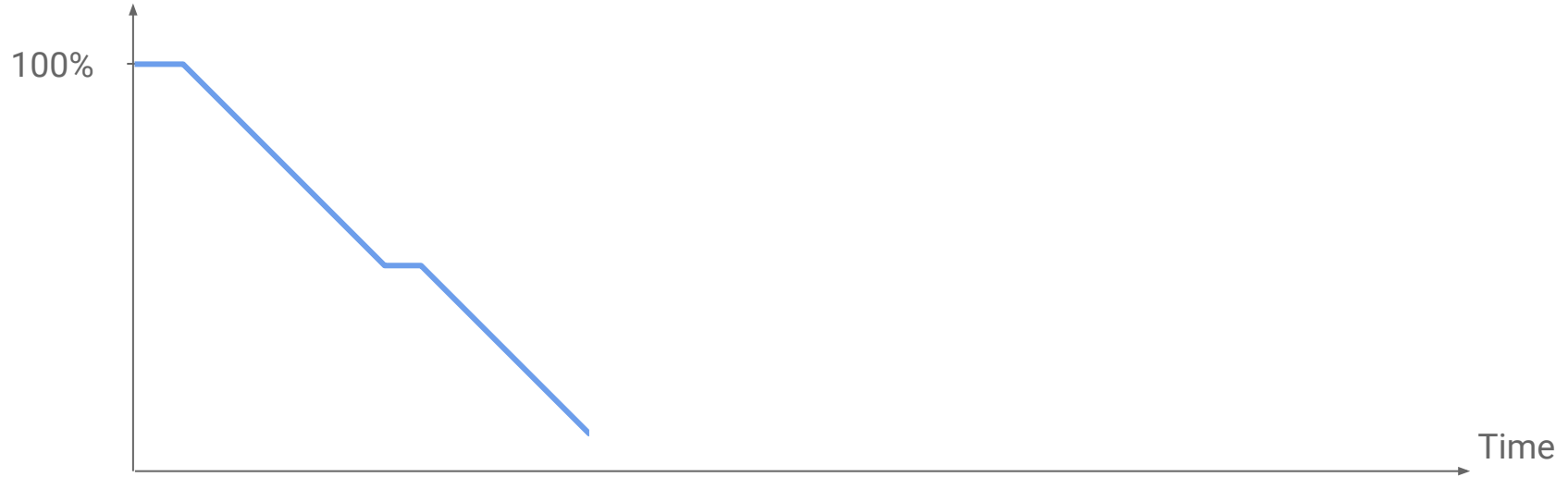
# Utilization function

Free time in window



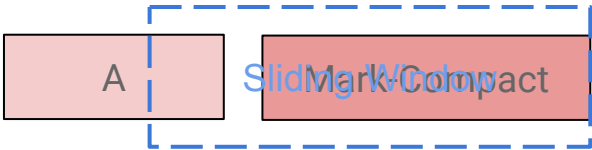
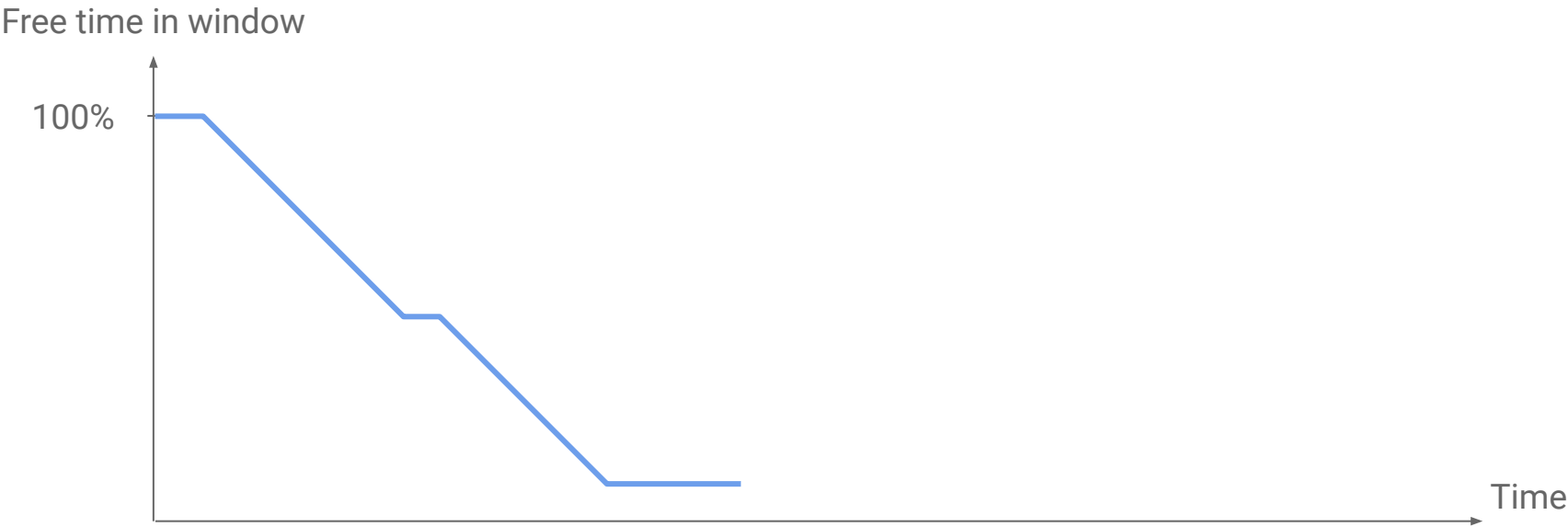
# Utilization function

Free time in window

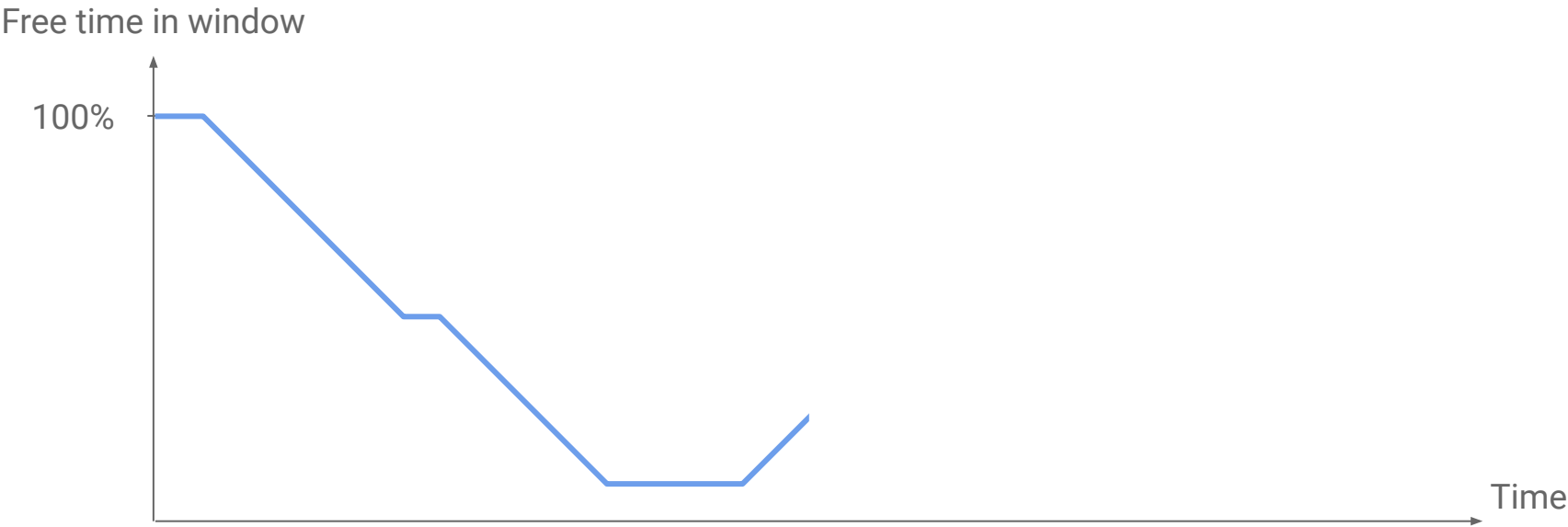




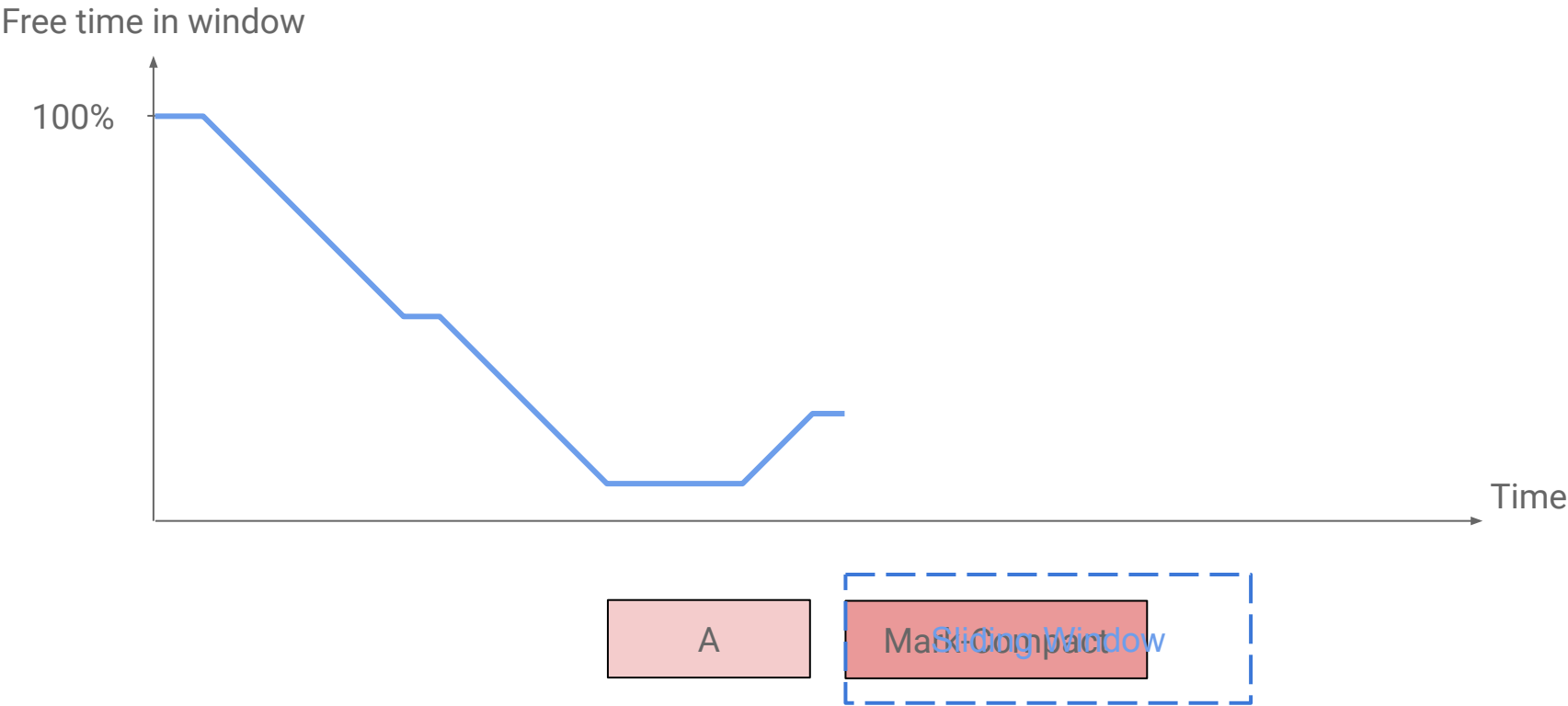
# Utilization function



# Utilization function

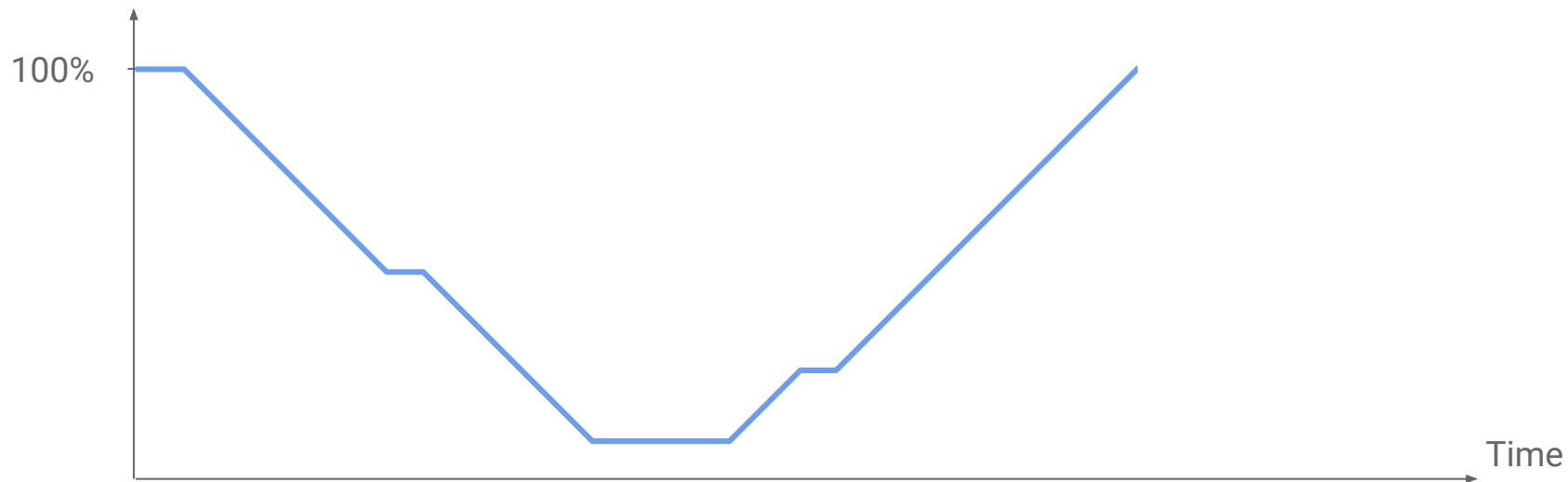


# Utilization function



# Utilization function

Free time in window



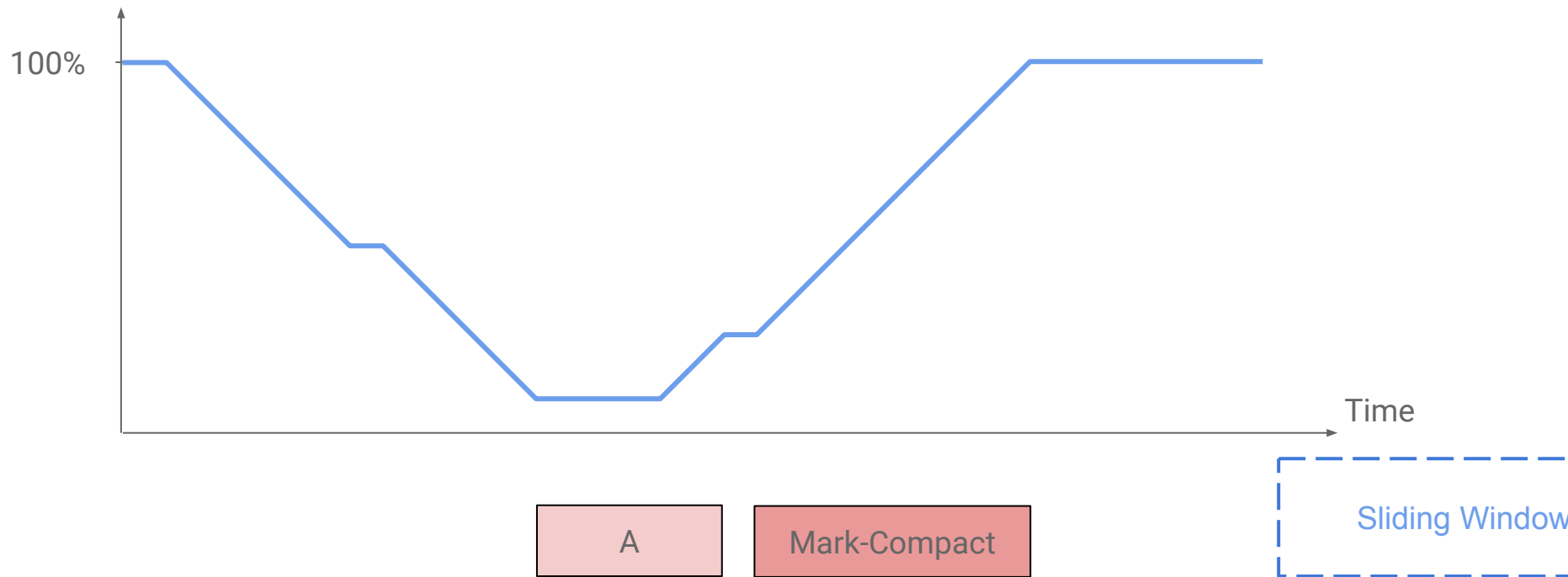
A

Mark-Compact

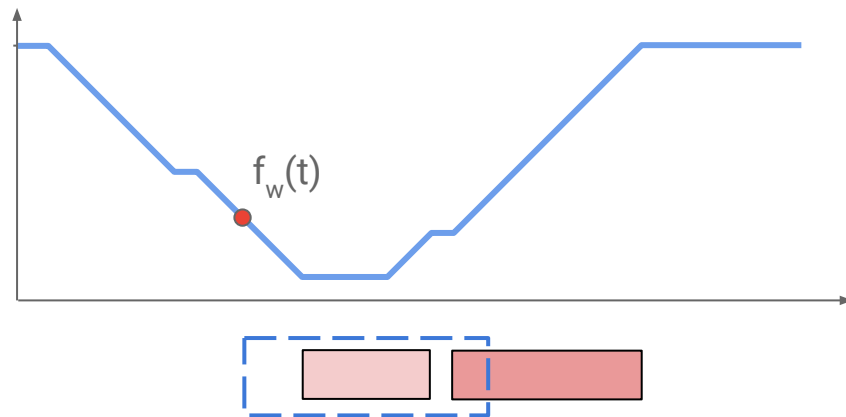
Sliding Window

# Utilization function

Free time in window

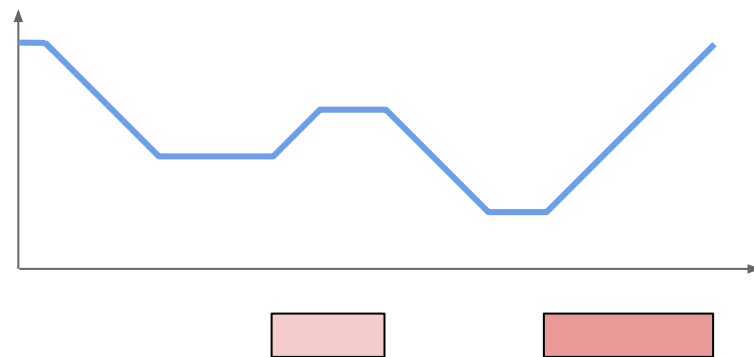
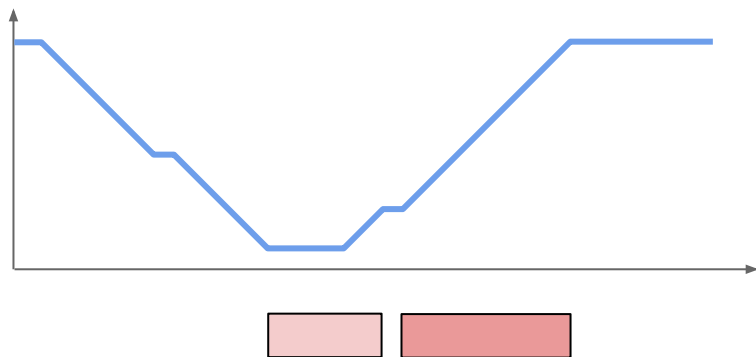


# Interpretations

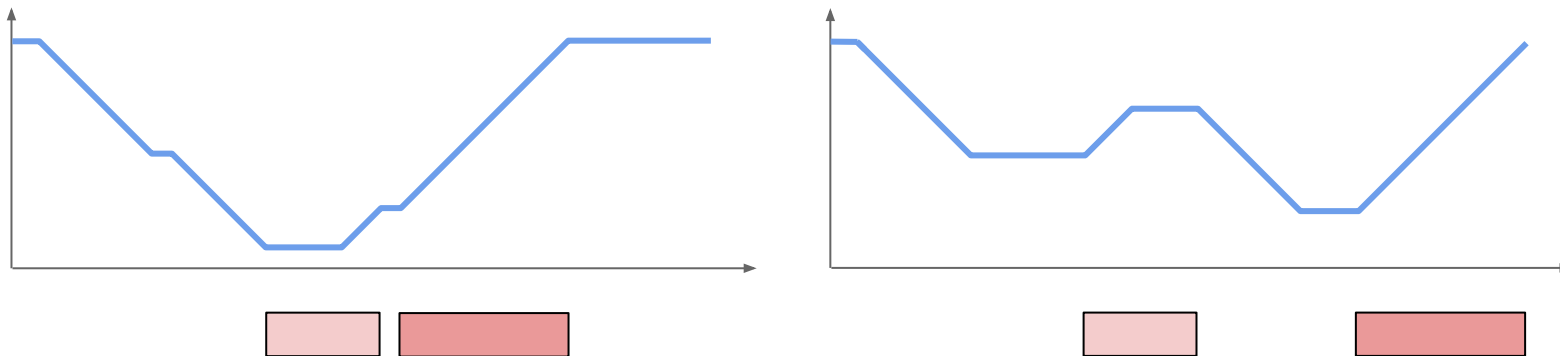


- $f_w(t)$  = fraction of time GC leaves to application in the time window  $[t, t+w)$ .
- $f_w(t)$  = probability that a high priority task arriving at any moment in  $[t, t+w)$  is not queued

# Utilization function



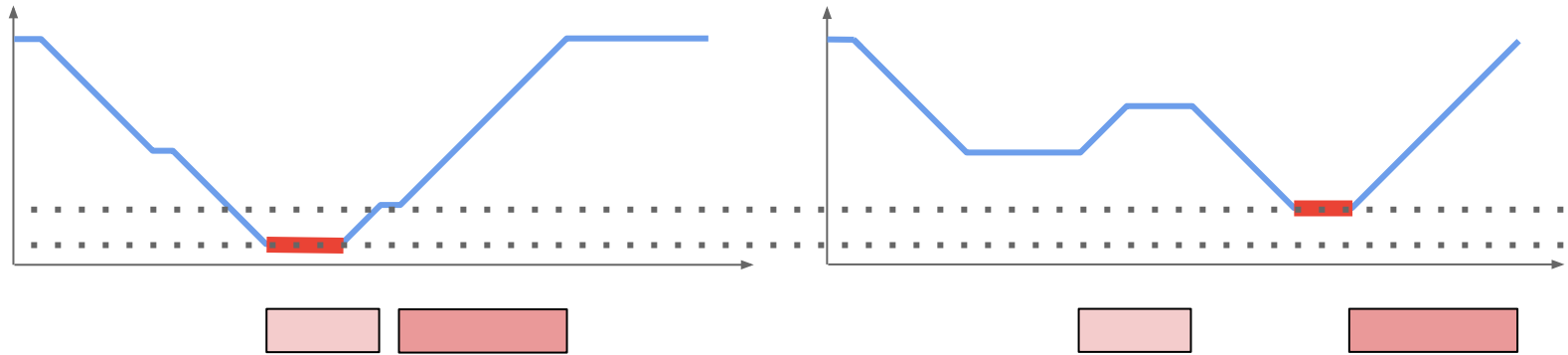
# Utilization function



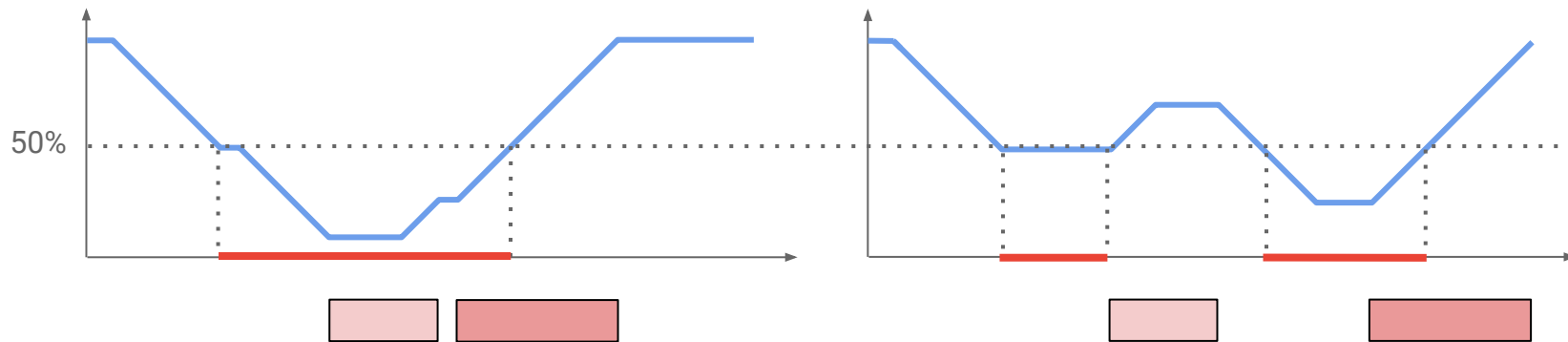
How to compare?



# Minimum mutator utilization

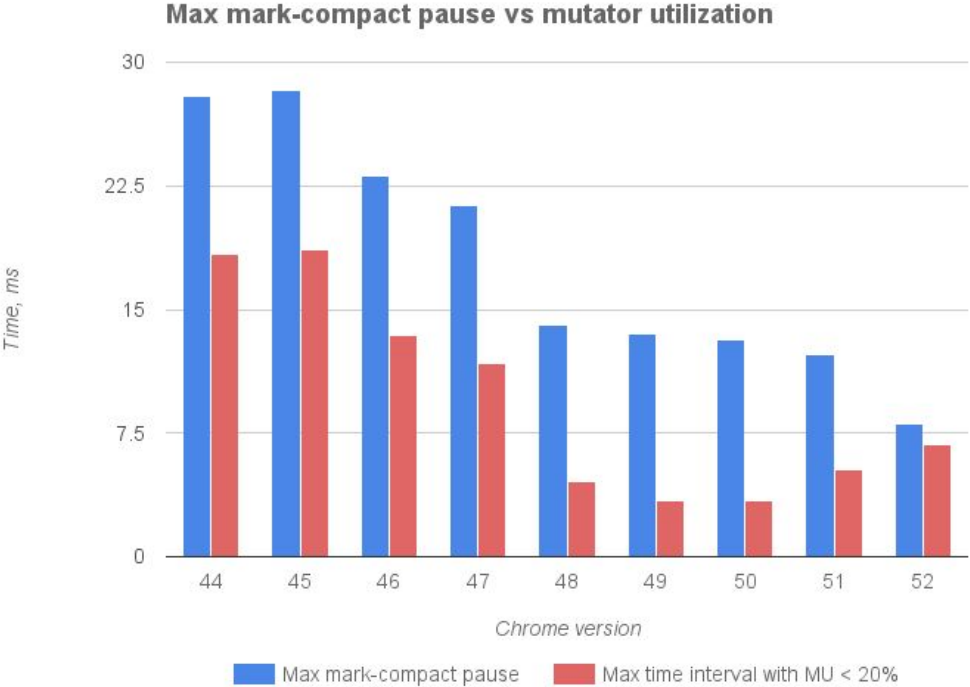


# Intervals with low mutator utilization



The largest interval where  $f(x) \leq 50\%$  for all points in the interval.

# Evaluation



# What we learned

- Incremental development of Orinoco - good idea!
  - Stability feedback via Canary channel
  - Performance feedback via UMA and telemetry
- Pause time distribution is not the whole performance story
  - Density of GC events on the timeline is also important
  - Mutator utilization can be good indicator of latency impact

# Thanks!

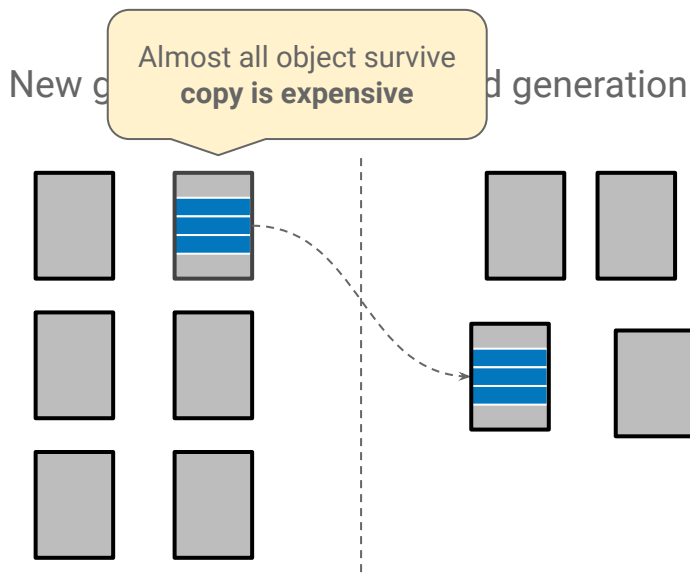
# Changes on the way

## Experimental

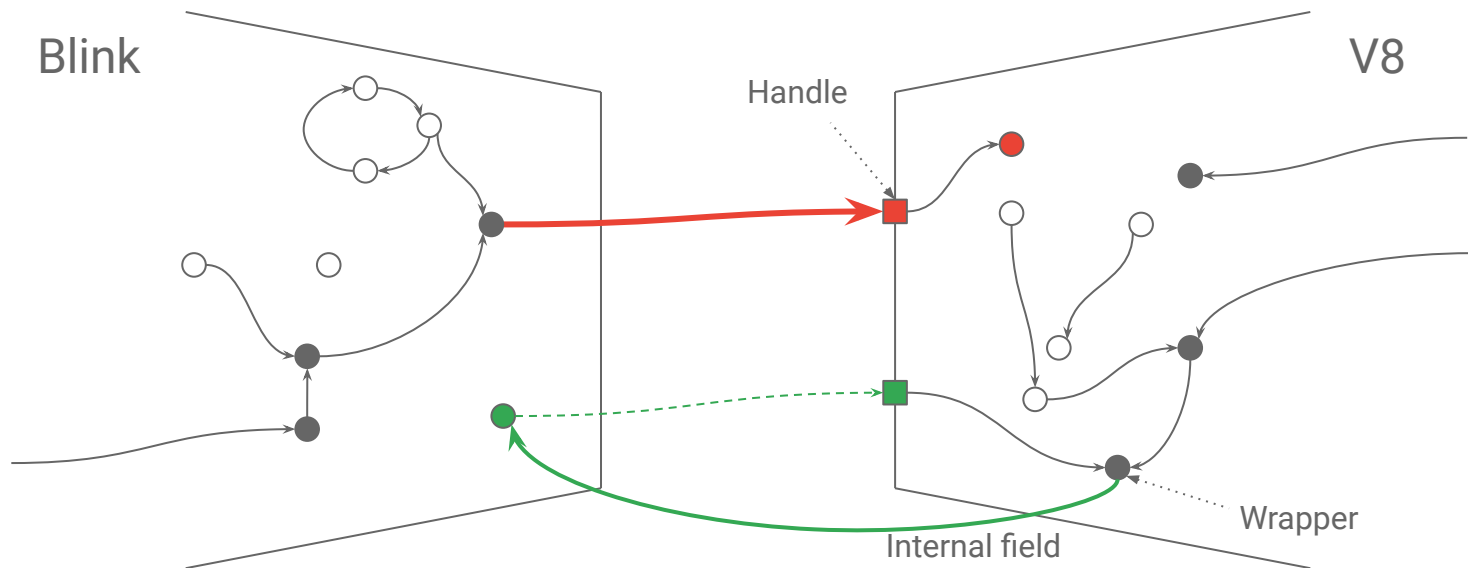
(behind flags)

# Unified Heap

- Allow pages to move from young to old generation without copying



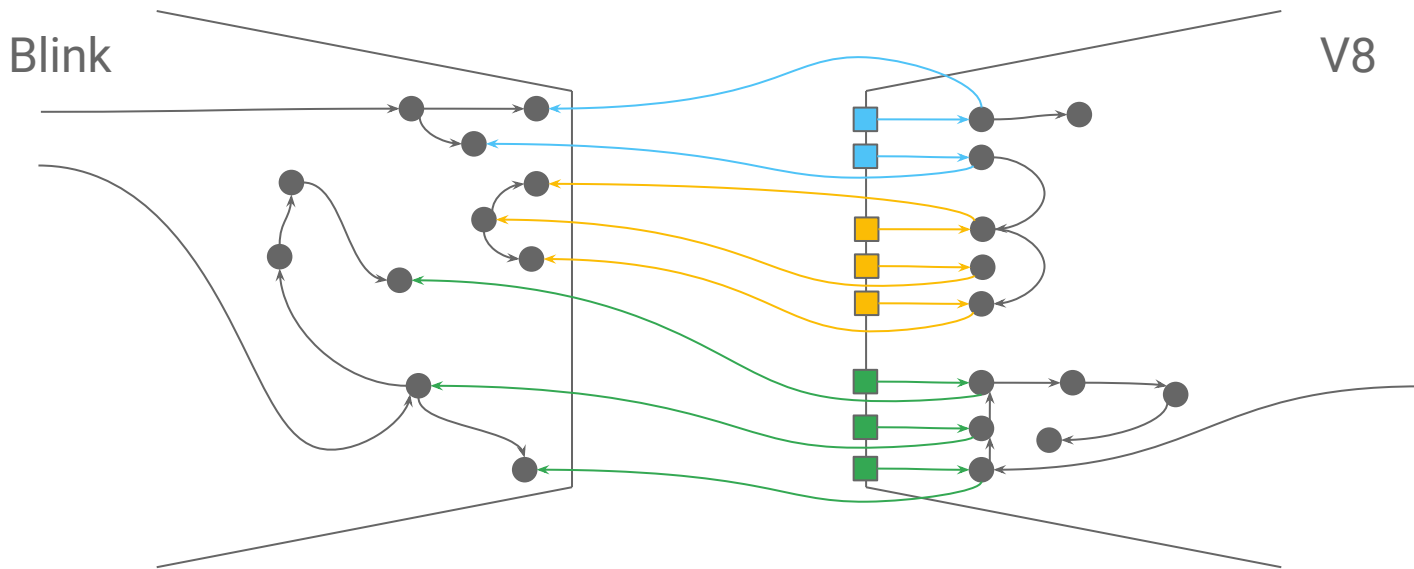
# Blink ⇔ V8 - Connecting the Worlds



Handles keep V8 objects alive  
Wrapper keep Blink objects alive



# Current Approach: Object Grouping



- Handles know which group they belong too (rule-based)
- Groups keep objects alive

# Becoming friends with Blink: Tracing of wrappers

